

## 5 Reference

For your reference, here is the syntax, large-step operational semantics, and typing rules of SIMPL that should be the basis of your type-checker implementation.

### 5.1 Syntax of SIMPL

commands	$c ::= \text{skip} \mid c_1; c_2 \mid v := e \mid \text{if } e \text{ then } c_1 \text{ else } c_2$ $\mid \text{while } e \text{ do } c \mid \text{int } v \mid \text{bool } v$
expressions	$e ::= n \mid \text{true} \mid \text{false} \mid v \mid e_1 \text{ aop } e_2 \mid e_1 \text{ bop } e_2 \mid e_1 \leq e_2 \mid !e$
arithmetic operators	$\text{aop} ::= + \mid - \mid *$
boolean operators	$\text{bop} ::= \&\& \mid \mid\mid$
variable names	$v$
integer constants	$n$
types	$\tau ::= \text{int} \mid \text{bool}$
stores	$\sigma : v \mapsto \mathbb{Z} \cup \{T, F\}$
typing contexts	$\Gamma : v \mapsto \tau \times \{T, F\}$

### 5.2 Large-step Semantics of SIMPL

#### 5.2.1 Commands

$$\begin{array}{ll}
 \langle \text{skip}, \sigma \rangle \Downarrow \sigma & (1) \\
 \langle \text{int } v, \sigma \rangle \Downarrow \sigma & (2) \\
 \langle \text{bool } v, \sigma \rangle \Downarrow \sigma & (3) \\
 \frac{\langle c_1, \sigma \rangle \Downarrow \sigma_2 \quad \langle c_2, \sigma_2 \rangle \Downarrow \sigma'}{\langle c_1; c_2, \sigma \rangle \Downarrow \sigma'} & (4) \\
 \frac{\langle e, \sigma \rangle \Downarrow n}{\langle v := e, \sigma \rangle \Downarrow \sigma[v \mapsto n]} & (5) \\
 \frac{\langle e, \sigma \rangle \Downarrow p}{\langle v := e, \sigma \rangle \Downarrow \sigma[v \mapsto p]} & (6) \\
 \frac{\langle e, \sigma \rangle \Downarrow T \quad \langle c_1, \sigma \rangle \Downarrow \sigma'}{\langle \text{if } e \text{ then } c_1 \text{ else } c_2, \sigma \rangle \Downarrow \sigma'} & (7) \\
 \frac{\langle e, \sigma \rangle \Downarrow F \quad \langle c_2, \sigma \rangle \Downarrow \sigma'}{\langle \text{if } e \text{ then } c_1 \text{ else } c_2, \sigma \rangle \Downarrow \sigma'} & (8) \\
 \frac{\langle \text{if } e \text{ then } (c; \text{while } e \text{ do } c) \text{ else skip}, \sigma \rangle \Downarrow \sigma'}{\langle \text{while } e \text{ do } c, \sigma \rangle \Downarrow \sigma'} & (9)
 \end{array}$$

### 5.2.2 Expressions

$$\langle n, \sigma \rangle \Downarrow n \quad (10)$$

$$\langle \text{true}, \sigma \rangle \Downarrow T \quad (11)$$

$$\langle \text{false}, \sigma \rangle \Downarrow F \quad (12)$$

$$\langle v, \sigma \rangle \Downarrow \sigma(v) \quad (13)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow n_1 \quad \langle e_2, \sigma \rangle \Downarrow n_2}{\langle e_1 \leq e_2, \sigma \rangle \Downarrow n_1 \leq n_2} \quad (14)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow p \quad \langle e_2, \sigma \rangle \Downarrow q}{\langle e_1 \&\& e_2, \sigma \rangle \Downarrow p \wedge q} \quad (15)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow p \quad \langle e_2, \sigma \rangle \Downarrow q}{\langle e_1 \mid\mid e_2, \sigma \rangle \Downarrow p \vee q} \quad (16)$$

$$\frac{\langle e, \sigma \rangle \Downarrow p}{\langle !e, \sigma \rangle \Downarrow \neg p} \quad (17)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow n_1 \quad \langle e_2, \sigma \rangle \Downarrow n_2}{\langle e_1 + e_2, \sigma \rangle \Downarrow n_1 + n_2} \quad (18)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow n_1 \quad \langle e_2, \sigma \rangle \Downarrow n_2}{\langle e_1 - e_2, \sigma \rangle \Downarrow n_1 - n_2} \quad (19)$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow n_1 \quad \langle e_2, \sigma \rangle \Downarrow n_2}{\langle e_1 * e_2, \sigma \rangle \Downarrow n_1 n_2} \quad (20)$$

## 5.3 Typing Rules for SIMPL

### 5.3.1 Commands

$$\Gamma \vdash \text{skip} : \Gamma \quad (21)$$

$$\frac{v \notin \Gamma^{\leftarrow}}{\Gamma \vdash \text{int } v : \Gamma[v \mapsto (\text{int}, F)]} \quad (22)$$

$$\frac{v \notin \Gamma^{\leftarrow}}{\Gamma \vdash \text{bool } v : \Gamma[v \mapsto (\text{bool}, F)]} \quad (23)$$

$$\frac{\Gamma \vdash c_1 : \Gamma_2 \quad \Gamma_2 \vdash c_2 : \Gamma'}{\Gamma \vdash c_1; c_2 : \Gamma'} \quad (24)$$

$$\frac{\Gamma \vdash e : \tau \quad \Gamma(v) = (\tau, p)}{\Gamma \vdash v := e : \Gamma[v \mapsto (\tau, T)]} \quad (25)$$

$$\frac{\Gamma \vdash e : \text{bool} \quad \Gamma \vdash c_1 : \Gamma_1 \quad \Gamma \vdash c_2 : \Gamma_2}{\Gamma \vdash \text{if } e \text{ then } c_1 \text{ else } c_2 : \Gamma} \quad (26)$$

$$\frac{\Gamma \vdash e : \text{bool} \quad \Gamma \vdash c : \Gamma_1}{\Gamma \vdash \text{while } e \text{ do } c : \Gamma} \quad (27)$$

### 5.3.2 Expressions

$$\Gamma \vdash n : \text{int} \quad (28)$$

$$\Gamma \vdash \text{true} : \text{bool} \quad (29)$$

$$\Gamma \vdash \text{false} : \text{bool} \quad (30)$$

$$\frac{\Gamma(v) = (\tau, T)}{\Gamma \vdash v : \tau} \quad (31)$$

$$\frac{\Gamma \vdash e_1 : \text{int} \quad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 \text{ aop } e_2 : \text{int}} \quad (32)$$

$$\frac{\Gamma \vdash e_1 : \text{bool} \quad \Gamma \vdash e_2 : \text{bool}}{\Gamma \vdash e_1 \text{ bop } e_2 : \text{bool}} \quad (33)$$

$$\frac{\Gamma \vdash e_1 : \text{int} \quad \Gamma \vdash e_2 : \text{int}}{\Gamma \vdash e_1 \leq e_2 : \text{bool}} \quad (34)$$

$$\frac{\Gamma \vdash e : \text{bool}}{\Gamma \vdash !e : \text{bool}} \quad (35)$$