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CODE/SLOT: CSE2002 / CII + DE1

REG NO: 20BCE10093

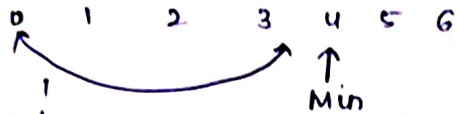
FACULTY : Ms. Meenakshi Choudhary

(a1) Selection sort

- * It uses the brute force approach.
- * It is a Inplace sorting algorithm: It takes constant amount of extra memory. Amount of extra memory is directly proportional to the size of the array.
- * PRINCIPLE: Sorts the elements in an array by finding minimum element in each pass from unsorted part and keeps it in the beginning. This divides the array into sorted and unsorted part.

RUN \Rightarrow

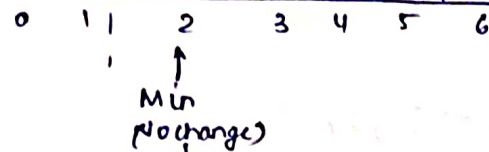
77	34	22	73	1	4	98
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1	34	22	73	77	4	98
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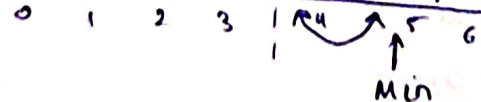
1	4	22	73	77	34	98
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1	4	22	73	77	34	98
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1	4	22	34	77	73	98
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1	4	22	34	73	77	98
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\rightarrow Sorted array

- * Select the minimum value and swap it from the first element of the unsorted array.

Pseudocode \Rightarrow

Selsort (A, n)

```
{
  for i  $\leftarrow$  0 to n-2
    {
      imin  $\leftarrow$  i
      for j  $\leftarrow$  i+1 to n-1
        {
          if (A[j] < A[imin])
            imin  $\leftarrow$  j
        }
      temp  $\leftarrow$  A[i]
      A[i]  $\leftarrow$  A[imin]
      A[imin]  $\leftarrow$  temp
    }
}
```

\rightarrow all are sorted by penultimate loop.

\rightarrow swap

PROGRAM

Void selsort (int A[], int n)

```
{
  for (int i = 0; i < n-1; i++)
  {
    int imin = i;
    for (int j = i+1; j < n; j++)
    {
      if (A[j] < A[imin])
        imin = j;
    }
    int temp = A[i];
    A[i] = A[imin];
    A[imin] = temp;
  }
}
```

}

int main ()

```
{
  int A[] = {7, 34, 22, 73, 1, 4, 98}
  selsort (A, 7)
  for (int i = 0; i < 7; i++)
  {
    cout << A[i] << " ";
  }
}
```

OUTPUT \Rightarrow 1 4 22 34 73 77 98 \Rightarrow (Time complexity = $O(n^2)$)

(b)

Insertion sort

- * Not the best sorting method, but better than selection and bubble sort.
- * Uses brute force approach.

* RUN \Rightarrow

array iterates from $arr[1]$ to $arr[n]$

Compare the current element with its predecessor.

If the element is smaller, compare it to elements before.

Move the bigger elements to make room for swapped elements

eg

77	34	22	73	1	4	98
----	----	----	----	---	---	----

\uparrow

value (stored) \rightarrow creates a hole (Any value can be chosen)

4	77	34	22	73	1	98
---	----	----	----	----	---	----

\uparrow

value (creates a hole, shift the greater elements towards right)

1	4	77	34	22	73	98
---	---	----	----	----	----	----

\uparrow

value

1	4	73	77	34	22	98
---	---	----	----	----	----	----

\uparrow

value

1	4	22	73	77	34	98
---	---	----	----	----	----	----

\uparrow

value

1	4	22	34	73	77	98
---	---	----	----	----	----	----

\rightarrow Sorted.

