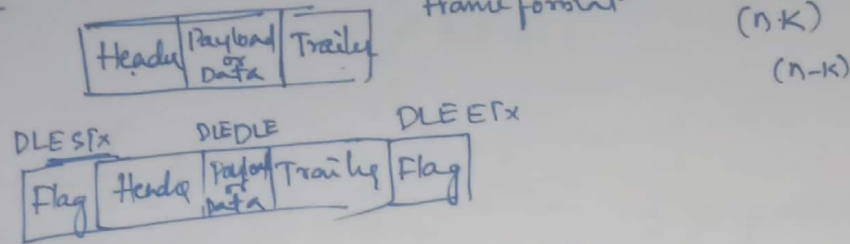


Datalink layer

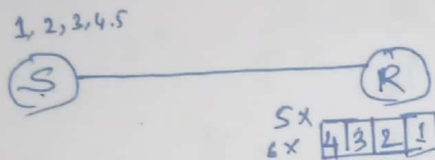
Function of DL

- 1) Framing ✓
- 2) flow control ✓
- 3) Error Control ✓

I Framing

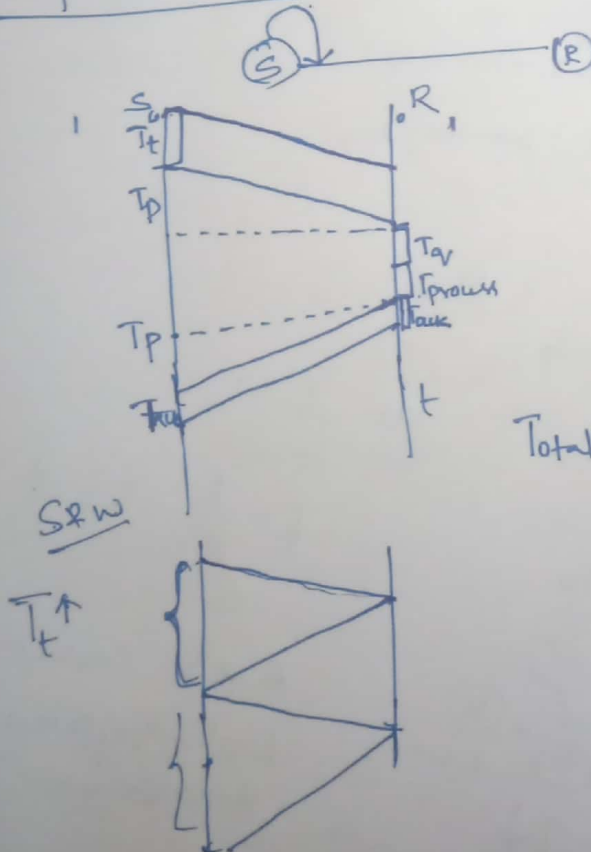
- 1) Bit Oriented → 01111110 → flag
- 2) Byte Oriented  
 ↳ DLE STX      DLE ETX

01111110  
Bit stuffing

Flow Control

load > Capacity of Buffer  
Receiver can't handle those packets

- Stop and wait protocol
- Sliding window protocol
- Go back N  
Selective repeat

Stop and wait
 $T_t = \text{Transmission time}$ 

$$T_t = \frac{L}{B}$$

 $T_p = \text{Propagation time}$ 

$$T_p = \frac{d}{v}$$

 $T_q \text{ or } T_{\text{Buffer}}$ 
 $T_{\text{process}}$ 

$$\text{Total Time} = T_{\text{total}} = T_t + T_p + T_q + T_{\text{process}} + T_{\text{ack}} + T_{\text{pau}}$$

$$= T_t + T_p + T_{\text{ack}} + T_p$$

$$= T_t + 2 * T_p + T_{\text{ack}}$$

$$T_{\text{total}} = T_t + 2 * T_p$$

$$\begin{aligned}\text{Efficiency } \eta &= \frac{\text{Useful time}}{\text{Total Cycle time}} \\ &= \frac{T_t}{T_t + 2 * T_p} \\ &= \frac{1}{1 + 2 * \frac{T_p}{T_t}}\end{aligned}$$

$$\text{Efficiency } \eta = \frac{1}{1 + 2a} \quad a = \frac{T_p}{T_t} \quad \text{Windowing sending } 1+2a \quad \square$$

$$\begin{aligned}\text{Throughput or Effective Bw or Bw utilization} &= \frac{L \text{ (Packet size)}}{T_t + 2 * T_p} \\ &= \frac{\frac{L}{B} * B}{T_t + 2 * T_p} \\ &= \frac{T_t}{T_t + 2 * T_p} * B \\ &= \eta * B \\ &= 0.5 * 100 \text{ Mbps}\end{aligned}$$

