

(A Constituent College of Somaiya Vidyavihar University) **Department of Electronics & Computer Engineering** 



<b>Course Name:</b>	Analysis of Algorithms	Semester:	IV
<b>Date of Performance:</b>	12 / 02 / 2024	Batch No:	EXCP B1
<b>Faculty Name:</b>	Prof. Payal Varangoankar	Roll No:	16014022096
Faculty Sign & Date:		Grade/Marks:	

## **Experiment No: 2**

**Title:** Binary search Finding Minimum and Maximum.

## **Aim and Objective of the Experiment:**

To learn the divide and conquer strategy of solving the problems of different types

#### COs to be achieved:

CO2: Describe various algorithm design strategies to solve different problems.

#### Theory:

#### **Historical Profile:**

Finding maximum and minimum or Binary search are a few problems that are solved with the divide-andconquer technique. This is one the simplest strategies that works on dividing the problem to the smallest possible level.

Binary Search is an extremely well-known instance of the divide-and-conquer paradigm. Given an ordered array of n elements, the basic idea of binary search is that for a given element, "probe" the middle element of the array. Then continue in either the lower or upper segment of the array, depending on the outcome of the probe until the required (given) element is reached.

#### **New Concepts to be learned:**

- 1. Number of comparisons
- 2. Application of algorithmic design strategy to any problem
- 3. Classical problem-solving vs. Divide-and-Conquer problem-solving.

Semester: IV Academic Year: 2023-24 Analysis of Algorithms Laboratory



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#### Code:

#### **BINARY SEARCH ITERATIVE:**

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
int binary_search(int A[], int key, int imin, int imax, int *iterations)
    *iterations = 0;
    while (imax >= imin) {
        (*iterations)++;
        int imid = (imin + imax) / 2;
        if (A[imid] == key) {
            return imid;
        } else if (A[imid] < key) {</pre>
    return -1;
int main() {
    int n, value, iterations;
    int arr[n];
    for (int i = 0; i < n; i++) {
        printf("%d ", arr[i]);
```



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```
l1 = clock();
int index = binary_search(arr, value, 0, n - 1, &iterations);
l2 = clock();

if (index != -1) {
    printf("Value found at index: %d\n", index);
} else {
    printf("Value not found\n");
}

double t1 = 1000000 * ((double)(l2 - l1)) / CLOCKS_PER_SEC;
printf("TIME (in micro seconds) : %f \n", t1);
printf("Number of iterations: %d\n", iterations);

return 0;
}
```

#### **BINARY SEARCH RECURSIVE:**

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>

void binary_search(int A[], int key, int imin, int imax, int *iterations) {
    if (imax < imin) {
        printf("KEY NOT FOUND");
    } else {
        int imid = (imin + imax) / 2;
        (*iterations)++;
        if (A[imid] < key) {
            binary_search(A, key, imid + 1, imax, iterations);
        } else if (A[imid] > key) {
            binary_search(A, key, imin, imid - 1, iterations);
        } else {
            printf("\nKEY %d FOUND AT %d position", A[imid], imid);
        }
    }
}

int main() {
    int main() {
    int n, value;
    clock_t 11, 12;
```



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```
printf("Enter value of n: ");
scanf("%d", &n);

int arr[n];
printf("Original Array: ");
for (int i = 0; i < n; i++) {
    arr[i] = i;
    printf("%d ", arr[i]);
}
printf("Nn");

printf("Enter value you want to search for: ");
scanf("%d", &value);

11 = clock();
int iterations = 0;
binary_search(arr, value, 0, n - 1, &iterations);
12 = clock();

double t1 = ((double)(12 - 11)) / CLOCKS_PER_SEC;
printf("TIME : %f \n", t1);
printf("Number of iterations: %d\n", iterations);
return 0;
}</pre>
```

