

EP2200 Course Project

Project I - Error Control in Relay Networks

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1 AF RELAYING WITH END-TO-END ARQ

The queuing model for the RS and MS nodes - M/M/1.

1) The end-to-end error probability of packet transmission ($p_{e,e2e}$):

$$p_{e,e2e} = 1 - \prod_{k=1}^{r+1} (1 - p_{(e,k)}) \quad (1)$$

The packet arrival rate at the first relay (λ_1) :

$$\begin{aligned} \lambda_1 &= \lambda_{r+1} \cdot p_{e,e2e} + \lambda \\ \lambda_2 &= \lambda_1 \\ \lambda_3 &= \lambda_3 \\ &\dots \\ \lambda_{r+1} &= \lambda_r \end{aligned}$$

We can write,

$$\lambda_{r+1} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)})}$$

Therefore,

$$\lambda_1 = \lambda_{r+1} \cdot p_{e,e2e} + \lambda = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)})} \quad (2)$$

2) For RS:

$$\rho_i^{AF} = \frac{\lambda_i^{AF}}{\mu_i^{AF}} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)}) \cdot \mu_{AF}}$$

For M/M/1 systems, the waiting time is $W = \frac{1}{\mu - \lambda} - \frac{1}{\mu}$.
The average queuing delay is:

$$W_{RS} = \frac{\lambda}{\mu_{AF}^2 \prod_{k=1}^{r+1} (1 - p_{(e,k)}) - \lambda \cdot \mu_{AF}} \quad (3)$$

For M/M/1 systems, the average number in the system is $N = \frac{\rho}{1-\rho}$. The average number of packets (waiting or under processing) is:

$$N_{RS} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)}) \cdot \mu_{AF} - \lambda} \quad (4)$$

For MS:

$$\rho_{MS} = \frac{\lambda_{r+1}}{\mu_{MS}} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)}) \cdot \mu_{MS}}$$

Average queuing delay is:

$$W_{MS} = \frac{\lambda}{\mu_{MS}^2 \prod_{k=1}^{r+1} (1 - p_{(e,k)}) - \lambda \cdot \mu_{MS}} \quad (5)$$

Average number of packets (waiting or under processing) is:

$$N_{MS} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{(e,k)}) \cdot \mu_{MS} - \lambda} \quad (6)$$

3) The average amount of time between the production of a packet and its successful reception at the MS is equal to the total of all the time spent in the RS and the MS.

$$\bar{T} = \sum_{i=1}^r (W_{RS} + \frac{1}{\mu_i^{AF}}) + W_{MS} + \frac{1}{\mu_{MS}}$$

The time spent in each RS is also the same because a packet is only decoded and examined on the final hop, when the MS gets it. The average end-to-end delay is:

$$\begin{aligned} \bar{T} &= \frac{r \cdot \prod_{k=1}^{r+1} (1 - p_{(e,k)})}{\mu_{AF} \cdot \prod_{k=1}^{r+1} (1 - p_{(e,k)}) - \lambda} \\ &\quad + \frac{\prod_{k=1}^{r+1} (1 - p_{(e,k)})}{\mu_{MS} \cdot \prod_{k=1}^{r+1} (1 - p_{(e,k)}) - \lambda} \quad (7) \end{aligned}$$

2 DF RELAYING WITH END-TO-END OR HOP-BY-HOP ARQ

A. DF Relaying With End-To-End ARQ

1) The packet arrival rate at each RS and MS (λ_j):

$$\lambda_1 = \lambda + \lambda_1 \cdot p_{e,1} + \lambda_2 \cdot p_{e,2} + \dots + \lambda_{r+1} \cdot p_{e,r+1}$$

$$= \lambda + \lambda \cdot \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{k=1}^{j=1} (1 - p_{(e,k)})]$$

$$= \frac{\lambda}{1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{k=1}^{i=1} (1 - p_{(e,i)})]}$$

Therefore,

$$\lambda_i = \lambda_1 \prod_{i=1}^{k=0} (1 - p_{(e,k)}) = \frac{\prod_{i=1}^{k=1} (1 - p_{(e,k)})}{1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^{k=1} (1 - p_{(e,j)})]} \quad (8)$$

2) For RS:

