

Non-Linear Programming

Convex Programming Problems

Anand Jacob Abraham

Department of Industrial & Systems Engineering
IIT Kharagpur

1. Convex set and convex/concave functions
2. Convex Programming Problems and its Properties

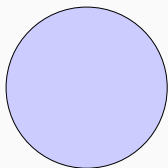
Convex set and convex/concave functions

Definition (Convex set)

The set S is convex if $\alpha \mathbf{x} + (1 - \alpha) \mathbf{y} \in S \forall \alpha \in (0, 1)$ and $\forall \mathbf{x}, \mathbf{y} \in S$.

- Basic geometric intuition of a convex set is that if you pick any distinct vectors in the set, it is possible to draw a straight line segment which can be drawn between these points without the line segment ever leaving the set.

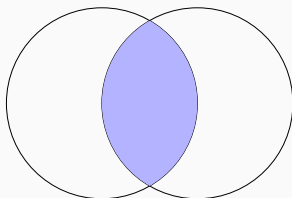
Convex sets examples



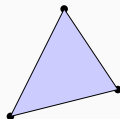
(a) Convex set (disk)



(b) Half-space



(c) Intersection of convex sets



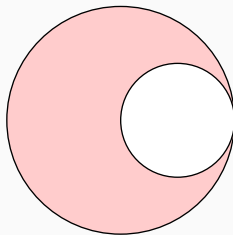
(d) Convex hull

Figure 1: Examples of convex sets

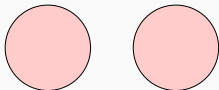
Non-convex sets examples



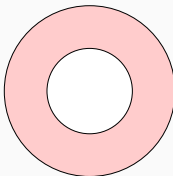
(a) Concave polygon



(b) Crescent shape



(c) Disconnected set



(d) Annulus (ring)

Figure 2: Examples of non-convex sets

Hypograph and Epigraph

The pair $(\mathbf{x}, l(\mathbf{x}))$ is called the graph of a function.

Definition (Hypograph)

The hypograph of a function $l: \mathcal{D} \mapsto \mathbb{R}$ is the locus of points/vectors $[\mathbf{x}, \beta]^T \in \mathcal{D} \times \mathbb{R} : \beta \leq l(\mathbf{x})$. Mathematically denoted by $\text{hyp}(l)$.

Definition (Epigraph)

The epigraph of a function $l: \mathcal{D} \mapsto \mathbb{R}$ is the locus of points/vectors $[\mathbf{x}, \beta]^T \in \mathcal{D} \times \mathbb{R} : \beta \geq l(\mathbf{x})$. Mathematically, denoted by $\text{epi}(l)$

Hypograph and Epigraph: Geometrically

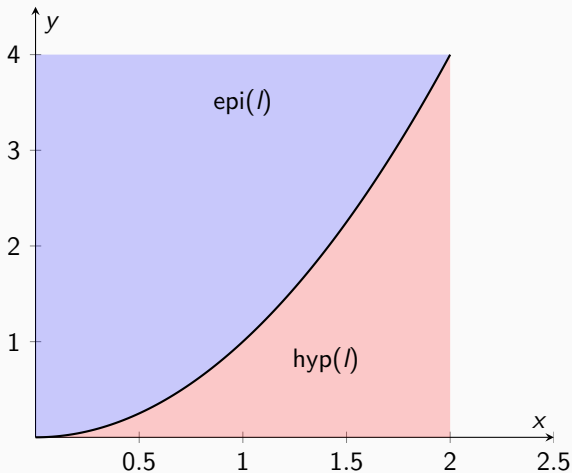


Figure 3: Graph, $\text{epi}(f)$, $\text{hyp}(f)$ of x^2 on $[0, 2]$

Convex and concave functions

Definition (Convex function)

A function $f: \mathcal{D} \mapsto \mathbb{R}$ is convex if $\text{epi}(f)$ is a convex set.

Definition (Concave function)

A function $f: \mathcal{D} \mapsto \mathbb{R}$ is concave if its $\text{hyp}(f)$ is a convex set.

