ocailleu willi caM

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	Linked list
	Suppose that in our OS we need to make
	an array of 10 MB size but the memory
	available is 10 MB but is not contiguous & hence we can't use arrays here but linked list can work here & hence it can work
	list converse due arrays here but linked
	on non-contiguous memory locations.
	The string your membry xocations.
2)	At run-time, we can do creation & deletion
	The state of the s
	would be no wastage of memory which was there in case of arrays.
	was there in case of average.
3)	In anyone inconting
	In arrays, insertion of element takes  O(n) time complexity but in linked list  insertion (an he done is (2011)
	insurtion can be done in O(1) time complexity
	I at the acities
	only.
4)	They a list by a comment of
	of linked list but consultating in case
	There is no concept of indexing in case of linked list but concept of address is used here.
	Definition
	hinked list can be defined as collection
	of nodes.  Node
	1 2
	1-data
	27 address of next node
11	

	Simple linked list
	104
	3 104 -> 4 -> Null
	class Node {
	int data ;
	Node* next;
	2
	As we create an integer pointer by int *,
	We create pointer to Node by Node*.
	Types of linked list Singly linked list
1)	Singly linked list
	$(1) \longrightarrow (2) \longrightarrow (3) \longrightarrow X$
	1 × 2 111 (B
5)	Circular linked list
	1 Ail to the second
3)	Doubly linked list ( frev & next pointers)
	$\times \leftarrow 1 \longrightarrow 2 \longrightarrow \times$
-	
4)	Doubly circular linked list (prev & next
	pointers)
	(1) (3)
1)	
Ote-	Linked list is Hindi & this is a magical
	line but don't tell ins to any
	Interviewer. Also linked list is a linear
	data structure as only one descendant is there.
	. 71 ATTICLE

Creation of Linked List (Better method is
also there)
Class Node {
public:
int data : "
Node* next;
Node() {
this -data = 0;
this I next = NULL;
<u>9</u>
Node (int data) {
this -data = data;
this I next = NULLi
3;
Node* first = 200 Node
- TOW # SP(D) A = 10 A 1 A 1 A 1 A 1 A 1 A 1 A 1 A 1 A 1
Node * third = new Node (3);
first -) next = second;
second + next = third;
-
Printing a linked list
1) Print the data of convent node
2) Move pointer forward.
Printing a linked list  1) Print the data of convert node.  2) Move pointer forward.  3) Stop When we reach null.
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	void print (Node* & head) {
	// Temperoscy pointer
	Node* temp = head;
	while (temp!=NULL) {//Step-3
	11 Steb-1
	cout << temp → data << " ";
	//Step-2 temp = temp = next;
	3
	3
	Output
	1 2 3
	Meaning of temp = temp - next
	head
	head $\uparrow$ $\uparrow$ $\uparrow$ $\uparrow$ $\uparrow$
	temp temp next
	temp = temp + next
	= head + next
	II described in the
	Hence we are moving forward in the
	linked list.
	1.11 to d temp 2
	Why we created temp?  It is a good practice not to change the
	bointing to head & hience we move temp forward leaving head at same place.
	pointing to head at same blace.
	Torward reaving head of sain
X7	
,	
	$\frac{1}{1+\exp \theta}$
	temp - next - data = 2 temp - next - next - data = 3
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	Betler way of creating linked list
(1)	Insertion at head
	Initial scenario -> (10) -> X Inserting 20 at the head
l)	Creating node with value 20
	$\begin{array}{cccccccccccccccccccccccccccccccccccc$
	newNode head
ર)	new Node to be connected to head
	$\begin{array}{c} \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \longrightarrow X \\ \nearrow \bigcirc \bigcirc \bigcirc \bigcirc \longrightarrow X \\ \end{array}$
3)	NewNode head Update head
	$\begin{array}{c} \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \rightarrow \times \\ \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \rightarrow \times \\ \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \rightarrow \times \\ \bigcirc \bigcirc$
	void insert (Node * & head, int data) { //Step-1
)	Node * new Node = new Node (data)
	newNode + next = head; //Step-3
	head = newNode;
	Suppose now we inserted 20,30,40 &50

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	using the function, then the linked list would become.
	pecome.
	$(50)\rightarrow (40)$ $(50)$
	$\begin{array}{c} (SO) \rightarrow (4O) \rightarrow (3O) \rightarrow (2O) \rightarrow (1O) \rightarrow \times \\ (SO) \rightarrow (4O) \rightarrow (3O) \rightarrow (2O) \rightarrow (1O) \rightarrow \times \\ (SO) \rightarrow (4O) \rightarrow (3O) \rightarrow (2O) \rightarrow (1O) \rightarrow \times \\ (SO) \rightarrow (4O) \rightarrow (3O) \rightarrow (2O) \rightarrow (1O) \rightarrow \times \\ (SO) \rightarrow (4O) \rightarrow $
	Printing linked list will give the output as 50 40 30 20 10
	<u>us 50 70 30 20 10</u>
Val	
NOE	nead - Starting point of linked list
	tail = ending point of linked list
	point = node here.
( ) )	
(1)	Insertion at tail
	Initial scenario -> (10)-1 X
	Inserting 20 at tail
,	d i u
- 1)	Creating node with value 20
	(tail
	$(20) \rightarrow \times$ $(10) \rightarrow \times$
	I new Noode head
3)	new Node to be connected to the tail
	l tail
	$(10) \rightarrow (20) \rightarrow \times$
	head ?
3)	Update tail
	$(10) \rightarrow (20) \rightarrow \times$
	heod I tail
	Void insert (Node * & head, Node * & tail,
	int data) {

