

Instruction Graph Statics

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1 Validity

p **valid** means that the Program p is a valid program.

$$\frac{(\mathbf{V}(s, c) :: vs, U) \text{ defined} \quad (\mathbf{V}(s, c) :: vs, \emptyset, s, U) \text{ connected}}{\mathbf{P}(\mathbf{V}(s, c), vs) \text{ valid}}$$

2 Defined

We let $U \subseteq \mathbb{Z}$ be a subset of the integers.

(vs, U) **defined** means that the Vertices vs define exactly the set U of vertex indices.

$$\frac{}{(\text{nil}, \{ \}) \text{ defined}} \qquad \frac{(\text{vs}, U) \text{ defined} \quad n \notin U}{(\mathbf{V}(n, c) :: vs, U \cup \{n\}) \text{ defined}}$$

3 Connected

We let $U \subseteq \mathbb{Z}$ be a subset of the integers.

(vs, U_v, n, U) **connected** means that the vertex represented by n is connected to each vertex represented by an index in U of the vertices in vs , where U_v is the set of vertex indices of vertices already visited.

$$\begin{array}{c}
\frac{n \in U_v}{(vs, U_v, n, \emptyset) \text{ connected}} \qquad \frac{\mathbf{V}(n, \text{end}) \in vs \quad n \notin U_v}{(vs, U_v, n, \{n\}) \text{ connected}} \\
\\
\frac{\mathbf{V}(n, \text{do } a \text{ then } n') \in vs \quad (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v}{(vs, U_v, n, U \cup \{n\}) \text{ connected}} \\
\\
\frac{\mathbf{V}(n, \text{do } a \text{ until } cnd \text{ then } n') \in vs \quad (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v}{(vs, U_v, n, U \cup \{n\}) \text{ connected}} \\
\\
\frac{\mathbf{V}(n, \text{if } cnd \text{ then } n' \text{ else } n'') \in vs \quad (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad (vs, U_v \cup U \cup \{n\}, n'', U') \text{ connected} \quad n \notin U_v}{(vs, U_v, n, U' \cup \{n\}) \text{ connected}} \\
\\
\frac{\mathbf{V}(n, \text{goto } n') \in vs \quad (vs, U_v \cup \{n\}, n', U) \text{ connected} \quad n \notin U_v}{(vs, U_v, n, U \cup \{n\}) \text{ connected}}
\end{array}$$