

# Data Flow Testing





# Learning objectives

1. Understand why data flow criteria have been designed and used
2. Recognize and distinguish basic DF criteria
  - All DU pairs, all DU paths, all definitions, ...
3. Understand how the infeasibility problem impacts data flow testing

# Motivation

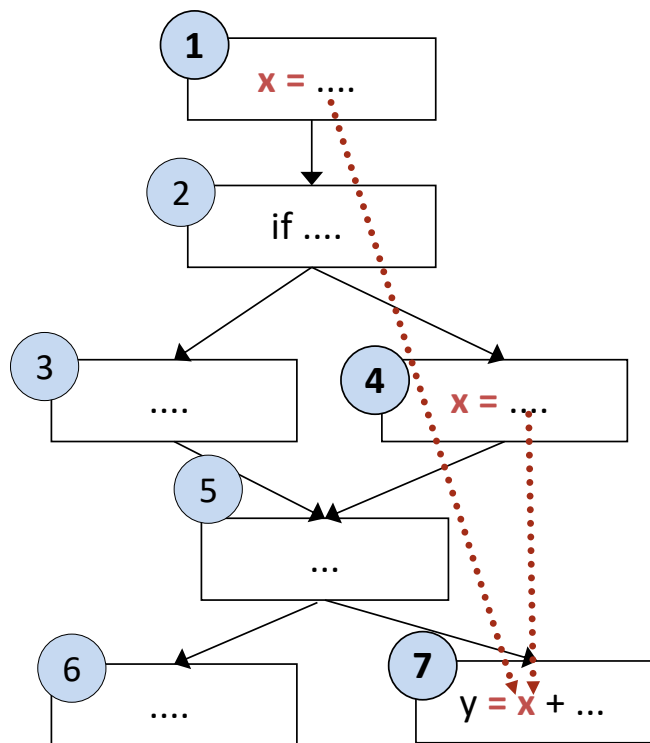
- Huge gap between **path** and **branch** coverage models
  - Path model is too strong
    - It is exhaustive but it still may miss significant test cases
    - It is often too time consuming
    - Path-based criteria require impractical number of test cases
      - And only a few paths uncover additional faults, anyway
  - Branch model is weak
    - Not exhaustive
    - May miss significant test cases
- Need another approach to distinguish “important” paths

# Data flow testing

- Instead of focus on control flow of a program
- Focus on data flow of a program
  - Program reads variables, assigns new values to variables and performs computations
  - One can visualize the flow of data values from one statement to another
- Data Flow Testing focuses on
  - the points at which variables change value
  - and the points at which variables are read
- Intuition:
  - Statements interact through *data flow*
  - Value computed in one statement, used in another
  - Bad value computation revealed only when it is used

# Data flow testing - 2

- Value of  $x$  at 7 could be computed at 1 or at 4
- Bad computation at 1 or 4 could be revealed only if they are used at 7



# Data operation categories

- **(d)** Defined, Created, Initialized
  - It assigned a value to the variable
  - A variable is **defined** when it:
    - appears in a data declaration
    - is assigned a new value
    - is a file that has been opened
    - is dynamically allocated
    - ...
- **(k)** Killed, Undefined, Released
  - Variable is deallocated at the statement fragment
  - The value and the location of the variable become unbound
- **(u)** Used:
  - The value of the variable is used in an expression

## (u) Use operation

- A variable **v** is **used** in a statement **s** when its value is applied in an expression belonging to that statement
- Two types of **use**, depending on the type of the expression:
  - **predicate use (p-use)**: **v** appears in a predicate expression of **s**
    - if ( $x > 5$ ) ...
  - **computational use (c-use)**: **v** appears in a *computation* expression of **s**
    - $y = 5 * x$ ;
      - c-use of **x** and def of **y**
- A variable is used for a computation (**c**) when it appears on the RHS (sometimes even the LHS in case of array indices) of an assignment statement.
- A variable is used in a predicate (**p**) when it appears directly in that predicate.

# Use and definition - 1

- Variables are defined by assigning values to them and are used in expressions:
  - $x = y + z$ 
    - defines variable  $x$  and uses (**c-use**) variables  $y$  and  $z$
  - `scanf ("%d %d", &x, &y)`
    - defines variables  $x$  and  $y$
  - `printf ("Output: %d \n", x + y)`
    - uses variables (**c-use**)  $x$  and  $y$
  - `if (x > 0)`
    - Uses variable (**p-use**)  $x$
- A parameter  $x$  passed to a function
  - *call-by-value*, is considered as *a use (c-use)* of  $x$
  - *call-by-reference*, is considered as *a definition and use (c-use)* of  $x$



# Use and definition - 2

- Variables can be used and re-defined in the same statement:
  - On both sides of an assignment
    - `x = x + 5;`
    - `x *= 5;`
  - As a call by reference parameter in function call
    - `increment( &y );`

# Use and definition – Arrays

- Arrays are tricky
- Example:  
    `int a[10];`  
    `a[i] = x + y;`
- Two approaches for second statement
  1. The second statement defines *a* and uses *i*, *x*, and *y*
  2. Or *second statement defines a[i]* and not the entire array *a*
    - The choice of whether to consider the entire array *a* as defined or the specific element *depends upon how stringent the requirement for coverage analysis is.*

# Example: Use and definition

```

1.  read (x, y);
2.  z = x + 2;
3.  if (z < y)
4    w = x + 1;
   else
5.  y = y + 1;
6.  print (x, y, w, z);

```

What are the *definitions* and *uses* for this program?

<i>Lines</i>	<i>Def</i>	<i>C-use</i>	<i>P-use</i>
1	x, y		
2	z	x	
3			z, y
4	w	x	
5	y	y	
6		x, y, w, z	

# Data flow testing: Two approaches

- Data flow testing can be performed at two conceptual levels.
  - Static data flow testing
  - Dynamic data flow testing
- Static data flow testing
  - Analyze source code
  - Do not execute code
  - Identify potential defects, commonly known as **data flow anomaly**
- Dynamic data flow testing
  - Involves actual program execution
  - Bears similarity with control flow testing
    - Identify paths to execute them
    - Paths are identified based on **data flow testing criteria**

# Data flow anomaly

- Anomaly: It is an abnormal way of doing something.
  - Example 1: The second definition of x overrides the first.  
x = f1(y);  
x = f2(z);
- Three types of abnormal situations with using variable.
  - Type 1: Defined and then defined again
    - Action sequence **dd**
  - Type 2: Undefined but referenced
    - Action sequence **-u**
  - Type 3: Defined but not referenced
    - Action sequence **dk**

# Data flow anomaly

- Type 1 - Defined and then defined again
  - Four interpretations of example
    - First statement is redundant.
    - First statement has a fault -- the intended one might be:  $w = f1(y)$ .
    - The second statement has a fault – the intended one might be:  $v = f2(z)$ .
    - There is a missing statement in between the two:  $v = f3(x)$ .
  - Note: It is for the programmer to make the desired interpretation.
- Type 2 - Undefined but referenced
  - Example:
    - $x = y - w$ ; //  $w$  has not been defined by the programmer
  - Two interpretations
    - The programmer made a mistake in using  $w$
    - The programmer wants to use the compiler assigned value of  $w$
- Type 3 - Defined but not referenced
  - Example: Consider  $x = f(x, y)$ . If  $x$  is not used subsequently, we have a Type 3 anomaly.

