COCCINELLE: FINDING BUGS IN THE LINUX KERNEL

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Who Am I?

- Linux Kernel developer at Oracle and associated with RGSoC as a co-organizer
- Working in kernel security engineering group
- Interested in many different subsystems of the Linux Kernel

Prerequisites

- Source code of the Linux kernel
- Latest version of the Coccinelle
 - Either install it from the package manager [Coccinelle is available with around 10 linux distros including Fedora, Ubuntu, Debian, ArchLinux etc.].
 - Or build it from the source. (https://github.com/coccinelle/coccinelle)

Code Maintenance Issues

• Software evolution:

- Refactoring code to use newer APIs
- Need to find all parts of the code that need updating
- Process should be fast, reliable and systematic
- However, things are never straightforward

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Software robustness:

- Are the programmers following the standards?
- Is the code accounting for all errors that can take place?
- Is the written code overly defensive?

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- Are the programmers following the standards?
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- Is the written code overly defensive?

The Human Factor:

Copy-pasting the code and mistakes can always happen

Example 1: Software evolution - code refactoring

Bit masking is preferrably done using the BIT macro

```
- BUILD_BUG_ON(max >= (1 << 16));
+ BUILD_BUG_ON(max >= (BIT(16)));
```

Example 2: Some useless code

```
diff --git a/drivers/staging/dgnc/dgnc tty.c b/drivers/staging/dgnc/dgnc tty.c
index 0e903dc..a4e6c9e 100644
--- a/drivers/staging/dgnc/dgnc ttv.c
+++ b/drivers/staging/dgnc/dgnc_tty.c
@@ -2286,7 +2286,6 @@ static inline int dgnc get mstat(struct channel t *ch)
 static int dgnc_get_modem_info(struct channel_t *ch, unsigned int __user *value
 {
        int result;
        int rc;
        if (!ch || ch->magic != DGNC_CHANNEL_MAGIC)
                return -ENXIO;
@@ -2296,9 +2295,7 @@ static int dgnc_get_modem_info(struct channel_t *ch,
        if (result < 0)</pre>
                return -ENXIO;
        rc = put_user(result, value);
        return rc;
        return put_user(result, value);
 }
```

Example 3: Security related issue

```
static long hwdep_read_locked(...)
         if (copy_to_user(buf, &event, count))
                  return -EFAULT;
          . . .
static long hwdep_read(...)
{
         while (...) {
                  spin_lock_irq (...);
          }
         if (efw->dev_lock_changed)
                  count = hwdep_read_locked(...);
          spin_unlock_irq (...);
          \dots \dots
```

Example 3: Security related issue

```
static long hwdep_read_locked(...)
         if (copy_to_user(buf, &event, count))
                 return -EFAULT;
         . . .
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         }
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                 count = hwdep_read_locked(...);
         spin_unlock_irq(...);
```

Commit id=04b2d9c9c319277ad4fbbb71855c256a9f4d5f98

Some interesting projects/use cases

- Kees cook's recent timer API changes
- Thomas Gleixner's talk on refactoring the Linux Kernel: https://www.youtube.com/watch?v=mxxicJZ8cis
- Kernel backporting project by Luis and Jmake by Lip6 researchers

Coccinelle

- Program matching and transformation tool
- Independent of the compilation process
- Very intuitive patch like style
- Used by several communities:
 - Linux Kernel: 5K+ patches
 - QEMU: 200+ patches
 - systemd: 80+ patches

Semantic Patch Language (SmPL)

- Abstract C-like grammar
- Independent of the compilation process
- Metavariables are used to abstract over sub-terms in code
 - If an expression matches within a pattern, it can be tracked throughout its presence in the code e.g. variable names, typedefs
- "..." is used to abstract over code sequences
 - Used as don't care
 - Variants are used as syntactic sugar for + and ? in regular expressions
- Lines can be annotated with {-,+,*}
 - Transformations are described using patch-like style (-/+)
 - Matching employs *

Example: Using BIT macro

• Bit masking is preferrably done using the BIT macro

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```
- 1 << 16
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• Code we should focus on for building a semantic patch:

```
- 1 << 16
+ BIT(16)
```

• Is 16 important here?

