CONSIDERATIONS FOR INTEGRATING VIRTUAL THREADS IN A JAVA FRAMEWORK: A QUARKUS EXAMPLE IN A RESOURCE-CONSTRAINED ENVIRONMENT

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- 1 Context
- 2 Integration
- 3 The experiment
- 4 Experimental protoco
- 5 Results

THE END OF THE MONOLITH

- · difficult to maintain
- · dependency hell
- · reboot for every change

- · sub-optimal deployment
- · limit scalability
- technology lock-in

from

Microservices: yesterday, today, and tomorrow - Dragoni & al

MICROSERVICES

- difficult to maintain
- split in smaller entities
- · dependency hell
- · less dependencies per service
- reboot for every change
- · reboot only impacted service

- sub-optimal deployment
- · per-service deployment
- limits scalability
- · scale each service
- technology lock-in
- per service technology

Remark

Transitioning to cloud becomes a thing

MAKE MICROSERVICES CHEAP

resource-efficiency ce-efficiency

resource-efficiency resource-efficiency

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SCALING AND REACTIVE PROGRAMMING

Problem

communications over the network are unreliable and slow

Solutions

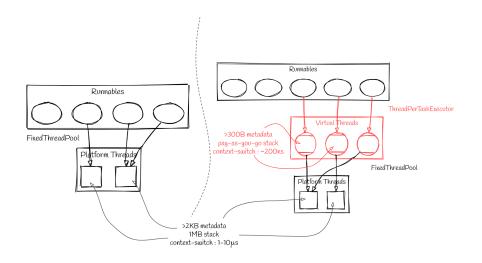
efficient and timely communication between services historically brought by non-blocking, asynchronous abstractions in the OS

Non-blocking abstractions

- asynchronous callbacks
- · promises and futures
- · reactive streams

- coroutines
- · async functions
- light threads

FINALLY, VIRTUAL THREADS



CODE COMPARISON - BLOCKING

```
var names = getAll();
var quotes = getQuotes(names.size());
for(int i=0; i < names.size();i ++){
   names.get(i).surname+= "-"+quotes.get(i);
}
return names;</pre>
```

CODE COMPARISON - ASYNCHRONOUS CALLBACKS

```
getAll( names => {
    getQuotes(names.size(), quotes => {
        for(int i=0; i < names.size();i ++){
            names.get(i).surname+= "- "+quotes.get(i);
        }
        //continuation
    })
});</pre>
```

CODE COMPARISON - REACTIVE STREAMS

```
var names = getAll().memoize().indefinitely();
var quotes = names.onItem().transformToUni(list ->
         getQuotes(list.size()));
return Uni.combine().all()
   .unis(names,quotes).asTuple()
   .onItem().transform(tuple -> {
      var nList=tuple.getItem1();
      //can await it since it is already resolved
      var qList = tuple.getItem2();
      for(int i=0; i < namesList.size();i ++){</pre>
         nList.get(i).surname += " - "+qList.get(i);
      return namesList;
   });
```

CODE COMPARISON - VIRTUAL THREADS

```
var names = getAll();
var quotes = getQuotes(names.size());
for(int i=0; i < names.size();i ++){
   names.get(i).surname+= "-"+quotes.get(i);
}
return names;</pre>
```

ASYNCHRONOUS PROGRAMMING ABSTRACTIONS

We should do (as wise programmers aware of our limitations) our utmost best to shorten the conceptual gap between the static program and the dynamic process, to make the correspondence between the program (spread out in text space) and the process (spread out in time) as trivial as possible.