$x = f1(y);$   
 $x = f2(z);$ 

 ~~$x = f1(y);$~~   
 ~~$x = f2(z);$~~   
 $x = f2(z);$

# Dynamic data flow testing

- A family of coverage models
- Select paths by analyzing the program's data flow in order to explore sequences of events related to the status of data objects
- *E.g.*, Pick enough paths to assure that:
  - Every definition of a variable was used at least once
  - All uses of a definition have been exercised
- This family leads to test path selection strategies that fill the **gap** between **path** and **branch** coverage models

# Overview of dynamic data flow testing

- A program manipulates/uses variables in several ways
  - Initialization, assignment of variables
  - And then used in a computation and/or condition
- Motivation for data flow testing?
  - One should not feel confident that a variable has been **assigned the correct value**, if no test causes the execution of a **path** from the point of assignment to a point where the value is **used**.
- The above motivation indicates that **certain kinds of paths** should be executed in data flow testing

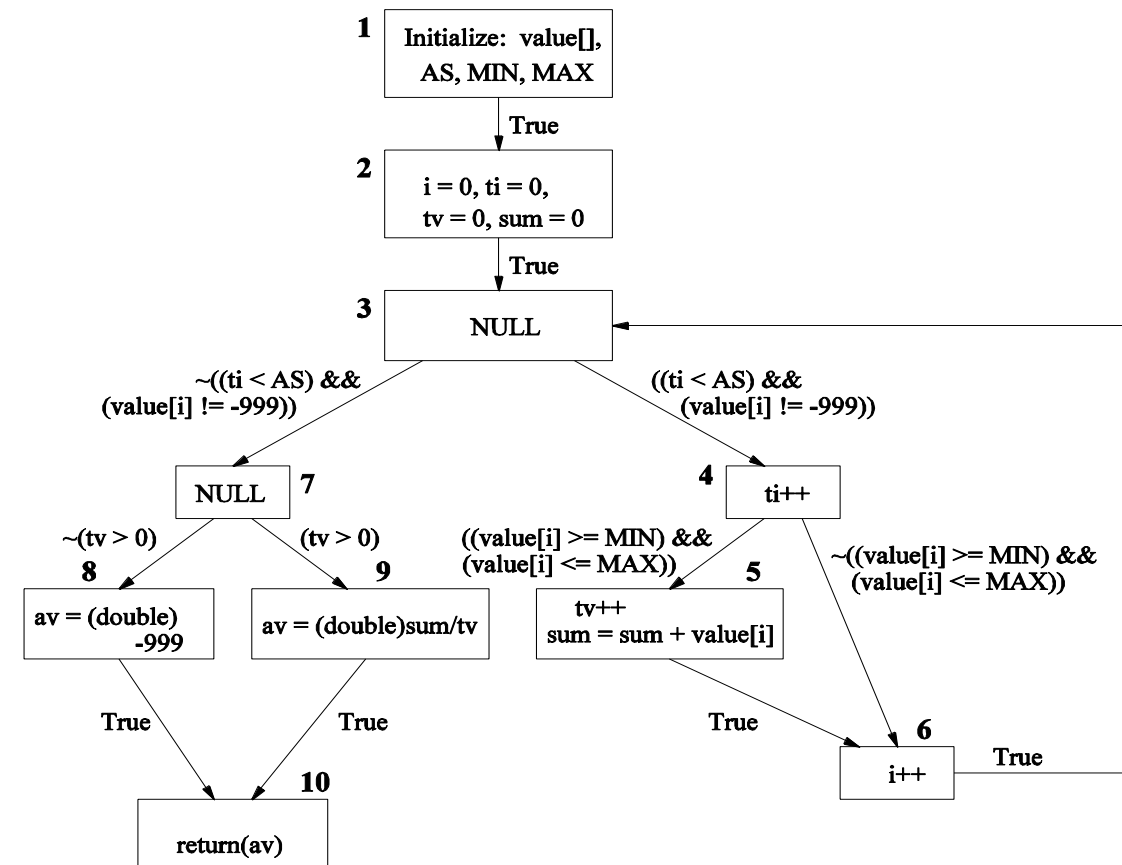


# Overview of dynamic data flow testing

- Data flow testing is outlined as follows:
  1. Draw a data flow graph from program **P**
    - Similar to control flow graph of **P**
      - All nodes, edges and paths of CGF are preserved
      - Can also consider one statement per node (makes analysis simpler)
  2. For each variable, classify each node as *defining* or *usage node*
  3. Select one or more data flow testing criteria
    - All-uses, all-defs, ...
  4. Identify paths in the data flow graph satisfying the testing criteria
  5. Compute input values for each path
    - Derive path predicate expressions from the selected paths
    - Solve the path predicate expressions to derive test inputs

# Definition clear path

- **Definition clear path (dc-path):**  
A path  $(i - n_1 - \dots - n_m - j)$ ,  $m \geq 0$ , is called a definition clear path (def-clear path) with respect to variable  $v$  if  $v$  has been neither defined nor undefined in nodes  $n_1 - \dots - n_m$ 
  - def-clear path w.r.t.  $tv$  (node 2)
    - $(2 - 3 - 4 - 5)$
    - $(2 - 3 - 4 - 6)$
    - $(2 - 3 - 4 - 6 - 3 - 4 - 6 - 3 - 4 - 5)$
    - $(2 - 3 - 4 - 5 - 6 - 3 - 4 - 5)$  ✗
  - def-clear path w.r.t.  $tv$  (node 5)
    - $(5 - 6 - 3 - 4 - 5)$



