

Workshop on Public Key Infrastructure and Red Hat Certificate System

Asha Akkiangady, Huzaifa Sidhpurwala, Niranjana Mallapadi

Contents

1	Introduction to Public Key Cryptography	1
1.1	Encryption	2
1.1.1	Symmetric Encryption	2
1.1.2	Asymmetric Encryption	3
1.2	Message Digest	4
1.2.1	Message Authentication Code	4
1.3	Algorithms	4
1.4	Protocols	7
1.4.1	SSL	7
2	Introduction to Public Key Infrastructure	10
2.1	Common Terms used in PKI	10
2.2	Detailed look on certificate/CRL	10
2.2.1	Certificates	10
2.2.2	Basic Certificate Fields	11
2.2.3	Extensions	13
2.2.4	Revocation	18
2.2.5	CRL Fields	19
2.2.6	CRL Extensions	20
2.3	Exercises	21
3	Common Terms Used in RHCS	24
4	Red Hat Certificate System	25
4.1	Overview	25
4.2	Certificate Manager	25
4.2.1	Introduction	25
4.2.2	Installation	25
4.2.3	Key Features	25
4.2.4	Architecture	26
4.2.5	Interfaces	26
4.2.6	Features	26
4.2.7	Instance Layout	31
4.2.8	Exercises	31
4.3	Key Recovery Authority	31
4.3.1	Introduction	31
4.3.2	Installation	31

4.3.3	Installation Layout	32
4.3.4	Key Features	33
4.3.5	Architecture	33
4.3.6	Interfaces	33
4.3.7	Features	34
4.3.8	Exercises	36
4.4	Online Certificate Status Protocol	36
4.4.1	Introduction	36
4.4.2	Installation	36
4.4.3	Interfaces	36
4.4.4	Exercises	37
4.5	SmartCard Management	37
4.5.1	Introduction	37
4.5.2	Architecture	38
4.5.3	Common Terms Used with regard to Smartcard	38
4.5.4	Token Key Service	39
4.5.5	Token Processing Service	39

1 Introduction to Public Key Cryptography

Before we start discussing about public key cryptography, we will in general discuss about how system communicate and what are the various threat models that are associated with the communication medium and what are the tools to overcome them. [3]

Example1: The most common protocol used to communicate between 2 systems is TCP/IP. TCP/IP allows information to be sent from one system to another system directly or through many intermediate systems.

Below are some of the threat models associated with above communication:

- **Eavesdropping:**

Information remains intact, But it's privacy is compromised, For example: some one could learn credit card number, record a sensitive conversation or intercept a classified information.

- **Tampering:**

Information in transit is changed or replaced and then sent to the recipient. For example: Some one could alter an order of goods or change a persons resume.

- **Impersonation:**

Information passes to a person who poses as intended recipient. Impersonation can take 2 forms:

- **Spoofing:**

A person can pretend to be someone else, For example, a person can pretend to have email address *joe@example.net*, or computer can identify itself as a site called *www.example.net*, when it is not. This type of impersonation is called spoofing.

– **Misrepresentation:**

A person or organization can misrepresent itself, For example, suppose the site *www.example.net* pretends to be a furniture store when it's just a site which accepts payments but never sends any goods.

Example2: Consider Alice , Bob and attacker are the parties , where alice and bob want to communicate to each other and Attacker is a threat to the communication. [3]

- Alice can send a postcard to bob to communicate, but this method is very weak as any random eavesdropper could read the postcard.
- Alice could write a letter and put in an envelope and send it to bob, but the attacker could open the envelope and read the letter, both confidentiality and integrity of the message is lost.
- Alice could seal the envelope with wax , but this method too is inefficient as Attacker could read the letter and seal it as it is without bob knowing that it was read by attacker
- Alice could write a letter and put it in a safe which has 2 keys and send one key before to Bob , and send the safe across to bob, this ensures confidentiality, integrity but practically this is not implementable. Considering the number of messages that has to be transferred , it's impractical to implement the mail safe method.

To mitigate above threat models, we will look in to cryptology as one of the tools of the trade.

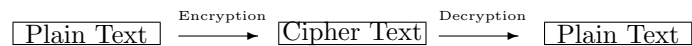
Cryptology: Cryptology is the theory of designing the various algorithms we use to provide security

Cryptography: Cryptography is the study of using these algorithms to secure systems and protocols.

1.1 Encryption

An encryption algorithm takes some data(called plaintext) and converts it to cipher-text under control of a key. cipher text contains random data which makes no sense without the key.

A key is a short random string (8-24 bytes):



When a message is encrypted and received , we cannot say if it's not tampered with. An encryption is strong when it can determine the number of possible keys. The attacker tries each key one at a time until he finds a key that produces a plausible decryption. The security of the algorithm should solely depend on the secrecy of the key. The algorithm should not need to be secret.

If the attacker knows the plaintext corresponding to the ciphertext it's called "*known plaintext attack*" .

An attack where attacker doesn't know the plaintext , it's called "*ciphertext-only attack*". *Ex:* If the attacker knows that plaintext is ASCII , so any decryption which uses non-ascii characters must be using the wrong key.

1.1.1 Symmetric Encryption

When Sender and recipient share the same key(which must be kept secret) is referred to as *Symmetric Key Cryptograph* or also referred to as *Secret Key Cryptography* as opposed to *Public Key Cryptography* citation required?.

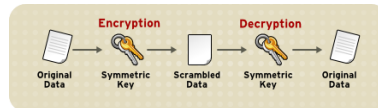


Figure 1: Symmetric Encryption [2]

1.1.2 Assymetric Encryption

- First started in stanford university by Whitfield Diffie and Martin Hellman
- The most commonly used implementation of PKC are based on algorithm based on algorithm patented by RSA data security.
- Each public key is published & private key is kept secret. Data encrypted with public key can be decrypted only with private key.
- In general, to send encrypted data to someone, we encrypt with public key & the person encrypt receiving the encrypted data decrypts with private key
- Compared to symmetric key encryption, public key encryption requires more computation & therefore not always appropriate for large amounts of data
- It is possible to use public-key encryption to send a symmetric key, which can then be used to encrypt additional data.

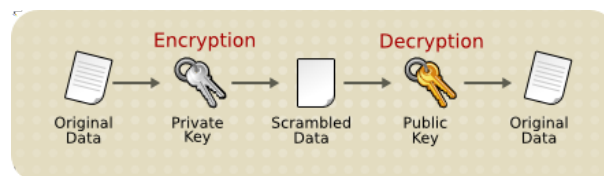


Figure 2: Assymmetric Encryption

- Reverse of the above figure also happens i.e. encrypt with private key and decrypt with public-key. But not useful for sensitive information
- There's a problem with the above method, how would the parties get each other's public-key ? If we send the keys through electronically, then the attacker can tamper, while they are in transit to the receiver.
- When the 2 parties want to communicate, the attacker can intercept the keys & instead send his own key to each other, thus each party encrypts to him and he re-encrypts it to the real recipient. This is called *man-in-the-middle attack*

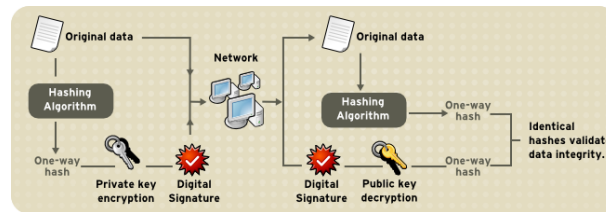


Figure 3: Digital Signature

1.2 Message Digest

A message digest is simply a function that takes as an input an arbitrary message and outputs a fixed length string which is characteristic of the message. The important property here is irreversibility. It's extremely difficult to compute a message from the given digest. Property of the message digest:

- For a digest to be secure, it must be difficult to generate any of the message that digests to the same value. You have to search a message space of proportional size of the digest in order to find a matching message text
- It should be difficult to produce 2 messages M and M' such that they have the same digest. This property is called collision-resistance. It turns out that the strength of any message digest against finding collision is only half the size of the digest., so a 128-bit digest is only 64 bits strong against collisions.

1.2.1 Message Authentication Code

Consider *Alice* and *Bob* share a key and *Alice* wants to send a message to *Bob*. The message can be encrypted, and send it across to *Bob*, but we are not sure that the encrypted message would be tampered and also not sure if *Alice* was the one who sent the encrypted message. So we use a new tool called *MAC*. *MAC* is a digest algorithm, but with a key, So the MAC is dependent on both the key and the message being MACed.

1.3 Algorithms

- **RSA**
 - RSA is the public-key algorithm widely used in Public-key Cryptography
 - Invented by Ron-Rivest, Adi Shamir and Len Adelman (RSA). Each user has a public-key and a private-key
 - The public-key can be freely distributed and the private-key must be kept secret.
 - Brief Explanation:
 - * Generate 2 prime numbers (3, 17), call them p and q
 - * $n = \text{modulus}(p * q)$
 - * p and q are kept secret

- * Difficulty is in factoring m , so that we get p and q
- * Take another number "e" which is called public exponent, which is a prime number, it's usually small prime numbers 3, 17 or 65537
- * Compute $d = e^{-1} \bmod ((p-1)(q-1))$
- * Public Key = (e, n)
- * Private Key = (d, n)
- * Message M is encrypted with public key and decrypted with private key
- * RSA assumes that M is a number, So we need a convention to convert strings in to numbers.

• DSA

- National Institute of standards and technology (NIST) published the Digital signature algorithm in the Digital signature standard.
- Compared to RSA where it can be used for both encryption and digital signature
- In DSA the signature generation is faster than signature verification

• RC4

- RC4 is a stream cipher.
- Assume we have a function (f) which produces 1 byte of data. This output is called keystream (ks)
- The function (f) takes a encryption key as an input to generate keystream
- Without the key we can't predict keystream
- So combine each byte with one byte of plain text which is our ciphertext
- $C[i]$ denotes the i th unit of ciphertext
- where a unit refers a 1 byte of data.
- $ks[i]$ refers the i th unit of keystream
- $M[i]$ refers to the i th unit of the message.
- $C[i] = ks[i] \text{ xor } M[i]$
- $m[i] = ks[i] \text{ xor } C[i]$
- Disadvantages:
 - * Assume we have used the same key to encrypt the Messages M and M'
 - * If the attacker learns M and can compute ks by simply computing $M \text{ xor } C$.
 - * Once the attacker knows KS he can generate M'
- RC4 was designed by Ron Rivest . RC4 is a variable key length cipher, with key can be anywhere between 8 to 2048 bytes long
- SSL/TLS protocol use RC4 with 128-bit (16 bytes) key length

- RC4 is extremely fast.