Average queuing delay $W_{RS}^i = \frac{1}{\mu_{DF} - \lambda_i} - \frac{1}{\mu_{DF}}$:

$$\lambda \cdot \prod_{i=1}^{k=1} (1 - p_{(e,k)}) \div \mu_{DF}^2 [1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]] - \lambda \mu_{DF} \prod_{i=1}^{k=1} (1 - p_{(e,k)}) \quad (9)$$

As the utilization here is:

$$\rho_{RS}^i = \frac{\lambda_i}{\mu_{DF}} = \frac{\lambda \cdot \prod_{k=1}^{i-1} (1 - p_{(e,k)})}{\mu_{DF} - \mu_{DF} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]}$$

Hence, average number of packets at every RS is $N_{RS}^i = \frac{\rho_{RS}^i}{1 - \rho_{RS}^i}$:

$$\lambda \cdot \prod_{k=1}^{i-1} (1 - p_{(e,k)}) \div \mu_{DF} - \mu_{DF} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})] - \lambda \prod_{k=1}^{i-1} (1 - p_{(e,k)}) \quad (10)$$

For MS:

Average queuing delay $W_{MS} = \frac{1}{\mu_{MS} - \lambda_{r+1}} - \frac{1}{\mu_{MS}}$:

$$\lambda \cdot \prod_{k=1}^r (1 - p_{(e,k)}) \div \mu_{MS}^2 [1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]] - \lambda \mu_{MS} \prod_{k=1}^r (1 - p_{(e,k)}) \quad (11)$$

As the utilization here is:

$$\rho_{MS} = \frac{\lambda_{r+1}}{\mu_{MS}} = \frac{\lambda \cdot \prod_{k=1}^r (1 - p_{(e,k)})}{\mu_{MS} - \mu_{MS} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]}$$

Hence, average number of packets at every MS is $N_{MS} = \frac{\rho_{MS}}{1 - \rho_{MS}}$:

$$\lambda \cdot \prod_{k=1}^r (1 - p_{(e,k)}) \div \mu_{MS} - \mu_{MS} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})] - \lambda \prod_{k=1}^r (1 - p_{(e,k)}) \quad (12)$$

3) The average end-to-end delay, that is, the average time from the generation of a packet until it is received successfully at the MS (\bar{T}).

Here, we know that,

$$\bar{T} = \sum_{i=1}^r (W_{RS}^i + \frac{1}{\mu_{DF}}) + W_{MS} + \frac{1}{\mu_{MS}}$$

The average end-to-end delay is \bar{T} :

$$\sum_{i=1}^r \frac{1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]}{\mu_{DF} - \mu_{DF} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})] - \lambda \prod_{k=1}^{i-1} (1 - p_{(e,k)})} + \frac{1 - \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})]}{\mu_{MS} - \mu_{MS} \sum_{k=0}^r [p_{e,k+1} \cdot \prod_{j=1}^k (1 - p_{(e,j)})] - \lambda \prod_{k=1}^r (1 - p_{(e,k)})} \quad (13)$$

B. DF Relaying With Hop-By-Hop ARQ

1) Packet arrival rate at each RS and MS (λ_j):

$$\lambda_1 = \lambda + \lambda_1 \cdot p_{e,1}$$

$$\lambda_2 = \lambda_1 \cdot (1 - p_{e,1}) + \lambda_2 \cdot p_{e,2}$$

...

$$\lambda_r = \lambda_{r-1} \cdot (1 - p_{e,r-1}) + \lambda_r \cdot p_{e,r}$$

$$\lambda_{r+1} = \lambda_r \cdot (1 - p_{e,r}) + \lambda_{r+1} \cdot p_{e,r+1}$$

From above, we get equations for λ_i and λ_{r+1} :

$$\lambda_i = \frac{\lambda}{1 - p_{e,i}}$$

$$\lambda_{r+1} = \frac{\lambda}{1 - p_{e,r+1}} \quad (14)$$

2) For RS: Average queuing delay $W_{RS}^i = \frac{1}{\mu_{RS} - \lambda_i} - \frac{1}{\mu_{RS}}$:

$$W_{RS}^i = \frac{\lambda}{\mu_{DF} \cdot [\mu_{DF} \cdot (1 - p_{e,i}) - \lambda]} \quad (15)$$

