



## C interfaces to GALAHAD LPB

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Mon May 1 2023



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# Chapter 1

## GALAHAD C package lpb

### 1.1 Introduction

#### 1.1.1 Purpose

This package uses a primal-dual interior-point method to solve the **linear programming problem**

$$\text{minimize } q(x) = g^T x + f$$

subject to the general linear constraints

$$c_i^l \leq a_i^T x \leq c_i^u, \quad i = 1, \dots, m,$$

and the simple bound constraints

$$x_j^l \leq x_j \leq x_j^u, \quad j = 1, \dots, n,$$

where the vectors  $g$ ,  $w$ ,  $x^0$ ,  $a_i$ ,  $c^l$ ,  $c^u$ ,  $x^l$ ,  $x^u$  and the scalar  $f$  are given. Any of the constraint bounds  $c_i^l$ ,  $c_i^u$ ,  $x_j^l$  and  $x_j^u$  may be infinite. Full advantage is taken of any zero coefficients in the matrix  $A$  whose rows are the transposes of the vectors  $a_i$ .

#### 1.1.2 Authors

N. I. M. Gould, STFC-Rutherford Appleton Laboratory, England.

C interface, additionally J. Fowkes, STFC-Rutherford Appleton Laboratory.

Julia interface, additionally A. Montoison and D. Orban, Polytechnique Montréal.

#### 1.1.3 Originally released

August 2018, C interface September 2021.

### 1.1.4 Terminology

The required solution  $x$  necessarily satisfies the primal optimality conditions

$$(1a) \quad Ax = c$$

and

$$(1b) \quad c^l \leq c \leq c^u, \quad x^l \leq x \leq x^u,$$

the dual optimality conditions

$$(2a) \quad g = A^T y + z$$

where

$$(2b) \quad y = y^l + y^u, \quad z = z^l + z^u, \quad y^l \geq 0, \quad y^u \leq 0, \quad z^l \geq 0 \quad \text{and} \quad z^u \leq 0,$$

and the complementary slackness conditions

$$(3) \quad (Ax - c^l)^T y^l = 0, \quad (Ax - c^u)^T y^u = 0, \quad (x - x^l)^T z^l = 0 \quad \text{and} \quad (x - x^u)^T z^u = 0,$$

where the vectors  $y$  and  $z$  are known as the Lagrange multipliers for the general linear constraints, and the dual variables for the bounds, respectively, and where the vector inequalities hold component-wise.

### 1.1.5 Method

Primal-dual interior point methods iterate towards a point that satisfies these conditions by ultimately aiming to satisfy (1a), (2a) and (3), while ensuring that (1b) and (2b) are satisfied as strict inequalities at each stage. Appropriate norms of the amounts by which (1a), (2a) and (3) fail to be satisfied are known as the primal and dual infeasibility, and the violation of complementary slackness, respectively. The fact that (1b) and (2b) are satisfied as strict inequalities gives such methods their other title, namely interior-point methods.

The method aims at each stage to reduce the overall violation of (1a), (2a) and (3), rather than reducing each of the terms individually. Given an estimate  $v = (x, c, y, y^l, y^u, z, z^l, z^u)$  of the primal-dual variables, a correction  $\Delta v = \Delta(x, c, y, y^l, y^u, z, z^l, z^u)$  is obtained by solving a suitable linear system of Newton equations for the nonlinear systems (1a), (2a) and a parameterized "residual trajectory" perturbation of (3); residual trajectories proposed by Zhang (1994) and Zhao and Sun (1999) are possibilities. An improved estimate  $v + \alpha \Delta v$  is then used, where the step-size  $\alpha$  is chosen as close to 1.0 as possible while ensuring both that (1b) and (2b) continue to hold and that the individual components which make up the complementary slackness (3) do not deviate too significantly from their average value. The parameter that controls the perturbation of (3) is ultimately driven to zero.

The Newton equations are solved by applying the GALAHAD matrix factorization package SBLS, but there are options to factorize the matrix as a whole (the so-called "augmented system" approach), to perform a block elimination first (the "Schur-complement" approach), or to let the method itself decide which of the two previous options is more appropriate. The "Schur-complement" approach is usually to be preferred when all the weights are nonzero or when every variable is bounded (at least one side), but may be inefficient if any of the columns of  $A$  is too dense.

Optionally, the problem may be pre-processed temporarily to eliminate dependent constraints using the GALAHAD package FDC. This may improve the performance of the subsequent iteration.

### 1.1.6 Reference

The basic algorithm is a generalisation of those of

Y. Zhang (1994), On the convergence of a class of infeasible interior-point methods for the horizontal linear complementarity problem, SIAM J. Optimization 4(1) 208-227,

and

G. Zhao and J. Sun (1999). On the rate of local convergence of high-order infeasible path-following algorithms for the  $P_*$  linear complementarity problems, Computational Optimization and Applications 14(1) 293-307,

with many enhancements described by

N. I. M. Gould, D. Orban and D. P. Robinson (2013). Trajectory-following methods for large-scale degenerate convex quadratic programming, Mathematical Programming Computation 5(2) 113-142.

