

Ch 2. Basic Structures: Sets, Functions

Ch 9. Relations

Sets

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Acknowledgement

- [Rosen 19] Kenneth H. Rosen, for Discrete Mathematics & Its Applications (8th Edition), Lecture slides
- [Hunter 11] David J. Hunter, Essentials of Discrete Mathematics, 2nd Edition, Jones & Bartlett Publishers, 2011, Lecture Slides

Ch 2. Basic Structures: Sets, Functions

2.1 Sets

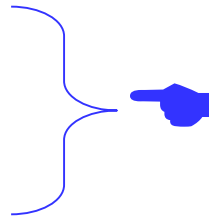
2.2 Set Operations

2.3 Functions

2.4 Sequences and Summations

2.5 Cardinality of Sets

2.6 Matrices



Sets

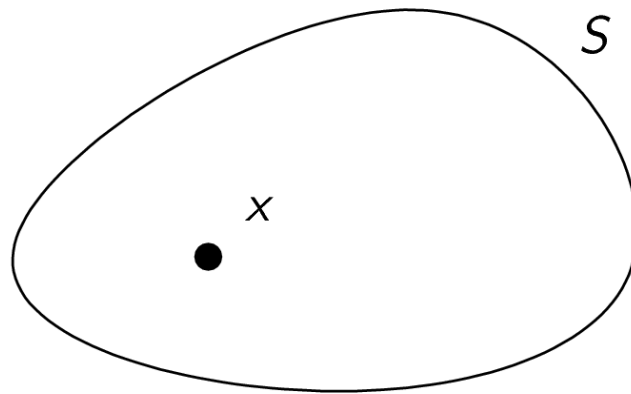
1. Set Membership
2. Subset & Equality Relationship
3. Constructing New Sets from Existing Sets
4. Set Identities
5. Discussion on Set Comprehension

1. Set Membership

$x \in S$ denotes that x is contained in S .
or x is an element of S
or x is a member of S

$x \notin S$ denotes that x is not an element of S

$x \in S$ denotes that x is contained in S .



Set Examples

$\emptyset, \{ \}$: the empty set, the null set
 $\{\text{Larry, Joe, Moe}\}$

\mathbb{N} : the set of natural numbers including 0

\mathbb{N}^+ : the set of natural numbers excluding 0

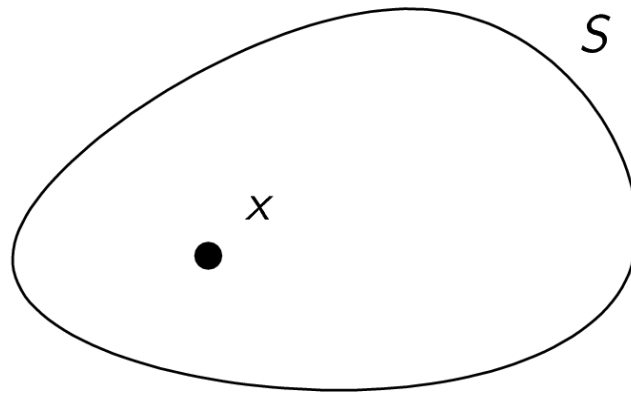
\mathbb{Z} : the set of integers

\mathbb{Q} : the set of rational numbers

\mathbb{R} : the set of real numbers

Defining Sets

$x \in S$ denotes that x is contained in S .

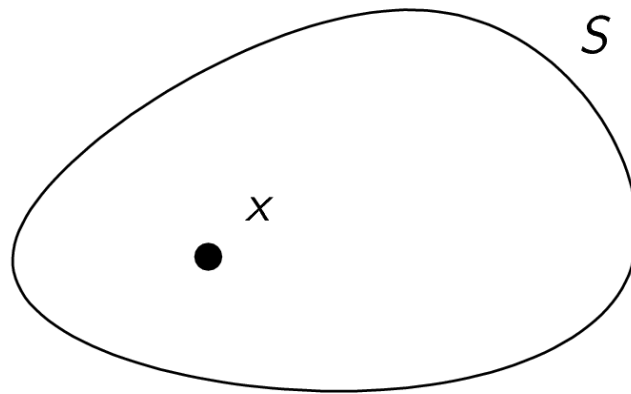


Two ways to define sets:

(1) Listing elements of a set: $S = \{x_1, x_2, \dots, x_n\}$

What is the
limitation of this
approach?

