

THE AMPL INTERFACE TO CONSTRAINT PROGRAMMING SOLVERS

Victor Zverovich, Robert Fourer



AMPL Optimization



The 13th INFORMS Computing Society Conference (ICS)
January 6th - 8th, 2013, Santa Fe, New Mexico, USA

WHY AMPL?

- **AMPL** is a popular algebraic modeling language:
 - used in businesses, government agencies, and academic institutions (over 100 courses in 2012)
 - large community
(> 1,300 members in **AMPL Google Group** alone)
 - the most popular input format on **NEOS**
(> 200,000 or 57% submissions in 2012)
- AMPL supports a wide range of problem types: linear, mixed integer, quadratic, second-order cone, nonlinear, complementarity problems and more.

CONSTRAINT PROGRAMMING IN AMPL

- **Constraint programming** (CP) allows natural formulation of many combinatorial optimization problems.
- AMPL has supported CP since early 2000s when the paper Extending an Algebraic Modeling Language to Support Constraint Programming [1] was published.
- CP solvers connected to AMPL:
 - **IBM/ILOG CP Optimizer** (ilogcp)
 - **Gecode**

SUPPORTED CP CONSTRUCTS

- Logical operators: `and`, `or`, `not`; iterated `exists`, `forall`
- Conditional operators: `if-then`, `if-then-else`, `==>`, `==> else`, `<==`, `<==>`
- Counting operators: `iterated count`, `atmost`, `atleast`, `exactly`, `numberof`
- Pairwise operators: `alldiff`

See <http://www.ampl.com/NEW/LOGIC/index.html>

AMPL SOLVER LIBRARY

AMPL Solver Library (ASL) is an open source library for connecting solvers to AMPL.

- Available from:
 - Netlib: <http://www.netlib.org/ampl/>
 - GitHub: <https://github.com/vitaut/ampl>
- Includes drivers for several solvers:
 - CPLEX
 - Gurobi
 - MINOS
 - ...

AMPL SOLVER INTERFACES

- C interface:
 - described in **Hooking Your Solver to AMPL**
 - used by most solvers
- C++ interface (new):
 - a very thin wrapper around the C interface
 - type-safe, no casts needed when working with expression trees
 - easy to use, less boilerplate code
 - efficient
 - used by **ilogcp** and **gecode**

PERFORMANCE

Problem	Input+Conv Time	C API Time	C++ API Time	C/C++ Ratio
assign1	2.001	0.0156233	0.0154757	1.009537533
balassign1	2.106	0.086136	0.085996	1.0016279827
flowshp1	0.529	0.000594689	0.000598812	0.9931147004
flowshp2	0.784	0.000593329	0.00059615	0.9952679695
magic	0.558	0.00276411	0.0028042	0.9857035875
mapcoloring	0.557	8.6488E-007	8.6538E-007	0.9994222191
money	0.472	0.000241843	0.000238001	1.0161427893
nqueens	1.073	8.6449E-007	8.6593E-007	0.998337048
sched1	0.595	0.0172613	0.0174135	0.9912596549
sched2	0.684	0.0142709	0.0146368	0.9750013664
party1	33.688	3.88768	3.94516	0.9854302487
party2	126.194	1.20008	1.21847	0.9849073018
sudokuHard	0.661	8.6618E-007	8.6361E-007	1.0029758803
sudokuVeryEasy	0.726	8.6547E-007	8.6383E-007	1.0018985217
				0.9957590574

Time is in milliseconds.

SOLVER WORKFLOW

1. Get a problem instance in AMPL form
2. Process solver options
3. Convert the problem from AMPL to solver form
4. Solve the problem
5. Return solution(s)

Step 3 is the most interesting especially for a CP solver, because it has to deal with expression trees. Other steps are easy to implement.