- **Block Ciphers**

- **DES**

- * The data to be encrypted is processed in blocks of bytes (8 or 16).
- * Each possible plain text block corresponds to a row in the table
- * So to encrypt a block you find a column corresponding to the key, run down the table to find the row corresponding to the block you want to encrypt.
- * If the data is huge, it's very difficult to manage, so we define a function that does the computation
- * The idea is to simulate a random table
- * we have a key(k), and a data block . and we define 2 function E for encryption ,and D for decryption
- * $C=E(K,M)$
- * $M=D(k,c)$
- * When we have large messages , we do in Electronic code book mode.
 - Break the message up to block sized-chunks.
 - individually encrypt it using encryption algorithm
 - $C[i]=E(k,M[i])$
 - $M[i]=D(K,C[i])$
- * DES was desinged by IBM , it's patented but freely available
- * DES is a 64 block cipher with 56-bit key . The data is encrypted in blocks of 8 bytes and a key space of 56 bits.
- * DES is used with public-key techniques.
- * Disadvantages:
 - if $M[i]$ and $M[j]$ are same, then we get the same $C[j]$ and $C[k]$. Attacker can get the pattern
 - Cipher block chaining mode, The encryption of each plain text block $M[i]$ depends on cipher text of previous block $C[i-1]$.
 - That can be accomplished by XORing previous Cipher text block $C[i-1]$ Xor $M[i]$ before encryption

- **3DES**

- * Run the DES algorithm 3 times.
- * It's used in Encrypt-Decrypt-Encrypt (EDE mode).
- * Message is encrypted with key-1, Decrypted with Key-2 and Encrypted with key-3.
- * 3DES is 3 times slower than DES.

- **RC2**

- * RC2 is a block cipher invented by Ron Rivest,
- * RC2 is a variable-length cipher , with variable length key. It uses 64 bit block size.

- **AES**

- * Advanced Encryption Standard uses a minimum of 128 bits and 3 key lengths, 128, 192 and 256 bits

- **Digest Algorithms**

- The two most popular Digest Algorithms are MD5, which was designed by Ron Rivest and SHA-1 by NIST.
 - MD5 and SHA share a common ancestor MD4 also designed by Ron Rivest

1.4 Protocols

1.4.1 SSL

- Secure Socket Layer (SSL) is a protocol that provides a secure channel between machines
- It has facilities for protecting data in transit and identifying machine with which it is begin communicated
- Protocol is transparent , so it can be run on any protocol.
- SSL has gone through many versions and currently is culminating with the adoption by IETF as Transport Layer Security
- SSL was originally designed for world wide web , but also was intended as a unifying solution to all the other communications (web, mail, news traffic)
- The current version of SSL is tlsv2 which fixes a lot of security problems

- **Overview of SSL Protocol**

- Primary goal of SSL is to provide privacy and reliability between 2 communicating applications
 - The protocol is composed of 2 layers : **Handshake protocol** and **record protocol**
 - In brief Handshake protocol provides
 - * allows server & client to authenticate to each other
 - * Negotiate encryption algorithm or Cryptography keys.
 - Record protocol provides:
 - * Confidentiality
 - * Authenticity
 - * replay protection

- **Handshake Protocol**

- Client sends a list of algorithms it's willing to support along with a random number used as input to a key generation proces
 - Server chooses out of that list and sends it back along with a certificate containing servers public-key.

- certificate also provides the server's identity for authentication purpose and the server supplies a random number which is used as a part of key generation process
- Client verifies the servers certificate and extracts the server's public-key from the certificate
- Client then generates a random secret called *pre_master_secret* and encrypts this with server's public-key
- Client sends this encrypted *pre_master_secret* to the server
- Client and the Server independently compute the encryption and MAC keys from *the pre_master_secret* and client and server's random values
- Client sends a MAC of all the handshake messages to the server
- The server sends a MAC of all the handshake messages to the client

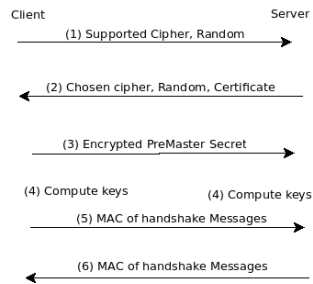


Figure 4: SSL Handshake

-
- Step:1 corresponds to single SSL Handshake message, Client Hello
- Step:2 corresponds to SSL Handshake messages
 - * First is a ServerHello Algorithm preference, Certificate
 - * Send ServerHelloDone
- Step3 corresponds to ClientKeyExchange
- Step 5 and 6 correspond to the finished message. The Finished message is the first message that's the protected using Just-Negotiated algorithms.
- To protect the handshake from tampering, the content of the message is MAC of all the previous handshake messages
-
- **Resuming Session**
 - * client sends a clientHello with a Session-ID of the session to be resumed.
 - * server check it's session cache for a match
 - * If the match is found & the server is willing to re-establish the connection under specified session state , it will send a Server-Hello with same session-ID

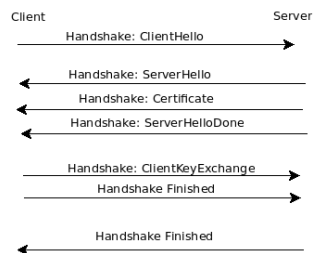


Figure 5: SSL Handshake

- * At this point both client and server must send change Cipher spec messages & proceed directly to finished messages
- * There is no exchange of certificates here.
- * Once the re-establishment is complete the client & server may begin exchange of application data
- * If the Session-ID match is not found the server generates a new session-ID & the client and Server have to perform a full handshake

• SSL Record Protocol

- The actual data transfer is accomplished by SSL Record Protocol
- Record Protocol works by breaking up the data stream to be transmitted in to a series of fragments, each of which is independently protected and transmitted.
- On the receiving end , each record is independently decrypted and verified.
- Before transmission , a MAC is computed on each record , This MAC is transferred along with the record
- The concatenated data and MAC are encrypted to form encrypted Payload. We attach a header to that encrypted payload which we refer to as a record

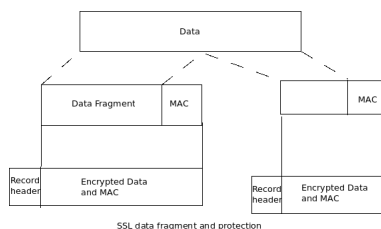


Figure 6: ssl data fragment and protection

-
- The record header provides the information for the other end to interpret the record. It contains:
 - * Content Type (Application data, alert messages, handshake message, change cipher spec)

- * length (how many bytes to read off wire)
- * SSL Version
- The *change_cipher_spec* message indicates a change in encryption and authentication of records
- Once the handshake is completed a new set of keys is negotiated, *change_cipher_spec* record is sent to indicate that those keys will now be used.

2 Introduction to Public Key Infrastructure

Below is a simplified architectural model of Public Key Infrastructure using X.509 (PKIX) Specifications

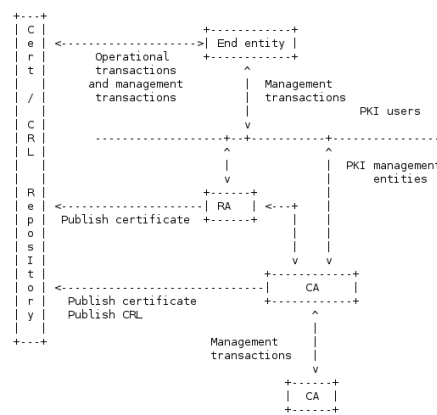


Figure 1 - PKI Entities

Figure 7: PKI Architectural Model

2.1 Common Terms used in PKI

- End Entity: User of the PKI certificates and/or end user system that is the subject of a certificate
- CA: Certificate Authority
- RA: Registration Authority, Optional System to which CA delegates certain management functions
- CRL issuers: A system that generates CRL
- repository: a system or collection of distributed systems that stores certificates and CRLs and services as a means of distributing these certificates and CRLs to end entities

2.2 Detailed look on certificate/CRL

2.2.1 Certificates

- Certificates are data structures that bind public key values to the subject. This binding is asserted by a trusted Certificate Authority.
- X.509 defines the standard certificate format and v3 is the latest version.
- Internet Privacy Enhanced Mail (PEM) RFC 1421 and 1422 also include specifications for PKI based on X.509 certs.
- Types of Certificates:
 - End-Entity
 - CA
- RFC 1422 defines hierarchical structure of CA's and there are three types of PEM CA
 - IPRA: Internet Policy Registration Authority, acts as Root of Certificate authority. IPRA operates under Internet Society Organization
 - PCA: Policy Certificate Authority, (Verisign, Digicert, etc) signed by IPRA
 - CA: Certificate Authorities signed by PCA (Organizational CA's)
- Policies used by CA:
 - IPRA certifies only PCA and not CA's or users cert
 - IPRA will make sure that the DN of the PCA is unique and will not certificate PCA's with similar DN
 - Certificates should not be issued to distinct entities under the same distinguished Name.
 - IPRA should not certify two PCA's with same DN
 - PCA's should not certify two CA's with same DN
 - CA's are expected to sign certificates only if the subject DN in the certificates is subordinate to the issuer CA DN.
- Types of Certificate Authorities
 - Cross-Certs: Where Issuer and Subject are different
 - Self-Issue: Where Issuer and Subject are same
 - * Self-Signed: Where key bound in to the certificate is same as the key used to sign the certificate

2.2.2 Basic Certificate Fields

- Version:
 - Describes the version of encoded certificate.
 - if extensions are used: **0x2(3)**
 - if extensions are not used but UniqueIdentifier is used: **0x1(2)**
 - if only basic fields are there: **0x0(1)**
 - Values: **0x2(3)**, **0x1(2)**, **0x0(1)**
- Serial Number:
 - Positive integer assigned by CA to each certificate
 - It must be unique for each Certificate given by CA
 - Can contain long integers (up to 20 Octets)
 - Values: Integers:
 - * **16694152257348400000**
 - * **0xe7ad8b07558a1727**
 - * **cd:ba:7f:56:f0:df:e4:bc:54:fe:22:ac:b3:72:aa:55**
- Signature:
 - Algorithm used by **CA** to sign the certificate
 - Value:
 - * **md2WithRSAEncryption**
 - * **sha1WithRSAEncryption**
- Issuer:
 - Identifies the entity that has signed and issued the certificate
 - **MUST** contain non-empty Distinguished Names
 - Names should confirm to X.501 standard
 - Generally contains Country, Organization, Common name, Serial-Number, province, State, title, Surname, Generation Qualifier(Jr, Sr).
 - Values:
 - * **C=US, O=VeriSign, Inc., OU=Class 1 Public Primary Certification Authority**
 - * **C=DE, ST=Bayern, L=Muenchen, O=Whatever it is, CN=IO::Socket::SSL Demo CA**
 - * **C=US, O=VeriSign, Inc., OU=Class 1 Public Primary Certification Authority**
- Validity:
 - The time interval during which the CA warrants that it will maintain status of the certificate