### 1.1.7 Call order

To solve a given problem, functions from the lpb package must be called in the following order:

- `lpb_initialize` - provide default control parameters and set up initial data structures
- `lpb_read_specfile` (optional) - override control values by reading replacement values from a file
- `lpb_import` - set up problem data structures and fixed values
- `lpb_reset_control` (optional) - possibly change control parameters if a sequence of problems are being solved
- `lpb_solve_lp` - solve the linear program
- `lpb_information` (optional) - recover information about the solution and solution process
- `lpb_terminate` - deallocate data structures

See Section 4.1 for examples of use.

### 1.1.8 Unsymmetric matrix storage formats

The unsymmetric  $m$  by  $n$  constraint matrix  $A$  may be presented and stored in a variety of convenient input formats.

Both C-style (0 based) and fortran-style (1-based) indexing is allowed. Choose `control.f_indexing` as `false` for C style and `true` for fortran style; the discussion below presumes C style, but add 1 to indices for the corresponding fortran version.

Wrappers will automatically convert between 0-based (C) and 1-based (fortran) array indexing, so may be used transparently from C. This conversion involves both time and memory overheads that may be avoided by supplying data that is already stored using 1-based indexing.

#### 1.1.8.1 Dense storage format

The matrix  $A$  is stored as a compact dense matrix by rows, that is, the values of the entries of each row in turn are stored in order within an appropriate real one-dimensional array. In this case, component  $n * i + j$  of the storage array `A_val` will hold the value  $A_{ij}$  for  $0 \leq i \leq m - 1$ ,  $0 \leq j \leq n - 1$ .

#### 1.1.8.2 Sparse co-ordinate storage format

Only the nonzero entries of the matrices are stored. For the  $l$ -th entry,  $0 \leq l \leq ne - 1$ , of  $A$ , its row index  $i$ , column index  $j$  and value  $A_{ij}$ ,  $0 \leq i \leq m - 1$ ,  $0 \leq j \leq n - 1$ , are stored as the  $l$ -th components of the integer arrays `A_row` and `A_col` and real array `A_val`, respectively, while the number of nonzeros is recorded as `A_ne = ne`.

#### 1.1.8.3 Sparse row-wise storage format

Again only the nonzero entries are stored, but this time they are ordered so that those in row  $i$  appear directly before those in row  $i+1$ . For the  $i$ -th row of  $A$  the  $i$ -th component of the integer array `A_ptr` holds the position of the first entry in this row, while `A_ptr(m)` holds the total number of entries. The column indices  $j$ ,  $0 \leq j \leq n - 1$ , and values  $A_{ij}$  of the nonzero entries in the  $i$ -th row are stored in components  $l = A\_ptr(i), \dots, A\_ptr(i+1)-1$ ,  $0 \leq i \leq m - 1$ , of the integer array `A_col`, and real array `A_val`, respectively. For sparse matrices, this scheme almost always requires less storage than its predecessor.





## Chapter 2

# File Index

### 2.1 File List

Here is a list of all files with brief descriptions:

<a href="#">galahad_lpb.h</a>	.....	??
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## Chapter 3

# File Documentation

### 3.1 galahad\_lpb.h File Reference

```
#include <stdbool.h>
#include <stdint.h>
#include "galahad_precision.h"
#include "galahad_cfunctions.h"
#include "galahad_fdc.h"
#include "galahad_sb1s.h"
#include "galahad_fit.h"
#include "galahad_roots.h"
#include "galahad_cro.h"
#include "galahad_rpd.h"
```

#### Data Structures

- struct [lpb\\_control\\_type](#)
- struct [lpb\\_time\\_type](#)
- struct [lpb\\_inform\\_type](#)

#### Functions

- void [lpb\\_initialize](#) (void \*\*data, struct [lpb\\_control\\_type](#) \*control, int \*status)
- void [lpb\\_read\\_specfile](#) (struct [lpb\\_control\\_type](#) \*control, const char specfile[])
- void [lpb\\_import](#) (struct [lpb\\_control\\_type](#) \*control, void \*\*data, int \*status, int n, int m, const char A\_type[], int A\_ne, const int A\_row[], const int A\_col[], const int A\_ptr[])
- void [lpb\\_reset\\_control](#) (struct [lpb\\_control\\_type](#) \*control, void \*\*data, int \*status)
- void [lpb\\_solve\\_lp](#) (void \*\*data, int \*status, int n, int m, const real\_wp\_ g[], const real\_wp\_ f, int a\_ne, const real\_wp\_ A\_val[], const real\_wp\_ c\_l[], const real\_wp\_ c\_u[], const real\_wp\_ x\_l[], const real\_wp\_ x\_u[], real\_wp\_ x[], real\_wp\_ c[], real\_wp\_ y[], real\_wp\_ z[], int x\_stat[], int c\_stat[])
- void [lpb\\_information](#) (void \*\*data, struct [lpb\\_inform\\_type](#) \*inform, int \*status)
- void [lpb\\_terminate](#) (void \*\*data, struct [lpb\\_control\\_type](#) \*control, struct [lpb\\_inform\\_type](#) \*inform)

### 3.1.1 Data Structure Documentation

#### 3.1.1.1 struct lpb\_control\_type

control derived type as a C struct

##### Examples

[lpbt.c](#), and [lpbtf.c](#).