#### **Max-Min Method:**

```
#include <stdio.h>
#include <time.h>

void MaxMin_linear(int a[], int n, int* maxl, int* minl, int* iterations) {
    *minl = *maxl = a[0];
    *iterations = 0;

for(int i = 1; i < n; i++) {
        (*iterations)++;
        if (a[i] >= *maxl) {
            *maxl = a[i];
        } else if (a[i] <= *minl) {
            *minl = a[i];
        }
}</pre>
```



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Department of Electronics & Computer Engineering



```
void MaxMin_dac(int a[], int i, int j, int* max, int* min, int* iterations) {
    if (i == j) {
        *min = *max = a[i];
    } else if (i == j - 1) {
        (*iterations)++;
        if (a[i] < a[j]) {</pre>
            *max = a[j];
            *min = a[i];
            *max = a[i];
            *min = a[j];
        int max1, min1;
        int mid = (i + j) / 2;
        MaxMin_dac(a, i, mid, max, min, iterations);
        MaxMin_dac(a, mid + 1, j, &max1, &min1, iterations);
        if (*min > min1) {
        *iterations += 2;
int main() {
    int maxl, minl, max, min, n, iterations_linear, iterations_dac;
    int arr[n];
    for (int i = 0; i < n; i++) {
        arr[i] = rand() % 10;
        printf("%d ", arr[i]);
```



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```
11 = \operatorname{clock}();
MaxMin_linear(arr, n, &maxl, &minl, &iterations_linear);
12 = clock();
d1 = clock();
d2 = clock();
double t1 = 10000000 * ((double)(12 - 11)) / CLOCKS_PER_SEC;
double t2 = 1000000 * ((double)(d2 - d1)) / CLOCKS_PER_SEC;
printf("LINEAR : \n");
printf("TIME (in microseconds) : %.2f \n", t1);
printf("DIVIDE AND CONQUER : \n");
printf("TIME (in microseconds): %.2f \n", t2);
return 0;
```

#### **Stepwise-Procedure / Algorithm:**

```
Algorithm Iterative Binary Search
```

int binary\_search(int A[], int key, int imin, int imax)

//The algorithm takes as parameters an array A[1...n], the search key and lower-higher index pair of the array.

Semester: IV

```
// Output- The algorithm returns index of the search key in the given array, if it's present. {
// continue searching while [imin, imax] is not empty
```

{
 // calculate the midpoint for roughly equal partition
 int imid = midpoint(imin, imax);

IF(A[imid] == key)

WHILE (imax >= imin)

// key found at index imid

Academic Year: 2023-24



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```
return imid:
   // determine which subarray to search
   ELSE If (A[imid] < key)
    // change min index to search upper subarray
    imin = imid + 1;
   ELSE
    // change max index to search lower subarray
    imax = imid - 1;
  }
 // key was not found
 RETURN KEY_NOT_FOUND;
The space complexity of Iterative Binary Search:
Algorithm Recursive Binary Search
int binary_search(int A[], int key, int imin, int imax)
//The algorithm takes as parameters an array A[1.. n], the search key and lower-higher index pair of
the array.
// Output- The algorithm returns index of the search key in the given array, if it's present.
 // test if array is empty
 IF (imax < imin)
  // set is empty, so return value showing not found
  RETURN KEY_NOT_FOUND;
 ELSE{
   // calculate midpoint to cut set in half
   int imid = midpoint(imin, imax);
   // three-way comparison
   IF (A[imid] > key)
    // key is in □ lower subset
    RETURN binary search(A, key, imin, imid-1);
   ELSE IF (A[imid] < key)
    // key is in \Box higher subset
    RETURN binary_search(A, key, imid+1, imax);
   ELSE
    // key has been found
    RETURN imid:
  }
}
Algorithm StraightMaxMin:
VOID StraightMaxMin (Type a[], int n, Type& max, Type& min)
// Set max to the maximum and min to the minimum of a[1:n].
  max = min = a[1];
```

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Semester: IV Academic Year: 2023-24



