



SMART CONTRACT AUDIT REPORT

for

Spot Protocol



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1 | Introduction

Given the opportunity to review the design document and related smart contract source code of the SPOT protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About SPOT

SPOT is a decentralized, inflation resistant store of value designed to be resilient in all market conditions. It has no reliance on centralized custodians, liquidations, or lenders of last resort. It has no feedback loops, no dependence on continual growth, and is free from bank runs. The system bends safely rather than breaking catastrophically in extreme market scenarios, and can forever resume its function without bailouts. SPOT can be held directly or rotated in as an alternative collateral asset to USDC within existing systems. It uses AMPL as the underlying unit of account, Buttonwood Tranche for collateral preparation, and onchain governance through the FORTH DAO. The basic information of the audited protocol is as follows:

Table 1.1: Basic Information of The SPOT Protocol

Item	Description
Client	Fragments, Inc.
Website	https://ampleforth.org
Type	EVM Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	June 3, 2023

In the following, we show the Git repository of reviewed files and the commit hash value used in

this audit.

- <https://github.com/ampleforth/spot.git> (901435c)

1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email (contact@peckshield.com).

Table 1.2: Vulnerability Severity Classification

Impact	High	Critical	High	Medium
	Medium	High	Medium	Low
	Low	Medium	Low	Low
		High	Medium	Low
		Likelihood		

1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [10]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the

Table 1.3: The Full List of Check Items

Category	Check Item
Basic Coding Bugs	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
Advanced DeFi Scrutiny	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
Additional Recommendations	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.



Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functionality that processes data.
Numeric Errors	Weaknesses in this category are related to improper calculation or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
Time and State	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
Error Conditions, Return Values, Status Codes	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper management of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
Business Logics	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the design and implementation of the `spot` protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	0	
Medium	0	
Low	4	
Informational	1	
Total	5	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in [Section 3](#).

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 4 low-severity vulnerabilities and 1 informational recommendation.

Table 2.1: Key Spot Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Low	Possible Costly Vault Tokens From Improper Initialization	Time And State	Resolved
PVE-002	Low	Improved Precision Calculation in <code>_computeRolloverAmt()</code>	Numeric Errors	Resolved
PVE-003	Low	Accommodation of Non-ERC20-Compliant Tokens	Code Practices	Resolved
PVE-004	Informational	Improved Validations Logic in RouterV1	Code Practices	Resolved
PVE-005	Low	Trust on Admin Keys	Security Features	Mitigated

Besides recommending specific countermeasures to mitigate these issues, we also emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for details.

3 | Detailed Results

3.1 Possible Costly Vault Tokens From Improper Initialization

- ID: PVE-001
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: RolloverVault
- Category: Time and State [6]
- CWE subcategory: CWE-362 [4]

Description

The `Spot` protocol allows users to deposit supported tranche tokens and get in return the perps to represent the overall share/ownership. It also has a `RolloverVault` that generates yield (from fees) by performing rollovers on `PerpetualTranche`. While examining the vault logic, we notice an issue that may unnecessarily make the vault token extremely expensive and bring hurdles (or even causes loss) for later depositors.

To elaborate, we show below the `deposit()` routine, which is used for participating users to deposit the collateral and get respective vault tokens in return. The issue occurs when the `RolloverVault` is being initialized under the assumption that it is empty.

```

280     function deposit(uint256 amount) external override nonReentrant whenNotPaused
281         returns (uint256) {
282             uint256 totalSupply_ = totalSupply();
283             uint256 notes = (totalSupply_ > 0) ? totalSupply_.mulDiv(amount, getTVL()) : (
                amount * INITIAL_RATE);
284
285             underlying.safeTransferFrom(_msgSender(), address(this), amount);
286             _syncAsset(underlying);
287
288             _mint(_msgSender(), notes);
289             return notes;
290         }

