Decentralized Algorithms for Generalized Nash Equilibrium Seeking in Peer-to-peer Electricity Markets

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1 Problem statement

Consider the same problem statement in problem_setup_v6.tex.

2 Algorithms

2.1 Semi-decentralized

The algorithm presented in this section corresponds to [1, Alg. 6B] applied on problem_setup_v6.tex. Consider the following choices for the step-sizes in Algorithm 1.

Assumption 1 (Step-size) Set the step sizes of prosumers and DSO as follows:

```
(i) \forall i \in \mathcal{N}: set A_i = \operatorname{diag}(\alpha_i^{pi}, \alpha_i^{st}, \alpha_i^{mg}, \{\alpha_{(i,j)}^{tr}\}_{j \in \mathcal{N}_i}) \otimes I_H, with \alpha_i^{pi}, \alpha_i^{st} > 1, \alpha_i^{mg} > 3 + N \max_{h \in \mathcal{H}} d_h, \beta_{(i,j)}^{tr} = \beta_{(i,j)}^{tr} < \frac{1}{2}, \forall j \in \mathcal{N}_i.
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(ii) Set A_{N+1} := \operatorname{diag}\left(\left\{\alpha_y^{\theta}, \alpha_y^{v}, \alpha_y^{tg}, \{\alpha_{(y,z)}^{p}, \alpha_{(y,z)}^{q}\}_{z \in \mathcal{B}_y}\right\}_{y \in \mathcal{B}}\right) \otimes I_H, with \alpha_y^{\theta}, \alpha_y^{v} > 0, \alpha_y^{tg} > 2, \alpha_{(y,z)}^{p} > 1 and \alpha_{(y,z)}^{q} > 0, \forall z \in \mathcal{B}_y, \forall y \in \mathcal{B}. Set \gamma^{mg} < \frac{1}{N}, \beta^{tg} < (|\mathcal{N}| + |\mathcal{B}|)^{-1} and \beta_y^{pb} < (1 + |\mathcal{N}_y| + |\mathcal{B}_y|)^{-1}, for all y \in \mathcal{B}.
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Algorithm 1 Euclid's algorithm

```
1: procedure Euclid(a, b)
                                                                                                        ▶ The g.c.d. of a and b
        r \leftarrow a \bmod b
        while r \neq 0 do
3:
                                                                                                \triangleright We have the answer if r is 0
            a \leftarrow b
4:
            b \leftarrow r
5:
            r \leftarrow a \bmod b
6:
        end while
7:
                                                                                                                    \triangleright The gcd is b
        return b
9: end procedure
```

Algorithm 2 Semi-decentralized GWE seeking for P2P Energy Markets

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1: procedure Initialization prosumers
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- for all prosumer $i \in \mathcal{N}$ do 2:
- Set the initial conditions: $u_i(0) \in \mathcal{U}_i$, $\mu_{(i,j)}^{\mathrm{tr}}(0) = \mathbf{0}$, $\forall j \in \mathcal{N}_i$. 3:
- Set the step sizes as in Assumption 1(i). 4:
- 5: end for
- 6: end procedure
- procedure Initialization DSO
- Set the initial conditions: $u_{N+1}(0) \in \mathcal{U}_{N+1}$, $\lambda^{\text{mg}} = \mathbf{0}$, $\mu^{\text{tg}} = \mathbf{0}$, $\mu^{\text{pb}}_{y}(0) = \mathbf{0}$, $\forall y \in \mathcal{B}$.
- Set the step sizes as in Assumption 1(ii).
- 10: end procedure
- 11: repeat
- for all prosumer $i \in \mathcal{N}$ do 12:
- procedure Strategy update 13:

14:
$$a_i(k) = \operatorname{col}\left(-\mu_y^{\operatorname{pb}}(k), -\mu_y^{\operatorname{pb}}(k), \begin{bmatrix} I_H \\ -I_H \end{bmatrix}^\top \lambda^{\operatorname{mg}}(k) + \mu^{\operatorname{tg}}(k), \left\{\mu_{(i,j)}^{\operatorname{tr}}(k)\right\}_{j \in \mathcal{N}_i}\right) \triangleright (\operatorname{aux. vector})$$

15:
$$u_{i}(k) = \operatorname{cor}(\begin{array}{c} \mu_{\tilde{y}}^{-}(k), & \mu_{\tilde{y}}^{-}(k), \begin{bmatrix} -I_{H} \end{bmatrix} \\ \chi^{-}(k) + \mu^{-}(k), \begin{bmatrix} \mu_{(i,j)} \\ \chi^{-}(k) \end{bmatrix} \\ u_{i}(k+1) = \begin{cases} \operatorname{argmin} & J_{i}(\xi, \sigma(k)) + a_{i}(k)^{\top} \xi + \frac{1}{2} \|\xi - u_{i}(k)\|_{A_{i}}^{2} \\ \operatorname{s.t.} & \xi \in \mathcal{U}_{i} \end{cases}$$
16: and procedure

- 16: end procedure
 - (1) Strategy update

$$u_i(k+1) = \begin{cases} \underset{\xi \in \mathbb{R}^{n_i}}{\operatorname{argmin}} & J_i(\xi, \sigma(k)) + a_i(k)^\top \xi + \frac{1}{2} \|\xi - u_i(k)\|_{A_i}^2 \\ \text{s.t.} & \xi \in \mathcal{U}_i \end{cases}$$

where

$$a_i(k) = \operatorname{col}\left(-\mu_y^{\operatorname{pb}}(k), -\mu_y^{\operatorname{pb}}(k), \begin{bmatrix} I_H \\ -I_H \end{bmatrix}^\top \lambda^{\operatorname{mg}}(k) + \mu^{\operatorname{tg}}(k), \left\{\mu_{(i,j)}^{\operatorname{tr}}(k)\right\}_{j \in \mathcal{N}_i}\right)$$

(2) Communication with trading partners and DSO:

(a)
$$\left\{ \begin{array}{l} p_i^{\text{mg}}(k+1) \\ b_i(k+1) := p_i^{\text{d}} - p_i^{\text{di}}(k+1) - p_i^{\text{st}}(k+1) \end{array} \right\} \longrightarrow \text{DSO},$$
(b)
$$p_{(i,j)}^{\text{tr}}(k+1) \longrightarrow \text{prosumer } j, \quad \text{for all } j \in \mathcal{N}_i.$$

(3) Dual variables (reciprocity constraints) update, $\forall j \in \mathcal{N}_i$:

$$\mu_{(i,j)}(k+1) = \mu_{(i,j)}(k) + \beta_{ij}^{\text{tr}} \left(2c_{(i,j)}^{\text{tr}}(k+1) - c_{(i,j)}^{\text{tr}}(k) \right).$$

where

$$c_{(i,j)}^{\text{tr}}(k+1) = p_{(i,j)}^{\text{tr}}(k+1) + p_{(j,i)}^{\text{tr}}x(k+1),$$

17: end for

DSO.

(1) Aggregate energy demand update

$$\begin{split} \sigma^{\text{mg}}(k+1) &= \sum_{i \in \mathcal{N}} p_i^{\text{mg}}(k+1) + \sum_{i \in \mathcal{B}} p_k^{\text{pd}} \\ \sigma^{\text{tg}}(k+1) &= \sum_{\sigma \in \mathcal{B}} p_y^{\text{tg}}(k+1) \end{split}$$

(2) Dual variable (grid constraints) update:

References

[1] G. Belgioioso and S. Grammatico, "Semi-decentralized generalized nash equilibrium seeking in monotone aggregative games," $arXiv\ preprint\ arXiv:2003.04031$, 2020.