# More dataflow terms and definitions

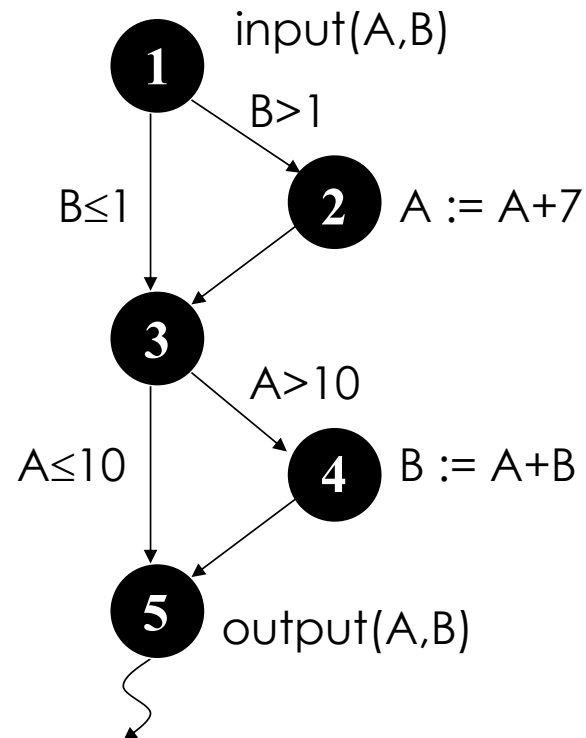
- A *definition-use pair* (“**du-pair**”) with respect to a variable **v** is a pair **(d,u)** such that
  - **d** is a node in the program’s flow graph at which **v** is defined,
  - **u** is a node or edge at which **v** is used *and*
  - there is at least one **def-clear** path *with respect to v* from **d** to **u**
  - In other words, there is at least one path  $(d, \dots, u)$  such that the value that is assigned to **v** at **d** is used at **u**
- DU-pair : A pair of nodes **(i, j)** (or **(l, <j, k>)**) such that a variable **v** is defined at **i** and that value is used at **j** (or **<j, k>**)
- Note that the definition of a du-pair does not require the existence of a **feasible** def-clear path from **d** to **u**

# Example 1

```

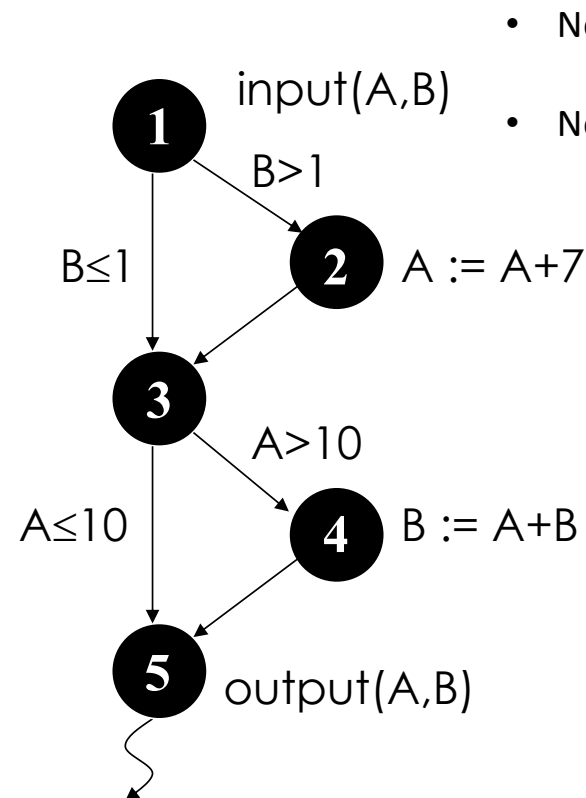
1. input(A,B)
   if (B > 1)
2.     A = A + 7
3. if (A > 10)
4.     B = A + B
5. output(A, B)
  
```

- Data flow testing
  - Compute the du-pairs and dc-paths for each variable



# Identifying DU-Pairs – Variable A

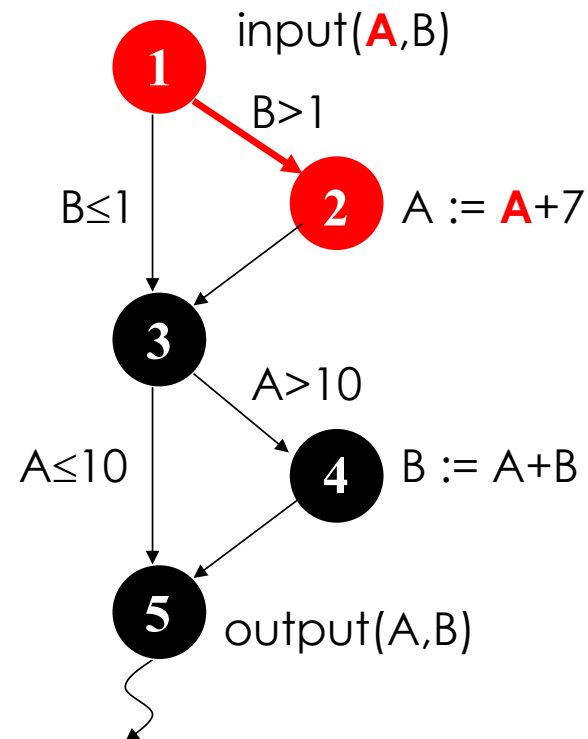
du-pair	dc-path
(1,2)	
(1,4)	
(1,5)	
(1,<3,4>)	
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



- Nodes defining A?
  - $\text{Def}(A) = \{1, 2\}$
- Nodes using A?
  - $\text{Use}\{A\} = \{2, 4, 5, \langle 3, 4 \rangle, \langle 3, 5 \rangle\}$

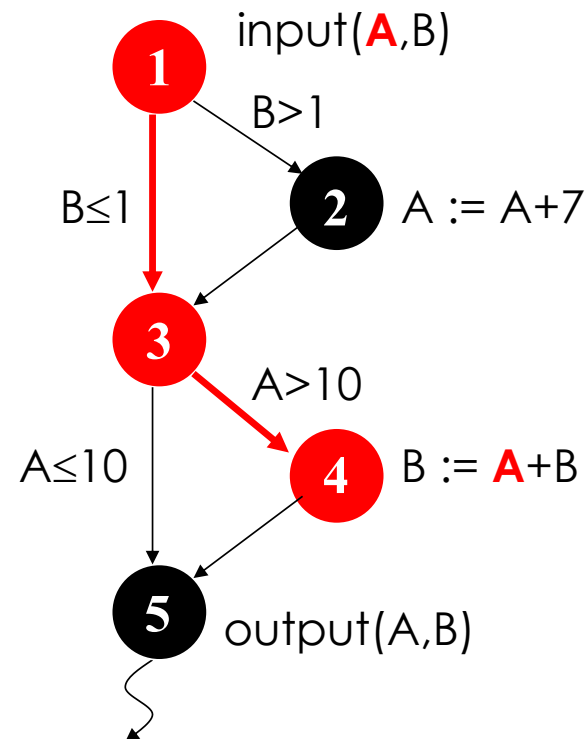
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	
(1,5)	
(1,<3,4>)	
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



