Discrete Structures— Graphs: Definitions

Dr. Ab Mosca (they/them)

Plan for Today

Graph definitions

Are any of the graphs below the same? If you said yes, which ones and why?

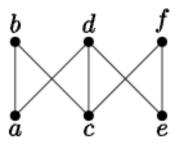
Warm Up: Visualizing

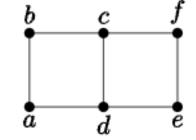
Warm Up:

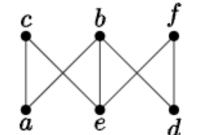
Visualizing

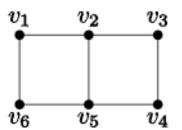
How about these with labels?

Are any of the graphs below the same? If you said yes, which ones and why?



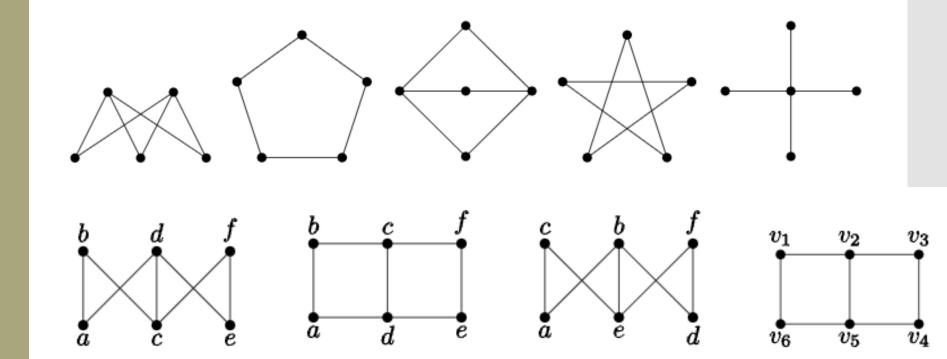






These examples are *drawings* of graphs.

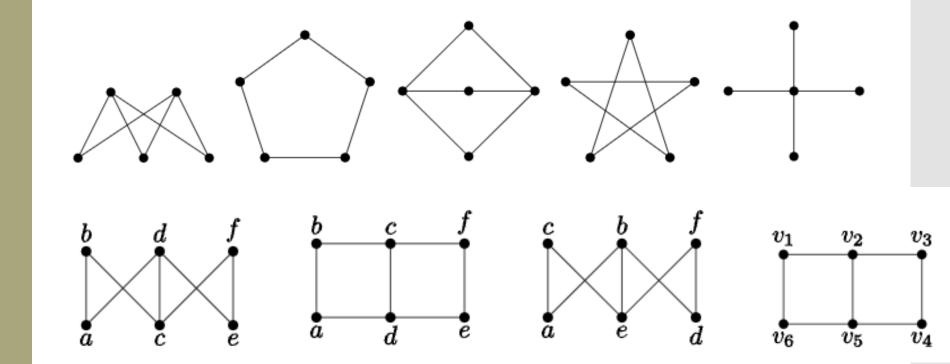




Mathematically, a **graph** is an ordered pair, G = (V, E), consisting of a nonempty set, V (called **vertices**), and a set E (called **edges**) of two-elements subsets of V.

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Definition



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Note: you may also hear vertices called *nodes*.

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Definition

Ex. Graph 1:

$$V = \{a, b, c, d, e\},\$$

$$E = \{\{a, b\}, \{a, c\}, \{a, d\}, \{a, e\}, \{b, c\}, \{d, e\}\}\}$$

Definition

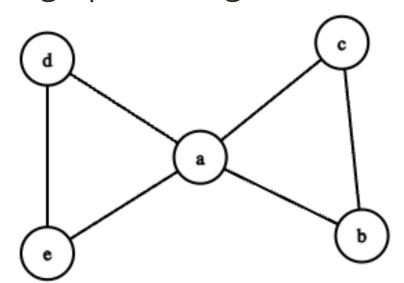
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To visualize this graph we might draw it:



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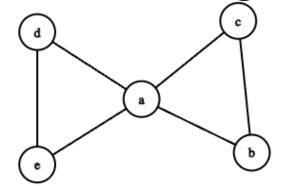
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Draw this graph:

$$V = \{v_1, v_2, v_3, v_4, v_5\}$$

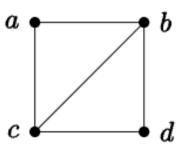
$$E = \{\{v_1, v_3\}, \{v_1, v_5\}, \{v_2, v_4\}, \{v_2, v_5\}, \{v_3, v_5\}, \{v_4, v_5\}\}$$

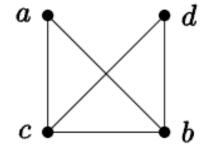
Two *graphs are equal* if their sets of vertices and edges are equal.

