

ASSIGNMENT COVER SHEET

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This coversheet must be attached to the front of your assessment

The assessment is due on 28/08/2023, 9:05 AM unless otherwise specified in the course outline.

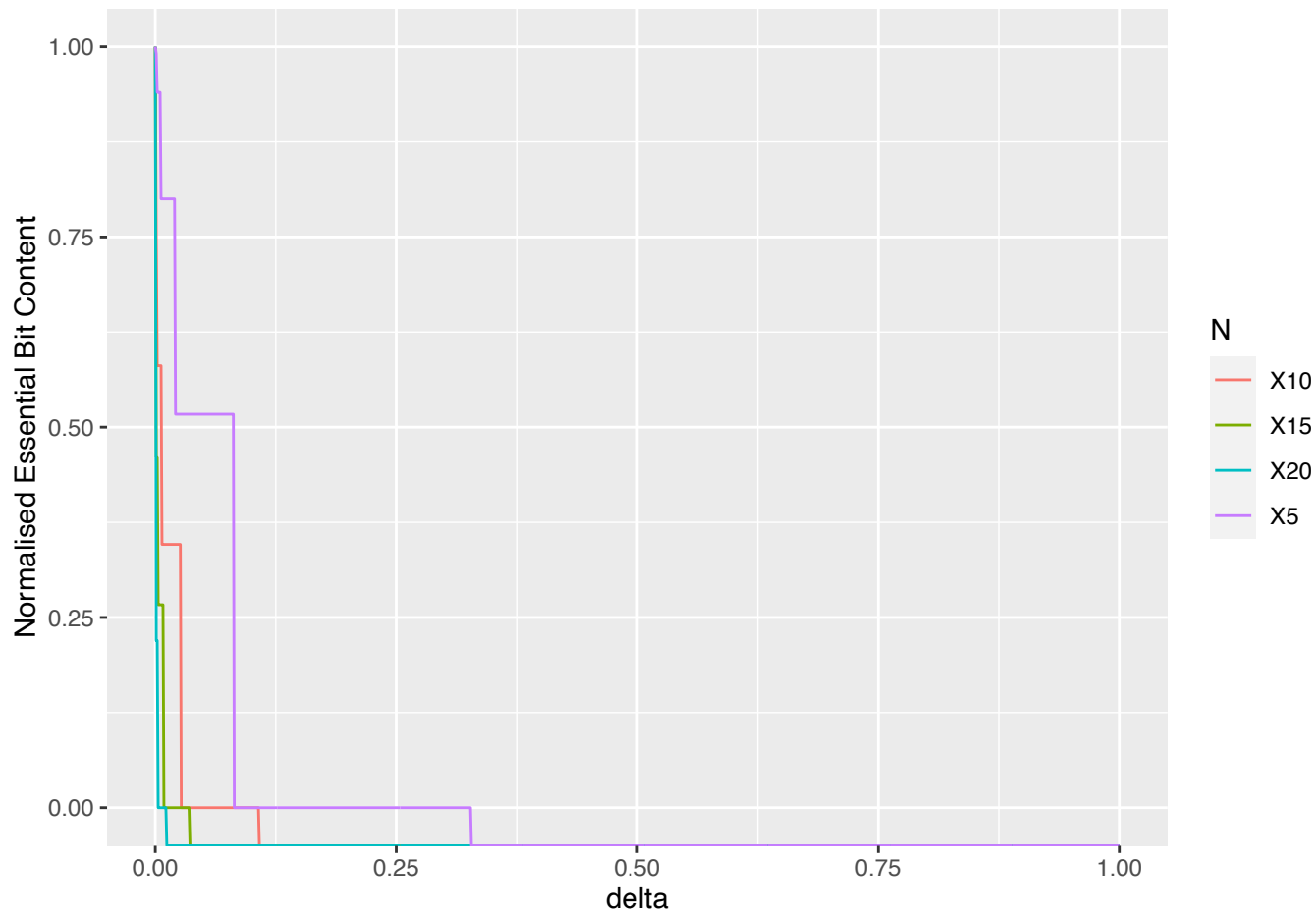
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Course Code	COMP6261
Course Name	Information Theory
Assignment Item	Assignment 2
Due Date	25/09/2023,
Date Submitted	23/09/2023