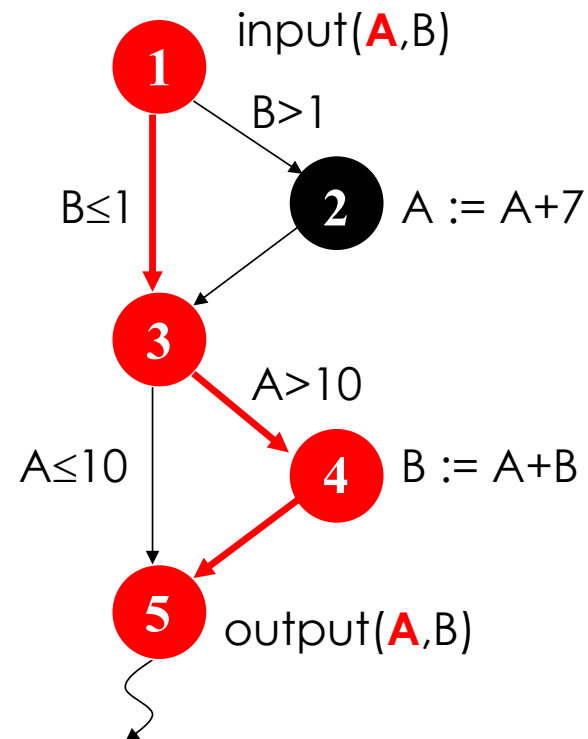
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
<b>(1,4)</b>	<b>&lt;1,3,4&gt;</b>
(1,5)	
(1,<3,4>)	
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



# Identifying DU-Pairs – Variable A

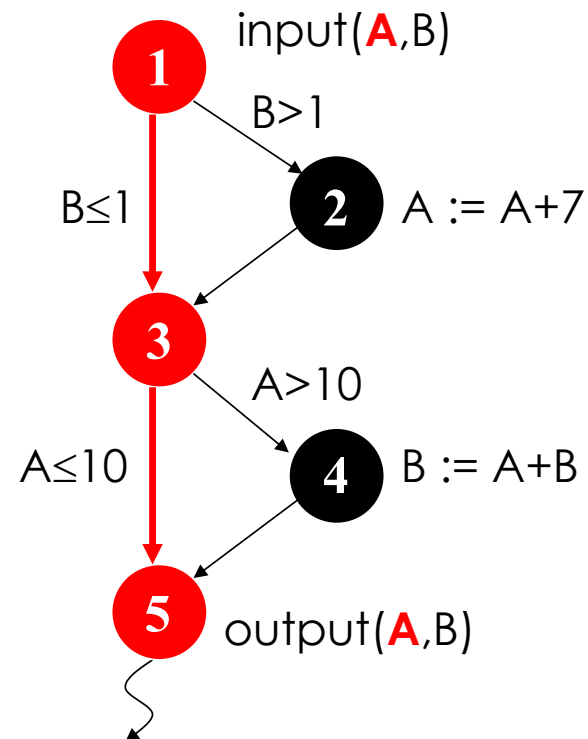
du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
<b>(1,5)</b>	<b>&lt;1,3,4,5&gt;</b>
(1,<3,4>)	
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	





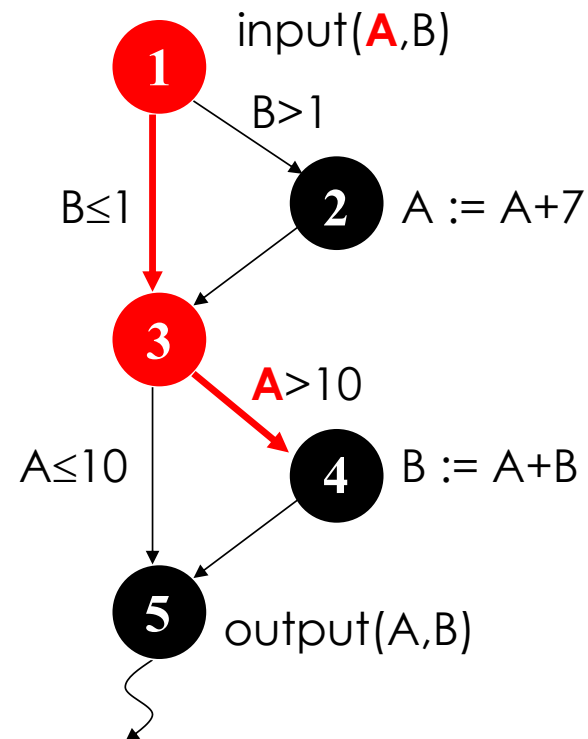
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
<b>(1,5)</b>	<1,3,4,5>
	<b>&lt;1,3,5&gt;</b>
(1,<3,4>)	
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



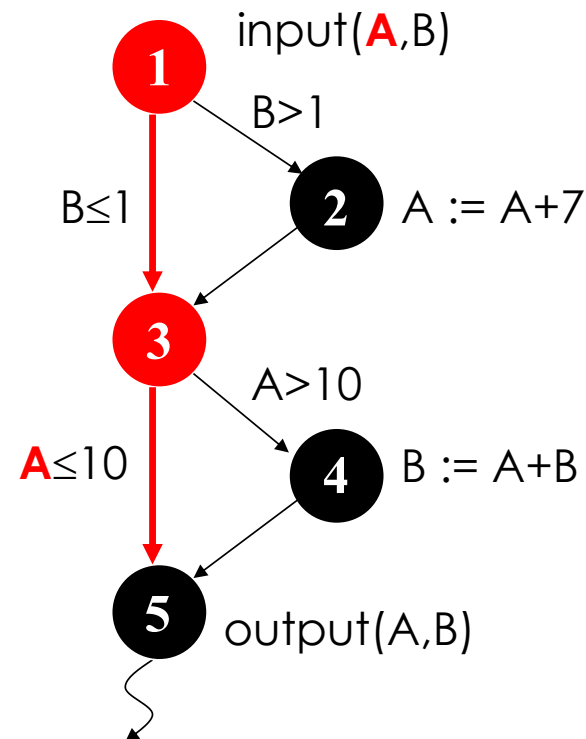
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
<b>(1,&lt;3,4&gt;)</b>	<b>&lt;1,3,4&gt;</b>
(1,<3,5>)	
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



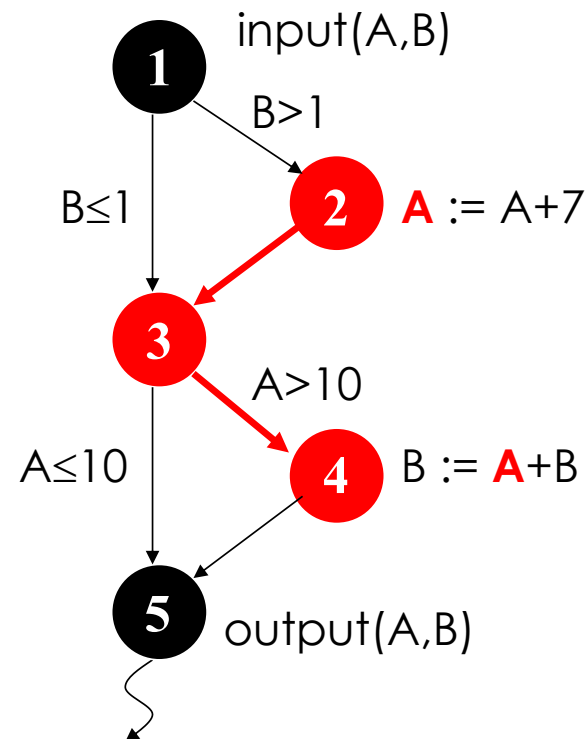
# Identifying DU-Pairs – Variable A

du-pair	path(s)
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
<b>(1,&lt;3,5&gt;)</b>	<b>&lt;1,3,5&gt;</b>
(2,4)	
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