Do we care about number of shifts?

```
if (opts & (1 << REISERFS_LARGETAIL))

if (opts & (BIT(REISERFS_LARGETAIL)))</pre>
```

Do we care about number of shifts?

```
- if (opts & (1 << REISERFS_LARGETAIL))
+ if (opts & (BIT(REISERFS_LARGETAIL)))</pre>
```

Use metavariables

```
@@
constant c;
@@
-1 << c
+BIT(c)
```

- Constant will capture numbers and defined constants
- What if we had something like

```
1 << (31 - inode->i_sb->s_blocksize_bits)
```

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```
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```

expression to the rescue

```
@@
expression E;
@@
-1 << E
+BIT(E)</pre>
```

```
Example: x->y = m->n + 1;
```

 Constant: match patterns on values and constants
 e.g. numbers like 2,3 and defined constants in a code

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```

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- Expression: match patterns on constants and complex subterms
 e.g. struct->elem, x-y, func(arg) etc

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- Constant: match patterns on values and constants
 e.g. numbers like 2,3 and defined constants in a code
- Expression: match patterns on constants and complex subterms
 e.g. struct->elem, x-y, func(arg) etc.
- Identifier: a structure field, a macro, a function, or a variable
- **Statement:** match patterns which do not return a value e.g. if, while, break etc

- Constant: match patterns on values and constants
 e.g. numbers like 2,3 and defined constants in a code
- Expression: match patterns on constants and complex subterms
 e.g. struct->elem, x-y, func(arg)
- Identifier: a structure field, a macro, a function, or a variable
- **Statement:** match patterns which do not return a value e.g. if, while, break etc
- **Type:** match patterns for the type of variables/functions e.gint, boolean, float etc

Transformation specification

- in the leftmost column for something to remove
- + in the leftmost column for something to add
- * in the leftmost column for something of interest
 - Cannot be used with + and -.
- Spaces, newlines that are irrelevant.

Spatch

- Coccinelle's command-line tool
- To check that your semantic patch is valid:

```
spatch --parse-cocci mysp.cocci
```

• To run your semantic patch:

```
spatch --sp-file mysp.cocci file.c
```

spatch --sp-file mysp.cocci --dir directory

Exercise 1

- Save the semantic patch to bitmask.cocci. [slide 12 and 14]
- Run it using spatch on any particular directory or on whole kernel.
 spatch --sp-file bitmask.cocci --dir directory
- Redirect results to an output file for an inspection.
- Is it ok to use BIT macro in every case? Should we want to restrict it for the files which are already using it?

Exercise 2

- Parentheses are not needed around the bitwise left shift operations like in u32 val = (1 << 31);.
- Write a semantic patch to remove these parentheses.
- Run the semantic patch over the directory drivers/net/wireless/.
- Some other cases to think about:
 - Extra parentheses around the function arguments
 - Using the same identifier on the left and right side of the assignment

Using BIT macro (Revisited)

Example:

```
diff -u -p a/arch/mips/pci/pci-mt7620.c b/arch/mips/pci/pci-mt7620.c
--- a/arch/mips/pci/pci-mt7620.c
+++ b/arch/mips/pci/pci-mt7620.c
@@ -37,11 +37,11 @@
#define PDRV SW SET
                                          BIT(23)
#define PPLL DRV
                                          0xa0
-#define PDRV SW SET
                                          (1 << 31)
-#define LC CKDRVPD
                                          (1 << 19)
-#define LC CKDRVOHZ
                                          (1 << 18)
-#define LC CKDRVHZ
                                          (1 << 17)
-#define LC CKTEST
                                          (1 << 16)
+#define PDRV SW SET
                                          (BIT(31))
+#define LC CKDRVPD
                                          (BIT(19))
+#define LC CKDRVOHZ
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+#define LC CKDRVHZ
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+++ b/arch/mips/pci/pci-mt7620.c
@@ -37,11 +37,11 @@
#define PDRV_SW_SET
                                          BIT(23)
#define PPLL DRV
                                          0xa0
-#define PDRV SW SET
                                          (1 << 31)
-#define LC CKDRVPD
                                          (1 << 19)
-#define LC CKDRVOHZ
                                          (1 << 18)
-#define LC CKDRVHZ
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-#define LC CKTEST
                                          (1 << 16)
+#define PDRV SW SET
                                          (BIT(31))
+#define LC CKDRVPD
                                          (BIT(19))
+#define LC CKDRVOHZ
                                          (BIT(18))
+#define LC CKDRVHZ
                                          (BIT(17))
+#define LC CKTEST
                                          (BIT(16))
```

 Would like to restrict the bitmask semantic patch to files that are already using the BIT macro?

