

Dynamic Programming

Minimum Edit Distance and Longest Increasing Subsequence

Introduction

- Scenario: spell checker handling misspelled words

Introduction

- When a spell checker encounters a possible misspelled word, it suggests possible words that are close to the misspelled word by looking at other words in its dictionary
- How does the spell checker determine which words are close to the misspelled word?

Introduction

- How do you measure how "close" one word is from another?

Alignments

- A natural measure of the distance between two strings is the extent to which they can be aligned or matched up
- An alignment is a way of writing the strings one above the other

Alignments

- Possible alignments of SNOWY and SUNNY

S - N O W Y
S U N N - Y

Cost: 3

- S N O W - Y
S U N - - N Y

Cost: 5

Edit distance

- Named as such because it can also be thought of as the minimum number of edits needed to transform first string into the second
- Also known as the sequence alignment problem

Edit distance

- The (minimum) edit distance between two strings is the cost of their best possible alignment
- Best possible alignment = an alignment that minimizes the cost
- aka Levenshtein distance, after Vladimir Levenshtein, who considered this distance in 1965

Edit distance

- **Input:** Two sequences $X[1..m]$ and $Y[1..n]$ (usually strings, but could also be numbers)
- **Output:** The edit distance and the best possible alignment of the two sequences

How can we solve the problem?



Use **Brute-force**
approach.

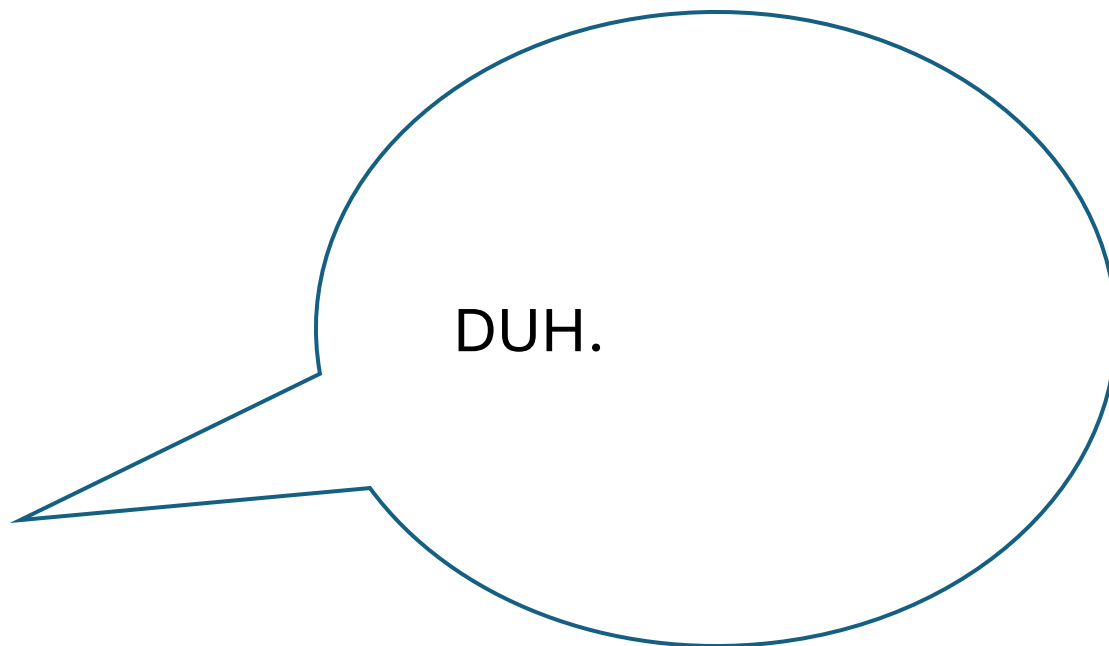
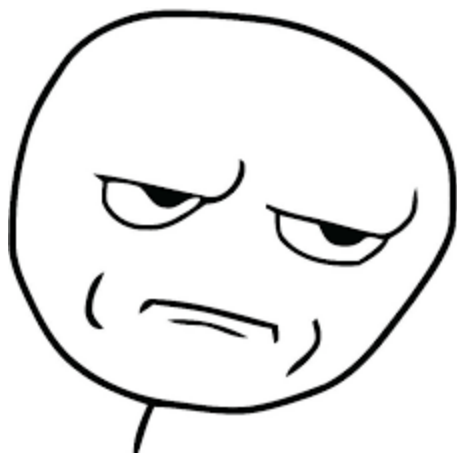
Brute Force Algorithm

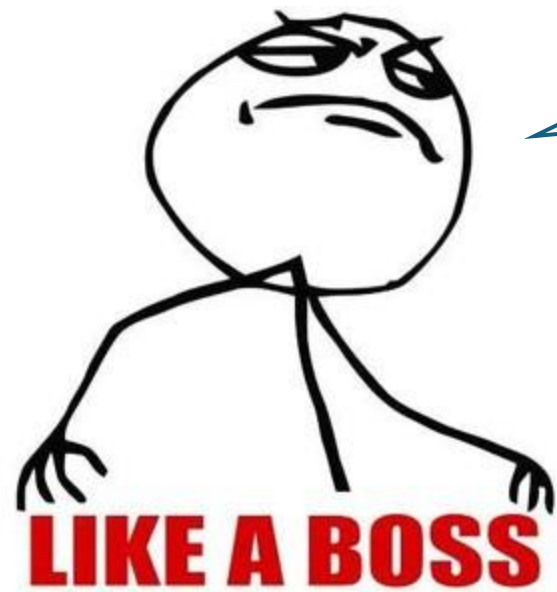
- Generate all possible alignments for the two sequences and compute their respective costs
- Return the alignment with the least cost

Inefficient

- Since there are so many possible alignments between two strings, it would be inefficient to search through all of them for the best one

Can we do better?





USE DP
APPROACH.

DP Solution

- First and most crucial question: What are the subproblems?
- We need to divide the problem into smaller subproblems such that it will have an optimal substructure

Subproblems

- We can look at some prefix of the first string $X[1..i]$ and some prefix of the second string $Y[1..j]$
- We'll call this subproblem **$E(i,j)$** : the edit distance of prefix $X[1..i]$ and prefix $Y[1..j]$

Defining $E(i, j)$

- We need to express $E(i, j)$ in terms of smaller subproblems
- What do we know about the best alignment of $X[1..i]$ and $Y[1..j]$?

Three cases

- We know that the rightmost column of the best alignment between $X[1..i]$ and $Y[1..j]$ can only be one of the three things:

X[i]

--

delete

--

Y[j]

insert

X[i]

Y[j]

substitution

3 operations

- Deletion
- Insertion
- Replacement
- Copy

Example

kitten  knitting

- Operation: insertion
- Knitten

Example

kitten \longrightarrow knitting

- Operation: Substitution
- Knittin

Example

kitten  knitting

- Operation: Insertion
- Knitting^g

Computing the solution

- The answers to all the subproblems form a $m \times n$ table
- What should be the order for solving these subproblems?

Two possible orderings:

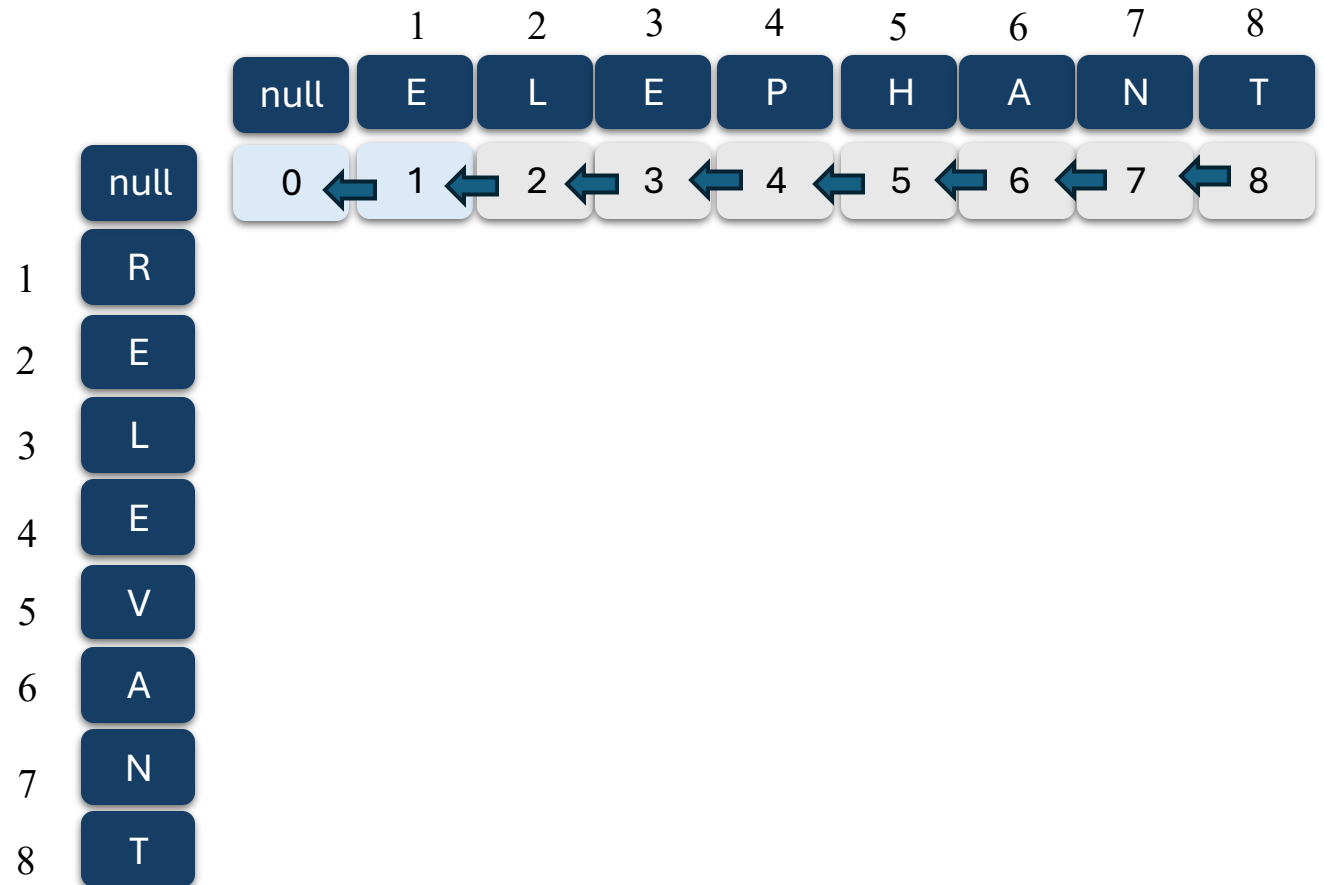
1. one row at a time moving from left to right;
2. one column at a time moving from top to bottom

Computing the solution

- What remains to be solved is the base case: the smallest subproblems
- In this case, our base cases are $E(0,*)$ and $E(*,0)$

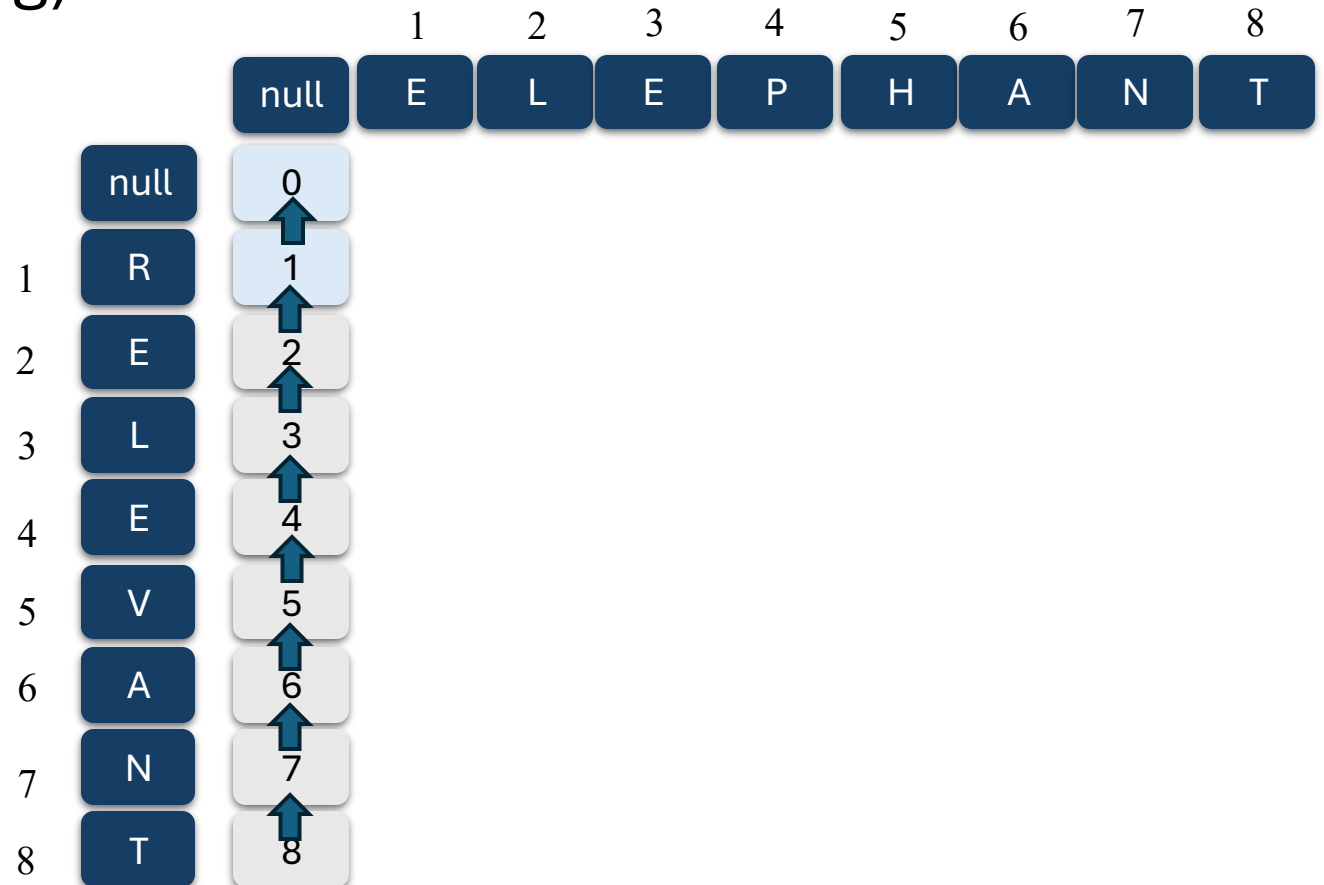
Computing the solution

- $E(0,j)$ is the edit distance between the 0-length prefix of X (empty string) and the first j letters of Y



Computing the solution

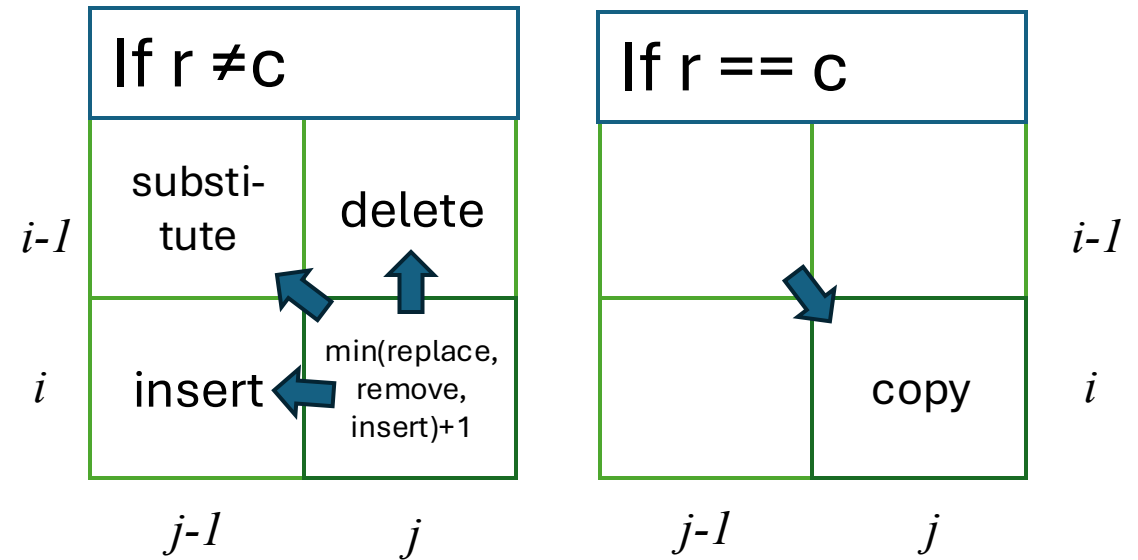
- $E(i,0)$ is the edit distance between the first i letters of X and the 0-length prefix of Y (empty string)



Defining $E(i, j)$

$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$

where $\text{diff}(i, j) = 0$ if $X[i] == Y[j]$, and 1 otherwise



Backtrace

- To construct the alignment of the X and Y, we need to keep a backtrace
- Idea is similar to constructing LCS

Backtrace

- Every time we enter a cell, remember where we came from
- Trace back path from the bottom-right corner and follow directions to read off the alignment

- Trace a path starting from the bottom-right corner (M,N) and follow the directions all the way back to
- the top-left corner (0,0); the directions along the way will tell you the alignment:

direction = \uparrow X
 deletion --

direction = \leftarrow --
 insertion Y

direction = \nwarrow X
 aligned Y

Backtrace

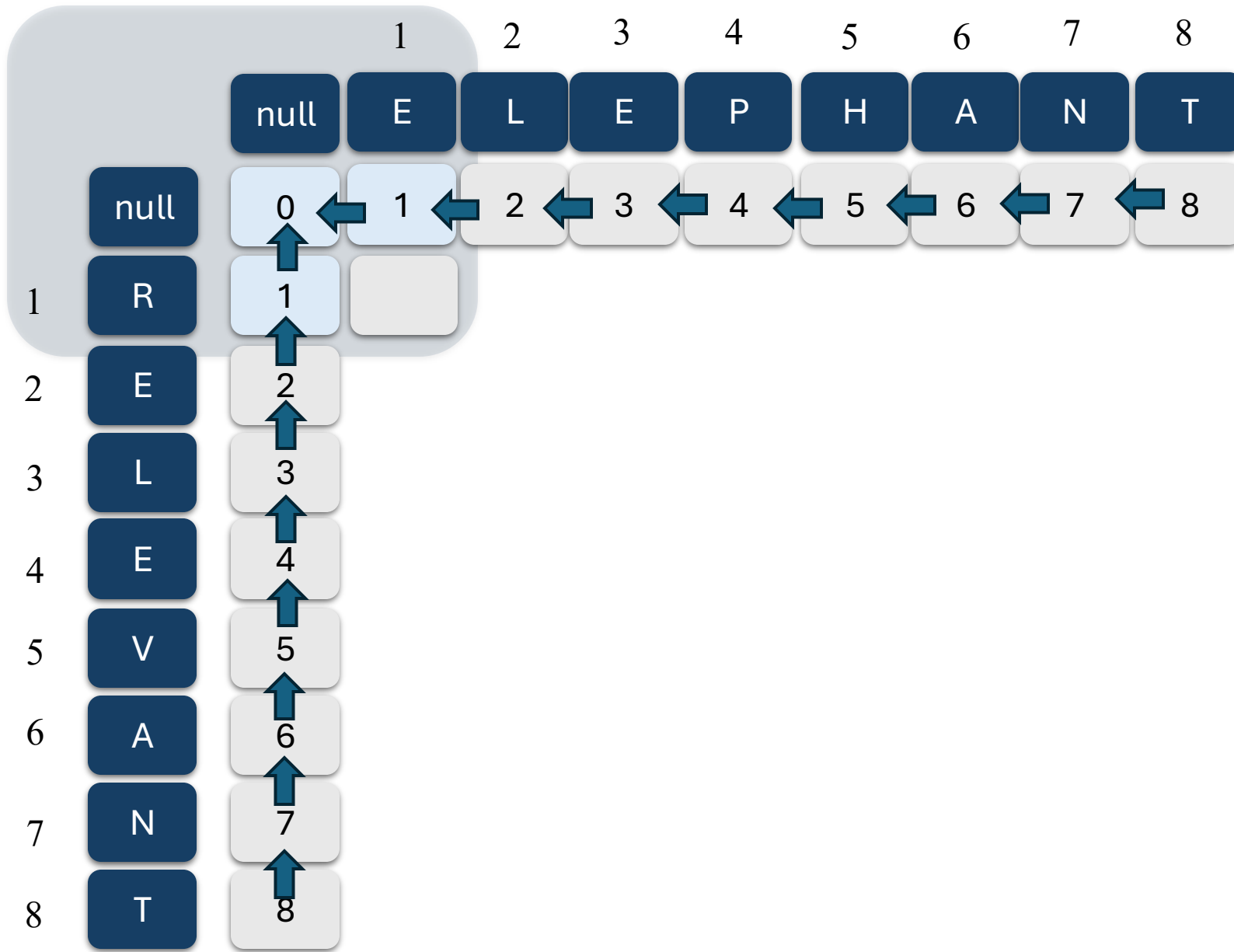
- Insertion: LEFT = $E(i, j-1)$
- Deletion: UP = $E(i-1, j)$
- Substitution: Diagonal = $E(i-1, j-1)$

Multiple Directions

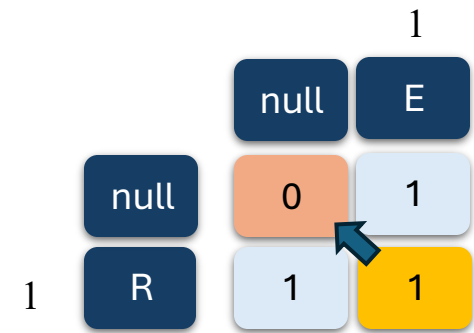
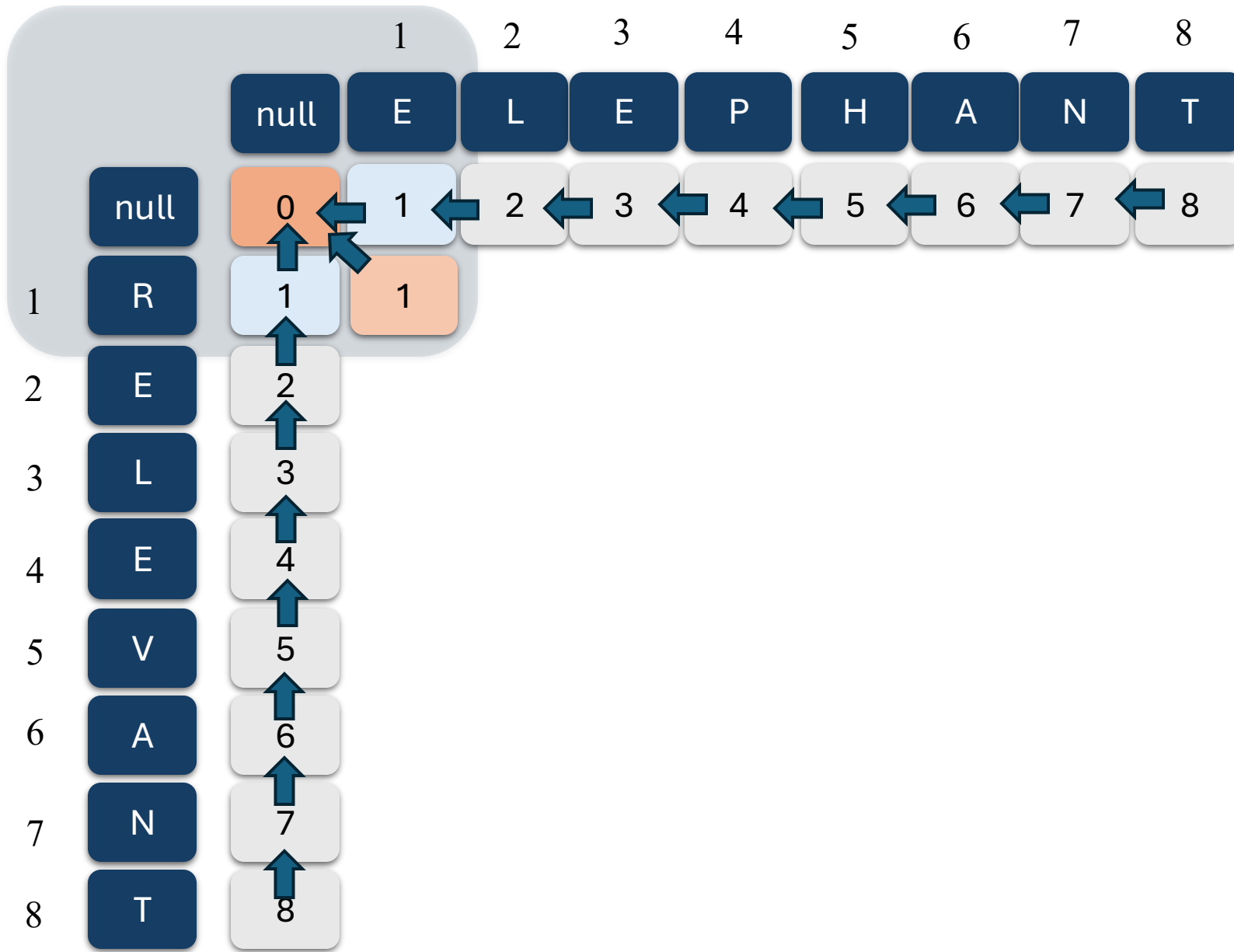
- In keeping backtrace directions, you can store more than 1 direction if there are ties
- Any path from bottom-right to the top-left is a sequence alignment

Computing the solution

- $E(0, j) = j$
- $E(l, 0) = i$

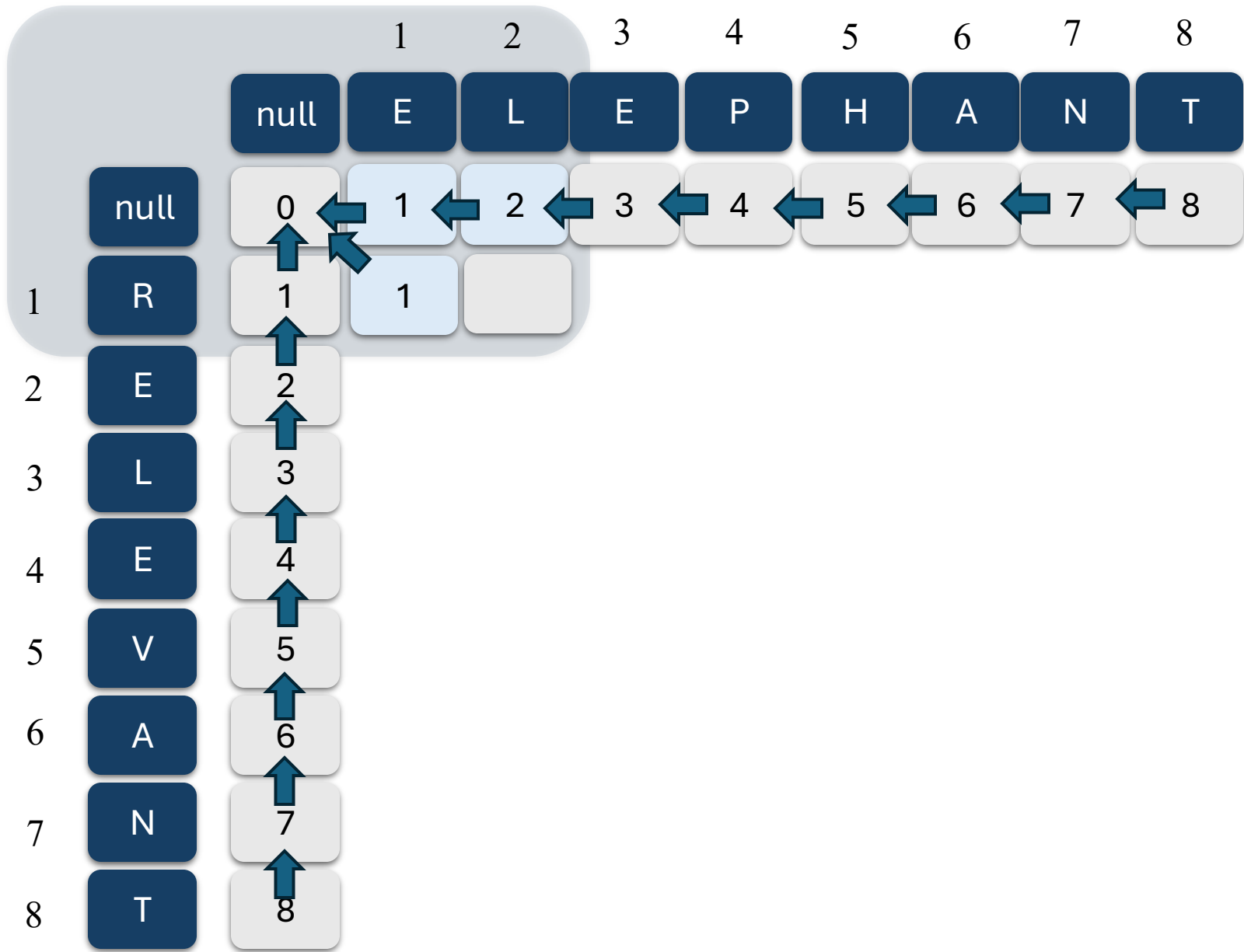


$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$

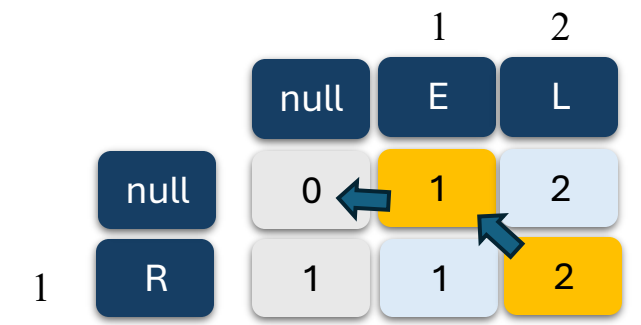
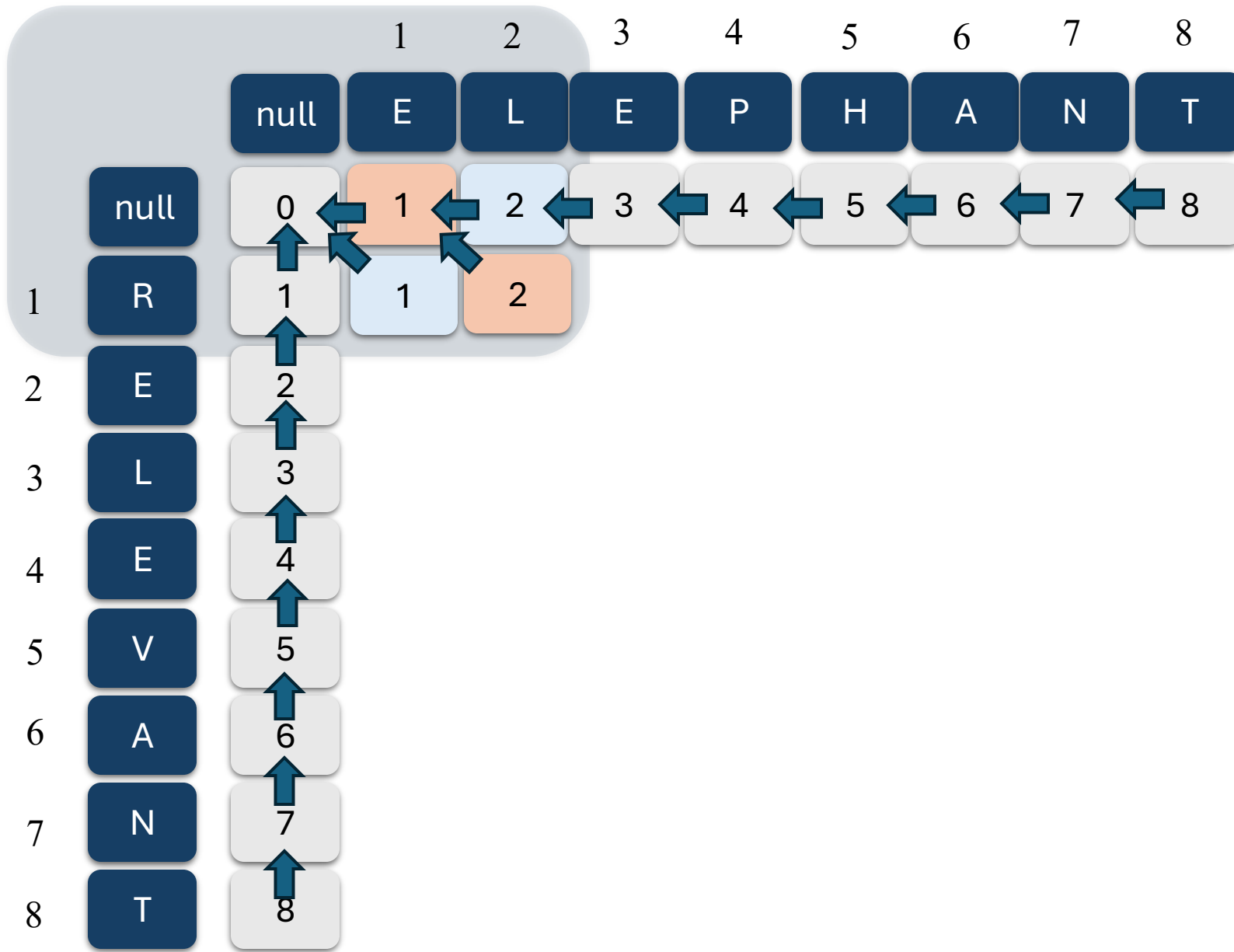


Replace R with E

$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$

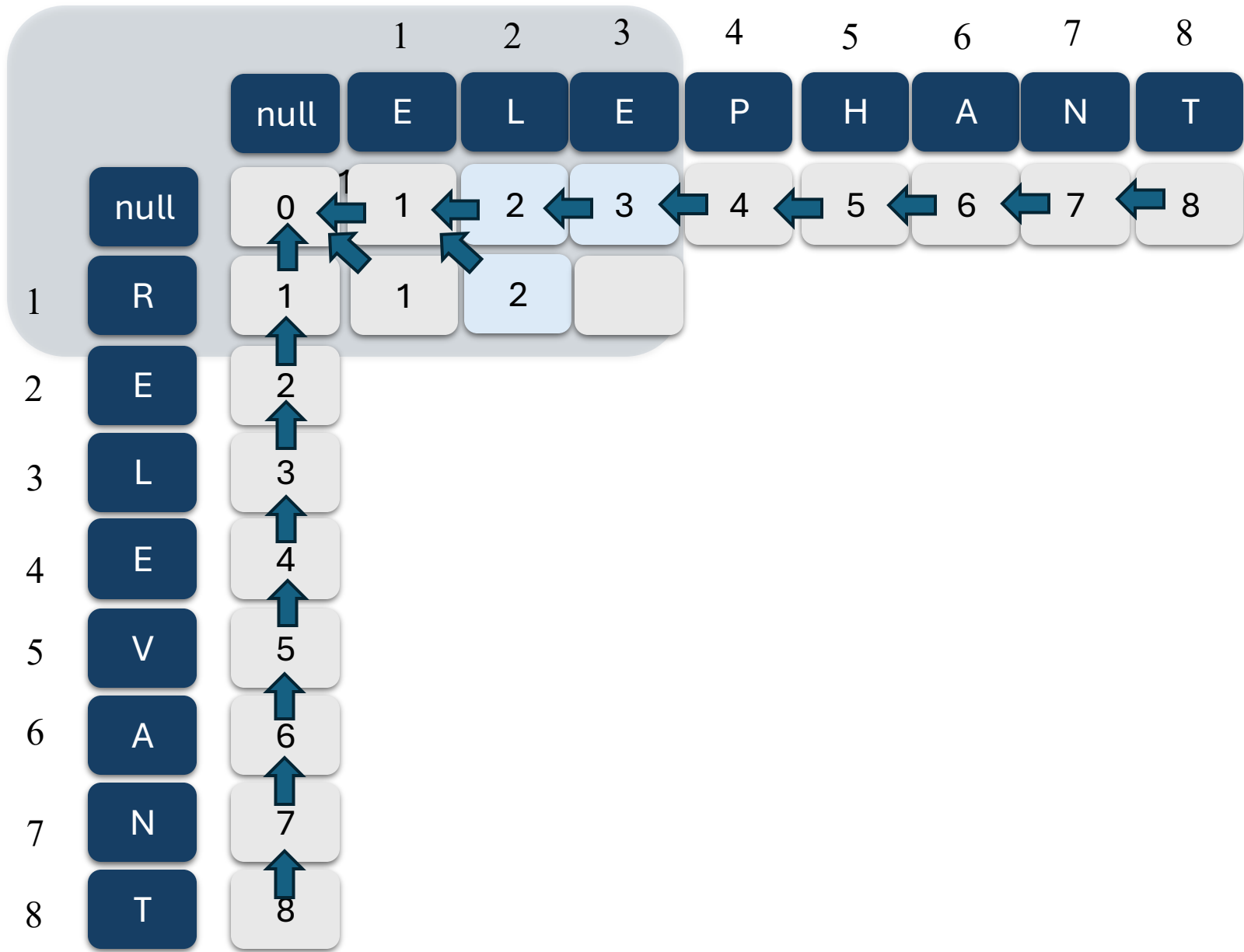


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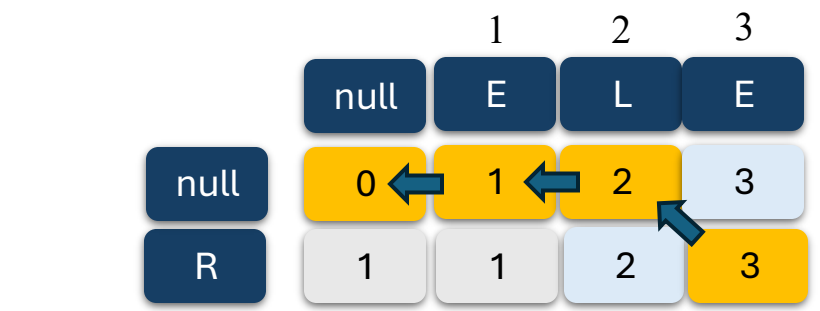
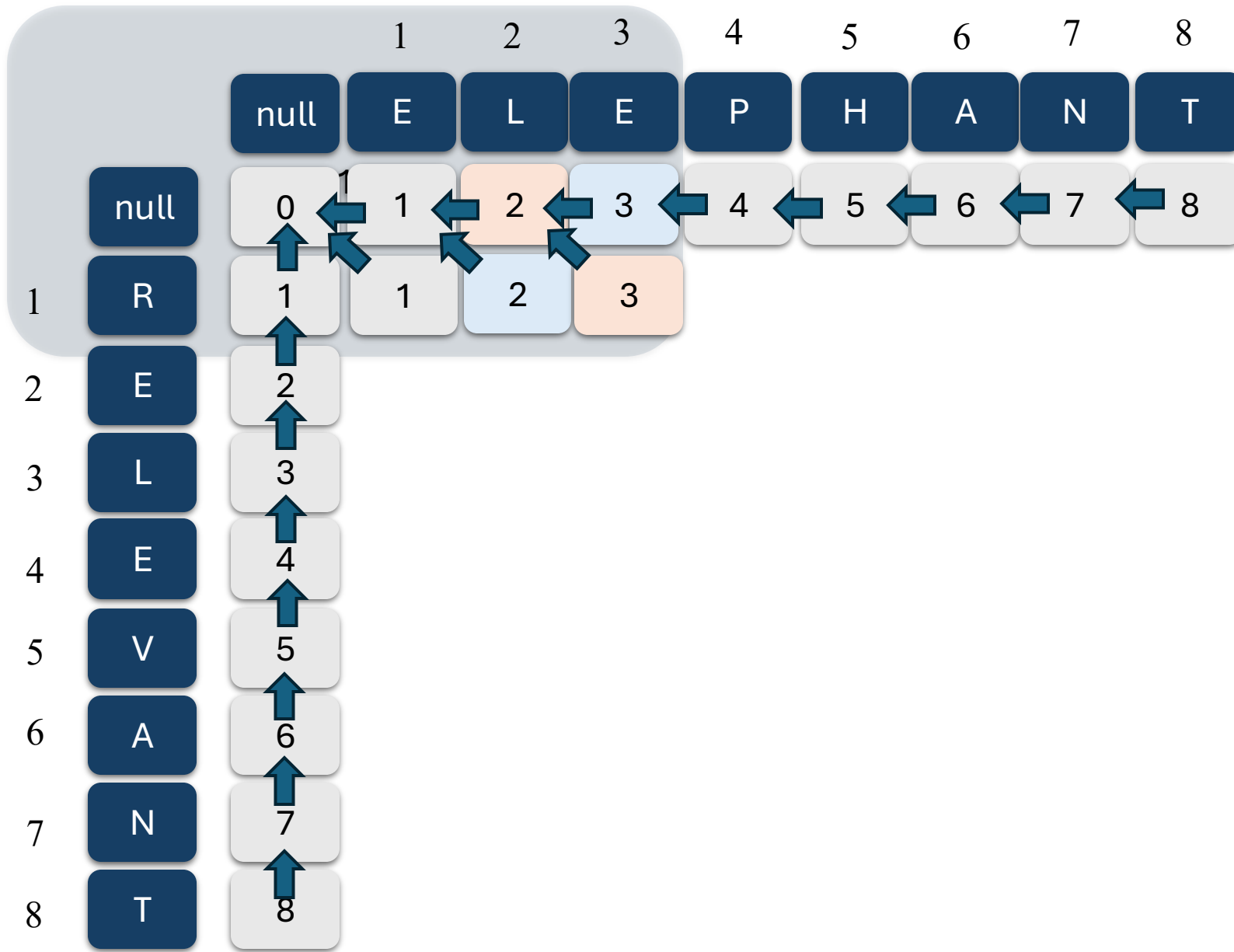