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```
FOR (int i=2; i \le n; i++)
IF (a[i]>max) then max = a[i];
IF (a[i] < min) min = a[i];
}
Algorithm: Recursive Max-Min
VOID MaxMin(int i, int j, Type& max, Type& min)
// A[1:n] is a global array. Parameters i and j are integers, 1 \le i \le j \le n.
//The effect is to set max and min to the largest and smallest values in a[i:j], respectively.
IF (i == j) \max = \min = a[i]; // Small(P)
ELSE IF (i == j-1) { // Another case of Small(P)
       IF (a[i] < a[j])
max = a[j]; min = a[i];
ELSE { max = a[i]; min = a[i];
ELSE {
           Type max1, min1;
// If P is not small divide P into subproblems. Find where to split the set.
int mid=(i+j)/2;
// solve the sub problems.
MaxMin(i, mid, max, min);
MaxMin(mid+1, j, max1, min1);
// Combine the solutions.
IF (max < max 1) max = max 1;
IF (min > min1) min = min1;
}
}
```



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Semester: IV



#### **Output:**

#### **BINARY METHOD ITERATIVE:**

Enter value of n: 5

Original Array: 0 1 2 3 4

Enter value you want to search for: 0

Value found at index: 0

TIME (in micro seconds): 2.000000

Number of iterations: 2

#### **BINARY METHOD RECURSIVE:**

Enter value of n: 5

Original Array: 0 1 2 3 4

Enter value you want to search for: 0

KEY 0 FOUND AT 0 position

TIME : 20.000000

Number of iterations: 2

#### **Max-Min METHOD:**

Enter value of n: 5

Original Array: 3 6 7 5 3

LINEAR :

MAX = 7, MIN = 3

TIME (in microseconds): 2.00

Number of iterations: 4 DIVIDE AND CONQUER :

MAX = 7, MIN = 3