Using BIT macro (Revisited)

Example:

```
-#define LC_CKDRVPD (1<<19)
-#define LC_CKDRVOHZ (1<<18)
+#define LC_CKDRVPD (BIT(19))
+#define LC_CKDRVOHZ (BIT(18))
```

Semantic patch:

```
@usesbit@
@@
BIT(...)

@depends on usesbit@
expression E;
@@
- 1 << E
+ BIT(E)</pre>
```

Isomorphism

- Coccinelle captures code as defined in your rule
- Valid variants of your defined pattern can exist
- Cumbersome to list them all in your rule/s
- Examples:
 - \blacksquare x == NULL and !x
 - sizeof(struct i) * e ande * sizeof(struct i)
- Isomorphisms can handle such variations
- Rules defining isomorphisms exist in standard.iso

Isomorphism Examples

Example 1:

```
Expression
@ is_null @
expression X;
@@
X == NULL <=> NULL == X => !X
```

Example 2:

```
Expression
@ drop_cast @
expression E;
pure type T;
@@
(T)E => E
```

Exercise 3

- To avoid code duplication or error prone code, the kernel provides macros such as DIV_ROUND_UP.
- The definition of the DIV_ROUND_UP goes like this:
 DIV_ROUND_UP(n,d) (((n) + (d) 1) / (d))
- Write the semantic patch for replacing the pattern (((n) + (d) 1) / (d)) with DIV_ROUND_UP.
- Redirect results to an output file for an inspection.

Exercise 4

- Consider the example of DIV_ROUND_UP.
- The macro is defined in linux/kernel.h. So, it depends on this header file.
- Expand the semantic patch you wrote in exercise 3 using 'depends'
 on'.
- Review the output given by updated semantic patch.

Example: setup_timer

 The function setup_timer combines the initialization of a timer with the initialization of the timer's function and data fields.

```
init_timer(&cf->timer);
cf->timer.function = omap_cf_timer;
cf->timer.data = (unsigned long) cf;
setup_timer(&cf->timer, omap_cf_timer, (unsigned long)cf);
```

- Why setup_timer?
- How Coccinelle can help here?

Warning: Do not send patches for this change in the kernel, refactoring of better timer interface is in process

setup_timer: case one

Example:

```
@@

@@
- init_timer(&cf->timer);
- cf->timer.function = omap_cf_timer;
- cf->timer.data = (unsigned long) cf;
+ setup_timer(&cf->timer, omap_cf_timer, (unsigned long)cf);
```

Semantic patch

```
@case_one@
expression e,func,da;
@@

- init_timer (&e);
+ setup_timer (&e, func, da);
- e.function = func;
- e.data = da;
```

setup_timer: case one

Semantic patch:

```
@case_one@
expression e,func,da;
@@

- init_timer (&e);
+ setup_timer (&e, func, da);
- e.function = func;
- e.data = da;
```

- Is this the only case where we can use setup_timer?
- Is it necessary that the call to init_and the initialization of the function and data fields always occur in the order shown in the example?

setup_timer: case two

Example:

```
- init_timer(&hose->err_timer);
- hose->err_timer.data = (unsigned long)hose;
- hose->err_timer.function = pcibios_enable_err;
+ setup_timer(&hose->err_timer, pcibios_enable_err, (unsigned long)hose);
```

Semantic patch:

```
@case_two@
expression e,func,da;
@@

- init_timer (&e);
+ setup_timer (&e, func, da);
- e.data = da;
- e.function = func;
```

setup_timer: comparing both cases

Case one:

```
@case_one@
expression e,func,da;
@@

- init_timer (&e);
+ setup_timer (&e, func, da);
- e.function = func;
- e.data = da;
```

Case two:

```
@case_two@
expression e,func,da;
@@

-init_timer (&e);
+setup_timer (&e, func, da);
-e.data = da;
-e.function = func;
```

Disjunctions

- A sequence of patterns between (... | ...).
- Patterns checked in order and the first that matches is chosen.
- Combining case one and case two in our example:

```
@case_one_and_two@
expression e, func, da;
@@

-init_timer (&e);
+setup_timer (&e, func, da);

(
-e.function = func;
-e.data = da;
|
-e.data = da;
-e.function = func;
)
```

Exercise 5

- Implement the semantic patches for both cases of the setup_timer.
 Compare the results.
- Implement the rule combining case one and case two using disjunction.
- Think about why do we need to use disjunctions? Can we use multiple rules?
- Check the results. Does it cover all the cases that were matched by the separate rules?
- Grep for the init_timer and check if the rule with disjunction covers everything?

setup_timer(Contd.)