Pseudocode \Rightarrow Insert (A, n)

{ for $i \leftarrow 1$ to $n-1$

{ value $\leftarrow A[i]$

hole $\leftarrow i$

while (hole > 0 && $A[\text{hole}-1] > \text{value}$)

{ $A[\text{hole}] \leftarrow A[\text{hole}-1]$

hole $\leftarrow \text{hole} - 1$

}

$A[\text{hole}] \leftarrow \text{value}$

}

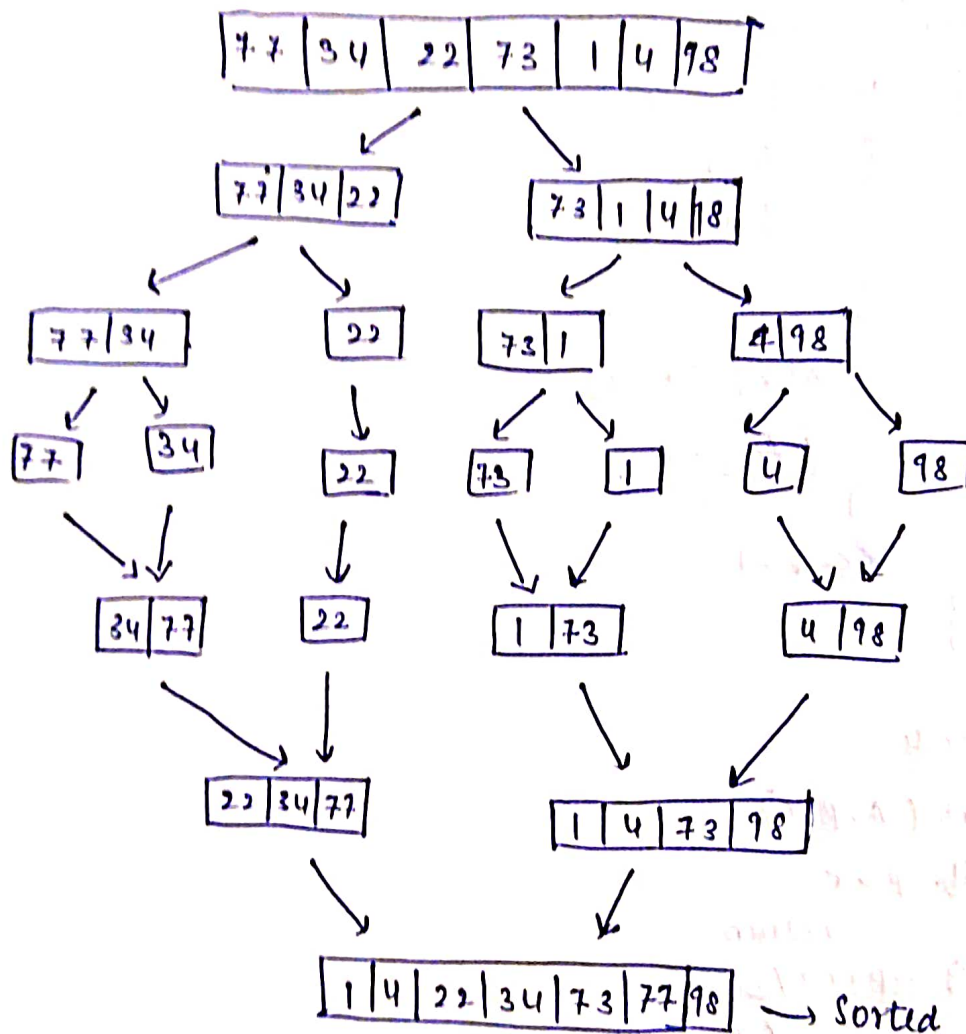
}

Time Complexity $\Rightarrow O(n^2)$

2(a) Merge Sort

- * Uses divide and conquer approach.
- * Works on the principle of breaking an array into smaller parts, and then arrange them in ascending order.

