



## **Design & Analysis of Algorithm (Lab)**

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[https://github.com/ananya438/DAALAB\\_ANANYA-590013832](https://github.com/ananya438/DAALAB_ANANYA-590013832)

## 0/1 knapsack (DP)

TUESDAY

\* 0/1 Knapsack (Dynamic Programming)

→ C = 8      P = 1 2 5 C

W = 2 3 4 5

We can either say  $x_i = 0/1$

$\sum p_i x_i$  should be maximum.

Time complexity  $O(2^n)$

Tabulation Method.

i	w	0	1	2	3	4	5	6	7	8
0	0	0	0	0	0	0	0	0	0	0
1	1	0	0	1	1	1	1	1	1	1
2	2	0	0	1	2	3	3	3	3	3
3	3	0	0	1	2	5	5	6	7	7
4	4	0	0	1	2	5	5	6	7	8
5	5	0	0	1	2	5	6	6	7	8

$V[i, w] = \max \{ V[i-1, w], V[i, w-w_i] + p_i \}$

$V[4, 1] = \max \{ V[3, 1], V[3, 1-3+6] \}$

2023



```
public class Knapsack {  
    static int knapsack(int W, int wt[], int val[], int n) {
```

O/P:

PROBLEMS    OUTPUT    DEBUG CONSOLE    TERMINAL    PORTS  
for (int i = 0; i <= n; i++) {  
PS C:\Users\nannu\Desktop\JAVA DSA\DAAs> cd "c:\Users\nannu\Desktop\JAVA DSA\DAAs\0\" ; if (\$?) { javac Knapsack.java  
● Maximum Profit =for (int w = 0; w <= W; w++) {  
○ PS C:\Users\nannu\Desktop\JAVA DSA\DAAs\0>  
if (i == 0 || w == 0)

## Complexity Analysis (Dynamic Programming)

Time Complexity     $dp[i][w] = \max(dp[i-1][w], dp[i-1][w - wt[i]] + val[i])$

$O(n \times W)$  → we fill a table of size  $n \times W$ .  
else

Space Complexity     $dp[i][w] = dp[i-1][w];$

$O(n \times W)$  → because of the 2D DP array.

```
return dp[n][W];  
}
```