

Instruction Graph Statics

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1 Validity

$$\frac{}{\emptyset \text{ valid}} \quad \frac{(\mathbf{vs}, \mathbf{U}) \text{ defined} \quad (\mathbf{vs}, \emptyset, \mathbf{s}, \mathbf{U}) \text{ connected}}{\mathbf{P}(\mathbf{vs}, \mathbf{s}) \text{ valid}}$$

2 Defined

$$\frac{}{(\mathbf{S}(\mathbf{V}(\mathbf{n}, \text{cont})), \{\mathbf{n}\}) \text{ defined}} \quad \frac{(\mathbf{vs}, \mathbf{U}) \text{ defined} \quad \mathbf{n} \notin \mathbf{U}}{(\mathbf{Cons}(\mathbf{V}(\mathbf{n}, \text{cont}), \mathbf{vs}), \mathbf{U} \cup \{\mathbf{n}\}) \text{ defined}}$$

3 Match

$$\frac{}{(\mathbf{S}(\mathbf{V}(\mathbf{n}, \text{cont})), \mathbf{n}, \mathbf{V}(\mathbf{n}, \text{cont})) \text{ match}}$$
$$\frac{(\mathbf{vs}, \mathbf{n}, \mathbf{V}(\mathbf{n}, \text{cont})) \text{ match}}{(\mathbf{Cons}(\mathbf{v}, \mathbf{vs}), \mathbf{n}, \mathbf{V}(\mathbf{n}, \text{cont})) \text{ match}}$$
$$\frac{}{(\mathbf{Cons}(\mathbf{V}(\mathbf{n}, \text{cont}), \mathbf{vs}), \mathbf{n}, \mathbf{V}(\mathbf{n}, \text{cont})) \text{ match}}$$

4 Connected

$$\frac{n \in U_v}{(vs, U_v, n, \emptyset) \text{ connected}} \quad \frac{(vs, n, V(n, \mathbf{End})) \text{ match} \quad n \notin U_v}{(vs, U_v, n, \{n\}) \text{ connected}}$$

$$\frac{\begin{array}{l} (vs, n, V(n, \mathbf{Do}(\mathbf{act}, n'))) \text{ match} \\ (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v \end{array}}{(vs, U_v, n, U \cup \{n\}) \text{ connected}}$$

$$\frac{\begin{array}{l} (vs, n, V(n, \mathbf{DoU}(\mathbf{act}, \mathbf{cond}, n'))) \text{ match} \\ (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v \end{array}}{(vs, U_v, n, U \cup \{n\}) \text{ connected}}$$

$$\frac{\begin{array}{l} (vs, n, V(n, \mathbf{Cond}(\mathbf{cond}, n', n''))) \text{ match} \\ (vs, U_v \cup \{n\}, n', U) \text{ connected} \\ (vs, U_v \cup U \cup \{n\}, n'', U') \text{ connected} \quad n \notin U_v \end{array}}{(vs, U_v, n, U \cup U' \cup \{n\}) \text{ connected}}$$

$$\frac{\begin{array}{l} (vs, n, V(n, \mathbf{GoTo}(n'))) \text{ match} \\ (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v \end{array}}{(vs, U_v, n, U \cup \{n\}) \text{ connected}}$$