Insert E
Replace R with L

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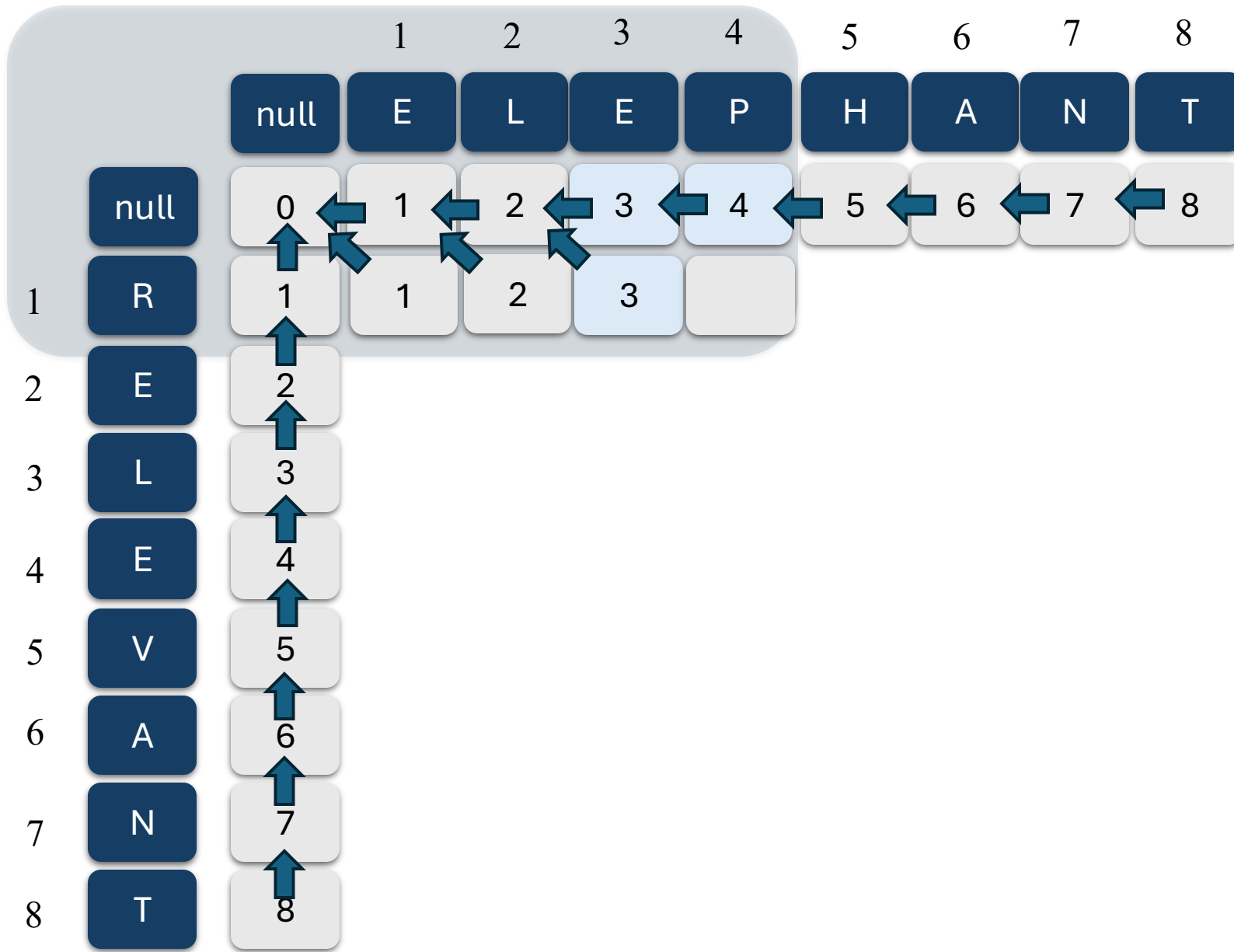


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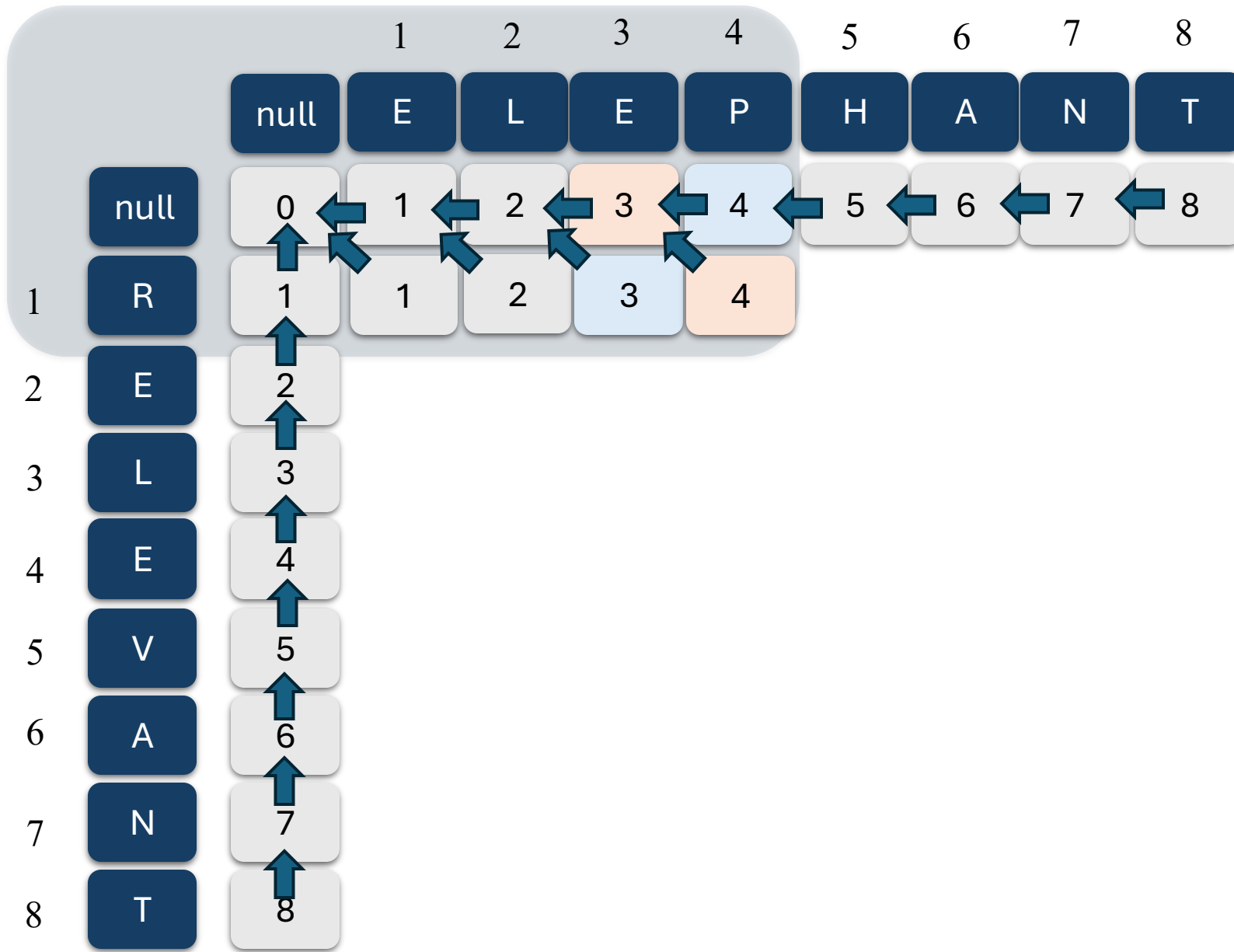


Insert E
 Insert L
 Replace R with E

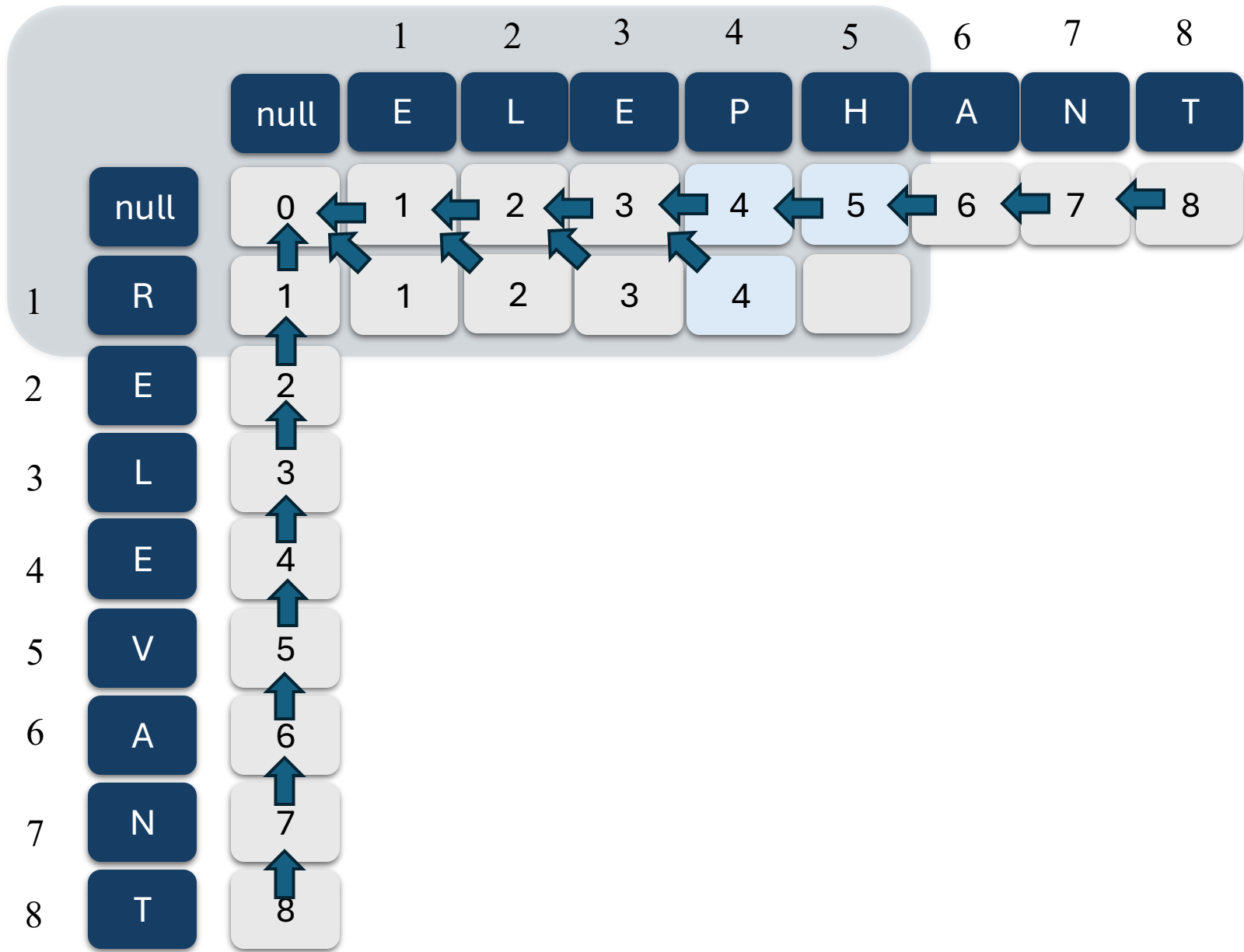
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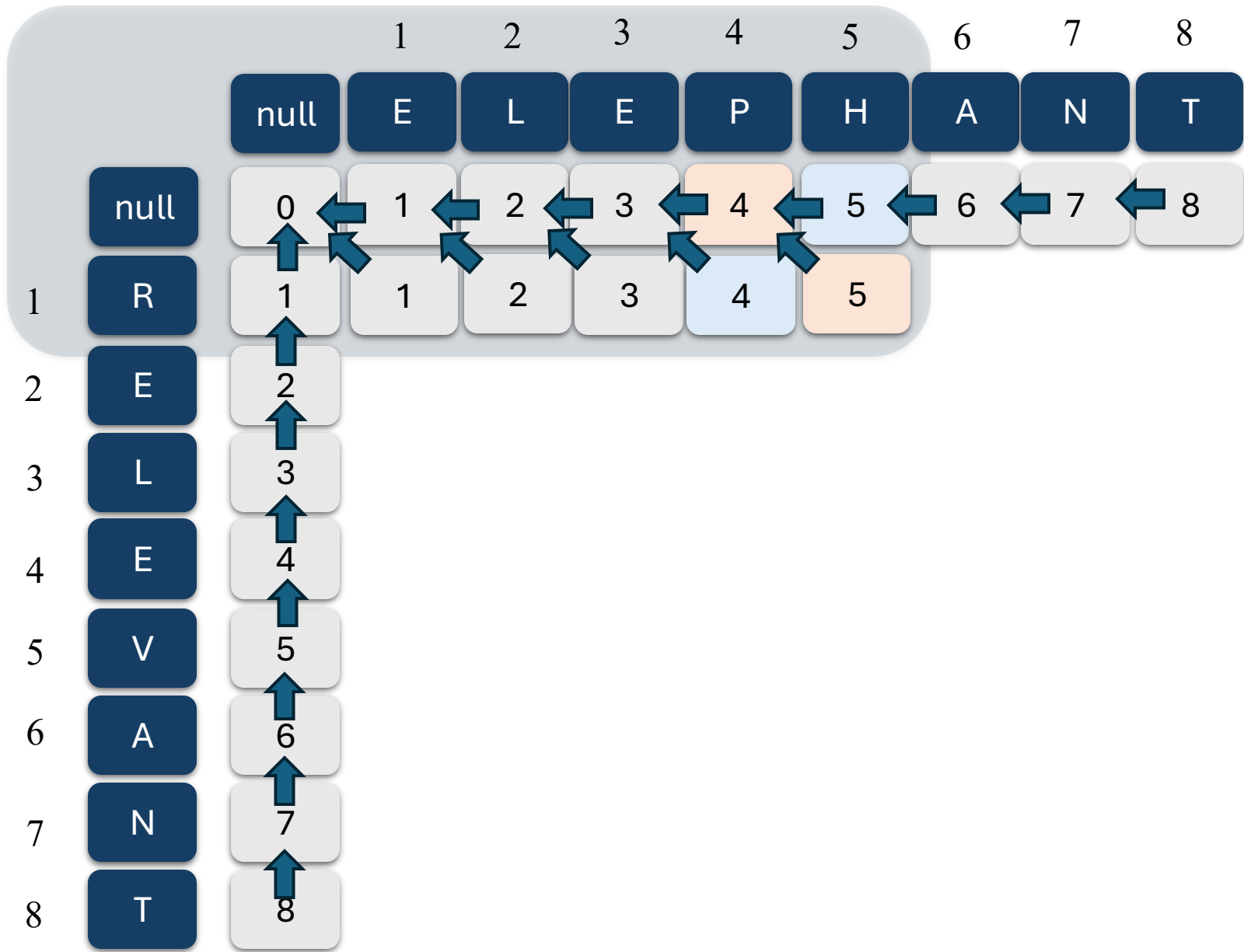
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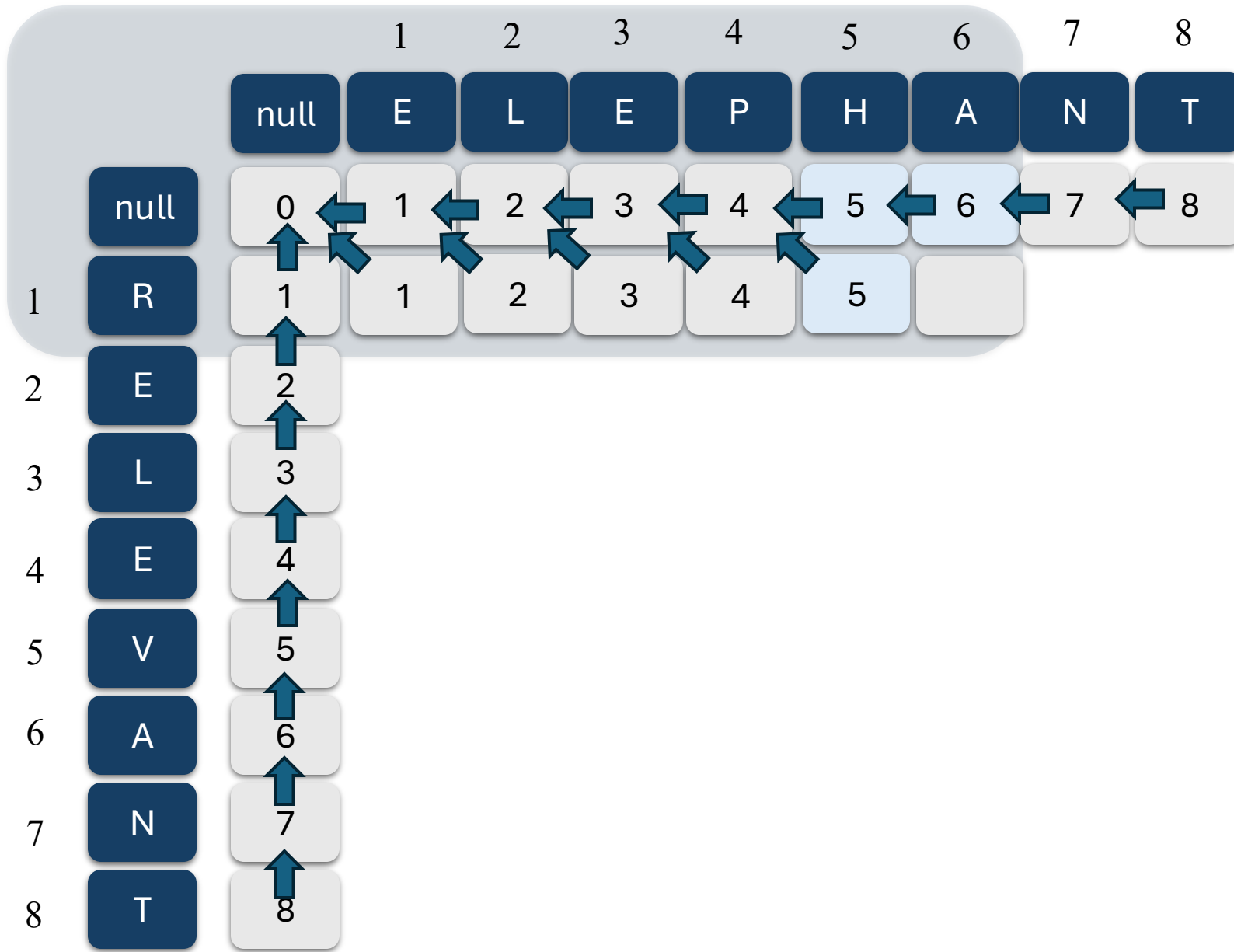
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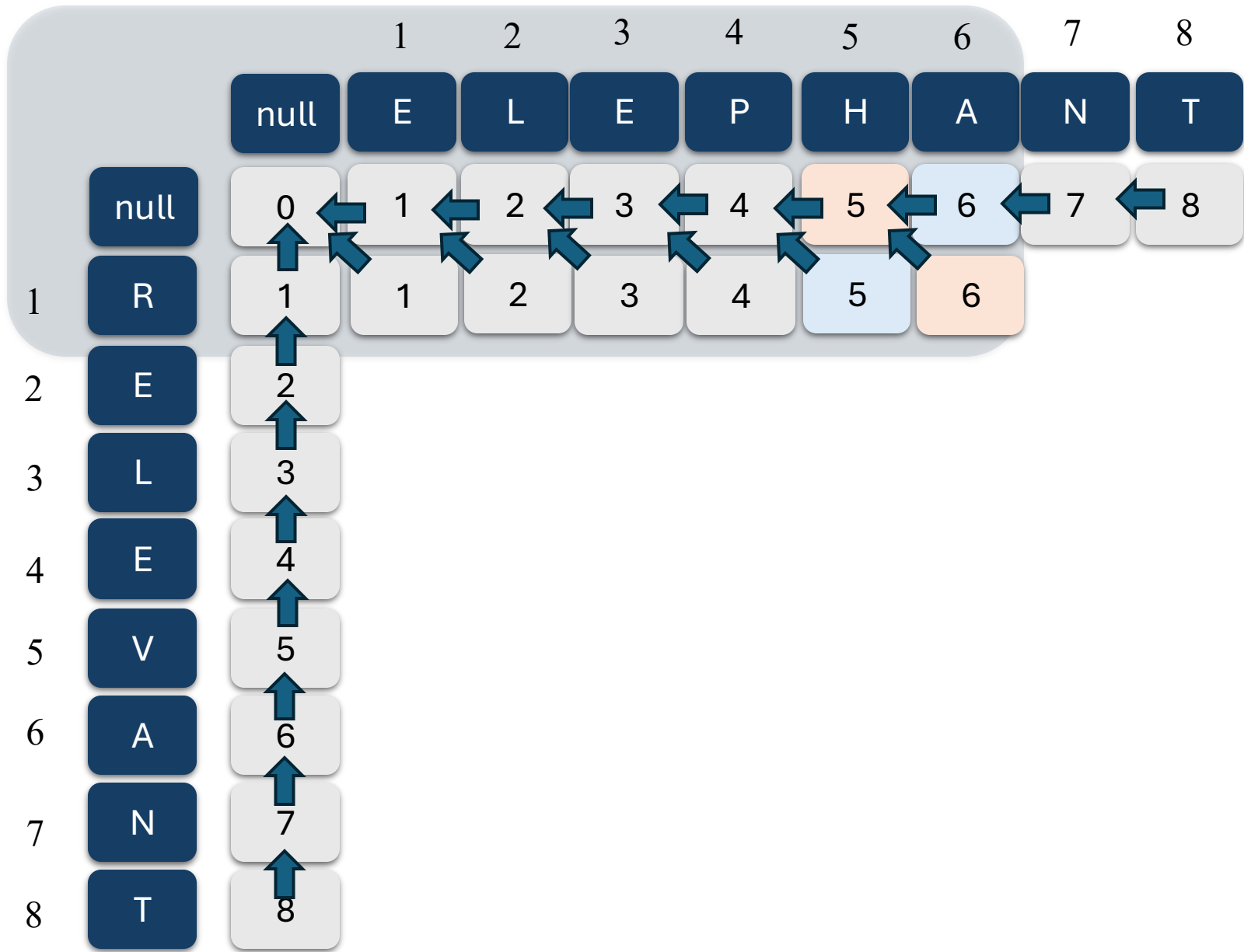
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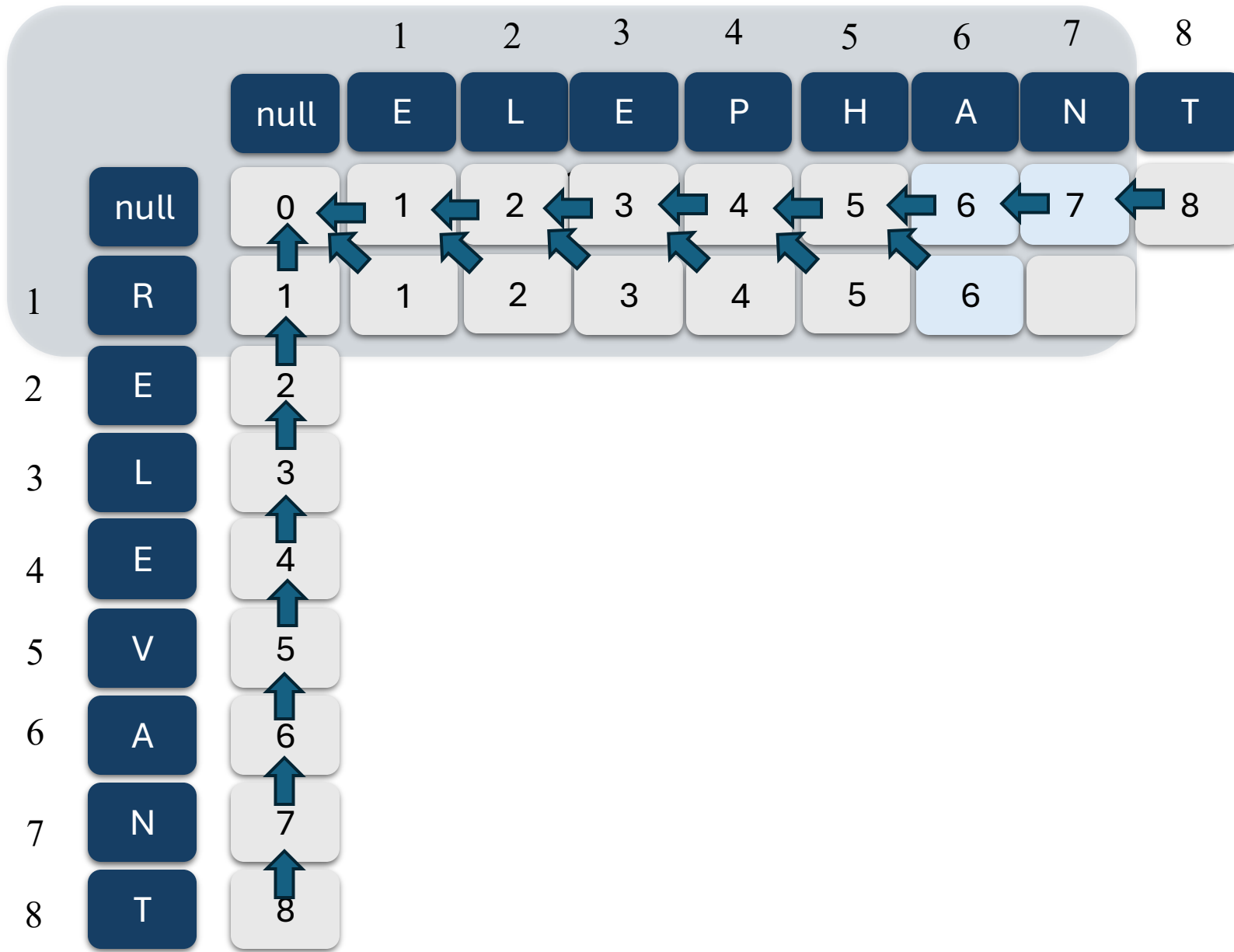
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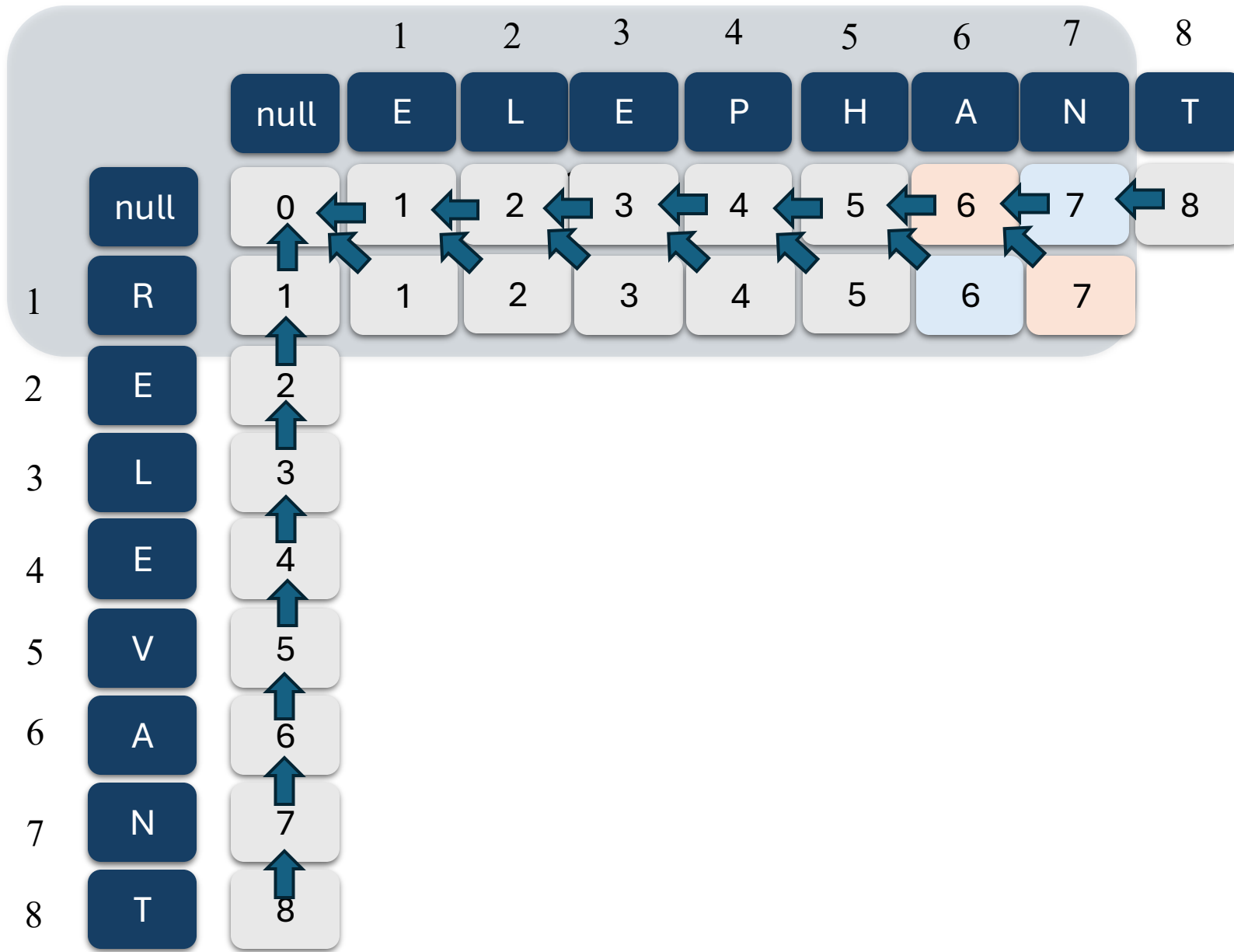
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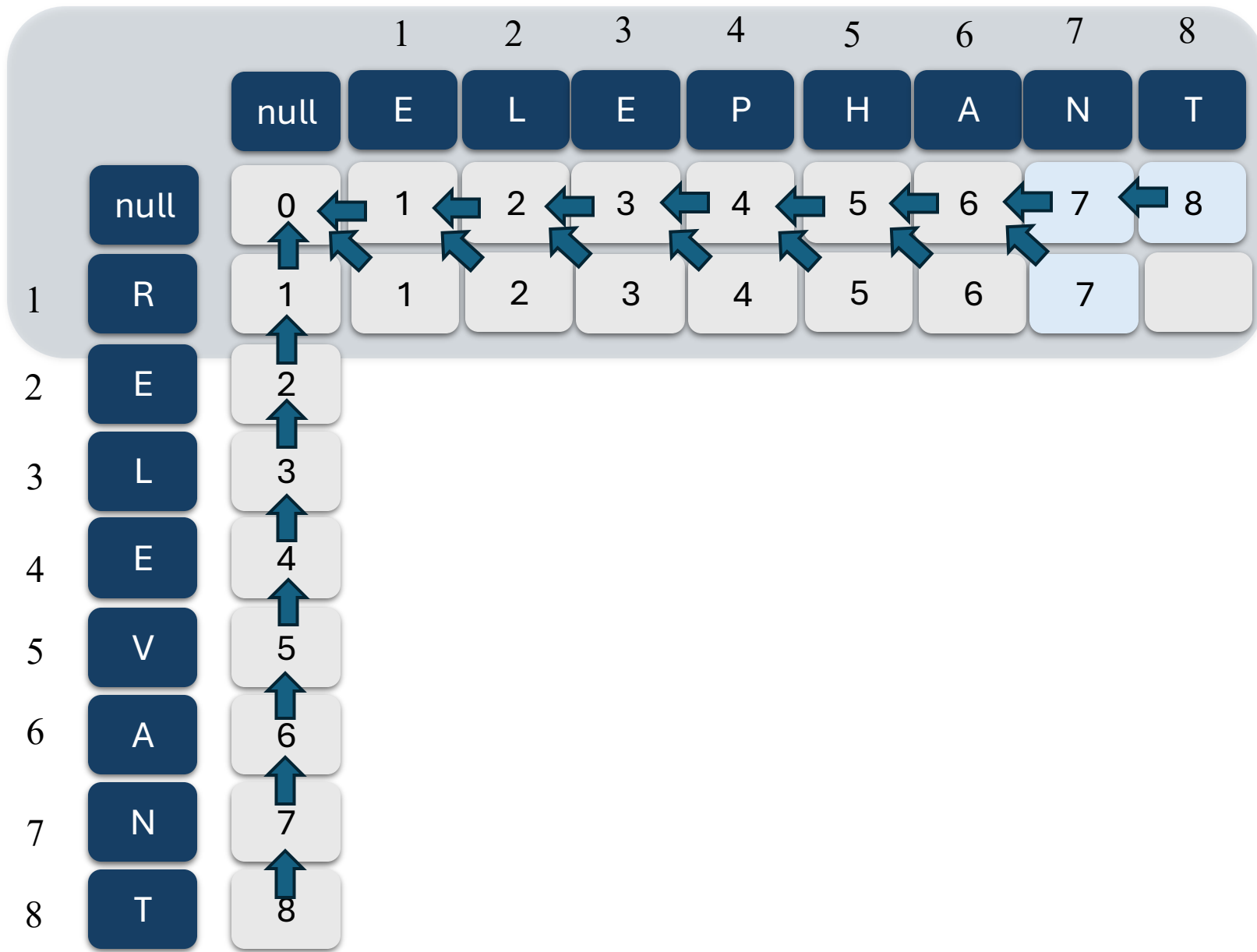
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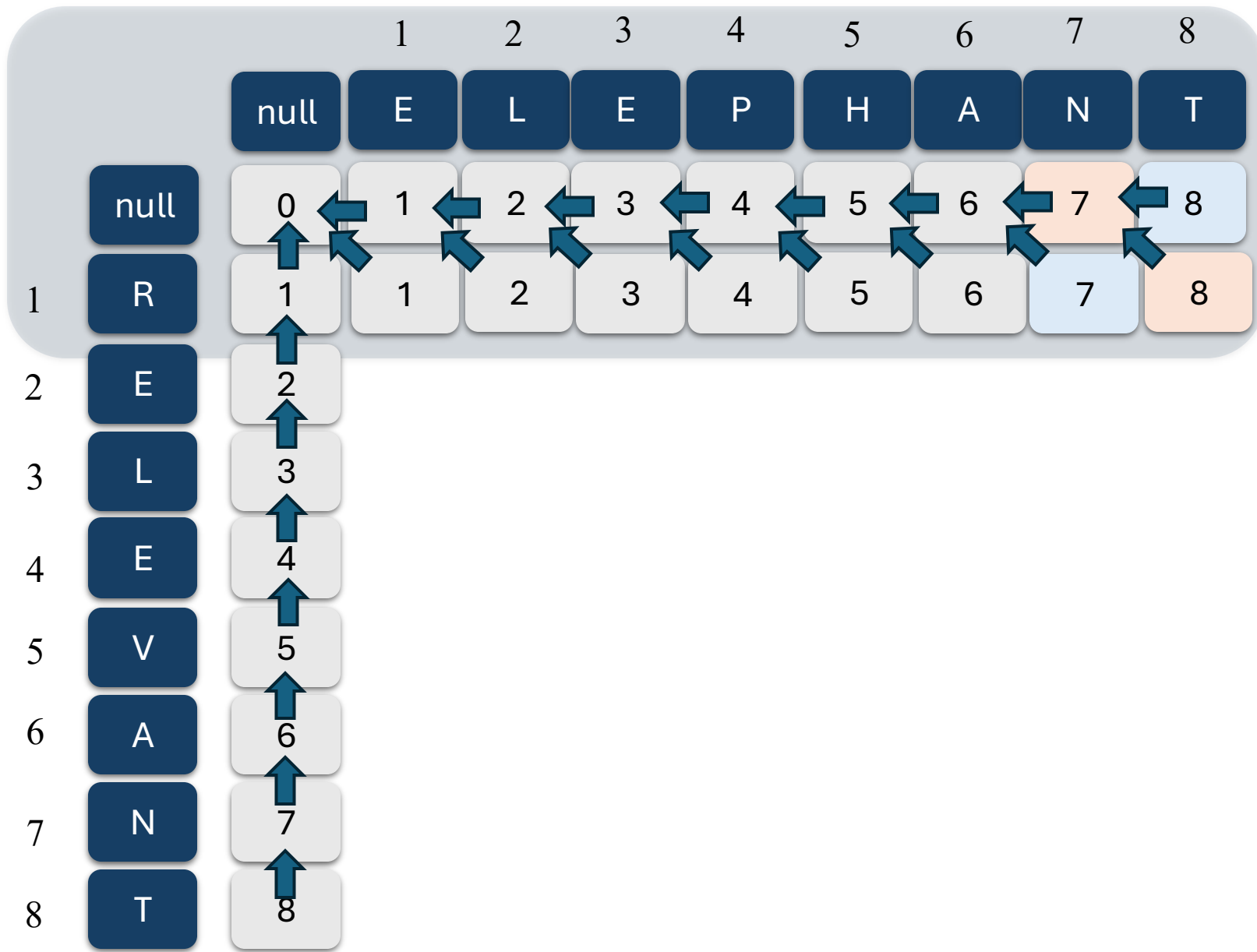
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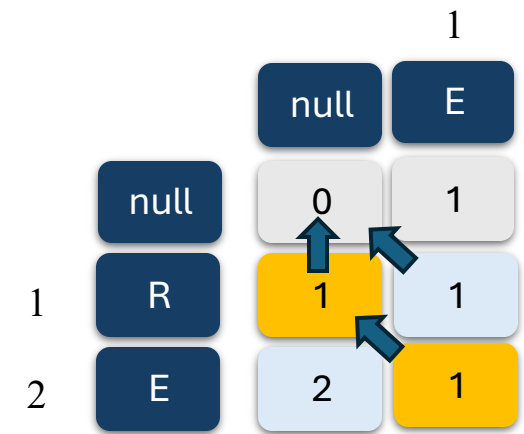
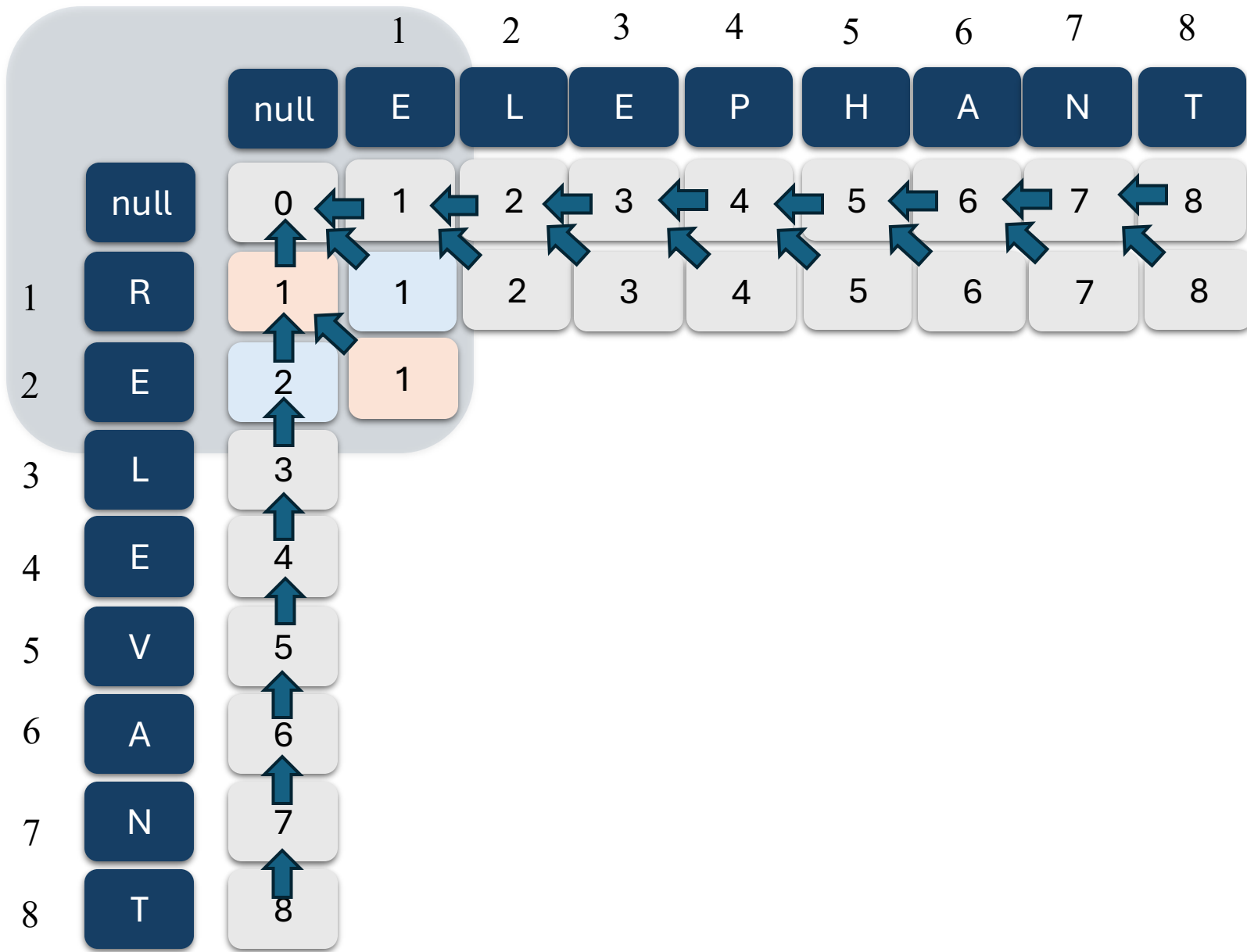
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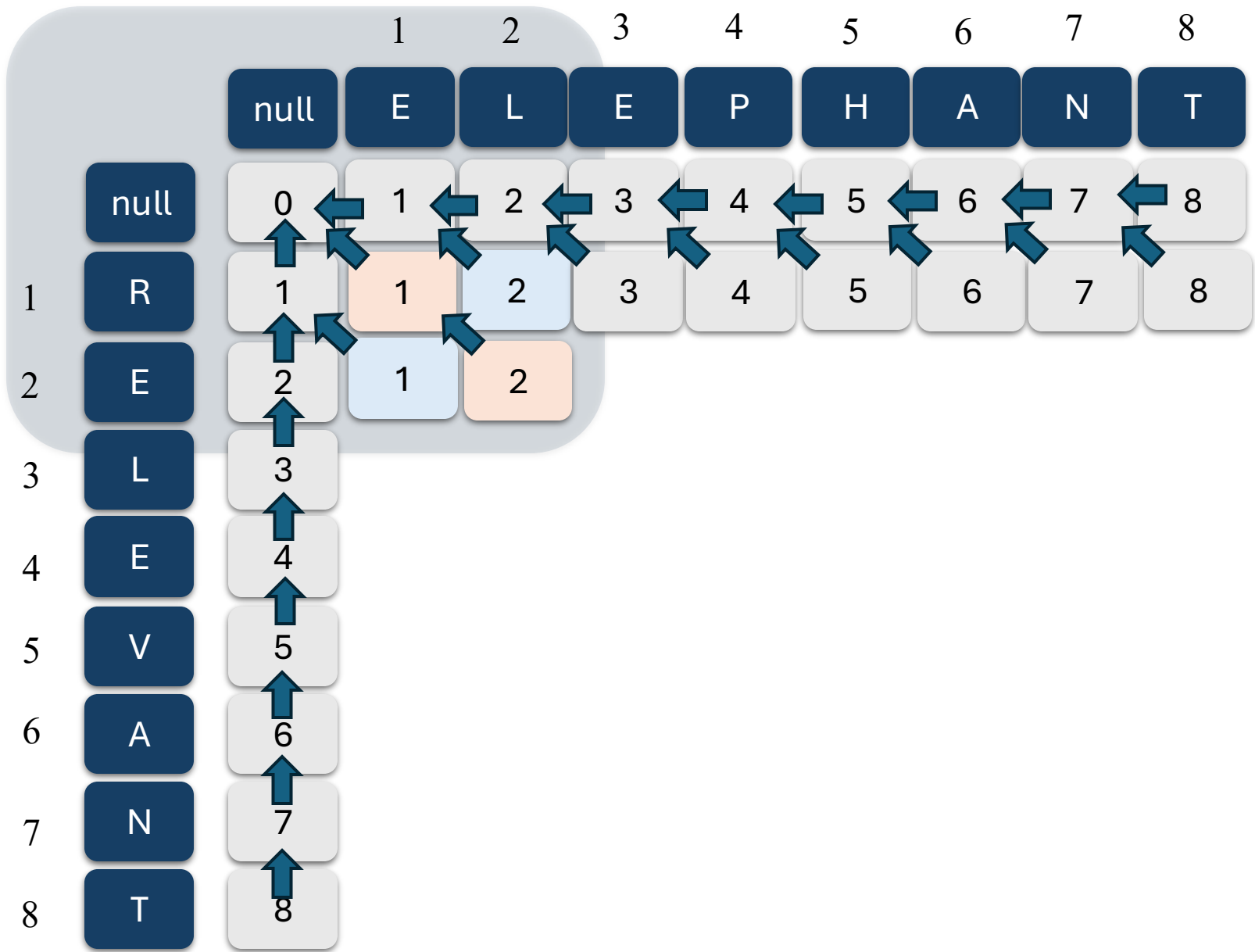