##### Data Fields

bool	f_indexing	use C or Fortran sparse matrix indexing
int	error	error and warning diagnostics occur on stream error
int	out	general output occurs on stream out
int	print_level	the level of output required is specified by print_level
int	start_print	any printing will start on this iteration
int	stop_print	any printing will stop on this iteration
int	maxit	at most maxit inner iterations are allowed
int	infeas_max	the number of iterations for which the overall infeasibility of the problem is not reduced by at least a factor .reduce_infeas before the problem is flagged as infeasible (see reduce_infeas)
int	muzero_fixed	the initial value of the barrier parameter will not be changed for the first muzero_fixed iterations
int	restore_problem	indicate whether and how much of the input problem should be restored on output. Possible values are <ul style="list-style-type: none"> <li>• 0 nothing restored</li> <li>• 1 scalar and vector parameters</li> <li>• 2 all parameters</li> </ul>
int	indicator_type	specifies the type of indicator function used. Possible values are <ul style="list-style-type: none"> <li>• 1 primal indicator: constraint active if and only if distance to nearest bound <math>\leq</math> .indicator_p_tol</li> <li>• 2 primal-dual indicator: constraint active if and only if distance to the nearest bound <math>\leq</math> .indicator_tol_pd * size of corresponding multiplier</li> <li>• 3 primal-dual indicator: constraint active if and only if distance to the nearest bound <math>\leq</math> .indicator_tol_tapia * distance to same bound at previous iteration</li> </ul>

## Data Fields

int	arc	<p>which residual trajectory should be used to aim from the current iteration to the solution</p> <ul style="list-style-type: none"> <li>• 1 the Zhang linear residual trajectory</li> <li>• 2 the Zhao-Sun quadratic residual trajectory</li> <li>• 3 the Zhang arc ultimately switching to the Zhao-Sun residual trajectory</li> <li>• 4 the mixed linear-quadratic residual trajectory</li> </ul>
int	series_order	the order of (Taylor/Puiseux) series to fit to the path data
int	sif_file_device	specifies the unit number to write generated SIF file describing the current problem
int	qplib_file_device	specifies the unit number to write generated QPLIB file describing the current problem
real_wp_	infinity	any bound larger than infinity in modulus will be regarded as infinite
real_wp_	stop_abs_p	the required absolute and relative accuracies for the primal infeasibility
real_wp_	stop_rel_p	see stop_abs_p
real_wp_	stop_abs_d	the required absolute and relative accuracies for the dual infeasibility
real_wp_	stop_rel_d	see stop_abs_d
real_wp_	stop_abs_c	the required absolute and relative accuracies for the complementarity
real_wp_	stop_rel_c	see stop_abs_c
real_wp_	prfeas	initial primal variables will not be closer than prfeas from their bound
real_wp_	dufeas	initial dual variables will not be closer than dufear from their bounds
real_wp_	muzero	the initial value of the barrier parameter. If muzero is not positive, it will be reset to an appropriate value
real_wp_	tau	the weight attached to primal-dual infeasibility compared to complementarity when assessing step acceptance
real_wp_	gamma_c	individual complementarities will not be allowed to be smaller than gamma_c times the average value
real_wp_	gamma_f	the average complementarity will not be allowed to be smaller than gamma_f times the primal/dual infeasibility
real_wp_	reduce_infeas	if the overall infeasibility of the problem is not reduced by at least a factor reduce_infeas over .infeas_max iterations, the problem is flagged as infeasible (see infeas_max)
real_wp_	obj_unbounded	if the objective function value is smaller than obj_unbounded, it will be flagged as unbounded from below.