- Consists sequence of 2 dates
 - * Date on which certificate validity begins
 - * Date on which certificate validity ends
- Validity period of a certificate is period from notBefore to notAfter(inclusive)
- Values:
 - * **Not Before: Jan 29 00:00:00 1996 GMT**
 - * **Not After : Aug 1 23:59:59 2028 GMT**
 - * Special Note:
 - Devices are given certificates where there is no expiration date
 - Value:
 - Certificate is to be used for entire lifetime of the device **Not After: 99991231235959Z.(represented in Generalized Time)**
- Subject:
 - Identifies the entry associated with public key stored in the subject public key field
 - If the subject is CA then it should be populated with same data as issuer
 - Names should confirm to X.501
 - Values:
 - * **C=US, ST=North Carolina, O=Fedora Project, OU=Fedora User Cert, CN=mrniranjan/emailAddress=niranjan@ashoo.in**
- Subject Public key Info:
 - This field is used to carry public key and identify the algorithm with which the key is used (RSA, DSA)
 - Supported Cryptographic Algorithms:
 - * **Rivest-Shamir-Adelman (RSA)**
 - * **Digital Signature Algorithm (DSA)**
 - * **Diffie-Hellman (DH)**
 - * **Elliptic Curve Digital Signature Algorithm (ECDSA)**
 - * **Key Encryption Algorithm (KEA)**
 - * **Elliptic Curve Diffie-Hellman (ECDH)**

2.2.3 Extensions

- This field only appears in v3 Certs
- This contains sequence of one or more Certificate Extensions
- Extensions provides method for associating additional attributes with public keys
- Managing relationship between CA's

- Can carry private extensions to carry information unique to their community
- Each Extension is either Critical / Non Critical
- Each Extension includes OID and ASN.1 DER (OCTET)
 - Example: DE:65:01:16:19:2E:51:E0:9A:51:1A:37:50:94:7D:39:29:2A:42:2C
- There cannot be duplicates of Extensions
- Default CRITICAL value is false
- If the Certificate is CA , then they should have below Extensions
 - Basic Key Identifier
 - Authoritative key Identifier
 - Basic Constraints
- **Authority Key Identifier:**
 - Provides a means of identifying the public key corresponding to the private key used to sign the certificate
 - This is required if the issuer has multiple signing keys
 - If the CA certificate is self signed Authority key identifier is skipped
 - Authority key Identifier helps in identifying the issuer certificate.
 - Values:
 - * **keyid:48:E6:68:F9:2B:D2:B2:95:D7:47:D8:23:20:10:4F:33:98:90:9F:D4**
- **Subject Key Identifier:**
 - Provides a means of identifying certificates that contain a particular public key
 - This extension is must for CA certificates
 - The value placed in this is same as Authority Key Identifier
 - For End-entity Certificates, this extension provides a means for identifying certificates containing particular key used in application
 - Value:
 - * **48:E6:68:F9:2B:D2:B2:95:D7:47:D8:23:20:10:4F:33:98:90:9F:D4**
- **Key Usage:**
 - Defines the purpose of the key contained in the certificate
 - This extension is used to restrict the usage of the key to be used for purpose other than what is defined in Key Usage
 - * **digitalSignature**
 - Public key should be used only to verify signatures on objects other than Public key certificates/CRL
 - * **nonRepudiation/contentCommitment**

- subject Public key is used to verify digital signatures other than signatures on public key(CRL)
- used to provide nonRepudiation Service that protects against signing entity falsely denying any public action
- * **keyEncipherment**
 - Public key is used for enciphering private key or secret keys
 - ,
 - Is used for encrypting symmetric content-decryption key or an assymetric private key
- * **dataEncipherment**
 - is used for enciphering the raw data without the use any cipher (very rarely used)
- * **keyAgreement**
 - is used for key Agreement,
 - when used Deffie-Hellman key is to be used for key management
- * **keyCertSign**
 - Public key is used for verifying the signatures on public keys
 - , if this is true then
- * **cRLSign**
 - when the subject public key is used for verifying signatures of CRL
- * **encipherOnly**
 - subject's public key may be used only for enciphering data while performing key agreement
- * **decipherOnly**
 - subject's public key may be used only for deciphering data while performing key agreement
- * Important Note:

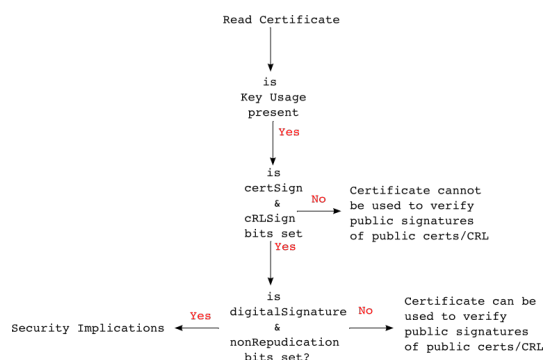


Figure 8: Key Usage restrictions

• Certificate Policies

- Contains sequence of one or more policy information terms each of which consists of OID, optional Queries
- OID Should not appear more than once
- for End-entity policy information terms indicate the policy under which the certificate has been issued and purposes for which the certificate must be used
- For CA the policy limits the policies for certificate paths that include this certificate
- if CA does not want to limit the set of policies for certification paths then it may assert **anyPolicy** with value **2.5.29.32.0**
- Policy Qualifiers:
 - * CPSNotice: Points URL to *certificate practice statement* document that describes the policy under which the subject was issued
 - * userNotice: text that describes the policy(not more than 200 characters)

- **Policy Mappings**

- This extension is used by CA certificates
- Lists one or more pairs of OID, which indicate that the corresponding policies of one CA are equivalent to policies of another CA [used in context of Cross pair certificates]
- **OID: 2.5.29.33**

- **Subject Alternative Name**

- Allows alternate names to be bound to subject of the certificate.
- The names specified in subjectAltName may be included in addition to or in place of the identity in the subject field of the certificate
- Defined subject names used:
 - * rfc822Name: IA5String
 - * otherName: OtherName
 - * dNSName: IA5String
 - * directoryName: Name
 - * URI: IA5string
 - * iPAddress: OCTET String
 - * UniformResourceIdentifier: IA5String

- **Issuer Alternative Name**

- Allows alternate names to be bound to certificate issuer
- Issuer alternative names are not processed as part of certification path validation
- Issuer Alternative name extension when present should be non-critical

- **Subject Directory Attributes**

- This extension is used to convey identification attributes of the subject
- Value:
 - * nationality of the subject

- **Basic Constraints**

- This extension identifies whether the subject of the certificate is CA ?
- It also identifies Maximum depth of valid certification paths that include this certificate
- Values:
 - * cA: (boolean): indicates whether certificate is CA cert
 - * pathLenConstraint: applicable if cA boolean is true. if present asserts keyCertSign bit.
 - * It gives maximum number of non-self-issued intermediate Certificates that may follow this certificate in a valid certification path. The last cert is End-entity certificate
 - * Ex: PathLength is 2:
 - RootCA - EE Cert
 - RootCA - subca1-EE Cert
 - RootCA - subca1-subca2-EE Cert
 - RootCA - subca1-subca2-subca3-EE Cert

- **Name Constraints**

- Used only in CA certificate
- Specifies name space within which all subject names in subsequent certificates in a certification path MUST be located.
- Restrictions apply to Subject Distinguished Name and subject alternative names.
- Restrictions do not apply for self-issued certs
- Restrictions are defined in terms of permitted or excluded name subtrees
- if name matches a restricted excluded subtree, it's invalid even though the name matches permitted subtrees.
- Examples:
 - * URI: constraint applies to host part of the name
 - * emailAddress: Example .example.org [indicates all emails with domain .example.com]
 - * DNS : pki.example.org [www.host1.pki.example.org would satisfy but host1.example.org will not]
 - * directoryName: Compares the DN attributes.

- **Policy Constraints**

- policy constraints extension can be used in certificates issued to CAs.

- Asserts policy related constraints
 - * inhibitPolicyMapping: if present policy mapping is to be inhibited while processing subsequent certificates
 - * requireExplicitPolicy: if present indicates subsequent certificates need to include an acceptable policy identifier

- **Extended Key Usage:**

- This extension is used to indicate one or more purposes for which certificates public key can be used:
 - * id-kp-serverAuth: TLS WWW server authentication
 - * id-kp-clientAuth: TLS WWW client authentication
 - * id-kp-codeSigning: Signing of downloadable executable code
 - * id-kp-emailProtection: Email protection
 - * id-kp-timeStamping: Binding the hash of an object to a time
 - * id-kp-OCSPSigning: Signing OCSP responses

- **CRL Distribution Points**

- cRLDistributionPoints extension is combination:
 - * distributionPoint
 - * reasons
 - * cRLIssuer

- **Authority Information Access**

- Indicates how to access information and services for the issuer of the certificate in which the extension appears
- Information and services include, On-line validation services, CA policy data
- This extension is added in both EE and CA cert

- **Subject Information Access**

- This extension indicates how to access information and services for the subject of the certificate in which the extension appears
- If the subject is CA, information and services may include certificate validation and CA policy data.
- If the subject is EE, information describes the type of services offered and how to access them

2.2.4 Revocation

When a certificate is issued, it is expected to be used for its entire validity period. Due to various circumstances, certificate can be invalidated like:

- Change of name
- Change of association with subject and CA (Employee left)
- Compromise or private key

- CA wants to revoke the certificate
- X.509 defines one method of revoking certificates, where CA periodically issues a signed data structure called Certificate revocation list (CRL).
- CRL is a time-stamped list identifying revoked certificates that is signed by CA or CRL issuer.
- This list is freely available through public repositories
- It is expected that the certificate system user not only verifies certificate validity, signature but also acquires latest CRL from public repositories and check against CRL
- CRL's are issued periodically (hourly, daily or weekly).
- CRLissuers or CA issue CRL
- CAs publish CRLs to provide status information about the certificates they issued
- CA may delegate this responsibility to another trusted authority
- Details of CRL
 - Each CRL ahas particular scope
 - CRL scope is the ste of certificates that could appear in a given CRL
 - Examples:
 - * all certificates issued by CA x
 - * all CA certificates that has been revoked by key compromise or CA compromise
 - * set of certificates base on arbitrary local information (all certificate issued to employees at location X)
 - CRL list all **unexpired** certificates, within it's scope that have been revoked for one or other reason
 - If the scope of the CRL includes one or more certificates issued by an entity other the CRL issuer it's called **indirect CRL**
 - CRL issuer may also generate delta CRL.
 - **Delta CRL** only lists those certificates whose revocation status has changed since the issuance of referenced complete CRL.
 - Referenced complete CRL is called **complete CRL**
 - Scope of the delta crl should be same as base CRL that it references
 - When CRL's are issued CRLs must be version 2 CRLs.