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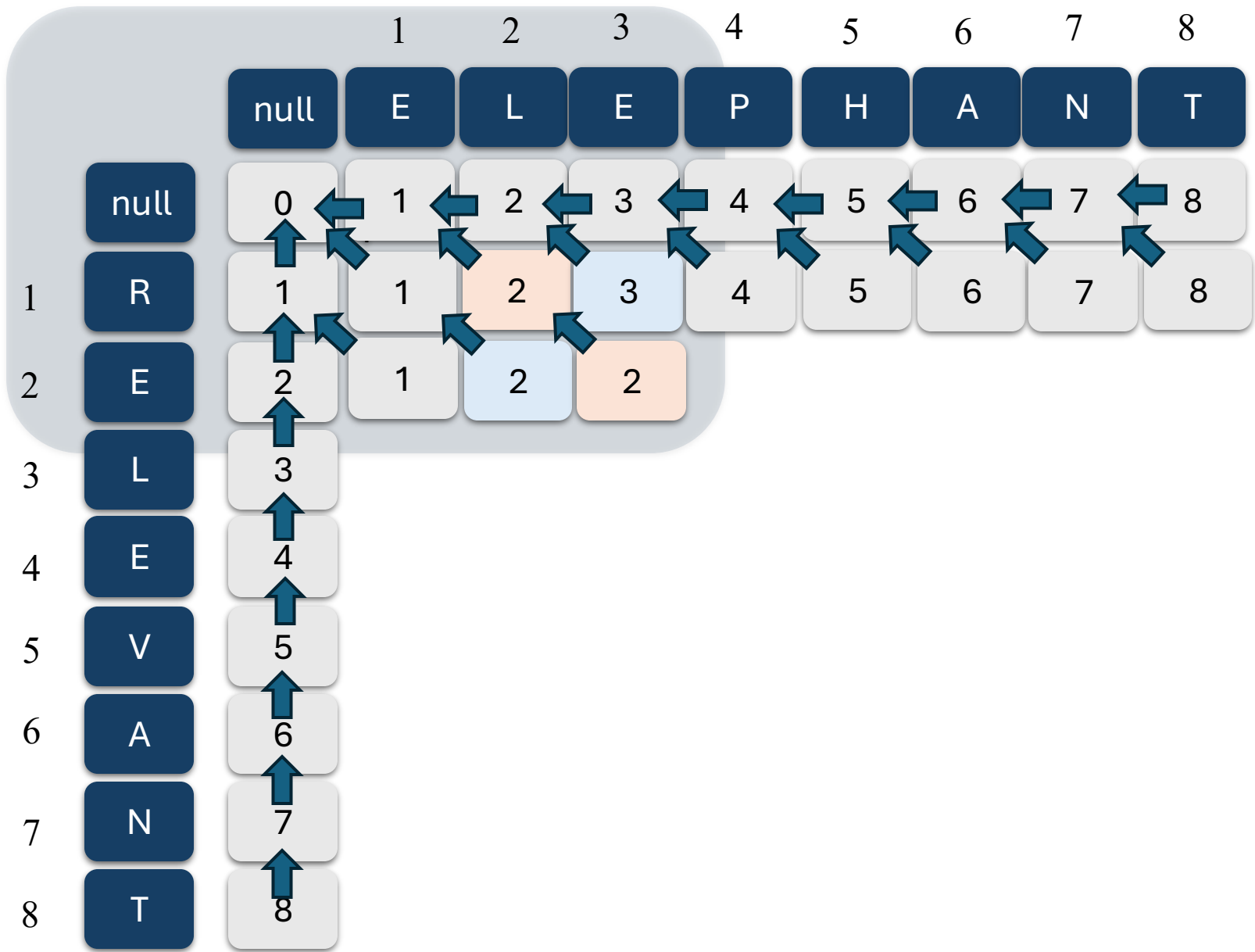


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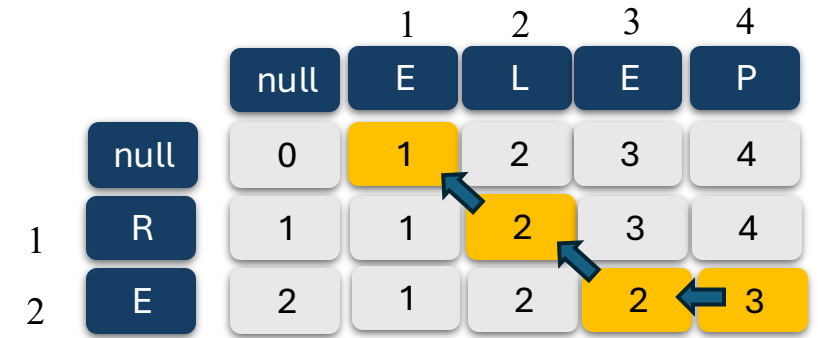
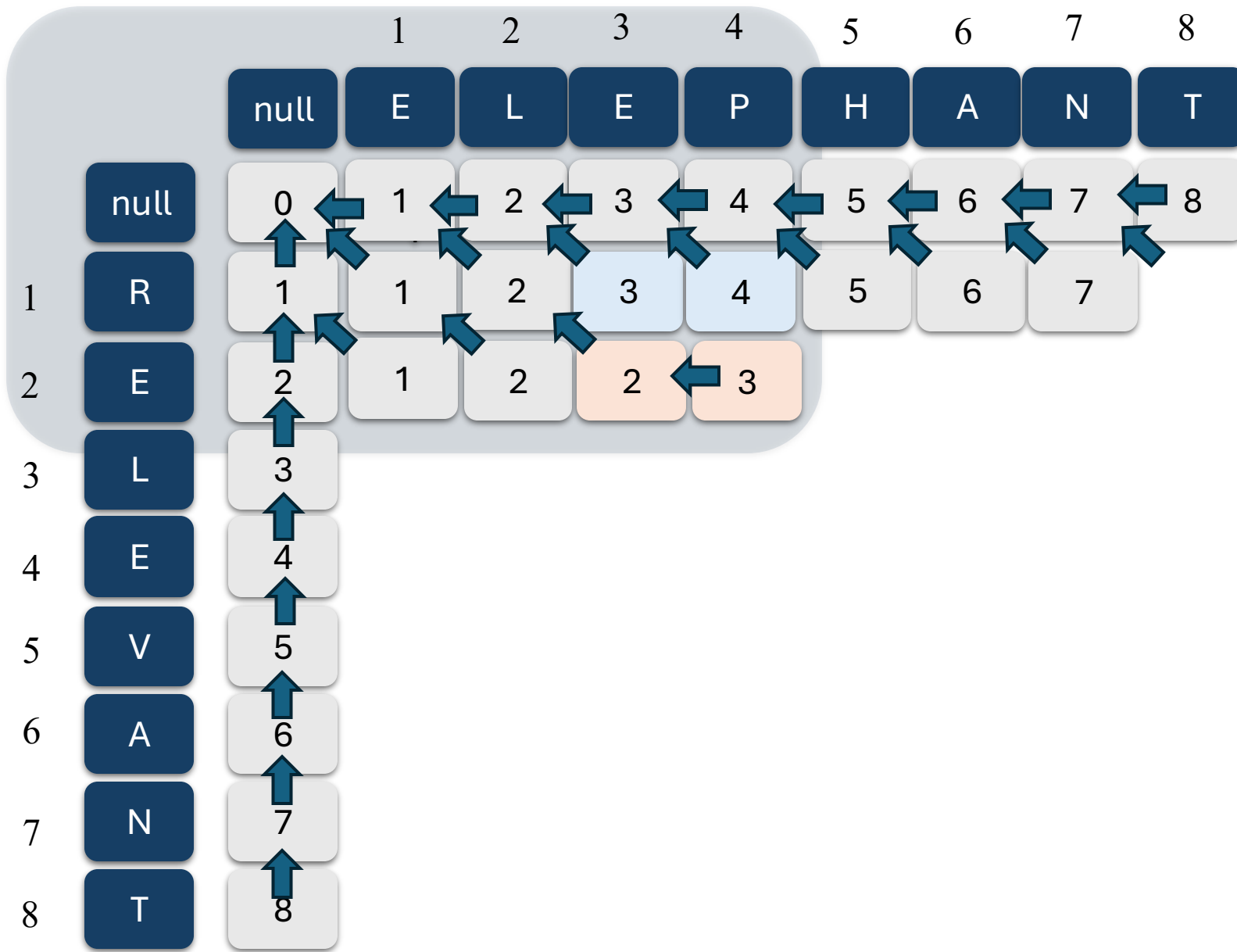
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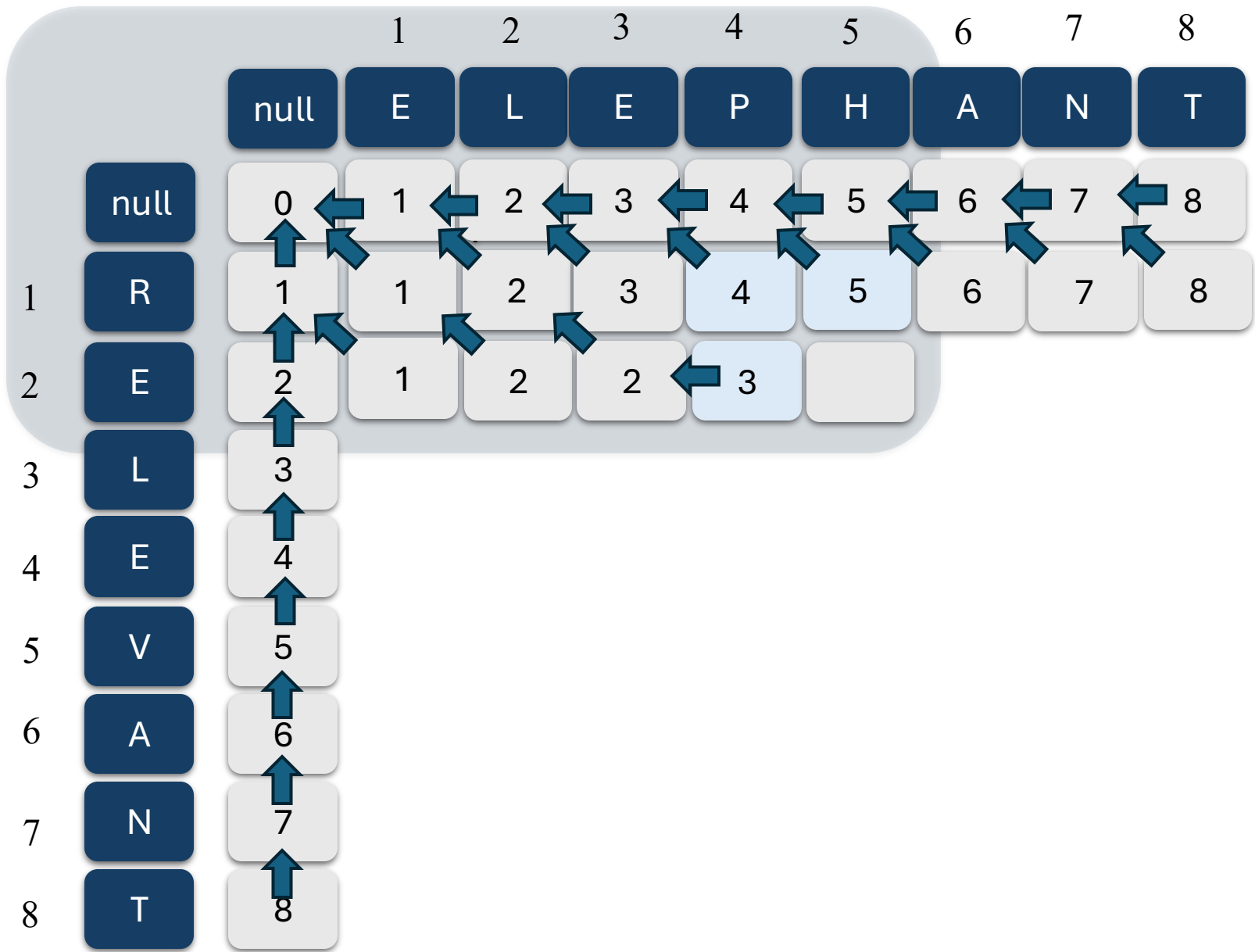


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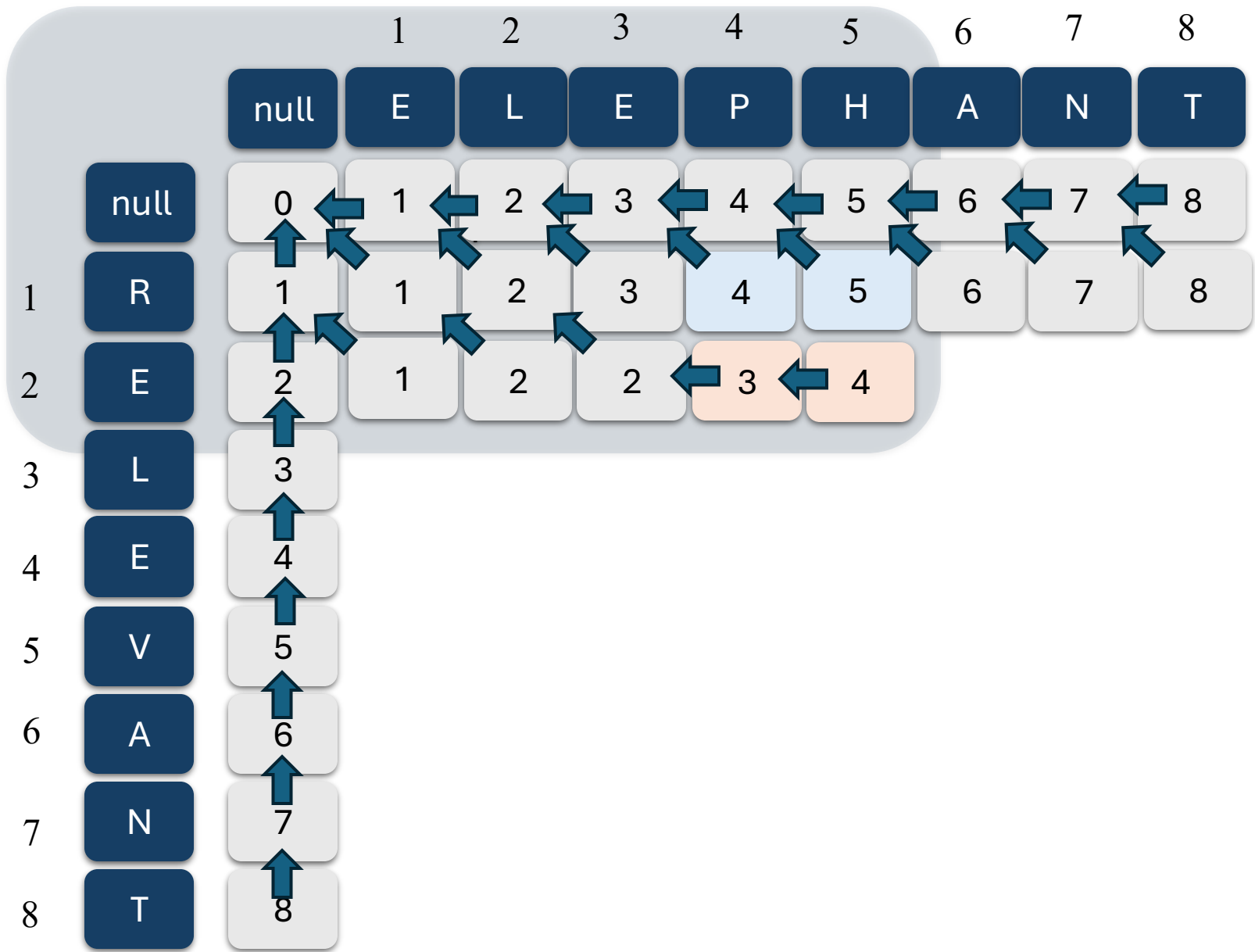


