## Why is a single depth buffer sufficient for this vulkan swapchain render loop?

Asked 4 years, 6 months ago Modified 7 months ago Viewed 4k times



**13** 



I was following the vulkan tutorial at <a href="https://vulkan-tutorial.com/">https://vulkan-tutorial.com/</a> and at the <a href="depth buffering chapter">depth buffering chapter</a>, the author Alexander Overvoorde mentions that "We only need a single depth image, because only one draw operation is running at once." This is where my issue comes in.



I've read many SO questions and articles/blog posts on Vulkan synchronization in the past days, but I can't seem to reach a conclusion. The information that I've gathered so far is the following:

Draw calls in the same subpass execute on the gpu *as if* they were in order, but only if they draw to the framebuffer (I can't recall exectly where I read this, it might have been a tech talk on youtube, so I am not 100% sure about this). As far as I understood, this is more GPU hardware behavior than it is Vulkan behaviour, so this would essentially mean that the above is true in general (including across subpasses and even render passes) - which would answer my question, but I can't find any clear information on this.

The closest I've gotten to getting my question answered is this <u>reddit comment</u> that the OP seemed to accept, but the justification is based on 2 things:

- "there is a queue flush at the high level that ensures previously submitted render passes are finished"
- "the render passes themselves describe what attachments they read from and write to as external dependencies"

I see neither any high level queue flush (unless there is some sort of explicit one that I cannot find for the life of me in the specification), nor where the render pass describes dependencies on its attachments - it describes the attachments, but not the dependencies (at least not explicitly). I have read the relevant chapters of the specification multiple times, but I feel like the language is not clear enough for a beginner to fully grasp.

I would also really appreciate Vulkan specification quotes where possible.

Edit: to clarify, the final question is: What synchronization mechanism guarantees that the draw call in the next command buffer is not submitted until the current draw call is finished?

graphics 3d vulkan

asked Jun 14, 2020 at 10:29

cluntraru

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This is not part of the question, but I'd also appreciate pointers on articles or preferably books on this sort of GPU behavior that's relevant to programming. I've ordered both Learning Vulkan by Parminder Singh and Vulkan Programming Guide by Graham Sellers, but they haven't arrived yet and don't seem to include too much on GPU hardware anyway (I might be wrong though). I unfortunately don't like the format of the Vulkan Cookbook by Pawel Lapinski that I read is one of the better options - I much prefer properly grasping the theory and doing things myself than following a "recipe" — cluntraru Jun 14, 2020 at 10:30

Does this answer your question? <u>Synchronization between</u> <u>drawcalls in Vulkan</u> – krOoze Jun 14, 2020 at 11:13

- Not exactly. It explains that draws in the same subpass are executed as if in order, but what about different subpasses (or almost identical command buffers submitted twice in a row in this case)? It mentions that those are governed by external dependencies, but the only external subpass dependency here would be the one used to synchronize with the imageAcquired semaphore that has no srcAccessMask and, as a result, should not actually wait on any color attachment stages since there is no resource it needs to access (correct me if I'm wrong). cluntraru Jun 14, 2020 at 11:27
- 1 Ran out of characters... in this case, wouldn't the following drawFrame() call run with no synchronization between the 2 and cause something like: render 1 starts, then render 2 starts, then render 1 finishes, then present 1 starts, then

render 2 finishes etc.? In this case, the depth buffer would be reused and presumably broken. — cluntraru Jun 14, 2020 at 11:30 ✓

@cluntraru: To answer your question would require analyzing the dependency graph of the tutorial code in question, which is a pretty big thing to ask. For all I know, there are explicit events or semaphore waits that prevent overlap between frames. I don't know if that's true, but I also don't know it isn't true. And the only way to find out would be to read the entire code of the tutorial. – Nicol Bolas Jun 14, 2020 at 14:00

## 4 Answers

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I'm afraid, I have to say that the Vulkan Tutorial is wrong. In its current state, it can not be guaranteed that there are no memory hazards when using only one single depth buffer. However, it would require only a very small change so that only one depth buffer would be sufficient.