eg 77, 34, 22, 73, 1, 4, 98



- * Divide the array from the middle into two halves
- * Call mergesort in first half and second half.
- * Merge the two halves together.
- * Find middle point to divide the array $\rightarrow m = l + (r - l) / 2$
- call for merge sort in first half \rightarrow call mergesort (arr, l, m)
- call for merge sort in second half \rightarrow call mergesort (arr, m+1, r)
- merge \rightarrow call merge (arr, l, m, r)

pseudo-code

MergeSort (A, p, r) :

if $p > r$

return

$q = (p+r) / 2$

MergeSort (A, p, q)

MergeSort (A, q+1, r)

Merge (A, p, q, r)

- (v) Datastructure :
- * Way to store data in a computer system for efficient use.
 - * It is a way of data organisation so that the function of data structure remains independent of implementation.
 - * The functional definition of a data structure is known as ADT (Abstract data type)
 - * A well designed datastructure helps in pursuing critical operations with limited memory and space.
 - * Data structure has two fundamental objectives :
- (i) ways to store data
 - (ii) Types of operations performed on it

The way the data is stored affects the performance of program.

eg arrays, stack, queues, trees etc.

Data structures are classified as:

- (i) Simple data structure
- (ii) Compound data structure
- (iii) Linear data structure → (stacks, queues, lists)
- (iv) Non-linear data structure → (tree, graphs)

LINEAR DATA STRUCTURE	NON LINEAR DATA STRUCTURE
<ul style="list-style-type: none"> * Sequentially connected * Single run, all elements traversable. * Easy to implement * Doesn't use memory efficiently 	<ul style="list-style-type: none"> * Hierarchically connected * elements present at various levels and need multiple runs. * difficult to understand and implement * Very efficient memory use
<u>Example</u> Arrays, lists, stacks, queues.	<u>Example</u> Heaps, trees, graphs

(3)

LINEAR SEARCH

- * Starts searching from $arr[0]$ and compares each element till searched element is not found.
- * Doesn't need to be a sorted array.
- * Can be applied on any data structure - array, linked list, etc.
- * Based on sequential approach
- * Used for small data sets.
- * Can be implemented on both, single and multi-dimensional array.

TIME COMPLEXITIES →

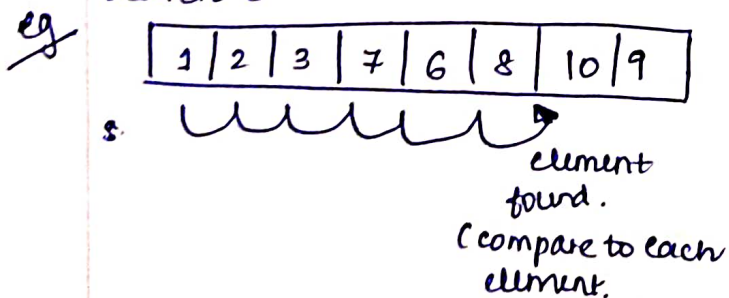
Worst case $\Rightarrow O(n)$

Best case $\Rightarrow O(1)$

Avg case $= O(n/2)$

- * Slow in speed.

Search 8



Codes \Rightarrow

```
void linearSearch(int obj[], int item, int start) {
    int start = 0;
    for (int i = start; i < obj.length; i++) {
        if (obj[i] == item)
            return i;
        return -1;
    }
}
```

BINARY SEARCH

It finds the position of the wanted element.

Sorted array is mandatory.

Applied only on data structures that have two way reversal.

- * Based on divide and conquer approach
- * Used on large data sets
- * Can be implemented only on a multidimensional array.

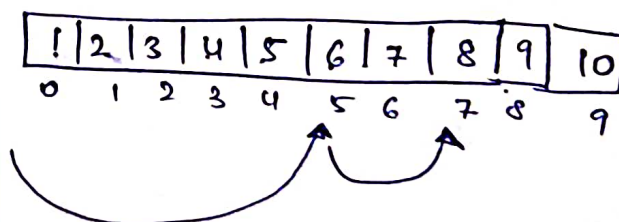
Worst case $\Rightarrow O(\log_2 n)$

Best case $\Rightarrow O(1)$

Average case $= O(\log n)$

- * Fast implementation

Search 8.