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What is the
limitation of this
approach?

Two ways to define sets:

- (1) Listing elements of a set: $S = \{x_1, x_2, \dots, x_n\}$
- (2) Denoting defining property of a set: $\{x \in S \mid x \text{ has property } p\}$
(= set comprehension)

How can we describe properties and relationships of variables for set comprehension ?

→ Use predicate logic

Example

The set of all even numbers
=

How can we describe properties and relations of variables for set comprehension ?

☛ Use predicate logic

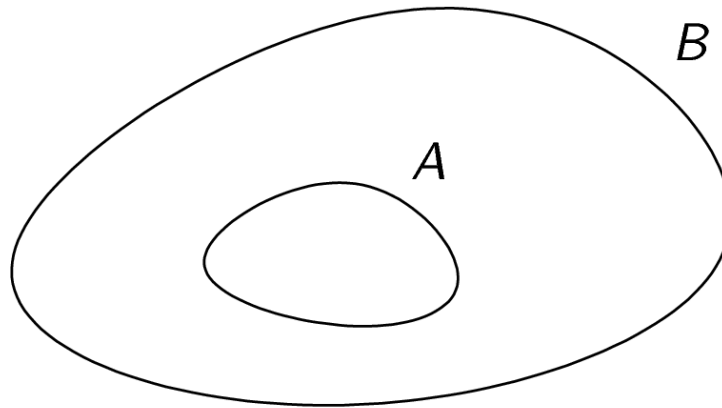
Example

The set of all even numbers

$$= \{x \in \mathbb{Z} \mid \exists y (y \in \mathbb{Z} \wedge x = 2 \times y) \}$$

2. Subset Relationship

“ $A \subseteq B$ ” means “ $(\forall x)(x \in A \rightarrow x \in B)$.”
definition



For example,

$$\mathbf{N \subseteq Z \subseteq Q \subseteq R.}$$

Set Equality

Definition 1. (Set Equality)

Two sets A and B are equal (written $A = B$) if and only if $A \subseteq B$ and $B \subseteq A$.

Definition 2. (Set Equality)

Two sets A and B are equal ($A = B$) if and only if $\forall x(x \in A \leftrightarrow x \in B)$ holds.

3. Constructing New Sets from Existing Sets

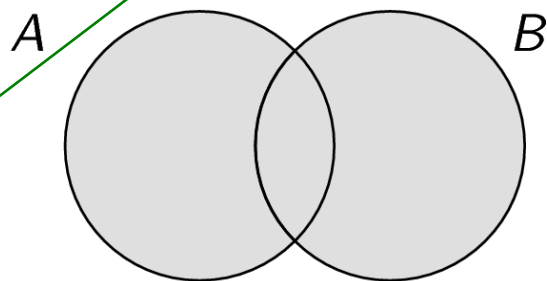
Union and Intersection

Definition of set union: $A \cup B = \{x \mid (x \in A) \vee (x \in B)\}$.

In predicate logic: $(\forall x)(x \in A \cup B \leftrightarrow (x \in A) \vee (x \in B))$

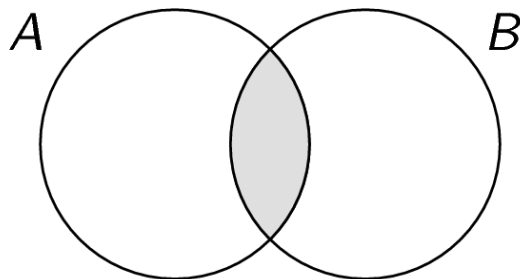
That is, rewriting the
definition in predicate logic

Operations being
defined



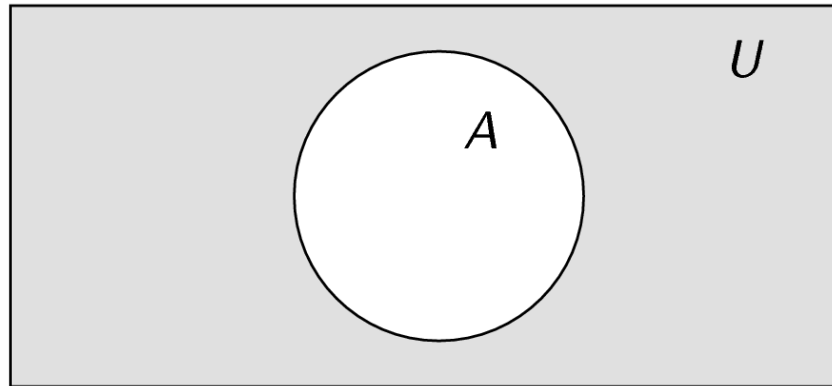
Definition of set intersection: $A \cap B = \{x \mid (x \in A) \wedge (x \in B)\}$.