## Data Fields

real_wp_	potential_unbounded	if $W=0$ and the potential function value is smaller than <code>potential_unbounded * number of one-sided bounds</code> , the analytic center will be flagged as unbounded
real_wp_	identical_bounds_tol	any pair of constraint bounds ( <code>c_l,c_u</code> ) or ( <code>x_l,x_u</code> ) that are closer than <code>identical_bounds_tol</code> will be reset to the average of their values
real_wp_	mu_lunge	start terminal extrapolation when <code>mu</code> reaches <code>mu_lunge</code>
real_wp_	indicator_tol_p	if <code>.indicator_type = 1</code> , a constraint/bound will be deemed to be active if and only if distance to nearest bound $\leq$ <code>.indicator_p_tol</code>
real_wp_	indicator_tol_pd	if <code>.indicator_type = 2</code> , a constraint/bound will be deemed to be active if and only if distance to nearest bound $\leq$ <code>.indicator_tol_pd * size of corresponding multiplier</code>
real_wp_	indicator_tol_tapia	if <code>.indicator_type = 3</code> , a constraint/bound will be deemed to be active if and only if distance to nearest bound $\leq$ <code>.indicator_tol_tapia * distance to same bound at previous iteration</code>
real_wp_	cpu_time_limit	the maximum CPU time allowed (-ve means infinite)
real_wp_	clock_time_limit	the maximum elapsed clock time allowed (-ve means infinite)
bool	remove_dependencies	the equality constraints will be preprocessed to remove any linear dependencies if true
bool	treat_zero_bounds_as_general	any problem bound with the value zero will be treated as if it were a general value if true
bool	just_feasible	if <code>.just_feasible</code> is true, the algorithm will stop as soon as a feasible point is found. Otherwise, the optimal solution to the problem will be found
bool	getdua	if <code>.getdua</code> is true, advanced initial values are obtained for the dual variables
bool	puiseux	decide between Puiseux and Taylor series approximations to the arc
bool	every_order	try every order of series up to <code>series_order</code> ?
bool	feasol	if <code>.feasol</code> is true, the final solution obtained will be perturbed so tha variables close to their bounds are moved onto these bounds
bool	balance_initial_complentarity	if <code>.balance_initial_complentarity</code> is true, the initial complemetarity is required to be balanced
bool	crossover	if <code>.crossover</code> is true, cross over the solution to one defined by linearly-independent constraints if possible
bool	space_critical	if <code>.space_critical</code> true, every effort will be made to use as little space as possible. This may result in longer computation time
bool	deallocate_error_fatal	if <code>.deallocate_error_fatal</code> is true, any array/pointer deallocation error will terminate execution. Otherwise, computation will continue

## Data Fields

bool	generate_sif_file	if .generate_sif_file is .true. if a SIF file describing the current problem is to be generated
bool	generate_qplib_file	if .generate_qplib_file is .true. if a QPLIB file describing the current problem is to be generated
char	sif_file_name[31]	name of generated SIF file containing input problem
char	qplib_file_name[31]	name of generated QPLIB file containing input problem
char	prefix[31]	all output lines will be prefixed by .prefix(2:LEN(TRIM(.prefix))-1) where .prefix contains the required string enclosed in quotes, e.g. "string" or 'string'
struct fdc_control_type	fdc_control	control parameters for FDC
struct sbls_control_type	sbls_control	control parameters for SBLS
struct fit_control_type	fit_control	control parameters for FIT
struct roots_control_type	roots_control	control parameters for ROOTS
struct cro_control_type	cro_control	control parameters for CRO

## 3.1.1.2 struct lpb\_time\_type

time derived type as a C struct

## Data Fields

real_wp_	total	the total CPU time spent in the package
real_wp_	preprocess	the CPU time spent preprocessing the problem
real_wp_	find_dependent	the CPU time spent detecting linear dependencies
real_wp_	analyse	the CPU time spent analysing the required matrices prior to factorization
real_wp_	factorize	the CPU time spent factorizing the required matrices
real_wp_	solve	the CPU time spent computing the search direction
real_wp_	clock_total	the total clock time spent in the package
real_wp_	clock_preprocess	the clock time spent preprocessing the problem
real_wp_	clock_find_dependent	the clock time spent detecting linear dependencies
real_wp_	clock_analyse	the clock time spent analysing the required matrices prior to factorization
real_wp_	clock_factorize	the clock time spent factorizing the required matrices
real_wp_	clock_solve	the clock time spent computing the search direction

## 3.1.1.3 struct lpb\_inform\_type

inform derived type as a C struct

## Examples

[lpbt.c](#), and [lpbtf.c](#).

## Data Fields

int	status	return status. See LPB_solve for details
int	alloc_status	the status of the last attempted allocation/deallocation
char	bad_alloc[81]	the name of the array for which an allocation/deallocation error occurred
int	iter	the total number of iterations required
int	factorization_status	the return status from the factorization
int64_t	factorization_integer	the total integer workspace required for the factorization
int64_t	factorization_real	the total real workspace required for the factorization
int	nfacts	the total number of factorizations performed
int	nbacts	the total number of "wasted" function evaluations during the linesearch
int	threads	the number of threads used
real_wp_	obj	the value of the objective function at the best estimate of the solution determined by LPB_solve
real_wp_	primal_infeasibility	the value of the primal infeasibility
real_wp_	dual_infeasibility	the value of the dual infeasibility
real_wp_	complementary_slackness	the value of the complementary slackness
real_wp_	init_primal_infeasibility	these values at the initial point (needed by GALAHAD_CLPB)
real_wp_	init_dual_infeasibility	see init_primal_infeasibility
real_wp_	init_complementary_slackness	see init_primal_infeasibility
real_wp_	potential	the value of the logarithmic potential function sum $-\log(\text{distance to constraint boundary})$
real_wp_	non_negligible_pivot	the smallest pivot which was not judged to be zero when detecting linearly dependent constraints
bool	feasible	is the returned "solution" feasible?
int	checkpointsIter[16]	checkpoints(i) records the iteration at which the criticality measures first fall below $10^{-i-1}$ , $i = 0, \dots, 15$ (-1 means not achieved)
real_wp_	checkpointsTime[16]	see checkpointsIter
struct lpb_time_type	time	timings (see above)
struct fdc_inform_type	fdc_inform	inform parameters for FDC
struct sbis_inform_type	sbis_inform	inform parameters for SBLS
struct fit_inform_type	fit_inform	return information from FIT
struct roots_inform_type	roots_inform	return information from ROOTS
struct cro_inform_type	cro_inform	inform parameters for CRO
struct rpd_inform_type	rpd_inform	inform parameters for RPD

