

Open system Interconnection (OSI)

PRACTICE PROBLEMS BASED ON CSMA / CD AND BACK OFF ALGORITHM

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Problem-01

Q1 After the k th consecutive collision, each colliding station waits for a random time chosen from the interval-

- ① $(0 \text{ to } 2^k) \times \text{RTT}$
- ② $(0 \text{ to } 2^k - 1) \times \text{RTT}$
- ③ $(0 \text{ to } 2^k - 1) \times \text{Maximum Propagation delay}$
- ④ $(0 \text{ to } 2^k - 1) \times \text{Maximum Propagation delay}$

Q2 In a CSMA / CD network running at 1 Gbps over 1 km cable with no repeaters, the signal speed in the cable is 200000 km/sec. What is minimum frame size?

● **Solution-** Given-

Bandwidth = 1 Gbps, Distance = 1 km, Speed = 200000 km/sec



- **Calculating Propagation Delay:-**

- Propagation delay (T_p) = Distance / Propagation speed
= 1 km / (200000 km/sec)
= 0.5×10^{-5} sec

- **Calculating Minimum Frame Size:-**

- Minimum frame size = 2 x Propagation delay x Bandwidth
= $2 \times 5 \times 10^{-6} \text{ sec} \times 10^9 \text{ bits per sec}$

Q3 A 2 km long broadcast LAN has 10^7 bps bandwidth and uses CSMA / CD. The signal travels along the wire at 2×10^8 m/sec. What is the minimum packet size that can be used on this network?

- 1 50 B
- 2 100 B
- 3 200 B
- 4 None of the above



- Solution- Given-

Distance = 2 km, Bandwidth = 10^7 bps Speed = 2×10^8 m/sec

- **Calculating Propagation Delay**

- Propagation delay (T_p) = Distance / Propagation speed
- = 2×10^3 m / (2×10^8 m/sec)

- **Calculating Minimum Frame Size**

- Minimum frame size = 2 x Propagation delay x Bandwidth
= 2×10^{-5} sec x 10^7 bits per sec



Q4 A and B are the only two stations on Ethernet. Each has a steady queue of frames to send. Both A and B attempts to transmit a frame, collide and A wins first back off race. At the end of this successful transmission by A, both A and B attempt to transmit and collide. The probability that A wins the second back off race is.

- ① 0.5
- ② 0.625
- ③ 0.75
- ④ 1.0

● **Solution-According to question, we have-**

● 1st Transmission Attempt-

- Both the stations A and B attempts to transmit a frame.
- A collision occurs.
- Back Off Algorithm runs.
- Station A wins and successfully transmits its 1st data packet.



- 2nd Transmission Attempt-
 - Station A attempts to transmit its 2nd data packet.
 - Station B attempts to retransmit its 1st data packet.
 - A collision occurs.
- Now,
- We have been asked the probability of station A to transmit its 2nd data packet successfully after 2nd collision.
- After the 2nd collision occurs
- **At Station A-**
 - 2nd data packet of station A undergoes collision for the 1st time.
 - So, collision number for the 2nd data packet of station A = 1.
 - Now, station A randomly chooses a number from the range $[0, 2^1 - 1] = [0, 1]$.
 - Then, station A waits for back off time and then attempts to retransmit its data packet.



- **At Station B-**

- 1st data packet of station B undergoes collision for the 2nd time.
- So, collision number for the 1st data packet of station B = 2.
- Now, station B randomly chooses a number from the range $[0, 2^2 - 1] = [0, 3]$.
- Then, station B waits for back off time and then attempts to retransmit its data packet.
- Following 8 cases are possible-



Station A	Station B	Remark
0	0	Collision
0	1	A wins
0	2	A wins
0	3	A wins
1	0	B wins
1	1	Collision
1	2	A wins
1	3	A wins

- From here, Probability of A winning the 2nd back off race = $5 / 8 = 0.625$.



Q5 Suppose nodes A and B are on same 10 Mbps Ethernet segment and the propagation delay between two nodes is 225 bit times. Suppose A and B send frames at $t=0$, the frames collide then at what time, they finish transmitting a jam signal. Assume a 48 bit jam signal.

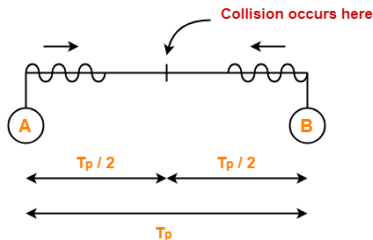
● **Solution-**

$$\begin{aligned}\text{Propagation delay } (T_p) &= 225 \text{ bit times} \\ &= 225 \text{ bit} / 10 \text{ Mbps} \\ &= 22.5 \times 10^{-6} \text{ sec} \\ &= 22.5 \mu\text{sec}\end{aligned}$$

● **At $t = 0$,**

- Nodes A and B start transmitting their frame.
- Since both the stations start simultaneously, so collision occurs at the mid way.
- Time after which collision occurs = Half of propagation delay.
- So, time after which collision occurs = $22.5 \text{ sec} / 2 = 11.25 \mu\text{sec}$.