Insert E
 Replace R with L
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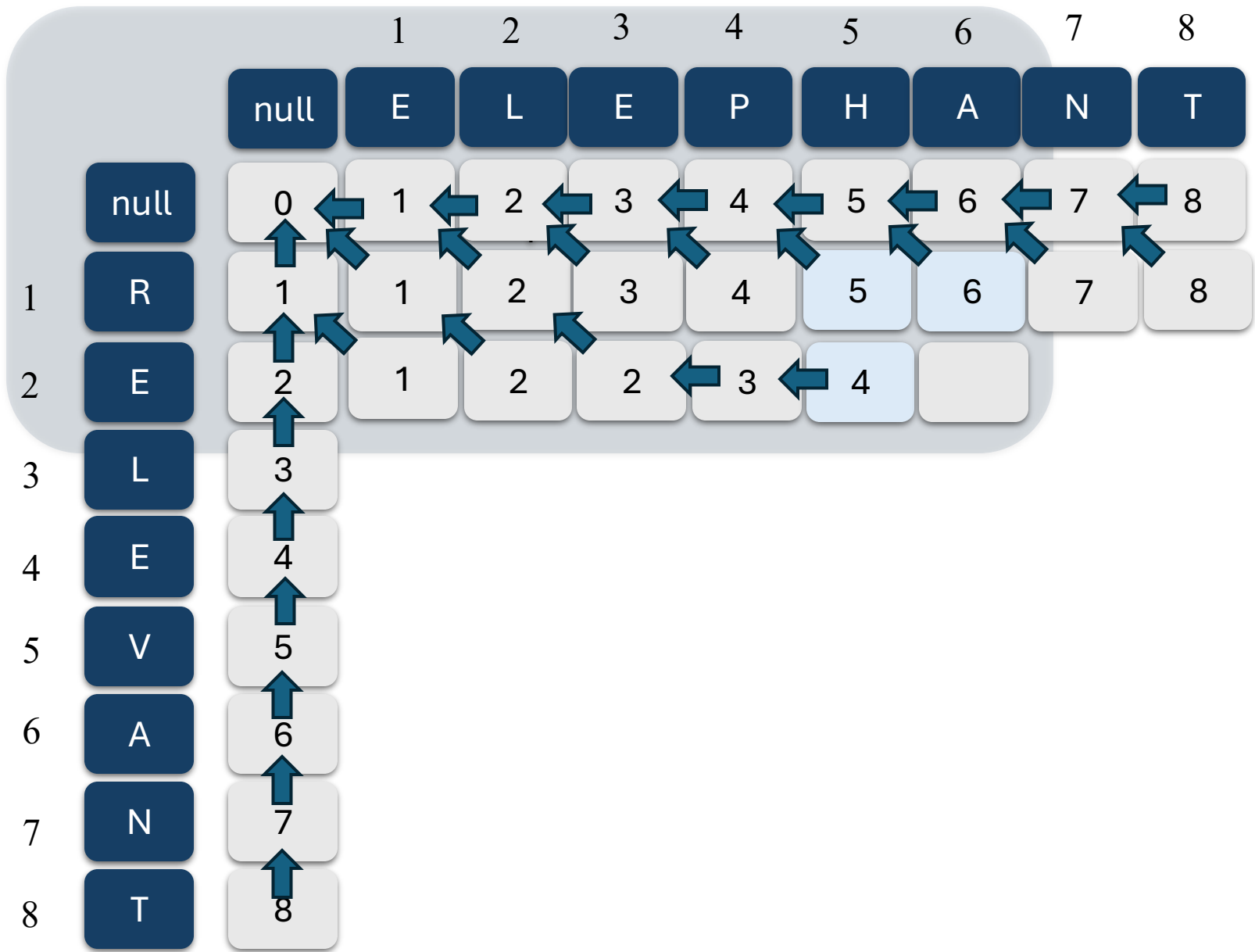
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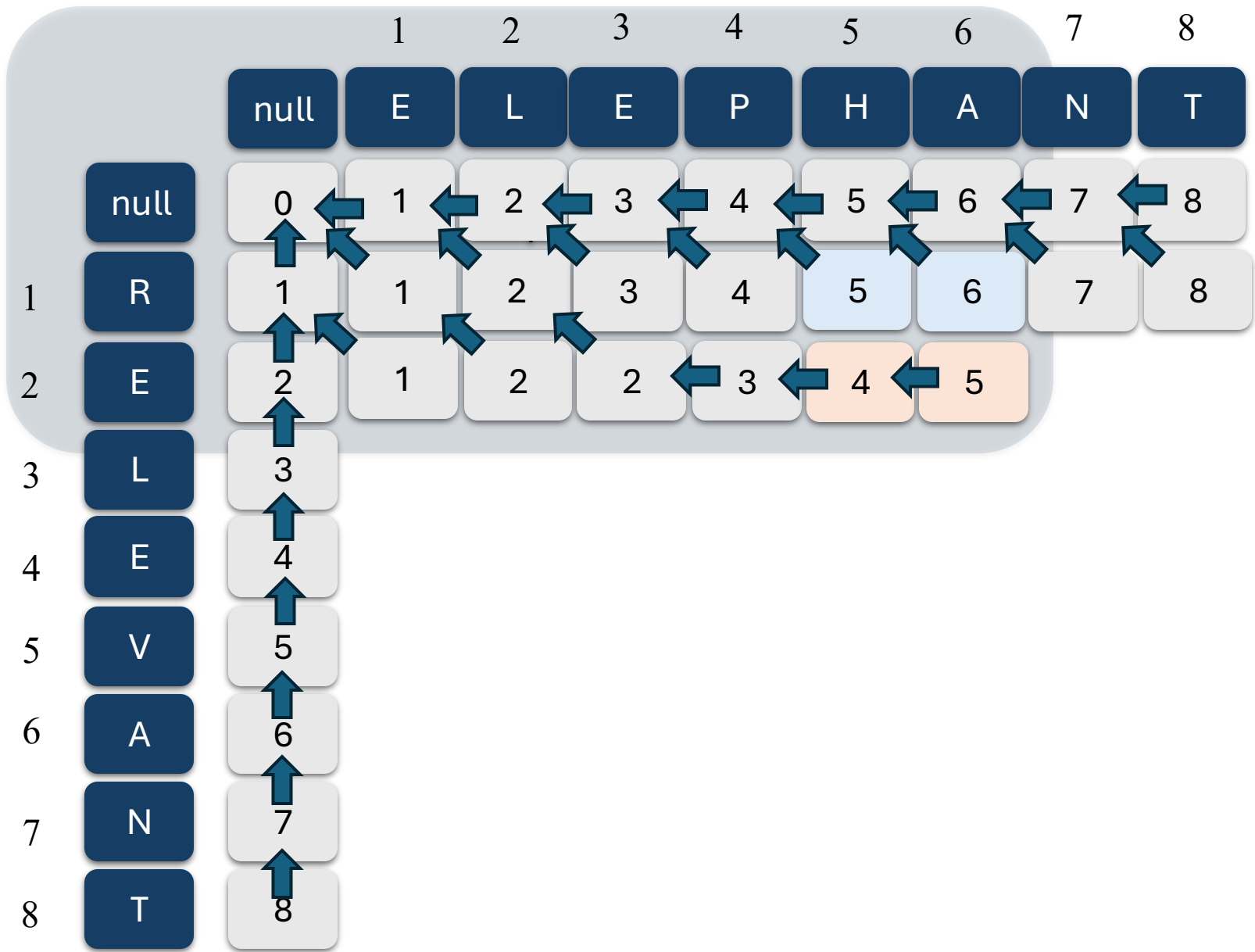
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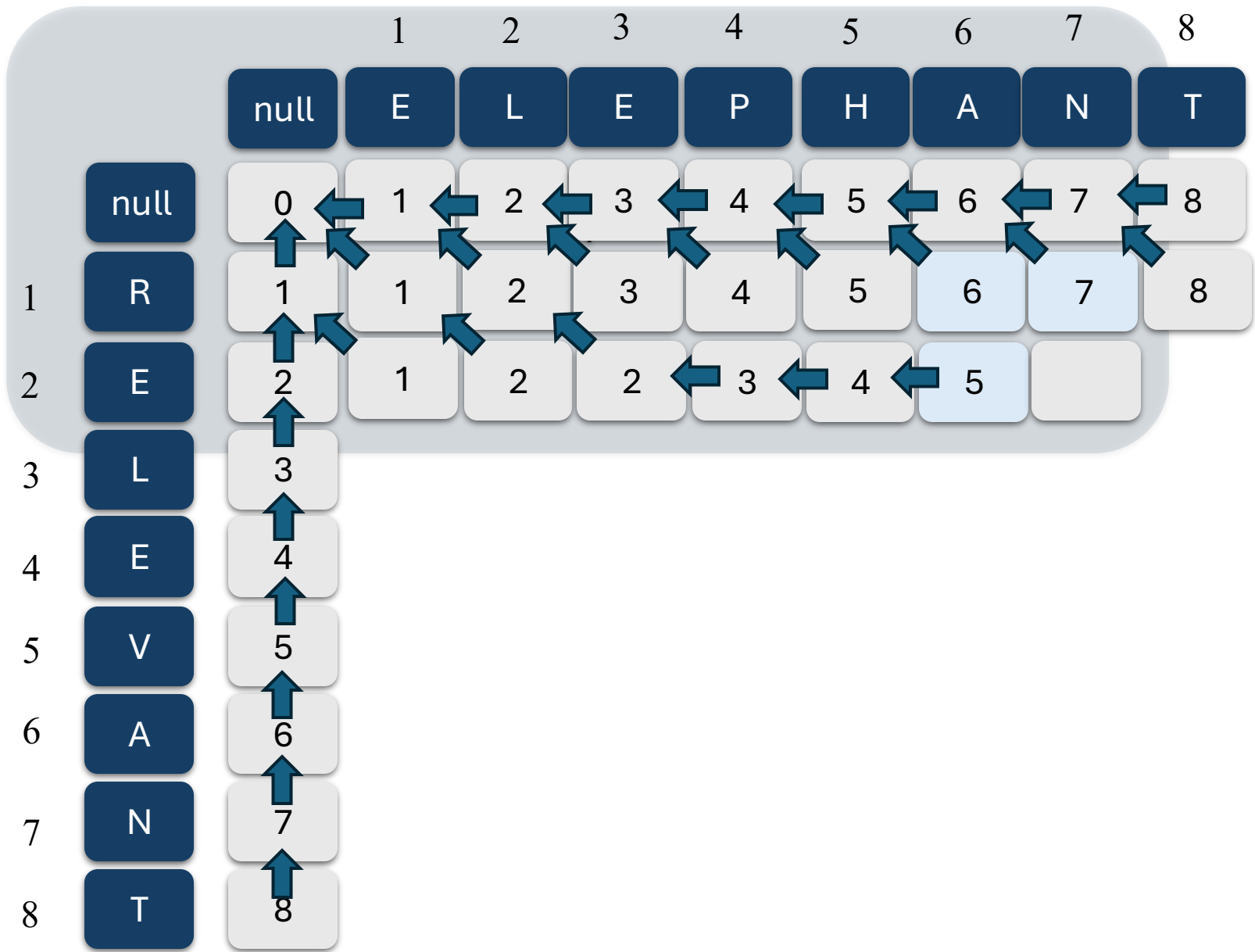
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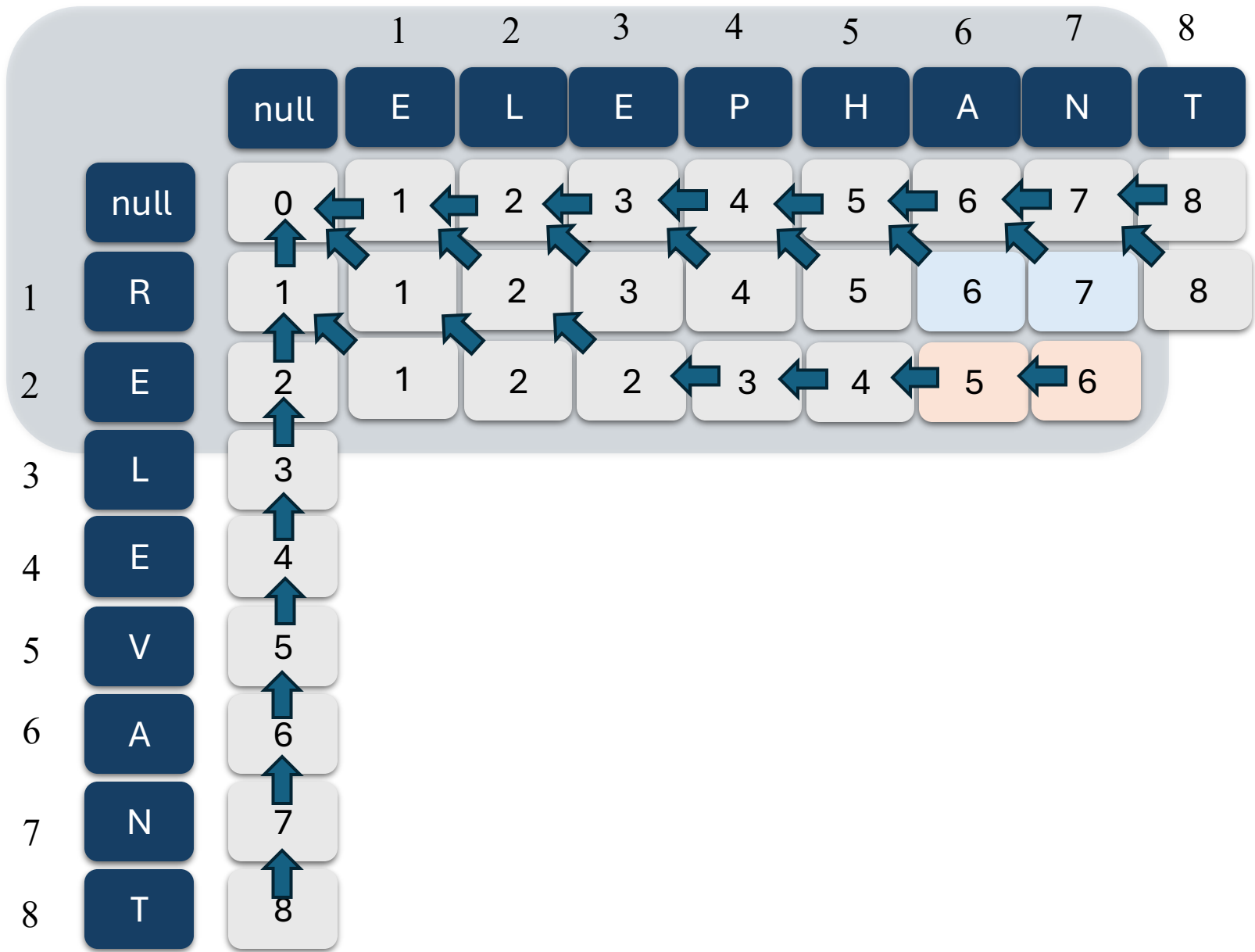
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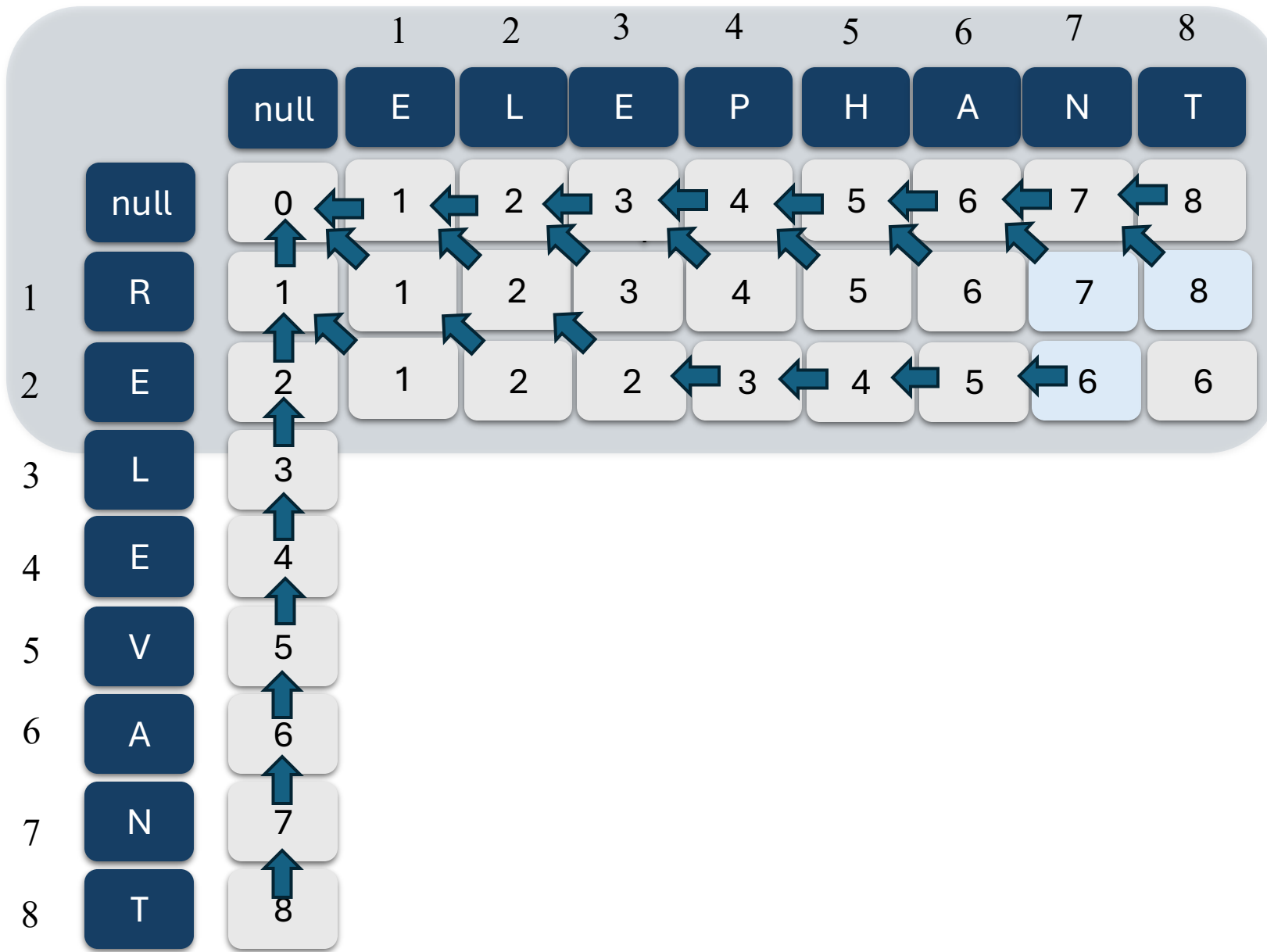
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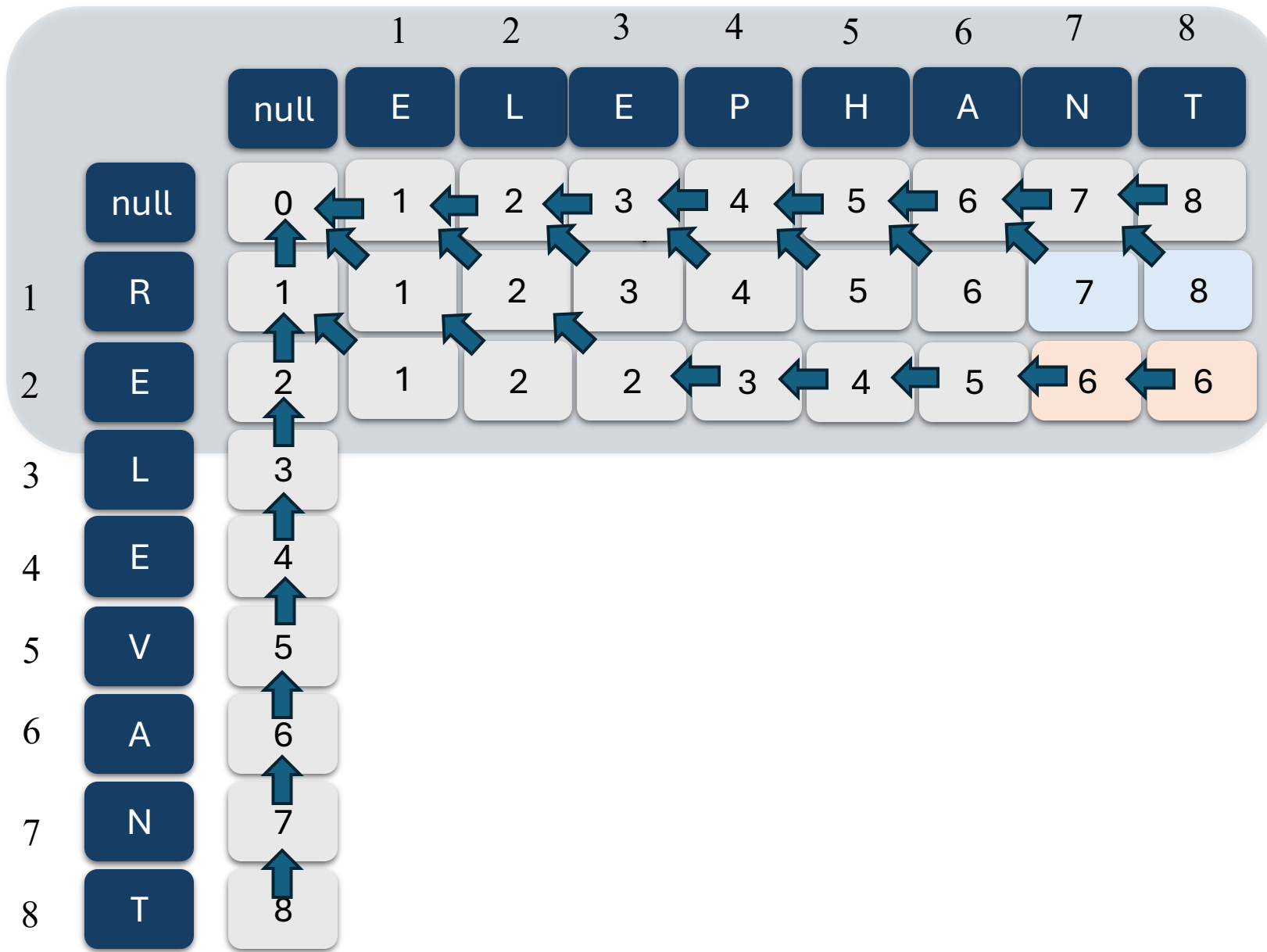
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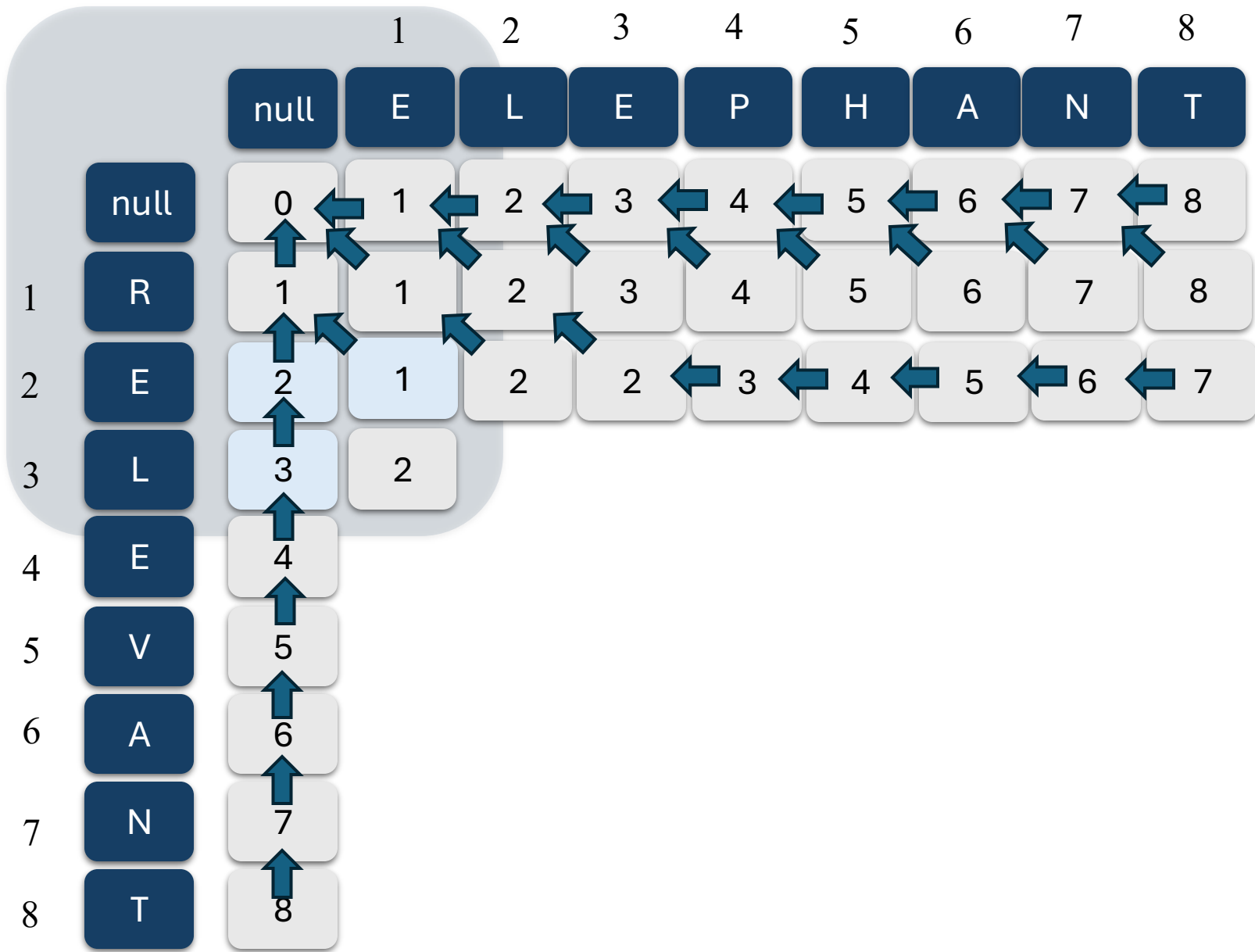
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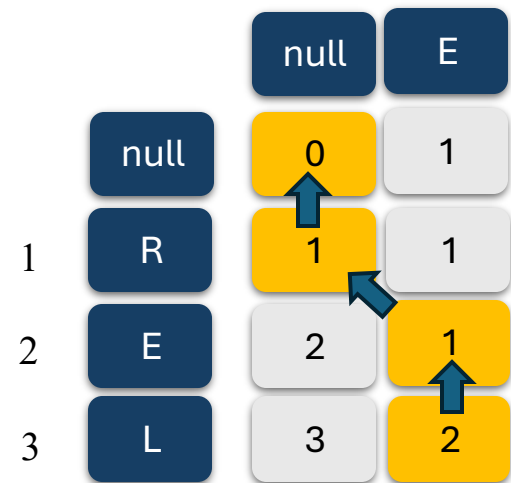
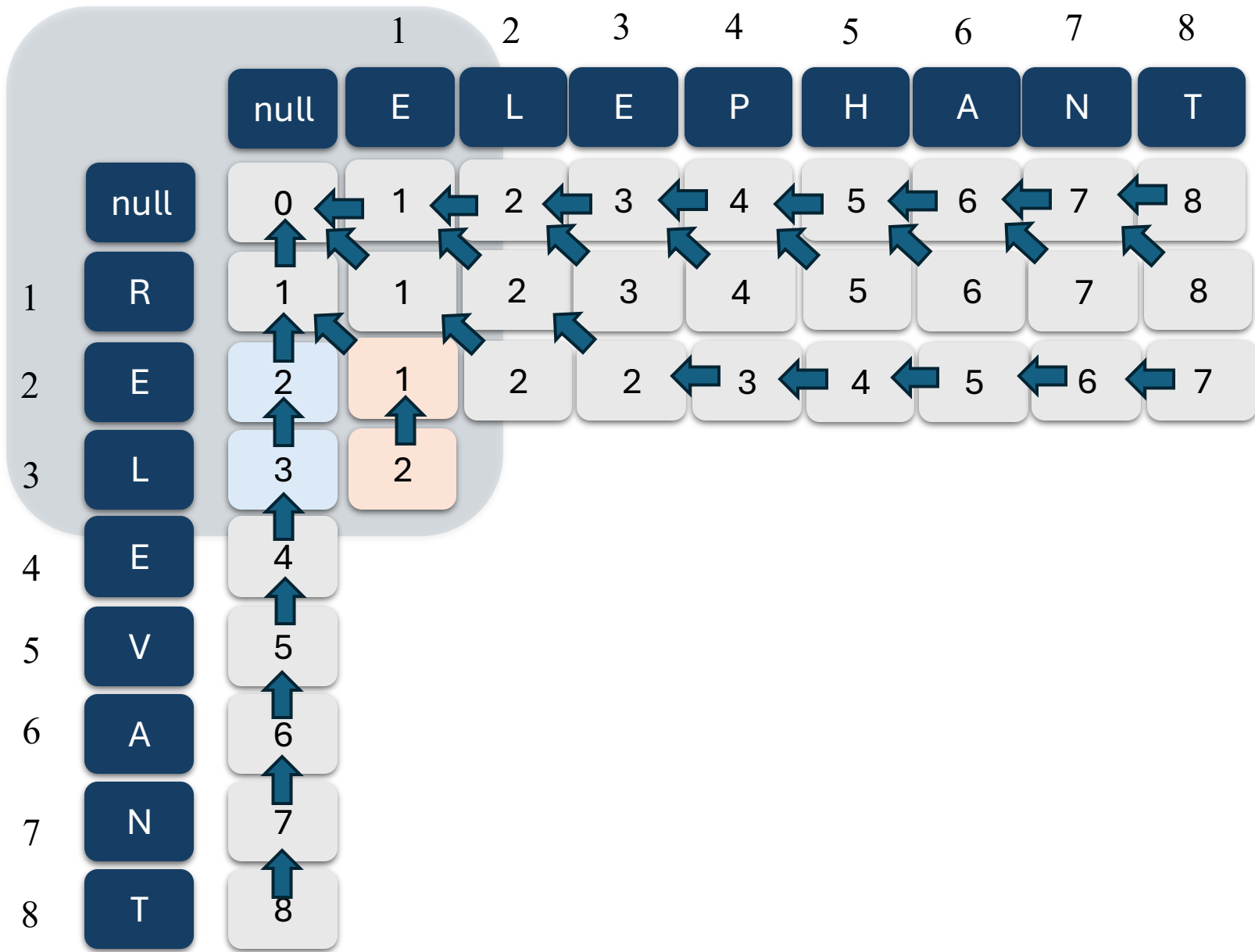
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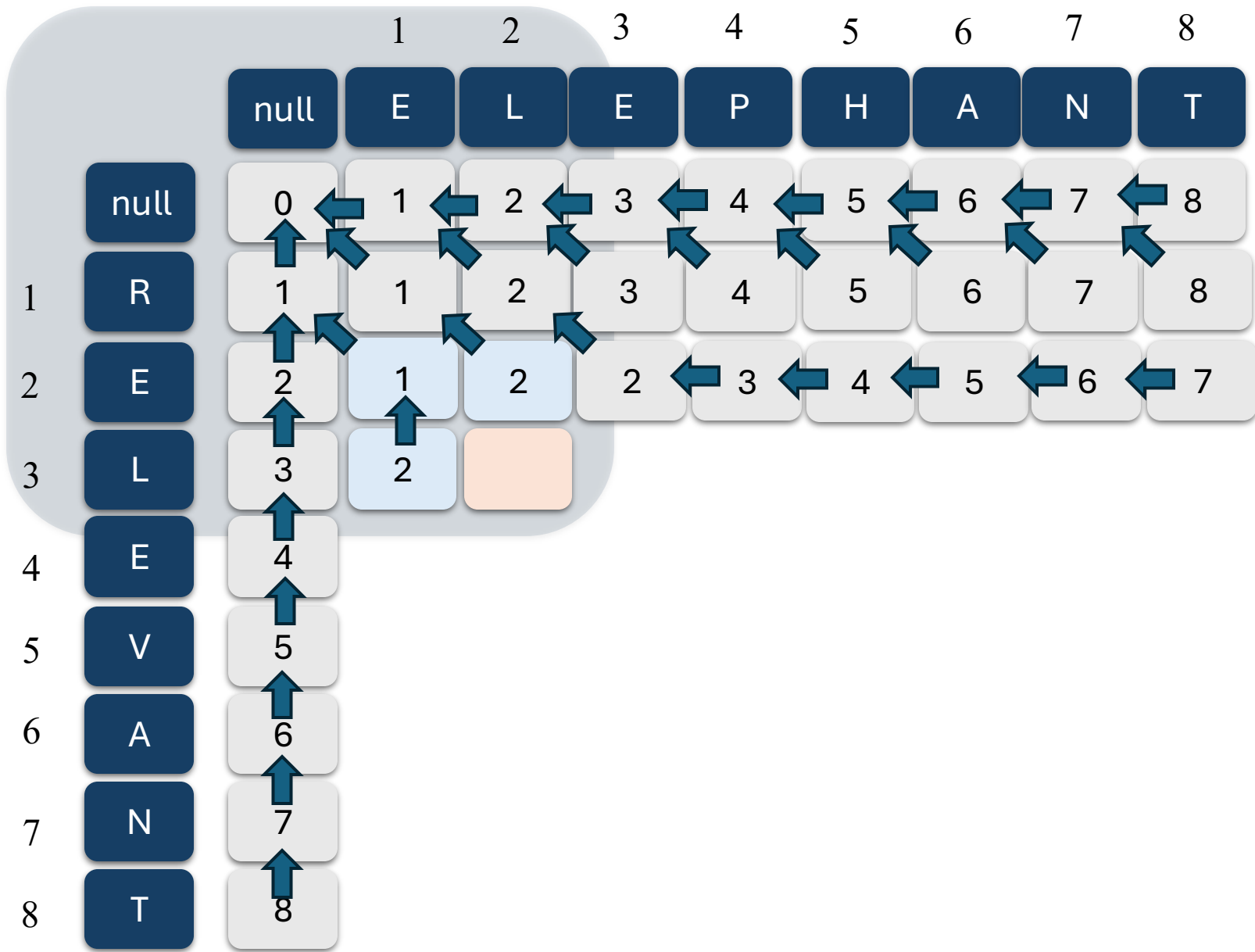


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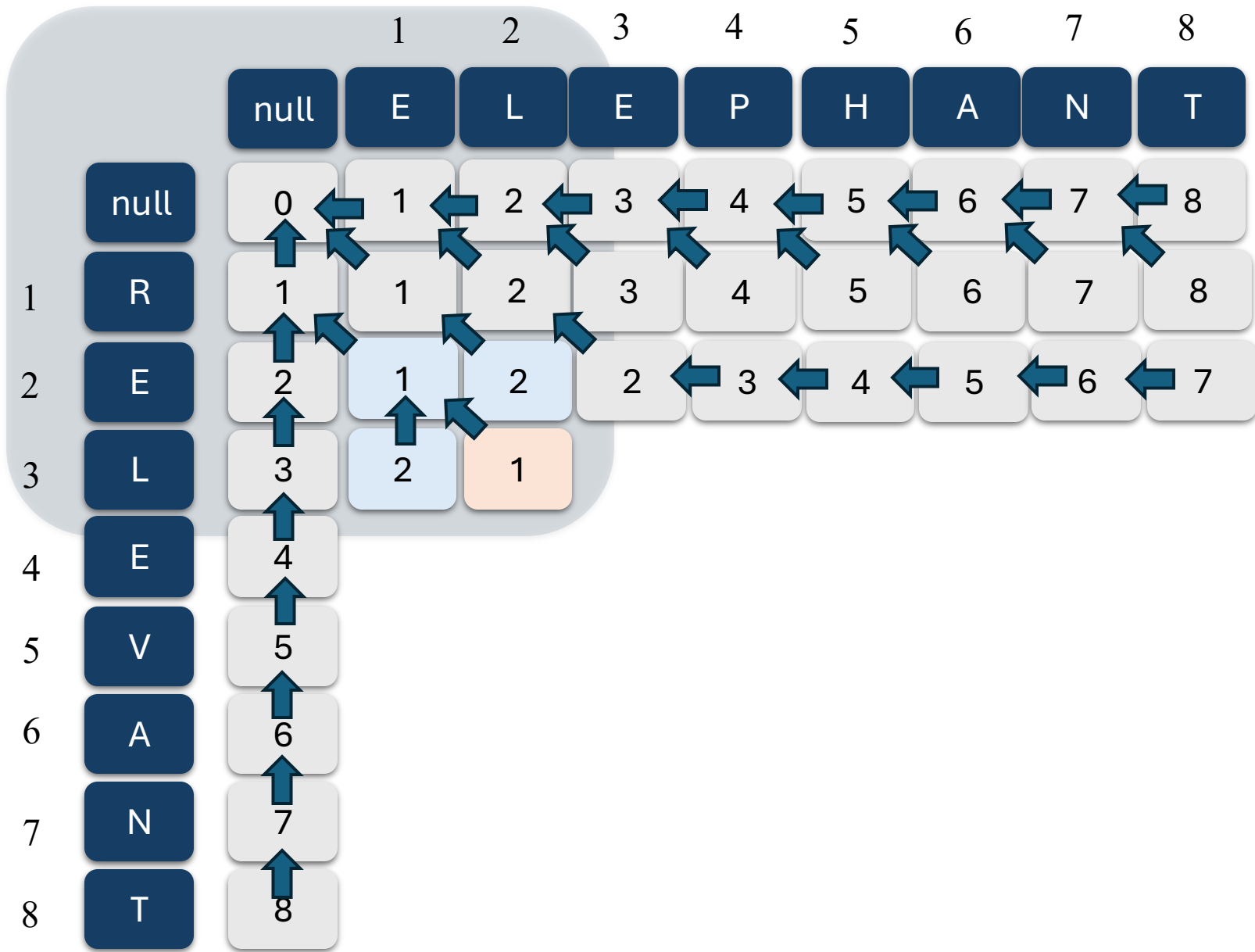