## 3.1.2 Function Documentation



### 3.1.2.1 lpb\_initialize()

```
void lpb_initialize (
    void ** data,
    struct lpb_control_type * control,
    int * status )
```

Set default control values and initialize private data

#### Parameters

in, out	<i>data</i>	holds private internal data
out	<i>control</i>	is a struct containing control information (see <a href="#">lpb_control_type</a> )
out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently): <ul style="list-style-type: none"> <li>• 0. The import was succesful.</li> </ul>

#### Examples

[lpbt.c](#), and [lpbtf.c](#).

### 3.1.2.2 lpb\_read\_specfile()

```
void lpb_read_specfile (
    struct lpb_control_type * control,
    const char specfile[] )
```

Read the content of a specification file, and assign values associated with given keywords to the corresponding control parameters. By default, the spcification file will be named RUNLPB.SPC and lie in the current directory. Refer to Table 2.1 in the fortran documentation provided in \$GALAHAD/doc/lpb.pdf for a list of keywords that may be set.

#### Parameters

in, out	<i>control</i>	is a struct containing control information (see <a href="#">lpb_control_type</a> )
in	<i>specfile</i>	is a character string containing the name of the specification file

### 3.1.2.3 lpb\_import()

```
void lpb_import (
    struct lpb_control_type * control,
    void ** data,
    int * status,
    int n,
```

```

int m,
const char A_type[],
int A_ne,
const int A_row[],
const int A_col[],
const int A_ptr[] )

```

Import problem data into internal storage prior to solution.

#### Parameters

in	<i>control</i>	is a struct whose members provide control paramters for the remaining procedures (see <a href="#">lpb_control_type</a> )
in, out	<i>data</i>	holds private internal data
in, out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are: <ul style="list-style-type: none"> <li>• 0. The import was succesful</li> <li>• -1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in inform.alloc_status and inform.bad_alloc respectively.</li> <li>• -3. The restrictions <math>n &gt; 0</math> or <math>m &gt; 0</math> or requirement that <i>A_type</i> contains its relevant string 'dense', 'coordinate' or 'sparse_by_rows' has been violated.</li> </ul>
in	<i>n</i>	is a scalar variable of type int, that holds the number of variables.
in	<i>m</i>	is a scalar variable of type int, that holds the number of general linear constraints.
in	<i>A_type</i>	is a one-dimensional array of type char that specifies the <a href="#">unsymmetric storage scheme</a> used for the constraint Jacobian, <i>A</i> . It should be one of 'coordinate', 'sparse_by_rows' or 'dense'; lower or upper case variants are allowed.
in	<i>A_ne</i>	is a scalar variable of type int, that holds the number of entries in <i>A</i> in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes.
in	<i>A_row</i>	is a one-dimensional array of size <i>A_ne</i> and type int, that holds the row indices of <i>A</i> in the sparse co-ordinate storage scheme. It need not be set for any of the other schemes, and in this case can be NULL.
in	<i>A_col</i>	is a one-dimensional array of size <i>A_ne</i> and type int, that holds the column indices of <i>A</i> in either the sparse co-ordinate, or the sparse row-wise storage scheme. It need not be set when the dense or diagonal storage schemes are used, and in this case can be NULL.
in	<i>A_ptr</i>	is a one-dimensional array of size $n+1$ and type int, that holds the starting position of each row of <i>A</i> , as well as the total number of entries, in the sparse row-wise storage scheme. It need not be set when the other schemes are used, and in this case can be NULL.

#### Examples

[lpbt.c](#), and [lpbtf.c](#).

### 3.1.2.4 lpb\_reset\_control()

```
void lpb_reset_control (
    struct lpb_control_type * control,
    void ** data,
    int * status )
```

Reset control parameters after import if required.

#### Parameters

in	<i>control</i>	is a struct whose members provide control paramters for the remaining prcedures (see <a href="#">lpb_control_type</a> )
in, out	<i>data</i>	holds private internal data
in, out	<i>status</i>	is a scalar variable of type int, that gives the exit status from the package. Possible values are: <ul style="list-style-type: none"> <li>• 0. The import was succesful.</li> </ul>

### 3.1.2.5 lpb\_solve\_lp()

```
void lpb_solve_lp (
    void ** data,
    int * status,
    int n,
    int m,
    const real_wp_ g[],
    const real_wp_ f,
    int a_ne,
    const real_wp_ A_val[],
    const real_wp_ c_l[],
    const real_wp_ c_u[],
    const real_wp_ x_l[],
    const real_wp_ x_u[],
    real_wp_ x[],
    real_wp_ c[],
    real_wp_ y[],
    real_wp_ z[],
    int x_stat[],
    int c_stat[] )
```

Solve the linear program.