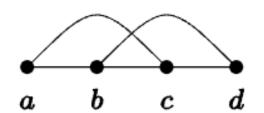
Drawings of the same graph can differ.

Ex.
$$(\{a,b,c,d,e\}, \{\{a,b\},\{a,c\},\{b,c\},\{b,d\},\{c,d\}\})$$

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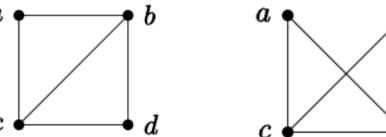


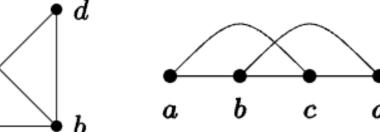
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Are these graphs equal?

$$G_1 = (\{a, b, c\}, \{\{a, b\}, \{b, c\}\})$$

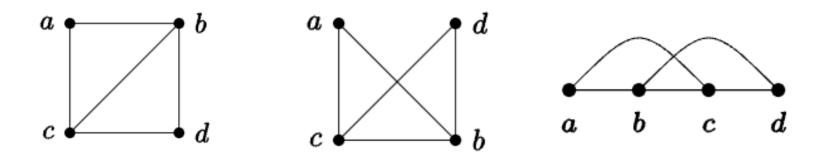
$$G_2 = (\{a, b, c\}, \{\{a, c\}, \{c, b\}\})$$

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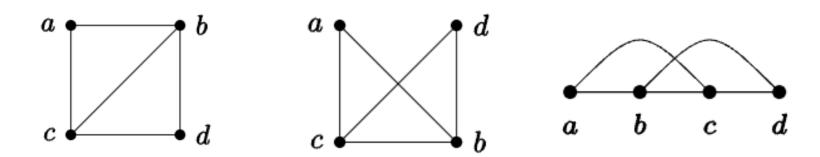
Are these graphs equal? $G_1 = \big(\{a, b, c\}, \big\{\{a, b\}, \{a, c\}, \{b, c\}\big\}\big)$ $G_2 = \big(\{u, v, w\}, \big\{\{u, v\}, \{u, w\}, \{v, w\}\big\}\big)$

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$$(\{a,b,c,d,e\}, \{\{a,b\},\{a,c\},\{b,c\},\{b,d\},\{c,d\}\})$$

To visualize this graph we might draw it:



Would you call these graphs the same? $G_1 = (\{a, b, c\}, \{\{a, b\}, \{a, c\}, \{b, c\}\})$ $G_2 = (\{u, v, w\}, \{\{u, v\}, \{u, w\}, \{v, w\}\})$

An *isomorphism* between two graphs G_1 and G_2 is a bijection $f: V_1 \to V_2$ between the vertices of the graphs such that $\{a,b\}$ is an edge in G_1 if and only if $\{f(a),f(b)\}$ is an edge in G_2 .

Two graphs are *isomorphic* if there is an isomorphism between them. In this case we write $G_1 \cong G_2$.

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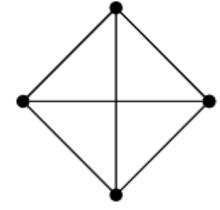
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A collection of isomorphic graphs is called an *isomorphism class*.

Ex. This graph with any labels for V

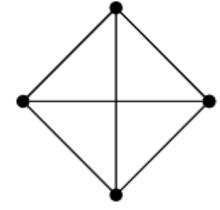


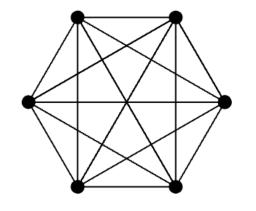
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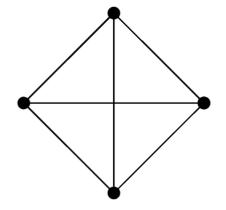
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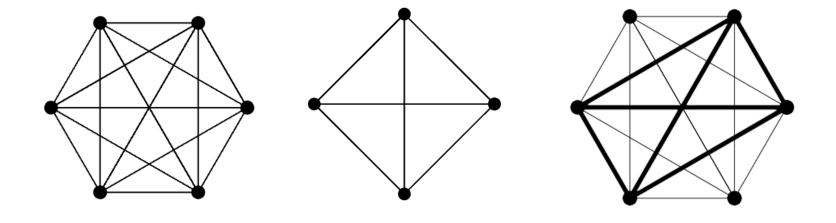
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Do you think these graphs are related? If so, how?