I declare that this work:

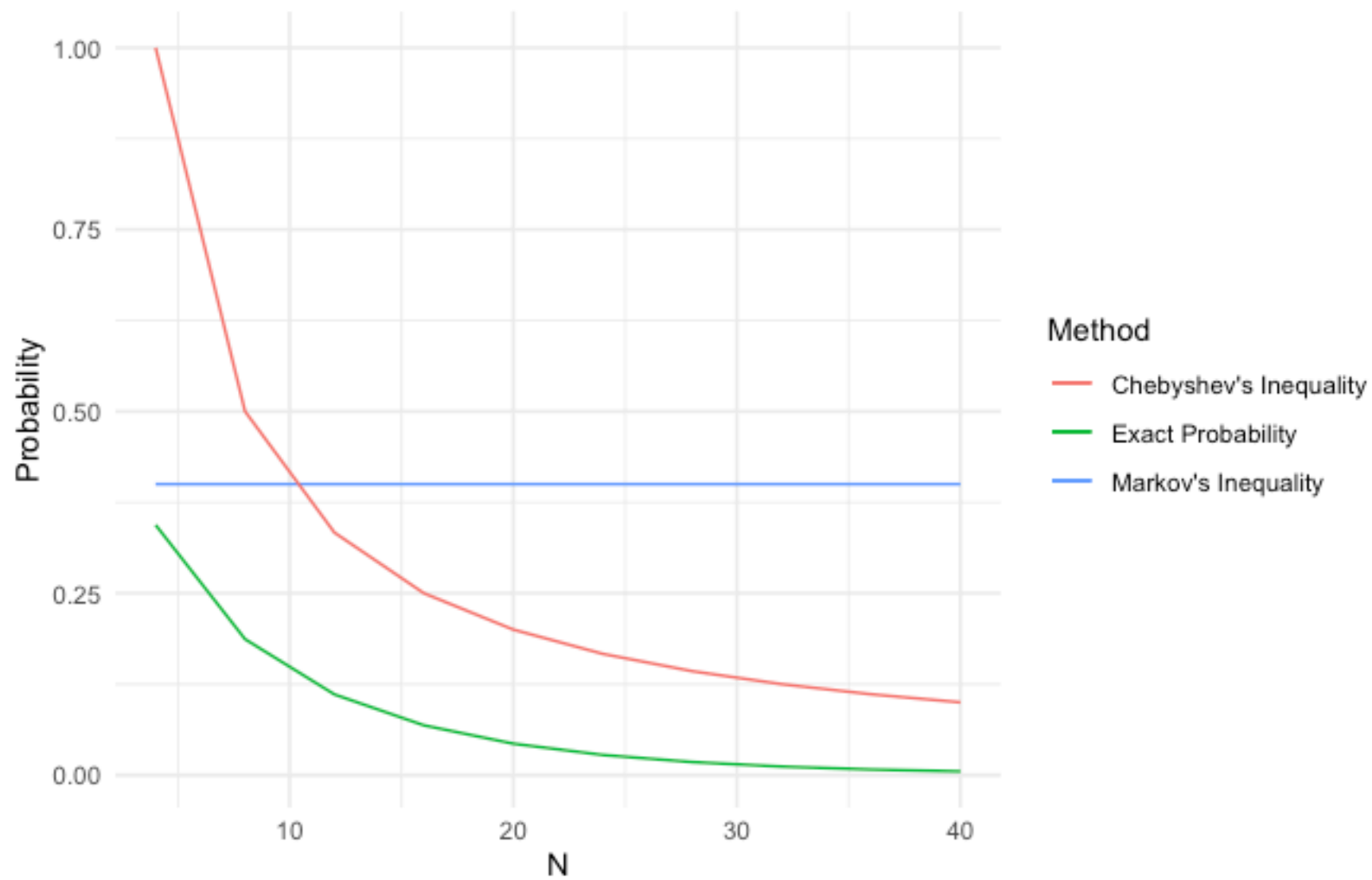
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Q3.e) Essential Bits given delta



