String-Related Algorithms / DS

- Rabin Karp
- KMP (Knuth-Morris-Pratt)
- Z-Algorithm
- Boyer-Moore
- Aho-Corasick
- Suffix
- Tries

Given a string s.

For every index i, find the length of longest substring from it which is equal to prefix of whole string.

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
x	x	y	z	x	x	y	z	w	x	x	y	z	x	x	y	z	x
					⇑								⇑				
					1								5				

prefix ⇒ start from 0

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
x	x	y	z	x	x	y	z	w	x	x	y	z	x	x	y	z	x
⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑	⇑
18	1	0	0	4	1	0	0	0	8	1	0	0	5	1	0	0	1

Z-array

j = 9 → 0

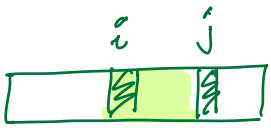
j = 10 → 1

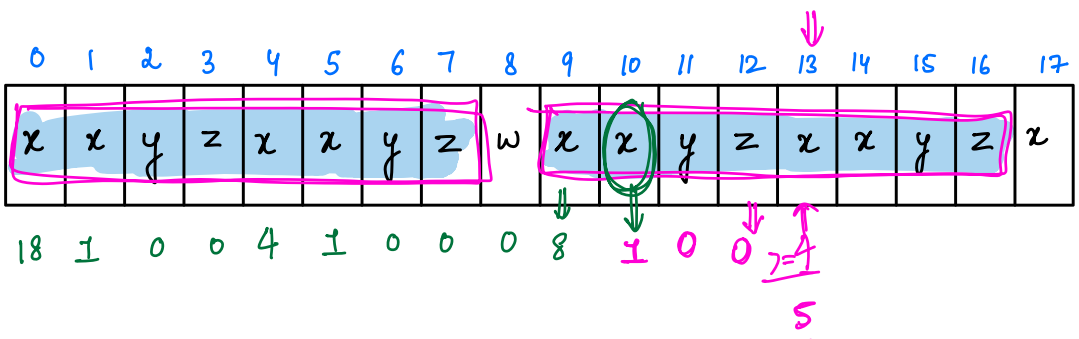
j → j - i

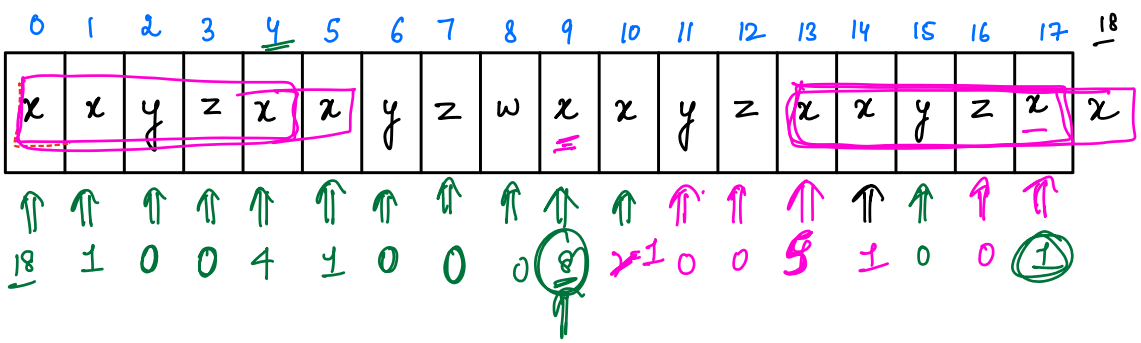
BoF:

```
z[0] = length
for( i = 1 ⟶ n)
{
    int j = i;
    while( j < N && s[j] == s[j-i])
    {
        j++;
    }
    z[i] = j-i;
}
```

