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School of Science
Master's Programme in Computer, Communication and Information Sciences

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Inferring Voting Networks in Online Elections

Master's Thesis
Espoo, March 14, 2020

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Supervisor: Professor Aristides Gionis
Advisor: Blank M.Sc. (Tech.)

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Chapter 1

Introduction

In recent years, researchers have become increasingly interested in understanding the behaviour of voters in social networks. Knowledge of the factors that motivate voters is of great importance in selecting successful policies or candidates. This is a classic problem and has been studied in the fields of game theory and political science [13, 19, 21]. More recently, there is a focus on using information from the network of voters to model their behaviour. This provides an insight into the interactions and effect of influence on voters in a community. For example, voting for bills in the United States Congress [12] or electing administrators in Wikipedia [2, 11, 15].

Votes can be represented as a *signed* network with positive or negative links. Finding groups using clustering techniques [1, 4, 18] and predicting signed links [5, 16, 17] in these networks is well researched. These approaches provide an ability to understand the group dynamics at play and predict votes and in such a network. However, they do not consider the iterative and chronological nature of the voting that takes place in these networks. In cases where research does focus on voter models, they rely on external features to build machine learning models that are task-specific and static [11, 12].

In this thesis, we propose a model that creates a local signed network consisting of the neighbours of the current voter and the preceding votes. It will then predict the vote, which when added to the network will comply the most with concepts of balance and status in signed networks. After all the votes are cast in a session, the model can be easily updated to improve quality and is, therefore, iterative and dynamic. The model is also flexible and can incorporate external features to build the local signed network of a voter. The results for Wikipedia administrator elections shows that our model outperforms machine learning based models and traditional signed link prediction solutions.

1.1 Thesis Outline

The rest of the thesis is organized in the following manner. We discuss the background relating to signed graphs, hierarchy in directed networks in Section 2. In Section 3 we describe the vote prediction problem and approaches to solving it. Section 4 provides a comprehensive view of Wikipedia and the election process for administrators. In Section 5 we explain the datasets used, construction of the model and evaluation criteria. We report our findings in Section 6 and discuss their implications. Finally, we conclude the thesis and present future work in Section 7.

Chapter 2

Graph Theory

In this chapter we will provide the fundamentals of the graph theory concepts required to understand the rest of the thesis. In Section 2.1 we cover the basic definitions, terminologies used to describe different types of graphs. Then we define a signed graph and discuss its unique properties in Section 2.2. We outline the theory of balance in signed networks and methods to measure it in Section 2.2.1. Next, we discuss the theory of status and illustrate the differences to balance in Section 2.2.2. Lastly, we explain techniques of finding hierarchies in directed networks and the concept of agony in Section 2.3.

2.1 Preliminaries

In this section we define the various types of graphs and their basic properties. The notation and terminologies used closely follows those used in Diestel [6]. Graphs are structures that describe relationships between entities. These entities are called *vertices* and entities related to one another are joined by edges. The terms graph, vertices and edges can be used interchangeably with *network*, *nodes* and *links* respectively.

Graphs can be classified broadly into two types based on whether the edges possess a direction or not. We now go on to define them in detail.

2.1.1 Undirected Graphs

An undirected graph is pair $G = (V, E)$, where V is the set of vertices and E is the set $E \subseteq \{(u, v) \mid u, v \in V\}$ of unordered pairs of vertices called edges. In this thesis we will deal with only *simple graphs*, i.e. no self loops $(u, v) \in V \times V$, $u \neq v$ and there is at most one edge between vertices u and v .

The number of the vertices in a graph is called the *order* of the graphs is denoted by $n = |G|$ and the *size* of a graph is the number of edges denoted by $m = \|G\|$ or $m = |E|$. A vertex u is *adjacent* to v if they are the end points of an edge, $(u, v) \in E$. All the vertices adjacent to vertex v is called the *neighbourhood* of v and is denoted by $N(v)$. The *degree* of a vertex v is the number of nodes adjacent to that vertex and is denoted by $d(v) = |N(v)|$.