2.2.5 CRL Fields

- **Version** All CRLs must be Version 2 CRLs
- **Signature Algorithm**
 - Contains algorithm identifier for the algorithm used by the CRL issuer to sign the certificatesList
 - Values:
 - * **SHA256withRSA - 1.2.840.113549.1.1.11**
- **Issuer Name**
 - Issuer Name identifies the entity that has signed and issued the CRL
 - Alternative name forms may also appear in the issuerAltName Extension
 - The issuer must contain a non-empty X.500 DN.
- **This Update**
 - This field indicates the issue date of this CRL.
 - Example:
 - * **This Update: Sunday, January 17, 2016 8:06:03 AM IST Asia/Kolkata**
- **Next Update**
 - This field indicates the date by which the next CRL will be issued.
 - The next CRL may be issued before the indicated date, but it will not be issued later than indicated date.
 - CRL issuers(CA) **SHOULD** issue CRLs with nextUpdate time equal to or later than all previous CRLs.
- **Revoked Certificates**
 - Contains list of certificates revoked by CA
 - Certificates revoked by CA are uniquely identified by their certificate serial Number.
 - Date on which revocation occurred is specified.

2.2.6 CRL Extensions

- **Authority Key Identifier**
 - This extension provides a means of identifying the public key corresponding to the private key used to sign the CRL. This identifier can be based on either the key identifier or issuer's name and serial number.
 - Example:
 - * Authority Key Identifier Example:

```

Reason: Key_Compromise
Extensions:
  Identifier: Authority Key Identifier - 2.5.29.35
  Critical: no
  Key Identifier:
    DF:B4:B0:0D:98:E2:EC:44:82:24:10:C3:D4:ED:BE:14:
    68:B1:FE:35
  Identifier: CRL Number - 2.5.29.20
  Critical: no
  Number: 2

```

Figure 9: CRL Authority Key identifier

- **Issuer Alternative Name**

- This extension allows additional identities to be associated with the issuer or CRL.

- **CRL Number**

- This is a non critical extension that specifies sequence number for a given CRL scope and issuer.
- This extension allows users to determine if a particular CRL supersedes another CRL.
- if a CRL issuer generates delta CRLs in addition to complete CRLs for a given scope, the complete CRLs and delta CRLs must share one numbering sequence.

- **Delta CRL Indicator**

- Delta CRL indicator is a critical CRL extension that identifies a CRL being a delta CRL.
- Delta CRLs contain updates to the revocation information previously distributed, rather than all the information that would appear in complete CRL.
- This helps in reducing network load and processing time in certain environments
- Example:

```

Reason: Remove_from_CRL
Extensions:
  Identifier: Authority Key Identifier - 2.5.29.35
  Critical: no
  Key Identifier:
    DF:B4:B0:0D:98:E2:EC:44:82:24:10:C3:D4:ED:BE:14:
    68:B1:FE:35
  Identifier: CRL Number - 2.5.29.20
  Critical: no
  Number: 7
  Identifier: Delta CRL Indicator - 2.5.29.27
  Critical: no
  Base CRL Number: 6
Signature:
  Algorithm: SHA256withRSA - 1.2.840.113549.1.1.11
  Signature:

```

Figure 10: Delta CRL Indicator Extension

*

- **Issuing Distribution Point**

- This is a critical CRL extension that identifies the CRL distribution point and scope for a particular CRL.

- It indicates whether CRL covers revocation of EE Certificates only, CA certificates etc.
- Example:

```

Reason: Key_Compromise
Extensions:
  Identifier: Authority Key Identifier - 2.5.29.35
  Critical: no
  Key Identifier:
    DF:B4:B0:0D:98:E2:EC:44:B2:24:10:C3:D4:ED:BE:14:
    68:B1:FE:35
  Identifier: CRL Number - 2.5.29.20
  Critical: no
  Number: 9
  Identifier: Issuing Distribution Point - 2.5.29.28
  Critical: yes
  Distribution Point:
    Full Name:
      URIName: http://pki2.example.org/mycrl
  Only Contains User Certificates: no
  Only Contains CA Certificates: no
  Indirect CRL: no

```

Figure 11: Issuing Distribution Point extension

*

2.3 Exercises

1. Setup: Create PKI Infrastructure using OpenSSL

```

$ mkdir -p
  ~/pki/ca/{certs,crl,newcerts,private}
$ cd ~/pki/ca
$ chmod 700 private
$ touch index.txt
$ touch index.txt.attr
$ echo 01 > serial
$ wget
  https://raw.githubusercontent.com/mrniranjan/pkiworkshop/master/openssl.cnf

```

2. Create a CA cert using Openssl

```

$ cd ~/pki/ca
$ openssl genrsa -out private/ca.key.pem
  2048 -config openssl.cnf
$ chmod 400 private/ca.key.pem
$ openssl req -config openssl.cnf -key
  private/ca.key.pem -new -x509 \
  -days 365 -extensions v3_ca -out
  certs/ca.cert.pem

```

3. Create a End User cert with CN=Pki User1

```

# Create a private key

$ cd ~/pki/ca

$ openssl genrsa -out private/pkiuser1.key
  2048

```

```
# Create Certificate Request

$ openssl req -new -key
  private/pkiuser1.key \
  -out pkiuser1.csr -config openssl.cnf

$ openssl ca -config openssl.cnf
  -extensions usr_cert \
  -days 375 -in pkiuser1.csr -out
  certs/pkiuser1.cert.pem
```

4. Create a User cert *cn=Pki User1* with adding extension "Authority Key Identifier" where extension adds the hash of CA public key
Add the below lines to **usr_cert** section of openssl.cnf

```
basicConstraints=CA:FALSE
subjectKeyIdentifier=hash
authorityKeyIdentifier=keyid
```

5. Create a user cert by adding followin extensions digitalSignature, nonRepudiation
Add the below lines to **user_cert** section of openssl.cnf

```
keyUsage = nonRepudiation, digitalSignature
```

6. create a server cert by adding following extensions:
Add the below lines to **server_cert** section of openssl.cnf

```
keyUsage = keyAgreement, digitalSignature
```

7. Create a user cert with subjectAltname having an email address
Add the below lines to **user_cert** section of openssl.cnf

```
subjectAltName = email:copy
```

8. create a CA cert with pathlength of 2 . Verify path validation by creating multiple CAs
Create a new section **v3_usr_ca** in openssl.cnf and create certs

```
subjectKeyIdentifier=hash
authorityKeyIdentifier=keyid
basicConstraints = critical,CA:true,pathlen:2
keyUsage = cRLSign, keyCertSign
```

9. Create a server cert with TLS server and client authentication
Create a new section **server_cert** in openssl.cnf with below contents

```

[ server_cert ]
basicConstraints=CA:FALSE
subjectKeyIdentifier=hash
authorityKeyIdentifier=keyid
keyUsage = digitalSignature, keyAgreement
extendedKeyUsage = serverAuth,clientAuth

```

10. create a server cert with subject CN=server2.example.org having subjectAltName *.example.org

Add below line to **server_cert** section and create a new section **alt_names**

```

[ server_cert ]
subjectAltName = @alt_names

[alt_names]
DNS.1 = *.example.org

```

11. Create a Name constraint where permitted DN attribute is "CN=pki*,OU=PKI,O=Example Org,C=US"

Add below lines to **v3_usr_ca**

```

[ v3_usr_ca ]
subjectKeyIdentifier=hash
authorityKeyIdentifier=keyid
basicConstraints = critical,CA:true,pathlen:0
keyUsage = cRLSign, keyCertSign
nameConstraints = permitted;dirName:dir_sect

[dir_sect]
C=IN
O=Example Org
OU=PKI
CN=*

```

3 Common Terms Used in RHCS

- Enrollment The point when the certificate is issued to the User/client is called enrollment
- Renewal When certificate is about to be expired or already expired and has CA requires to renew the certificate
- Dual Keys Conventional hierarchical PKI uses a single key pair one public key and one private key. According to VeriSign, Key Pairs are used for one or more of three basic purposes:
 - Encryption
 - Authentication

- non-repudiation

A single key used for multiple purposes violates non-repudiation. It is recommended that you assign two key pairs per person – one key pair for signing messages providing authentication and non-repudiation, and a second key pair for encryption. This allows someone to recover the encryption key and decrypt documents that were encrypted using it, without their gaining the ability to sign documents with that users private key as well (which could lead to forgery and violation of non-repudiation)

- ECC is a cryptographic system that uses elliptic curves to create keys for encrypting data. ECC creates cryptographically-stronger keys with shorter key lengths than RSA, which makes it faster and more efficient to implement. But it is not widely supported as RSA
- CRMF Certificate Request Message Format is a syntax use to convey a request for a certificate to Certificate Authority (CA). This request contains a public key, Proof of Possession(POP), Additional registration informtion combined with pop and public key.[4]
- PKCS10
- FIPS
- FIPS-140
- Token
 - Internal Token
 - External Token
- HSM

4 Red Hat Certificate System

4.1 Overview

4.2 Certificate Manager

4.2.1 Introduction

Certificate Manager is the first subsystem that needs to be configured in PKI Environment, Certificate Manager can be configured as RootCA, Subordinate CA

4.2.2 Installation

- RPM: **pki-ca**
- configuration: CA subsystem is configured using utility **pkispawn**.
pkispawn provides both interactive configuration or silent configuration by reading a configuration file.

pkispawn first reads **default.cfg** first and gets other deployment specific information through interactive method or through batch mode by reading a file.

