

Predicate Logic - Part II

f) a) Statements	Reasons	
$\neg(\exists x)(R(x) \wedge \beta(x))$	given	①
$(\forall x) \neg(R(x) \wedge \beta(x))$	Existential negation, 1	②
$(\forall x) (\neg R(x) \vee \neg \beta(x))$	De Morgan's Law, 2	③
$(\forall x) (R(x) \rightarrow \neg \beta(x))$	implication, 3	④

Now, we'll prove the other direction

Statements	Reasons
1) $(\forall x) (R(x) \rightarrow \neg \beta(x))$	given
2) $(\forall x) (\neg R(x) \vee \neg \beta(x))$	implication, 1
3) $(\forall x) \neg(R(x) \wedge \beta(x))$	De Morgan's Law, 2
4) $\neg(\exists x)(R(x) \wedge \beta(x))$	Existential negation, 3

Hence, we proved that $\neg(\exists x)(R(x) \wedge \beta(x)) \Leftrightarrow \Leftrightarrow (\forall x)(R(x) \rightarrow \neg \beta(x))$ or $S_1 \Leftrightarrow S_0$

b) There does not exist a right triangle that has an obtuse angle

c) A right triangle does not have an obtuse angle ■

2) a) Assume domain \rightarrow integers, that is, given domain is \mathbb{Z} , or all integers. $P(x) = x$ is even

$Q(x) = x$ is odd

$(\exists x)(P(x) \wedge Q(x))$ means there exist an integer x such that x is even and x is odd. However, we know integer can never be both even and odd \Rightarrow

$(\exists x)(P(x) \wedge Q(x))$ is false. $(\exists x)(P(x))$ means

there exist an integer x such that x is even, which is obviously true (take $x=2$) \Rightarrow $(\exists x)(P(x))$ is true

$(\exists x)(Q(x))$ means there exists an integer which is odd; definitely true (take integer as equal to 1) \Rightarrow

$(\exists x)(Q(x))$ is true. Hence, $(\exists x)P(x) \wedge (\exists x)Q(x) \Rightarrow$
 \Rightarrow is always True

whereas $(\exists x)(P(x) \wedge Q(x))$ is always False / This concludes these statements are not logically equivalent

b) $(\exists x)(P(x) \vee Q(x))$ means there exist an integer x such that x is even or x is odd, which is definitely true as we can pick $x=2$ for example

Since $P(x) \vee Q(x)$ is disjunction \Rightarrow $(\exists x)(P(x) \vee Q(x)) \rightarrow \text{True}$

$(\exists x)P(x) \rightarrow$ means there exist integer which is even; definitely true, as we can pick integer equal to 10

$(\exists x)P(x) \rightarrow \text{True}$ Similarly, $(\exists x)Q(x) \rightarrow$ means there exist x -integer for which x is odd \Rightarrow pick $x=1$ and it becomes true $(\exists x)Q(x) \rightarrow \text{True}$ Therefore, we've

$(\exists x)P(x) \vee (\exists x)Q(x) \rightarrow \text{True}$ Because both of them are True

This concludes our example from part (a) satisfying the equivalence $(\exists x)(P(x) \vee Q(x)) \Leftrightarrow (\exists x)P(x) \vee (\exists x)Q(x)$

3) a) Pick $x=0$ and $0^2 > 0$ is false \rightarrow meaning this inequality is not true for all real numbers x False

b) Choose $x=\sqrt{3}$, then $x^2-2=3-2=1$, meaning there exist x such that $x^2-2=1 \Rightarrow$ True

c) If there exist $x \in \mathbb{R}$ such that $x^2+2=1$, then $x^2+1=0$ ^{$x^2=-1$} for some real number x . However, we know there does not exist such solution in the domain of real numbers ($x=\pm i$ is solution, where $i=\sqrt{-1}$) It only exists in complex numbers \Rightarrow False

d) For any real number x , choose $y=4-x^2$. It's obvious that y will be real number and we've

$x^2 + y = 4$, meaning for all real numbers x , we can find real number y (particularly $y = 4 - x^2$ which is a real number) such that $x^2 + y = 4 \Rightarrow \boxed{\text{True}}$

e) Assume that there exist a real number y such that, for all real numbers x , $x^2 + y = 4$. As it's true for all $x \in \mathbb{R}$, we pick $x=0 \Rightarrow y=4$. But if we choose $x=1 \Rightarrow y+1=4$ implies $y=3$, which is contradiction as we had $y=4 \Rightarrow \boxed{\text{X}}$. Hence, there should not exist such real number y False ■