#### Parameters

in, out	<i>data</i>	holds private internal data
---------	-------------	-----------------------------

## Parameters

<code>in, out</code>	<code>status</code>	<p>is a scalar variable of type int, that gives the entry and exit status from the package. Possible exit are:</p> <ul style="list-style-type: none"> <li>• 0. The run was succesful.</li> <li>• -1. An allocation error occurred. A message indicating the offending array is written on unit control.error, and the returned allocation status and a string containing the name of the offending array are held in <code>inform.alloc_status</code> and <code>inform.bad_alloc</code> respectively.</li> <li>• -2. A deallocation error occurred. A message indicating the offending array is written on unit control.error and the returned allocation status and a string containing the name of the offending array are held in <code>inform.alloc_status</code> and <code>inform.bad_alloc</code> respectively.</li> <li>• -3. The restrictions <math>n &gt; 0</math> and <math>m &gt; 0</math> or requirement that <code>A_type</code> contains its relevant string 'dense', 'coordinate' or 'sparse_by_rows' has been violated.</li> <li>• -5. The simple-bound constraints are inconsistent.</li> <li>• -7. The constraints appear to have no feasible point.</li> <li>• -9. The analysis phase of the factorization failed; the return status from the factorization package is given in the component <code>inform.factor_status</code></li> <li>• -10. The factorization failed; the return status from the factorization package is given in the component <code>inform.factor_status</code>.</li> <li>• -11. The solution of a set of linear equations using factors from the factorization package failed; the return status from the factorization package is given in the component <code>inform.factor_status</code>.</li> <li>• -16. The problem is so ill-conditioned that further progress is impossible.</li> <li>• -17. The step is too small to make further impact.</li> <li>• -18. Too many iterations have been performed. This may happen if <code>control.maxit</code> is too small, but may also be symptomatic of a badly scaled problem.</li> <li>• -19. The CPU time limit has been reached. This may happen if <code>control.cpu_time_limit</code> is too small, but may also be symptomatic of a badly scaled problem.</li> </ul>
<code>in</code>	<code>n</code>	is a scalar variable of type int, that holds the number of variables
<code>in</code>	<code>m</code>	is a scalar variable of type int, that holds the number of general linear constraints.
<code>in</code>	<code>g</code>	is a one-dimensional array of size n and type double, that holds the linear term $g$ of the objective function. The j-th component of $g$ , $j = 0, \dots, n-1$ , contains $g_j$ .
<code>in</code>	<code>f</code>	is a scalar of type double, that holds the constant term $f$ of the objective function.
<code>in</code>	<code>a_ne</code>	is a scalar variable of type int, that holds the number of entries in the constraint Jacobian matrix $A$ .
<code>in</code>	<code>A_val</code>	is a one-dimensional array of size <code>a_ne</code> and type double, that holds the values of the entries of the constraint Jacobian matrix $A$ in any of the available storage schemes.
<code>in</code>	<code>c_l</code>	is a one-dimensional array of size m and type double, that holds the lower bounds $c^l$ on the constraints $Ax$ . The i-th component of <code>c_l</code> , $i = 0, \dots, m-1$ , contains $c_i^l$ .
<code>in</code>	<code>c_u</code>	is a one-dimensional array of size m and type double, that holds the upper bounds $c^u$ on the constraints $Ax$ . The i-th component of <code>c_u</code> , $i = 0, \dots, m-1$ , contains $c_i^u$ .
<code>in</code>	<code>x_l</code>	is a one-dimensional array of size n and type double, that holds the lower bounds $x^l$ on the variables $x$ . The j-th component of <code>x_l</code> , $j = 0, \dots, n-1$ , contains $x_j^l$ .

## Parameters

in	$x_u$	is a one-dimensional array of size $n$ and type double, that holds the upper bounds $x^l$ on the variables $x$ . The $j$ -th component of $x_u$ , $j = 0, \dots, n-1$ , contains $x_j^l$ .
in, out	$x$	is a one-dimensional array of size $n$ and type double, that holds the values $x$ of the optimization variables. The $j$ -th component of $x$ , $j = 0, \dots, n-1$ , contains $x_j$ .
out	$c$	is a one-dimensional array of size $m$ and type double, that holds the residual $c(x)$ . The $i$ -th component of $c$ , $i = 0, \dots, m-1$ , contains $c_i(x)$ .
in, out	$y$	is a one-dimensional array of size $n$ and type double, that holds the values $y$ of the Lagrange multipliers for the general linear constraints. The $j$ -th component of $y$ , $j = 0, \dots, m-1$ , contains $y_i$ .
in, out	$z$	is a one-dimensional array of size $n$ and type double, that holds the values $z$ of the dual variables. The $j$ -th component of $z$ , $j = 0, \dots, n-1$ , contains $z_j$ .
out	$x_{stat}$	is a one-dimensional array of size $n$ and type int, that gives the optimal status of the problem variables. If $x_{stat}(j)$ is negative, the variable $x_j$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.
out	$c_{stat}$	is a one-dimensional array of size $m$ and type int, that gives the optimal status of the general linear constraints. If $c_{stat}(i)$ is negative, the constraint value $a_i^T x$ most likely lies on its lower bound, if it is positive, it lies on its upper bound, and if it is zero, it lies between its bounds.