Codes \Rightarrow (assuming that the array is sorted.)

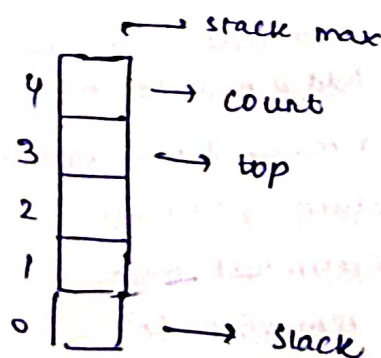
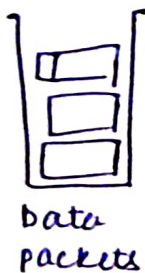
```
int bsearch(int a[], int s, int item) {
    int beg, mid, last;
    beg = 0;
    last = size - 1;
    while (beg <= last) {
        mid = (beg + last) / 2;
        if (item == a[mid])
            beg = mid + 1;
        else
            last = mid - 1;
    }
    return -1;
}
```


5) ADT (Abstract data type)

- * The abstract datatype is a special datatype, which is defined by a set of values and set of operations.
- * "Abstract" is used because we can use these datatypes for different operations. But the operation working is concealed from the user.
- * It is made up of primitive, existing datatypes, but operation logics are hidden.
- * So, basically, they are definitions of data and operations but do not have implementation details.

⇒ STACK ADT

- * Elements of same type arranged in a sequential order.
- * Operations:
 - Initialize() → initializing it to be empty
 - Push() → Insert an element
 - Pop() → Delete an element
 - isEmpty() → checks if stack is empty
 - isFull() → checks if stack is full
- * It uses LIFO data structure (Last in first out). The element which is placed last, is accessed first.
- * Data is not stored in each node, the pointer to the data is stored.
- * The program allocates memory for the data and address is passed to the stack ADT.



The head node and data node are encapsulated in ADT.
Head structure contains a pointer to the top and count the number of entries.


```

typedef struct node
{
    void *data_ptr;
    struct node *link;
}

```

Stack node;

```

typedef struct
{
    int count;
    Stacknode *top;
} Stack;

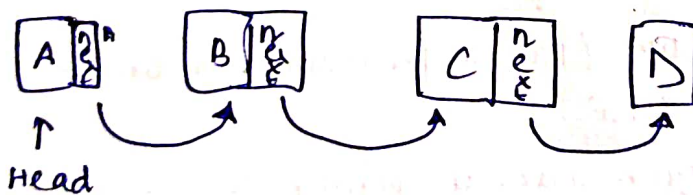
```

LINKED LIST

It is a type of ADT which has a collection of nodes, the nodes can be accessed sequentially..

No random access to a node.

Nodes are connected to the next node and with previous node, called links. When nodes are connected with only one (next) pointer, its called a singly linked list.



Main operations:

- Prepend → add a node at beginning
- append → Add a node at end
- pop() → remove a node from end
- head() → return first node
- tail() → return last node
- remove(Node) → remove node

def binary_search(arr, val, start, end):

if start == end:

if arr[start] > val:

return start;

else:

return start + 1;

if start > end

return start;

mid = (start + end) / 2

if arr[mid] < val:

return binary_search(arr, val, mid + 1, end)

elif (arr[mid] > val):

return binary_search(arr, val, start, mid - 1)

else:

return mid;

insertion.

sort(arr)

~~for i in range(1, len(arr))~~

for i = 1, i < length, i++

val = arr[i];

i = binary_search(arr, val, 0, i - 1)

arr = arr[0:i] + val + arr[i+1:]

return arr.

print → 84

=

from 81, 82, 56, 57, 30, 19, 84, 47, 42, 25.

THANK YOU MA'AM