pkispawn then passes this information to a java servlet which performs the configuration.

```
[root@pk12 ~]# pkispawn
IMPORTANT:
  Interactive installation currently only exists for very basic deployments!
  For example, deployments intent upon using advanced features such as:
    * Cloning,
    * Elliptic Curve Cryptography (ECC),
    * External CA,
    * Hardware Security Module (HSM),
    * Subordinate CA,
    * etc.,
  must provide the necessary override parameters in a separate
  configuration file.
  Run 'man pkispawn' for details.
Subsystem (CA/KRA/OCSP/TKS/TPS) [CA]:
Tomcat:
Instance [pki-tomcat]:
HTTP port [8080]:
Secure HTTP port [8443]:
AJP port [8009]:
Management port [8005]:
Administrator:
Username [caadmin]:
Password:
Verify password:
Import certificate (Yes/No) [N]?
Export certificate to [/root/.dogtag/pki-tomcat/ca_admin.cert]:
Directory Server:
Hostname [pk12.example.org]:
Use a secure LDAPs connection (Yes/No/Quit) [N]?
LDAP Port [389]:
Bind DN [cn=Directory Manager]:
Password:
Base DN [o=pki-tomcat-CA]:
Security Domain:
Name [example.org Security Domain]:
Begin installation (Yes/No/Quit)? Yes
```

Figure 12: Configuring CA subsystem using pkispawn

4.2.3 Key Features

- CA subsystem issues, renews, revokes Certificates, generates Certificate Revocation lists
- Publishes Certificates/CRL in form of files or can publish to LDAP or OCSP responder
- CA also has an inbuilt OCSP responder enabling OCSP-Compliant clients to query CA about revocation status of Certificate
- Some CA's can delegate some of it's responsibility to another Subordinate CA

4.2.4 Architecture

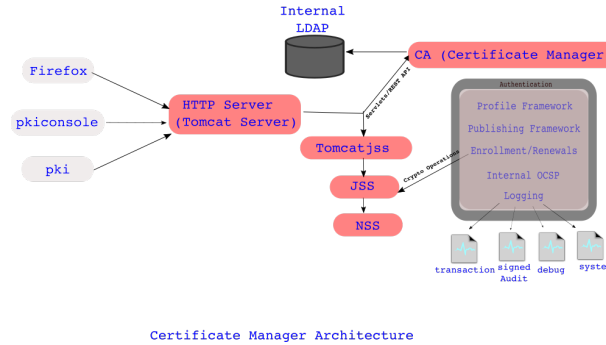


Figure 13: CA Subsystem Architecture

4.2.5 Interfaces

- End User Interface(Browser/CLI)
- Agent Interface(Browser/CLI)
- Admin interface(java console)

4.2.6 Features

- Enrollment:

End user Enrolls in the PKI infrastructure by submitting a Enrollment(certificate) request through End Entity Interface. This request can be submitted through 2 Methods:

- Browser
- CLI

There can be different kinds of Enrollment(Certificate) Request:

- Request for User, Server, SMIME, Dual Cert,.. certificate
- Request certificate if authentication through ldap, pin, cert etc.

Based on the above types, there are different certificate profiles associated with it. When end-entity(user) enrolls a certificate following events occur:

- The End-entity provides the information in one of the enrollment forms and submits a request
- The enrollment forms triggers the creation of public-key and private-key or dual-key pairs
- The End-entity provides authentication credentials before submitting the request, depending on the authentication type. This can be LDAP authentication, PIN-based authentication or certificate-based authentication.
- The request is submitted either to an agent-approved enrollment process or an automated process.

- * Agent-approved process requires no end-entity authentication, sends the request to the request queue in the agent-services interface.
 - * Automatic notification can be setup so an email can be sent to an agent any time a request appears in the queue
 - * The automated process, which involves end-entity authentication, process the certificate as soon as the end-entity successfully authenticates
- This form collects information about the end entity from the LDAP directory when the form is submitted
 - The profile associated with form determine the aspects of certificate that is issued. Depending upon the certificate profile the request is evaluated to determine if the request meets the constraints set.
 - The Certificate request is either rejected because it did not meet the certificate profile or authentication requirement, or a certificate is issued
 - The certificate is delivered to end-entity through HTML interface or email or certificate can be retrieved through Agents interface by serial number or request-ID.
 - The new certificate is stored in Certificate Managers internal database.

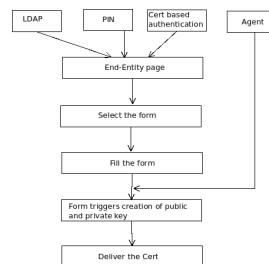


Figure 14: Certificate Enrollment

- Profiles:

Profile determine the content of a certificate. Certificate manager provides customizable framework to apply policies for incoming certificate requests and to control the input requests types and output certificate types

Certificate Profile define the following:

- Authentication Method
- Authorization Method
- Certificate content
- Constraints for the values of content
- Contents of input
- Output

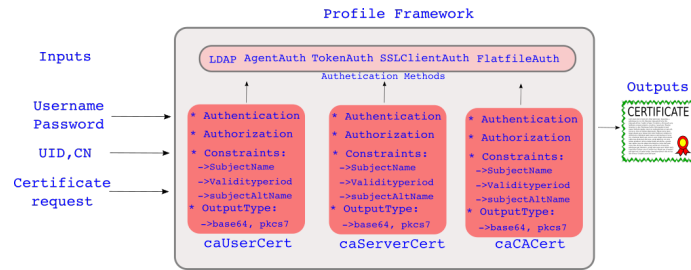


Figure 15: Certificate Profile Architecture

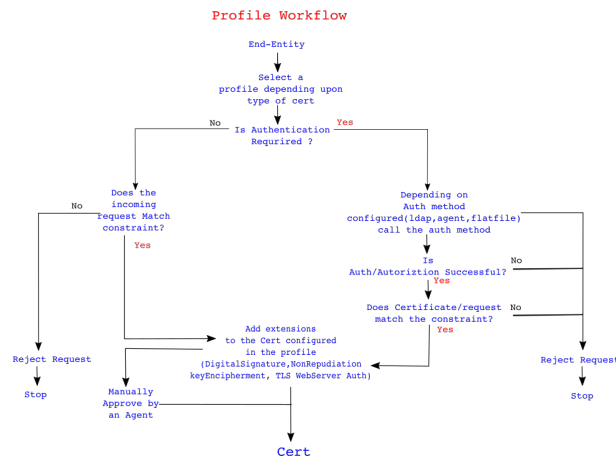


Figure 16: Certificate Profiles Workflow

Profile Workflow:

Each profile are defined in **.cfg** file located at **/var/lib/instance_name/profiles/ca** directory

Example **caUserCert** Profile:

The first part of the profile is the description and specifies whether profile is enabled or disabled, if enabled who enabled it.

```
desc=This certificate profile is for enrolling user certificates.
visible=true
enable=true
enableBy=admin
name=Manual User Dual-Use Certificate Enrollment
```

Figure 17: Profile Description

The second part of the profile describes the inputs:

```
auth.class_id=
input.list=1,12,13
input.11.class_id=keyGenInputImpl
input.12.class_id=subjectNameInputImpl
input.13.class_id=submitterInfoInputImpl
```

Figure 18: Profile inputs

- KeyGenInputImpl: specifies the key pair generation during the request submission, This provides if the request should be of type CRM-F/PKCS10, also provides dropdown specifying the key size.
- subjectNameInputImpl: specifies the subject Distinguished Name(DN) to be used in the cert. The subject DN can be constructed from *UID*, *Email*, *Common Name*, *Organizational Unit*, *Country*
- submitterInfoImpl: This input specifies three fields: *Requester Name*, *Requester email*, *Requester phone*

```
output.list=01
output.01.class_id=certOutputImpl
```

Figure 19: Profile output

Third part of the profile is output,

- certOutputImpl: The certificate output format *base64*, *pkcs7*, *prettyprint*

```
policyset.userCertSet.list=10,20,30,40,50,60
policyset.userCertSet.1.constraint.class_id=subjectNameConstraintImpl
policyset.userCertSet.1.constraint.name=Subject Name Constraint
policyset.userCertSet.1.constraint.params.pattern=UID.*
policyset.userCertSet.1.constraint.params.accept=true
policyset.userCertSet.1.default.class_id=userSubjectNameDefaultImpl
policyset.userCertSet.1.default.name=Subject Name Default
policyset.userCertSet.1.default.params.names
policyset.userCertSet.10.constraint.class_id=renewGracePeriodConstraintImpl
policyset.userCertSet.10.constraint.name=Renewal Grace Period Constraint
policyset.userCertSet.10.constraint.params.renewal.graceBefore=30
policyset.userCertSet.10.constraint.params.renewal.graceAfter=30
policyset.userCertSet.10.default.class_id=noDefaultImpl
policyset.userCertSet.10.default.name=no default
policyset.userCertSet.2.constraint.class_id=validityConstraintImpl
policyset.userCertSet.2.constraint.name=Validity Constraint
policyset.userCertSet.2.constraint.params.ranges=0
policyset.userCertSet.2.constraint.params.notBeforeCheck=false
policyset.userCertSet.2.constraint.params.notAfterCheck=false
policyset.userCertSet.2.default.class_id=validityDefaultImpl
policyset.userCertSet.2.default.name=Validity Default
policyset.userCertSet.2.default.params.ranges=0
policyset.userCertSet.2.default.params.starttime=0
policyset.userCertSet.3.constraint.class_id=keyConstraintImpl
policyset.userCertSet.3.constraint.name=key Constraint
policyset.userCertSet.3.constraint.params.keyTypes
policyset.userCertSet.3.constraint.params.keyParameters=1024,2048,3072,4096,nistp256,nistp384,nistp521
policyset.userCertSet.3.default.class_id=userKeyDefaultImpl
policyset.userCertSet.3.default.name=key default
policyset.userCertSet.4.constraint.class_id=noConstraintImpl
policyset.userCertSet.4.constraint.name=no Constraint
policyset.userCertSet.4.default.class_id=authorityKeyIdentifierDefaultImpl
policyset.userCertSet.4.default.name=Authority Key Identifier Default
policyset.userCertSet.5.constraint.class_id=noConstraintImpl
policyset.userCertSet.5.constraint.name=no Constraint
policyset.userCertSet.5.default.class_id=subjectInfoAccessDefaultImpl
policyset.userCertSet.5.default.name=Info Extension Default
```

Figure 20: Profile Policies

Last part of the profile is constraints, Policies like:

- validity of the cert
- renewal settings,
- key Usage Extensions
- User supplied extensions
- Publishing Certificate System provides customizable framework from CA's to publish.
 - Key Features
 - * Publish to a single repositories or multiple repositories
 - * Split locations by certificates/CRL
 - * Set individual rules for each type of certs/crl

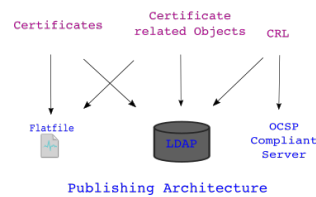


Figure 21: Publishing Architecture

- Publishing framework consists
 - * Publishers
 - * Mappers
 - * Rules
- Publishers: Publishers specify location to which certificates/CRL's are to be published. Example:
 - * To publish to a file, publishers specify the location of the publishing directory.
 - * To publish to LDAP, publishers specify the attribute in the directory that stores the cert/CRL.
 - * To publish to OCSP, we specify OCSP Server details.
- Rules: Rules define
 - * what is to be published and where ?
 - * What type of certs can be published to what location
 - * Set rules to publish certs to file and LDAP
 - * Set individual rules for each type of cert/rule
 - * There are rules for Files, LDAP and OCSP
- Mappers: Mappers are only used when publishing to LDAP
 - * Mappers construct the DN for an entry based on the information from the certificate or certificate request.
 - * Mappers use certificate or certificate request's subject name to construct the DN of the entry to which cert/certificate request/CRL has to be published

4.2.7 Instance Layout

4.2.8 Exercises

4.3 Key Recovery Authority

4.3.1 Introduction

Key Recovery Authority is an optional subsystem that is configured to archive private keys. Key Recovery Authority subsystem is by default installed after CA subsystem is installed and is dependent on CA for certain operations but can be installed as a standalone subsystem too.