In predicate logic: $(\forall x)(x \in A \cap B \leftrightarrow (x \in A) \wedge (x \in B))$



Complement

Define: $A' = \{x \in U \mid x \notin A\}$.

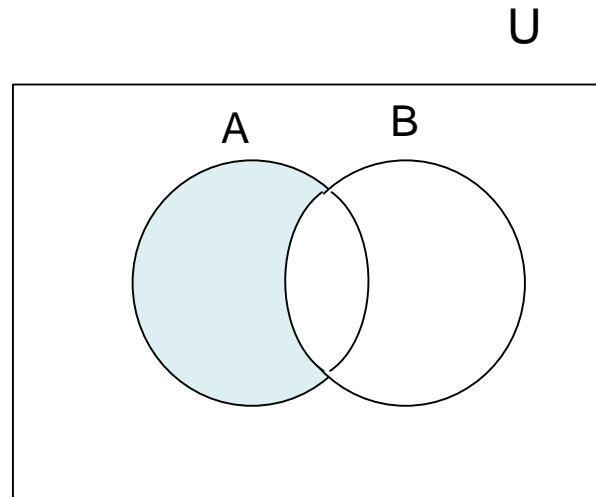


Note that we could also write $A' = \{x \in U \mid \neg(x \in A)\}$ to make the use of the \neg connective explicit.

Set Difference

$$A - B = \{x \in U \mid x \in A \wedge \neg(x \in B)\}$$

"A - B" is also written " $A \setminus B$ "



Examples

Let $A = \{1, 2, 3, 4, 5\}$, $B = \{4, 5, 6, 7, 8\}$, and suppose the universal set is $U = \{1, 2, \dots, 10\}$. Then

$$A \cup B =$$

$$A \cap B =$$

$$B' =$$

$$A - B =$$

Examples

Let $A = \{1, 2, 3, 4, 5\}$, $B = \{4, 5, 6, 7, 8\}$, and suppose the universal set is $U = \{1, 2, \dots, 10\}$. Then

$$A \cup B = \{1, 2, \dots, 8\}$$

$$A \cap B = \{4, 5\}$$

$$B' = \{1, 2, 3, 9, 10\}$$

$$A - B = \{1, 2, 3\}$$

Cartesian products and power sets

The *Cartesian product* $A \times B$ of two sets A and B is the set of all ordered pairs where the first item comes from the first set and the second item comes from the second set. Formally,

$$A \times B = \{(a, b) \mid a \in A \text{ and } b \in B\}.$$

The *power set* $\mathcal{P}(S)$ of the set S is the set of all subsets of S :

$$\mathcal{P}(S) = \{X \mid X \subseteq S\}.$$

Examples

Let $A = \{1, 2, 3, 4, 5\}$, $B = \{4, 5, 6, 7, 8\}$, and suppose the universal set is $U = \{1, 2, \dots, 10\}$. Then

$$(A \cap B) \times A =$$

$$\mathcal{P}(A \cap B) =$$

Examples

Let $A = \{1, 2, 3, 4, 5\}$, $B = \{4, 5, 6, 7, 8\}$, and suppose the universal set is $U = \{1, 2, \dots, 10\}$. Then

$$(A \cap B) \times A = \{(4, 1), (4, 2), (4, 3), (4, 4), (4, 5), (5, 1), \\ (5, 2), (5, 3), (5, 4), (5, 5)\}$$

$$\mathcal{P}(A \cap B) = \{\emptyset, \{4\}, \{5\}, \{4, 5\}\}$$

Set of Strings over an Alphabet

Let Σ be a set of symbols $\{a_1, a_2, \dots, a_n\}$, called an *alphabet*.

Let a string be a sequence of symbols. (We will defined this precisely later.)