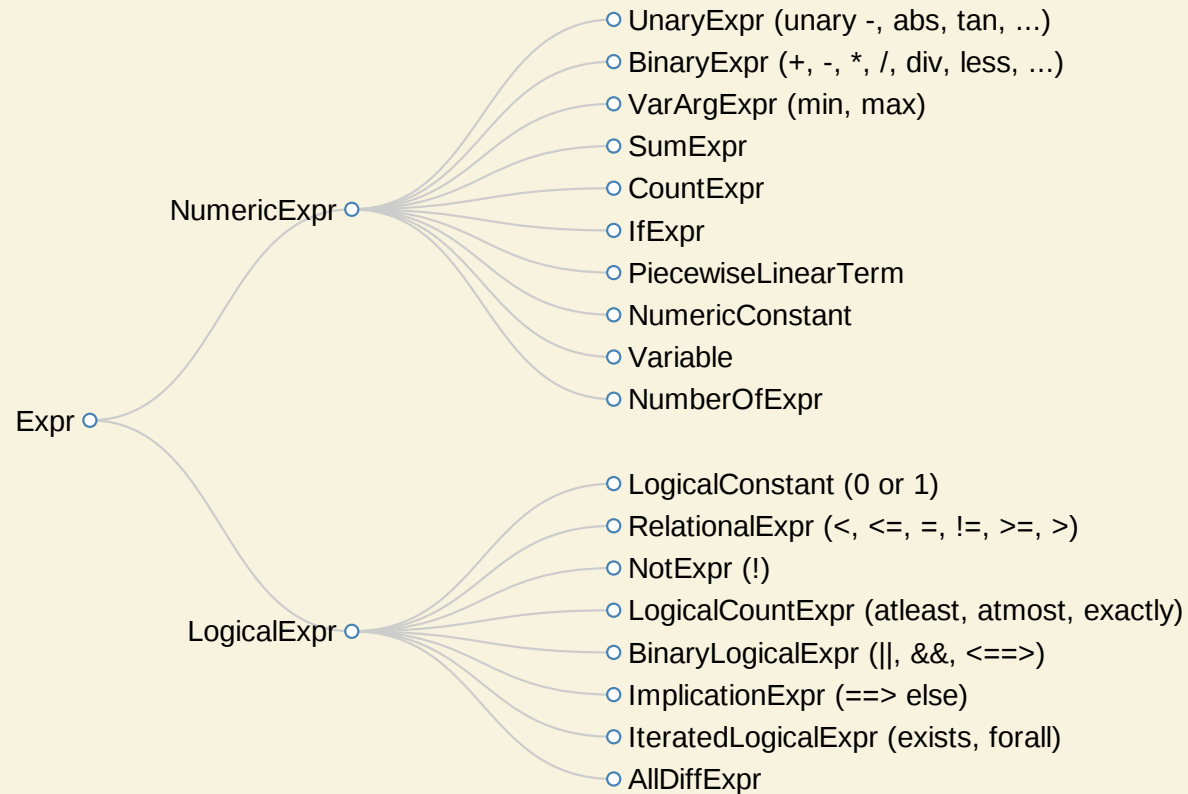
CHOICE OF LANGUAGE



Both IBM/ILOG CP Optimizer and Gecode use C++ for their main APIs. Therefore I'll give all examples in C++ with the new interface library.

However, everything discussed here is possible to do with the C API with a bit more work.

EXPRESSION TREES



WORKING WITH EXPRESSIONS

- `c.value()` returns the value of a numeric or logical constant `c`.
- `e.arg()` returns an argument of a unary expression `e`.
- `e.lhs()` and `e.rhs()` return arguments of a binary expression `e`.
- `Cast<ExprType>(e)` casts expression `e` to `ExprType` if possible, otherwise returns a null expression. Fast alternative to `dynamic_cast`.

ITERATING OVER ARGUMENTS

- Iterating over arguments of an expression e with variable number of arguments:

```
for (auto arg: e) // C++11  
    // use arg
```

or

```
for (SumExpr::iterator i = e.begin(); i != e.end(); ++i)  
    // use *i
```

- Works with VarArgExpr, SumExpr, CountExpr, NumberOfExpr, IteratedLogicalExpr and AllDiffExpr.

TREE TRAVERSAL WITH VISITORS

ilogcp

```
class IlogCPSolver :  
    public ExprVisitor<IlogCPSolver, IloExpr, IloConstraint> {  
public:  
    // Convert logical expressions.  
    // Convert numeric expressions.  
};
```

gecode

```
class NLToGecodeConverter :  
    private ExprVisitor<NLToGecodeConverter,  
                        Gecode::LinExpr, Gecode::BoolExpr> {  
    // Convert logical expressions.  
    // Convert numeric expressions.  
public:  
    LinExpr ConvertFullExpr(NumericExpr e) { return Visit(e); }  
    BoolExpr ConvertFullExpr(LogicalExpr e) {  
        // Process global constraints calling Visit(e)  
        // for nested expressions.  
    }  
};
```

CONVERTING NUMERIC EXPRESSIONS

```
IloExpr VisitNumericConstant(NumericConstant n) {  
    return IloExpr(env_, n.value());  
}  
  
IloExpr VisitVariable(Variable v) {  
    return vars_[v.index()];  
}  
  
IloExpr VisitPlus(BinaryExpr e) {  
    return Visit(e.lhs()) + Visit(e.rhs());  
}  
  
IloExpr VisitPow(BinaryExpr e) {  
    return IloPower(Visit(e.lhs()), Visit(e.rhs()));  
}  
  
IloExpr VisitSum(SumExpr e) {  
    IloExpr sum(env_);  
    for (auto arg: e) // C++11  
        sum += Visit(arg);  
    return sum;  
}
```