4.3.2 Installation

- RPM: pki-kra
- Configuration: KRA subsystem can be configured through pkispawn using interactive or to have a customize setup a separate file can be specified with pkispawn.
- KRA subsystem can share the same tomcat instance created by the CA subsystem. This allows KRA subsystem to run on the same ports as that of CA.

```
root@pki2 ~]#
root@pki2 ~]# pkispawn

IMPORTANT:

Interactive installation currently only exists for very basic deployments!

For example, deployments intent upon using advanced features such as:

* Cloning,
* Elliptic Curve Cryptography (ECC),
* External CA,
* Hardware Security Module (HSM),
* Subordinate CA,
* etc.,

must provide the necessary override parameters in a separate
configuration file.

Run 'man pkispawn' for details.

Subsystem (CA/KRA/OCSP/TKS/TPS) [CA]: KRA

Tomcat:
Instance [pki-tomcat]: FoobarKRA
HTTP port [8080]: 12080
Secure HTTP port [8443]: 12443
AJP port [8009]: 12009
Management port [8005]: 12005

Administrator:
Username [kraadmin]:
Password:
Verify password:
Import certificate (Yes/No) [Y]? N
Export certificate to [/root/.dogtag/FoobarKRA/kra_admin.cert]:

Directory Server:
Hostname [pki2.example.org]:
Use a secure LDAP connection (Yes/No/Quit) [N]?
LDAP Port [389]: 1389
Bind DN [cn=Directory Manager]:
Password:
Base DN [o=FoobarKRA-KRA]:

Security Domain:
Hostname [pki2.example.org]:
Secure HTTP port [8443]:
Name: Foobar Org
Username [caadmin]:
Password:

Begin installation (Yes/No/Quit)? Yes
```

Figure 22: pkispawn KRA subsystem

4.3.3 Installation Layout

- **/var/lib/pki/<instance>**: <Instance> here is the tomcat Instance name given during configuring subsystem
- **/var/lib/pki/<instance>/alias**: Is the NSS Database of the system where the subsystem system certs are stored. When a subsystem is configured, Certain certificates are created for certain tasks like Transport Cert for securely transporting private keys, storage cert for storing the private key, These system certificates reside in alias directory. List of system certs are mentioned below, All the below system certs are issued by CA during configuration of KRA.

- transportCert
- Server-Cert
- auditSigningCert
- storageCert
- subsystemCert
- **/var/lib/pki/<instance>/conf/**: Configuration directory containing tomcat configuration, password file containing internal database(ldap) password, and NSS Database password
- **/var/lib/pki/<instance>/kra/**: Contains KRA Subsystem configuration file CS.cfg. This file is the main configuration file for KRA subsystem
- **/var/lib/pki/<instance>/lib/**: Contains kra subsystem related jar files
- **/var/lib/pki/<instance>/bin/**: Tomcat related jar files

4.3.4 Key Features

- Key Recovery Authority when configured with with certificate Manager(CA) stores private keys as part of certificate issuance process
- Supports two key recovery modes:
 - Synchronous Recovery(Deprecated)
 - Asynchronous Recovery
- Supports Archiving passwords, symmetric keys, aka acts as a vault to store anything that needs to be stored securely and transported back securely
- Private keys are generally wrapped with a key called storage key and private keys while transporting are wrapped with transport keys, These keys can be rotated

4.3.5 Architecture

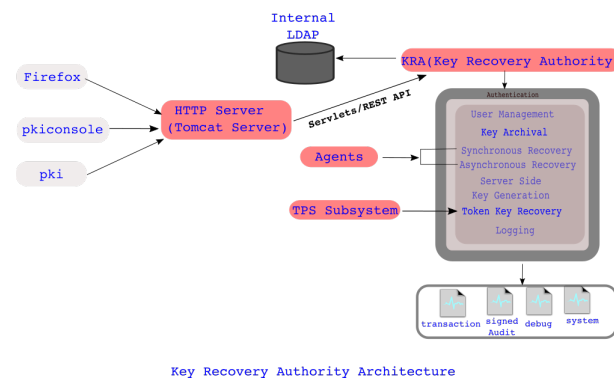


Figure 23: KRA Architecture

4.3.6 Interfaces

- End User Interface through **pki** cli
- Agent Interface(Browser/CLI)
- Admin Interface(Java console)

4.3.7 Features

- Key Archival
 - key archival is triggered when a user creates a certificate request through firefox(31.0 or older only)/pki/CRMFPopClient utility
 - The key pair generated as part of the request(public/Private key) is submitted through a defined profile.
 - While submitting the private key is wrapped(encrypted) with transport cert(public key) and submitted to CA
 - CA inturn retrieves the encrypted private key, decrypts with transport private key
 - Private key is then stored in interan ldap of kRA by encrypting again with Storage Cert
 - CA upon successful archival issues certificate

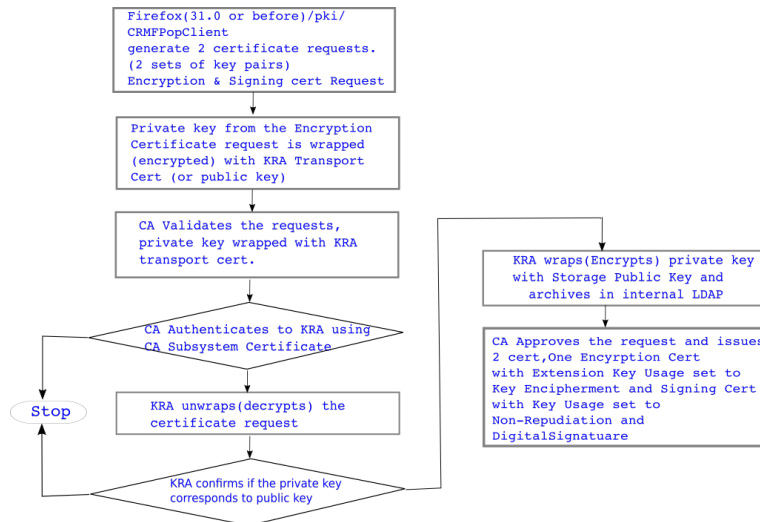


Figure 24: Key Archival Workflow

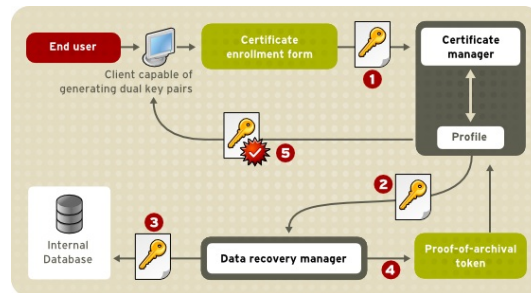


Figure 25: Key Archival Process

- Key Recovery
 - KRA supports Agent initiated key recovery
 - Designated Agents use the key recovery form on Agent services page to Approve/Reject key recovery
 - CS uses m-of-n ACL-based recovery scheme
 - CS uses it's ACL to ensure recovery agents are appropriately authenticated over SSL
 - ACL also requires that agents belong to specific recovery agent group
 - Recovery request is executed only when m of n agents have granted authorization

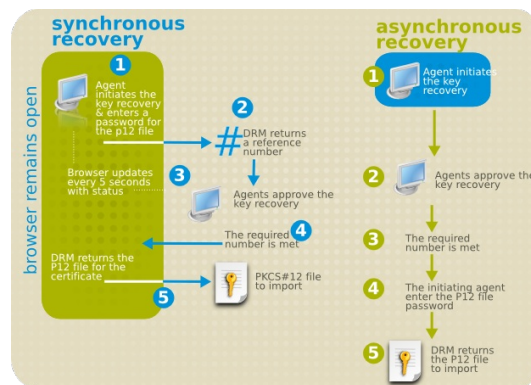


Figure 26: Key Recovery Process

- Archive Symmetric Keys, passwords, Asymmetric keys (Vault)
 - KRA has new REST interface to archive and recover
 - * Symmetric Keys
 - * Random Security Data
 - * Asymmetric Keys
 - This features is implemented through pki key CLI to generate, archive, retrieve keys.