Bounds and Exact Probabilities for $N/4$ or More Heads



Question 1Question 1(a)

$$1) p(X \geq a) \leq \frac{E[X]}{a} \text{ given } E[X] = 2000, a = 2400$$

$$p(X \geq \frac{2400}{2400}) \leq \frac{2000}{2400} \text{ or } \frac{5}{6} \quad \#$$

$$2) P(|X - \mu| \geq k \times \sigma) \leq \frac{1}{k^2} \text{ given } \mu = 2000, \sigma = 200 \text{ Hence, } k = \frac{2400 - 2000}{200} = 2$$

$$p(1600 \leq X \leq 2400) \geq 1 - \frac{1}{2^2} \text{ or } \frac{3}{4} \quad \#$$

Question 1(b)

$$1) p(X = \text{head}) < \frac{1}{8}, \text{ coin flipped } N \text{ times } N \in \mathbb{N}$$

$$p(X \geq \frac{N}{4}) \leq \frac{E[X^N]}{N/4} = \frac{1}{2} \quad \#$$

$$2) p(X \geq \frac{N}{4}) < \frac{1}{8}, \text{ coin flipped } N \text{ times}$$

$$\mu = \frac{N}{8}, \sigma^2 = \frac{7N}{64} \quad k = \frac{\frac{N}{4} - \frac{N}{8}}{\sqrt{\frac{7N}{64}}} = \frac{N}{\sqrt{7N}} = \sqrt{\frac{N}{7}}$$

$$p(X \geq \frac{N}{4}) \leq \frac{7}{N}$$

3) plot

Question 2: Markov ChainQuestion 2(a)

1) yes it's possible, for example let X be a random variable where $x \in \{0, 1\}$ with $P(X=0) = P(X=1) = \frac{1}{2}$ then let $Y = X$ and $Z = X$. Hence $I(X; Y) = I(X; Z) = H(X)$ since Y and Z are completely determined by X

2) yes. let X be random v. that $P(X=0) = P(X=1) = \frac{1}{2}$ and $Y = X$. Z is a variable that independent to Y . Hence, $I(X; Y) = H(X) = H(Y)$ and $I(X; Z) = 0$, then $I(X; Y) \geq I(X; Z)$



3) according to chain rule for mutual information :

$$I(X; Y, Z) = I(X; Y) + I(X; Z|Y) = I(X; Z) + I(X; Y|Z)$$

$$\text{Hence } I(X; Y) + I(X; Z|Y) = I(X; Z) + I(X; Y|Z)$$

according to markov chain assumption where $I(X; Z|Y) = 0$

$$\text{Proof: } I(X; Z|Y) = H(Z|Y) - H(Z)$$

$$= E \left[\log_2 \frac{p(X, Z|Y)}{p(X, Y)p(Z|Y)} \right] = E[\log_2 1] = 0$$

$$\text{Then } I(X; Y) = I(X; Z) + I(X; Y|Z) \text{ Hence}$$

$$I(X; Y) \geq I(X; Z)$$

4) according to the chain rule on previous question, we get

$$I(X; Y) + I(X; Z|Y) = I(X; Z) + I(X; Y|Z)$$

$$I(X; Y|Z) = I(X; Y) - I(X; Z) \text{ since } I(X; Z) \text{ must be equal or more than } 0 \text{ Hence } I(X; Y) \geq I(X; Z)$$

this because when we know Z , the uncertainty can only decrease or stay the same. In other words, the dependence between X and Y can't be increased by the observation of a downstream variable.