CONVERTING LOGICAL EXPRESSIONS

```
IloConstraint VisitLogicalConstant(LogicalConstant c) {  
    return IloNumVar(env_, 1, 1) == c.value();  
}  
  
IloConstraint VisitEqual(RelationalExpr e) {  
    return Visit(e.lhs()) == Visit(e.rhs());  
}  
  
IloConstraint VisitGreater(RelationalExpr e) {  
    return Visit(e.lhs()) > Visit(e.rhs());  
}  
  
IloConstraint VisitAnd(BinaryLogicalExpr e) {  
    return Visit(e.lhs()) && Visit(e.rhs());  
}  
  
IloConstraint IlogCPSolver::VisitExists(IteratedLogicalExpr e) {  
    IloOr disjunction(env_);  
    for (auto arg: e) // C++11  
        disjunction.add(Visit(arg));  
    return disjunction;  
}
```

HANDLING NUMBEROF

```
class IlogCPSolver {  
    // CreateVar is a functor that creates an IlogCP variable.  
    NumberOfMap<IloIntVar, CreateVar> numberofs_  
    ...  
};  
  
IloExpr IlogCPSolver::VisitNumberOf(NumberOfExpr e) {  
    NumericExpr value = e.value();  
    if (NumericConstant num = Cast<NumericConstant>(value))  
        return numberofs_.Add(num.value(), e);  
    IloExpr sum(env_);  
    IloExpr concert_value(Visit(value));  
    for (Expr arg: e)  
        sum += (Visit(arg) == concert_value);  
    return sum;  
}
```

NumberOfMap is a map from numberof expressions with the same argument lists to values and corresponding variables. Such expression can be converted to a single IloDistribute constraint.

EXPRESSION VISITOR

Copy the great architectures.
-- Edward Tufte



- Inspired by the AST visitor from the **Clang compiler frontend**
- Visitor design pattern with static instead of dynamic polymorphism
- Uses **curiously recurring template pattern**
- Very efficient: no virtual function calls, `Visit*` functions can be inlined

TWO-LEVEL CONVERSION

1. Top level - global constraints such as `alldiff` and possible optimizations for the case when expression value is not used
 - `ilogcp`: no extra work is necessary, the Concert interface does necessary processing
 - `gecode`: manual handling of `alldiff` in `ConvertFullExpr`
2. General case for nested expressions

Example:

```
s.t. c: alldiff ({j in 1..n} Row[j]+j);
```

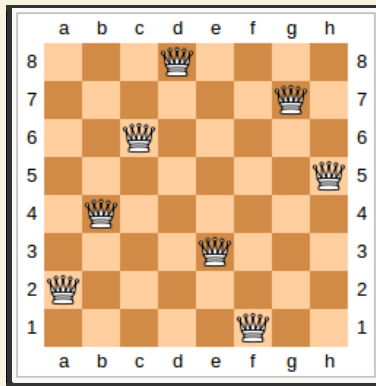
`alldiff (...)` - top level expression

`Row[j] + j` - subexpression

SUPPORTING MULTIPLE SOLVERS

- Separate hierarchies for logical and numeric expressions (ilogcp and gecode) are handled easily
- Possible to deal with more complex expression hierarchies, but with more efforts
- Not necessary to convert all expressions, solver will report an error when unhandled expression is encountered and exit gracefully. For example, gecode doesn't support many nonlinear expressions.

EXAMPLE



```
# Place n queens on an n by n board  
# so that no two queens can attack  
# each other (nqueens.mod).
```

```
param n integer > 0;  
var Row {1..n} integer >= 1 <= n;
```

```
subj to c1: alldiff ({j in 1..n} Row[j]);  
subj to c2: alldiff ({j in 1..n} Row[j]+j);  
subj to c3: alldiff ({j in 1..n} Row[j]-j);
```

More examples available at

<http://www.ampl.com/NEW/LOGIC/examples.html>

EXAMPLE - ILOGCP

```
ampl: model nqueens.mod;
ampl: let n := 20;
ampl: option solver ilogcp;
ampl: solve;
ilogcp 12.4.0: feasible solution
2898 choice points, 1286 fails
ampl: display Row;
Row [*] :=
  1  14
  2  16
  3   8
  4   6
  5  15
  6   3
  ...
 19  10
 20   5
;
```

EXAMPLE - GECODE

```
ampl: model nqueens.mod;  
ampl: let n := 20;  
ampl: option solver gecode;  
ampl: solve;  
gecode 3.7.3: feasible solution  
220 nodes, 105 fails  
ampl: display Row;  
Row [*] :=  
  1   1  
  2   3  
  3   5  
  4   7  
  5  14  
  6  10  
  ...  
19  15  
20  13  
;
```

LINKS

- GitHub repository:
<https://github.com/vitaut/ampl>
- C++ Solver API:
solvers/util
- Gecode AMPL solver:
solvers/gecode
- IlogCP AMPL solver:
solvers/ilogcp