- Transport Key Rotation
 - Transport Certs are created when KRA instance is configured using pkispawn
 - Transport Certs are issued by CA to KRA
 - Transport Certs public key is used to *encrypt* private key while it's being transport by CA to KRA
 - In Large organizations, this transport key needs to be rotated.
 - The feature here adds support for second Transport Key

4.3.8 Exercises

4.4 Online Certificate Status Protocol

4.4.1 Introduction

- Online Certificate Status Manager is an subsystem external to CA.
- This subsystem acts as an external OCSP which can be publicly accessible.
- Multiple CA's can publish CRL to a single OCSP
- Clients can verify certs using OCSP

4.4.2 Installation

- RPM: **pki-ocsp** OCSP subsystem can be configured through pkispawn using interactive or to have a customize setup a separate file can be specified with pkispawn

OCSP subsystem can share the same tomcat instance created by the CA subsystem which allows OCSP to use same ports as that of CA.

```

[root@pk12 ~]# pkispawn
IMPORTANT:

Interactive installation currently only exists for very basic deployments!

For example, deployments intent upon using advanced features such as:

    * Cloning,
    * Elliptic Curve Cryptography (ECC),
    * External CA,
    * Hardware Security Module (HSM),
    * Subordinate CA,
    * etc.,

must provide the necessary override parameters in a separate
configuration file.

Run 'man pkispawn' for details.
Subsystem (CA/KRA/OCSP/TKS/TPS) [CA]: OCSP
Tomcat:
Instance [pki-tomcat]: FooBarOCSP
HTTP port [8080]: 14080
Secure HTTP port [8443]: 14443
AJP port [8009]: 14009
Management port [8005]: 14005
Administrator:
Username [ocspadmin]:
Password:
Verify password:
Import certificate (Yes/No) [Y]? N
Export certificate to [/root/.dogtag/FooBarOCSP/ocsp_admin.cert]:
Directory Server:
Hostname [pk12.example.org]:
Use a secure LDAPS connection (Yes/No/Quit) [N]?
LDAP Port [389]: 1389
Bind DN [cn=Directory Manager]:
Password:
Base DN [o=FooBarOCSP-OCSP]:
Security Domain:
Hostname [pk12.example.org]:
Secure HTTP port [8443]:
Name: FooBar Org
Username [caadmin]:
Password:
Begin installation (Yes/No/Quit)? Yes

```

Figure 27: pkispawn KRA subsystem

4.4.3 Interfaces

- End User Interface through **pki cli**
- Agent Interface(Browser/CLI)
- Admin Interface(Java console)

4.4.4 Exercises

4.5 SmartCard Management

4.5.1 Introduction

- **What:** Smartcard contains an embedded microprocessor. This microprocess runs Java Virtual Machine(JVM) and provides environment to run applications on top of it.
 - Smartcard contains a chip with java on it.
 - it contains a JVM, Java Runtime Environment, and some API for applications to access the card.
- **Why:**
 - Smartcards are typically used by enterprise to provide identification and authentication

- Smartcards provide 2-Factor Authentication which is
 - * Something you have - smartcard
 - * Something you know - pin(password) to access the smartcard
- **Use cases:**
 - System Authentication using Smartcard (Linux/Windows System Authentication)
 - Secure VPN Access
 - Email Signature and Encryption
 - Disc Encryption
- **How:**
 - For enterprise to use smartcard a digital Certificate is issued to the smartcard
 - This certificate is used to identify the user of the smartcard
 - User uses digital certificate on the smartcard to perform 2-factor authentication
- **Mangement:**
 - We need a management software to manage certificates on the smart-card
 - * Enrolling certificate
 - * Formatting Smartcard
 - * What to do with the certificate if smartcard is lost
 - * What to do if the user of the smartcard is no longer in the company
 - To do the above operations, we have Token Management system in Certificate Services which comprises of the following:
 - * **TPS:** The Token Processing System (TPS) interacts with smart cards to help them generate and store keys and certificates for a specific entity, such as a user or device.
 - * **TKS:** The Token Key Service (TKS) generates, or derives, symmetric keys used for communication between the TPS and smart card.
 - * **CA:** The Certificate Authority (CA) creates and revokes user certificates stored on the smart card.
 - * **KRA:** Optionally, the Key Recovery Authority (KRA) archives and recovers keys for the smart card.
 - * **Enterprise Security Client:** is the conduit through which TPS communicates with each token over a secure HTTP channel (HTTPS), and, through the TPS, with the Certificate System.

4.5.2 Architecture

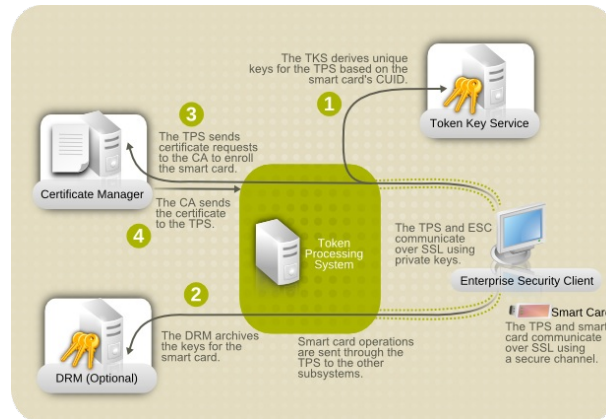


Figure 28: Token Management System

4.5.3 Common Terms Used with regard to Smartcard

- Secure Channel Protocol: A secure communication protocol and set of security protocol used to create a secure channel with token during TPS Operations.
- Applets: TPS communicates with an applet on the smartcard.
- APDU(Application Protocol Data Unit): Standard communication messaging protocol between a card accepting device(Reader) and a smart card
- SmartCards come with a cardmanager applet and a vendor applet or with only card manager applet.
- Coolkey: applet is the supported Card Manager applet
- Smartcard has cryptographic functionality supporting symmetric and Asymmetric Encryption(like RSA)
- CCID: CCID (chip card interface device) protocol is a USB protocol that allows a smartcard to be connected to a computer via a card reader using a standard USB interface, without the need for each manufacturer of smartcards to provide its own reader or protocol. [5]

4.5.4 Token Key Service

- Introduction
 - A subsystem in the token management system which derives specific, separate keys for every smart card based on the smart card APDUs and other shared information, like the token CUID. [1]
 - The TKS derives keys based on the token CCID, private information, and a defined algorithm. These derived keys are used by the TPS to format tokens and enroll, or process, certificates on the token(smartcard). [1]

- TKS and TPS have a sharedSecret to wrap sensitive info sent between TKS and TPS
- Features:
 - **Manage shared keys**
 - * Token Key Service (TKS) derives keys for the TPS to use.
 - * Based on Token CUID, and other token material , Derive keys that encrypt session between TPS and ESC
 - * Keys are derived based on Common Master Key which is known to TKS and existent on smartcards
 - **Generate Master Key**
 - * Master key is 3DES symmetric Key stored in software/Hardware
 - * Transport Key wraps the other keys derived(from master key) and used by TKS and TPS

4.5.5 Token Processing Service

- Introduction
 - TPS interacts with smart cards to help them generate and store keys and certificates for a specific entity
 - Smart card operations go through TPS and are forwarded to the appropriate subsystem for action, such as the Certificate Authority to generate certificates or the Key Recovery Authority to archive and recover key
- Features
 - Secure Channel
 - * The Token Key Service (TKS) generates the (token) keys (a set of 3 keys derived) from Master keys
 - * These keys are used by TPS to communicate with ESC(smartcard)
 - * TPS communicates with the TKS over SSL but TPS communicates with ESC over secure channel
 - * Every smartcard has 3 keys:
 - An **auth key** which is used for encryption and authentication: a session key dervied from the auth key each time a secure channel is opened.
 - A **MAC key** which is used for message authentication; like the auth keys, a session key is derived from the MAC key each time a secure channel is opened.
 - A **key encryption key (KEK)** which is to encrypt the session keys.

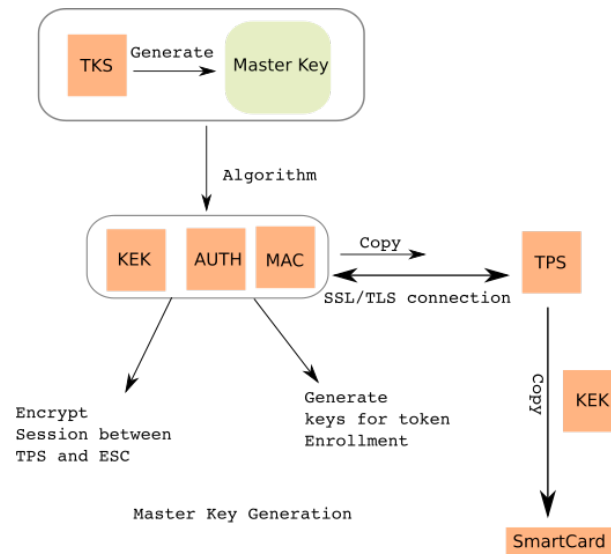
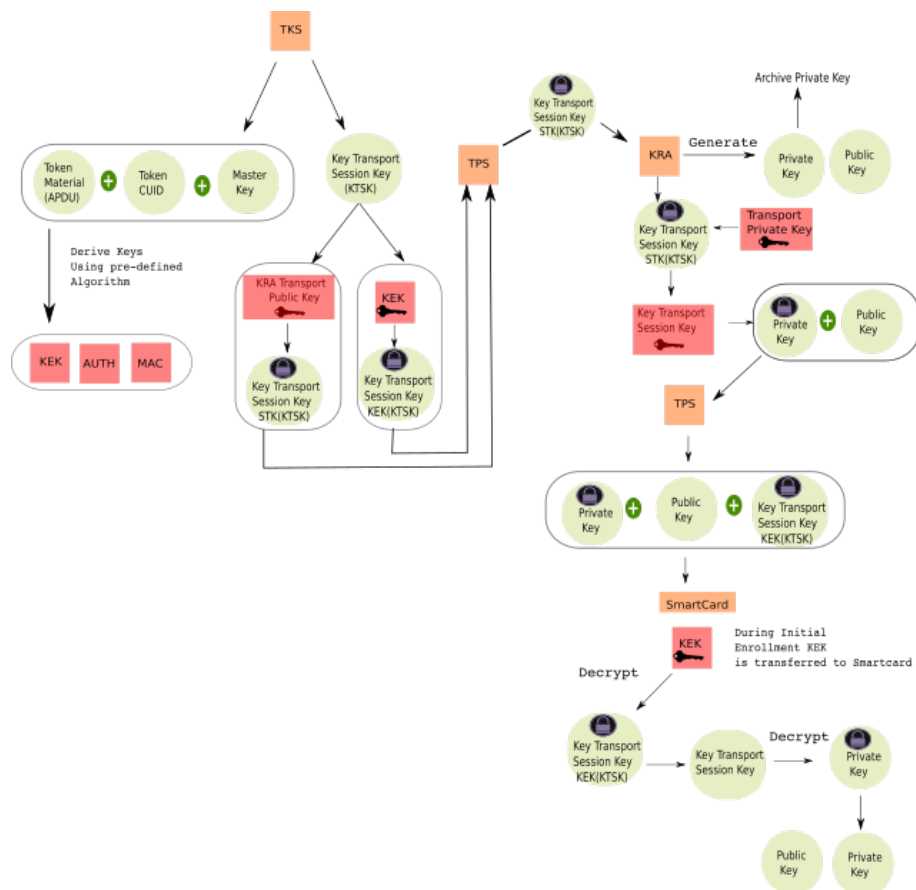


Figure 29: TKS Key Derivation

- Supported TPS Operations
 - * Formatting Smartcards
 - * Resetting PIN on smartcards
 - * Upgrading the applet for smartcard tokens
 - * Perform LDAP Authentication
 - * Managing token database
 - * Logging Token events
- Token Profiles
 - * Just like certificate profiles there are different token profiles to format different kind of tokens
 - * Token profiles define the below:
 - The steps to format and enroll the token
 - The configuration of the final enrolled token
 - Profiles also specify which LDAP server to be used for authentication and which CA should be used for enrolling certificates
 - List of Token profiles:

Token Type	Description
cleanToken	For operations for any blank token, without any other applied token types.
soKey	For operations for generating keys for security officer stations.
soCleanSOToken	For operations for blank tokens for security officer stations.
soKeyTemporary	For operations for temporary security officer tokens.
soCleanUserToken	For operations for blank user tokens for security officers.
soUserKey	For operations for security officer user tokens.
tokenKey	For operations for generating keys for uses with servers or devices.
userKey	For operations for regular user tokens.
userKeyTemporary	For operations for temporary user tokens.

