

Discrete Structures— Propositional Logic Pt. 2

Dr. Ab Mosca (they/them)

Plan for Today

- Boolean Algebra
- Propositions
- Quantifiers

One way we model truth values of statements is with **truth tables**.

conjunction

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

disjunction

P	Q	$P \vee Q$
T	T	T
T	F	T
F	T	T
F	F	F

implication

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
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biconditional negation

P	Q	$P \leftrightarrow Q$
T	T	T
T	F	F
F	T	F
F	F	T

P	$\neg P$
T	F
F	T

Warm Up

We also use truth tables to show logical equivalence. Two molecular statements P and Q are **logically equivalent** provided P is true precisely when Q is true.

Work with 1-2 other people to show that $\neg(P \wedge Q)$ is logically equivalent to $\neg P \vee \neg Q$.

De Morgan's Laws

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You just showed

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Notice that these laws show you how to “distribute” a negative into parentheses.

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(The same is true of De Morgan's Laws, or any logically equivalent statements; we can always replace one with another without changing the truth values of a statement.)

Double Negation

A double negation cancels itself out.

In other words, $\neg\neg P$ is logically equivalent to P .

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So far we know:

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 - $\neg(P \vee Q)$ is logically equivalent to $\neg P \wedge \neg Q$
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- Implications are disjunctions
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Practice: Use Boolean Algebra to show $\neg(P \rightarrow Q)$ and $P \wedge \neg Q$ are logically equivalent.

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Practice: Use Boolean Algebra to show $\neg((\neg P \wedge Q) \vee \neg(R \vee \neg S))$ and $(Q \rightarrow P) \wedge (S \rightarrow R)$ are logically equivalent.

Variables and Statements

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Now we can represent the statement above as:

$$P(n) \rightarrow \neg P(n + 7)$$

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Practice: Write the sentence "Primes greater than 2 are odd" symbolically.

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Notice that $P(n) \rightarrow \neg P(n + 7)$ is *not a statement*. We have not specified anything about n ; it can be any value.

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To make this predicate a statement (something that is always true or false), we need to quantify the variable, n .

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Ex.

$$\exists x(x < 0)$$

is read “There exists an x such that x is less than 0”

This statement says there exists a number less than 0.

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The ***universal quantifier*** is \forall and is read “for all” or “every”

Ex.

$$\forall x(x \geq 0)$$

is read “For all x , x is greater than or equal to 0”

This statement says every number is greater than or equal to 0.

Quantifiers

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The ***universal quantifier*** is \forall and is read “for all” or “every”

Practice: Write the sentence “All primes greater than 2 are odd” symbolically.

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Consider this statement: $\forall x \exists y (y < x)$

Let the domain of x, y be the ***natural numbers*** (0, 1, 2, ...)

What is the truth value of the statement?

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Consider this statement: $\exists x \forall y (y \geq x)$

Let the domain of x, y be the **natural numbers** (0, 1, 2, ...)

What is this statement in English?

What is the truth value of the statement?

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Negation of Quantifiers

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Practice: Rewrite the following statement without negation

$$\neg \exists x \forall y (x \leq y)$$

What is each version in English?