## Examples

[lpbt.c](#), and [lpbtf.c](#).

## 3.1.2.6 lpb\_information()

```
void lpb_information (
    void ** data,
    struct lpb_inform_type * inform,
    int * status )
```

Provides output information

## Parameters

in, out	$data$	holds private internal data
out	$inform$	is a struct containing output information (see <a href="#">lpb_inform_type</a> )
out	$status$	is a scalar variable of type int, that gives the exit status from the package. Possible values are (currently): <ul style="list-style-type: none"> <li>• 0. The values were recorded successfully</li> </ul>

## Examples

[lpbt.c](#), and [lpbtf.c](#).

### 3.1.2.7 lpb\_terminate()

```
void lpb_terminate (
    void ** data,
    struct lpb_control_type * control,
    struct lpb_inform_type * inform )
```

Deallocate all internal private storage

#### Parameters

in, out	<i>data</i>	holds private internal data
out	<i>control</i>	is a struct containing control information (see <a href="#">lpb_control_type</a> )
out	<i>inform</i>	is a struct containing output information (see <a href="#">lpb_inform_type</a> )

#### Examples

[lpbt.c](#), and [lpbtf.c](#).

## Chapter 4

# Example Documentation

### 4.1 lpbt.c

This is an example of how to use the package to solve a linear program. A variety of supported constraint matrix storage formats are shown.

Notice that C-style indexing is used, and that this is flagged by setting `control.f_indexing` to `false`.

```
/* lpbt.c */
/* Full test for the LPB C interface using C sparse matrix indexing */
#include <stdio.h>
#include <math.h>
#include "galahad_precision.h"
#include "galahad_cfunctions.h"
#include "galahad_lpb.h"
int main(void) {
    // Derived types
    void *data;
    struct lpb_control_type control;
    struct lpb_inform_type inform;
    // Set problem data
    int n = 3; // dimension
    int m = 2; // number of general constraints
    real_wp_ g[] = {0.0, 2.0, 0.0}; // linear term in the objective
    real_wp_ f = 1.0; // constant term in the objective
    int A_ne = 4; // Jacobian elements
    int A_row[] = {0, 0, 1, 1}; // row indices
    int A_col[] = {0, 1, 1, 2}; // column indices
    int A_ptr[] = {0, 2, 4}; // row pointers
    real_wp_ A_val[] = {2.0, 1.0, 1.0, 1.0}; // values
    real_wp_ c_l[] = {1.0, 2.0}; // constraint lower bound
    real_wp_ c_u[] = {2.0, 2.0}; // constraint upper bound
    real_wp_ x_l[] = {-1.0, - INFINITY, - INFINITY}; // variable lower bound
    real_wp_ x_u[] = {1.0, INFINITY, 2.0}; // variable upper bound
    // Set output storage
    real_wp_ c[m]; // constraint values
    int x_stat[n]; // variable status
    int c_stat[m]; // constraint status
    char st;
    int status;
    printf(" C sparse matrix indexing\n\n");
    printf(" basic tests of lp storage formats\n\n");
    for( int d=1; d <= 3; d++){
        // Initialize LPB
        lpb_initialize( &data, &control, &status );
        // Set user-defined control options
        control.f_indexing = false; // C sparse matrix indexing
        // Start from 0
        real_wp_ x[] = {0.0,0.0,0.0};
        real_wp_ y[] = {0.0,0.0};
        real_wp_ z[] = {0.0,0.0,0.0};
        switch(d){
            case 1: // sparse co-ordinate storage
                st = 'C';
                lpb_import( &control, &data, &status, n, m,
```

```

        "coordinate", A_ne, A_row, A_col, NULL );
    lpb_solve_lp( &data, &status, n, m, g, f,
        A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
        x_stat, c_stat );

    break;
printf(" case %li break\n",d);
case 2: // sparse by rows
    st = 'R';
    lpb_import( &control, &data, &status, n, m,
        "sparse_by_rows", A_ne, NULL, A_col, A_ptr );
    lpb_solve_lp( &data, &status, n, m, g, f,
        A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
        x_stat, c_stat );

    break;
case 3: // dense
    st = 'D';
    int A_dense_ne = 6; // number of elements of A
    real_wp_ A_dense[] = {2.0, 1.0, 0.0, 0.0, 1.0, 1.0};
    lpb_import( &control, &data, &status, n, m,
        "dense", A_ne, NULL, NULL, NULL );
    lpb_solve_lp( &data, &status, n, m, g, f,
        A_dense_ne, A_dense, c_l, c_u, x_l, x_u,
        x, c, y, z, x_stat, c_stat );

    break;
}
lpb_information( &data, &inform, &status );
if(inform.status == 0){
    printf("%c:%6i iterations. Optimal objective value = %5.2f status = %li\n",
        st, inform.iter, inform.obj, inform.status);
}else{
    printf("%c: LPB_solve exit status = %li\n", st, inform.status);
}
//printf("x: ");
//for( int i = 0; i < n; i++) printf("%f ", x[i]);
//printf("\n");
//printf("gradient: ");
//for( int i = 0; i < n; i++) printf("%f ", g[i]);
//printf("\n");
// Delete internal workspace
lpb_terminate( &data, &control, &inform );
}
}