```

Listing 3.1: `RolloverVault::deposit()`

```

320     function getTVL() public override returns (uint256) {
321         uint256 totalValue = 0;

322         // The underlying balance
323         totalValue += underlying.balanceOf(address(this));

324         // The deployed asset value denominated in the underlying
325         for (uint256 i = 0; i < _deployed.length(); i++) {
326             ITranche tranche = ITranche(_deployed.at(i));
327             uint256 trancheBalance = tranche.balanceOf(address(this));
328             if (trancheBalance > 0) {
329                 (uint256 collateralBalance, uint256 trancheSupply) = tranche.
330                     getTrancheCollateralization();
331                 totalValue += collateralBalance.mulDiv(trancheBalance, trancheSupply);
332             }
333         }

334         // The earned asset (perp token) value denominated in the underlying
335         uint256 perpBalance = perp.balanceOf(address(this));
336         if (perpBalance > 0) {
337             // The "earned" asset is assumed to be the perp token.
338             // Perp tokens are assumed to have the same denomination as the underlying
339             totalValue += perpBalance.mulDiv(IPerpetualTranche(address(perp)).
340                 getAvgPrice(), PERP_UNIT_PRICE);
341         }

342         return totalValue;
343     }

```

Listing 3.2: RolloverVault::getTVL()

Specifically, when it is being initialized (line 282), the minted amount is based on the given amount, which could be manipulatable by the malicious actor. As this is the first deposit, the current total supply equals 0 and the return amount is computed as `amount * INITIAL_RATE`. With that, the actor can further donate a huge amount of the underlying assets with the goal of making the perps extremely expensive.

An extremely expensive vault tokens can be very inconvenient to use as a small number of 1 Wei may denote a large value. Furthermore, it can lead to precision issue in truncating the vault tokens for deposited assets. If truncated to be zero, the deposited assets are essentially considered dust and kept by the pool without returning any pool tokens.

This is a known issue that has been mitigated in popular Uniswap. When providing the initial liquidity to the contract (i.e. when totalSupply is 0), the liquidity provider must sacrifice 1000 LP tokens (by sending them to `address(0)`). By doing so, we can ensure the granularity of the LP tokens is always at least 1000 and the malicious actor is not the sole holder. This approach may bring an additional cost for the initial liquidity provider, but this cost is expected to be low and acceptable.

Recommendation Revise current deposit logic to defensively calculate the vault token amount

when it is being initialized. An alternative solution is to ensure a guarded launch process that safeguards the first deposit to avoid being manipulated.

Status The issue has been resolved as the team ensures that there is a decimal offset between the notes and underlying denominations. Also, the team plans to follow a guarded launch process and observe rollovers through the vault for 4 bond cycles (i.e., 28 days) before the public launch.

3.2 Improved Precision Calculation in `_computeRolloverAmt()`

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: `PerpetualTranche`
- Category: Numeric Errors [8]
- CWE subcategory: CWE-190 [2]

Description

The lack of `float` support in `Solidity` may introduce another subtle, but troublesome issue: precision loss. In this section, we examine one possible precision loss source that stems from the rollover operations on `PerpetualTranche`.

In particular, we show below the related `_computeRolloverAmt()` routine. As the name indicates, this routine computes the amount of reserve tokens that can be rolled out for the given amount of tranches deposited. We notice that when the token out balance is not covered, the current logic re-calculates the values by initially setting `tokenOutAmt` to be `tokenOutAmtRequested` (lines 963-971). And the calculation may introduce possible precision loss.

```

925     function _computeRolloverAmt(
926         ITranche trancheIn,
927         IERC20Upgradeable tokenOut,
928         uint256 trancheInAmtAvailable,
929         uint256 tokenOutAmtRequested
930     ) private view returns (IPerpetualTranche.RolloverPreview memory) {
931         IPerpetualTranche.RolloverPreview memory r;

933         uint256 trancheInDiscount = computeDiscount(trancheIn);
934         uint256 trancheOutDiscount = computeDiscount(tokenOut);
935         uint256 trancheInPrice = computePrice(trancheIn);
936         uint256 trancheOutPrice = computePrice(tokenOut);
937         uint256 tokenOutBalance = _reserveBalance(tokenOut);
938         tokenOutAmtRequested = MathUpgradeable.min(tokenOutAmtRequested, tokenOutBalance
939             );

940         if (trancheInDiscount == 0 & trancheOutDiscount == 0 & trancheInPrice == 0
941             & trancheOutPrice == 0) {