Example:

```
init_timer (&np->timer);
np->timer.expires = jiffies + 1*HZ;
np->timer.data = (unsigned long) dev;
np->timer.function = rio_timer;
add_timer (&np->timer);
```

- Does previous rule covered all cases?
- Is it necessary that the call to init_timer and the initialization of the function & the data field always occurs in a contiguous manner?

Dots

Problem:

Sometimes it is necessary to search for multiple related code fragments.

Solution:

- Specify patterns consisting of the fragments of code separated by arbitrary execution paths.
- Specify constraints on the contents of those execution paths.

setup_timer: case three

Semantic patch:

```
@case_three@
expression e,func,da;
@@
- init_timer (&e);
+ setup_timer (&e, func, da);
...
- e.data = da;
- e.function = func;
```

Example:

```
init_timer (&np->timer);

setup_timer(&np->timer, rio_timer, (unsigned long)dev);

np->timer.expires = jiffies + 1*HZ;

np->timer.data = (unsigned long) dev;

np->timer.function = rio_timer;

add_timer (&np->timer);
```

Using dots

Semantic patch:

```
@case_three@
expression e,func,da;
@@
- init_timer (&e);
+ setup_timer (&e, func, da);
...
- e.data = da;
- e.function = func;
```

- '...' matches all possible execution paths from the pattern before to the pattern after
- The patterns before and after cannot appear in the region matched by "..." (shortest path principle).

Example: Compressing lines for immediate return

• In the following code last two lines could be compressed into one:

```
int bytes_written;
u16 link_speed;
link_speed = rtw_get_cur_max_rate(padapter) / 10;
bytes_written = snprintf(command, total_len, "LinkSpeed %d", link_speed);
return bytes_written;
```

Compressing lines for immediate return

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return bytes_written;
```

```
int bytes_written;
u16 link_speed;
link_speed = rtw_get_cur_max_rate(padapter) / 10;
return snprintf(command, total_len, "LinkSpeed %d", link_speed);
```

Dots: Compressing lines for immediate return

Example:

Semantic patch:

```
@@
expression r;
identifier f;
@@
-r = f(...)
+return
    f(...);
-return r;
```

Exercise 6

- Implement the rule for case three of setup_timer using dots. [Slide 40]
- Run the patch over the kernel code and investigate the result.
- Think about the case three like pattern for the case two.
- Implement the rule for those kind of patterns.
- Try to limit the number of rules.

Exercise 6(Contd.)

Example:

```
init_timer(&sharpsl_pm.ac_timer);
sharpsl_pm.ac_timer.function = sharpsl_ac_timer;

init_timer(&sharpsl_pm.chrg_full_timer);
sharpsl_pm.chrg_full_timer.function = sharpsl_chrg_full_timer;
```

- Is it even necessary that the initialization of the data field always occurs?
- Expand the semantic patch to include such cases.

Exercise 7

Example:

```
int bytes_written;
u16 link_speed;
link_speed = rtw_get_cur_max_rate(padapter) / 10;
return snprintf(command, total_len, "LinkSpeed %d", link_speed);
```

- Do we really need the variable bytes_written after compressing the lines?
- Expand the semantic patch[slide 44] to remove the variable along with compressing lines.

Hint: Ensure that the variable is not used anywhere else.

Using dots(Contd.)

Semantic patch:

```
@case_three@
expression e,func,da;
@@
- init_timer (&e);
+ setup_timer (&e, func, da);
...
- e.data = da;
- e.function = func;
```

- Check the properties of the matched statement sequence
- Does the rule look correct? Or do we need to ensure something?

Using dots with when

 Dots can be modified with a when clause, indicating a pattern that should not occur