Delete R
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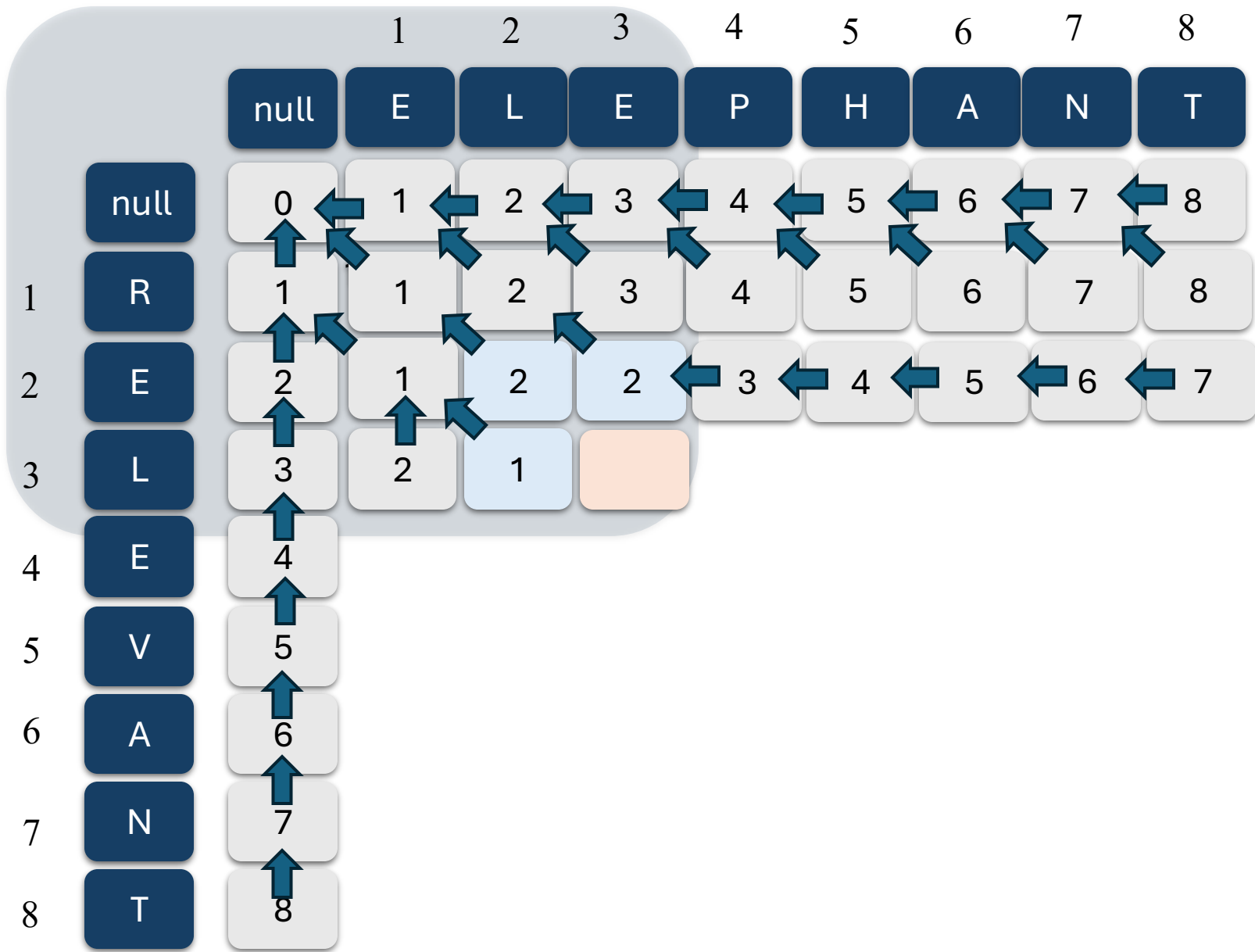
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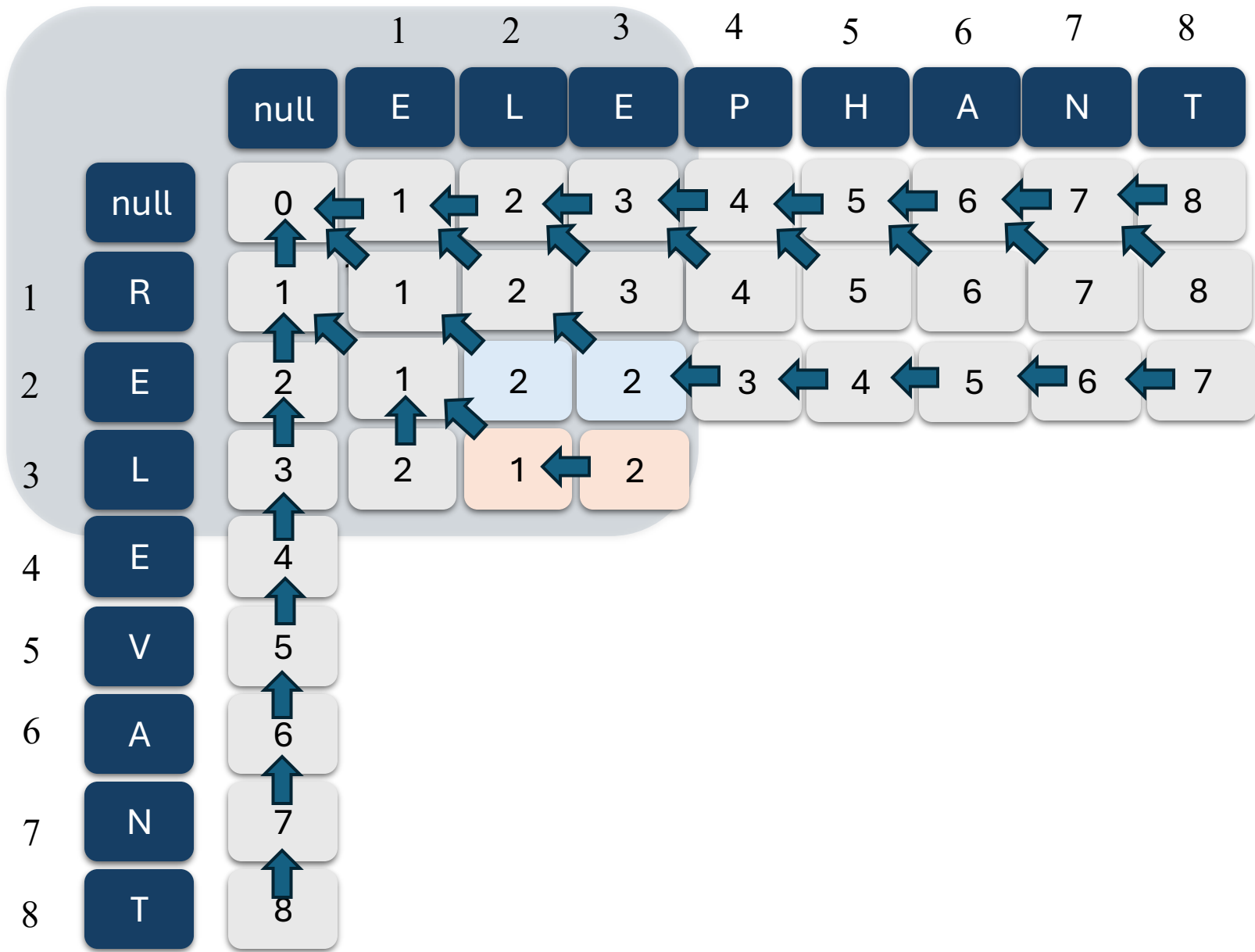
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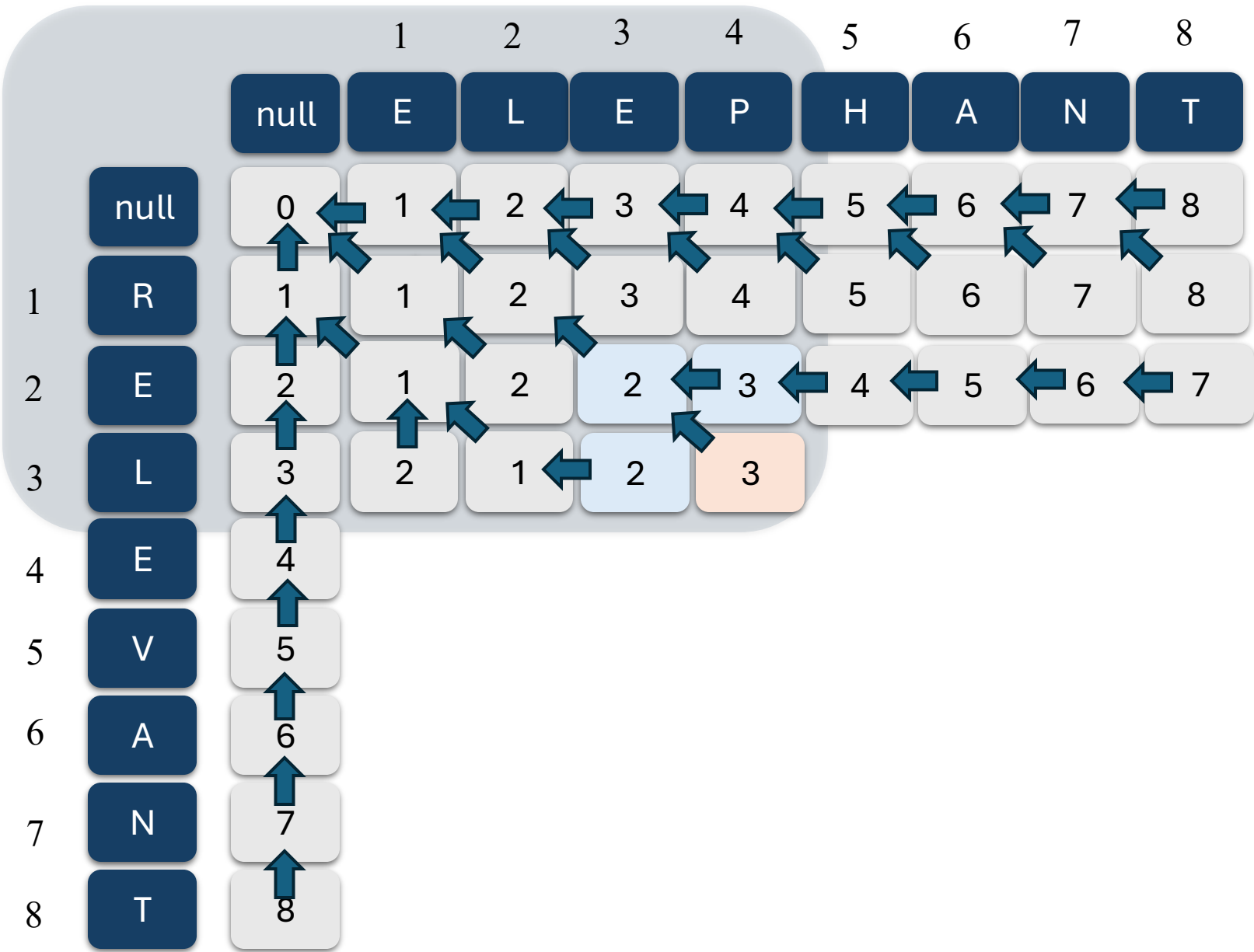
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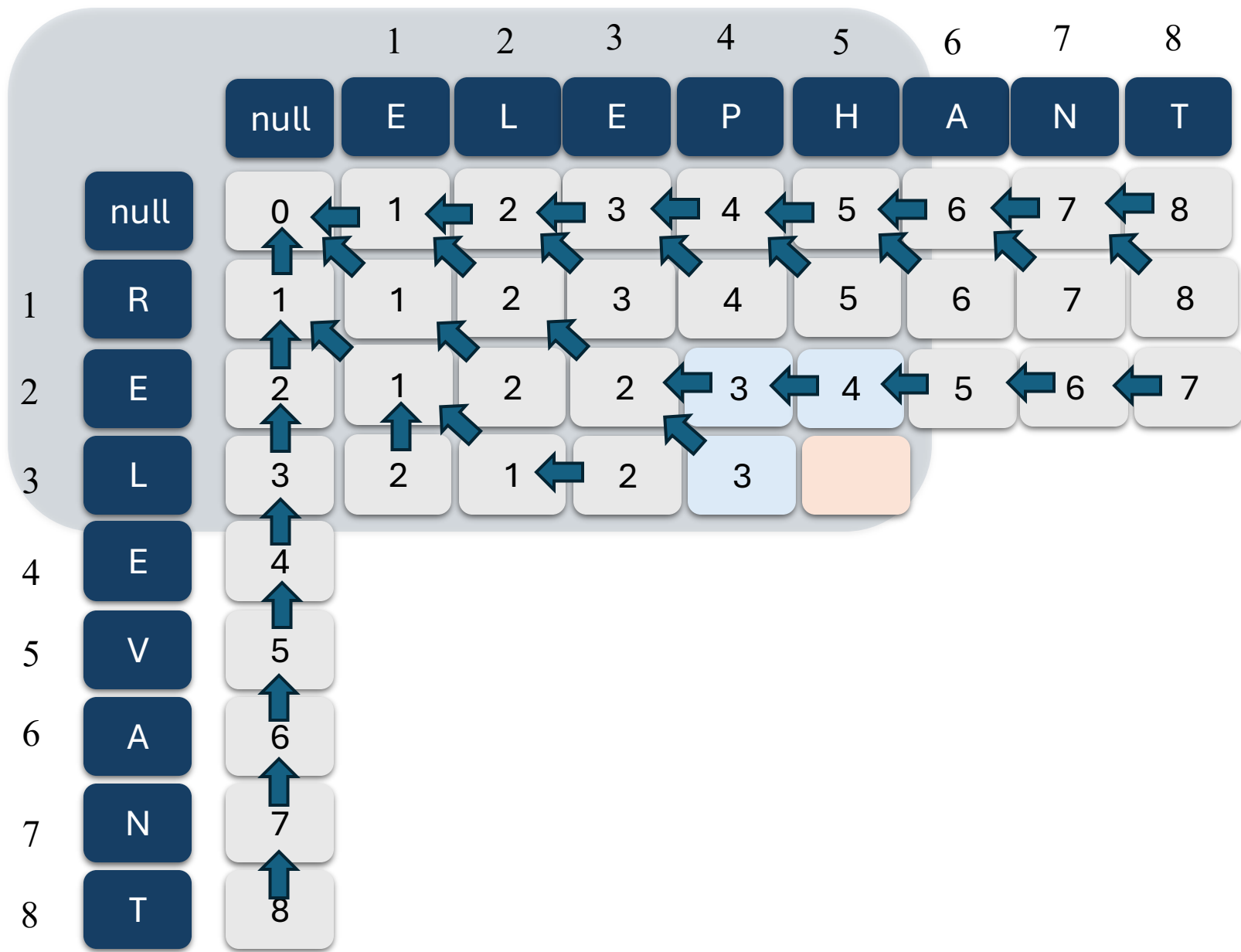
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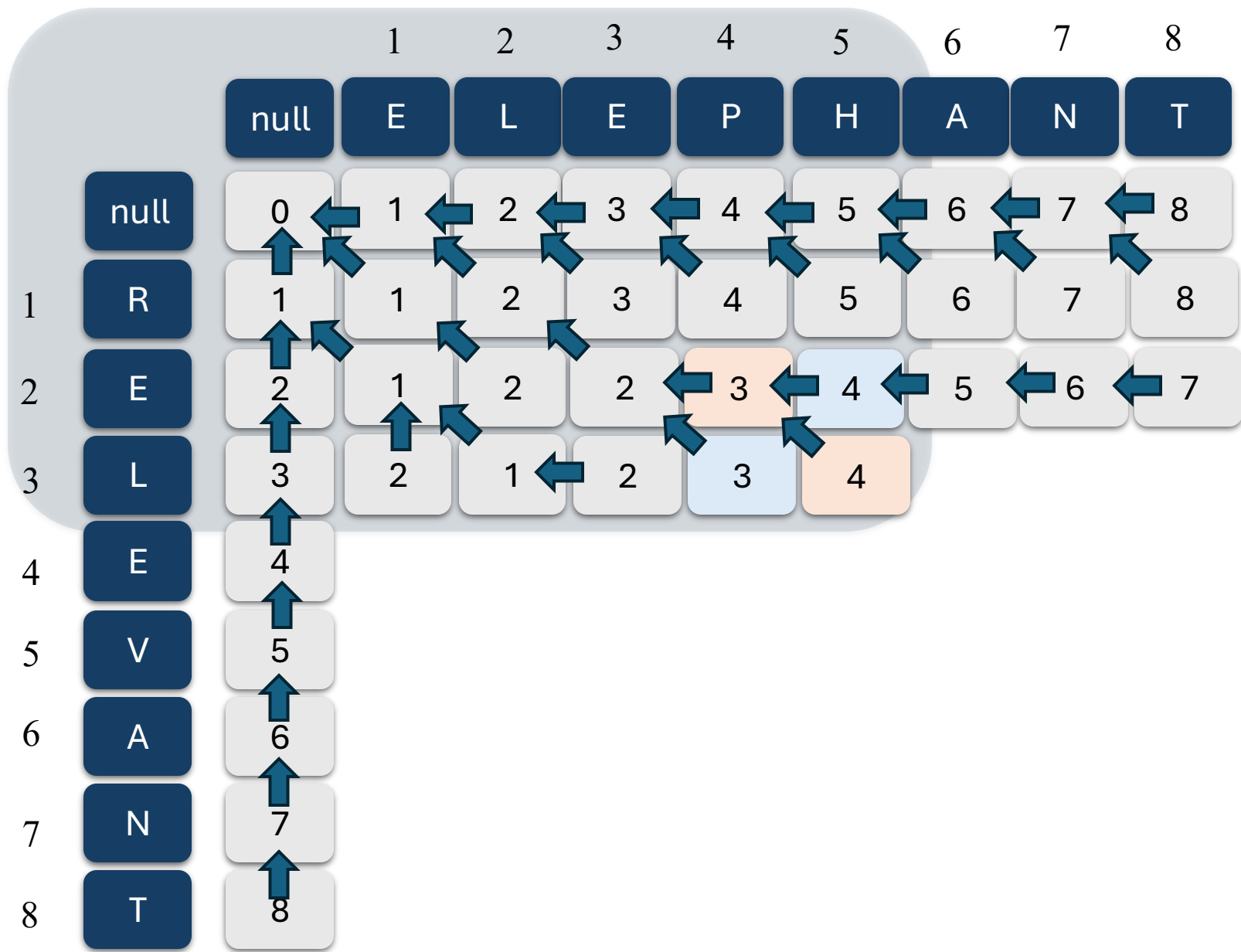
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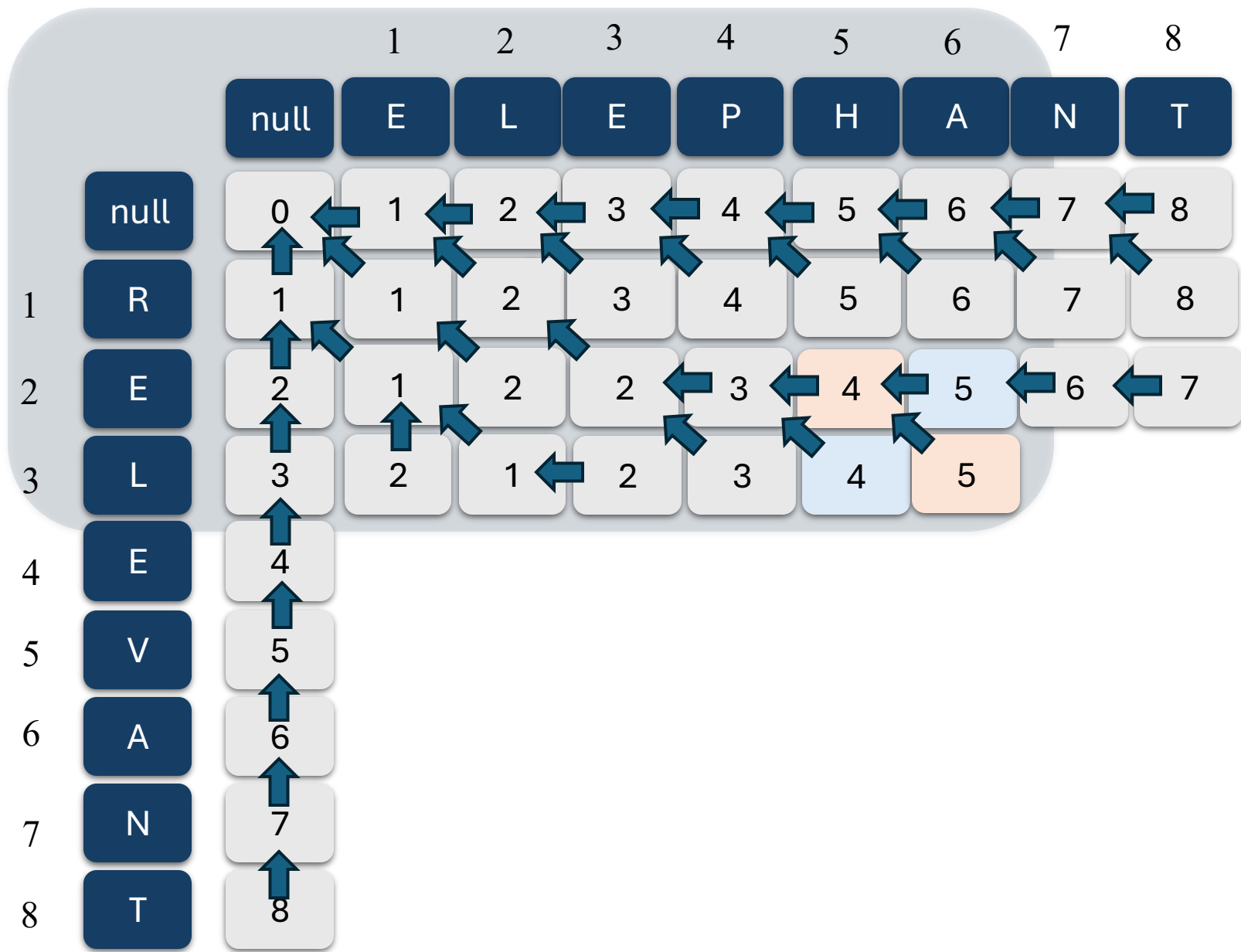
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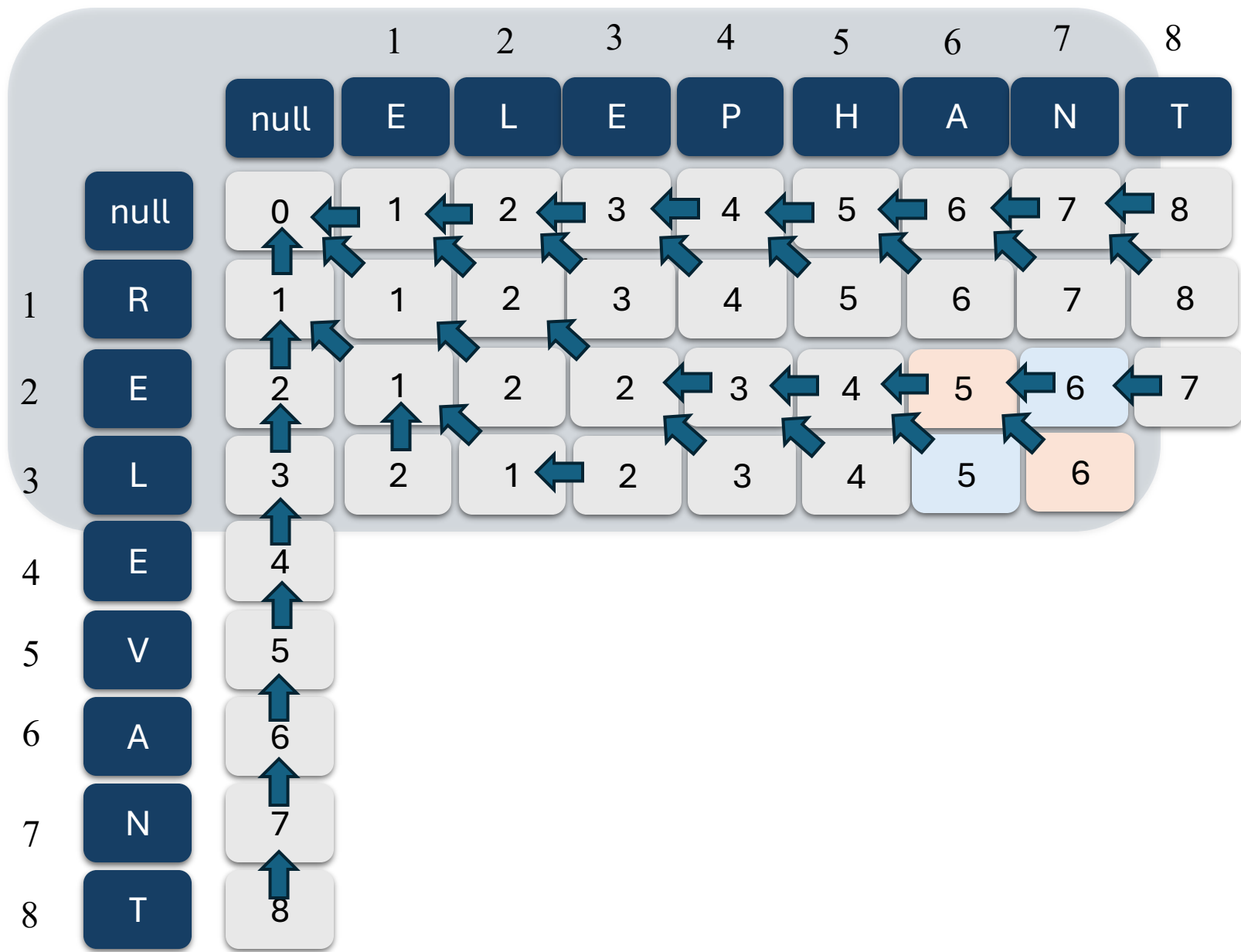
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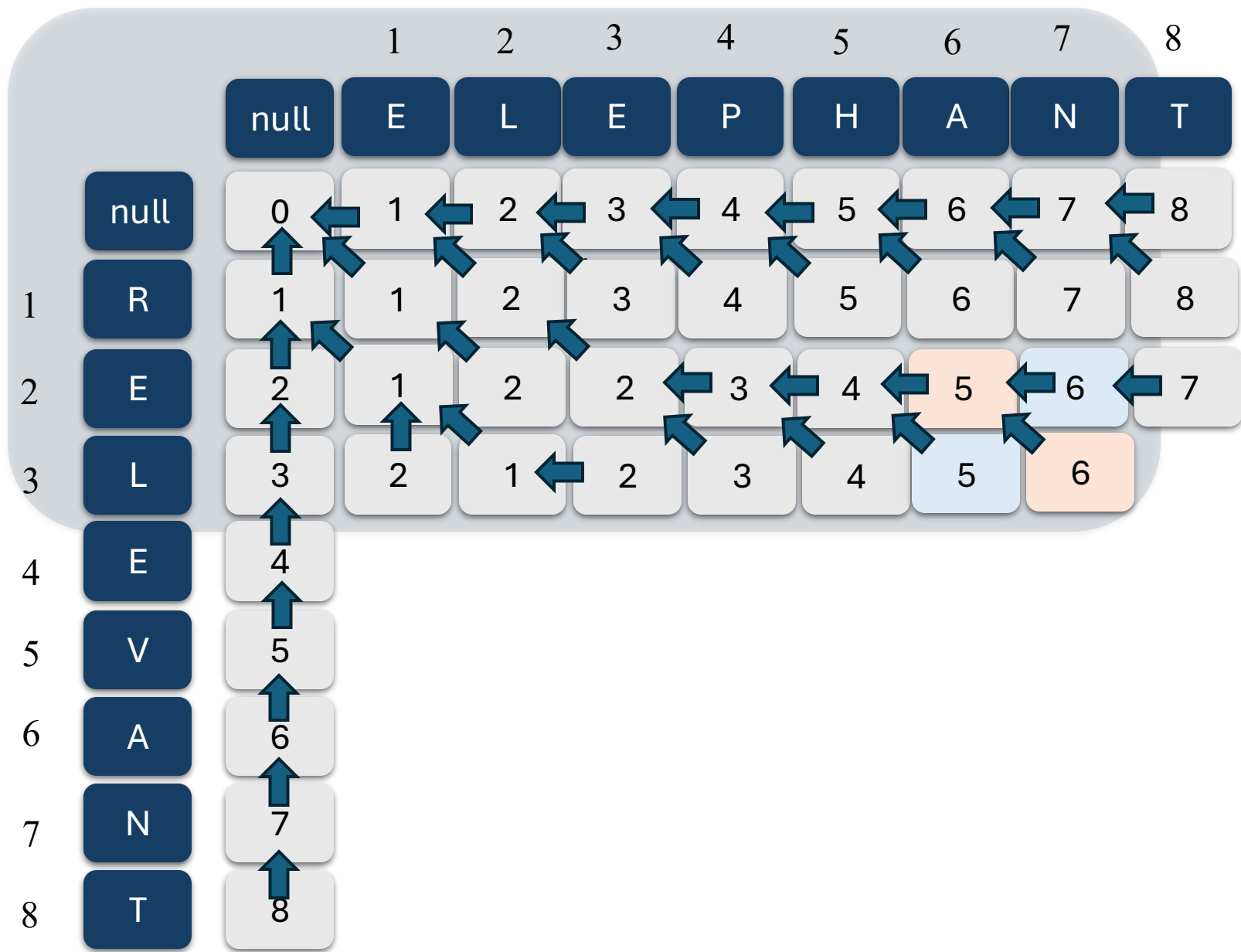
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



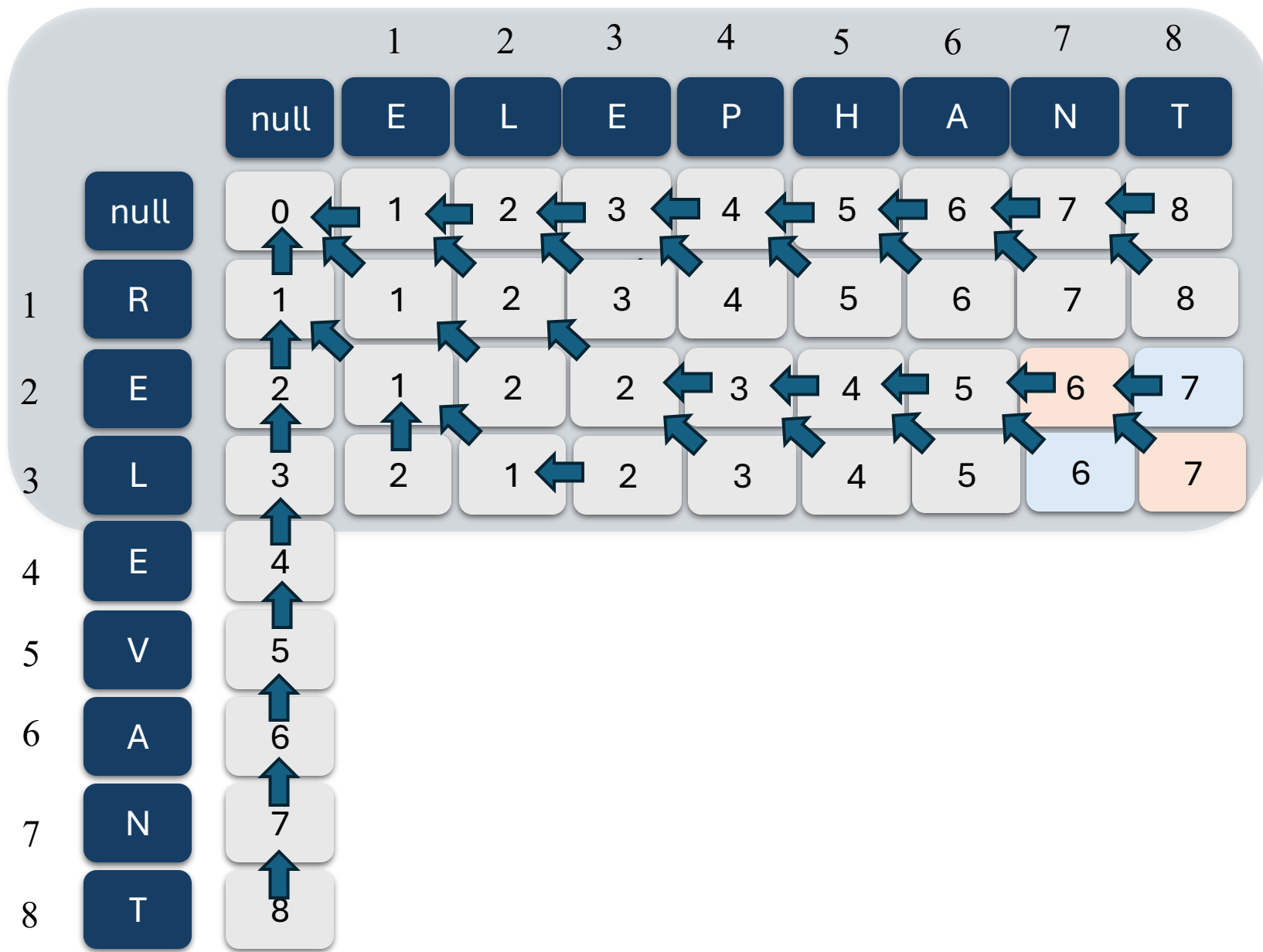
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



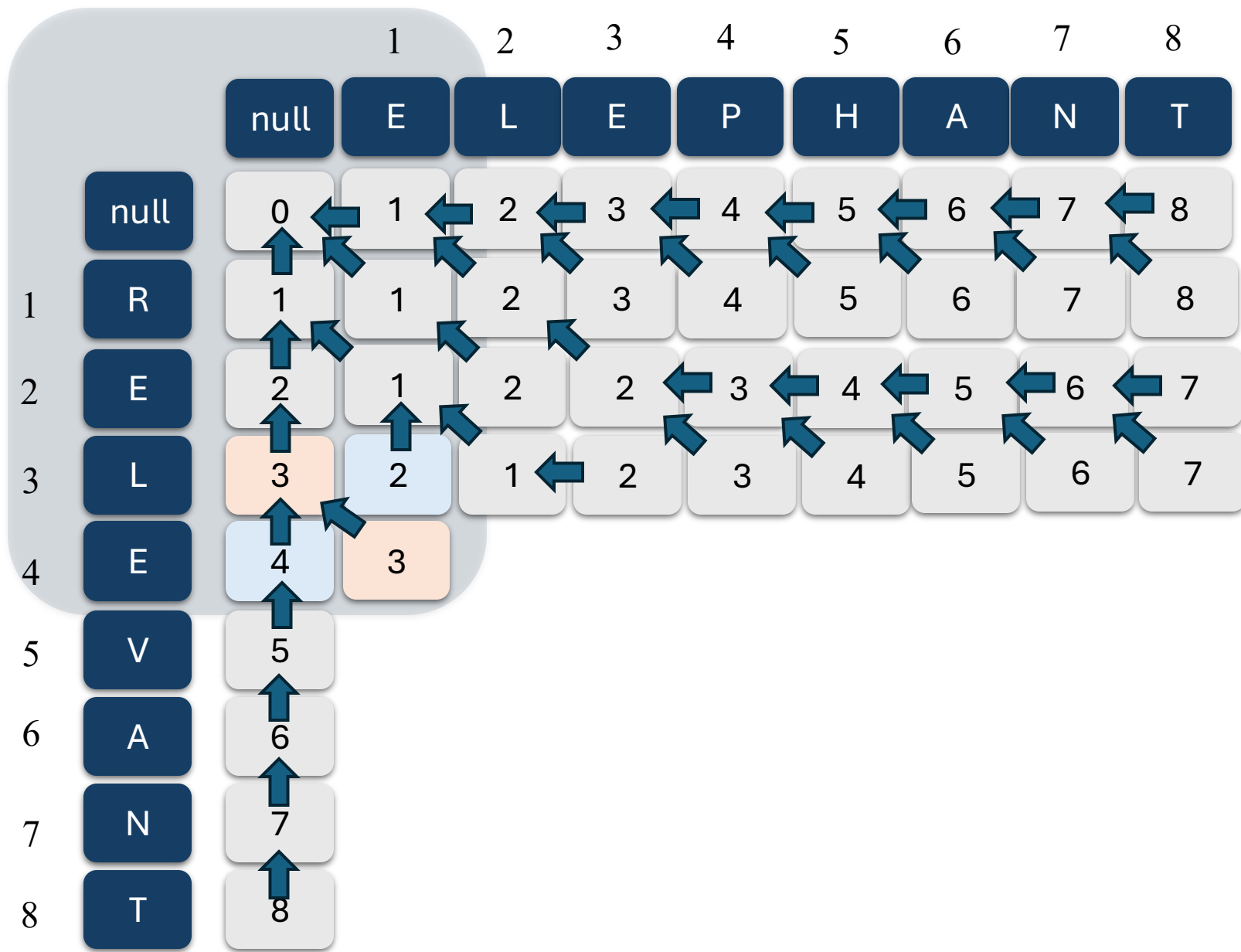
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



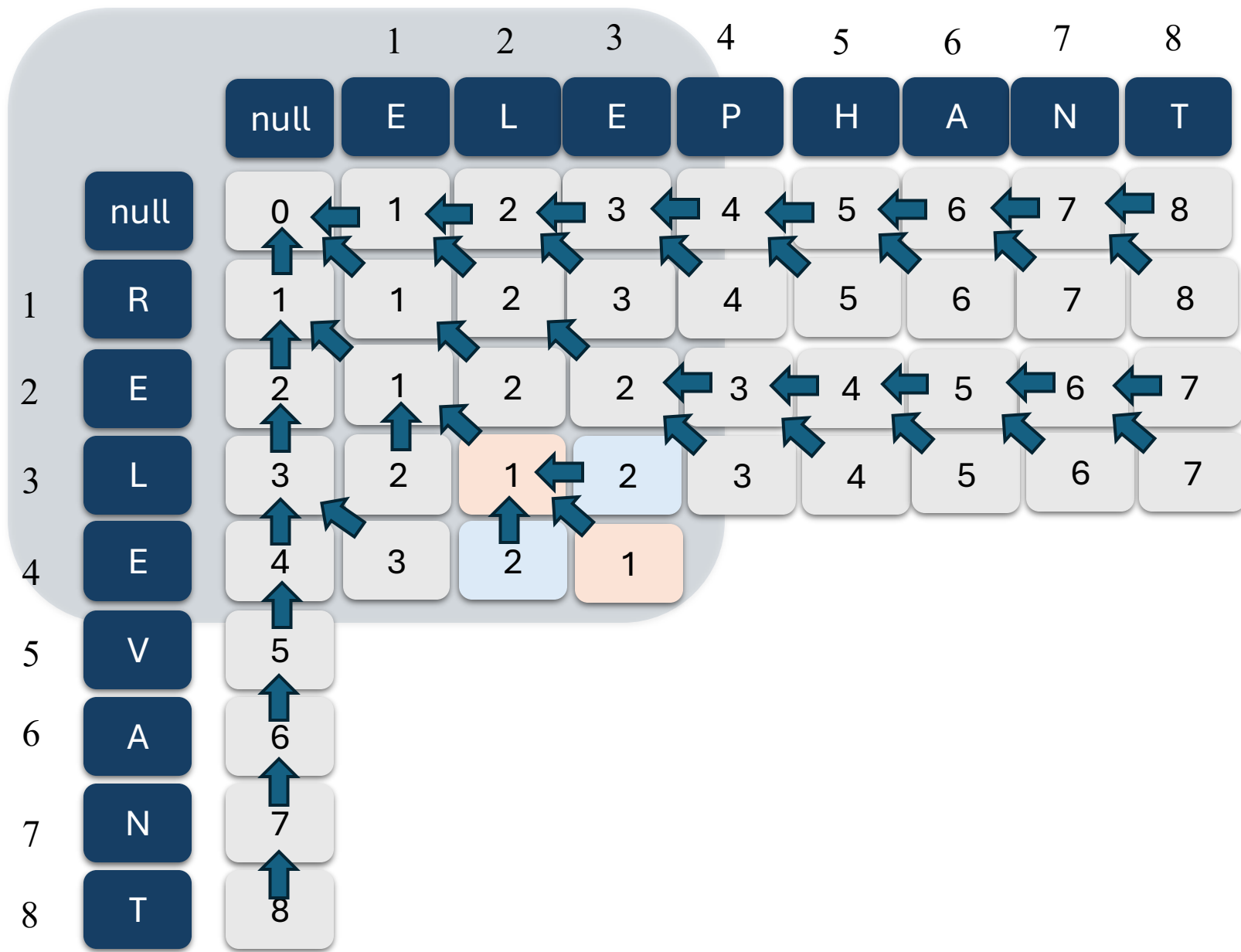
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



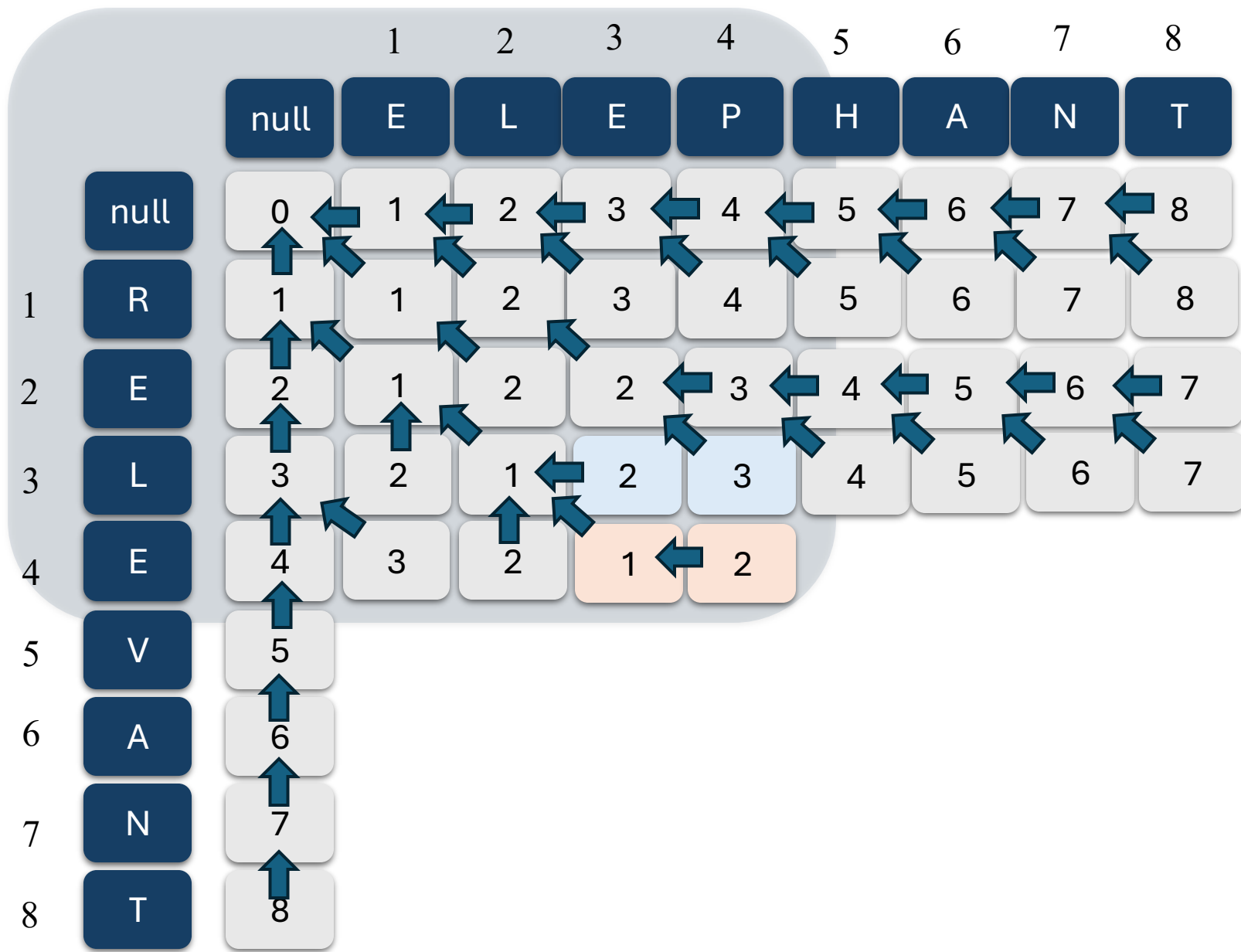
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



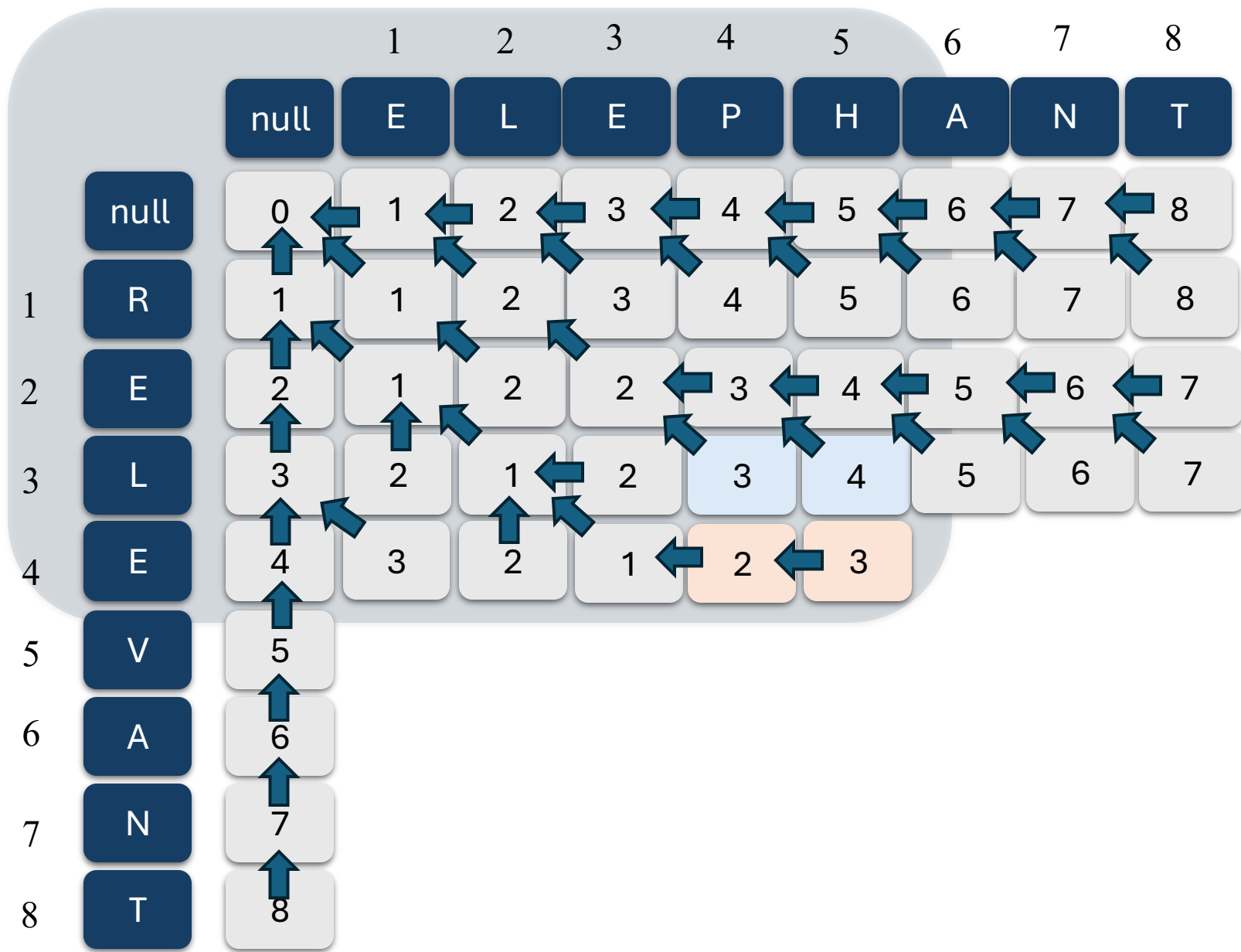
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



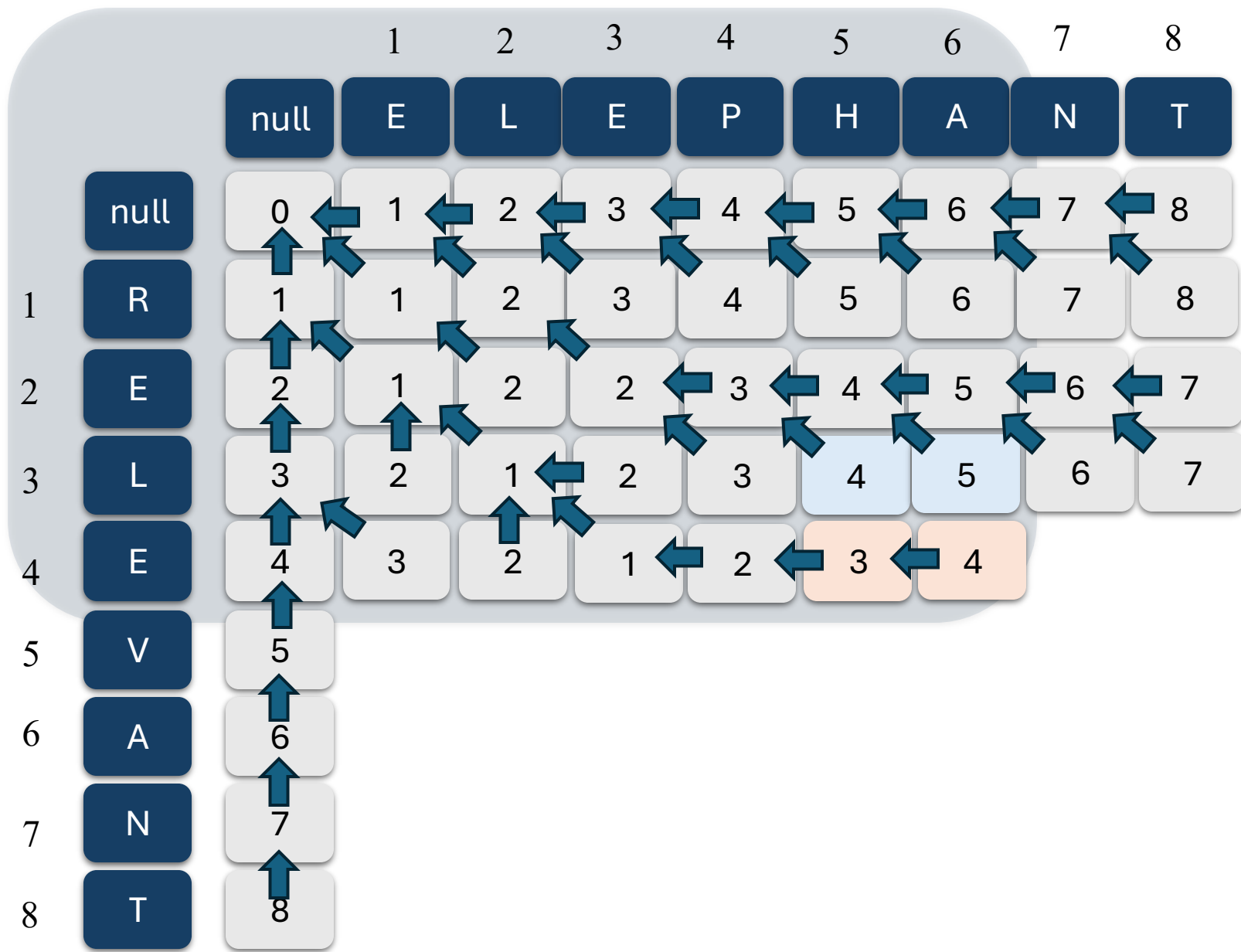
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