Question 2(b)

1) we know that $I(X; Y, Z) = H(X) - H(X|Y, Z)$ and $I(X; T) = H(X) - H(X|T)$ since $p(t|y, z, x) = p(t|y, z)$, according Markov, this means $H(X|Y, Z) \leq H(X|T)$, then we substitute and get $I(X; Y, Z) = H(X) - H(X|Y, Z) \geq H(X) - H(X|T) = I(X; T)$

2) In order for $I(X; Y, Z) = I(X; T)$ knowing T must give same amount of information with Y, Z or $H(X|Y, Z) = H(X|T)$. Hence, T must be a function of Y, Z that $p(t|y, z) = p(t|y, z, x)$

Question 2 (c)

1) By definition of mutual information, we get

$$I(X_1; X_2, \dots, X_n) = H(X_1) - H(X_1 | X_2, \dots, X_n)$$

$$\text{where } H(X_2 | X_2, \dots, X_n) = H(X_2 | X_2, \dots, X_{n-1}) +$$

$$H(X_2 | X_2, \dots, X_{n-2}, X_n) + \dots + H(X_2 | X_2, X_n) + H(X_2 | X_n)$$

$$\text{and we plug } H \text{ back, we get: } I(X_1; X_2, \dots, X_n) = H(X_1)$$

$$- H(X_1 | X_2, \dots, X_{n-1}) - H(X_1 | X_2, \dots, X_{n-2}, X_n) - \dots -$$

$$H(X_1 | X_2, X_n) - H(X_1 | X_n)$$

2) we can simplify expression above and get

$$I(X_1; X_2, \dots, X_n) = H(X_1) - H(X_1 | X_n)$$

Question 3: AEP

a) $H(x) = 0.8 \times \log_2\left(\frac{1}{0.8}\right) + 0.2 \times \log_2\left(\frac{1}{0.2}\right) = 0.7219 \text{ bits}$

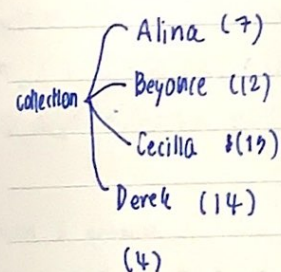
b) since there are two outcome and $N \in \{0, 1, 2, \dots, N\}$ the size of $A_X^N = 2^N$

c) since $|A_X| = 16$, then $H_0(X^4) = \log_2 |A_X| = 4$

d) $H(X^N) = N \sum_x p(x) \log p(x) = N H(x)$

e) As we increase N from small to large, the slope of the line become more flat meaning that it become less sensitive to change in error. When N get large, such sequence occupy most

Question 5: AEP of the probability mass, and are equally likely.



1) a. $\log_2(4) = 2 \text{ bits}$, $\{00, 01, 10, 11\}$

b. $\lceil \log_2(7) \rceil = 3 \text{ bits}$, $\{000, 001, 010, 011, 100, 101, 110\}$

c. $\lceil \log_2(48) \rceil = 6 \text{ bits}$

2) ~~unrelated~~ $H_0(x) = \lceil \log_2(4) \rceil = 2 \text{ bits}$

3) Ensemble $A = (x, A_x, P_x)$ where $A_x = \{ \text{Alina, Beyonce, Cecilia, Derek} \}$ and $P_x = \left(\frac{7}{48}, \frac{12}{48}, \frac{13}{48}, \frac{14}{48} \right)$ such that the first element of P_x is Alina and so on.

