What's up in the land of Linux kernel security!

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Hi! I'm Vaishali Thakkar.







What this talk is about?

History behind Linux as a OS

Clone of a UNIX operating system [early 1990s]

Core security model - Discretionary Access Control (DAC)

UNIX DAC Security Model

 DAC = Restricting access to objects based on the identity of subjects and/or groups to which they belong

 UNIX DAC = Allows the owner of an object (such as a file) to set the security policy for that object

 Superuser— an entity which bypasses Unix DAC policy for the purpose of managing the system.

Problems with UNIX DAC

Originally aimed at protection [rather than security] in multiuser systems

DAC does not protect against flawed or malicious code

Superuser == compromise on user's security policy

 Cannot express modern security requirements as lots of rights accessible by default

Problems with UNIX DAC

Users can invoke system services by switching to root user (setuid)

9 bits model (rwx per owner, group and others)

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Extension of UNIX DAC Features

POSIX ACCESS CONTROL LISTS

 Extension of abbreviated UNIX DAC ACLs, allows separate permissions for individual users and different groups

| Entry Type | Text form |
|--------------|--------------------|
| owner | user::rwx |
| Named user | user : name : rwx |
| Owning group | group : : rwx |
| Named group | group : name : rwx |
| Mask | mask :: rwx |
| Others | Other :: rwx |

POSIX Capabilities

Solution of a problem with superuser

An application requiring some privilege do not get all privileges

 A process has three sets of bitmaps called the inheritable(I), permitted(P), and effective(E) capabilities. Each capability is implemented as a bit in each of these bitmaps which is either set or unset.

Namespaces

Derived from the plan 9 operating system

Partitioning resources as seen by the process

Not a security feature but helps with implementing it

Namespaces

 Have been used to help implement multi-level security, where files are labeled with security classifications, and potentially entirely hidden from users without an appropriate security clearance

Not a security feature but helps with implementing it

Cryptography API

Used by kernel subsystems

 Provides support for a wide range of cryptographic algorithms and operating modes, including commonly deployed ciphers, hash functions, and limited support for asymmetric cryptography

 Key management subsystem for managing cryptographic keys within the kernel.

Cryptography API

 Who uses cryptography API: IPSec code, Disk encryption schemes, Kernel module signature verification

Support for hardware-based cryptographic features is growing too.

Network Security

 <u>Netfilter</u>: An IP network layer framework which hooks packets which pass into, through and from the system.

 <u>Iptables</u>: A module which implements an IPv4 firewalling scheme, managed via the userland iptables tool.

 <u>Ebtables</u>: Provides filtering at the link layer, and is used to implement access control for Linux bridges

Network Security

Arptables: provides filtering of ARP packets

 <u>IPSec</u>: A network protocol suite which authenticates and encrypts the packets of data sent over a network

Linux Security Modules [LSMs]

 Linux Security Modules (LSM) API implements hooks at all security-critical points within the kernel

 A user of the framework (LSM) can register with the API and receive callbacks from these hooks

 Was designed to provide the specific needs of everything needed to successfully implement MAC [Mandatory Access control]

SELinux

A LSM which provides a mechanism for supporting access control policies

 In SELinux, all objects on the system, are assigned security labels. All security-relevant interactions between entities on the system are hooked by LSM and passed to the SELinux module, which consults its security policy to determine whether the operation should continue.

Security policy is loaded from userland, and can be modified

Smack

 A LSM which was designed to provide a simple form of MAC security, in response to the relative complexity of SELinux

Works best with file-systems which supports extended attributes

It's a part of the Tizen security architecture

AppArmor

 Fundamentally different MAC scheme to SELinux and Smack, no direct labeling and security policy is applied to pathnames

 Allows the system administrator to restrict programs' capabilities with per-program profiles

 Also features a learning mode, where the security behavior of the application is observed and converted automatically into a security profile

TOMOYO

Another MAC scheme, which implements path based security

 Utilizes the learning mode similar to AppArmor where behavior of the system is observed to enhance the security policy

It records the trees of process invocation, described as domains

YAMA

Collection of DAC security enhancements from projects like Grsecurity

Enhanced restrictions on ptrace are implemented in YAMA

LoadPin

 Fairly new LSM, which ensures that all kernel loaded files are loaded from trusted device [dm-verity or CDROM]

 Allows systems that have a verified and/or unchangeable filesystem to enforce module and firmware loading restrictions without needing to sign the files individually.