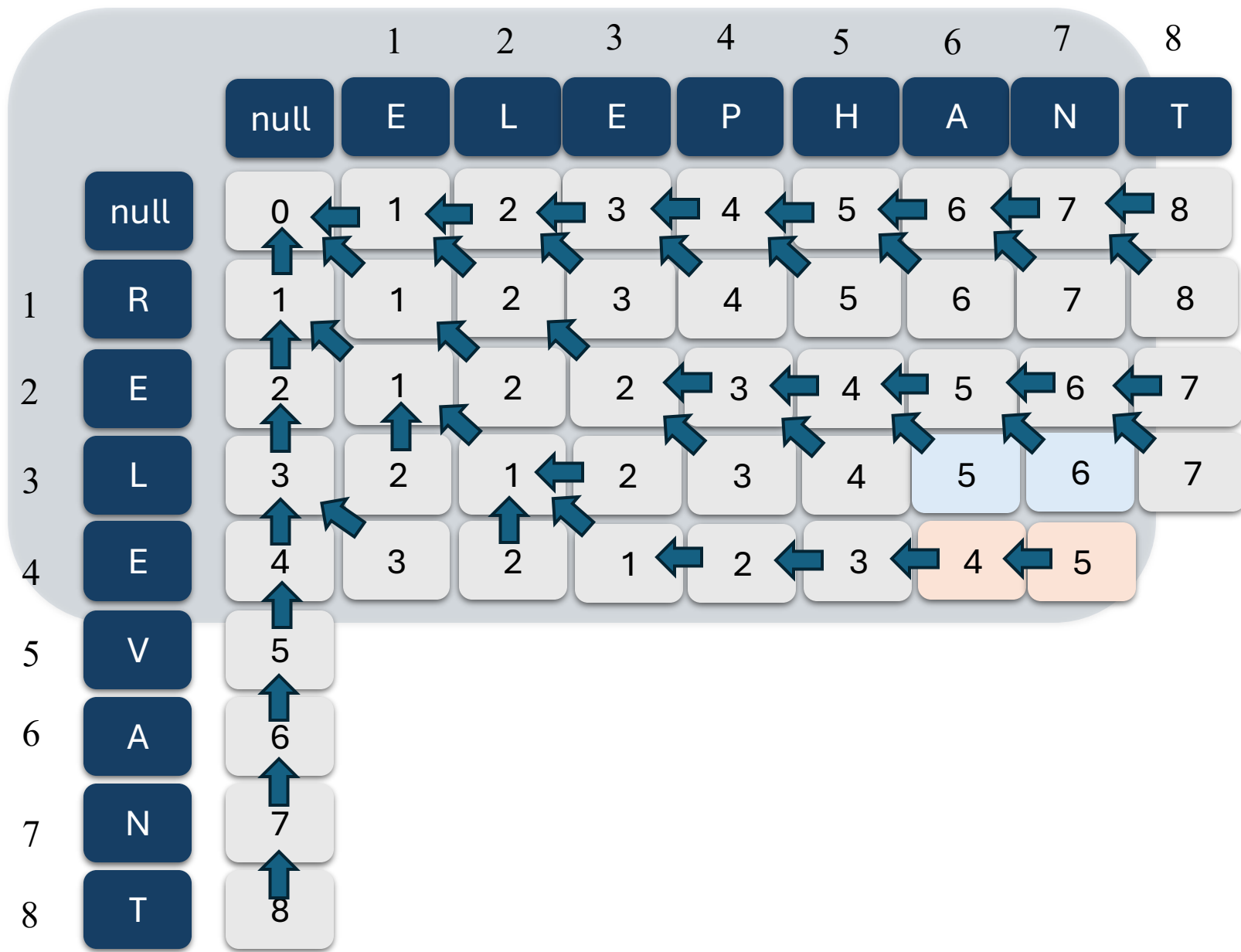
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



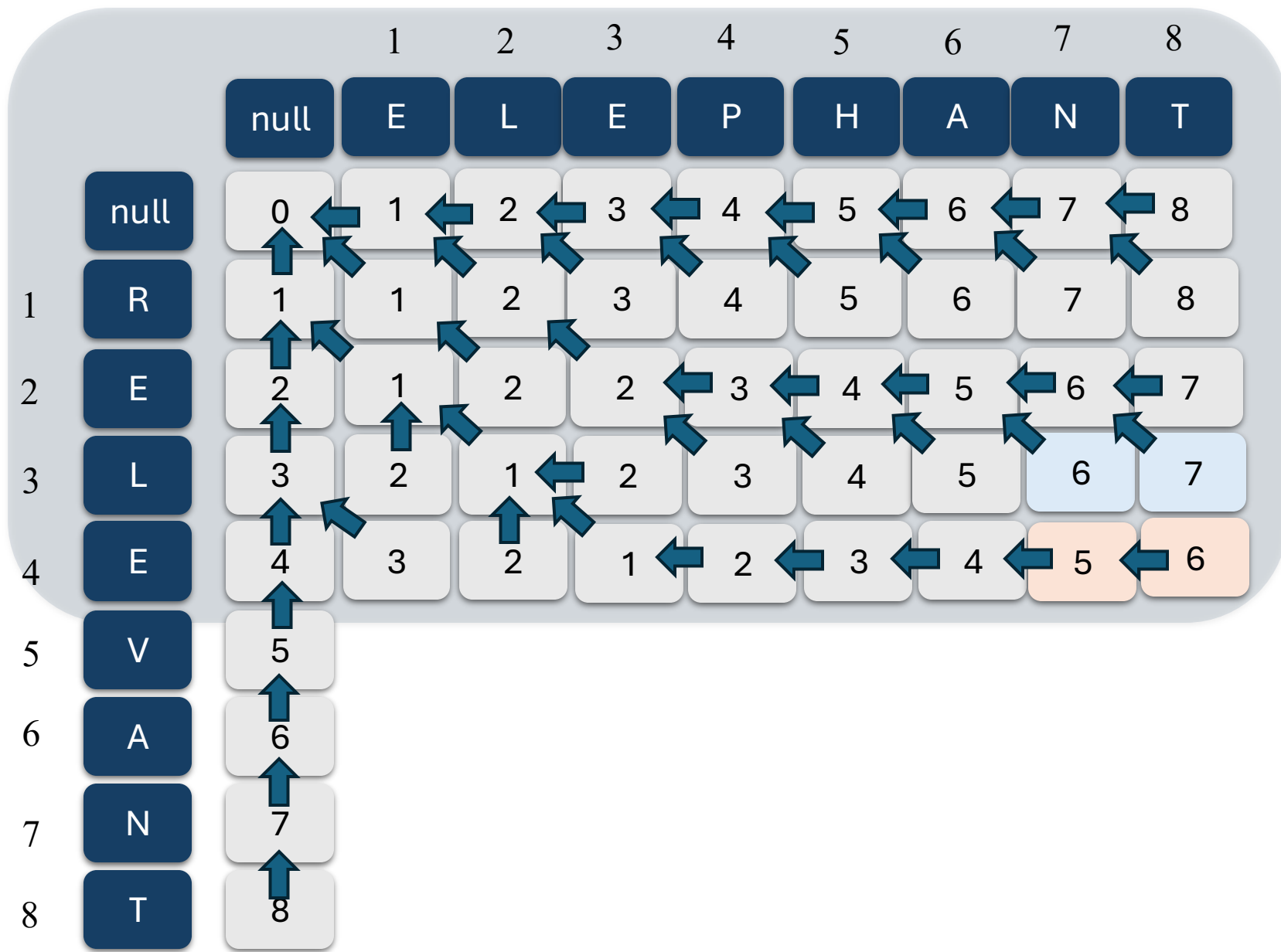
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



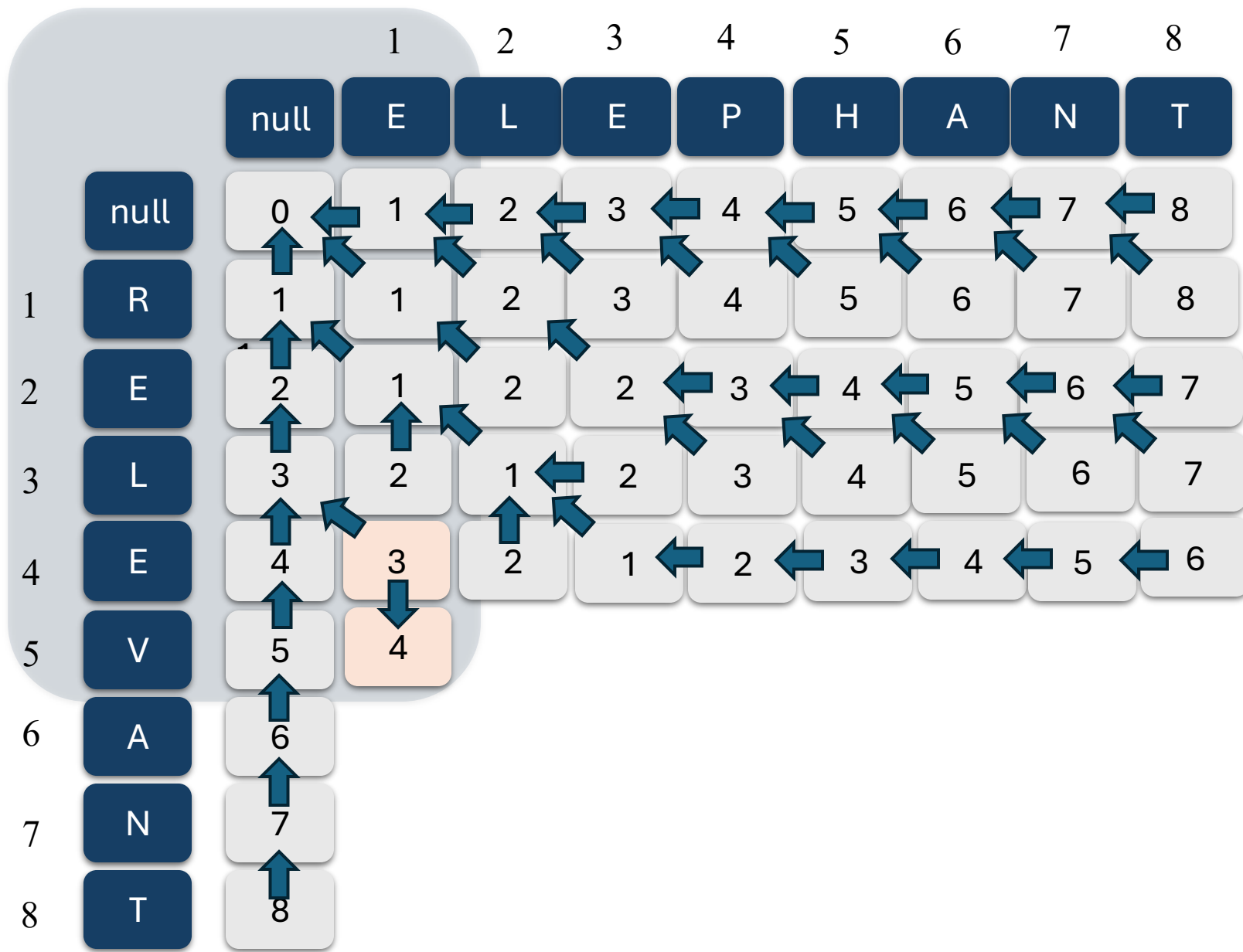
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



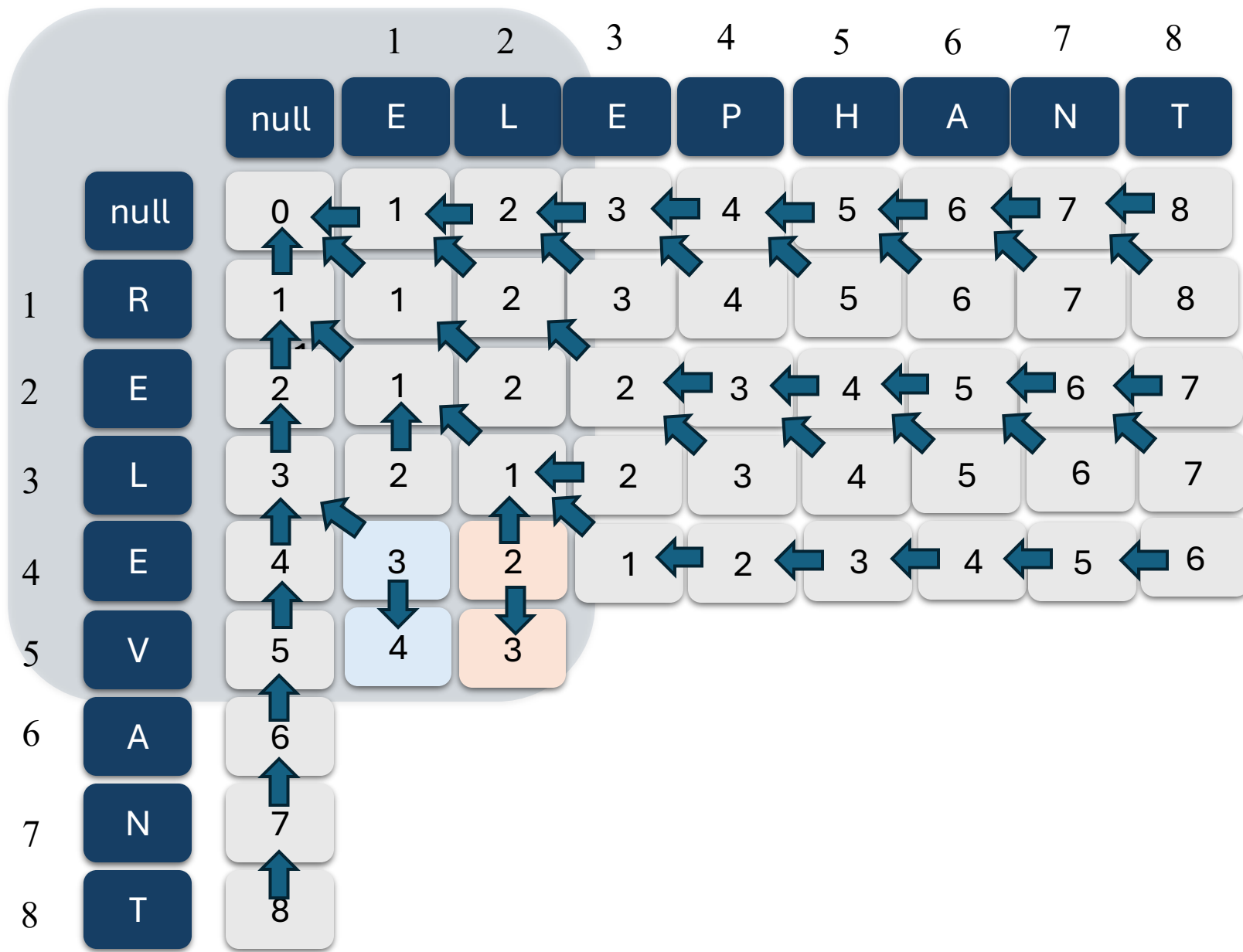
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



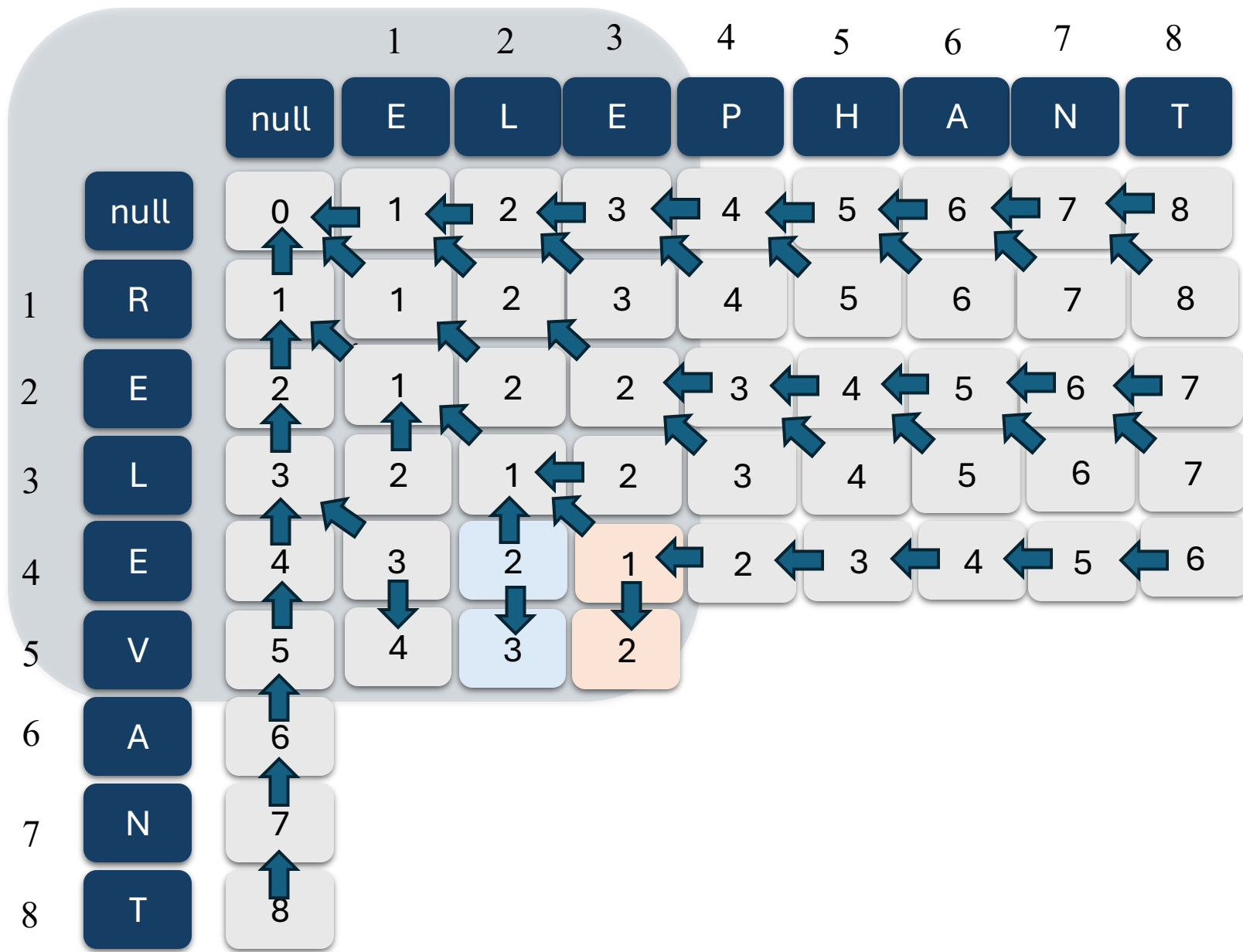
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



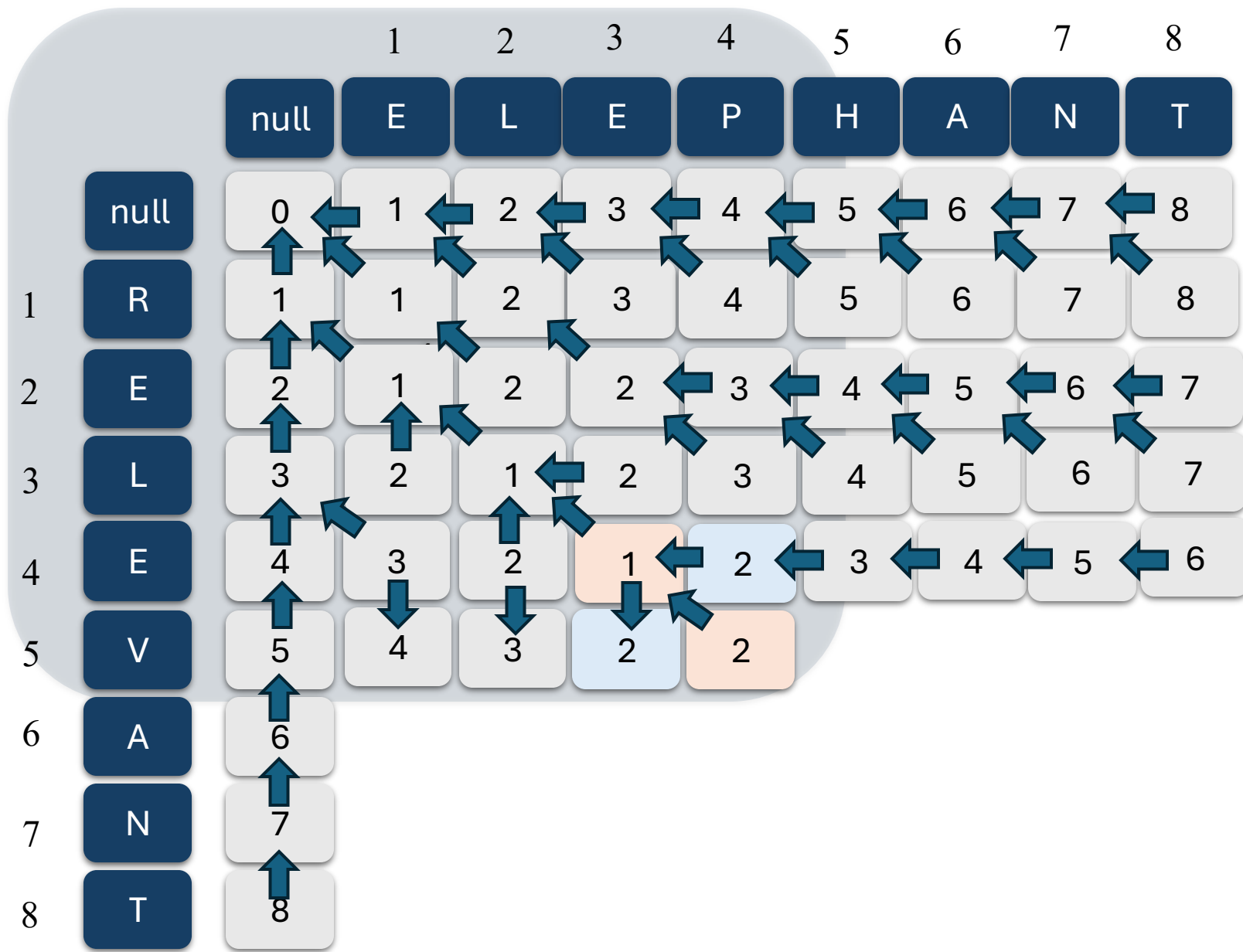
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



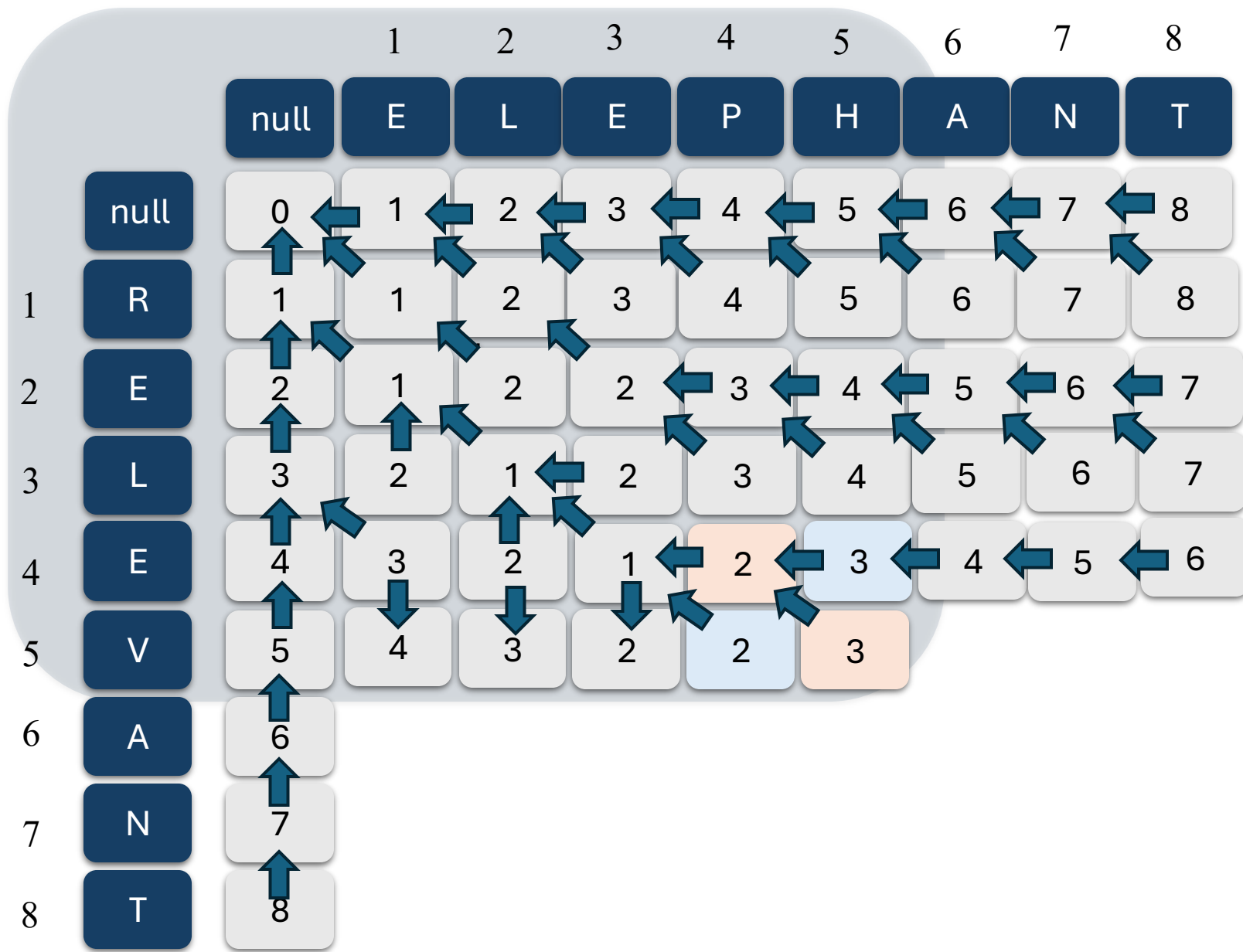
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



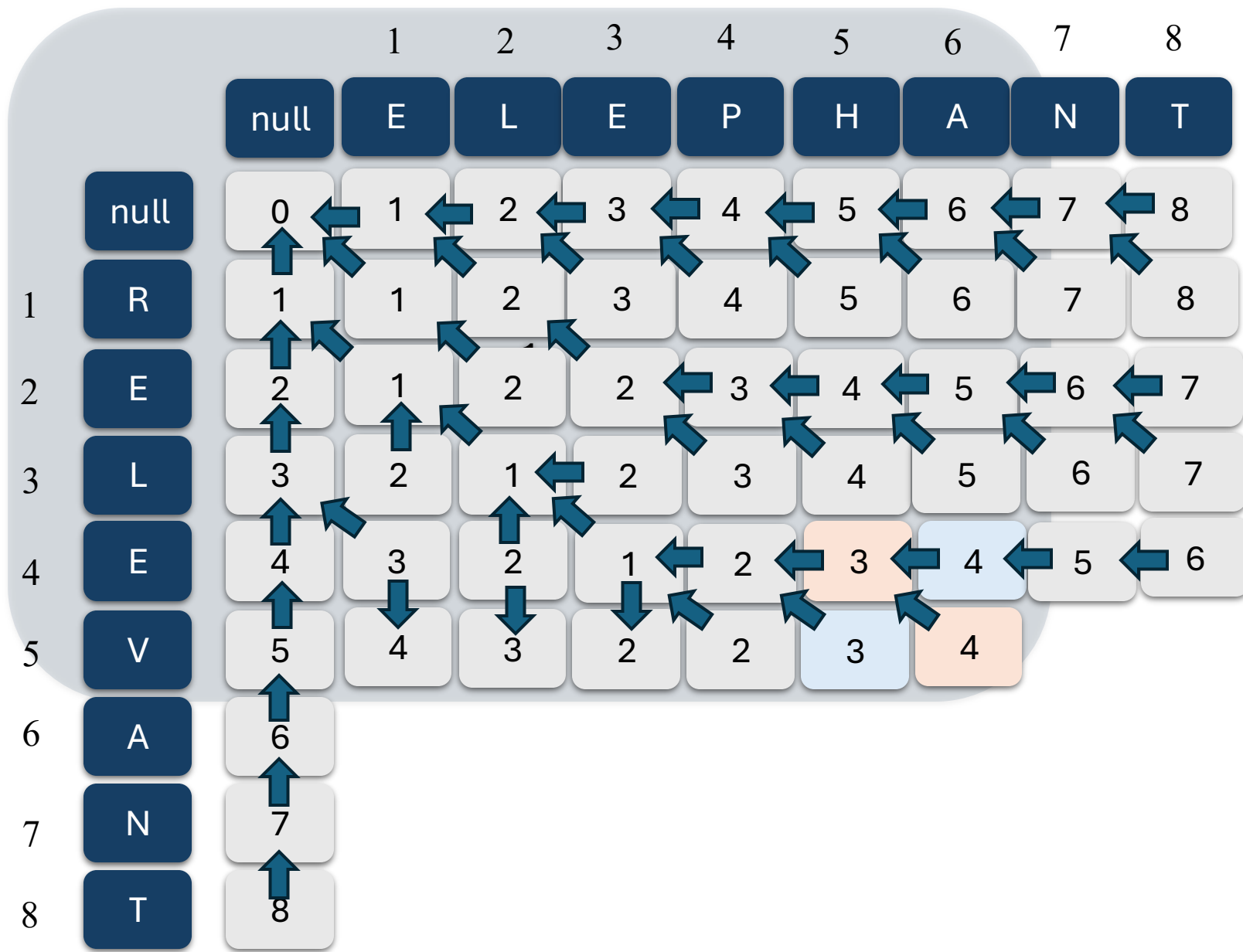
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



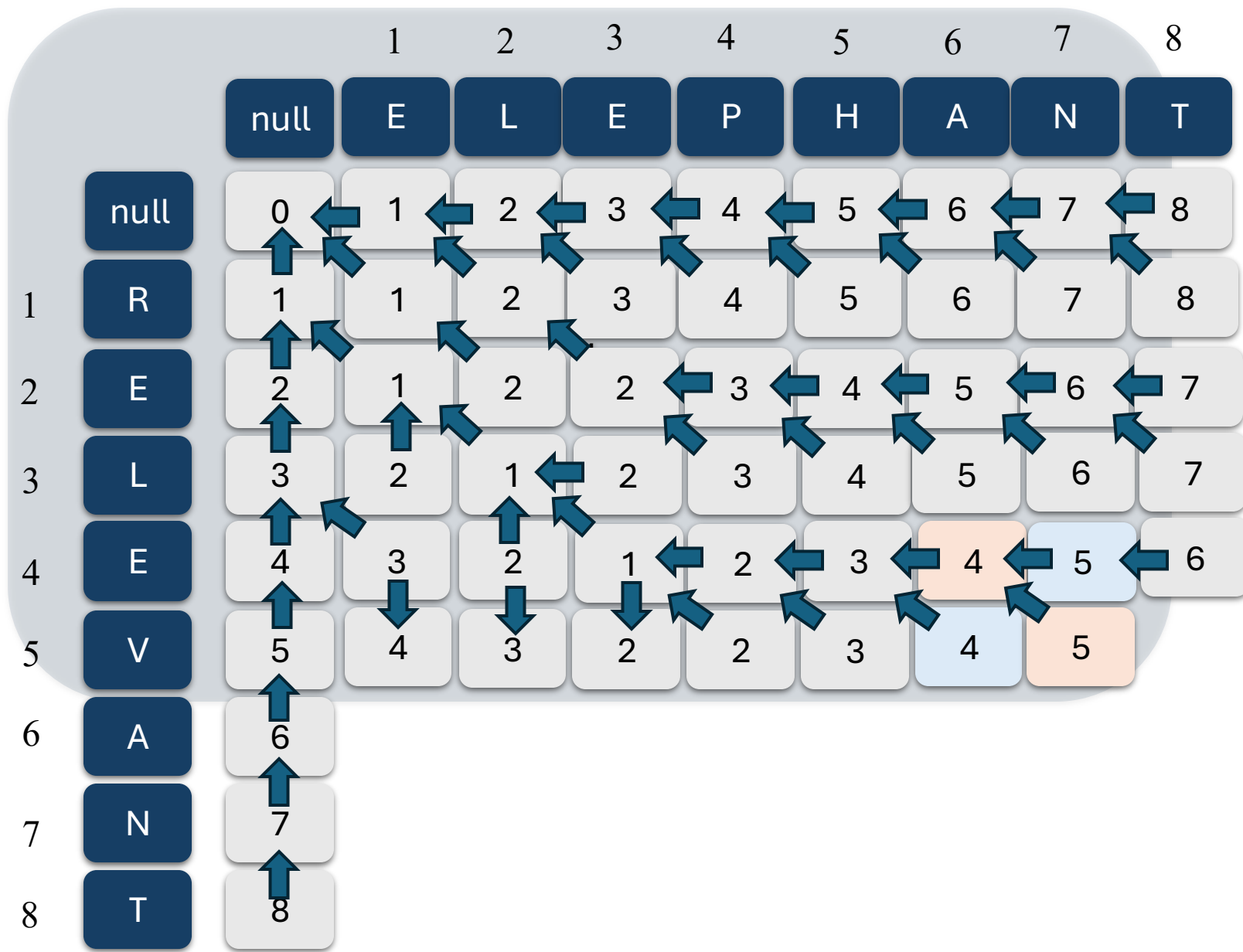
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



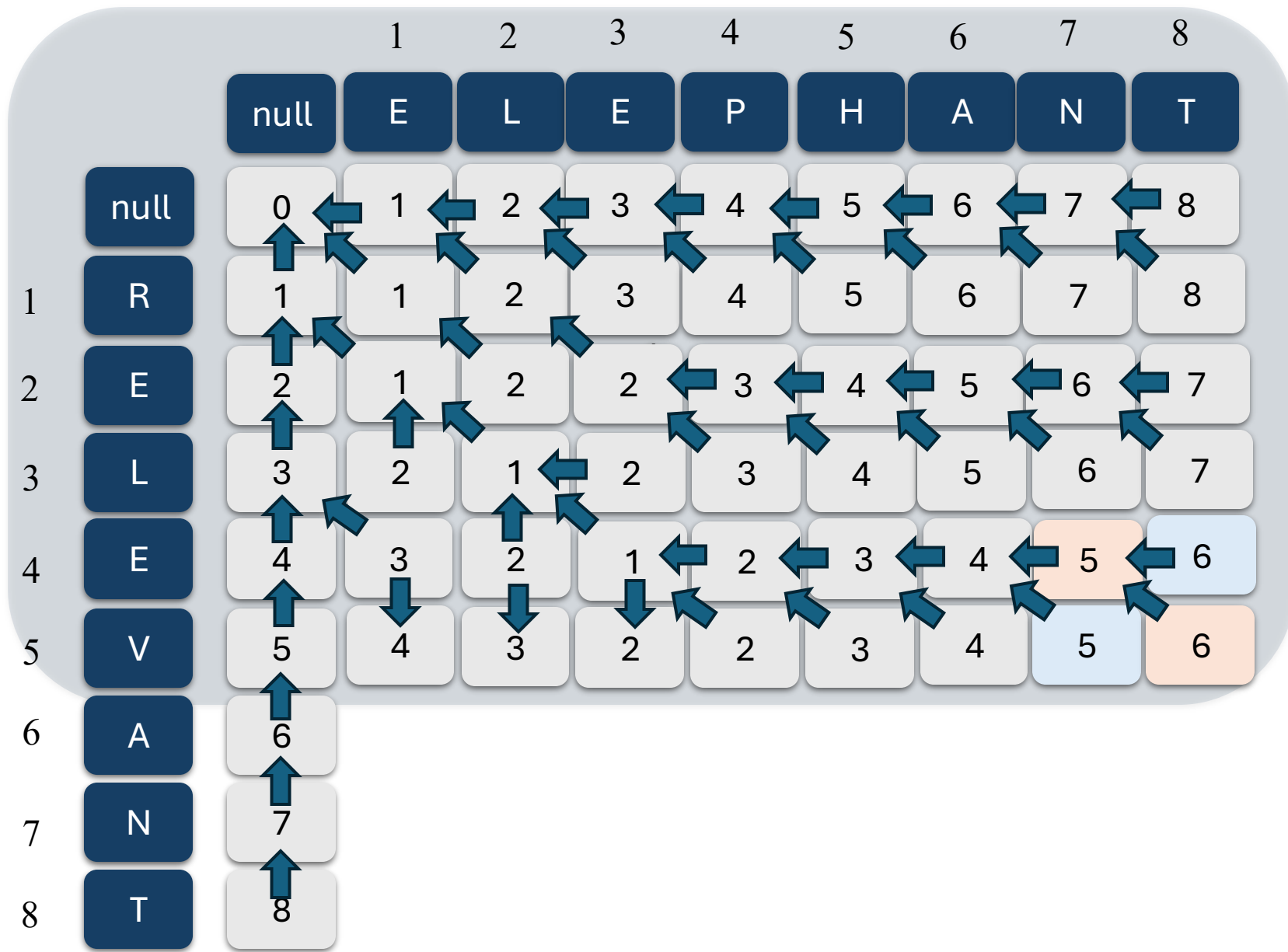
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