```
@case_three@
expression e1, e2, e3, e4, func, da;
@@

-init_timer(&e1);
+setup_timer(&e1, func, da);

... when != func = e2
    when != da = e3

-e1.data = da;
-e1.function = func;
```

when

- Keyword used to indicate conditions on execution path
- As seen before, controls the behavior of "..."
- Can be coupled with:
 - strict: force condition on every execution path (including failures)
 - forall: force condition on every execution path (excluding failures)
 - exists: is there an execution path that matches the pattern?
 - any: allow the patterns specified...
 - conditions specified by the user

More use of dots

- Two possible modifiers to the control flow for ellipses:
 - 1. <...P...> indicates that matching the pattern in between the ellipses is optional
 - 2. <+...P...+> indicates that the pattern in between the ellipses must be matched at least once, on some control-flow path.
 - The + is intended to be reminiscent of the + used in regular expressions.

More use of dots(Contd.)

Example:

```
@r@
@@
-if (...) {
<+...
   return ...;
   ...+>
}
```

Meaning:

• To remove all ifs that contain at least one return.

More use of dots(Contd.)

Example:

```
@r@
@@
-if (...) {
<...
   return ...;
   ...>
}
```

Meaning:

• To remove all ifs

Exercise 8

- 1. Implement the example of 'compression of lines for the immediate return problem'.
- 2. The semantic patch for removing unused variables only matches a variable declaration when the declaration does not initialize the variable.
- 3. Extend the complete semantic patch so that it also removes unused variables that are initialized to a constant.

Exercise 9

In the following code, when x has any pointer type, the cast to u8 *, or to any other pointer type is not needed.

```
kfree((u8 *)x);
```

- Write a semantic patch to remove such casts.
- Consider generalizing your semantic patch to functions other than kfree.
- Are there any patterns that can benefit from using disjunctions?

Coccicheck

- A Coccinelle-specific target which is defined in the top level Makefile.
- Four basic modes
 - Patch mode
 - Context mode
 - Org mode
 - Report mode
- Default output: Report mode
- Command that can be used for specifying particular mode: make coccicheck MODE=patch

- Four basic modes
 - Patch mode: proposes a fix when possible.

- Four basic modes
 - Context mode:
 - 1. highlights lines of interest and their context in a diff-like style.
 - 2. Lines of interest are indicated with '-'.

```
@@ -582,8 +580,7 @@ static int iss_net_configure(int index,
return 1;
}

-     init_timer(&lp->tl);
-     lp->tl.function = iss_net_user_timer_expire;
-     setup_timer(&lp->tl, iss_net_user_timer_expire, OUL);
return 0;
```

- Four basic modes
 - Org mode: Generates a report in the Org mode format of Emacs.

```
* TODO [[view:/home/linux-next/linux/arch/sh/drivers/pci/common.c::face=ovl-fa::cole=12] [Use setup_timer function.]]
[[view:/home/linux-next/linux/arch/sh/drivers/pci/common.c::face=ovl-face1::li
[/home/linux-next/linux/arch/sh/drivers/pci/common.c::109]]

* TODO [[view:/home/linux-next/linux/arch/sh/drivers/pci/common.c::face=ovl-fa::cole=12] [Use setup_timer function.]]
[[view:/home/linux-next/linux/arch/sh/drivers/pci/common.c::face=ovl-face1::li
[/home/linux-next/linux/arch/sh/drivers/pci/common.c::115]]
```

Four basic modes

 Report mode: Generates a list in the following format file:line:column-column: message

```
/home/linux-next/linux/arch/sh/drivers/pci/common.c:108:2-12: Use setup_timer /home/linux-next/linux/arch/sh/drivers/pci/common.c:114:2-12: Use setup_timer /home/linux-next/linux/arch/sh/drivers/push-switch.c:81:1-11: Use setup_timer /home/linux-next/linux/arch/x86/kernel/pci-calgary_64.c:1010:1-11: Use setup_t line 1011. /home/linux-next/linux/arch/powerpc/oprofile/op_model_cell.c:682:1-11: Use set line 683.
```

setup_timer again

Problem:

- What if init_timer is called in one function and data field is initialized in another function?
- Will it be safe to use setup_timer in that case?

Solution:

How about giving warning in such cases?

setup_timer again

We need two rules to match both parts

Semantic patch:

```
@r1@
identifier f;
@@
f(...) { ...
  init_timer(...)
  . . .
@r2@
identifier g;
struct timer_list t;
expression e;
@@
g(...) { ...
t.data = e
 . . .
```

setup_timer again

• We want to match 2 different functions. So, let's avoid function name overriding.

Semantic patch:

```
@r1 exists@
identifier f;
@@
f(...) { ...
  init_timer(...)
  . . .
@r2 exists@
identifier g != r1.f;
struct timer_list t;
expression e;
@@
g(...) { ...
t.data = e
 . . .
```

Position variables

- Position metavariables can be used to store the position of any token, for later matching or printing.
- In the case of setup_timer we want to use the position of init_timer so that Coccinelle can give warning at such code.

Position variables

Example:

```
@r1 exists@
identifier f;
position p;
@@
f(...) { ...
  init_timer@p(...)
  . . .
@r2 exists@
identifier g != r1.f;
struct timer_list t;
expression e8;
@@
g(...) { ...
  t.data = e8
```

Embedding python script

- Coccinelle can embed Python code. Python code is used inside special SmPL rule annotated with script:python.
- Python rules inherit metavariables, such as identifier or token positions, from other SmPL rules.
- The inherited metavariables can then be manipulated by Python code.

Python script with the warning

Example:

```
@r1 exists@
identifier f;
position p;
@@
f(...) { ...
  init_timer@p(...)
@r2 exists@
identifier g != r1.f;
struct timer_list t;
expression e;
@@
g(...) { ...
  t.data = e
@script:python depends on r2@
p << r1.p;
@@
print "Data field initialized in another function. Dangerous to use
setup_timer %s:%s" % (p[0].file,p[0].line)
```

Python script without printing warning

Example:

```
@r1 exists@
identifier f;
position p;
@@
f(...) { ...
  init_timer@p(...)
@r2 exists@
identifier g != r1.f;
struct timer_list t;
expression e;
@@
g(...) { ...
  t.data = e
@script:python depends on r2@
p << r1.p;
cocci.include_match(False)
```

Exercise 10

- When searching for things, rather than transforming them, it may be useful to generate the output in a variety of formats. This can be done using the interface to python (ocaml is also available).
- Position variables are useful in this context, because they provide the file name and line number of various program elements.

Exercise 10 (Contd.)

Consider the following patch discussed earlier:

```
@@
expression r;
identifier f;
@@
-r = f(...)
+return
    f(...);
-return r;
```

 Following python code is intended to print the file name and line numbers of the assignment and erroneous test, respectively:

```
@script:python@
p1 << r.p1; // inherit a metavariable p1 from rule r
p2 << r.p2; // inherit a metavariable p2 from rule r
@@
print p1[0].file, p1[0].line, p2[0].line</pre>
```

Exercise 10 (Contd.)

Do this:

- Create a semantic patch consisting of the original patch rule shown on the previous page followed by the python code specified in the last slide.
- Give name r to the rule and remove the transfromation.
- Add position variables p1 and p2.
- Attach position variables to the relevant code.
- Test the semantic patch and investigate the results.

Exercise 11

- We have seen that * can be used to highlight items of interest.
- Repeat the previous exercise, this time without using python, but instead annotate the original code pattern with * rather than performing transformations.
- How is the result different than the result produced when using python?

Exercise 12

- Implement the setup_timer case with the python code.
- Combine all rules in a single script and then try to run it. Observe how output changes.
- Try to reorder the rules in a semantic patch and then observe the changes.
- Do we also need a rule for the immediate call of init_timer, intialization of data and function fields? If yes, then why? If no, then why?

Hint: Consider performance and speed of the semantic patch.

Feature summary

- Metavariables and Isomorphisams
- Different uses of ...
- When
- Named rules and metavariable inheritance
- Position variables
- Scripting through Python/Ocaml
- Different modes for the Coccinelle script

Useful links

- Source code of the Coccinelle: "https://github.com/coccinelle/coccinelle"
- Grammar and features: "http://coccinelle.lip6.fr/docs/options.pdf"
- Documentation: "Documentation/coccinelle.txt"
- Project: "http://coccinelle.lip6.fr/"
- Spgen: "https://github.com/coccinelle/coccinelle/tree/master/tools/spgen"

THANK YOU!