# PROJECT 1

Advanced Real Time Systems - ECE 5550G

The Blink AnalogRead program implements two periodic tasks:

TaskBlink

```
    Human Readable Name : Blink
    Task Stack Size : 128
    Task Priority : 2
    Task Period : 1s
```

Code Snippet :

```
// TaskBlink Function Definition
void TaskBlink(void *pvParameters) // This is a task.
{
   (void) pvParameters;

   // Initialize digital LED_BUILTIN on pin 13 as an output.
   pinMode(LED_BUILTIN, OUTPUT);

   for (;;)
   {
      digitalWrite(LED_BUILTIN, HIGH);
      vTaskDelay( 1000 / portTICK_PERIOD_MS ); // wait for one second digitalWrite(LED_BUILTIN, LOW);
      vTaskDelay( 1000 / portTICK_PERIOD_MS ); // wait for one second
   }
}
```

#### TaskAnalogRead

o Human Readable Name : AnalogRead

o Task Stack Size: 128

o Task Priority: 1

o Task Period: 15ms

o Code Snippet:

The modified **AnalogRead Blink** program:

- For each task we get and store the task handle.
- Using the <u>task handle</u>, we get the assigned <u>task name</u>.
- Every time the task is executed the task name is printed.

The task periods of the TaskBlink and TaskAnalogRead have been changed to 500 ms and 100 ms respectively to help in debugging and printing the task name in Serial Monitor.

Code Snippet for TaskBlink:

Code Snippet for TaskAnalogRead:

```
void TaskAnalogRead(void *pvParameters)
{
   (void) pvParameters;

   const TaskHandle_t task_handle = xTaskGetCurrentTaskHandle();
   const char *task_name = pcTaskGetName(task_handle);

   for (;;)
   {
        // Print the task name
        Serial.print(task_name);
        Serial.print(" : ");

        // read the input on analog pin 0:
        int sensorValue = analogRead(A0);
        // print out the value you read:
        Serial.println(sensorValue);

        // Wait for 100 ms
        vTaskDelay(pdMS_T0_TICKS(100));
    }
}
```

#### Output of the Serial Monitor for Question 2:

```
Output
        Serial Monitor ×
Message (Enter to send message to 'Arduino Mega or Mega 2560' on 'COM7')
20:32:39.035 -> AnalogR : 290
20:32:39.165 -> AnalogR : 300
20:32:39.255 -> AnalogR : 295
20:32:39.376 -> AnalogR : 289
20:32:39.465 -> AnalogR : 298
20:32:39.592 -> Blink
20:32:39.592 -> AnalogR : 289
20:32:39.686 -> AnalogR : 287
20:32:39.776 -> AnalogR : 296
20:32:39.905 -> AnalogR : 286
20:32:39.986 -> AnalogR : 287
20:32:40.125 -> AnalogR : 294
20:32:40.125 -> Blink
20:32:40.225 -> AnalogR : 283
20:32:40.305 -> AnalogR : 289
20:32:40.445 -> AnalogR : 293
20:32:40.535 -> AnalogR : 282
20:32:40.625 -> AnalogR : 290
```

Output of the Serial Monitor:

```
Output
        Serial Monitor ×
Message (Enter to send message to 'Arduino Mega or Mega 2560' on 'COM7')
21:33:45.169 -> ---- SETUP SECTION ----
21:33:45.169 -> ---- END OF SECTION ----
21:33:45.169 ->
21:33:45.169 -> Running : Task2
21:33:45.169 -> Blocking : Task2
21:33:45.169 -> Running : Task1
21:33:45.169 -> Blocking : Task1
21:33:45.169 -> Running : Task3
21:33:45.169 -> Blocking : Task3
21:33:47.388 -> Running : Task1
21:33:47.388 -> Blocking : Task1
21:33:48.476 -> Running : Task2
21:33:48.476 -> Blocking : Task2
21:33:49.577 -> Running : Task1
21:33:49.577 -> Blocking : Task1
21:33:50.656 -> Running : Task3
21:33:50.656 -> Blocking : Task3
21:33:51.757 -> Running : Task2
21:33:51.757 -> Blocking : Task2
21:33:51.757 -> Running : Task1
```

#### From the output we can observe:

- Each of the tasks are running periodically at the specified intervals.
- Whenever the tasks have the same arrival time, the highest priority task is executed first followed by the next highest priority task and so on. An example of such an instant is at the timestamp 21:33:45.169, where all three tasks have the same arrival time but are executed based on their priorities.
- If a lower priority task is already running, it is preempted by a task which has a higher priority.

