



Design & Analysis of Algorithm (Lab)

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B-33

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<https://github.com/ananya438/DAALAB> ANANYA-590013832

Implement Longest Common Subsequence Problem.

```
public class LCSDynamicProgramming {

    public static String findLCS(String text1, String text2) {

        int m = text1.length();
        int n = text2.length();

        int[][] c = new int[m + 1][n + 1];

        for (int i = 1; i <= m; i++) {
            for (int j = 1; j <= n; j++) {
                if (text1.charAt(i - 1) == text2.charAt(j - 1)) {
                    c[i][j] = c[i - 1][j - 1] + 1;
                } else {
                    c[i][j] = Math.max(c[i - 1][j], c[i][j - 1]);
                }
            }
        }

        System.out.println("Length of LCS: " + c[m][n]);

        return reconstructLCS(c, text1, text2, m, n);
    }

    private static String reconstructLCS(int[][] c, String text1, String text2, int i, int j) {
        if (i == 0 || j == 0) {
            return "";
        }

        if (text1.charAt(i - 1) == text2.charAt(j - 1)) {
            return reconstructLCS(c, text1, text2, i - 1, j - 1) + text1.charAt(i - 1);
        } else {
            if (c[i - 1][j] > c[i][j - 1]) {
                return reconstructLCS(c, text1, text2, i - 1, j);
            } else {
                return reconstructLCS(c, text1, text2, i, j - 1);
            }
        }
    }
}
```

```

public static void main(String[] args) {
    String S1 = "GTCGTCGGAATACTGTC";
    String S2 = "GGTCGTCGGAATTGAC";

    String lcsResult = findLCS(S1, S2);
    System.out.println("LCS: " + lcsResult);
}
}

```

Complexity Analysis for LCS (Dynamic Programming)

The complexity is mainly determined by the size of the 2D table created for the dynamic programming solution.

Time Complexity

The time complexity is $O(m \cdot n)$.

Dominant Factor: The main part of the algorithm is filling the 2D DP table.

Process: This task requires two nested loops: one loop for the m length of the first string and one for the n length of the second string.

Work per Step: Inside the loops, only simple, constant-time operations (like comparison and addition) are performed.

The final step to build the actual sequence takes only $O(m+n)$ time, which is much faster (dominated by) the $O(m \cdot n)$ table-filling time.

Space Complexity

The space complexity is $O(m \cdot n)$.

Primary Factor: The algorithm requires an auxiliary 2D array (the DP table, c).

Size: This table is of size $(m+1) \times (n+1)$ to store the lengths of the LCS for all possible prefixes.

Necessity: This space is essential because the algorithm must store the results of all smaller subproblems to build the final solution.

O/P:

```

● PS C:\Users\nannu\Desktop\JAVA DSA\DAA> cd "c:\Users\nannu\Desktop\JAVA DSA\DAA\" ; if ($?) {
    Length of LCS: 14
    LCS: GTCGTCGGAATTGC
○ PS C:\Users\nannu\Desktop\JAVA DSA\DAA>

```