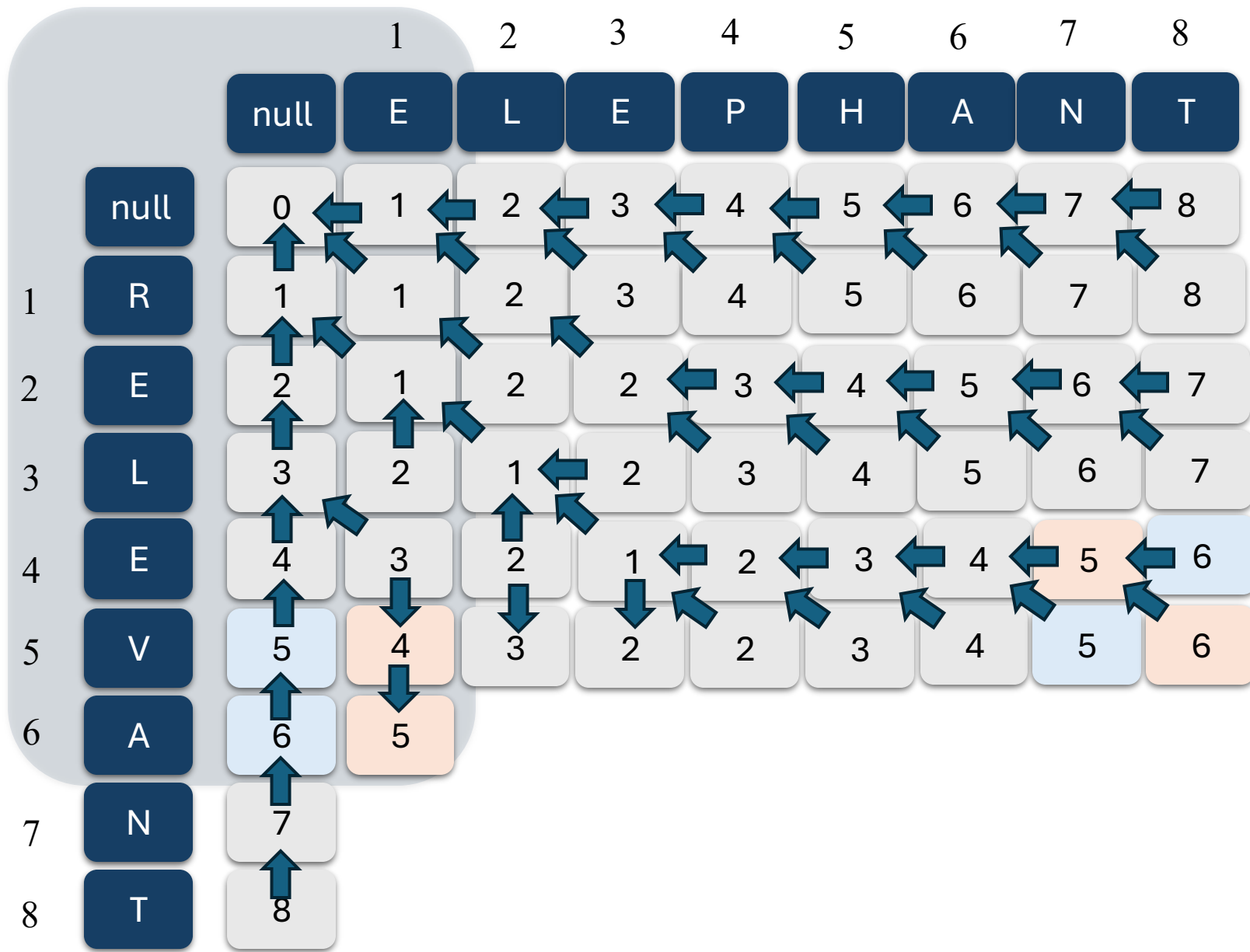
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



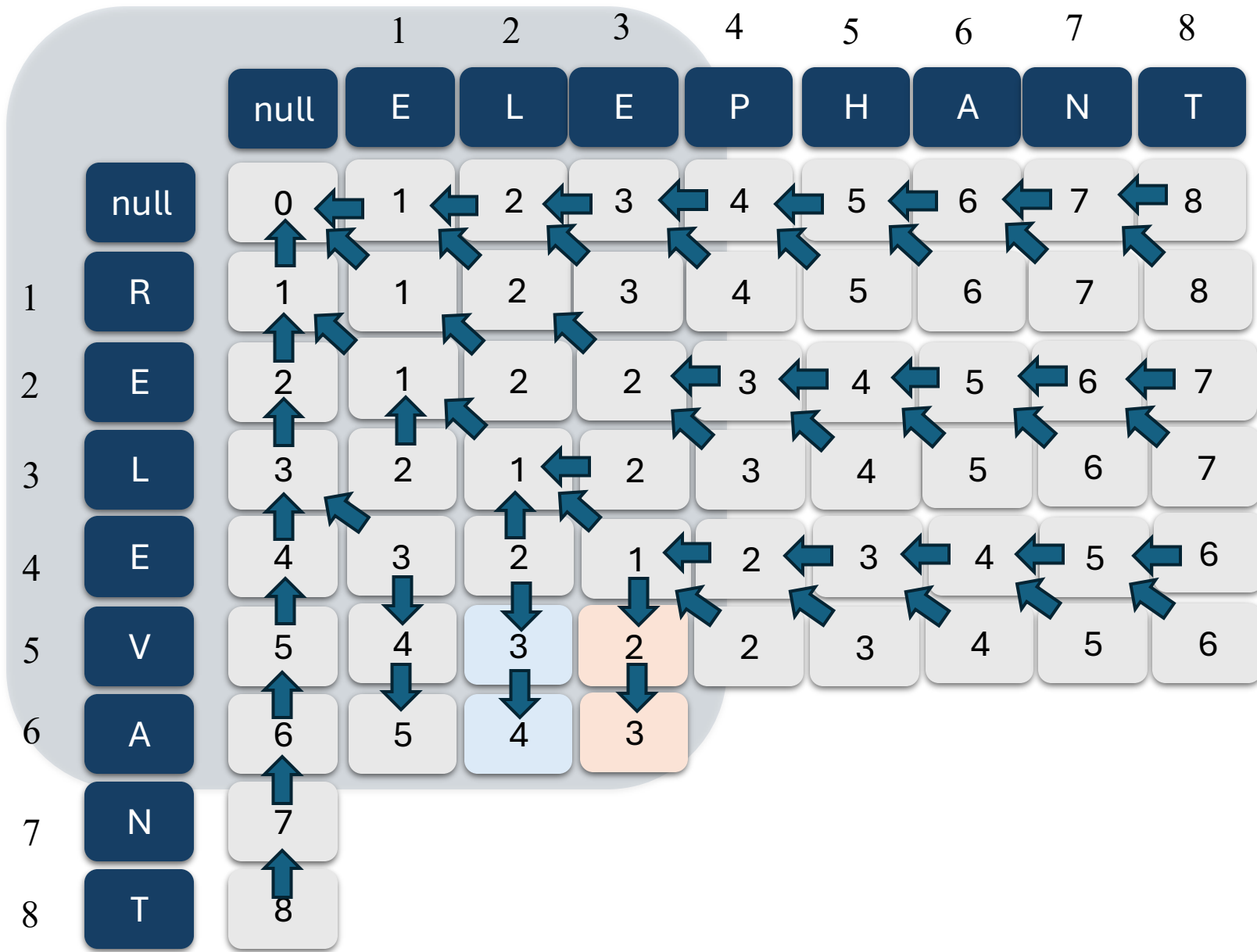
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



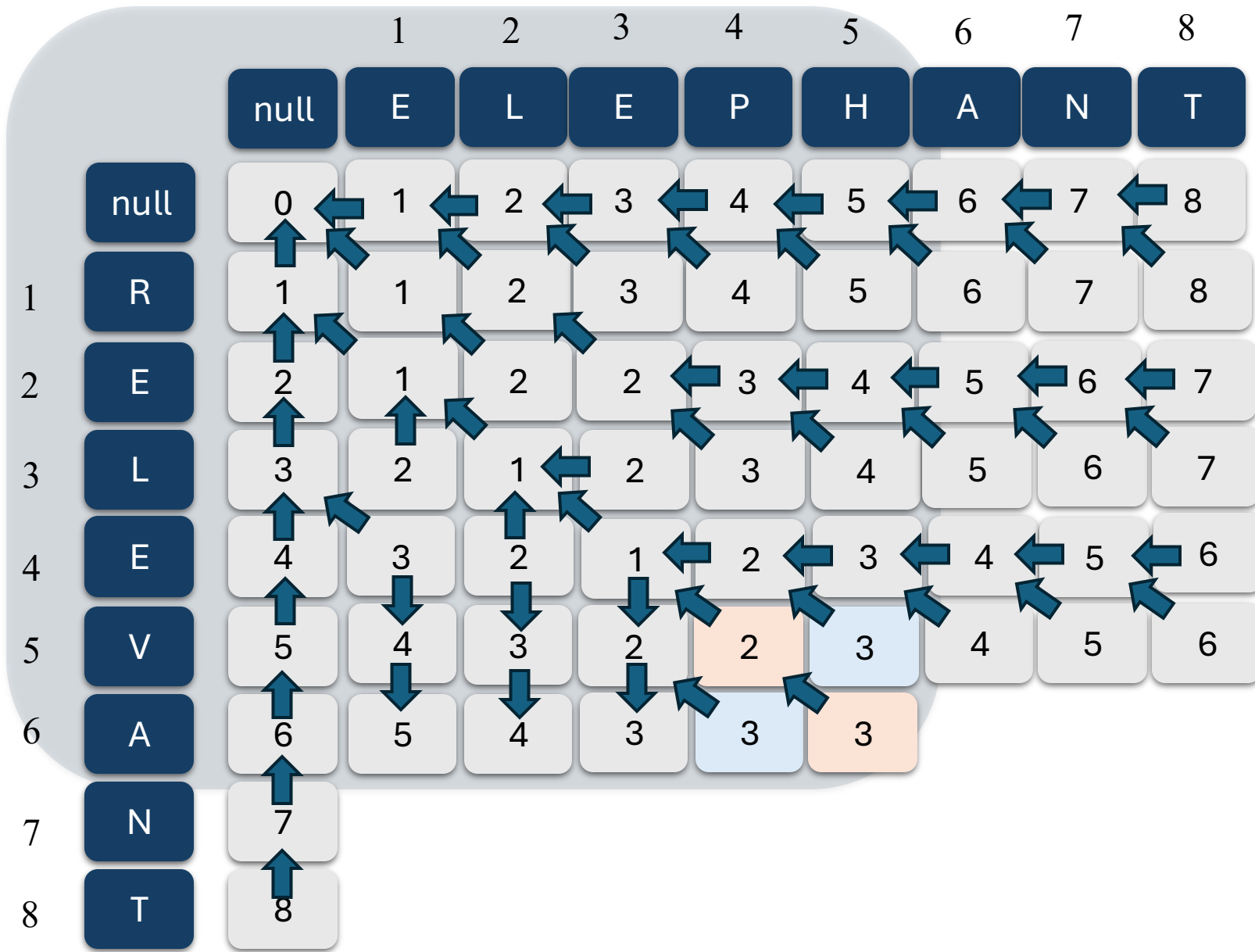
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



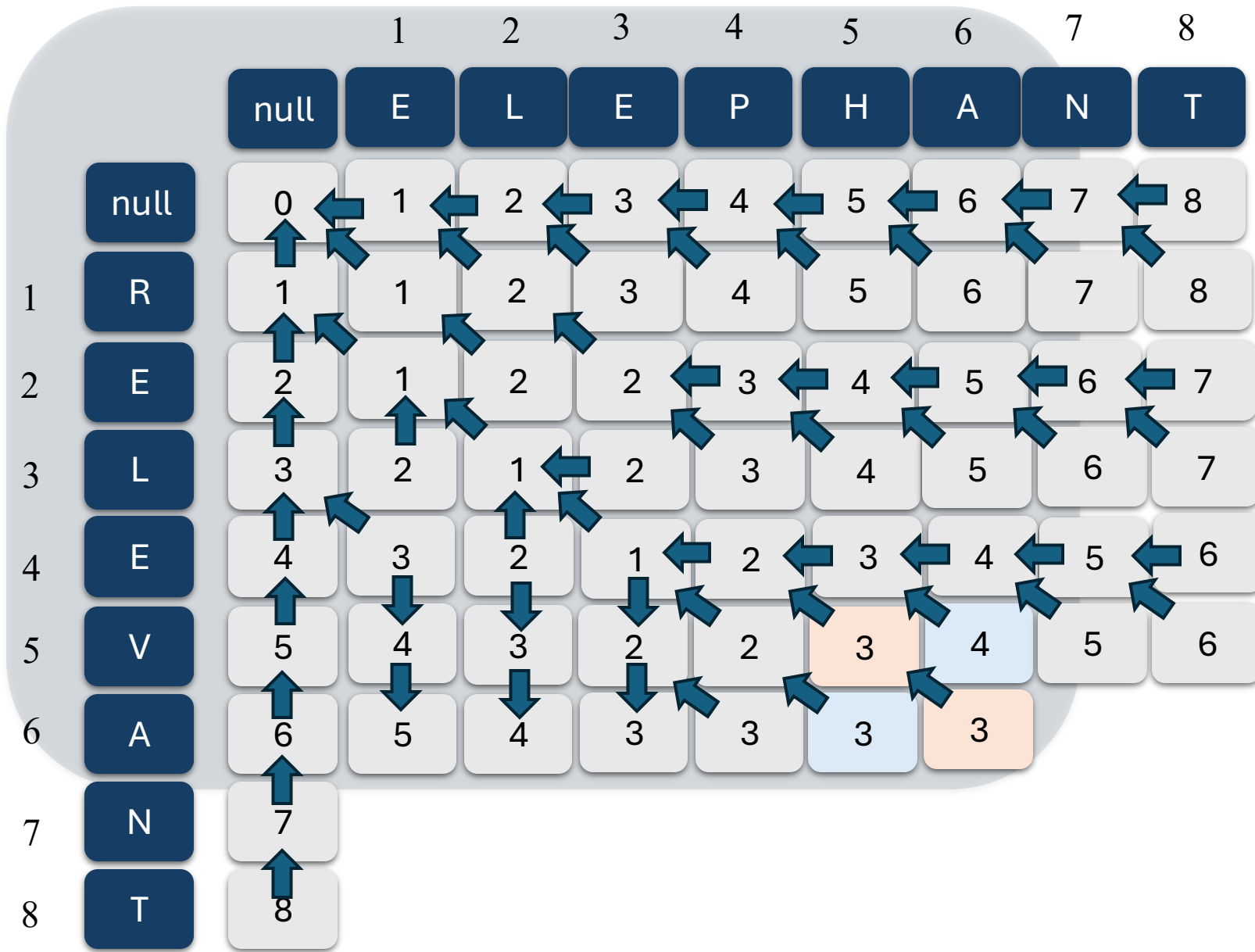
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



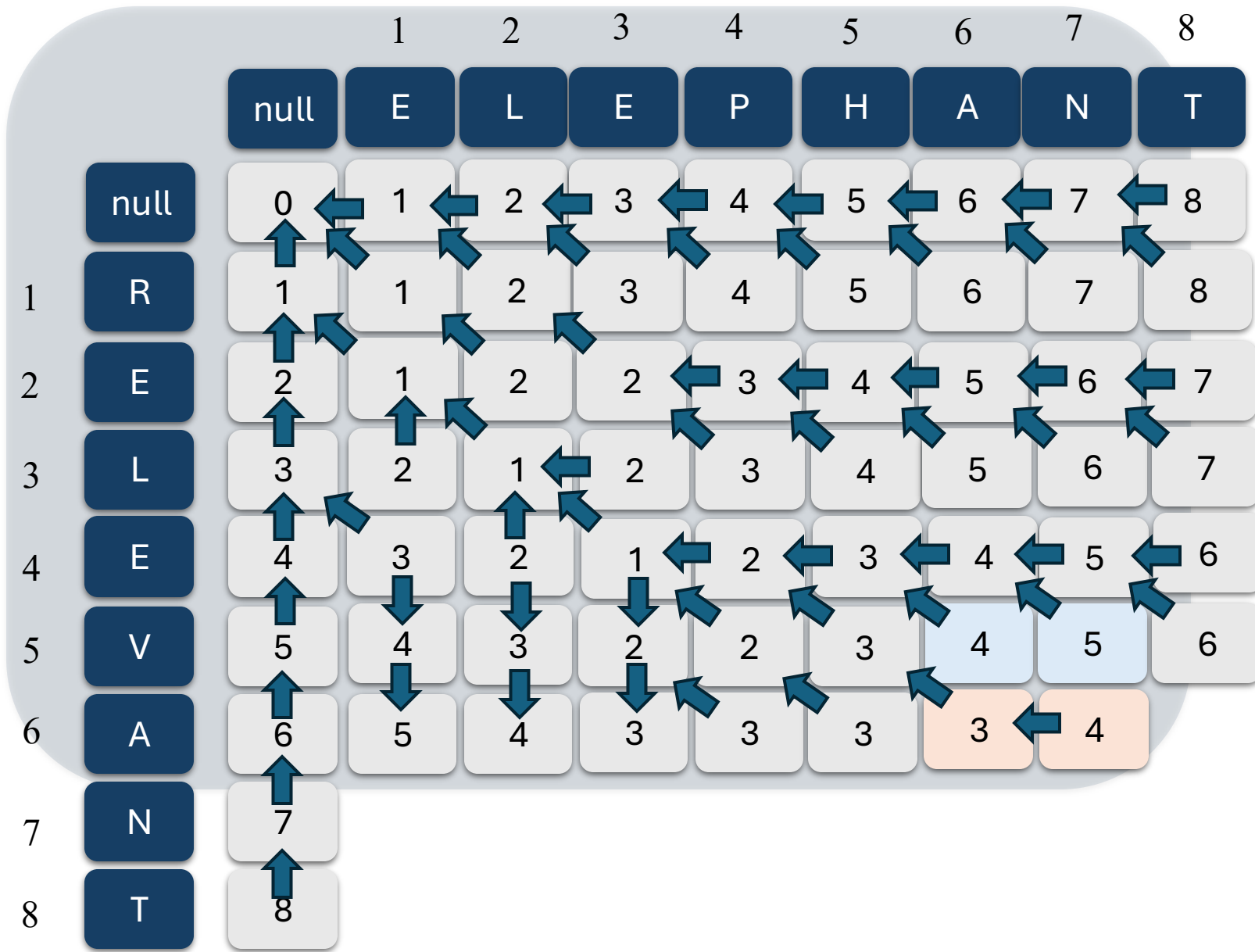
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



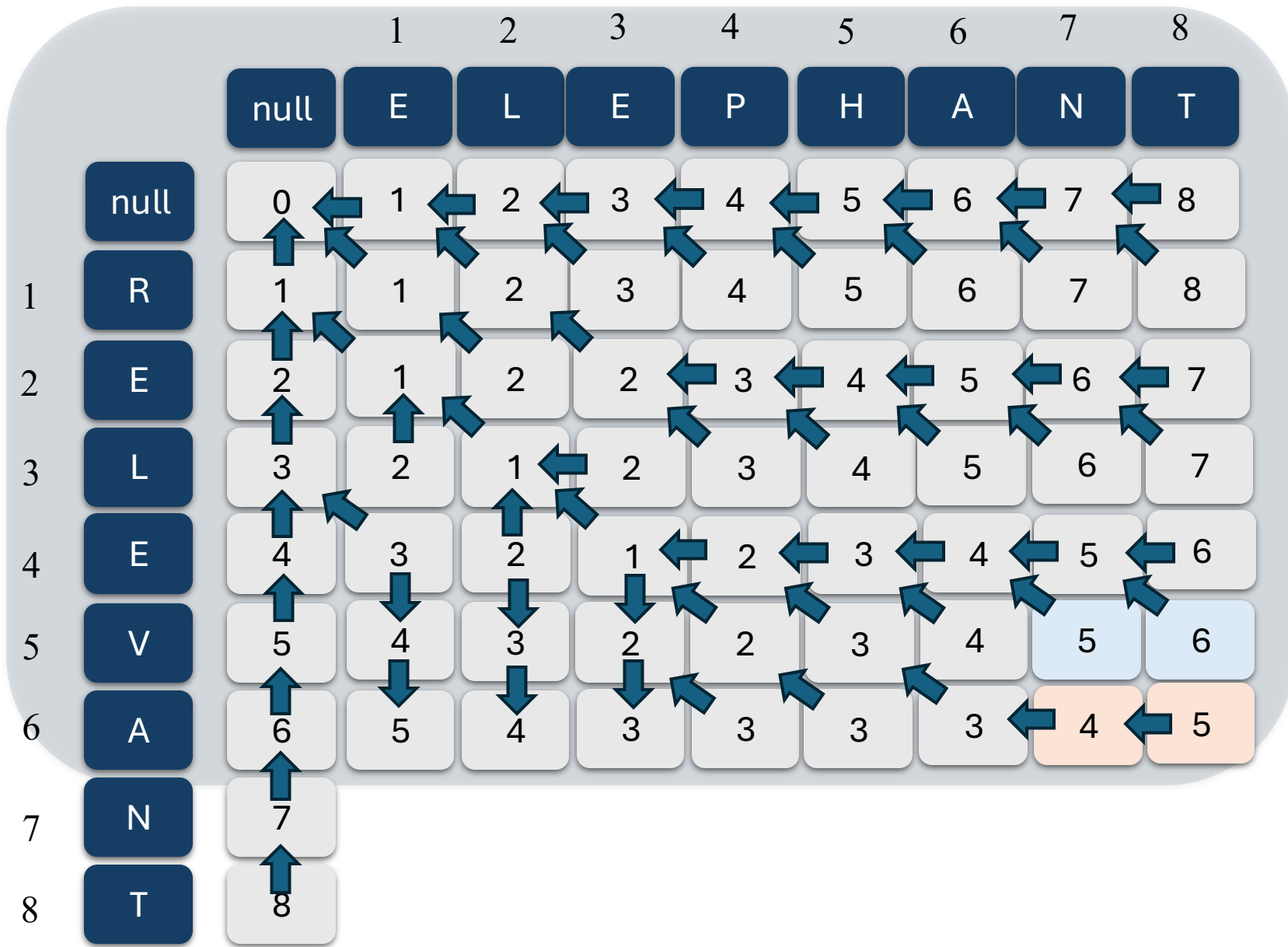
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



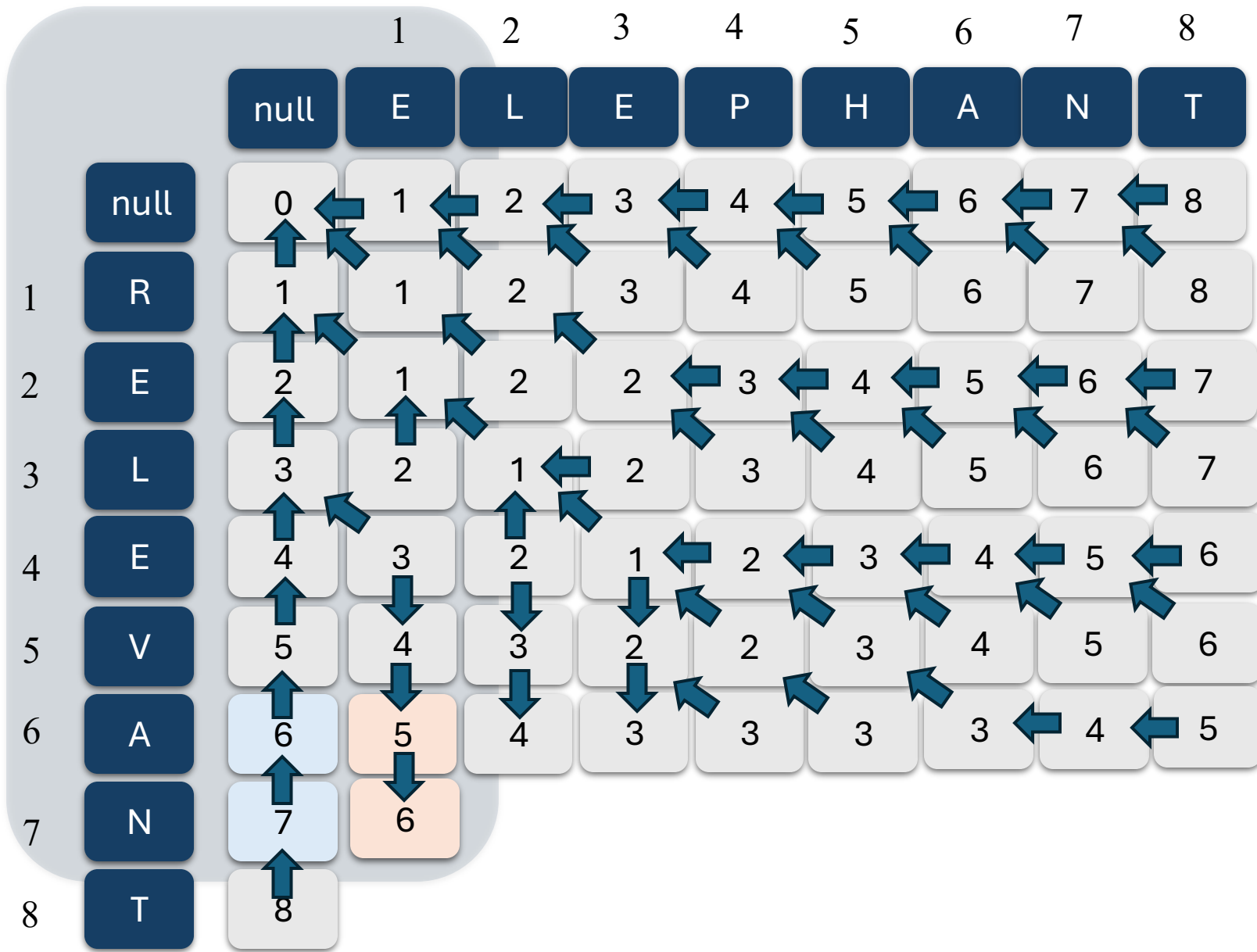
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



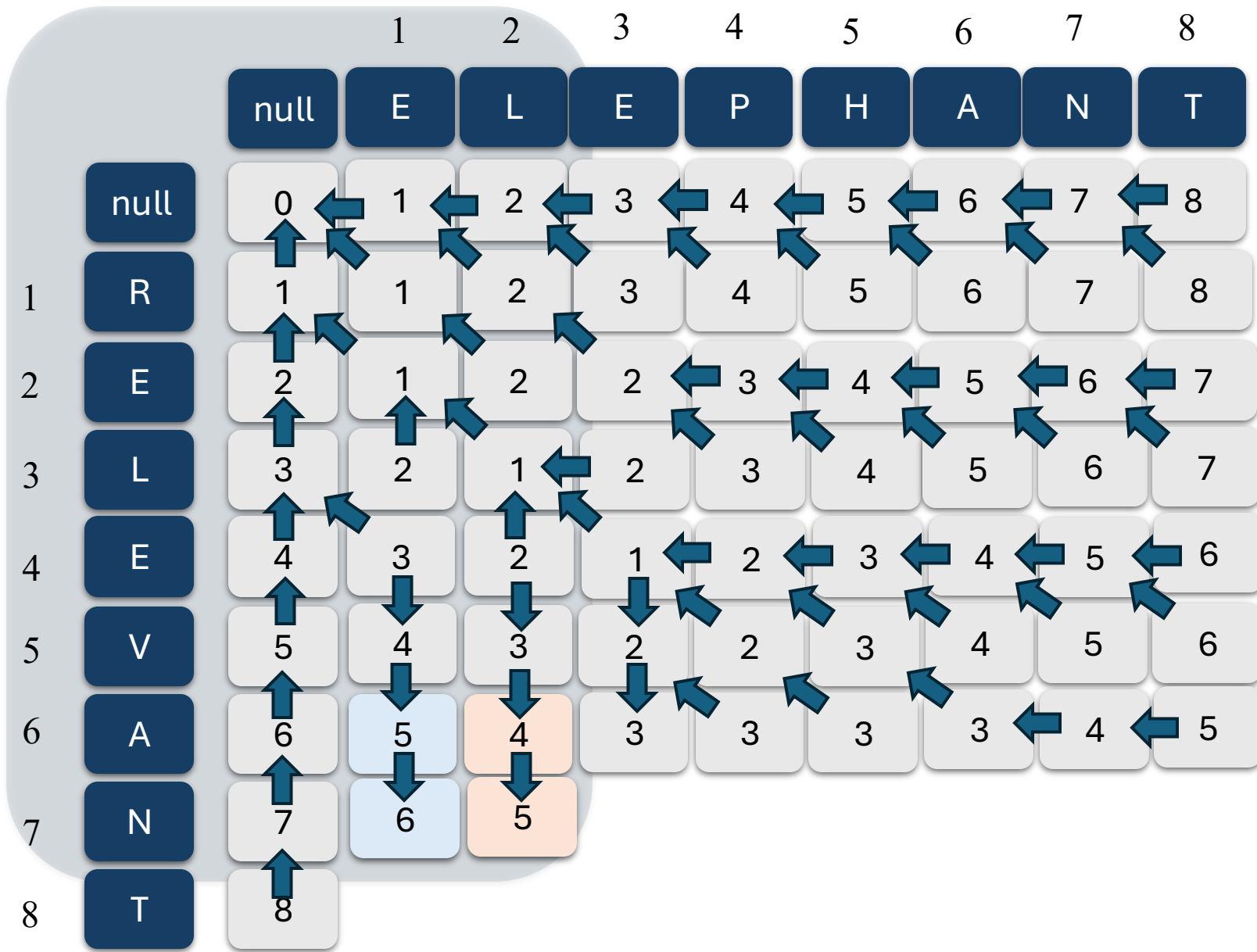
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



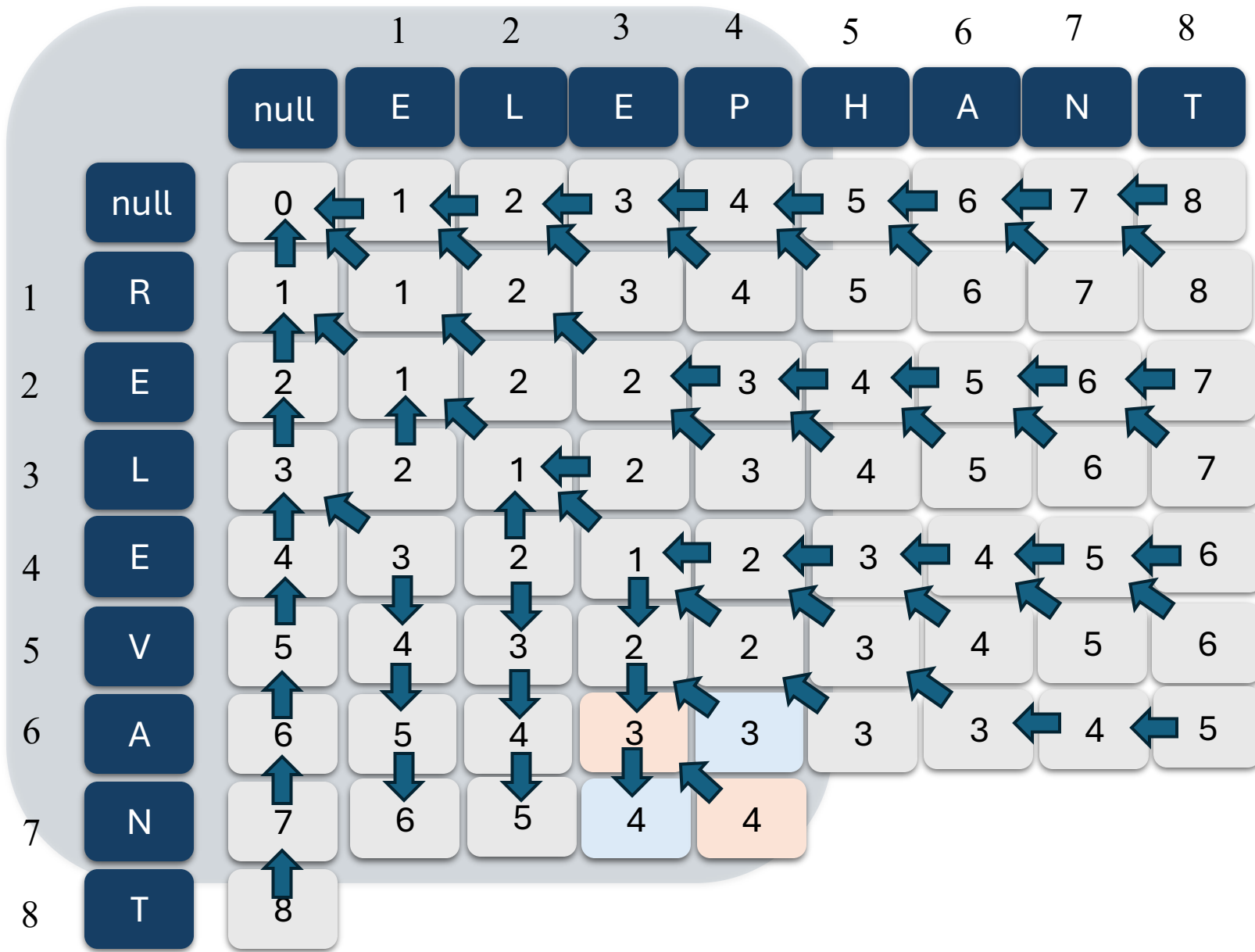
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



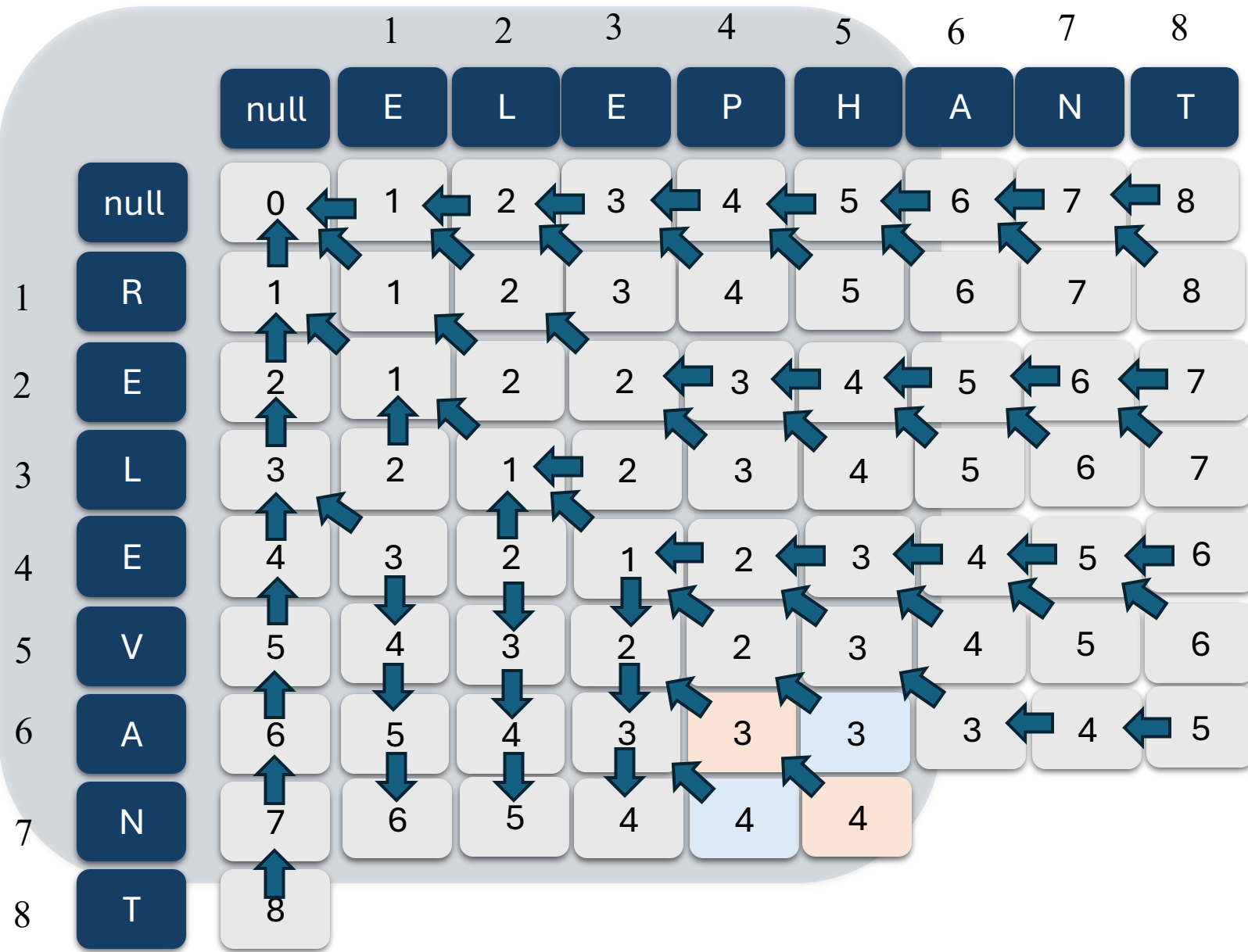
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



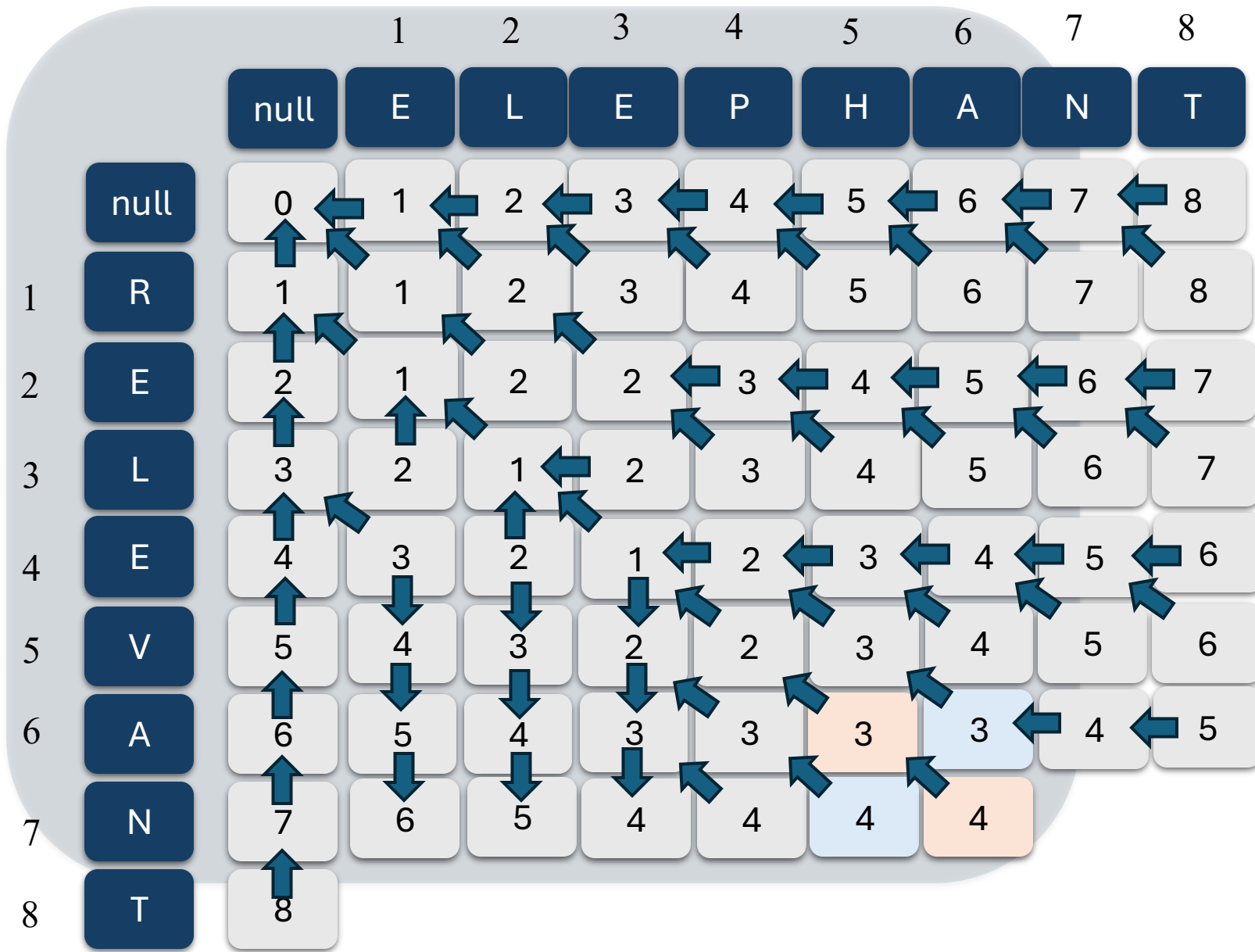
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