(thus we don't require the same real number for each real number x to make the equation become true)

$$4) a) (\forall x) (\neg P(x) \rightarrow (\exists y) (P(y) \wedge Q(y, x)))$$

b)	Statements	Reasons	
	$\neg(\forall x) (\neg P(x) \rightarrow (\exists y) (P(y) \wedge Q(y, x)))$	Given	①
	$(\exists x) \neg(\neg P(x) \rightarrow (\exists y) (P(y) \wedge Q(y, x)))$	Universal negation, 1	②
	$(\exists x) \neg(\neg P(x)) \vee (\exists y) (P(y) \wedge Q(y, x))$	Implication, 2	③
	$(\exists x) \neg(\neg P(x)) \vee (\exists y) \neg(P(y) \wedge Q(y, x))$	Double negation, 3	④
	$(\exists x) (\neg P(x) \wedge \neg(\exists y) (P(y) \wedge Q(y, x)))$	De Morgan's Law, 4	⑤
	$(\exists x) (\neg P(x) \wedge (\forall y) \neg(P(y) \wedge Q(y, x)))$	Existential negation, 5	⑥
	$(\exists x) (\neg P(x) \wedge (\forall y) (\neg P(y) \vee \neg Q(y, x)))$	De Morgan's Law, 6	⑦
	$(\exists x) (\neg P(x) \wedge (\forall y) (P(y) \rightarrow \neg Q(y, x)))$	Implication, 7	⑧
Hence, we found that $\rightarrow (\exists x) (\neg P(x) \wedge (\forall y) (P(y) \rightarrow \neg Q(y, x)))$			

c) There is a non-prime x such that no prime y divides it ■ (There is a non-prime x such that for all primes y , y does not divide x)

5) We proceed the proof by contradiction. Let x, m, n be ^{distinct} lines, and suppose $x \parallel m$ and $m \parallel n \Rightarrow$ Suppose, to the contrary, that x is not parallel to n . By the definition of "parallel", this means there exist a point x on both x and n . Since $x \parallel m$, x should not be on m (otherwise, x is on m and x is on x would mean x and m are not parallel).

According to Axiom 2, it says there exist a unique line such that x is on the line and that line is parallel to m . However, that contradicts the fact x and n are both ^{distinct} lines which contain point x and both are parallel to $m \Rightarrow \boxed{X}$

Hence, our assumption was wrong $\Rightarrow \boxed{x \parallel n}$ ■

6) a) $(\forall x)(\exists y) P(y, x)$

b) Using universal & existential negation rules,
 $\neg(\forall x)(\exists y) P(y, x) \Leftrightarrow (\exists x) \neg(\exists y) P(y, x) \Leftrightarrow (\exists x)(\forall y) \neg P(y, x)$

Hence, we found negation of $(\forall x)(\exists y) P(y, x)$ expression from (a) $\Rightarrow \boxed{(\exists x)(\forall y) \neg P(y, x)}$

2) Some number is greater than or equal to all numbers

There is a number that all numbers are less than or equal to that number ■

†) \leftrightarrow -intro
$$\frac{P \rightarrow Q, Q \rightarrow P}{P \leftrightarrow Q}$$

The rule is that wherever instances of " $P \rightarrow Q$ " and " $Q \rightarrow P$ " appear on lines of a proof, " $P \leftrightarrow Q$ " can validly be placed on a subsequent line

The "Biconditional introduction" rule may be written in sequent notation: $(P \rightarrow Q), (Q \rightarrow P) \vdash (P \leftrightarrow Q)$ where \vdash is a metalogical symbol meaning that $P \leftrightarrow Q$ is a syntactic consequence when $P \rightarrow Q$ and $Q \rightarrow P$ are both in a proof

\leftrightarrow -elim
$$\frac{P \leftrightarrow Q}{P \rightarrow Q}$$

\leftrightarrow -elim
$$\frac{P \leftrightarrow Q}{Q \rightarrow P}$$

The rule is that wherever an instance of " $P \leftrightarrow Q$ " appears on a line of a proof, either " $P \rightarrow Q$ " or " $Q \rightarrow P$ " can be replaced on a subsequent line