Let's analyze the relevant steps of the code that are performed within drawFrame.



We have two different queues: presentQueue and graphicsQueue, and MAX\_FRAMES\_IN\_FLIGHT concurrent frames. I refer to the "in flight index" with cf (which stands for currentFrame = (currentFrame + 1) % MAX\_FRAMES\_IN\_FLIGHT). I am using sem1 and sem2 to represent the different arrays of semaphores and fence for the array of fences.

The relevant steps in pseudocode are the following:

```
vkWaitForFences(..., fence[cf], ...);
vkAcquireNextImageKHR(..., /* signal when done: */
sem1[cf], ...);
vkResetFences(..., fence[cf]);
vkQueueSubmit(graphicsQueue, ...
    /* wait for: */ sem1[cf], /* wait stage: *,
COLOR ATTACHMENT OUTPUT ...
    vkCmdBeginRenderPass(cb[cf], ...);
      Subpass Dependency between EXTERNAL -> 0:
          srcStages = COLOR_ATTACHMENT_OUTPUT,
          srcAccess = 0,
          dstStages = COLOR_ATTACHMENT_OUTPUT,
          dstAccess = COLOR_ATTACHMENT_WRITE
      vkCmdDrawIndexed(cb[cf], ...);
      (Implicit!) Subpass Dependency between 0 ->
EXTERNAL:
          srcStages = ALL_COMMANDS,
          srcAccess =
COLOR ATTACHMENT WRITE | DEPTH STENCIL WRITE,
          dstStages = BOTTOM_OF_PIPE,
          dstAccess = 0
    vkCmdEndRenderPass(cb[cf]);
    /* signal when done: */ sem2[cf], ...
    /* signal when done: */ fence[cf]
);
vkQueuePresent(presentQueue, ... /* wait for: */
sem2[cf], ...);
```

The draw calls are performed on one single queue: the graphicsQueue. We must check if commands on that graphicsQueue could theoretically overlap.

Let us consider the events that are happening on the graphicsQueue in chronological order for the first two frames:

```
img[0] -> sem1[0] signal -> t|...|ef|fs|lf|co|b ->
sem2[0] signal, fence[0] signal
```

```
img[1] \rightarrow sem1[1] signal \rightarrow t|...|ef|fs|lf|co|b \rightarrow sem2[1] signal, fence[1] signal
```

where  $t \mid ... \mid ef \mid fs \mid lf \mid co \mid b$  stands for the different pipeline stages, a draw call passes through:

```
t ... TOP_OF_PIPE
ef ... EARLY_FRAGMENT_TESTS
fs ... FRAGMENT_SHADER
lf ... LATE_FRAGMENT_TESTS
co ... COLOR_ATTACHMENT_OUTPUT
b ... BOTTOM_OF_PIPE
```

While there **might** be an implicit dependency between sem2[i] signal -> present and sem1[i+1], this only applies when the swap chain provides only one image (or if it would always provide the same image). In the general case, this can not be assumed. That means, there is nothing which would delay the **immediate** progression of the subsequent frame after the first frame is handed over to present. The fences also do not help because after fence[i] signal, the code waits on fence[i+1], i.e. that also does not prevent progression of subsequent frames in the general case.

What I mean by all of that: The second frame starts rendering **concurrently** to the first frame and there is nothing that prevents it from accessing the depth buffer concurrently as far as I can tell.

## The Fix:

If we wanted to use only a single depth buffer, though, we can fix the tutorial's code: What we want to achieve is that the ef and If stages wait for the previous draw call to complete before resuming. I.e. we want to create the following scenario:

```
img[0] -> sem1[0] signal -> t|...|ef|fs|lf|co|b ->
sem2[0] signal, fence[0] signal
img[1] -> sem1[1] signal ->
t|...|____|ef|fs|lf|co|b -> sem2[1] signal,
fence[1] signal
```

where \_ indicates a wait operation.