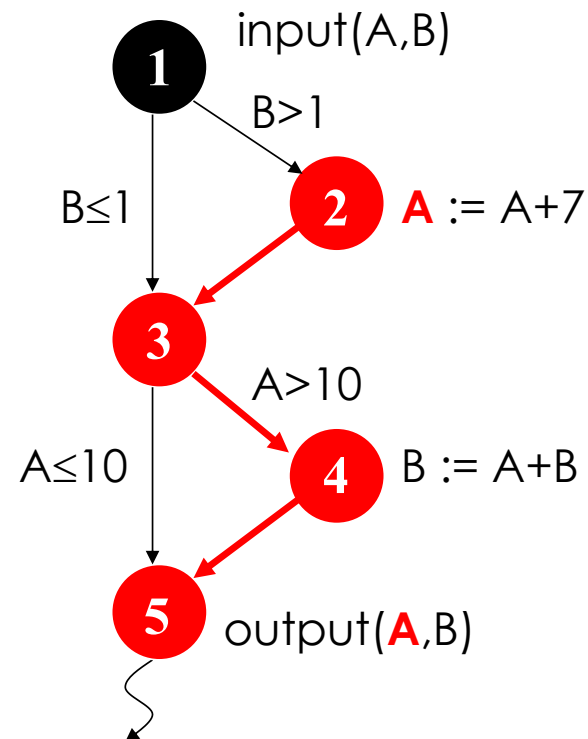
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
<b>(2,4)</b>	<b>&lt;2,3,4&gt;</b>
(2,5)	
(2,<3,4>)	
(2,<3,5>)	



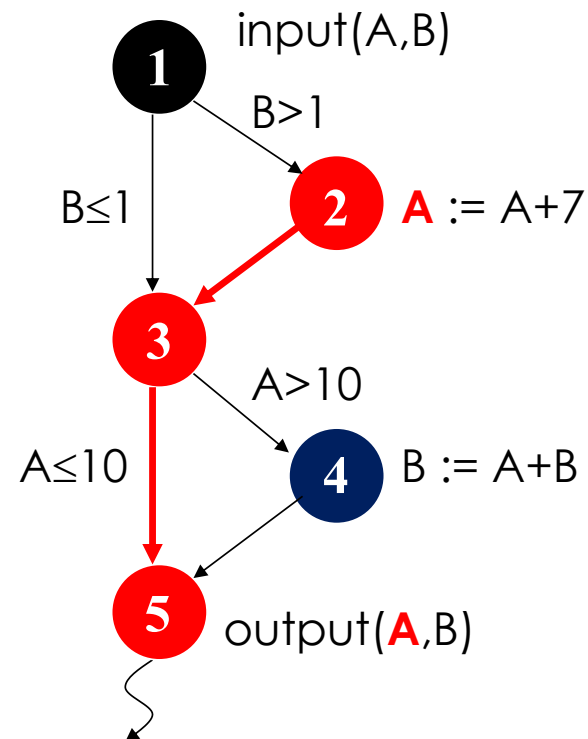
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
<b>(2,5)</b>	<b>&lt;2,3,4,5&gt;</b>
(2,<3,4>)	
(2,<3,5>)	



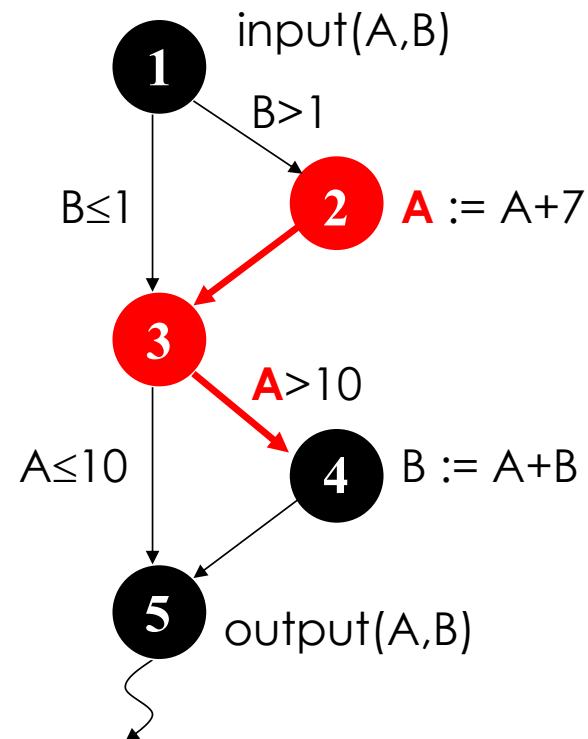
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
<b>(2,5)</b>	<2,3,4,5>
	<b>&lt;2,3,5&gt;</b>
(2,<3,4>)	
(2,<3,5>)	



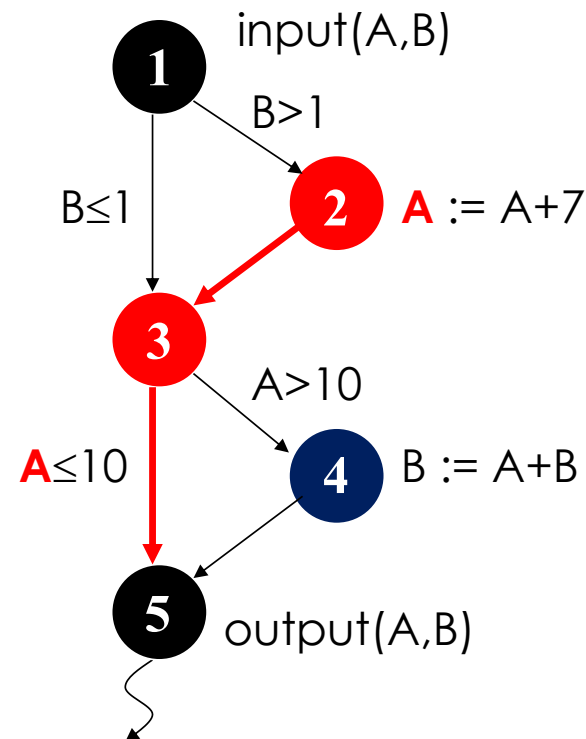
# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
(2,5)	<2,3,4,5>
	<2,3,5>
<b>(2,&lt;3,4&gt;)</b>	<b>&lt;2,3,4&gt;</b>
(2,<3,5>)	