```

```

941         r.remainingTrancheInAmt = trancheInAmtAvailable;
942         return r;
943     }

945     r.trancheInAmt = trancheInAmtAvailable;
946     uint256 stdTrancheInAmt = _toStdTrancheAmt(trancheInAmtAvailable,
        trancheInDiscount);

948     // Basic rollover:
949     // (stdTrancheInAmt . trancheInPrice) = (stdTrancheOutAmt . trancheOutPrice)
950     uint256 stdTrancheOutAmt = stdTrancheInAmt.mulDiv(trancheInPrice,
        trancheOutPrice);
951     r.trancheOutAmt = _fromStdTrancheAmt(stdTrancheOutAmt, trancheOutDiscount);

953     // However, if the tokenOut is the mature tranche (held as naked collateral),
954     // we infer the tokenOut amount from the tranche denomination.
955     // (tokenOutAmt = collateralBalance * trancheOutAmt / matureTrancheBalance)
956     bool isMatureTrancheOut = _isMatureTranche(tokenOut);
957     r.tokenOutAmt = isMatureTrancheOut
        ? tokenOutBalance.mulDiv(r.trancheOutAmt, _matureTrancheBalance)
958       : r.trancheOutAmt;

961     // When the token out balance is NOT covered:
962     // we fix tokenOutAmt = tokenOutAmtRequested and back calculate other values
963     if (r.tokenOutAmt > tokenOutAmtRequested) {
964         r.tokenOutAmt = tokenOutAmtRequested;
965         r.trancheOutAmt = isMatureTrancheOut
            ? _matureTrancheBalance.mulDiv(r.tokenOutAmt, tokenOutBalance)
966           : r.tokenOutAmt;
967         stdTrancheOutAmt = _toStdTrancheAmt(r.trancheOutAmt, trancheOutDiscount);
968         stdTrancheInAmt = stdTrancheOutAmt.mulDiv(trancheOutPrice, trancheInPrice);
969         r.trancheInAmt = _fromStdTrancheAmt(stdTrancheInAmt, trancheInDiscount);
970     }
971 }

973     r.perpRolloverAmt = (stdTrancheOutAmt * trancheOutPrice).mulDiv(totalSupply(),
        _reserveValue());
974     r.remainingTrancheInAmt = trancheInAmtAvailable - r.trancheInAmt;
975     return r;
976 }

```

Listing 3.3: PerpetualTranche::_computeRolloverAmt()

Specifically, we notice the calculation of the resulting `stdTrancheInAmt = stdTrancheOutAmt.mulDiv(trancheOutPrice, trancheInPrice)` (line 969) takes a floor division, which apparently favors the trading user. The design here is better improved to take a ceiling division so that the computation is chosen in favor of the protocol. Note that the resulting precision loss may be just a small number, but it plays a critical role when certain boundary conditions are met.

Recommendation Revise the above calculations to better mitigate possible precision loss.

Status The issue has been fixed by this commit: 323acd9.

3.3 Accommodation of Non-ERC20-Compliant Tokens

- ID: PVE-003
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: RolloverVault, RouterV1
- Category: Coding Practices [7]
- CWE subcategory: CWE-1126 [1]

Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the `approve()` routine and analyze possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below. On its entry of `approve()`, there is a requirement, i.e., `require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)))`. This specific requirement essentially indicates the need of reducing the allowance to 0 first (by calling `approve(_spender, 0)`) if it is not, and then calling a second one to set the proper allowance. This requirement is in place to mitigate the known `approve()/transferFrom()` race condition (<https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729>).

```

194  /**
195   * @dev Approve the passed address to spend the specified amount of tokens on behalf
        of msg.sender.
196   * @param _spender The address which will spend the funds.
197   * @param _value The amount of tokens to be spent.
198   */
199   function approve(address _spender, uint _value) public onlyPayloadSize(2 * 32) {

201       // To change the approve amount you first have to reduce the addresses '
202       // allowance to zero by calling 'approve(_spender, 0)' if it is not
203       // already 0 to mitigate the race condition described here:
204       // https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729
205       require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)));

207       allowed[msg.sender][_spender] = _value;
208       Approval(msg.sender, _spender, _value);
209   }

