HB+Trie

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Data Management Data Structures

Agenda

- 1. Motivations
- 2. Overview
- 3. Implementation
- 4. Performance
- 5. Possible improvements
- 6. Discussion

Motivations

Motivations

Variable-length sized keys

Disadvantages with B+ tree or LSM-tree:

- Fanout degree decreases if key length increases
- Tree Height grows to maintain the same capacity
- Benefit of prefix B+ tree becomes limited for randomly distributed keys
- B+ tree nodes are randomly scattered on disk when it ages

HB+ trie stands for *Hierarchical B+ tree based trie*

Characteristics:

- Key space is divided into buckets. Every bucket has its own HB+ trie
- High disk throughput due to append-only disk layout
- Disk updates are delayed with a Write buffer index

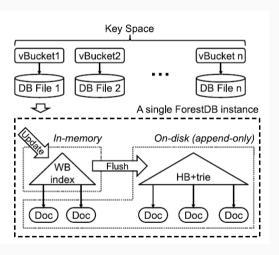


Figure 1: Architecture

HB+ trie stands for *Hierarchical B+ tree based trie*

Characteristics:

- Key space is divided into buckets. Every bucket has its own HB+ trie
- High disk throughput due to append-only disk layout
- Disk updates are delayed with a Write buffer index
- Fixed size chunking of the key
- $\bullet\,$ Every unique chunk has a dedicated B+ tree

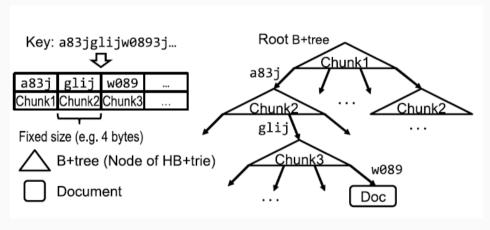


Figure 2: Chunking

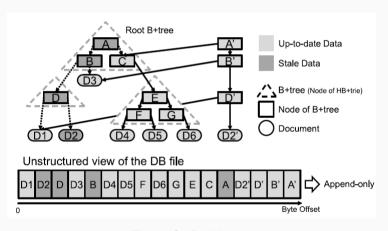


Figure 3: Disk layout

Implementation

Implementation

- Using 16 byte chunks for keys
- Each page frame holds a complete B+ subtree.
- \bullet Storing pageld in the leaf to reference a B+ subtree

Implementation - Chunking

```
func createChunkFromKey(key []byte) (*[16]byte, *[]byte) {
    chunkedKey := [16]byte{}
    var trimmedKey []byte
    if len(kev) > 16 {
        trimmedKey = make([]byte, 0, len(key)-16)
        copy(chunkedKey[:], key[:16])
        trimmedKey = key[16:]
    } else {
        trimmedKey = make([]byte, 0, len(key))
       copy(chunkedKey[:], key[:])
        trimmedKev = kev
    return &chunkedKey, &trimmedKey
```

Implementation - Insert

```
func (hbt *HBTrieInstance) insert(key []byte, value uint64, bpt *bptree.BPlusTree) error {
    chunkedKey, trimmedKey := createChunkFromKey(key)
        subTree, err := hbt.createSubTree(bpt, *chunkedKey)
       if err != nil {
            return err
       return hbt.insert(*trimmedKey, value, subTree)
    } else {
        success, err := bpt.Insert(*chunkedKey, value)
        if success {
           return nil
       return err
```

Implementation - Search

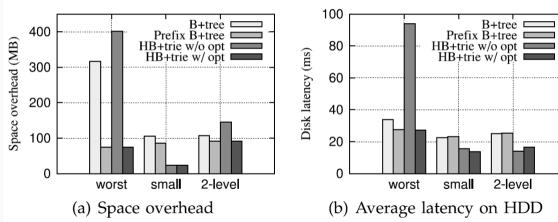
```
• • •
func (hbt *HBTrieInstance) search(bpt *bptree BPlusTree, key []byte) (uint64, []byte,
                                                                       *bptree BPlusTree, error) {
    chunkedKey, trimmedKey := createChunkFromKey(key)
    if err != nil {
        return 0, key, bpt, err
    if val. IsTree {
        subbpt := bptree.LoadBplusTree(hbt.pool, val.Value)
        return hbt.search(subbpt, *trimmedKev)
    } else {
        return val. Value, key, bpt, nil
```

Implementation - Page frame

```
• • •
type Page struct {
   Id uint64 // 8 byte
   Dirty bool // 1 byte
   prev *Page // 8 byte
   next *Page // 8 byte
type Node struct {
                   uint64
                   uint64
                    [120]uint64 // 960 byte
   NumberOfChildren uint64 // 8 byte
   NumberOfEntries uint64
type Entry struct {
   IsTree bool
   Key [16]byte // keys are chunks of 16 bytes
   Value uint64 // values are pointers to subsequent b+ trees
```

Performance

Performance



small 2-level (b) Average latency on HDD

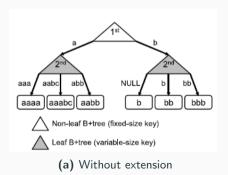
Possible improvements

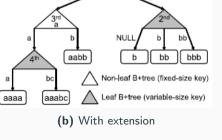
Possible improvements

HB+ trie is not a balanced structure

• Leads to key skew under specific key pattern

To address this issue, Leaf B+ tree extension is proposed





Possible improvements

Performance of Range scans are bad in comparison with B+ tree.

Write Buffer index improves the write throughput and lowers write amplifications

Discussion