# Identifying DU-Pairs – Variable A

du-pair	dc-path
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
(2,5)	<2,3,4,5>
	<2,3,5>
(2,<3,4>)	<2,3,4>
<b>(2,&lt;3,5&gt;)</b>	<b>&lt;2,3,5&gt;</b>

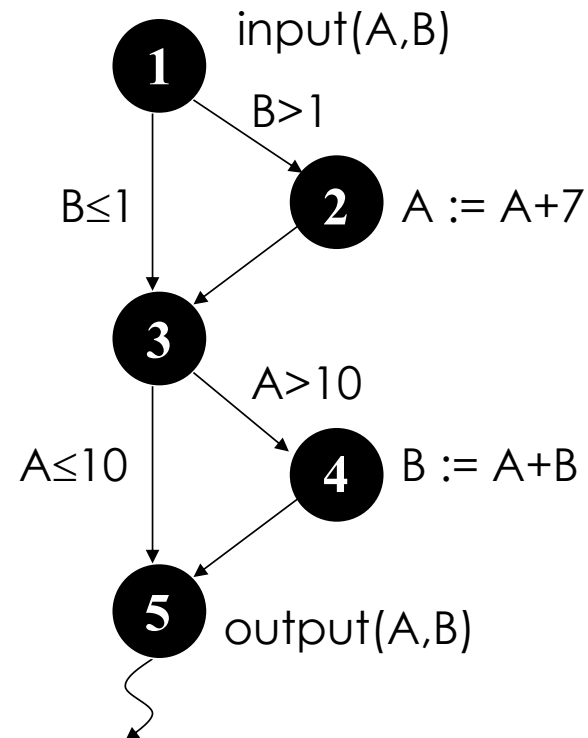




# Identifying DU-Pairs – Variable B

- Nodes defining B?
  - $\text{Def}(B) = \{1, 4\}$
- Nodes using B?
  - $\text{Use}\{B\} = \{4, 5, \langle 1, 2 \rangle, \langle 1, 3 \rangle\}$

du-pair	dc-path
(1,4)	$\langle 1, 2, 3, 4 \rangle$
	$\langle 1, 3, 4 \rangle$
(1,5)	$\langle 1, 2, 3, 5 \rangle$
	$\langle 1, 3, 5 \rangle$
(1, $\langle 1, 2 \rangle$ )	$\langle 1, 2 \rangle$
(1, $\langle 1, 3 \rangle$ )	$\langle 1, 3 \rangle$
(4,5)	$\langle 4, 5 \rangle$



# Frankl and Weyuker's Data flow coverage criteria

- Seven data flow testing criteria
  - **All-defs**
  - **All-uses**
  - All-c-uses
  - All-p-uses
  - All-p-uses/some-c-uses
  - All-c-uses/some-p-uses
  - **All-du-paths**
- The family of data flow criteria requires that the test data execute definition-clear paths from each node containing a definition of a variable to specified nodes containing *c-use* and edges containing *p-use* of that variable.

# All-Defs coverage criterion

- *All-Defs*

For **every** program variable **v**,  
at least one def-clear path from **every** definition of  
**v**  
to at least one c-use or one p-use of **v** must be covered

➤ Meaning: All definitions get used at least once

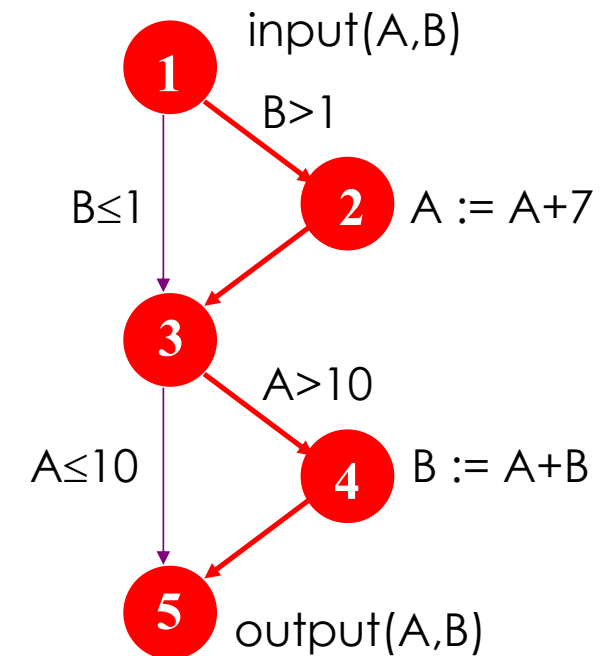
# All-Defs coverage criterion - Example

- Consider a test case executing complete path  $\langle 1, 2, 3, 4, 5 \rangle$ 
  - **Complete path:** A complete path is a path from the entry node to the exit node
  - Corresponds to the entry-exit node of CF models
- Identify **all def-clear paths covered (ie subsumed)** by this path for each variable
- Are **all definitions for each variable** associated with at least one of the subsumed def-clear paths?

Def-Clear Paths subsumed by **<1,2,3,4,5>** for Variable A and B

du-pair for A	dc-path
(1,2)	<1,2> ✓
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4> ✓
(2,5)	<2,3,4,5> ✓
	<2,3,5>
(2,<3,4>)	<2,3,4> ✓
(2,<3,5>)	<2,3,5>

du-pair for B	dc-path
(1,4)	<1,2,3,4> ✓
	<1,3,4>
(1,5)	<1,2,3,5>
	<1,3,5>
(4,5)	<4,5> ✓
(1,<1,2>)	<1,2> ✓
(1,<1,3>)	<1,3>



- Since **<1,2,3,4,5>** covers at least one def-clear path from every definition of A or B to at least one c-use or p-use of A or B, **All-Defs** coverage is achieved with a single test case

# All-Uses coverage criterion

- All-Uses coverage criterion:

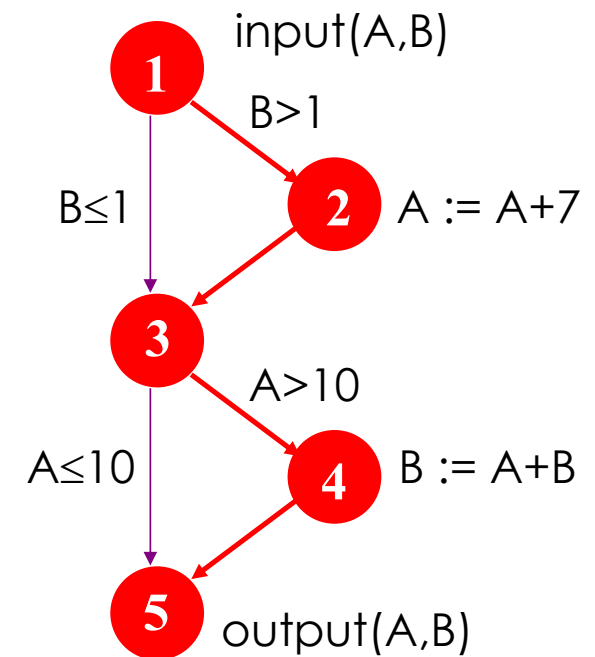
For **every** program variable **v**,  
**at least one def-clear path** from **every definition** of **v**  
to **every c-use** and **every p-use** of **v** must be covered