100 Mbps  
50%

Factor Effecting Throughput (or) Bw utilization

$$= \frac{T_t}{T_t + 2 * T_p} * B$$

$$= \frac{1}{1 + 2a}$$

$$= \frac{1}{1 + 2 * \frac{T_p}{T_t}}$$

$$= \frac{1}{1 + 2 * \frac{d}{v} * \frac{B}{L}}$$

$$= \frac{d \uparrow}{L \uparrow} \quad \frac{v \downarrow}{\eta \uparrow}$$

$$a = \frac{T_p}{T_t} \quad T_t = \frac{L}{B} \quad T_p = \frac{d}{v}$$

B & v. constant  
L → length of packet or size  
d → distance

A link has a transmission speed of  $10^6$  bits/sec. If we use data packet size of

- 1) calculate the efficiency for a stop & wait protocol for  $T_t = 1 \text{ ms}$  &  $T_p = 1 \text{ ms}$

$$\eta_t = \frac{1}{1+2a} \times 100 \quad a = \frac{T_p}{T_t} = \frac{1 \text{ ms}}{1 \text{ ms}} = 1 \quad T_p = 1 \text{ ms}$$

$$= \frac{1}{1+2} = \frac{1}{3} = 33\% \quad 67\%$$



Stop & wait protocol is good for LAN  
not for WAN

- 2) If the efficiency have to be 50% In case of s&w what is the relationship b/w.  $T_t$  &  $T_p$

$$\eta = \frac{T_t}{T_t + 2 \times T_p}$$

$$\frac{T_t}{T_t + 2 \times T_p} = 0.5 = 1/2$$

$$2 T_p = T_t + 2 \times T_p$$

$$T_t > 2 \times T_p$$

$$T_t = \frac{L}{B}$$

$$\frac{L}{B} > 2 \times T_p$$

$$\boxed{L > 2 \times T_p \times B}$$

If  $B = 4 \text{ Mbps}$   $T_p = 10 \text{ ms}$ . what is the length of the packet for an efficiency of 50%.

$$L > 2 \times T_p \times B$$

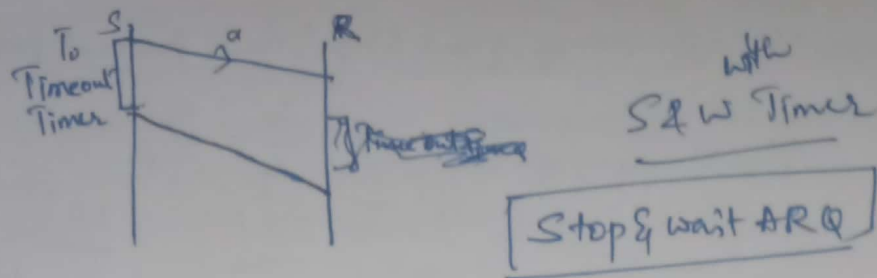
$$L > 2 \times 10 \times 10^{-3} \times 4 \times 10^6$$

$$\boxed{L > 8 \times 10^3 \text{ bits}}$$

$$\text{Throughput} = \eta \times B$$

$$= \frac{1}{2} \times 4 \text{ Mbps}$$

$$= \underline{\underline{2 \text{ Mbps}}}$$



In S&W protocol If we want to send 10 packets from S to R out of which every 4th packet being transmitted is lost  
How many packets we are sending

1 2 3 4 5 6 7 8 9 10  
↑

4 5 6 7 8 9 10  
↑

7 8 9 10  
↑

13 packets

10

Suppose that the stop and wait protocol is used on a link with a bitrate of 64 Kbits per second and 20 milliseconds Propagation delay. Assume that the transmission time for acknowledgement and processing time at the nodes are negligible. Calculate the minimum frame size in bytes to achieve a link Utilization of 50%.

$$L \geq 2 \times T_p \times B$$

$$L \geq 2 \times 20 \times 10^{-3} \times 64 \times 10^3$$

$$L \geq 2 \times \frac{20}{1000} \times 64 \times 10^3$$

$$L \geq 2560$$

$$\frac{2560}{8} \text{ Bytes}$$

$$= \underline{320 \text{ bytes}}$$



A link has a transmission speed of  $10^6$  bps. If user data packets are 8000 bytes each. Assume that the acknowledgement has negligible transmission delay and its propagation delay is the same as the data propagation delay also assume that processing delay at nodes are negligible. The efficiency of the stop and wait protocol in this setup is exactly 25%. Calculate the value of one-way propagation delay in milliseconds.