The edges of an undirected graph can also have an associated value. This value can indicate the distance or similarity between a pair of vertices. These values are called *weights* and the corresponding graph is called a *weighted undirected graph*. Therefore, a weighted graph is defined as a triple $G = (V, E, w)$, where $w : E \rightarrow \mathbb{R}^+$ is a function that maps an edge e to a positive real weight $w(e)$. Now an *unweighed graph* is simply a weighted graph where the function w is defined as: if $e \in E$ then $w(e) = 1$ else $w(e) = 0$. The degree of a vertex v in a weighed graph is the sum of the weights to all the neighbours of v and is defined as $d(v) = \sum_{u \in N(v)} w((u, v))$. An example of a weighted undirected graph is shown in Figure 2.1.

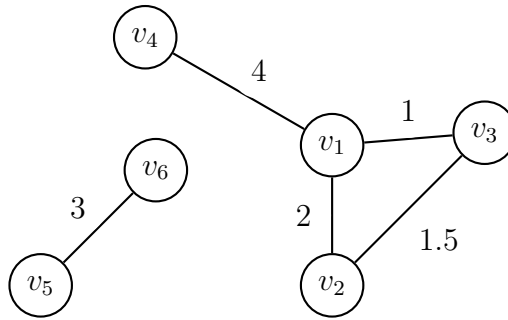


Figure 2.1: An example of a weighted undirected graph

2.1.2 Directed Graphs

The main distinction regarding a *directed graph* (or *digraph*) is that the edges are ordered pairs, i.e. $(u, v) \neq (v, u)$. Therefore, a directed graph has a similar definition: a pair $G = (V, E)$, where V is the set of vertices and E is the set of *ordered* pairs of vertices. Now given an edge $e = (u, v)$ we can define a source function $\text{src} : E \rightarrow V$ such that $\text{src}(e) = u$ and a destination function $\text{dest} : E \rightarrow V$ where $\text{dest}(e) = v$. These functions classify the vertices in an edge e as either the source or the destination. In this thesis, we deal only with *simple directed graphs*, i.e. no self-loops and there can be at most one edge from u to v .

As the edges now have an inherent direction we can define the *successors* and *predecessors* of a node v . A vertex u is called the *successor* of a node v if there exists a directed edge from v to u , therefore the set of successors for a vertex v can be defined as $S(v) = \{u \mid (v, u) \in E\}$. A *predecessor* of a node v is a vertex u such that there exists a directed edge from u to v , the set of predecessors for a vertex v can be defined as $P(v) = \{u \mid (u, v) \in E\}$. We now a vertex u that is either a successor or a predecessor of a vertex v can be called a neighbour of the vertex v . Therefore, we define the *neighbourhood* of a vertex v as the set of vertices in the union of successors and predecessor, i.e $N(v) = S(v) \cup P(v)$. This definition is also compatible with undirected graphs because if $(u, v) \in E$ then $(v, u) \in E$.

Directed graphs can also have values associated with each directed edge called a *weight*. A *weighted directed graph* can be defined as a triple $G = (V, E, w)$, where the weight function $w : E \rightarrow \mathbb{R}^+$ that maps each edge e to a weight $w(e)$. The indegree of a vertex v is defined as the sum of the edge weights from the predecessors of v and is denoted as $d_{\text{in}}(v) = \sum_{u \in P(v)} w((u, v))$. Similarly, the outdegree of a vertex v is defined as the sum of the edge weights to the successors of v and is denoted by $d_{\text{out}}(v) = \sum_{u \in S(v)} w((v, u))$. Figure 2.2 shows an example of a weighted directed graph.