As the utilization here is:

$$\rho_{RS}^i = \frac{\lambda_i}{\mu_{DF}} = \frac{\lambda}{\mu_{DF} \cdot (1 - p_{e,i})}$$

Hence, average number of packets at every RS is $N_{RS}^i = \frac{\rho_{RS}^i}{1 - \rho_{RS}^i}$:

$$N_{RS}^i = \frac{\lambda}{\mu_{DF} (1 - p_{e,i}) - \lambda} \quad (16)$$

For MS: Average queuing delay $W_{MS} = \frac{1}{\mu_{MS} - \lambda_i} - \frac{1}{\mu_{MS}}$:

$$W_{MS} = \frac{\lambda}{\mu_{MS} \cdot [\mu_{MS} \cdot (1 - p_{e,r+1}) - \lambda]} \quad (17)$$

As the utilization here is:

$$\rho_{MS} = \frac{\lambda_{r+1}}{\mu_{MS}} = \frac{\lambda}{\mu_{MS} \cdot (1 - p_{e,r+1})}$$

Hence, average number of packets at MS, $N_{MS} = \frac{\rho_{MS}}{1 - \rho_{MS}}$:

$$N_{MS} = \frac{\lambda}{\mu_{MS} \cdot (1 - p_{e,r+1}) - \lambda} \quad (18)$$

3) Here, we can write that

$$\bar{T} = \sum_{i=1}^r \frac{1 - p_{e,i}}{\mu_{DF} (1 - p_{e,i}) - \lambda} + \frac{1 - p_{e,r+1}}{\mu_{MS} (1 - p_{e,r+1}) - \lambda}$$

Average end-to-end delay \bar{T} :

$$\bar{T} = \sum_{i=1}^r \frac{1 - p_{e,i}}{\mu_{DF} (1 - p_{e,i}) - \lambda} + \frac{1 - p_{e,r+1}}{\mu_{MS} (1 - p_{e,r+1}) - \lambda} \quad (19)$$

3 STABILITY REGION

A. For AF relaying with end-to-end ARQ

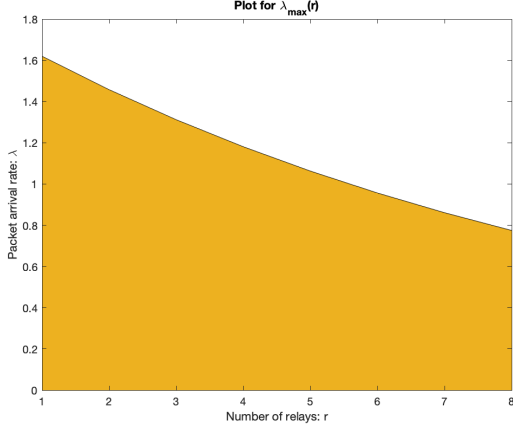


Figure 1: Plot of $\lambda_{max}(r)$ for AF relaying with end-to-end ARQ

In this case, $\lambda \leq 2 * 0.9^{r+1}$, so as r tends to infinity $2 * 0.9^{r+1}$ will tend to 0. As a result, $\lambda_{max}(r)$ will tend to 0 as well.

B. For DF relaying with end-to-end ARQ

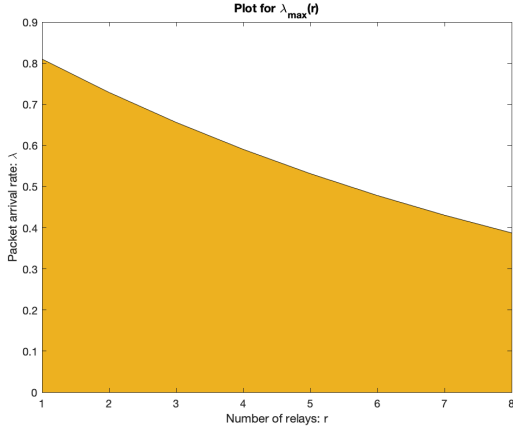


Figure 2: Plot of $\lambda_{max}(r)$ for DF relaying with end-to-end ARQ

4 ARRIVAL RATE-DELAY CHARACTERISTICS

A. For AF relaying with end-to-end ARQ

$$T = \frac{5 * 0.9^5}{2 * 0.9^5 - \lambda}$$

In this case, $(\lambda \leq 0.9^{r+1}) \cap (\lambda \leq 1.8)$, so when r tends to infinity $2 * 0.9^{r+1}$ tends to 0. As a result, $\lambda_{max}(r)$ will tend to 0 as well.

