



Design & Analysis of Algorithm (Lab)

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[https://github.com/ananya438/DAALAB ANANYA-590013832](https://github.com/ananya438/DAALAB_ANANYA-590013832)

Kruskal's Spanning TREE ALGO

1. Kruskal's Algorithm:

(a) Working Principle:

⇒ This is a greedy algorithm used to find Minimum Spanning Tree, of a connected, weighted graph.

STEP 1: Sort all edges of the graph in increasing order of weights.

STEP 2: Initialize an empty set for the MST.

STEP 3: Pick the smallest edge from sorted list & check if it forms a cycle in MST.

STEP 4: If no cycle is formed, include the edge in MST.

STEPS: Repeat steps 3 & 4 until the MST contains exactly $(V-1)$ edges, where V = no. of vertices.

The algorithm ensures that the total weight of MST is minimized.

EXAMPLE:

5. Kruskal's Algo (MST):

Find (MST) using it.

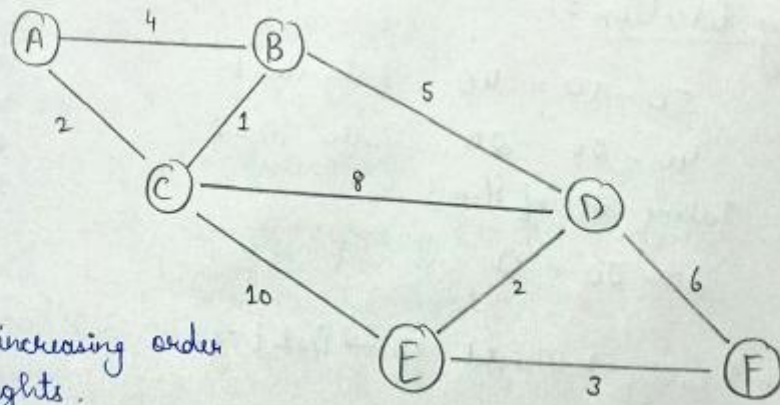
Vertices = $\{A, B, C, D, E, F\}$

Edges with weights

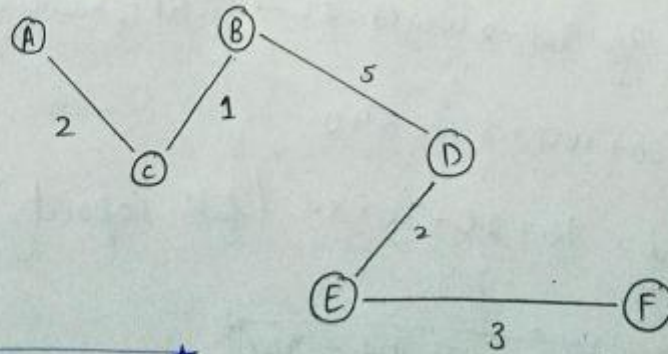
$(A, B, 4)$ $(A, C, 2)$, $(B, C, 1)$, $(B, D, 5)$, $(C, D, 8)$, $(C, E, 10)$,

$(D, E, 2)$, $(D, F, 6)$, $(E, F, 3)$.

* Original Weighted Graph:



- ① Sort all edges increasing order of their weights.
- ② Pick edges one by one from sorted list, & add them to MST only if they don't form cycle, until MST contains $(V-1)$ edges.



Total Min. Weight = 13 ★

```
import java.util.*;

class Edge implements Comparable<Edge> {
    int src, dest, weight;

    Edge(int s, int d, int w) { src = s; dest = d; weight = w; }

    public int compareTo(Edge o) { return this.weight - o.weight; } }

class Subset { int parent, rank; }

public class KruskalMST {
    int V, E; Edge[] edges;

    KruskalMST(int v, int e) { V = v; E = e; edges = new Edge[E]; }

    int find(Subset[] subsets, int i) {
        if (subsets[i].parent != i)
            subsets[i].parent = find(subsets, subsets[i].parent);
        return subsets[i].parent; }
}
```

```

void union(Subset[] subsets, int x, int y) {
    int xr = find(subsets, x), yr = find(subsets, y);

    if (subsets[xr].rank < subsets[yr].rank) subsets[xr].parent = yr;
    else if (subsets[xr].rank > subsets[yr].rank) subsets[yr].parent = xr;
    else { subsets[yr].parent = xr; subsets[xr].rank++; }
}

void kruskalMST() {
    Arrays.sort(edges);
    Edge[] result = new Edge[V - 1];
    Subset[] subsets = new Subset[V];

    for (int v = 0; v < V; v++) { subsets[v] = new Subset(); subsets[v].parent = v; }

    int e = 0, i = 0, total = 0;

    while (e < V - 1 && i < E) {
        Edge next = edges[i++];

        int x = find(subsets, next.src), y = find(subsets, next.dest);

        if (x != y) { result[e++] = next; union(subsets, x, y); } }

    for (i = 0; i < e; i++) {
        System.out.println(result[i].src + " - " + result[i].dest + " : " + result[i].weight);
        total += result[i].weight; }

    System.out.println("Total weight of MST = " + total) }

public static void main(String[] args) {
    int V = 6, E = 9;

    KruskalMST g = new KruskalMST(V, E);

    g.edges[0] = new Edge(0, 1, 4);
    g.edges[1] = new Edge(0, 2, 2);
    g.edges[2] = new Edge(1, 2, 1);
    g.edges[3] = new Edge(1, 3, 5);
    g.edges[4] = new Edge(2, 3, 8);
    g.edges[5] = new Edge(2, 4, 10);
    g.edges[6] = new Edge(3, 4, 2);
    g.edges[7] = new Edge(3, 5, 6);
    g.edges[8] = new Edge(4, 5, 3);

    g.kruskalMST();
}

```

}

O/P:

```
PS C:\Users\nannu\Desktop\JAVA DSA\JAVA\First lectures> & 'C:\Pr  
nannu\AppData\Roaming\Code\User\workspaceStorage\0af2579802541dcb  
1 - 2 : 1  
0 - 2 : 2  
3 - 4 : 2  
4 - 5 : 3  
1 - 3 : 5  
Total weight of MST = 13
```

Time Complexity

$O(E \log E)$, where E is the number of edges. This is because the algorithm is dominated by the time it takes to sort all the edges.

It can also be written as $O(E \log V)$, where V is the number of vertices, because in a connected graph, $\log E$ is on the same order as $\log V$.

Space Complexity

$O(V + E)$. This is because the algorithm needs to store all the edges and the Disjoint Set Union (Union-Find) data structure to keep track of the vertices.