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	// Step-1 Node * new Node = new Node (data); // Step-2
	tail - next = new Node; //Step-3 tail = new Node;
	9 9 finitially head & tail was initialized
	to NULL, then we need to handle this explicitly as when we are creating the new node, initialize the head & tail with that node & this would be when we are creating the 1st node.
	if (head == NULL) {  head = new Node;  tail = new Node;  3
	Betler way & easy to understand for handling empty linked list case.
	if (head = = NULL) {  Node * new Node = new Node (Data);  head = new Node;  tail = new Node;  return; // As node has been created;  3
<u>(iii)</u>	Inserting at specified position.
(i <i>ii</i> )	

.1	
1)	Check for empty linked list case.
3)	or non-empty men
	1) Traverse to that position.
	(ii) (reate a node
	PC
1	$\bigcirc \rightarrow \bigcirc \times \bigcirc \rightarrow \times$
	2nd 1st
	newNode
	(iii) new Node + next = c ;
	(iv) p → next = new Node;
	What if we do prinext = new Node first
	and then update new Node - next to cury
	but by this, we will be losing the track
	of further linked list.
	Some edge cases  pos = = 0 → insert At Head ? Reusing funct.
1)	pos = = 0 - insert At Head ( Rewsing funct.
5)	pos = = len + insert At Tail J
	length of linked list int find Length (Node * & head) {
	int find Length (Node * & head)?
	Int len = 1 j
	Node * temp = head;
	While (temp - next ] = NULL) {
	temp = temp - next;
	len + + j
	3
	return len;
	3

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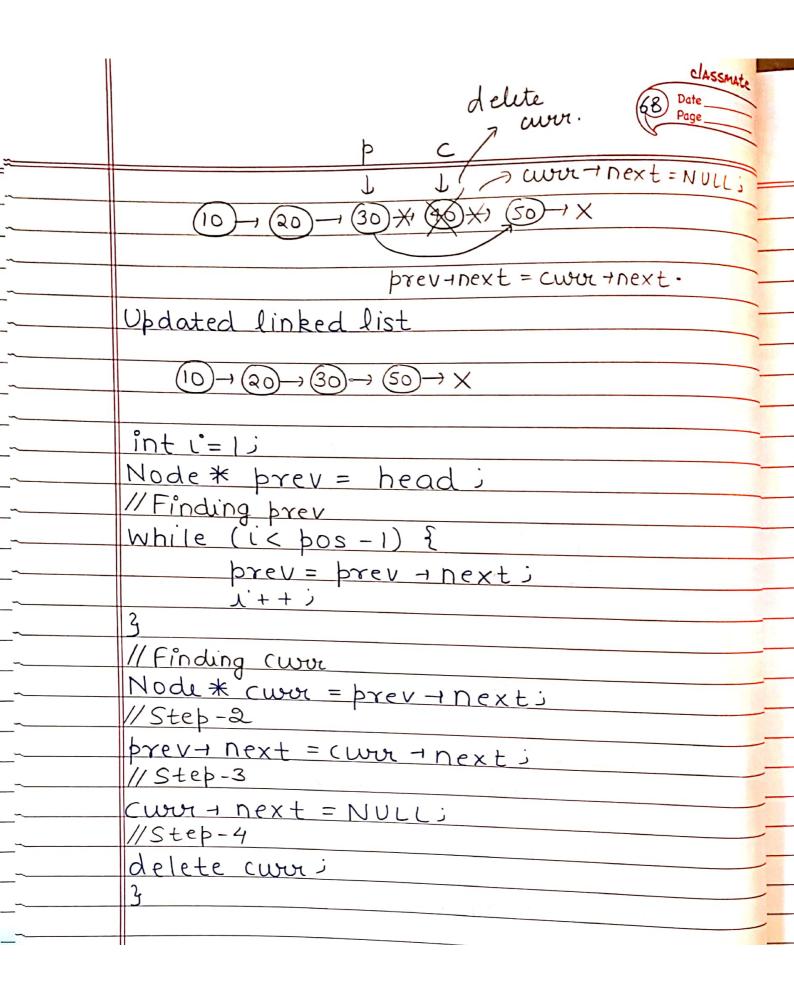
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	rage
	Deletion operation in Linked list
(i)	Deleting head of linked list
	10-> 20-> 30-> X temp= head
	ead = head Inext
	temp - next = null; delete temp
	$(20) \rightarrow (30) \rightarrow X$
(ii)	Deleting tail in linked list
_ +1	10 -> (30) -> x temp = tail
	Find prev
	prev-i next = null tail = prev delete temp
	(10)-1 (20)-1 X
	Void deletell(int pos, Node * & head, Node * 4 tail) {
	// Empty Linked List if (head = = NULL) {
	3
	// Delete head if (pos ==1)
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```
Node * temp = head;
head = head + next; //s-1
      temp - next = NULLills-2
delete temp ills-3
      // Deleting tail
int len = findLength (head);
if (pos = = len) {
// Find prev (S-1)
            int (= 1)
            Node* prev = head i
           while (i < pos-1) {
                 prev= prev nexti
         115-2
        prev + next = NULLi
        //5-3
         tail = previ
       //5-4___
        delete temp;
(iii) Deleting middle node (specified position)
2) Do prev - next = cwor - next >
3) Cwri-next = NULL;
4) de lete como
```



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