- Requires all du-pairs are exercised at least once
- Meaning: Every computation and branch directly affected by a definition is exercised

Does **<1,2,3,4,5>** achieves All-Uses for variables A and B?

du-pair for A	dc-path
(1,2)	<1,2> ✓
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4> ✓
(2,5)	<2,3,4,5> ✓
	<2,3,5>
(2,<3,4>)	<2,3,4> ✓
(2,<3,5>)	<2,3,5>

du-pair for B	dc-path
(1,4)	<1,2,3,4> ✓
	<1,3,4>
(1,5)	<1,2,3,5>
	<1,3,5>
(4,5)	<4,5> ✓
(1,<1,2>)	<1,2> ✓
(1,<1,3>)	<1,3>

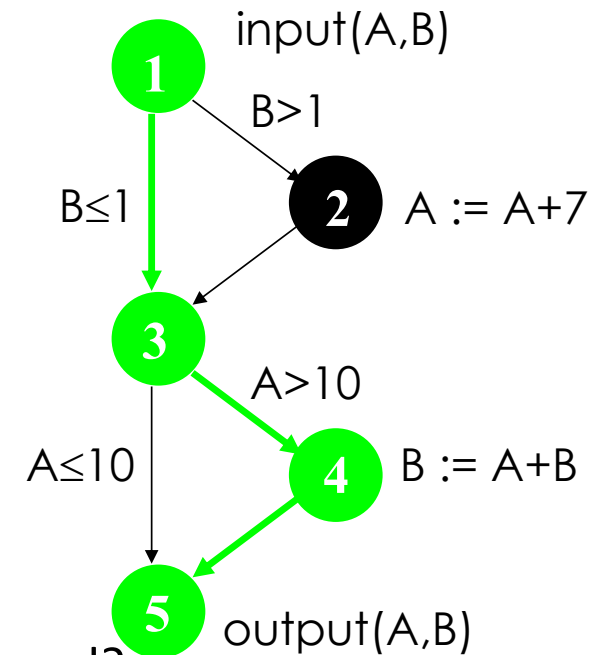


- No! For example (1,4) and (1,5) for A are not covered
- Consider additional test cases executing paths:
  - **<1,3,4,5>**
  - **<1,2,3,5>**

## Def-Clear paths subsumed by $\langle 1, 3, 4, 5 \rangle$

du-pair for A	dc-path
(1,2)	$\langle 1, 2 \rangle$ ✓
(1,4)	$\langle 1, 3, 4 \rangle$ ✓
(1,5)	$\langle 1, 3, 4, 5 \rangle$ ✓
	$\langle 1, 3, 5 \rangle$
(1, $\langle 3, 4 \rangle$ )	$\langle 1, 3, 4 \rangle$ ✓
(1, $\langle 3, 5 \rangle$ )	$\langle 1, 3, 5 \rangle$
(2,4)	$\langle 2, 3, 4 \rangle$ ✓
(2,5)	$\langle 2, 3, 4, 5 \rangle$ ✓
	$\langle 2, 3, 5 \rangle$
(2, $\langle 3, 4 \rangle$ )	$\langle 2, 3, 4 \rangle$ ✓
(2, $\langle 3, 5 \rangle$ )	$\langle 2, 3, 5 \rangle$

du-pair for B	dc-path
(1,4)	$\langle 1, 2, 3, 4 \rangle$ ✓
	$\langle 1, 3, 4 \rangle$ ✓
(1,5)	$\langle 1, 2, 3, 5 \rangle$
	$\langle 1, 3, 5 \rangle$
(4,5)	$\langle 4, 5 \rangle$ ✓✓
(1, $\langle 1, 2 \rangle$ )	$\langle 1, 2 \rangle$ ✓
(1, $\langle 1, 3 \rangle$ )	$\langle 1, 3 \rangle$ ✓



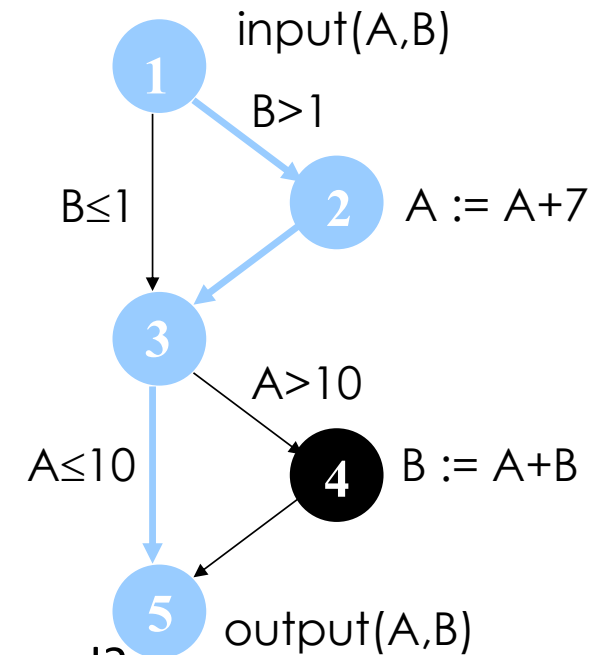
- Is the All-Uses coverage achieved?
- No! Du-pairs not exercised yet:
- Variable A : du-pair (1,  $\langle 3, 5 \rangle$ ) and (2,  $\langle 3, 5 \rangle$ )
- Variable B: du-pair (1, 5)



# Def-Clear paths subsumed by $\langle 1,2,3,5 \rangle$

du-pair for A	dc-path
(1,2)	$\langle 1,2 \rangle$ ✓ ✗
(1,4)	$\langle 1,3,4 \rangle$ ✓
(1,5)	$\langle 1,3,4,5 \rangle$ ✓
	$\langle 1,3,5 \rangle$
(1,<3,4>)	$\langle 1,3,4 \rangle$ ✓
(1,<3,5>)	$\langle 1,3,5 \rangle$
(2,4)	$\langle 2,3,4 \rangle$ ✓
(2,5)	$\langle 2,3,4,5 \rangle$ ✓
	$\langle 2,3,5 \rangle$ ✗
(2,<3,4>)	$\langle 2,3,4 \rangle$ ✓
(2,<3,5>)	$\langle 2,3,5 \rangle$ ✗

du-pair for B	dc-path
(1,4)	$\langle 1,2,3,4 \rangle$ ✓
	$\langle 1,3,4 \rangle$ ✓
(1,5)	$\langle 1,2,3,5 \rangle$ ✗
	$\langle 1,3,5 \rangle$
(4,5)	$\langle 4,5 \rangle$ ✓ ✓
(1,<1,2>)	$\langle 1,2 \rangle$ ✓ ✗
(1,<1,3>)	$\langle 1,3 \rangle$ ✓