C. For DF relaying with hop-by-hop ARQ

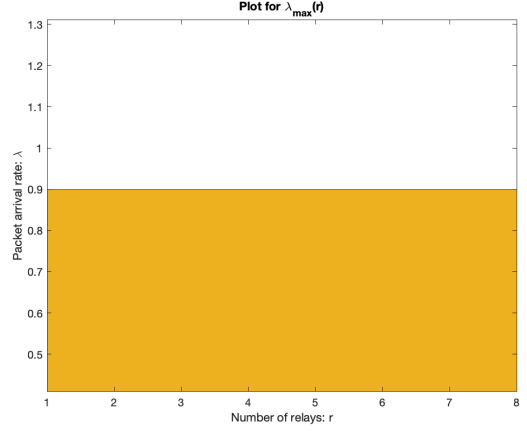
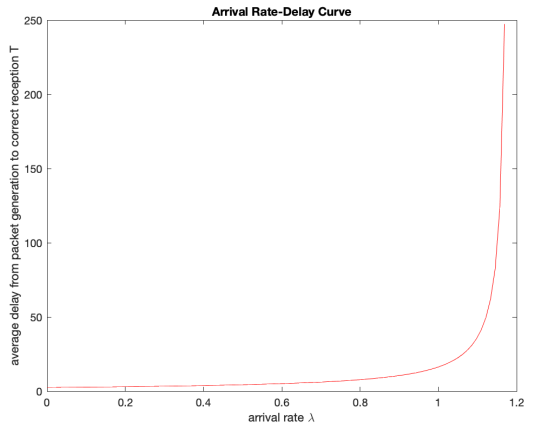


Figure 3: Plot of $\lambda_{max}(r)$ for DF relaying with hop-by-hop ARQ

In this case, $(\lambda \leq 0.9) \cap (\lambda \leq 1.8)$, as r tends to infinity, there is no influence on $\lambda_{max}(r)$



3 Figure 4: Plot of arrival-rate delay for AF relaying with end-to-end ARQ

From this plot we can observe that, as $p_{e,j}$ increases, the curve shifts to the left.

B. For DF relaying with end-to-end ARQ

$$T = \sum_{i=1}^4 \frac{0.9^5}{0.9^5 - \lambda * 0.9^{i-1}} + \frac{0.9^5}{2 * 0.9^5 - 0.9^4 * \lambda}$$

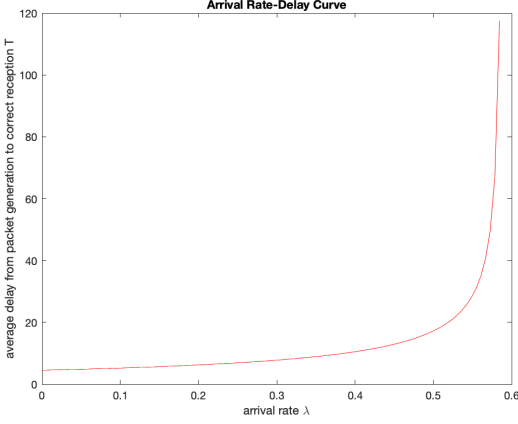


Figure 5: Plot of arrival-rate delay for DF relaying with end-to-end ARQ

From this plot we can observe that, as $p_{e,j}$ increases,

5 AF OR DF ?

A. Part I

We are to consider $p_{e,j} = 1, 2, \dots, r+1$ for this part. In this case, we can formulate the average system delays as follows: For AF relaying with end-to-end ARQ:

$$T_1 = 5$$

For DF relaying with end-to-end ARQ:

$$T_2 = \sum_{i=1}^4 \frac{(1-p)^5}{(1-p)^5 - 0.5 * (1-p)^{i+4}}$$

$$+ \frac{(1-p)^5}{2 * (1-p)^5 - 0.5 * (1-p)^9}$$

For DF relaying with hop-by-hop ARQ:

$$T_3 = \frac{26}{3}$$

From this, we can plot a graph for $0.05 \leq j \leq 0.95$.

the curve shifts to the left.