Properties of convex function

1. $f: \mathcal{D} \mapsto \mathbb{R}$ is a convex function if
$$f(\alpha \mathbf{x} + (1 - \alpha) \mathbf{y}) \leq \alpha f(\mathbf{x}) + (1 - \alpha) f(\mathbf{y}) \forall \mathbf{x}, \mathbf{y} \in \mathcal{D}.$$
2. $f: \mathcal{D} \mapsto \mathbb{R}$ is a strict convex function if
$$f(\alpha \mathbf{x} + (1 - \alpha) \mathbf{y}) < \alpha f(\mathbf{x}) + (1 - \alpha) f(\mathbf{y}) \forall \mathbf{x}, \mathbf{y} \in \mathcal{D}.$$
3. $f: \mathcal{D} \mapsto \mathbb{R}$ is a concave function if
$$f(\alpha \mathbf{x} + (1 - \alpha) \mathbf{y}) \geq \alpha f(\mathbf{x}) + (1 - \alpha) f(\mathbf{y}) \forall \mathbf{x}, \mathbf{y} \in \mathcal{D}.$$
4. $f: \mathcal{D} \mapsto \mathbb{R}$ is a strict concave function if
$$f(\alpha \mathbf{x} + (1 - \alpha) \mathbf{y}) > \alpha f(\mathbf{x}) + (1 - \alpha) f(\mathbf{y}) \forall \mathbf{x}, \mathbf{y} \in \mathcal{D}.$$
5. If f_1, f_2 are concave functions, then $f_1 + f_2$ is also concave.
6. If $f: \mathcal{D} \mapsto \mathbb{R}$ is strictly convex, then $-f$ is strictly concave.
7. If a function $f: \mathcal{D} \mapsto \mathbb{R}$ is convex/concave, then the set \mathcal{D} is a convex set.

Alternate definitions for concave function

- $f: \mathcal{D} \mapsto \mathbb{R}$ is a concave function if and only if the Hessian of the function is negative semi-definite.
- $f: \mathcal{D} \mapsto \mathbb{R}$ is a concave function if and only if
$$f(\alpha \mathbf{u} + (1 - \alpha) \mathbf{v}) \geq \alpha f(\mathbf{u}) + (1 - \alpha) f(\mathbf{v}) \forall \mathbf{u}, \mathbf{v} \in \mathcal{D}$$
- Similarly, we can get results for convex functions also.

Convex Programming Problems and its Properties

Convex Programming Problem

Definition (Convex Programming Problem)

A Convex Programming Problem (CPP) is an optimization problem

$\max_{\mathbf{x} \in S} f(\mathbf{x})$ where f is concave or convex and S is a convex set.

- If the objective function is concave and the problem is of maximization type with convex constraint set, some books/literature calls this type of problems as *Concave Maximization Problems* (CMP).

We now show three important results (theorems) for CPPs/CMP. There are some intermediate Lemmas as well (used to derive other main results).

Result 1: Local maximizers are global

Theorem

\mathbf{x}^* is a local maximizer to a CPP $\max_{\mathbf{x} \in S} f(\mathbf{x})$ if and only if it is a global maximizer.

Proof.

Suppose \mathbf{x}^* is a local maximizer which is not a global maximizer.

$$\exists \mathbf{y} \in S : f(\mathbf{y}) > f(\mathbf{x}^*)$$

, then as $\mathbf{y} \in S$ and due to convexity of S and concaveness of f , we may write

$$f(\alpha \mathbf{y} + (1 - \alpha) \mathbf{x}^*) \geq \alpha f(\mathbf{y}) + (1 - \alpha) f(\mathbf{x}^*) \forall \alpha \in (0, 1)$$

$$\implies f(\alpha \mathbf{y} + (1 - \alpha) \mathbf{x}^*) \geq \alpha f(\mathbf{y}) + (1 - \alpha) f(\mathbf{x}^*) > f(\mathbf{x}^*) \forall \alpha \in (0, 1)$$

$$\implies f(\alpha \mathbf{y} + (1 - \alpha) \mathbf{x}^*) > f(\mathbf{x}^*) : \alpha > 0, \alpha \rightarrow 0$$

