

Asn 5

3.1)

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4) procedure largestDifference( $a_1 - a_2, a_2 - a_1, \dots, a_{n-1} - a_n$ )
    largestDifference := 0
    for  $i := 1$  to  $(n-1)$ 
        difference :=  $a_i - (a_{i+1})$ 
        if difference > largestDifference
            largestDifference := difference
    {largestDifference has desired value}
    
```

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10) procedure  $x^n$  ( $x \cdot x$  ( $n$  times))
    procedure  $x^{-n}$  ( $\frac{1}{x} \cdot \frac{1}{x}$  ( $n$  times))
        answer := 1
        if  $n := +$  (positive int)
            for  $i := 1$  to  $n$ 
                answer := answer  $\cdot x$ 
            else if  $n := -$  (negative int)
                for  $i := 1$  to  $(-1 \cdot n)$ 
                    answer := answer  $\cdot \frac{1}{x}$ 
            else
                answer := 1
    {answer has desired value}
    
```

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16) procedure smallest( $a_1, a_2, a_3, a_4, \dots, a_n$ )
    smallest :=  $a_1$ 
    for  $i := 2$  to  $n$ 
        if smallest >  $a_i$ 
            smallest :=  $a_i$ 
    {smallest has desired value}
    
```

28) procedure quadsearch(x : integer, a_1, a_2, \dots, a_n : increasing ints)

$i := 1$

$j := n$

while $i < j$

$m := \lfloor (i+j)/4 \rfloor$

if $x > a_m$

$i := m+1$

else

$j := m$

if $x = a_i$

location := i

else

location := 0

{location is the subscript of the term equal to x , or 0 if x isn't found}

50) procedure insertion sort(a_1, a_2, \dots, a_n : real numbers with $n \geq 2$)

for $j := 2$ to n

$i := j-1$

while $a_j < a_i$

$i := j-1$

$m := a_j$

for $k := 0$ to $j-i-1$

$a_{j-k} = a_{j-k-1}$

$a_i = m$

{ a_1, a_2, \dots, a_n sorted}

$$c) 2+3+4+\dots, t_n =$$

$$\frac{n^2+n-2}{2}$$

$$d) n+n-1+\dots+2 =$$

$$\frac{n^2+n-2}{2}$$

b) 3, 2, 4, 5, 1, 6

$2 < 3$

2, 3, 4, 5, 1, 6

473, 472

2, 3, 4, 5, 1, 6

574, 573, 572

1, 2, 3, 4, 5, 6

$1 < 5, 1 < 4, 1 < 3, 1 < 2$

1, 2, 3, 4, 5, 6

675, 674, 673, 672, 671

56/ 20 coins:

greedy algorithm = 1: 2 cent, 1: nickel, 3: pennies
(5 coins)

fewest coins = 2: dimes
(2 coins)

So, Greedy Alg doesn't
produce fewest coins

1) for $i := 2$ to n (N) $n-1$ (since $i=2$)
if smallest $\geq a_i$ (N)
[$i = n+1$] \rightarrow to exit (1) $2(n-1) + 1 = \boxed{2n-1}$

2) for $i := 1$ to 4 (4) $4^3 + 2 = O(1)$
for $j := 1$ to 4 (4)
if $a_j > a_{j+1}$ (4)
swap a_j with a_{j+1} (+2 to exit)

4) Yes, because you are exponentially cutting
your search time for k .

8) a) $y=3, i=1, y=7, i=2, y=15$

b) n multiplications and n additions

16) $O(\log n)$