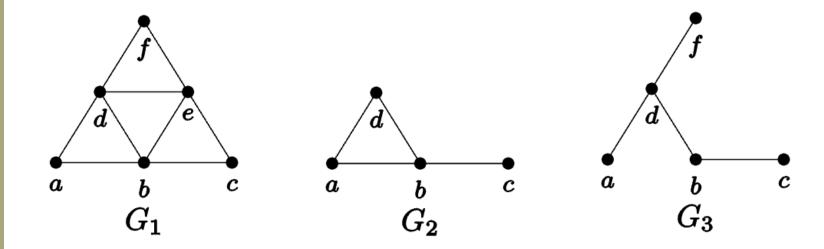


We say that G' = (V', E') is a **subgraph** of G = (V, E), and write $G' \subseteq G$, provided $V' \subseteq V$ and $E' \subseteq E$.

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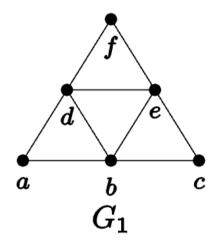
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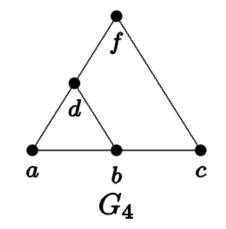
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Is G_4 a subgraph of G_1 ?

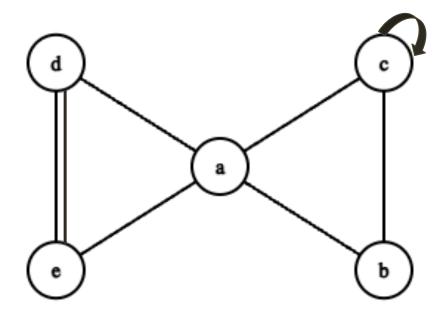
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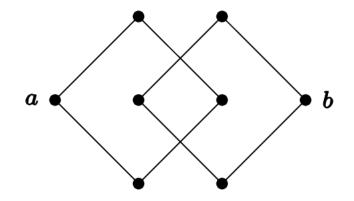
A *multigraph*, is a graph that allows double (or more) edges, and for a vertex to be connected to itself.

Special Graphs



A **connected** graph is a graph where you can get from any vertex to any other vertex by following some path of edges.

Special Graphs



Is this graph connected?

A graph that includes all possible edges is called **complete**. In other words, a graph is complete if every pair of vertices is connected by an edge.

 K_n is the complete graph on n vertices.

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Special Graphs

Each vertex in K_n is adjacent to n-1 other vertices.

How many edges does K_n have?

This can be written symbolically as: $\sum_{v \in V} d(v) = 2e$

Handshake Lemma

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Handshake Lemma

Given a **degree sequence** for a graph (a list of every degree of every vertex in the graph), we can use the handshake lemma to find the number of edges in the graph.

How many vertices and edges must the graph with the degree sequence (4, 4, 3, 3, 3, 2, 1) have?

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We can generalize the previous example into a proposition:

In any graph, the number of vertices with odd degree must be even.

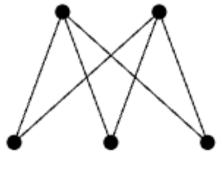
Proposition

Prove this proposition.

We say a graph is **bipartite** if the vertices can be divided into two sets, *A* and *B*, with no two vertices in *A* adjacent and no two vertices in *B* adjacent. The vertices in *A* can be adjacent to some or all of the vertices in *B*.

Bipartite

If each vertex in A is adjacent to all the vertices in B, then the graph is a **complete bipartite graph**, and gets the special name: $K_{m,n}$ where |A| = m, |B| = n.

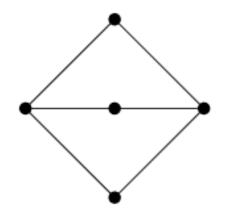


 $K_{2,3}$

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Is this graph bipartite?

 K_n : The complete graph on n vertices.

 $K_{m,n}$: The complete bipartite graph with sets of m and n vertices.

Named Graphs

 C_n : The cycle on n vertices, just one big loop.

 P_n : The path on n+1 vertices (so n edges), just one long path.