In order to achieve this, we would have to add a barrier that prevents subsequent frames performing the <code>EARLY\_FRAGMENT\_TEST</code> and <code>LATE\_FRAGMENT\_TEST</code> stages at the same time. There is only one queue where the draw calls are performed, so only the commands in the <code>graphicsQueue</code> require a barrier. The "barrier" can be established by using the subpass dependencies:

```
srcAccess =
DEPTH STENCIL ATTACHMENT WRITE,
          dstStages =
EARLY_FRAGMENT_TEST | LATE_FRAGMENT_TEST,
          dstAccess =
DEPTH_STENCIL_ATTACHMENT_WRITE|DEPTH_STENCIL_ATTACHM
      vkCmdDrawIndexed(cb[cf], ...);
      (Implicit!) Subpass Dependency between 0 ->
EXTERNAL:
          srcStages = ALL_COMMANDS,
          srcAccess =
COLOR_ATTACHMENT_WRITE | DEPTH_STENCIL_WRITE,
          dstStages = BOTTOM_OF_PIPE,
          dstAccess = 0
    vkCmdEndRenderPass(cb[cf]);
    /* signal when done: */ sem2[cf], ...
    /* signal when done: */ fence[cf]
);
vkQueuePresent(presentQueue, ... /* wait for: */
sem2[cf], ...);
```

This should establish a proper barrier on the graphicsQueue between the draw calls of the different frames. Because it is an EXTERNAL -> 0 -type subpass dependency, we can be sure that renderpass-external commands are synchronized (i.e. sync with the previous frame).

Update: Also the wait stage for sem1[cf] has to be changed from COLOR\_ATTACHMENT\_OUTPUT to EARLY\_FRAGMENT\_TEST. This is because layout transitions happen at vkCmdBeginRenderPass time: after the first synchronization scope (srcStages and srcAccess) and before the second synchronization scope (dstStages and dstAccess). Therefore, the swapchain image must be

available there already so that the layout transition happens at the right point in time.

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edited Jun 16, 2020 at 9:58

answered Jun 15, 2020 at 22:55



Thank you for the complete answer! I think it might be worth clarifying that the old subpass dependency is not deleted (still needed for semaphore sync), and that there are now 2 from EXTERNAL to 0, just to avoid potential confusion for anyone else reading this. Also, shouldn't dstAccess in the fix also include DEPTH\_STENCIL\_ATTACHMENT\_READ, not just WRITE? — cluntraru Jun 16, 2020 at 6:35

Regarding the "old subpass dependency": Do you mean that COLOR\_ATTACHMENT\_OUTPUT ->

COLOR\_ATTACHMENT\_OUTPUT would still be necessary? The batch's COLOR\_ATTACHMENT\_OUTPUT stages must wait for sem1[cf] to signal anyways. I do not think that such a dependency must be included in the subpass dependencies, unless I am overlooking something. – j00hi Jun 16, 2020 at 7:55

Regarding the two subpass dependencies: Are you referring to the second " (Implicit!) " subpass dependency? – j00hi Jun 16, 2020 at 7:55

Regarding the dstAccess: Yes, you are right. A depth test always performs read and write access. The read access must see the cleared depth buffer values already.

Synchronizing with DEPTH\_STENCIL\_ATTACHMENT\_WRITE only is not sufficient. I'll update it to

DEPTH\_STENCIL\_ATTACHMENT\_WRITE|DEPTH\_STENCIL\_ATT ACHMENT\_READ . Thanks for catching that! – j00hi Jun 16, 2020 at 7:57

@DanielMarques of course that would be an option---maybe even a good one since it theoretically allows to parallelize more. It's just that the question was about using a single depth buffer, which could be advantageous in extremely memory-limited settings, for example. – j00hi Sep 28, 2021 at 7:16



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No, rasterization order does not (per specification) extend outside a single subpass. If multiple subpasses write to the same depth buffer, then there should be a <code>VkSubpassDependency</code> between them. If something outside a render pass writes to the depth buffer, then there should also be explicit synchronization (via barriers, semaphores, or fences).