TIME (in microseconds): 1.00

Number of iterations: 6

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Academic Year: 2023-24





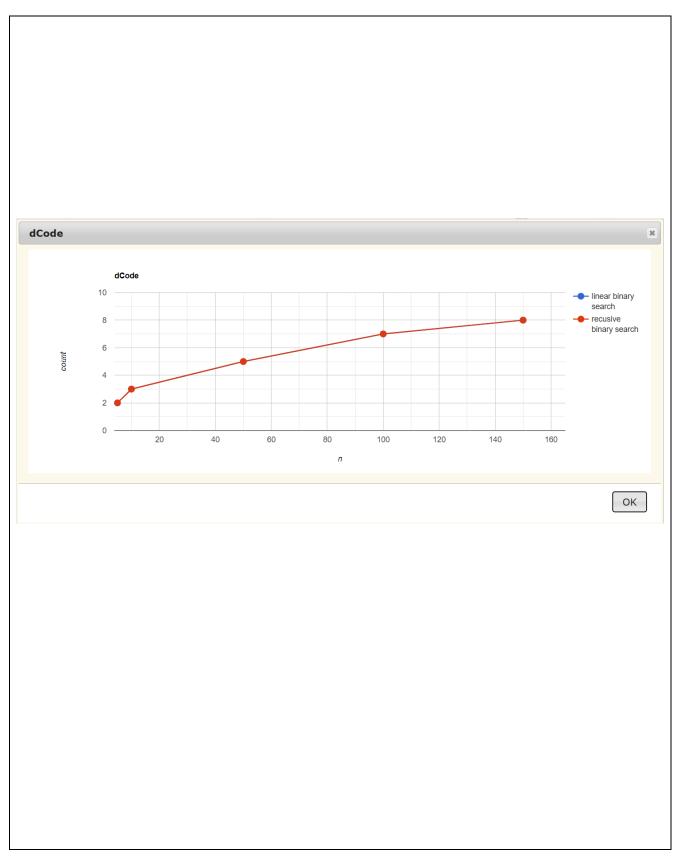


## **Observation Table: BINARY SEARCH METHOD:** abscissa x --+ ordinate y or f(x) 1 5 2 1 3 2 10 3 50 5 7 4 100 5 150 8 6 ★ ORIGINS AUTOMATICALLY CALCULATED ○ SET TO (0,0) ★ HORIZONTAL X-AXIS NAME n ★ VERTICAL Y-AXIS NAME COUNT ★ LEGEND linear binary search ★ DOT SIZE 10 ★ DISPLAY MODE ○ ONLY POINTS Draw line between points ▶ PLOT See also: Function Equation Finder - Y-Intercept 2ND PLOT (ON THE SAME GRAPH) abscissa x → ordinate y or f(x) 1 2 5 1 3 2 10 5 3 50 4 100 7 5 150 8 6 ★ LEGEND recusive binary sea...



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Semester: IV

Academic Year: 2023-24





#### (A Constituent College of Somaiya Vidyavihar University) **Department of Electronics & Computer Engineering** MAX MIN METHOD: abscissa x → ordinate y or f(x) 1 5 1 9 2 10 3 100 49 99 4 150 5 6 ★ ORIGINS AUTOMATICALLY CALCULATED ○ SET TO (0,0)

Semester: IV

- ★ HORIZONTAL X-AXIS NAME n
- ★ VERTICAL Y-AXIS NAME COUNT
- ★ LEGEND linear max min
- ★ DOT SIZE 10
- ★ DISPLAY MODE ONLY POINTS
  - DRAW LINE BETWEEN POINTS

► PLOT

See also: Function Equation Finder - Y-Intercept 2ND PLOT (ON THE SAME GRAPH)

	abscissa x →	ordinate y or f(x) †				
1	5	6				
2	10	14				
3	50	80				
4	100	162				
5						
★ LEGEND recursive max min						

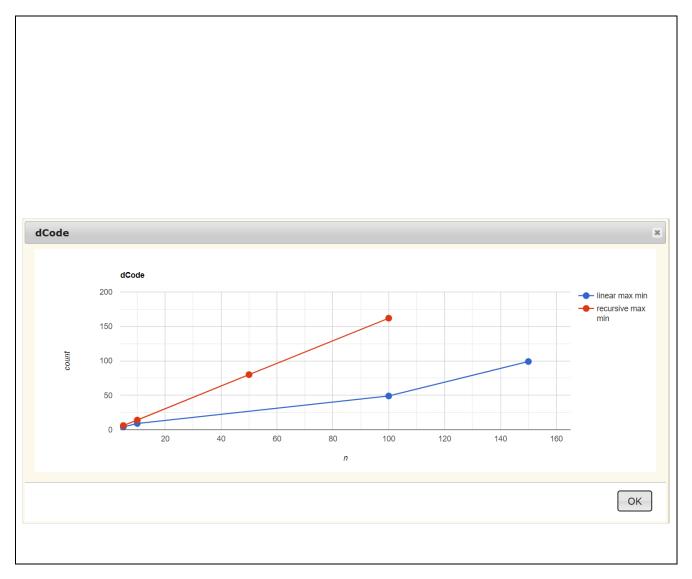
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Academic Year: 2023-24



