spelling

word ⟹ valid words

$O(n * l) * q$

Autocomplete

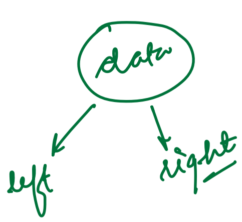
Trie → N-aray Tree

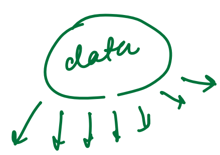
└ prefix-tree

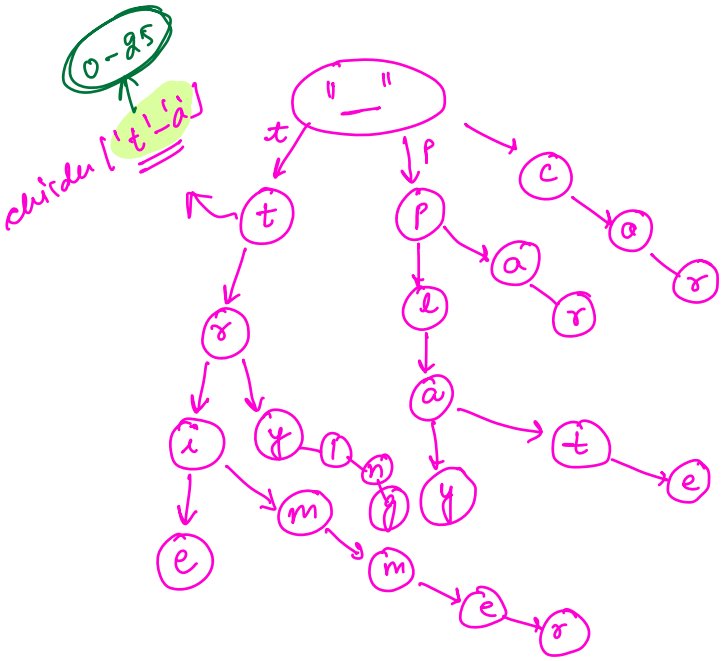
→ information retrieval

Tree like data-structure which stores the data from top to bottom

Trie	Try	Trim	play
plate	car	par	trimmer
trying	pla		








```
class Node{
    char data;
    bool isEnd;
    Node childer[26];
}
```

fals

null

'a' — 0 — 'a'-'a'
'b' — 1 — 'b'-'a'
'c' — 2 — 'c'-'a'
'd' — 3

search(trap) : Traverse your string & for every character traverse your trie

(tri)
(pla)

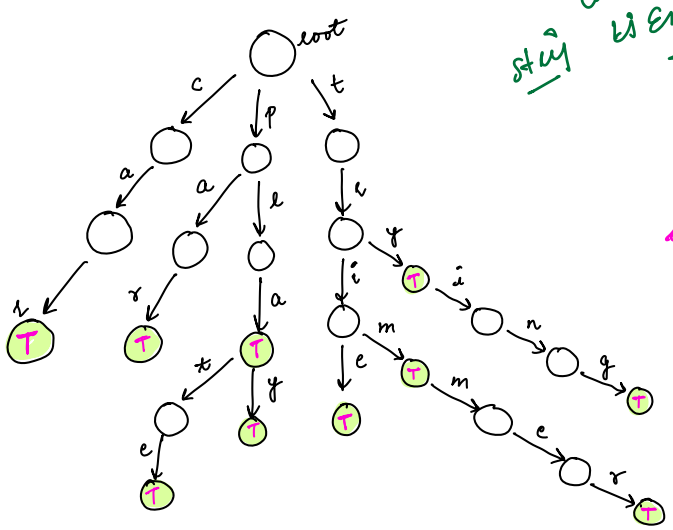
you don't find one or any of the character → stop Not found

find all characters → isEnd = true

the last character of the string is End is false

trie

autocomplete



spell

search ⇒ T.C : O(l)

S.C : n * l * 26

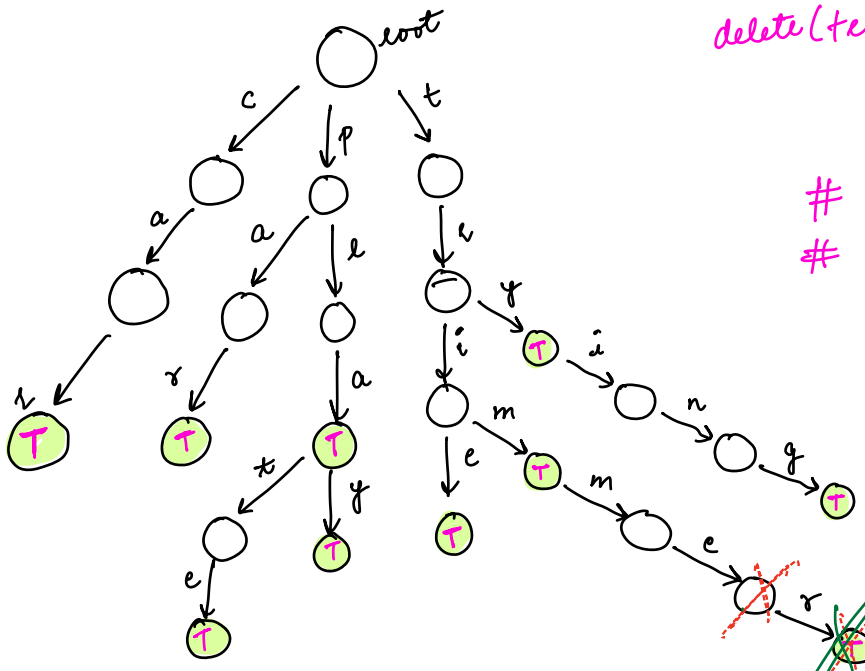
word is present multiple times!

```
class Node{
    char data;
    bool isEnd;        int freq;
    Node children[26];
}

void insert(root, word)
{
    curr = root;
    int l = word.length();
    for(int i=0; i<l; i++)
    {
        int idx= word[i]-'a';
        if( curr.children[idx] ==null)
        {
        curr.children[idx]  = new Node( word[i]);
        }
        curr= curr.children[idx];
    }
    curr.freq++;        // curr.isEnd = true;
}
```

```
bool search( root, word)
{
        curr = root;
        int l = word.length();
        for( int i=0;  i<l; i++)
        {
              int idx= word[i]-'a';
              if( curr.children[idx] ==null)
                      return false;
              curr = curr.children[idx];
        }
        if( curr.isEnd) return true; <== curr.freq>0
          return false;
}
```