FWIW I think the vulkan-tutorial sample is non-conformant. At least I do not see anything that would prevent a memory hazard on the depth buffer. It seems that the depth buffer should be duplicated to MAX\_FRAMES\_IN\_FLIGHT, or explicitly synchronized.

The sneaky part about undefined behavior is that wrong code often works correctly. Unfortunately making sync proofs in the validation layers is little bit tricky, so for now only thing that remains is to simply be careful.

Futureproofing the answer:

What I do see is conventional WSI semaphore chain

(used with vkAnquireNextImageKHR and
vkQueuePresentKHR) with imageAvailable and
renderFinished semaphores. There is only one subpass
dependency with

VK\_PIPELINE\_STAGE\_COLOR\_ATTACHMENT\_OUTPUT\_BIT, that is chained to the imageAvailable semaphore. Then there are fences with MAX\_FRAMES\_IN\_FLIGHT == 2, and fences guarding the individual swapchain images. Meaning two subsequent frames should run unimpeded wrt each other (except in the rare case they acquire the same swapchain image). So, the depth buffer seems to be unprotected between two frames.

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edited Jun 15, 2020 at 22:45

answered Jun 15, 2020 at 17:15



This what I suspected as well, leading to my posting the question. Thank you for putting in the time to look through the code, I know it was a lot to ask. – cluntraru Jun 15, 2020 at 19:55

If Vulkan is so complex even experts disagree on simple things (like writing to a depth buffer correctly) then that's an indication Vulkan just needs a do over. The fact incorrect code often still runs (at least on most platforms) makes the situation even worse. – JustinBlaber Jun 28 at 13:40



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Yes, I also spent some time trying to figure out what was meant by the statement "We only need a single depth image, because only one draw operation is running at once."







That didn't make sense to me for a triple buffered rendering setup where work is submitted to the queues until MAX\_FRAMES\_IN\_FLIGHT is reached - there's no guarantee that all three aren't running at once!

Whilst the single depth image worked OK, triplicating everything so each frame uses a fully independent set of resources (blocks and all) would seem to be the safest design and yielded identical performance under test.

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answered May 6, 2022 at 23:32



How the "triplicating everything ... " design fares in terms of memory, in a more complex rendering scenario: multipass rendering for cascade shadow for example where you have 4 layered depth textures ... is that feasable? – Avi Nov 25, 2023 at 10:22

Excellent question!! Since posting, my render engine has evolved a bit - I have linked list based OIT transparency up and running now which needs around 256MB to handle three to four layers of full-screen transparency. I'm certainly not triplicating that... At the moment with full 3D PBR and OIT transparency workflows as well as 2D ortho, I'm running around 665 MB of image buffers. What I'm thinking though is to pull it back to double buffering in terms of frame generation, but maintaining a triple buffered presentation.

- MarvinWS Nov 26, 2023 at 11:29



I believe the tutorial's code is correct and needs only a single depth buffer to function, for the following reasons:









Command buffer submissions to a single queue
 respect submission order. While drivers are free to
 reorder commands in the absence of synchronization
 mechanisms such as semaphores and fences,
 submission order becomes significant once such
 mechanisms are adopted:

"Command buffer submissions to a single queue respect submission order and other implicit ordering guarantees, but otherwise may overlap or execute out of order. Other types of batches and queue submissions against a single queue (e.g. sparse memory binding) have no implicit ordering constraints with any other queue submission or batch. Additional explicit ordering constraints between queue submissions and individual batches can be expressed with semaphores and fences."