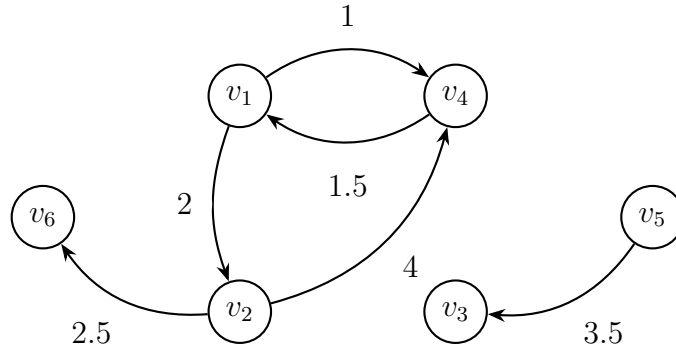


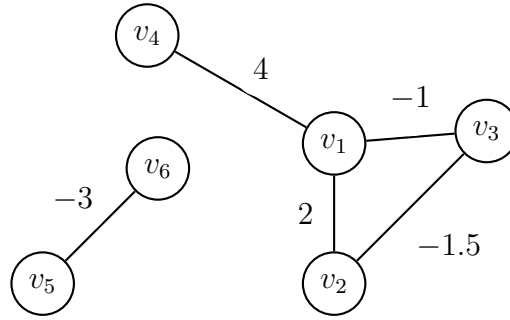
Figure 2.2: An example of a weighted directed graph

2.2 Signed Graphs

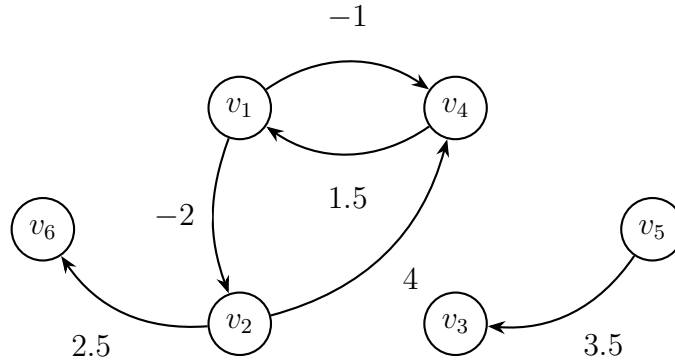
The simple weighted graphs we have defined so far only have non-negative edge weights that can represent similarity or closeness. In the 1950s social psychologists found it desirable to express liking, disliking or indifference in social interactions. This was formalized by Harary [8] using graphs with weights $(-1, 0, 1)$. These graphs are therefore called *signed graphs*, where a

negative edge weight can denote dissimilarity between a pair of vertices. In this thesis we will use notations and terms similar to Gallier [7], Kunegis et al. [14], Hou [10] and Zaslavsky [20].

A signed graph is a triple $G = (V, E, w)$, where V is the set of vertices, E is the set of pair of vertices and the weight function $w : E \rightarrow \mathbb{R}$. The weight function now takes an edge e and maps it to a signed weight $w(e)$. We can partition the edges into positive and negative edges, $E = E^+ \cup E^-$, where $E^+ = \{e \mid w(e) > 0\}$ and $E^- = \{e \mid w(e) < 0\}$. Similar to Zaslavsky [20], we consider undirected signed graphs where and directed signed graphs as distinct and separate entities. We can see some examples of signed graphs in Figure 2.3.



(a) A undirected signed graph



(b) A directed signed graph

Figure 2.3: Examples of Signed Graphs

We can now proceed to define a few more terms. The degree of a vertex v is now the sum of the absolute edge weight of its neighbours is called the *signed degree* and is defined as

$$\bar{d}(v) = \sum_{u \in N(v)} |w((u, v))|$$

This can also be extended to *signed indegree* and *signed outdegree* denoted by $\overline{d}_{\text{in}}(v)$ and $\overline{d}_{\text{out}}(v)$ and defined as

$$\overline{d}_{\text{in}}(v) = \sum_{u \in P(v)} |w((u, v))|$$

$$\overline{d}_{\text{out}}(v) = \sum_{u \in S(v)} |w((v, u))|$$

We can create a $n \times n$ square weight matrix W , where each entry w_{ij} is defined as

$$w_{ij} = \begin{cases} w((v_i, v_j)) & \text{if } (v_i, v_j) \in E \\ 0 & \text{if } (v_i, v_j) \notin E \end{cases}$$