So we can conclude that FreeRTOS is running a priority-based scheduler.

The objective is to modify the code such that the lowest priority job is selected instead of the highest priority job.

This can be accomplished by <u>inverting</u> all the <u>arithmetic</u>, <u>comparison</u> and <u>sorting</u> <u>operations</u>, and changing the initial value of the variables associated with selecting a job for execution. By making these changes to the tasks.cpp file the lowest priority job will be treated as a highest priority job and vice-versa.

The list of changes made to the task.h:

- Defined two macro flags in the task.h:
  - P\_TEST: When set to 1, the priority selection order is inverted by applying the changes described below to tasks.cpp file and task.h file. When set to 0, the original code behavior is restored.
  - **P\_DEBUG**: When set to 1, Debug output is printed on the Serial Monitor. When set to 0, the debug is not displayed.
- The macro definition tskidle\_PRIORITY is changed from tskidle\_PRIORITY ((UBaseType\_t) @U) to tskidle\_PRIORITY ((UBaseType\_t) (configMAX\_PRIORITIES 1U))

The list of changes made to the tasks.cpp file:

- Included the "Arduino.h" file to perform debugging via Serial.println()
- In the taskRECORD\_READY\_PRIORITY(uxPriority) macro definition, (uxPriority) > uxTopReadyPriority is changed to (uxPriority) < uxTopReadyPriority</p>
- In the taskSELECT\_HIGHEST\_PRIORITY\_TASK() macro definition, ++uxTopPriority is changed to --uxTopPriority
- The initial value of the uxTopReadyPriority variable is changed to uxTopReadyPriority = configMAX\_PRIORITIES 1U
- The initial value of the uxTopUsedPriority variable is changed to uxTopUsedPriority = 0
- The macro invocation listSET\_LIST\_ITEM\_VALUE(&( pxNewTCB->xEventListItem ),

  ( TickType\_t ) configMAX\_PRIORITIES ( TickType\_t ) uxPriority) is changed to listSET\_LIST\_ITEM\_VALUE(&( pxNewTCB->xEventListItem ), ( TickType\_t ) uxPriority) so that items are sorted in order with respect to inverted priority value.
- The comparison statement (pxCurrentTCB->uxPriority >= pxNewTCB->uxPriority) is changed to (pxCurrentTCB->uxPriority <= pxNewTCB->uxPriority)

- The comparison statement (pxCurrentTCB->uxPriority > pxNewTCB->uxPriority) is changed to (pxCurrentTCB->uxPriority < pxNewTCB->uxPriority)
- The comparison statement (uxNewPriority > uxCurrentBasePriority) is changed to (uxNewPriority < uxCurrentBasePriority)
- The comparison statement (uxNewPriority >= pxCurrentTCB->uxPriority) is changed to (uxNewPriority <= pxCurrentTCB->uxPriority)
- The comparison statement (pxTCB->uxPriority >= pxCurrentTCB->uxPriority) is changed to (pxTCB->uxPriority <= pxCurrentTCB->uxPriority)
- In the Idle task creation statement, the priority is changed from portPRIVILEGE\_BIT to tskidle\_priority
- Many comparison statements changes can be seen in the tasks.cpp file.

#### Summary of the changes:

- So we can observe that wherever priority is being compared we have to reverse or invert the logic so that the task selection gets inverted.
- Event Lists are always sorted based on the priority order, so these lines also have to be changed to maintain the order of the lists.
- The Idle task is given the lowest priority value when it is created, so the priority of the idle task needs to be made the highest value when the priority is inverted.

### Observations in the output:

- The task which had lowest priority in Q3 (i.e Task3) has the highest priority now and the task which had the highest priority in Q3 (i.e Task2) has the lowest priority now, and so Task3 is executed first when the tasks arrive at the time.
- If a lower priority task is already running, it is preempted by a task which has a higher priority.