The "Biconditional elimination" rule may be written in sequent notation: $(P \leftrightarrow Q) \vdash (P \rightarrow Q)$ and $(P \leftrightarrow Q) \vdash (Q \rightarrow P)$ where \vdash is a metalogical symbol meaning that $P \rightarrow Q$, in the 1st case, and $Q \rightarrow P$ in the other are syntactic consequences of $P \leftrightarrow Q$ in some logical system

8) The given diagram contains 4 Baddas, which are thus the 4 inner points that are touched by 4 line segments each

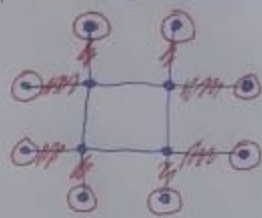
We see that each Badda is hit by exactly 4 Bings, because each of 4 inner points are touched by 4 line segments each \Rightarrow Axiom 1 is satisfied

Now, we observe each Bing is not hit by exactly 2 Baddas, because there are 8 line segments that touch only 1 Bing \Rightarrow Axiom 2 is not satisfied

We see that 2 Bings hit the same two Baddas, as there's no more than 1 line segment between each pair of points \Rightarrow Axiom 3 is satisfied

Finally, the diagram contains at least 1 Bing, since diagram has 4 Bings \Rightarrow Axiom 4 is satisfied

Because axiom 2 is not satisfied \Rightarrow this is not a model for the system



For Axiom 2, these 8 points provide evidence that not every Bing is hit by exactly 2 Baddas. Individually, these 8 points prove that those Bings are hit by exactly one Baddas (individually).

For Axiom 3, it's just saying if two Baddas intersect at Bing q , then there's no other Bing hit by those two Baddas. It is essentially saying no other intersection, which is obvious because

if 2 lines intersect, they have only 1 common point and it infers that there is no other line hit by x and y . If 2 lines have a common point, it should not have other common point, implying for distinct saddles x and y , each hitting line $\varphi \Rightarrow$ no other lines are hit by both x and y

$$g) \quad \neg (\exists y) (H(y) \wedge \text{IsTailPop}(x, y))$$

$$b) \text{ Horses are animals: } (\forall y) (H(y) \rightarrow A(y))$$

Horses' tails are tails of animals:

$$(\forall x)(\forall y)((H(y) \wedge \text{IsTailPop}(x, y)) \rightarrow (A(y) \wedge \text{IsTailPop}(x, y)))$$

c) Inference in (b) is valid, meaning that argument is valid

$$(\forall y) (H(y) \rightarrow A(y))$$

$$(\forall x)(\forall y)((H(y) \wedge \text{IsTailPop}(x, y)) \rightarrow (A(y) \wedge \text{IsTailPop}(x, y)))$$

d) We can formalize our predicate logic inference using natural deduction as follows:

$$\frac{\forall y (H(y) \rightarrow A(y))}{H(y) \rightarrow A(y)} \text{ } \forall\text{-elim}$$

$$\frac{[H(y) \wedge \exists x \text{TaifOP}(x, y)]^*}{H(y)} \text{ } \wedge\text{-elim}$$

$$H(y) \rightarrow A(y)$$

$$H(y)$$

$$\rightarrow\text{-elim}$$

$$A(y)$$

$$\exists x \text{TaifOP}(x, y)$$

$$A(y) \wedge \exists x \text{TaifOP}(x, y)$$

$$(H(y) \wedge \exists x \text{TaifOP}(x, y)) \rightarrow (A(y) \wedge \exists x \text{TaifOP}(x, y))$$

$$\forall y ((H(y) \wedge \exists x \text{TaifOP}(x, y)) \rightarrow (A(y) \wedge \exists x \text{TaifOP}(x, y)))$$

$$(\forall x (\forall y ((H(y) \wedge \exists x \text{TaifOP}(x, y)) \rightarrow (A(y) \wedge \exists x \text{TaifOP}(x, y))))$$

Note: $[A]^*$ stands for an assumption A discharged by the rule \star
 $[A]$ means that the hypotheses A is discharged

3) a) Note: If "x is a tail of a horse" means that x is a tail of every horse, then our predicate logic expression would be $(\forall y)(H(y) \rightarrow \text{IsTailOf}(x, y))$
However, we accepted that x is a tail of some horse, because of "a" article. Therefore, we implemented existential quantifier in first place