The string of length 0 is called the *empty string* or the null string and is denoted λ ($= \varepsilon = ""$).

Σ^0 : set of strings of length 0, i.e. $\{\lambda\}$

Σ^1 : set of strings of length 1, i.e. $\{a_1, a_2, \dots, a_n\} = \Sigma$

...

Σ^2 : set of strings of length 2, i.e. $\{a_1a_1, a_2a_1, a_2a_1, a_2a_2, \dots, a_na_n\}$

...

Σ^k : set of strings of length k

Σ^* : the set of all strings of finite length with symbols from Σ , i.e.

$$\cup \{\Sigma^k \mid k \in \mathbb{N}\} = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots \cup \Sigma^k \cup \dots$$

Cardinality of Set

Cardinality of a set S is the number of elements in S and written $|S|$.

Example

(1) If $|A| = a$ and $|B| = b$, what is $|A \times B|$?

(2) What is $|\mathcal{P}(S)|$?

Cardinality of Set

Cardinality of a set S is the number of elements in S and written $|S|$.

Example

(1) If $|A| = a$ and $|B| = b$, what is $|A \times B|$? $a \times b$

(2) What is $|\mathcal{P}(S)|$? $2^{|S|}$

4. Set Identities

Identities from propositional logic

The inference rules for propositional logic give identities for set theory. For example, the addition inference rule

$$q \vee p \Leftrightarrow p \vee q$$

allows us to prove the identity

$$B \cup A = A \cup B$$

in set theory.

2nd Definition of Set Equality

When A and B are sets,
 $A = B$ iff $\forall x(x \in A \leftrightarrow x \in B)$.

Proof that $B \cup A = A \cup B$

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Proof that $B \cup A = A \cup B$

Let $x \in B \cup A$. Then, by definition of \cup , $x \in B \vee x \in A$, which is equivalent to $x \in A \vee x \in B$, by commutativity of \vee . Then the last formula implies $x \in A \cup B$, by definition of \cup . So $B \cup A \subseteq A \cup B$. Similarly, $A \cup B \subseteq B \cup A$. Thus $B \cup A = A \cup B$.

TABLE 1 Set Identities.

<i>Identity</i>	<i>Name</i>
$A \cap U = A$ $A \cup \emptyset = A$	Identity laws
$A \cup U = U$ $A \cap \emptyset = \emptyset$	Domination laws
$A \cup A = A$ $A \cap A = A$	Idempotent laws
$\overline{\overline{A}} = A$	Complementation law
$A \cup B = B \cup A$ $A \cap B = B \cap A$	Commutative laws
$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$ $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$	Associative laws
$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$ $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$	Distributive laws

Source:
[Rosen 19] p.136

<i>Identity</i>	<i>Name</i>
$\overline{A \cap B} = \overline{A} \cup \overline{B}$ $\overline{A \cup B} = \overline{A} \cap \overline{B}$	De Morgan's laws
$A \cup (A \cap B) = A$ $A \cap (A \cup B) = A$	Absorption laws
$A \cup \overline{A} = U$ $A \cap \overline{A} = \emptyset$	Complement laws

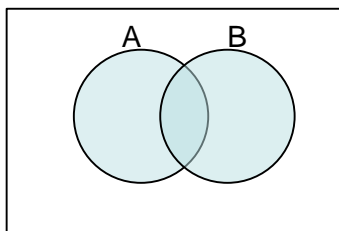
Theorem (Set Theoretic version of De Morgan's laws).

Let A and B be sets. Then

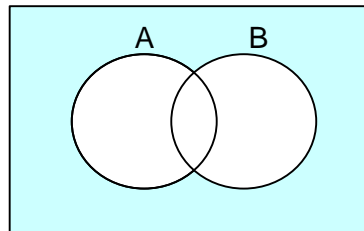
1 $(A \cup B)' = A' \cap B'$

2 $(A \cap B)' = A' \cup B'$

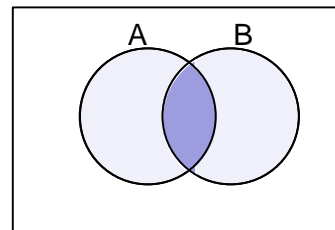
[1a]



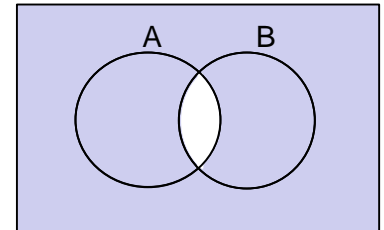
[1b]



[2a]



[2b]



Prove Part [1] of the theorem.