C. For DF relaying with hop-by-hop ARQ

$$T = \frac{3.6}{0.9 - \lambda} + \frac{0.9}{1.8 - \lambda}$$

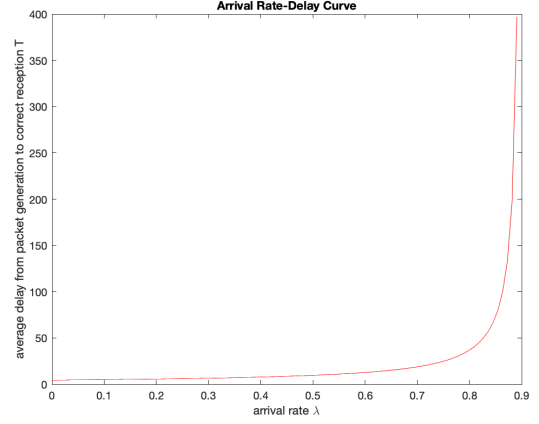


Figure 6: Plot of arrival-rate delay for DF relaying with hop-by-hop ARQ

From this plot we can observe that, as $p_{e,j}$ increases, the curve shifts to the left.

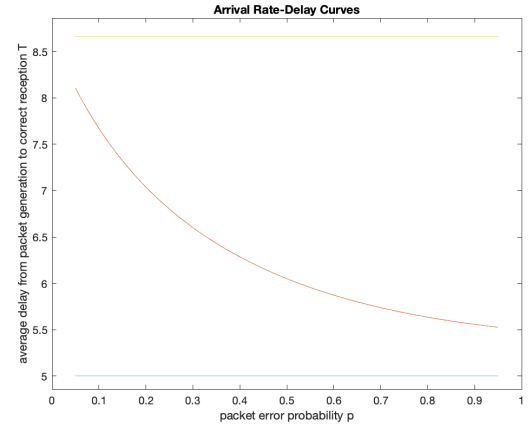


Figure 7: Plots of arrival-rate delay for the three systems

B. Part II

In this case, we can formulate the average system delays as follows:

For AF relaying with end-to-end ARQ:

$$T = \frac{4}{k-1} + 1$$

For DF relaying with end-to-end ARQ:

$$T = \sum_{i=1}^4 \frac{(0.9)^5}{(0.9)^5 - 0.5 * (0.9)^{i+4}} + \frac{(0.9)^5}{2 * (0.9)^5 - 0.5 * (0.9)^9}$$

For DF relaying with hop-by-hop ARQ:

$$T = \frac{26}{3}$$

From this, we can plot a graph where k is ranging from 1 to 4.

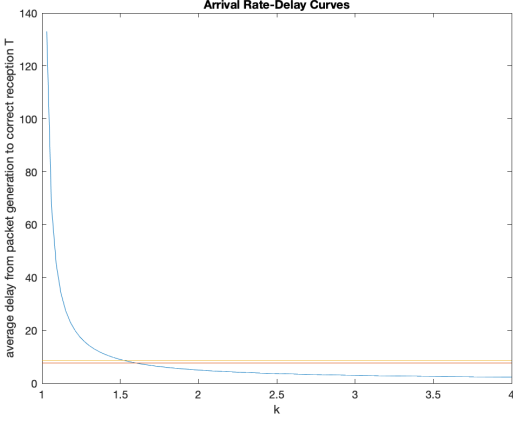


Figure 8: Plots of arrival-rate delay for the three systems

6 PROBABILISTIC ARQ

A. The P-ARQ queuing network model

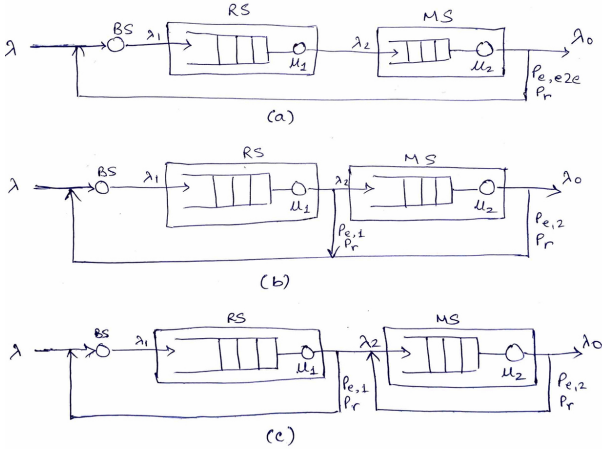


Figure 9: (a) For AF relaying with end-to-end ARQ, (b) For DF relaying with end-to-end ARQ, (c) For DF relaying with hop-by-hop ARQ