$$\eta = \frac{T_t}{T_t + 2 \times T_p} = 1/4$$

$$4T_t = T_t + 2 \times T_p$$

$$3T_t = 2 \times T_p$$

$$T_p = \frac{3 \times 8000}{10^6}$$

$$\underline{T_p = 24 \text{ ms}}$$

$$T_t = \frac{L}{B}$$

$$= \frac{8000 \times 8}{1000000}$$

$$T_t = 64 \text{ ns}$$

A channel has a bit rate of 40 kbps and one-way propagation delay of 20 ms. The channel uses stop and wait protocol. The transmission time of the acknowledgement frame is negligible. Calculate the minimum frame size required to get channel efficiency of at least 50%.

$$L \geq 2 \times T_p \times BW$$

$$L \geq 2 \times \frac{20}{1000} \times 40 \times 10^3$$

$$\underline{L \geq 1600 \text{ bits}}$$

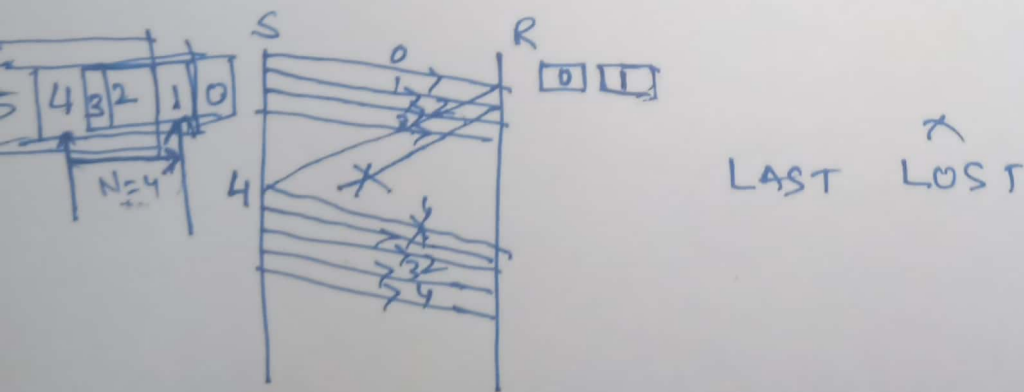
Consider two hosts 'x' and 'y' connected by a single direct link of rate  $10^6$  bits/sec. The distance between hosts is 10000 km. The Propagation speed is along the link is  $2 \times 10^8$  m/s. Host 'x' sends a file of 50000 bytes as one large message to host 'y' continuously. Let the transmission & Propagation delay be 'P' milliseconds & 'q' milliseconds respectively. Then Calculate the values of 'P' and 'q'.

## Sliding Window Protocol

- ↳ Go Back N
- ↳ Selective Repeat

1+2a

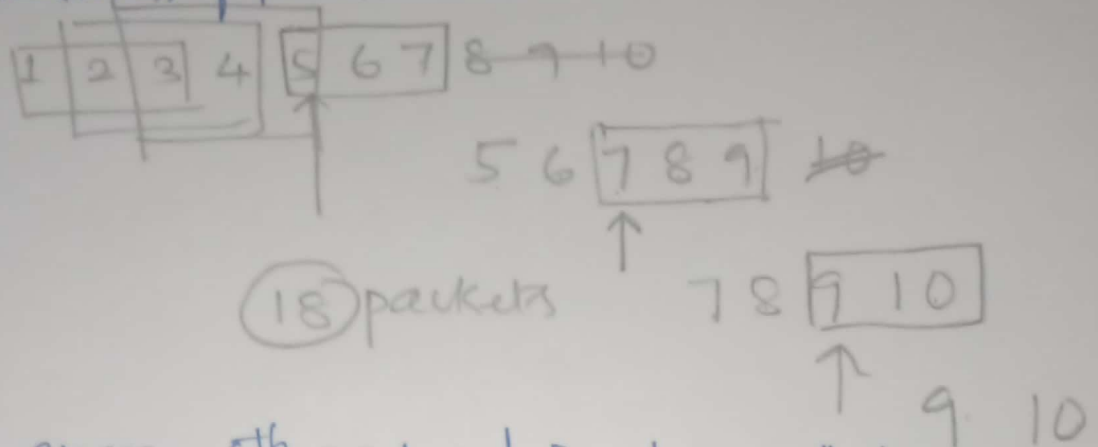
GB 4



In GB 3 if every 5<sup>th</sup> packet being transmitted is in total 10 packets. How many <sup>total</sup> transmissions of R

1 2 3 4 5 6 7 8 9 10

In GB3, If every 5<sup>th</sup> packet being transmitted is lost, in total of 10 packets  
 How <sup>many</sup> total transmissions of packets are done.



In GB4 If every 6<sup>th</sup> packet being transmitted is lost  
 It is desired to transmit 10 packets. How many total transmissions of packets are done.