```

## 4.2 lpbtfc

This is the same example, but now fortran-style indexing is used.

```

/* lpbtfc.c */
/* Full test for the LPB C interface using Fortran sparse matrix indexing */
#include <stdio.h>
#include <math.h>
#include "galahad_precision.h"
#include "galahad_cfunctions.h"
#include "galahad_lpb.h"
int main(void) {
    // Derived types
    void *data;
    struct lpb_control_type control;
    struct lpb_inform_type inform;
    // Set problem data
    int n = 3; // dimension
    int m = 2; // number of general constraints
    real_wp_ g[] = {0.0, 2.0, 0.0}; // linear term in the objective
    real_wp_ f = 1.0; // constant term in the objective
    int A_ne = 4; // Jacobian elements
    int A_row[] = {1, 1, 2, 2}; // row indices
    int A_col[] = {1, 2, 2, 3}; // column indices
    int A_ptr[] = {1, 3, 5}; // row pointers
    real_wp_ A_val[] = {2.0, 1.0, 1.0, 1.0 }; // values
    real_wp_ c_l[] = {1.0, 2.0}; // constraint lower bound
    real_wp_ c_u[] = {2.0, 2.0}; // constraint upper bound
    real_wp_ x_l[] = {-1.0, - INFINITY, - INFINITY}; // variable lower bound
    real_wp_ x_u[] = {1.0, INFINITY, 2.0}; // variable upper bound
    // Set output storage
    real_wp_ c[m]; // constraint values
    int x_stat[n]; // variable status
    int c_stat[m]; // constraint status
    char st;
    int status;
    printf(" Fortran sparse matrix indexing\n\n");
    printf(" basic tests of lp storage formats\n\n");
}

```



```

for( int d=1; d <= 3; d++){
    // Initialize LPB
    lpb_initialize( &data, &control, &status );
    // Set user-defined control options
    control.f_indexing = true; // Fortran sparse matrix indexing
    // Start from 0
    real_wp_ x[] = {0.0,0.0,0.0};
    real_wp_ y[] = {0.0,0.0};
    real_wp_ z[] = {0.0,0.0,0.0};
    switch(d){
        case 1: // sparse co-ordinate storage
            st = 'C';
            lpb_import( &control, &data, &status, n, m,
                "coordinate", A_ne, A_row, A_col, NULL );
            lpb_solve_lp( &data, &status, n, m, g, f,
                A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
                x_stat, c_stat );
            break;
        printf(" case %li break\n",d);
        case 2: // sparse by rows
            st = 'R';
            lpb_import( &control, &data, &status, n, m,
                "sparse_by_rows", A_ne, NULL, A_col, A_ptr );
            lpb_solve_lp( &data, &status, n, m, g, f,
                A_ne, A_val, c_l, c_u, x_l, x_u, x, c, y, z,
                x_stat, c_stat );
            break;
        case 3: // dense
            st = 'D';
            int A_dense_ne = 6; // number of elements of A
            real_wp_ A_dense[] = {2.0, 1.0, 0.0, 0.0, 1.0, 1.0};
            lpb_import( &control, &data, &status, n, m,
                "dense", A_ne, NULL, NULL, NULL );
            lpb_solve_lp( &data, &status, n, m, g, f,
                A_dense_ne, A_dense, c_l, c_u, x_l, x_u,
                x, c, y, z, x_stat, c_stat );
            break;
    }
    lpb_information( &data, &inform, &status );
    if(inform.status == 0){
        printf("%c:%6i iterations. Optimal objective value = %5.2f status = %li\n",
            st, inform.iter, inform.obj, inform.status);
    }else{
        printf("%c: LPB_solve exit status = %li\n", st, inform.status);
    }
    //printf("x: ");
    //for( int i = 0; i < n; i++) printf("%f ", x[i]);
    //printf("\n");
    //printf("gradient: ");
    //for( int i = 0; i < n; i++) printf("%f ", g[i]);
    //printf("\n");
    // Delete internal workspace
    lpb_terminate( &data, &control, &inform );
}
}

```