Audit Subsystem

 Was first designed to meet government certification requirements, now used by LSMs and other security components

Helps to track security relevant information

Seccomp

A mechanism which restricts access to system calls by processes

 Reduce the attack surface of the kernel by preventing applications from entering system calls they don't need

 The original seccomp code, also known as "mode 1", provided access to only four system calls: read, write, exit, and sigreturn

Seccomp - bpf

 An arbitrary specification of which system calls are permitted for a process, and integration with audit logging

Was developed for use as part of the Google Chrome OS.

Integrity management subsystem

Used to maintain the integrity of files on the system

 Integrity Measurement Architecture component performs runtime integrity measurements of files using cryptographic hashes, comparing them with a list of valid hashes

 Dm-verity module: Device mapper target which manages file integrity at block level

Is this level of security sufficient?

What are the possible solutions?

Bug Fixing Using Tools

Wide scope of research projects and static/dynamic analysis tools

Useful only if it is used regularly to detect the security issues

Automatic testing helps, but have limitations

• Sparse:

- Written by Linus Torvalds, provides a set of annotations designed to convey semantic information about types
- Warns about unsupported operations or type mismatches with restricted integer types
- Warns about any non-static variable or function definition that has no previous declaration.

• Smatch:

- Written by Dan Carpenter, more than 3000 bug fixes so far
- Warns about issues like null pointer dereference, error pointer dereference, uninitialized data, information leak, some cases of use after free, double free, unnecessary/missing null check etc

Coccinelle:

- Pattern matching and transformation tool, developed by Julia Lawall,
 more than 4000 bug fixes so far
- Handles few security issues like null pointer dereference, use of sleeping functions under locks, use after free, few locking related bugs etc

GCC and GCC plugins:

- New GCC versions [6 and 7] has added many new warning options, though sometimes they are added on case basis and might not be able to handle all kind of cases for a particular bug classes
- GCC plugins helps with handling specific kind of bug classes at compiler level, without adding code in compiler itself
- As of now, there are 5 gcc plugins added in the Linux kernel
- Plugins like randomizing structures layout at compile time or detecting any structures that contain __user attributes and makes sure it is being fulling initialized
- helps with making the attack surface harder

• Fuzzers:

- <u>Trinity</u>: Developed by dave Jones, helps with OOPS, locking related bugs and memory leaks etc
- <u>Syzkaller</u>: Developed by Dmitry Vyukov and a team [Google], helps with {resource, memory, information} leaks, deadlocks etc
- Some others: AFL [American Fuzzy Loop], Address sanitizer, Thread sanitizer etc

Is Fixing bugs sufficient?

Kernel Self Protection [KSPP]

Idea: Bugs have longer lifetime, kernel should be able to protect itself

Kill classes of bugs instead of individual bug

Current focus on upstreaming grsecurity/PAX features

More information on Kees cook's blog:
 https://kernsec.org/wiki/index.php/Feature_List

Conclusion

 We have come far from UNIX security but there is always a scope of more research and improvement at the kernel level.

 With the advancement of technology and wide variety of requirements, security is no longer a buzzword.

Resources

- LWN: http://lwn.net/Security
- Kernel Self Protection project:
 https://kernsec.org/wiki/index.php/Kernel_Self_Protection_Project
- LSM mailing list and kernel-hardening mailing list
- Kernel security summit [2016]:
 https://www.youtube.com/playlist?list=PLbzoR-pLrL6pq6qCHZUuhbXsTsyz1
 https://www.youtube.com/playlist?list=PLbzoR-pLrL6pq6qCHZUuhbXsTsyz1
 https://www.youtube.com/playlist?list=PLbzoR-pLrL6pq6qCHZUuhbXsTsyz1