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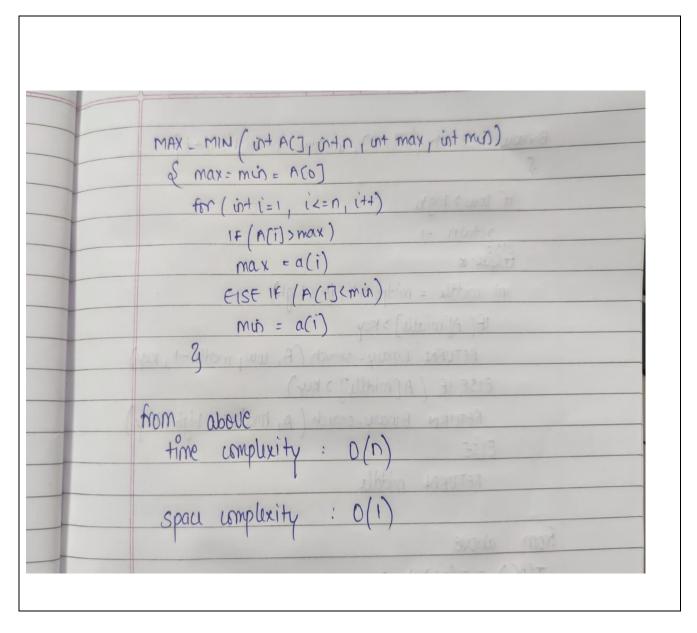
Calculations:
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Binary - Search (int AC), int low, int high; int tey)
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return -
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ELSE IF (A[middle] > Kuy)
RETURN Binary-search (A, middlett, high, key)
EISE (a) a viviginas grant
RETURN middle
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$T(n _2) = T(n _4) + 2c$ $T(n _2) = T(n _2) + 2c$
T(n 4) = T(n 8) + 3C
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2 <sup>k</sup>
K = log_n
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time complixity o (log_n)
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Academic Year: 2023-24



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$\alpha$		4.0	
Cal	CH	latio	ons:

The space complexity of Recursive Binary Search:

 $O(\log_2 n)$ 

The Time complexity of Binary Search:

O(log<sub>2</sub> n)

The space complexity of Max-Min:

 $O(\log_2 n)$ 

Time complexity for Max-Min:

O(n)

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Semester: IV Academic Year: 2023-24



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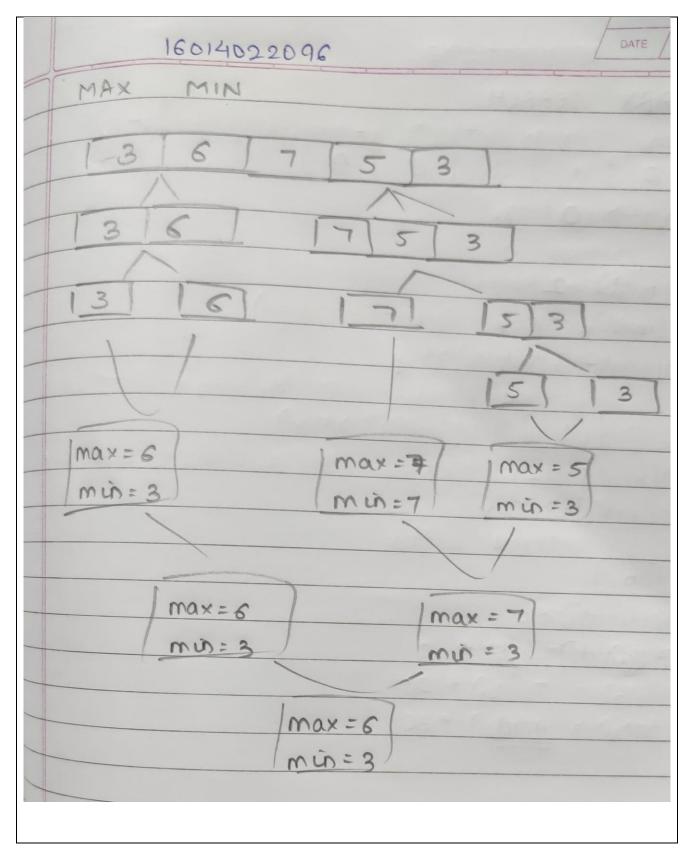
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Solve the problems	theoretically which was implem	nented during practical	
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bin	ary	search	algorithm	n and	to	max	min	alg	gorit	.hm.		

Semester: IV

**Signature of faculty in-charge with Date:** 

Academic Year: 2023-24