Edgar. J. Dijkstra, 1968

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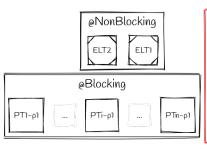
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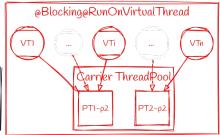
THREE INTEGRATION STRATEGIES

Strategy	Pros	Cons
Forking worker model	Simple, fits virtual threads model	Context switches
Using event-loop as carrier	No context-switch, Fewer threads overall	Potential deadlocks
Modifying Netty event- loops to be virtual threads	Integration done at the Netty level, Netty- based frameworks would benefit from it	stream, unpredictable

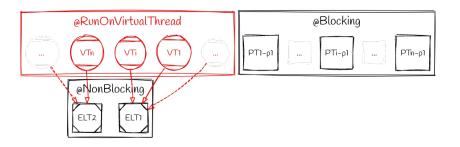
Table: Comparison of the different Quarkus-virtual-threads options

FORKING THE WORKER MODEL

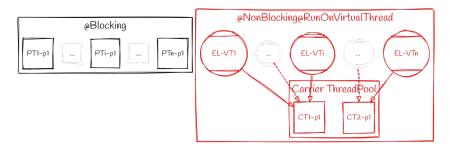




REUSING EVENT-LOOPS



"VIRTUALIZING" NETTY EVENT-LOOPS



DEADLOCK: ROLES OF THE EVENT-LOOP

- 1. event-loop for the entire non-blocking inner-workings of the framework,
- 2. event-loop for the non-blocking endpoint handlers,
- 3. work as a carrier thread for virtual threads

DEADLOCK SITUATION

SCHEMA

DEADLOCK: SHARING LOCKS

Conclusion

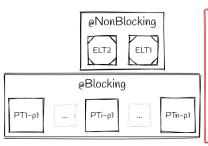
The event-loop can't reuse locks as a carrier

INTEGRATION STRATEGIES SUMMARY

Strategy	Pros	Cons
Forking worker model	Simple, fits virtual threads model	Context switches
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Modifying Netty event- loops to be virtual threads	Integration done at the Netty level, Netty- based frameworks would benefit from it	stream, unpredictable

Table: Comparison of the different Quarkus-virtual-threads options

FINAL CHOICE: FORKING THE WORKER MODEL



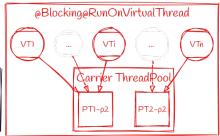


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THE EXPERIMENT

Goal

Measure how performance of the application is affected by replacing reactive endpoints with virtual-threads offloading, in conditions close to

Hypothesis

Quarkus-virtual-threads should perform better than **Quarkus-blocking** but not as well as **Quarkus-reactive**

THE ENVIRONMENT

Limited resouces

- 512MB memory
- 0.5 vCPU
- 256MB heap

limited resouces

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ELEMENTS OF THE EXPERIMENT

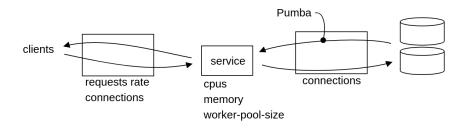
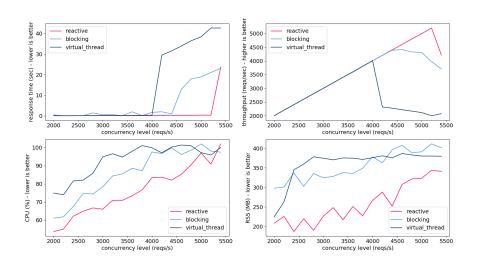


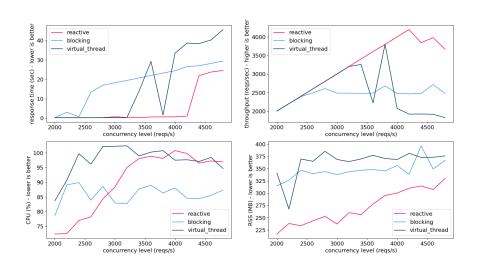
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RESULTS WITHOUT DELAY



RESULTS WITH 200MS DELAY



CONCLUSIONS

performance	Quarkus-blocking < Quarkus-virtual-threads < 0
resource efficiency	X