- **At $t = 11.25$ sec,**

- After collision occurs at $t = 11.25$ sec, collided signals start travelling back.
- Collided signals reach the respective nodes after time = Half of propagation delay
- Collided signals reach the respective nodes after time = $22.5 \text{ sec} / 2 = 11.25 \text{ sec}$.
- Thus, at $t = 22.5$ sec, collided signals reach the respective nodes.



- **At $t = 22.5$ sec,**
 - As soon as nodes discover the collision, they immediately release the jam signal.
 - Time taken to finish transmitting the jam signal = 48 bit time = 48 bits / 10 Mbps = 4.8 sec.
- Thus,
 - Time at which the jam signal is completely transmitted
= 22.5 sec + 4.8 sec
= 27.3 sec or 273 bit times



Q6 Suppose nodes A and B are attached to opposite ends of the cable with propagation delay of 12.5 ms. Both nodes attempt to transmit at $t=0$. Frames collide and after first collision, A draws $k=0$ and B draws $k=1$ in the exponential back off protocol. Ignore the jam signal. At what time (in seconds), is A's packet completely delivered at B if bandwidth of the link is 10 Mbps and packet size is 1000 bits.

● **Solution-** Given-

Propagation delay = 12.5 ms, Bandwidth = 10 Mbps, Packet size = 1000 bits

● **Time At Which Collision Occurs-**

- Collision occurs at the mid way after time = Half of Propagation delay
= $12.5 \text{ ms} / 2$
= 6.25 ms

Thus, collision occurs at time $t = 6.25 \text{ ms}$.



● Time At Which Collision is Discovered-

- Collision is discovered in the time it takes the collided signals to reach the nodes = Half of Propagation delay
= $12.5 \text{ ms} / 2$
= 6.25 ms

Thus, collision is discovered at time $t = 6.25 \text{ ms} + 6.25 \text{ ms} = 12.5 \text{ ms}$.

● Scene After Collision-

- After the collision is discovered,
- Both the nodes wait for some random back off time.
- A chooses $k=0$ and then waits for back off time = $0 \times 25 \text{ ms} = 0 \text{ ms}$.
- B chooses $k=1$ and then waits for back off time = $1 \times 25 \text{ ms} = 25 \text{ ms}$.
- From here, A begins retransmission immediately while B waits for 25 ms.



• **Waiting Time For A-**

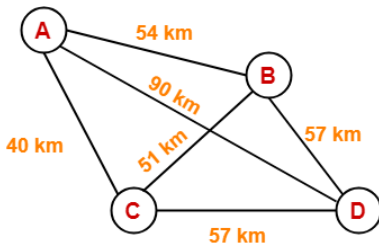
- After winning the back off race, node A gets the authority to retransmit immediately.
- But node A does not retransmit immediately.
- It waits for the channel to clear from the last bit aborted by it on discovering the collision.
- Time taken by the last bit to get off the channel = Propagation delay = 12.5 ms.
- So, node A waits for time = 12.5 ms and then starts the retransmission.
- Thus, node A starts the retransmission at time $t = 12.5 \text{ ms} + 12.5 \text{ ms} = 25 \text{ ms}$.



• Time Taken in Delivering Packet To Node B-

- Time taken to deliver the packet to node B = Transmission delay + Propagation delay
= $(1000 \text{ bits} / 10 \text{ Mbps}) + 12.5 \text{ ms}$
= $100 \text{ s} + 12.5 \text{ ms}$
= $0.1 \text{ ms} + 12.5 \text{ ms}$
= 12.6 ms
- Thus, At time $t = 25 \text{ ms} + 12.6 \text{ ms} = 37.6 \text{ ms}$, the packet is delivered to node B.

Q7 The network consists of 4 hosts distributed as shown below-



- Assume this network uses CSMA / CD and signal travels with a speed of 3×10^5 km/sec. If sender sends at 1 Mbps, what could be the minimum size of the packet?

- ① 600 bits
- ② 400 bits
- ③ 6000 bits
- ④ 1500 bits

- **Solution-**

- CSMA / CD is a Access Control Method.
 - It is used to provide the access to stations to a broadcast link.
 - In the given network, all the links are point to point.
 - So, there is actually no need of implementing CSMA / CD.
 - Stations can transmit whenever they want to transmit.
- In CSMA / CD, The condition to detect collision is-
Packet size $\geq 2 \times (\text{distance} / \text{speed}) \times \text{Bandwidth}$



- To solve the question,
 - We assume that a packet of same length has to be used in the entire network.
 - To get the minimum length of the packet, what distance we should choose?
 - To get the minimum length of the packet, we should choose the minimum distance.
 - But, then collision would be detected only in the links having distance less than or equal to that minimum distance.
 - For the links, having distance greater than the minimum distance, collision would not be detected.
 - So, we choose the maximum distance so that collision can be detected in all the links of the network.
- So, we use the values-
 Distance = 90 km, Speed = 3×10^5 km/sec, Bandwidth = 1 Mbps
 - Substituting these values, we get-
 Minimum size of data packet = $2 \times (90 \text{ km} / 3 \times 10^5 \text{ km per sec}) \times 1 \text{ Mbps}$
 = 600 bits



Thank You