b. $\log_2(256) = 8 \text{ bits}$

c. smallest is 0

d. largest is 0.3958

4) a. $H(x) = - \left(\frac{7}{48} \log_2 \frac{7}{48} + \frac{12}{48} \log_2 \frac{12}{48} + \frac{13}{48} \log_2 \frac{13}{48} + \frac{14}{48} \log_2 \frac{14}{48} \right)$

$$= 1.95$$

b. $|T_N^0| \approx 2^{100(1.95+0.1)} \approx 2^{100 \times 2.05}$

c. no, it's not possible because the minimum of bit of uniform code with zero loss is 1.95 which is the minimum for uniform code and 1.5 bits is less than 1.95

```

# Function for toss coin
library(ggplot2)
coinTossOutcomes <- function(N, p) {
  # Define the sample space and probabilities
  sample_space <- c("H", "T")
  probabilities <- c(p, 1 - p)

  # Generate all possible outcomes for N coin tosses
  all_outcomes <- expand.grid(replicate(N, sample_space, simplify = FALSE))

  # Initialize a vector to store the probabilities of each outcome
  outcome_probabilities <- numeric(nrow(all_outcomes))

  # Calculate the probability of each outcome
  for (i in 1:nrow(all_outcomes)) {
    outcome <- as.character(unlist(all_outcomes[i, ]))
    prob <- prod(ifelse(outcome == "H", probabilities[1], probabilities[2]))
    outcome_probabilities[i] <- prob
  }

  # Add the probabilities to the data frame
  all_outcomes$Probability <- outcome_probabilities

  # Combine the toss outcomes into a single string for easier viewing
  all_outcomes$X <- do.call(paste0, all_outcomes[, 1:N])

  # Create the final table
  final_table <- all_outcomes[, c("X", "Probability")]

  # Normalize table
  return(final_table)
}

# 1. N = 5
N <- 5
p <- 0.8
result <- coinTossOutcomes(N, p)
x_axis = seq(0,1,0.001) ;print(x_axis)
y = c()
for (x in x_axis) {
  events = sum(result$Probability > x) # >= (1-x)
  y_axis = log2(events)
}

```

```

y = c(y,y_axis)

}
final_result = data.frame(y,x_axis)
final_result$normal = final_result$y/N
a = final_result$normal

# 2. N = 10
rm(final_table)
rm(result)
N <- 10
p <- 0.8
result <- coinTossOutcomes(N, p)
x_axis = seq(0,1,0.001) ;print(x_axis)
y = c()
for (x in x_axis) {
  events = sum(result$Probability > x) # >= (1-x)
  y_axis = log2(events)
  y = c(y,y_axis)
}
final_result = data.frame(y,x_axis)
final_result$normal = final_result$y/N
b = final_result$normal


# 3. N = 15
rm(final_table)
rm(result)
N <- 15
p <- 0.8
result <- coinTossOutcomes(N, p)
x_axis = seq(0,1,0.001) ;print(x_axis)
y = c()
for (x in x_axis) {
  events = sum(result$Probability > x) # >= (1-x)
  y_axis = log2(events)
  y = c(y,y_axis)
}
final_result = data.frame(y,x_axis)
final_result$normal = final_result$y/N
c = final_result$normal

```

```

# 4. N = 20
rm(final_table)
rm(result)
N <- 20
p <- 0.8
result <- coinTossOutcomes(N, p)
x_axis = seq(0,1,0.001) ;print(x_axis)
y = c()
for (x in x_axis) {
  events = sum(result$Probability > x) # >= (1-x)
  y_axis = log2(events)
  y = c(y,y_axis)
}
final_result = data.frame(y,x_axis)
final_result$normal = final_result$y/N
d = final_result$normal

# 3. Combine all
df <- data.frame(x_axis = final_result$x_axis, "5" = a, "10" = b, "15" = c, "20" = d)

# Reshape the data to long format
df_long <- tidyr::gather(df, variable, value, -x_axis)

# Generate the plot
ggplot(df_long, aes(x = x_axis, y = value, color = variable)) +
  geom_line() +
  labs(title = "Q3.e) Essential Bits given delta",
       x = "delta",
       y = "Normalised Essential Bit Content") +
  scale_color_discrete(name = "N")

```