- TPS-TKS-Smartcard Workflow



- TKS Configuration

- 43

be specified with pkispawn

- * TKS subsystem can share the same tomcat instance created by the CA subsystem which allows TKS to use same ports as that of CA

```
[root@pk12 ~]#
[root@pk12 ~]# pkispawn

IMPORTANT:

Interactive installation currently only exists for very basic deployments!

For example, deployments intent upon using advanced features such as:

* Cloning,
* Elliptic Curve Cryptography (ECC),
* External CA,
* Hardware Security Module (HSM),
* Subordinate CA,
* etc..

must provide the necessary override parameters in a separate
configuration file.

Run 'man pkispawn' for details.

Subsystem (CA/KRA/OCSP/TKS/TPS) [CA]: TKS

Tomcat:
Instance [pk1-tomcat]: FoobarTKS
HTTP port [8000]: 16000
Secure HTTP port [8443]: 16443
AJP port [8009]: 16009
Management port [8005]: 16005

Administrator:
Username [tkadmin]:
Password:
Verify password:
Import certificate (Yes/No) [Y]? N
Export certificate to [/root/.dogtag/FoobarTKS/tks_admin.cert]:

Directory Server:
Hostname [pk12.example.org]:
Use a secure LDAPS connection (Yes/No/Quit) [N]?
LDAP Port [389]: 1389
Bind DN [cn=Directory Manager]:
Password:
Base DN [o=FoobarTKS-TKS]:

Security Domain:
Hostname [pk12.example.org]:
Secure HTTP port [8443]:
Name: Foobar Org
Username [caadmin]:
Password:

Begin installation (Yes/No/Quit)? Yes

Log file: /var/log/pki/pki-tks-spawn,20160125182959.log
Installing TKS into /var/lib/pki/FoobarTKS.
Storing deployment configuration into /etc/sysconfig/pki/tomcat/FoobarTKS/tks/deployment.cfg.
```

Figure 32: Interactive installation of TKS using pkispawn

– Creating sharedSecret

- * TKS generates a sharedSecret to wrap sensitive information passed over TPS.
- * This sharedSecret is to be created on TKS NSS database using tool called **tkstool**

```
[root@pk12 conf]#
[root@pk12 conf]#
[root@pk12 conf]# cat password.conf
internal=26642654353
replicationdb=1526134020
[root@pk12 conf]# cd alias/
[root@pk12 alias]# tkstool -T -d /var/lib/pki/FoobarTKS/alias/ -n sharedSecret
Enter Password or Pin for "NSS Certificate DB":
```

Figure 33: NSS Internal database password

```
Generating the first session key share . . .

first session key share:  73FD 68D0 0D08 E6F1
                        B949 9BD9 EF64 D369

first session key share KCV: 1392 0088

(1) Write down and save the value for this first session key share.
(2) Write down and save the KCV value for this first session key share.

Type the word "proceed" and press enter to continue (or ^C to break): █
```

Figure 34: tkstool Usage

```

Generating the second session key share . . .

second session key share:  FB0D 5888 A073 9857
                          4F72 C8D0 E8D4 349F

second session key share KCV:  827F 9805

(1) Write down and save the value for this second session key share.
(2) Write down and save the KCV value for this second session key share.

Type the word "proceed" and press enter to continue (or ^C to break): █

```

Figure 35: First Session Key

```

Generating the third session key share . . .

third session key share:  FE31 7E30 891C D64C
                          D92A 2B94 15F4 917F

third session key share KCV:  7C0B AB95

(1) Write down and save the value for this third session key share.
(2) Write down and save the KCV value for this third session key share.

Type the word "proceed" and press enter to continue (or ^C to break): █

```

Figure 36: Second Session Key

```

Generating first symmetric key . . .
Generating second symmetric key . . .
Generating third symmetric key . . .
Extracting transport key from operational token . . .

transport key KCV:  6F68 C5A0

Storing transport key on final specified token . . .
Naming transport key "sharedSecret" . . .
Successfully generated, stored, and named the transport key!
[root@pki2 alias]# █

```

Figure 37: Third Session Key

- * tkstool utility prints out the key shares and KCV values for each of the three session keys that are generated. These are necessary to import the **sharedSecret** to TPS
- TPS Configuration
 - * RPM: **pki-tps**
 - * Configuration: TPS subsystem can be configured through pkispawn using interactive or to have a customize setup a separate file can be specified with pkispawn
 - * TPS subsystem can share the same tomcat instance created by the CA subsystem which allows TPS to use same ports as that of CA

```

[root@pki2 ~]#
[root@pki2 ~]# pkispawn

IMPORTANT:

Interactive installation currently only exists for very basic deployments!

For example, deployments intent upon using advanced features such as:

    * Cloning,
    * Elliptic Curve Cryptography (ECC),
    * External CA,
    * Hardware Security Module (HSM),
    * Subordinate CA,
    * etc.,

must provide the necessary override parameters in a separate
configuration file.

Run 'man pkispawn' for details.

Subsystem (CA/KRA/OCSP/TKS/TPS) [CA]: TPS
Tomcat:
Instance [pki-tomcat]: FoobarTPS
HTTP port [8080]: 18080
Secure HTTP port [8443]: 18443
AJP port [8009]: 18009
Management port [8005]: 18005
Administrator:
Username [tpsadmin]:
Password:
Verify password:
Import certificate (Yes/No) [Y]? N
Export certificate to [/root/.dogtag/FoobarTPS/tps_admin.cert]:
Directory Server:
Hostname [pki2.example.org]:
Use a secure LDAP connection (Yes/No/Quit) [N]?
LDAP Port [389]: 1389
Bind DN [cn=Directory Manager]:
Password:
Base DN [o=FoobarTPS-TPS]:
Security Domain:
Hostname [pki2.example.org]:
Secure HTTP port [8443]:
Name: Foobar Org
Username [caadmin]:
Password:
External Servers:
CA URL [https://pki2.example.org:18443]: https://pki2.example.org:8443
TKS URL [https://pki2.example.org:18443]: https://pki2.example.org:16443
Enable server side key generation (Yes/No) [N]?
Authentication Database:
Hostname [pki2.example.org]:
Port [389]: 1389
Base DN: dc=example,dc=org

```

Figure 38: Interactive installation of TKS using pkispawn

```

Begin installation (Yes/No/Quit)? Yes
Log file: /var/log/pki/pki-tps-spawn.20160128044654.log
Installing TPS into /var/lib/pki/FoobarTPS.
Storing deployment configuration into /etc/sysconfig/pki/tomcat/FoobarTPS/tps/deployment.cfg.
Notice: Trust flag u is set automatically if the private key is present.

=====
INSTALLATION SUMMARY
=====

Administrator's username:          tpsadmin
Administrator's PKCS #12 file:
    /root/.dogtag/FoobarTPS/tps_admin_cert.p12

To check the status of the subsystem:
    systemctl status pki-tomcatd@FoobarTPS.service

To restart the subsystem:
    systemctl restart pki-tomcatd@FoobarTPS.service

The URL for the subsystem is:
    https://pki2.example.org:18443/tps

PKI instances will be enabled upon system boot
=====

```

Figure 39: Continuation of pkispawn

- * Import SharedSecret
 - Import the sharedSecret created by TKS so that TKS and TPS can wrap sensitive information
 - sharedSecret needs to be imported in to TPS NSS Database

```

[root@pki2 conf]# pwd
/var/lib/pki/FoobarTPS/conf
[root@pki2 conf]# cat password.conf
internal=548782292471
internaldb=Secret123
replicationdb=-352748034
[root@pki2 conf]# tkstool -I -d /var/lib/pki/FoobarTPS/alias/ -n sharedSecret
Enter Password or Pin for "NSS Certificate DB":

```

Figure 40: Get NSS Password of TPS NSS database

```

Enter the first session key share . . .

Type in the first session key share (or ^C to break):

    first session key share:      73FD 68D0 0D08 E6F1
                                B949 9BD9 EF64 D3B9

Type in the corresponding KCV for the first session key share (or ^C to break):

    first session key share KCV:  1392 8DB8

Verifying that this session key share and KCV correspond to each other . . .

Congratulations, the first session key share KCV value entered CORRESPONDS
to the first session key share value entered!

Type the word "proceed" and press enter to continue (or ^C to break):

```

Figure 41: tkstool to import first session key

```

Enter the second session key share . . .

Type in the second session key share (or ^C to break):

    second session key share:    FB0D 5868 AD73 3857
                                5EF2 CED0 E3DC 34DF

Type in the corresponding KCV for the second session key share (or ^C to break):

    second session key share KCV: 827F 9805

Verifying that this session key share and KCV correspond to each other . . .

Congratulations, the second session key share KCV value entered CORRESPONDS
to the second session key share value entered!

Type the word "proceed" and press enter to continue (or ^C to break):

```

Figure 42: tkstool to import second session key

```

Enter the third session key share . . .

Type in the third session key share (or ^C to break):

    third session key share:     FE31 7FD0 831C D64C
                                D92A 3B54 15F4 917F

Type in the corresponding KCV for the third session key share (or ^C to break):

    third session key share KCV: 7C0B AB35

Verifying that this session key share and KCV correspond to each other . . .

Congratulations, the third session key share KCV value entered CORRESPONDS
to the third session key share value entered!

Type the word "proceed" and press enter to continue (or ^C to break): proceed

```

Figure 43: tkstool to import third session key

* TKS Instance Layout

References

- [1] Red Hat. Red hat certificate services 8.0 admin guide, 2015.
- [2] Red Hat. Symmetric encryption, 2015.
- [3] Eric Rescorla. *SSL and TLS Designing and Building Secure System*. Addison-Wesely, 2001.
- [4] J. Schaad. Internet X.509 Public Key Infrastructure Certificate Request Message Format (CRMF). RFC 4211, RFC Editor, September 2005.
- [5] Wikipedia. Ccid (protocol), 2015.