 Commands split across different command buffers and submitted to the queue in the order they were recorded will execute <u>as if those commands were</u> recorded to a single command buffer: "Command buffer boundaries, both between primary command buffers of the same or different batches or submissions as well as between primary and secondary command buffers, do not introduce any additional ordering constraints. In other words, submitting the set of command buffers (which can include executing secondary command buffers) between any semaphore or fence operations execute the recorded commands as if they had all been recorded into a single primary command buffer, except that the current state is reset on each boundary."

 A wait operation defined by a semaphore passed to vkQueueSubmit() will block execution of all commands appearing later in submission order, including commands recorded to other command buffers:

"The second synchronization scope includes every command submitted in the same batch. In the case of vkQueueSubmit, the second synchronization scope is limited to operations on the pipeline stages determined by the destination stage mask specified by the corresponding element of pWaitDstStageMask. Also, in the case of vkQueueSubmit, the second synchronization

scope additionally includes all commands that occur later in submission order."

• Execution dependencies for a render pass help ensure that commands which certain other commands depend on will finish before the dependent commands begin execution. The render pass in the Vulkan Tutorial establishes such a happens-before constraint on commands submitted prior to the vkcmdBeginRenderPass() which the program calls every frame:

"If srcSubpass is equal to VK SUBPASS EXTERNAL, the first synchronization scope includes commands that occur earlier in submission order than the vkCmdBeginRenderPass used to begin the render pass instance. Otherwise, the first set of commands includes all commands submitted as part of the subpass instance identified by srcSubpass and any load, store, or multisample resolve operations on attachments used in srcSubpass. In either case, the first synchronization scope is limited to operations on the pipeline stages determined by the source stage mask specified by srcStageMask."

Imagine the following sequence of commands has been submitted to the queue:

```
// frame 1
vkAcquireNextImageKHR(); // signals
imageAvailableSemaphores[0].
RecordRenderPass();
vkQueueSubmit(); // waits on
imageAvailableSemaphores[0] at
COLOR ATTACHMENT OUTPUT stage.
vkQueuePresentKHR();
// frame 2
vkAcquireNextImageKHR(); // signals
imageAvailableSemaphores[1].
RecordRenderPass();
vkQueueSubmit(); // waits on
imageAvailableSemaphores[1] at
COLOR_ATTACHMENT_OUTPUT stage.
vkQueuePresentKHR();
// frame 3
vkAcquireNextImageKHR(); // signals
imageAvailableSemaphores[0].
RecordRenderPass();
vkQueueSubmit(); // waits on
imageAvailableSemaphores[0] at
COLOR ATTACHMENT OUTPUT stage.
vkQueuePresentKHR();
```

In addition to the semaphores referenced above there is also a fence between each frame to ensure that command buffers aren't modified while in their *pending* state, but these fences are not responsible for ensuring that the depth buffer isn't accessed simultaneously by two different frames so I'm ignoring them for the purposes of this analysis.

The first vkQueueSubmit() command will cause the driver to wait for imageAvailableSemaphores[0] to be signaled before moving on to the COLOR\_ATTACHMENT\_OUTPUT stage.

Because the wait operation blocks execution according to submission order, it will also prevent execution of the drawing and presentation commands in frames 2 and 3 until <code>imageAvailableSemaphores[0]</code> has been signaled in the first frame. The execution dependency on the render pass further prevents draw commands for subsequent passes from executing before the first pass has reached the <code>COLOR\_ATTACHMENT\_OUTPUT</code> stage.

Frames are therefore guaranteed to be executed in order (frame 2 won't be rendered before frame 1), and each frame will have finished writing to the depth buffer (in the LATE\_FRAGMENT\_TESTS stage preceding the COLOR\_ATTACHMENT\_OUTPUT stage) before subsequent frames can write to the depth buffer.

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edited Apr 26 at 17:51

answered Dec 22, 2023 at 21:58