deletion



delete (teamwee)

node completes a word
nodes has more children

stack (r)

e
m
m
i
r
t

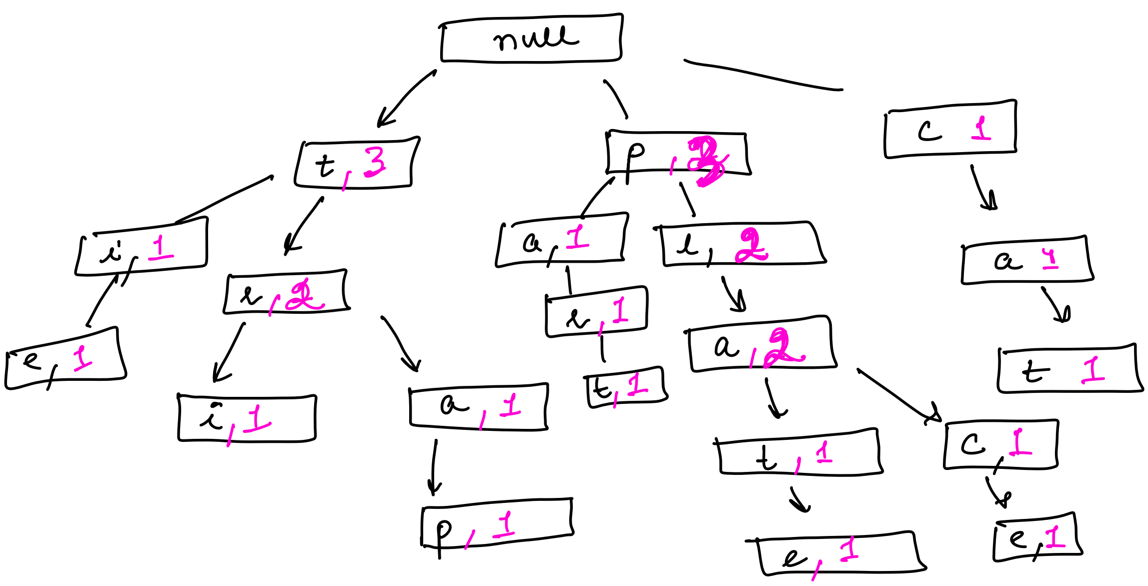
e.children['r'-'a'] = null

only store that node which is definitely not going to be deleted.

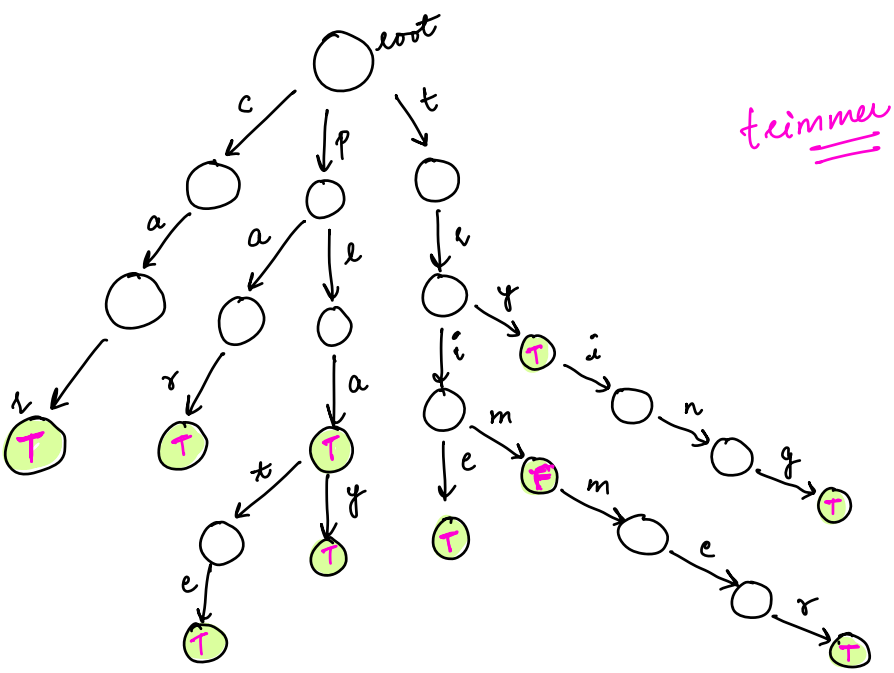
Q Find shortest unique prefix to represent each word.

Note: Assume that no word is prefix of another.
In other words, the representation is always possible.

tri	trap	plate	cat	part	place	tie
⇓	⇓	⇓	⇓	⇓	⇓	⇓
tri	tra	plat	c	pa	plac	tie



# of a string is prefix in how many differ. words



teimmer