This means that there exists a vector $\alpha \mathbf{y} + (1 - \alpha) \mathbf{x}^*$ which in the neighborhood of \mathbf{x}^* which is better than \mathbf{x}^{*1} . This contradicts the local optimality of \mathbf{x}^* . Which means that if \mathbf{x}^* is locally optimal then it also has to be globally optimal in a CPP.

The only if part of the proof is trivial.



¹This point can be seen visualized as a point on the line connecting \mathbf{x}^* and \mathbf{y} on the set S . Every point on this line guarantees a better objective function value owing to the concavity of the function. For example, the line is feasible and contains \mathbf{x}^* and

A less important short result

Lemma

If $f: \mathcal{D} \mapsto \mathbb{R}$ is concave then $\forall \mathbf{u}, \mathbf{v} \in \mathcal{D}$

$$f(\mathbf{v}) \leq f(\mathbf{u}) + (\mathbf{v} - \mathbf{u})^T \nabla f(\mathbf{u})$$

Proof.

As the function is concave in \mathcal{D} , we have

$$f(\alpha \mathbf{v} + (1 - \alpha) \mathbf{u}) \geq \alpha f(\mathbf{u}) + (1 - \alpha) f(\mathbf{v}) \forall \alpha \in (0, 1).$$

On simplifying the above expression, we get the following.

$$f(\mathbf{u} + \alpha(\mathbf{v} - \mathbf{u})) \geq f(\mathbf{u}) + \alpha(f(\mathbf{v}) - f(\mathbf{u}))$$

Upon rearrangement we get

$$f(\mathbf{u}) + \frac{f(\mathbf{u} + \alpha(\mathbf{v} - \mathbf{u})) - f(\mathbf{u})}{\alpha} \geq f(\mathbf{v})$$

Applying the limit $\alpha \rightarrow 0$, we get the following required result.

$$f(\mathbf{u}) + (\mathbf{v} - \mathbf{u})^T \nabla f(\mathbf{u}) \geq f(\mathbf{v})$$

□

Result 2: Sufficiency of first order condition in Unconstrained optimization

Theorem

If \mathbf{x}^* is a stationary point for the optimization problem $\max_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x})$, then \mathbf{x}^* is a global maximizer.

Proof.

As \mathbf{x}^* is a stationary point $\nabla f(\mathbf{x}^*) = \mathbf{0}$. Assume that \mathbf{x}^* is not the global maximizer, then $\exists \mathbf{y} \in \mathbb{R}^n : f(\mathbf{y}) > f(\mathbf{x}^*)$. Using the previous lemma 13 we can write the following.

$$f(\mathbf{y}) \leq f(\mathbf{x}^*) + (\mathbf{y} - \mathbf{x}^*)^T \nabla f(\mathbf{x}^*)$$

As \mathbf{x}^* is a stationary point $\nabla f(\mathbf{x}^*) = \mathbf{0}$ we get the following $f(\mathbf{y}) \leq f(\mathbf{x}^*)$. However, this contradicts our earlier assumption that \mathbf{x}^* is not a global maximizer. Therefore \mathbf{x}^* is a global maximizer. \square

Result 3: Sufficiency of KKT

Theorem

If $f: \mathbb{R}^n \mapsto \mathbb{R}$ is concave over the convex set

$S = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}$ and $\mathbf{x}^ \in S$ satisfies the KKT conditions below*

1. $\nabla f(\mathbf{x}^*) = Dh(\mathbf{x}^*)^T \lambda^* + Dg(\mathbf{x}^*)^T \mu^*.$
2. $\mu^* \mathbf{g}(\mathbf{x}) = 0.$
3. $\mu^* \geq \mathbf{0}.$

then \mathbf{x}^ is a global maximizer.*