The signed degree matrix \overline{D} is a diagonal matrix consisting of the signed vertex degrees, $\overline{D} = \text{diag}(\overline{d}(v_1), \dots, \overline{d}(v_n))$. We can now define the *signed Laplacian*, \overline{L} as

$$\overline{L} = \overline{D} - W$$

The signed Laplacian along with results from spectral analysis of signed graphs [10, 14], will be useful for balance theory of signed graphs.

2.2.1 Balance Theory

Heider [9] proposed in 1940's that when there are either positive relations (friendship, love, support) and negative relations (enmity, hate, oppose) in a group, the group tends towards *balance*. Balance was defined as there being all positive relations or one positive and two negative relations for a group of 3 people. Harary and Cartwright [3] generalized this notion of balance by using signed graphs. This can be seen in Figure 2.4 where solid edges are positive and dashed edges are negative. Social psychology also showed that balanced triads T_1 and T_2 mirror aphorisms such as "the friend of my friend is also a friend" and "the enemy of my enemy is a friend" respectively.

The concept of balance can be generalized to state that an undirected signed graph $G = (V, E, w)$ is balanced iff every cycle in G has an even number of negative edges. This leads to result from Harary [8] that states that if a graph G is balanced, then there is a partition of the vertices $V = V_1 \cup V_2$ such that edges within the vertices of each set is positive and edges between the sets are negative. This means that we can have a bipartite style graph where edges between the two sets of vertices are negative and every cycle has an even number of negative edges.

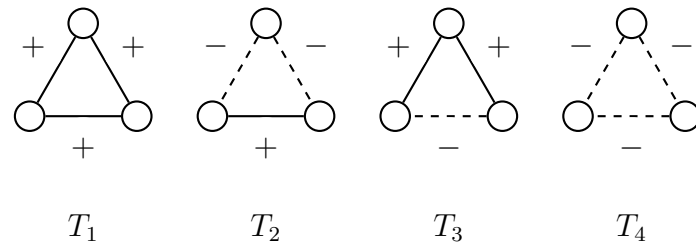


Figure 2.4: Triads in undirected signed graphs. T_1 and T_2 are *balanced* triads as they have even number of negative edges. T_3 and T_4 are *unbalanced* as they have odd number of negative edges.

2.2.2 Status Theory

- Describe the nature of the directed setting
- Illustrate the differences to Balance theory
- Mention existing ways to measure violation to status in a network

2.3 Hierarchy in directed networks

- Discuss the hierarchy in DAGs
- Explain concept of Agony
- Provide existing algorithms to find the ranking of nodes.

Chapter 3

Vote Prediction

In this section we first provide the motivation of choosing independent vote prediction as our target and the differences from predicting the result of an election. Next we describe the available techniques and methods to predict individual votes or signed edges in a network and how it relates to the problem at hand. We then provide two novel methods of using user information long with past election results to predict votes.

3.1 Result versus Vote Prediction

- Discuss existing result predictions schemes
- Discuss the limitations in understanding voting dynamics through just predicting results
- Describe the process as an information cascade, discuss the potential Game Theory settings
- Show the two parts of the problem from an information cascading perspective
 - Who is going to vote next
 - How they are going to vote
- Discuss the assumptions in usual Independent Cascade (IC) models
- Explain the difficulty of both aspects in the domain of an election
- Motivate the selection of the problem as an **Independent Vote Prediction**

3.2 Signed Edge Prediction

- Discuss the existing edge predictions work
- Directly using signed triads as features
- Using triads along with network features
- Using user information and interaction data for predicting votes and/or elections
- The main drawbacks in these methods when considering an election setting

3.3 Linear Combination of Graphs

- Describe the linear combination of graphs derived from user and election data
- Explain topic similarity, follows network, interaction networks and other features
- How it can also incorporate signed features as additional features in prediction

3.4 Local Signed Network

- Explain the concept of the local signed network for a particular user
- Motivate the definition with respect to elections and influence
- Describe how to use balance and status theory to predict the vote
- Clarify the differences to signed edge prediction efforts
- Mention Agony as a way to measure status compliance here?