De Morgan's Laws: proof

Proof.

We will prove part 1. (Part 2 is similar.) Let $x \in (A \cup B)'$. In other words, if $P(x)$ is the statement " $x \in A$ " and $Q(x)$ is the statement " $x \in B$," we are starting with the assumption $\neg(P(x) \vee Q(x))$. By De Morgan's laws for propositional logic, this is equivalent to $\neg P(x) \wedge \neg Q(x)$, which, in the language of set theory, is the same as $x \in A' \cap B'$. We have shown that

$$x \in (A \cup B)' \Leftrightarrow x \in A' \cap B',$$

hence the sets $(A \cup B)'$ and $A' \cap B'$ are equal. □

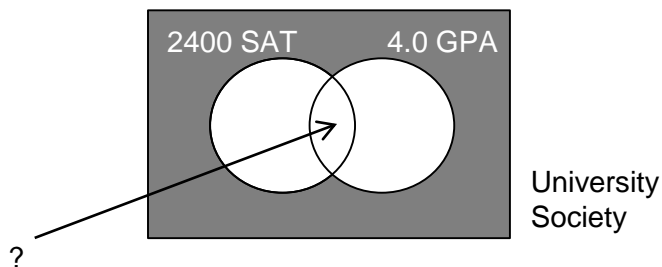
Inclusion-exclusion principle

For finite sets A and B , $|A \cup B| = |A| + |B| - |A \cap B|$

Example

The Master of the Universe Society at a certain college accepts members who have 2400 SAT's or 4.0 GPA's in high school. Of the 11 members of the society, 8 had 2400 SAT's and 5 had 4.0 GPA's. How many members had both 2400 SAT's and 4.0 GPA's?

Solution



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Solution

$11 = 8 + 5 - |A \cap B|$, so there are two.

5. Discussion on Set Comprehension

Consider the following:

The barber is the “one who shaves all those, and those only, who do not shave themselves,” The question is, does the barber shave himself?

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Can you describe it with set comprehension?

Let A be the set of such barbers. So

$$\begin{aligned} A &= \{x \mid x \text{ shaves } y \text{ iff } y \text{ does not shave } y\} \\ &= \{x \mid \forall y (x \text{ shaves } y \leftrightarrow \neg(y \text{ shaves } y)) \} \end{aligned}$$

If $z \in A$, does z shave z ?

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If $z \in A$, does z shave z ?

Let $z \in A$. Then

$$\forall y (x \text{ shaves } y \leftrightarrow \neg(y \text{ shaves } y))$$

and by the \forall -elim rule

$$z \text{ shaves } z \leftrightarrow \neg(z \text{ shaves } z)$$

which is always false. Therefore $z \notin A$.

What does this mean?

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What does this mean?

$$A = \emptyset.$$

☛ Called *the Barber's paradox* (Not a genuine paradox. It is only paradoxical.)

Is $\{1\} \notin \{1\}$ true? Or is the set $\{1\}$ an element of itself?

Is it a “meaningful” statement?

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Is it a “meaningful” statement?

Consider the following set comprehension.

$$A = \{x \mid x \notin x\}$$

Is $A \in A$ true?

☛ If true, then $A \in A$, hence $A \notin A$ must be true.

If false, then $A \notin A$, hence $A \in A$ must be true.

Neither true nor false.

Is $\{1\} \notin \{1\}$ true? Or is the set $\{1\}$ an element of itself?

Is it a “meaningful” statement?

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Is $A \in A$ true?

If true, then $A \in A$, hence $A \notin A$ must be true.

If false, then $A \notin A$, hence $A \in A$ must be true.

Neither true nor false.

☛ Russell's paradox

- ☛ Should not allow all possible descriptions for set comprehension.
- ☛ Restrict the underlying domains of variables.

Quiz 10-1

Let S be a set.

[1] What is the smallest subset of S ?

[2] What is the largest subset of S ?

[3] What is the smallest element of $P(S)$?

[4] What is the largest element of $P(S)$?

[5] What is the smallest subset of $P(S)$?

[6] What is the largest subset of $P(S)$?