T.C : $O(n^2)$







case II: inside the pink

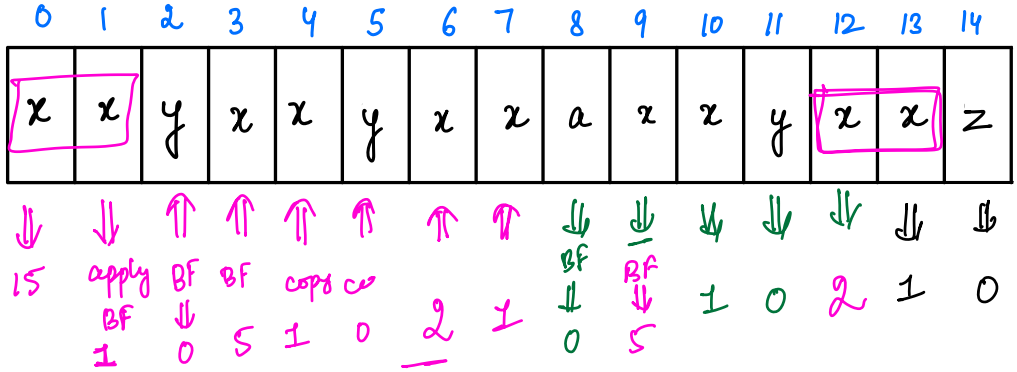
1) no of elements within pink → copy

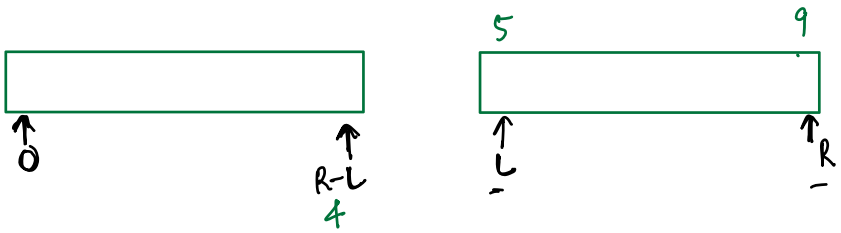
case I: when you don't know anything → out of pink box → BF

2) touch the boundary
3) crossed the boundary

} → BF → not pink box

whenever you'll brute force → update the pink box

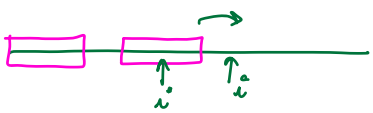


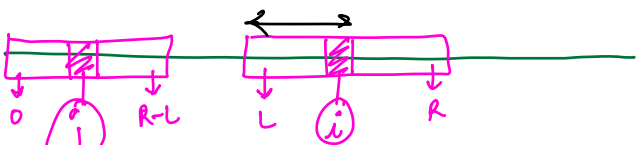


z[0] = length of the array

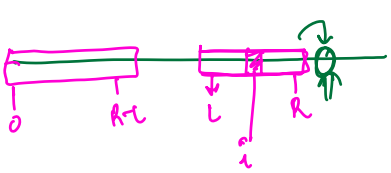
L = 0 , R = 0

```
for( i = 1 ⟶ n)
{
    // outside the segment
    if( i > R)
    {
        L = i; R = i
        while( R < N && s[R] == s[R-L])
        {
            R++;
        }
        z[i] = R-L;
        R--;
    }
    else {
        int j = i-L;        // correspo ele
        if(z[j] < R-i+1)
        {   z[i] = z[j]; }
        else
        {   L = i; R++;
            while( R < N && s[R] == s[R-L]){ R++; }
            z[i] = R-L; R--;
        }
    }
```





= i-L



T.C : O(n)

i : 1 ⟶ N : O(N)

L : 1 ⟶ N : O(N)

R : 1 ⟶ N : O(N)

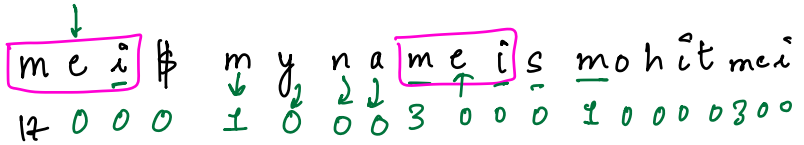
3N = O(3N)

Q Given a string S and a pattern t. Find if there exists a substring in S which matches the pattern t.

S : my name is mohit

t : mei

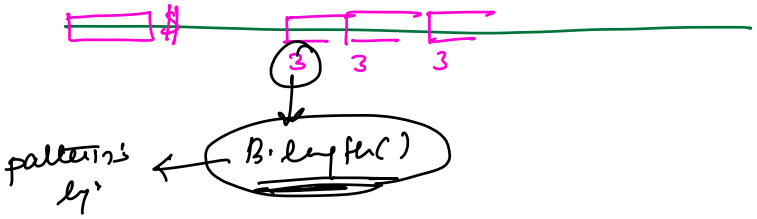
mei$myname ismohit



T : O(n+m)

S.C : O(n+m)

Q Given two strings A & B. Find the no of times B occurs in A as a substring.



Q Given a binary string, find all the cyclic permutations of the string such that

B = string

You string ↖ B^ cyclic permutation = 0.

B = 1010

B = 1010 → 0101 — 1010 0101

B^ cyclic = 0

ans = 2

a b c d ← b c d a ← c d a b ← d a b c ← a b c d

abcd $ a b c d a b c d

1010 $ 10101010

X0200 0 40404020

③ - 1 = ②

a b c d ↖
↳ b c d a
c d a b
d a b c

abcdabcd