

**a. Write a pseudocode.**

```
function next_palindrome
    pass in: number
    number += 1 # 1
    digits = list of extracted digits # d or O(log n)

    # indices
    i = 0 # 1
    j = length of digits list - 1 # 1

    # while i is not equal to j
    while i < j: # d/2 or O(log n)
        if digits[j] > digits[i]: # 1
            digits[j - 1] += 1 # 1

            # carry propagation
            x = j - 1 # 1
            while digits[x] >= 10: # O(1) amortized
                digits[x] % 10
                x += 1
                digits[x] += 1

            last digit = first digit # 1
            i += 1 # 1
            j -= 1 # 1

    return concatenated digits # O(log n)
```

**b. Analyse the performance of the code given the structure of an input.****I. What is the best-case complexity of your program? What is the structure of the input that results in the best case?**

$O(d)$  or  $O(\log n)$ , where  $d$  is the number of digits in the number. The number of digits in a number can also be calculated using  $\log_{10}(\text{number}) + 1$ , hence the  $O(\log n)$  time complexity. Extraction of the digits has a time complexity of  $O(d) \rightarrow O(\log n)$ . Looping through half the digits (while  $i < j$ ) is  $O(d/2) \rightarrow O(\log n)$ . Constructing the integer back is also  $O(d) \rightarrow O(\log n)$ . Therefore, the best-case complexity of my program is  $O(\log n)$ .

**II. What is the worst case complexity of the program? What is the structure of the input that results in the worst case?**

$O(d)$  or  $O(\log n)$ . Extraction of the digits has a time complexity of  $O(d) \rightarrow O(\log n)$ . Two halves of a list are essentially compared and equalized (while  $i < j$ ) until the

middle of the list is reached. This prevents us from brute forcing a palindrome check from the input number until the next smallest palindrome. The case of a carry propagation is run for a total max of all digits  $O(\log n)$ , that is  $O(1)$  amortized throughout the outer loop. Constructing the integer back is also  $O(d) \rightarrow O(\log n)$ . Therefore, the worst-case complexity of my program is  $O(\log n)$ .

- c. Check the correctness and performance of your program by testing it on a set of inputs.

```
lab1_final.py X
17 # extract list of digits out of an integer
18 def extract_digits(n: int) -> list[int]:
19     return list(map(int, str(n)))
20
21 # concat list of digits back to an integer
22 def make_int(digits: list[int]) -> int:
23     return int("".join(map(str, digits)))
24
25 def next_palindrome(n: int) -> int:
26     n += 1
27     digits = extract_digits(n)
28
29     i = 0
30     j = len(digits) - 1
31     while i < j:
32         if digits[j] > digits[i]:
```

PROBLEMS OUTPUT DEBUG CONSOLE PORTS TERMINAL

\CMSC 142-Lab1\output\small\_output' and 'C:\Users\ASUS\Desktop\BSCS 3 res-and-Algorithms\2025\CMSC 142-Lab1\output\full\_output'.

PS C:\Users\ASUS\Desktop\BSCS 3 - 2nd Sem\CMSC 142\Data-Structures-a

python lab1checker.py

Run checker for: (1) Small Input / (2) Full Input ?

Choice: 1

Checking output for mistakes...

Number of mistakes: 0

Final result: 5/5 = 100.00% accuracy

Time elapsed: 0.0037450790405273438 seconds

```
PS C:\Users\ASUS\Desktop\BSCS 3 - 2nd Sem\CMSC 142\Da
python lab1checker.py
Run checker for: (1) Small Input / (2) Full Input ?
Choice: 2

Checking output for mistakes...

Number of mistakes: 0
Final result: 100001/100001 = 100.00% accuracy
Time elapsed: 0.7273902893066406 seconds
```