Chapter 4

Wikipedia

In this section we provide an overview of how Wikipedia is structured, the hierarchy that exists withing editors. We then explain the election process of getting administrator rights in Wikipedia.

4.1 Structure and hierarchy in Wikipedia

4.2 Elections in Wikipedia

- Explain Editors and Administrators in Wikipedia
- Describe the Request for Administrator(RfA) process
- Discuss general trends and patters
- Mention research interest and possible current works?

Chapter 5

Experiments

In this section we first describe the datasets that will be used in building our vote prediction models. Then we discuss the various linear and graphical models that we consider and their implementations details. Lastly we define the metrics and other means of evaluating the models and the results.

5.1 Datasets

- Maybe a short description of existing SNAP datasets and their limitations
- The details of the *Wiki-RfA* data and the *User-Contribution* datasets

5.2 Graphs

- Discuss the process of extraction of the various graphs discussed in the previous sections
- **Agree Graphs and Follows Graph**, where we measure the degree to which one user agrees and follows another user in previous elections
- **Topic similarity** from the top 100 articles edited for each user and the pairwise Jaccard similarity
- **Talk and Interaction graphs**, measures communication between users on their respective user talk pages
- **Signed Graphs**, triad encoding and extracting the triad counts for each voter

5.3 Models

5.3.1 Linear Combination of Graphs

- Discuss the various linear models considered for Graph Combinations
 - Linear Regression
 - Support Vector Classifier
 - Extreme Gradient Boosting (XGBOOOOST)
- Discuss how each graph contributes features and the problem is a linear classification problem

5.3.1.1 Iterative Mode

- Discuss the motivation behind an iterative model versus a static prediction model
- Describe how balance is derived from the Agree Graph in a local signed network
- Discuss how the Agree graph is updated in terms of Balance
- Describe how status is derived from the Follows graph in a local signed network
- Discuss how the Follows graph is updated after every election
- Describe how to make the predictions
 - Deterministic : just decide based on eigen value or agony as support or oppose
 - Probabilistic : provide a probability for predicting a support vote

5.4 Evaluation

- Discuss the issues with the imbalance in the datasets
- Illustrate the issues with pure measures of accuracy
- Define Precision, Recall and Macro F1 score
- Discuss ROC AUC and Precision Recall curves for probability based predictions

Chapter 6

Results and Discussion

In this section we will present the results of the models and discuss their implications.

6.1 Linear Combination of Graphs

- Present results for each linear classifier
- Discuss the different splits of the dataset to check for robustness and chronological consistency
- Show the feature importances and discuss their relevance
- Compare the raw accuracy versus the macro f1 scores
- Highlight the difficulty of predicting negative votes

6.2 Local Signed Network

- Present the Iterative Balance model results
- Discuss quality of predictions using evaluation metrics
- Mention the difference between deterministic and probabilistic prediction accuracies
- Explain the Iterative Status model results
- Discuss the issues with local model of status and the potential reasons for lower score and quality

6.3 Comparison

- Compare results from signed edge prediction and Iterative signed models
- Discuss Static Linear combination predictions versus Iterative signed predictions
- Discuss the assumptions used in the models and limitations

Chapter 7

Conclusions and Future Work

- Explain the quality of results with the election perspective
- Future work is to extend this to other election settings and investigate generality of this approach
- Possible future work in congressional voting data
- Can also tackle the other problem in information cascade theory of how to predict who is most likely to vote next
- This can lead to a complete model of election dynamics and could incorporate elements of game theory and network inference

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