- Is the All-Uses coverage achieved?
- None of the three test cases covers the du-pair (1,<3,5>) for variable **A**
  - All-Uses Coverage is not achieved
  - Need additional test case

# More data flow coverage criteria

- Given the set **V** of variables of program **P**
- **All-P-Uses:**
  - For every variable  $v$  in  $V$ , at least one dc-path from every definition of  $v$  to every P-use of  $v$  must be covered
- **All-C-Uses:**
  - For every variable  $v$  in  $V$ , at least one dc-path from every definition of  $v$  to every C-use of  $v$  must be covered

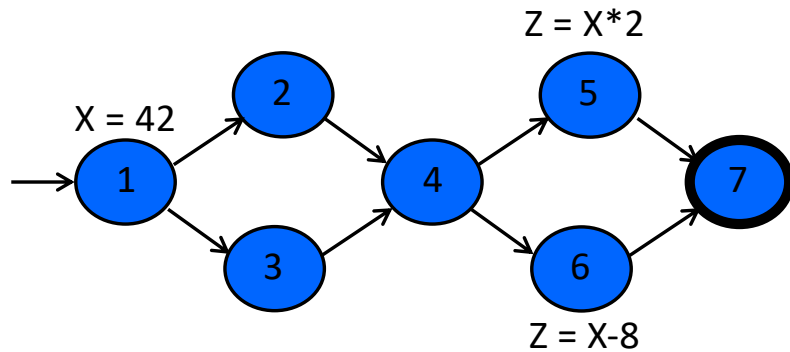
## More data flow coverage criteria - 2

- Given the set **V** of variables of program **P**
- **All-P-Uses/Some-C-Uses:**
  - For every variable  $v$  in  $V$ , at least one dc-path from every definition of  $v$  to every P-use of  $v$  must be covered
  - If a definition of  $v$  has no P-uses, at least one dc-path from the definition of  $v$  to a C-use of  $v$  must be covered
- **All-C-Uses/Some-P-Uses:**
  - For every variable  $v$  in  $V$ , at least one dc-path from every definition of  $v$  to every C-use of  $v$  must be covered
  - If a definition of  $v$  has no C-uses, at least one dc-path from the definition of  $v$  to a P-use of  $v$  must be covered

# All-DU-Paths coverage criterion

- **Simple path:** A simple path is a path in which all nodes, except possibly the first and the last, are distinct
  - 1-2-3-4-5-9 and 1-3-4-5-2-1 are simple paths
  - 1-2-3-4-5-3-7 and 1-2-3-4-1-6-5-8 are not simple paths
- A path  $\langle n_1, n_2, \dots, n_j, n_k \rangle$  is a **du-path** wrt a variable  $v$  if
  1.  $v$  is defined at node  $n_1$  and
  2. there is a **c-use** of  $v$  at node  $n_k$  or a **p-use** of  $v$  at edge  $\langle n_j, n_k \rangle$  and
  3.  $\langle n_1, n_2, \dots, n_j, n_k \rangle$  is a def-clear **simple** path
- **All DU-Paths criterion:**  
for every program variable  $v$ ,  
every du-path  
from every definition of  $v$  to every c-use and every p-use of  $v$  must be covered

# Data flow testing example



- Nodes defining X?
  - Def(X) = {1}
- Nodes using X?
  - Use{X} = {5, 6}

du-pair	dc-path	
(1, 5)	< 1, 2, 4, 5 >	✓
	< 1, 3, 4, 5 >	✓
(1, 6)	< 1, 2, 4, 6 >	✓
	< 1, 3, 4, 6 >	✓

Determine a minimal test suite that achieves 100% coverage for:

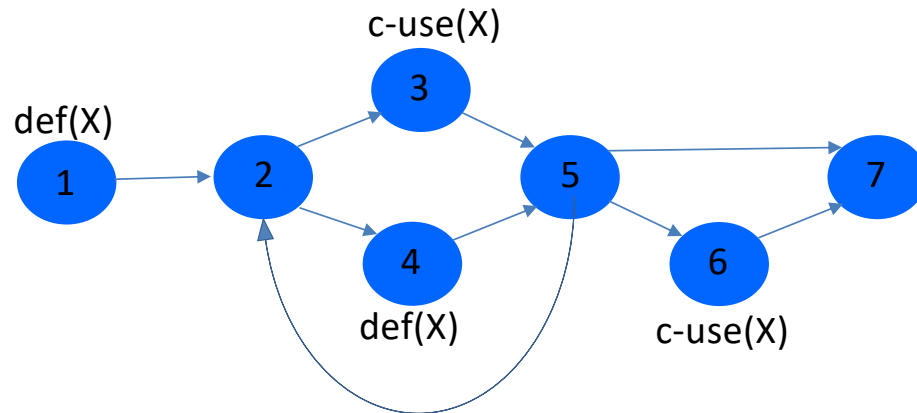
- All-defs
- All-uses
- All-du-paths

All-defs for X
[ 1, 2, 4, 5, 7 ]

All-uses for X
[ 1, 2, 4, 5, 7 ]
[ 1, 2, 4, 6, 7 ]

All-du-paths for X
[ 1, 2, 4, 5, 7 ]
[ 1, 2, 4, 6, 7 ]
[ 1, 3, 4, 5, 7 ]
[ 1, 3, 4, 6, 7 ]

# Data Flow Testing Example - 2



Determine a minimal test suite that achieves 100% coverage for:

- All-defs
- All-uses
- All-du-paths

All-defs for X	All-uses for X	All-du-paths for X	All-du-paths for X
[ 1, 2, 3, 5, 2, 4, 5, 6, 7 ]	[ 1, 2, 3, 5, 6, 7 ] [ 1, 2, 4, 5, 2, 3, 5, 6, 7 ]	<del>[ 1, 2, 3, 5, 6, 7 ] [ 1, 2, 4, 5, 2, 3, 5, 6, 7 ] [ 1, 2, 4, 5, 6, 7 ]</del>	[ 1, 2, 3, 5, 6, 7 ] [ 1, 2, 4, 5, 2, 3, 5, 2, 4, 5, 6, 7 ]

du-pair	dc-path	Du-path
(1, 3)	< 1, 2, 3 >	✓
(1, 6)	< 1, 2, 3, 5, 6 >	✓
	< 1, (2, 3, 5)+, 6 >	✗
(4, 3)	< 4, 5, 2, 3 >	✓
(4, 6)	< 4, 5, 2, 3, 5, 6 >	✗
	< 4, (5,2,3)+, 5, 6 >	✗
	< 4, 5, 6 >	✓

# Relationship between strategies