```

Listing 3.4: USDT Token Contract

Because of that, a normal call to `approve()` is suggested to use the safe version, i.e., `safeApprove()`. In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. Similarly, there is a safe version of `transfer()` as well, i.e., `safeTransfer()`.

```

38  /**
39   * @dev Deprecated. This function has issues similar to the ones found in
40   * {IERC20-approve}, and its usage is discouraged.
41   *
42   * Whenever possible, use {safeIncreaseAllowance} and
43   * {safeDecreaseAllowance} instead.
44   */
45   function safeApprove(
46       IERC20 token,
47       address spender,
48       uint256 value
49   ) internal {
50       // safeApprove should only be called when setting an initial allowance,
51       // or when resetting it to zero. To increase and decrease it, use
52       // 'safeIncreaseAllowance' and 'safeDecreaseAllowance'
53       require(
54           (value == 0) || (token.allowance(address(this), spender) == 0),
55           "SafeERC20: approve from non-zero to non-zero allowance"
56       );
57       _callOptionalReturn(token, abi.encodeWithSelector(token.approve.selector,
58           spender, value));
59   }

```

Listing 3.5: SafeERC20::safeApprove()

In current implementation, if we examine the RouterV1::_checkAndApproveMax() routine that is designed to check if the spender has sufficient allowance. If not, there is a need to approve the maximum possible amount. To accommodate the specific idiosyncrasy, there is a need to make use of safeApprove() twice: the first time resets the allowance to be 0 and the second time sets the intended amount (line 307).

```

300   function _checkAndApproveMax(
301       IERC20Upgradeable token,
302       address spender,
303       uint256 amount
304   ) private {
305       uint256 allowance = token.allowance(address(this), spender);
306       if (allowance < amount) {
307           token.safeApprove(spender, type(uint256).max);
308       }
309   }

```

Listing 3.6: RouterV1::_checkAndApproveMax()

Recommendation Accommodate the above-mentioned idiosyncrasy about ERC20-related approve().

Status The issue has been fixed by this commit: 89b50a5.

3.4 Improved Validations Logic in RouterV1

- ID: PVE-004
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: RouterV1
- Category: Coding Practices [7]
- CWE subcategory: CWE-1126 [1]

Description

To facilitate the user interaction, the `Spot` protocol provides a `RouterV1` contract to dry-run and batch multiple operations. While examining one specific function `trancheAndDeposit()`, we notice it can benefit from improved validation.

Specifically, this function is proposed to tranche the collateral by using the current deposit bond and then depositing individual tranches to mint perp tokens. It currently takes four arguments and the first two are `perp` and `bond`. And there is an implicit assumption that the given `bond` needs to be the active deposit bond of `perp`. As a result, we can either validate this assumption or simply obtain the `bond` directly from `perp` (without passing it as an argument).

```

97     function trancheAndDeposit(
98         IPerpetualTranche perp,
99         IBondController bond,
100         uint256 collateralAmount,
101         uint256 feePaid
102     ) external afterPerpStateUpdate(perp) {
103         BondTranches memory bt = bond.getTranches();
104         IERC20Upgradeable collateralToken = IERC20Upgradeable(bond.collateralToken());
105         IERC20Upgradeable feeToken = perp.feeToken();
106
107         // transfers collateral & fees to router
108         collateralToken.safeTransferFrom(msg.sender, address(this), collateralAmount);
109         if (feePaid > 0) {
110             feeToken.safeTransferFrom(msg.sender, address(this), feePaid);
111         }
112
113         // approves collateral to be trached
114         _checkAndApproveMax(collateralToken, address(bond), collateralAmount);
115
116         // tranches collateral
117         bond.deposit(collateralAmount);
118         ...
119     }

```

Listing 3.7: RouterV1::trancheAndDeposit()

Recommendation Validate the given `bond` argument in the above routine to be the `deposit bond` of the given `perp`. Note another routine `trancheAndRollover()` shares the same issue.