B. Performance evaluation

1) P-ARQ - AF relaying with end-to-end ARQ, T_{AF-ETE}^P :

$$\frac{1 - (1 - p_{e,1})[1 - (1 - p_{e,1})(1 - p_{e,2})]p_r}{\mu_{AF} - \mu_{AF}p_r[1 - (1 - p_{e,1})(1 - p_{e,2})] - \lambda} + \frac{1 - (1 - p_{e,1})[1 - (1 - p_{e,1})(1 - p_{e,2})]p_r}{\mu_{MS} - \mu_{MS}p_r[1 - (1 - p_{e,1})(1 - p_{e,2})] - \lambda} \quad (20)$$

2) P-ARQ - DF relaying with end-to-end ARQ, T_{DF-ETE}^P :

$$\frac{1 - p_{e,1}p_r - (1 - p_{e,1})p_{e,2}p_r}{\mu_{DF} - \mu_{DF}p_{e,1}p_r - \mu_{DF}(1 - p_{e,1})p_{e,2}p_r - \lambda} + \frac{(1 - p_{e,1})(1 - p_{e,1}p_r - (1 - p_{e,1})p_{e,2}p_r)}{\mu_{MS} - \mu_{MS}p_{e,1}p_r - \mu_{MS}(1 - p_{e,1})p_{e,2}p_r - \lambda(1 - p_{e,1})} \quad (21)$$

From the above graphs we can observe and therefore conclude that in the interval $1 \leq k \leq 1.6$, the most favorable choice would be DF relaying with end-to-end ARQ and in the interval of $1.6 \leq k \leq 4$, the system with AF relaying with end-to-end ARQ shows an advantage.

3) P-ARQ - DF relaying with hop-by-hop ARQ, T_{DF-HBH}^P :

$$\frac{1 - p_{e,1}p_r}{\mu_{DF} - \mu_{DF}p_{e,1}p_r - \lambda} + \frac{(1 - p_{e,2}p_r)(1 - p_{e,1}p_r)}{\mu_{MS}(1 - p_{e,2}p_r)(1 - p_{e,1}p_r) - (1 - p_{e,1})\lambda} \quad (22)$$

4) With the given information we can plot a graph and rewrite the above equations as follows:

$$T_{AF-ETE}^P = \frac{2 - 0.342 * p_r}{1.8 - 0.38 * p_r}$$

$$T_{DF-ETE}^P = \frac{1 - 0.19 * p_r}{0.8 - 0.19 * p_r} + \frac{0.9 - 0.171 * p_r}{1.82 - 0.38 * p_r}$$

$$T_{DF-HBH}^P = \frac{1 - 0.1 * p_r}{0.8 - 0.1 * p_r} + \frac{(1 - 0.1 * p_r)^2}{2 * (1 - 0.1 * p_r)^2 - 0.18}$$

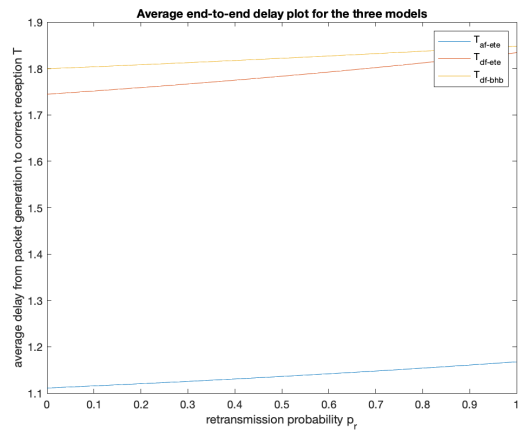


Figure 10: End-to-end Delays for the three models