Status The issue has been resolved as the `perp` contract validates whether the tranches entering the system are part of the current deposit bond.

3.5 Trust Issue of Admin Keys

- ID: PVE-005
- Severity: Low
- Likelihood: Low
- Impact: Medium
- Target: Multiple Contracts
- Category: Security Features [5]
- CWE subcategory: CWE-287 [3]

Description

In the `spot` protocol, there is a special administrative account, i.e., `owner`. This `owner` account plays a critical role in governing and regulating the system-wide operations (e.g., configure various settings and execute privileged operations). It also has the privilege to control or govern the flow of assets within the protocol contracts. In the following, we examine the privileged account and the related privileged accesses in current contracts.

```

392     function pause() public onlyKeeper {
393         _pause();
394     }

396     /// @notice Unpauses deposits, withdrawals and rollovers.
397     /// @dev NOTE: ERC-20 functions, like transfers will always remain operational.
398     function unpause() public onlyKeeper {
399         _unpause();
400     }

402     /// @notice Updates the reference to the keeper.
403     /// @param newKeeper The address of the new keeper.
404     function updateKeeper(address newKeeper) public virtual onlyOwner {
405         address prevKeeper = keeper;
406         keeper = newKeeper;
407         emit UpdatedKeeper(prevKeeper, newKeeper);
408     }

410     /// @notice Updates the authorized roller set.
411     /// @dev CAUTION: If the authorized roller set is empty, all rollers are authorized.
412     /// @param roller The address of the roller.
413     /// @param authorize If the roller is to be authorized or unauthorized.
414     function authorizeRoller(address roller, bool authorize) external onlyOwner {
415         if (authorize && !_rollers.contains(roller)) {

```

```

416         _rollers.add(roller);
417     } else if (!authorize && _rollers.contains(roller)) {
418         _rollers.remove(roller);
419     } else {
420         return;
421     }

423     emit UpdatedRollerAuthorization(roller, authorize);
424 }

426 /// @notice Update the reference to the bond issuer contract.
427 /// @param bondIssuer_ New bond issuer address.
428 function updateBondIssuer(IBondIssuer bondIssuer_) public onlyOwner {
429     if (address(bondIssuer_) == address(0)) {
430         revert UnacceptableReference();
431     }
432     if (address(_reserveAt(0)) != bondIssuer_.collateral()) {
433         revert InvalidCollateral(bondIssuer_.collateral(), address(_reserveAt(0)));
434     }
435     bondIssuer = bondIssuer_;
436     emit UpdatedBondIssuer(bondIssuer_);
437 }

```

Listing 3.8: Example Privileged Operations in `PerpetualTranche`

We understand the need of the privileged functions for proper contract operations, but at the same time the extra power to these privileged accounts may also be a counter-party risk to the contract users. Therefore, we list this concern as an issue here from the audit perspective and highly recommend making these privileges explicit or raising necessary awareness among protocol users.

Recommendation Promptly transfer the privileged account to the intended DAO-like governance contract. All changes to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status The issue has been mitigated as the team clarifies that the `owner` account will be managed by a multisig and eventually handed over to `Forth` DAO control.

4 | Conclusion

In this audit, we have analyzed the design and implementation of the SPOT protocol, which is a decentralized, inflation resistant store of value designed to be resilient in all market conditions. The system bends safely rather than breaking catastrophically in extreme market scenarios, and can forever resume its function without bailouts. SPOT can be held directly or rotated in as an alternative collateral asset to USDC within existing systems. It uses AMPL as the underlying unit of account, Buttonwood Tranche for collateral preparation, and onchain governance through the FORTH DAO. The current code base is clearly organized and those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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