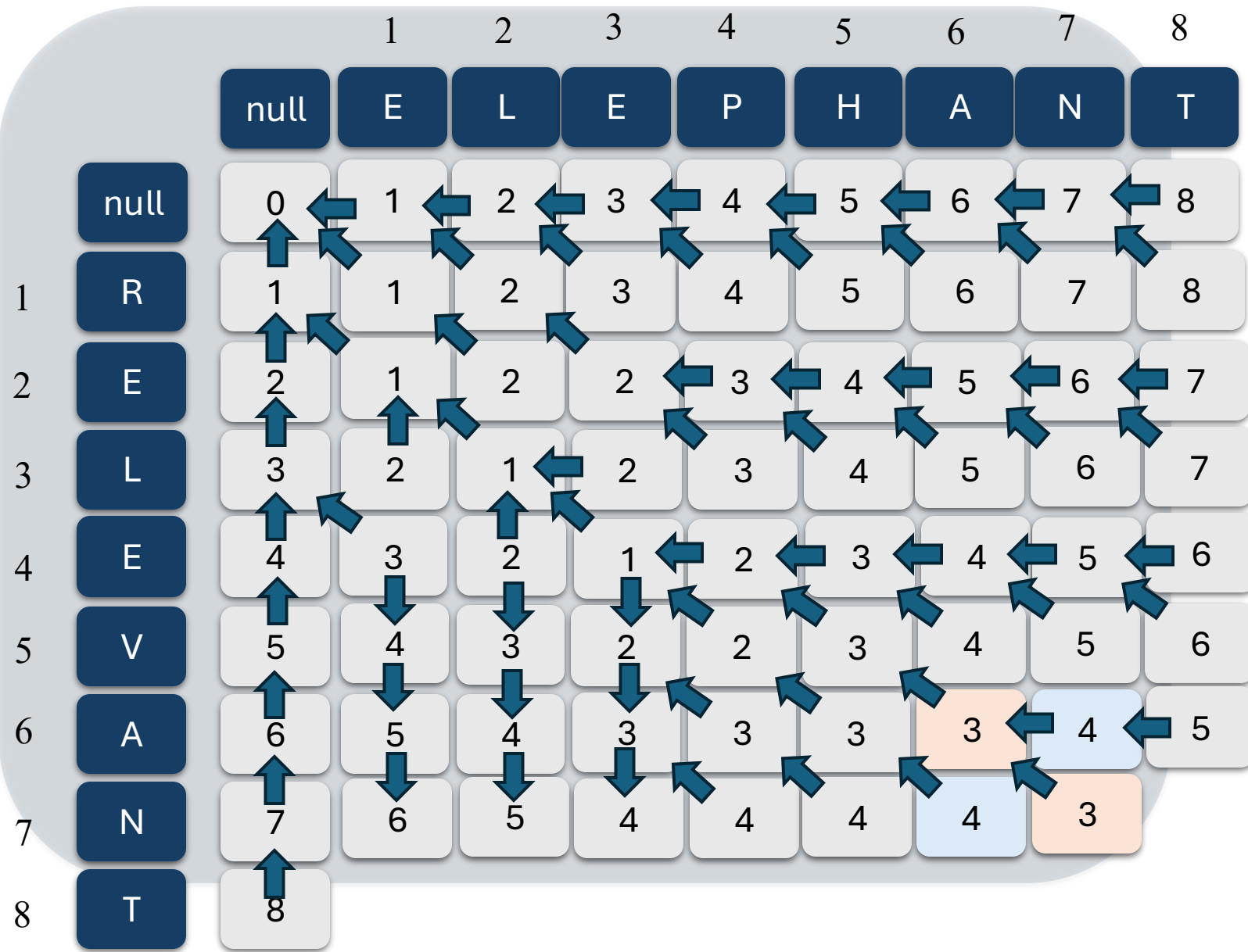
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



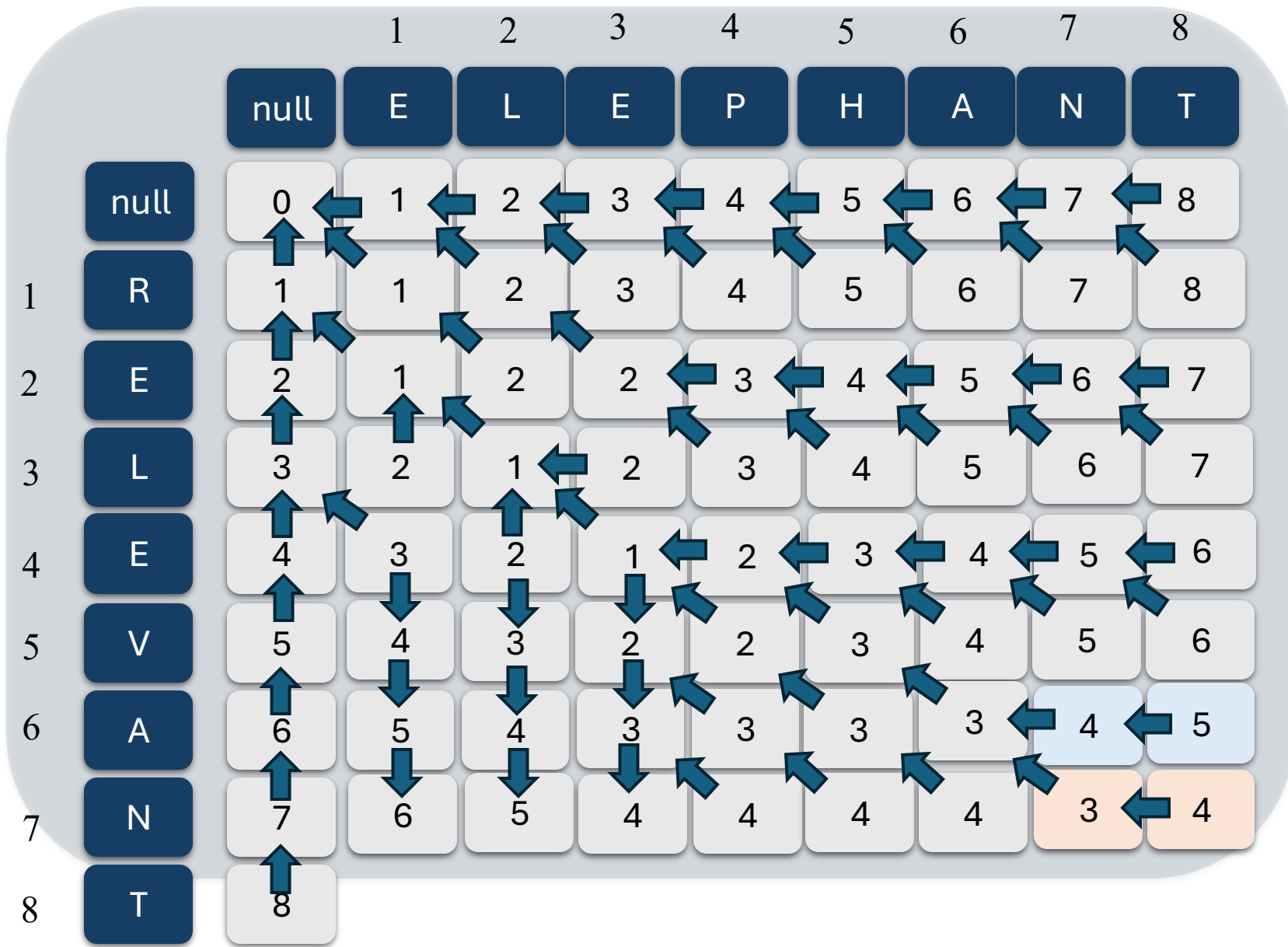
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



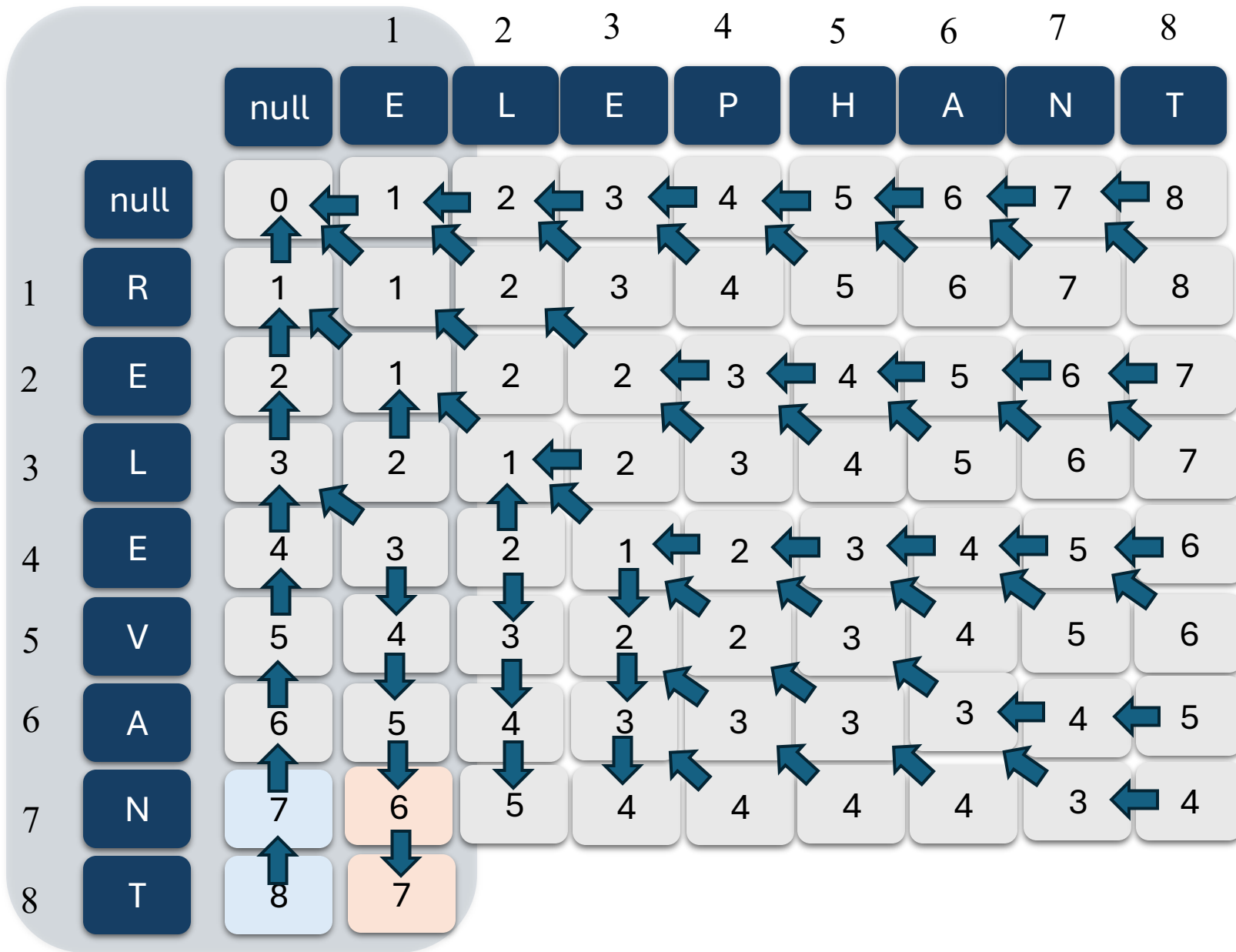
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



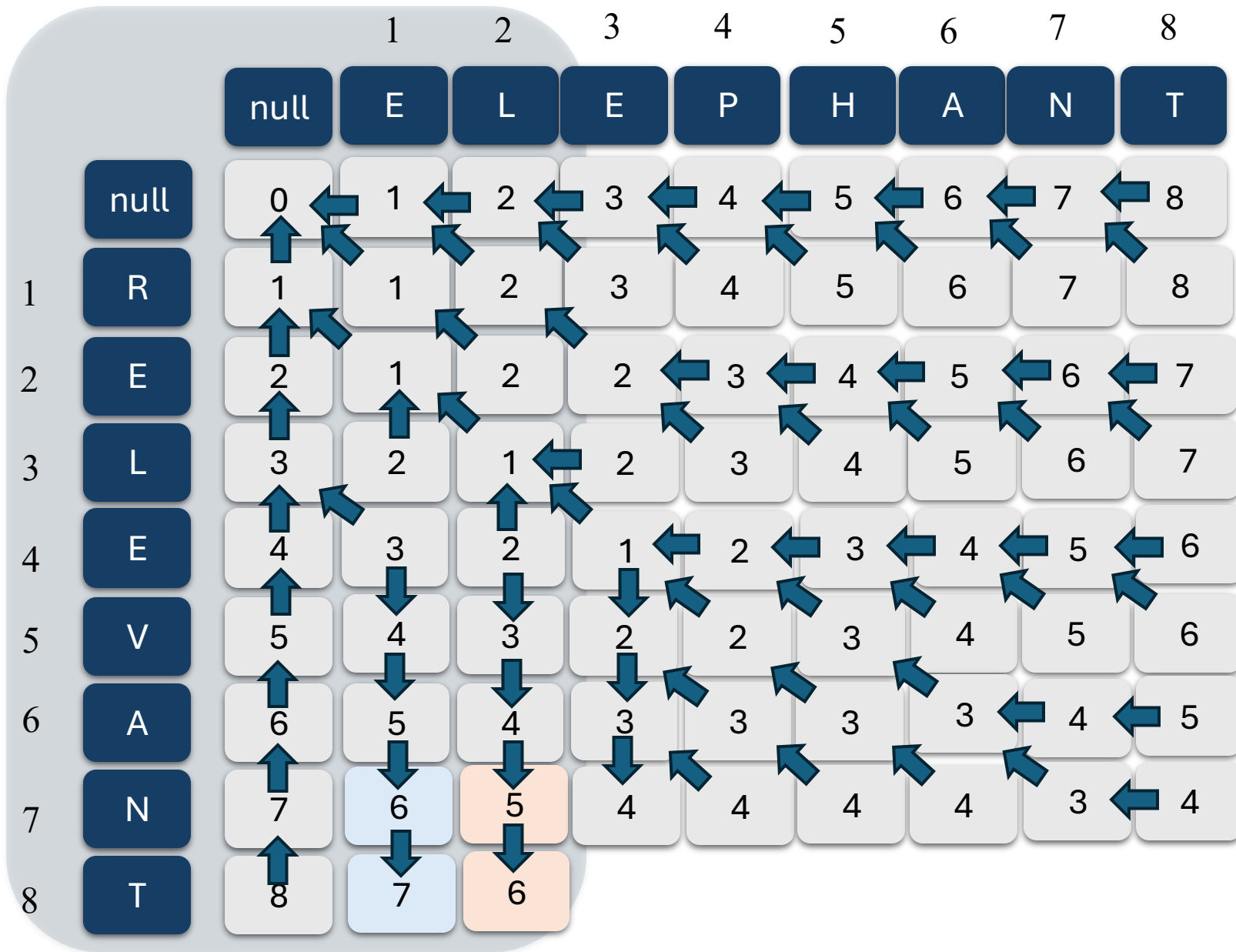
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



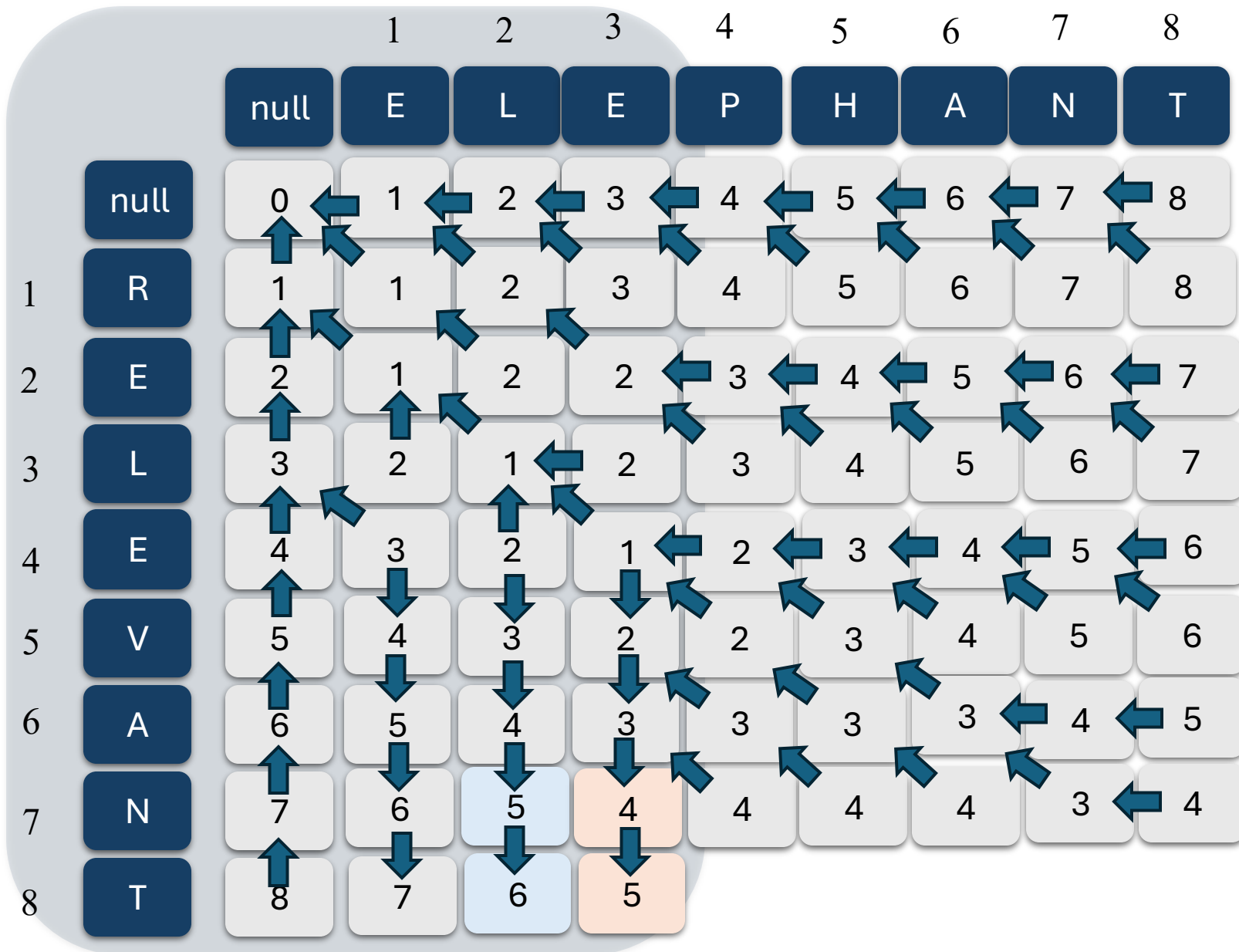
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



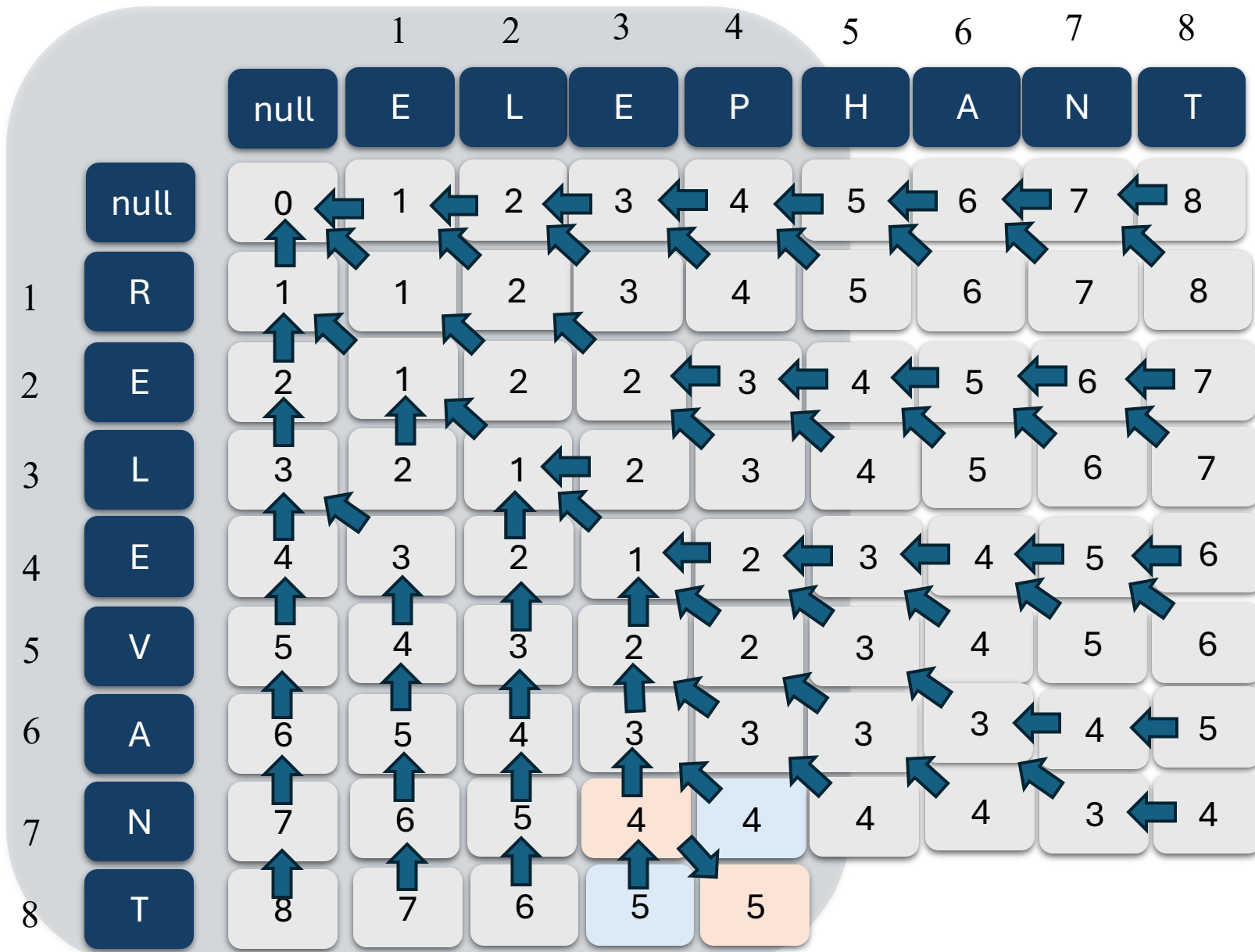
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



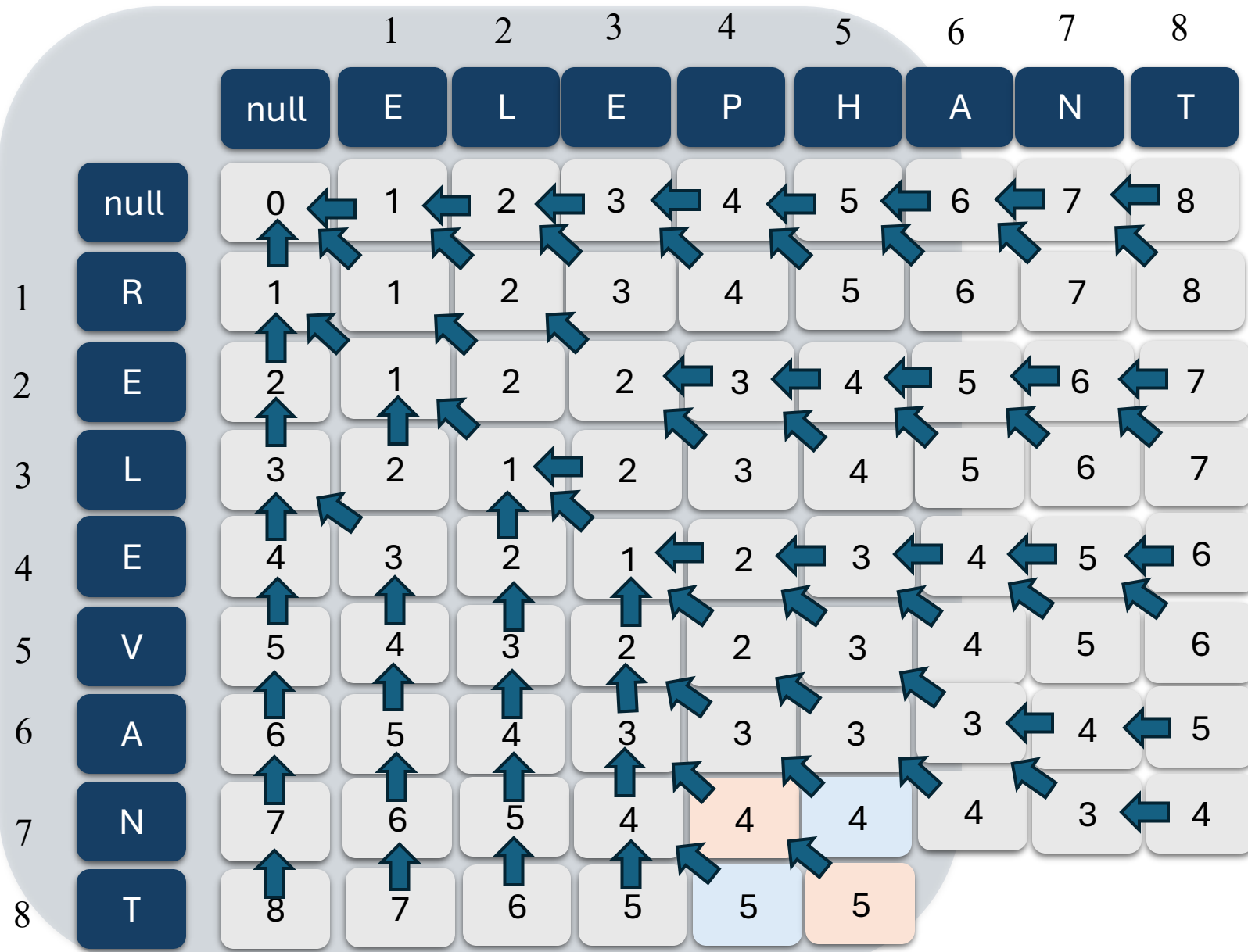
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



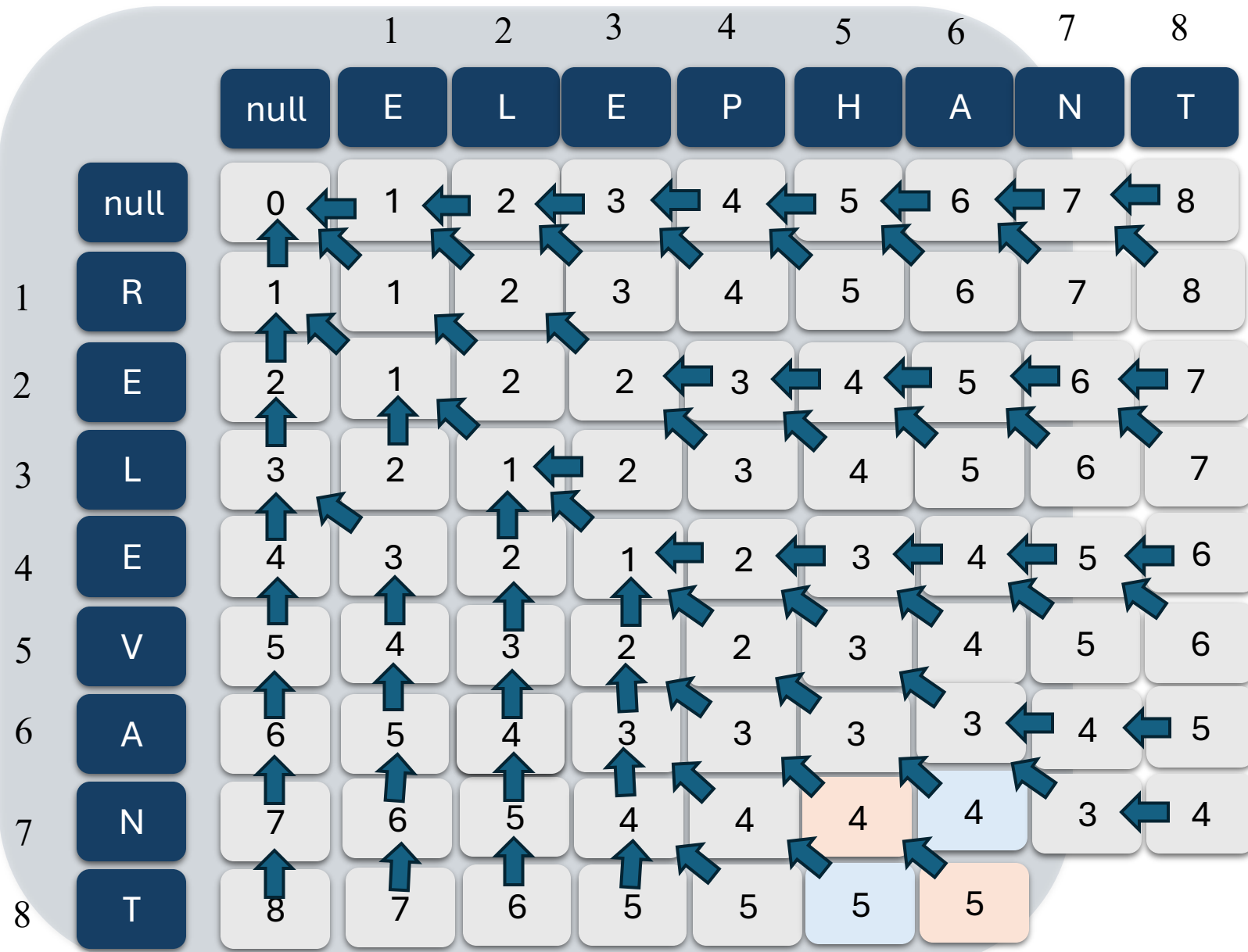
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



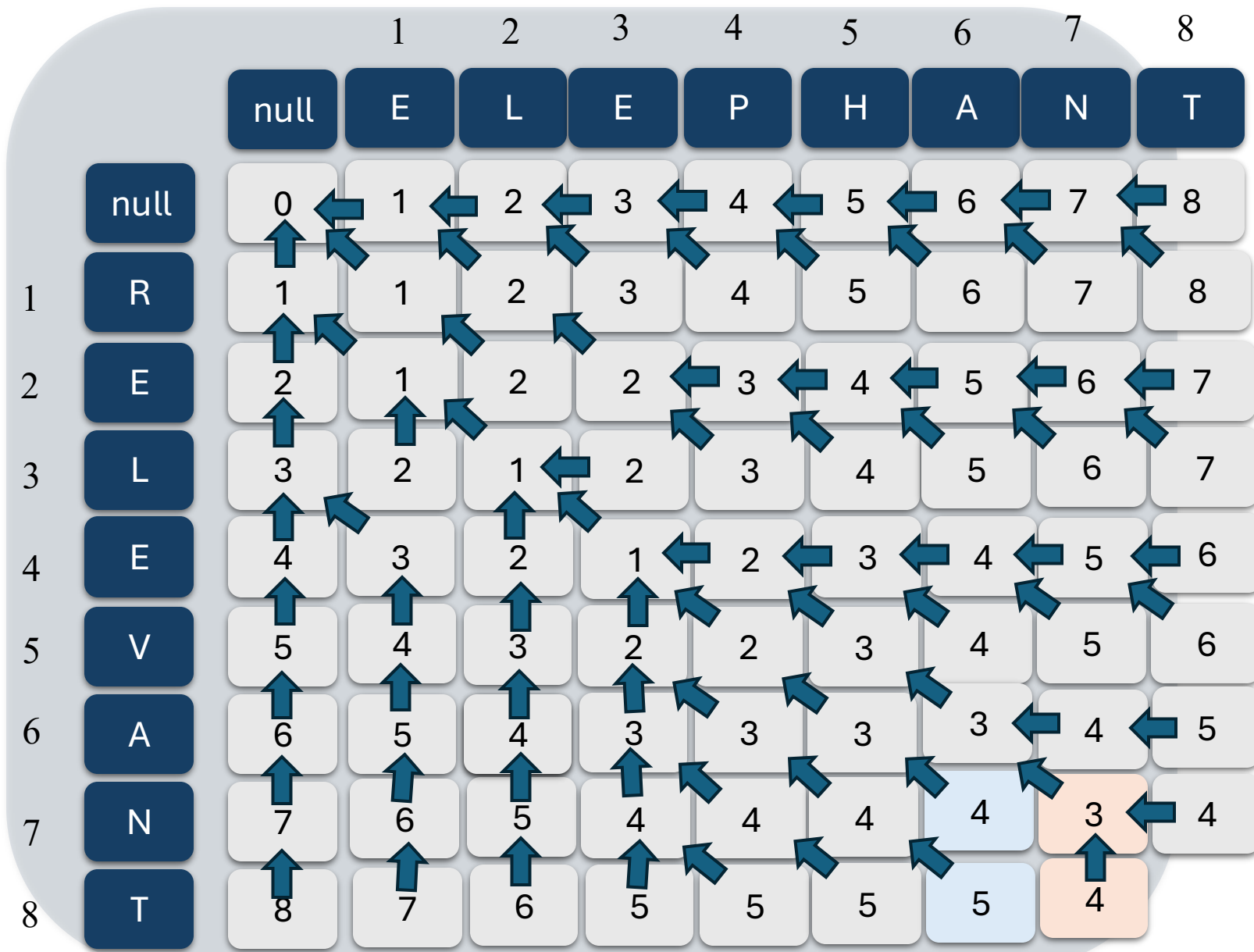
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



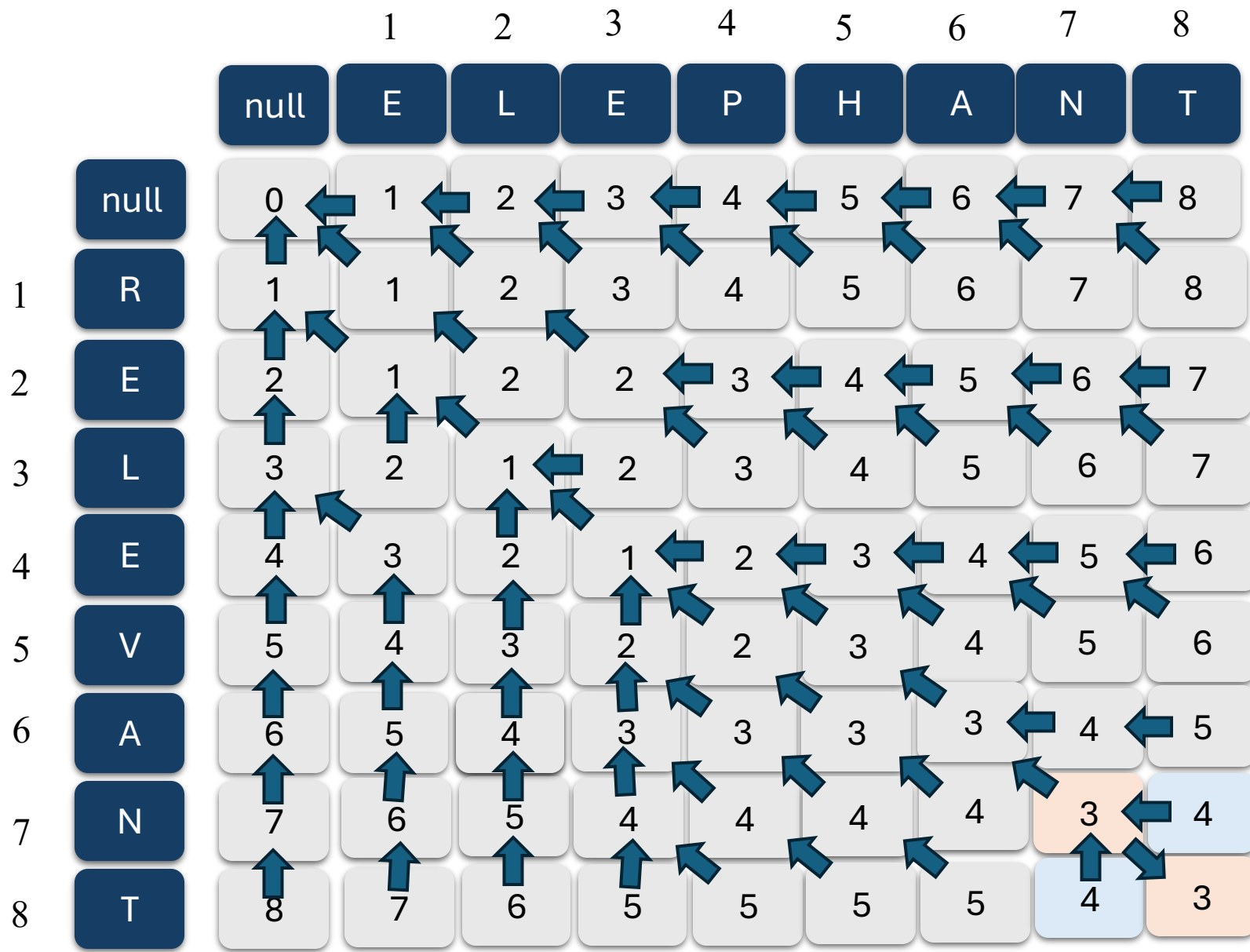
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



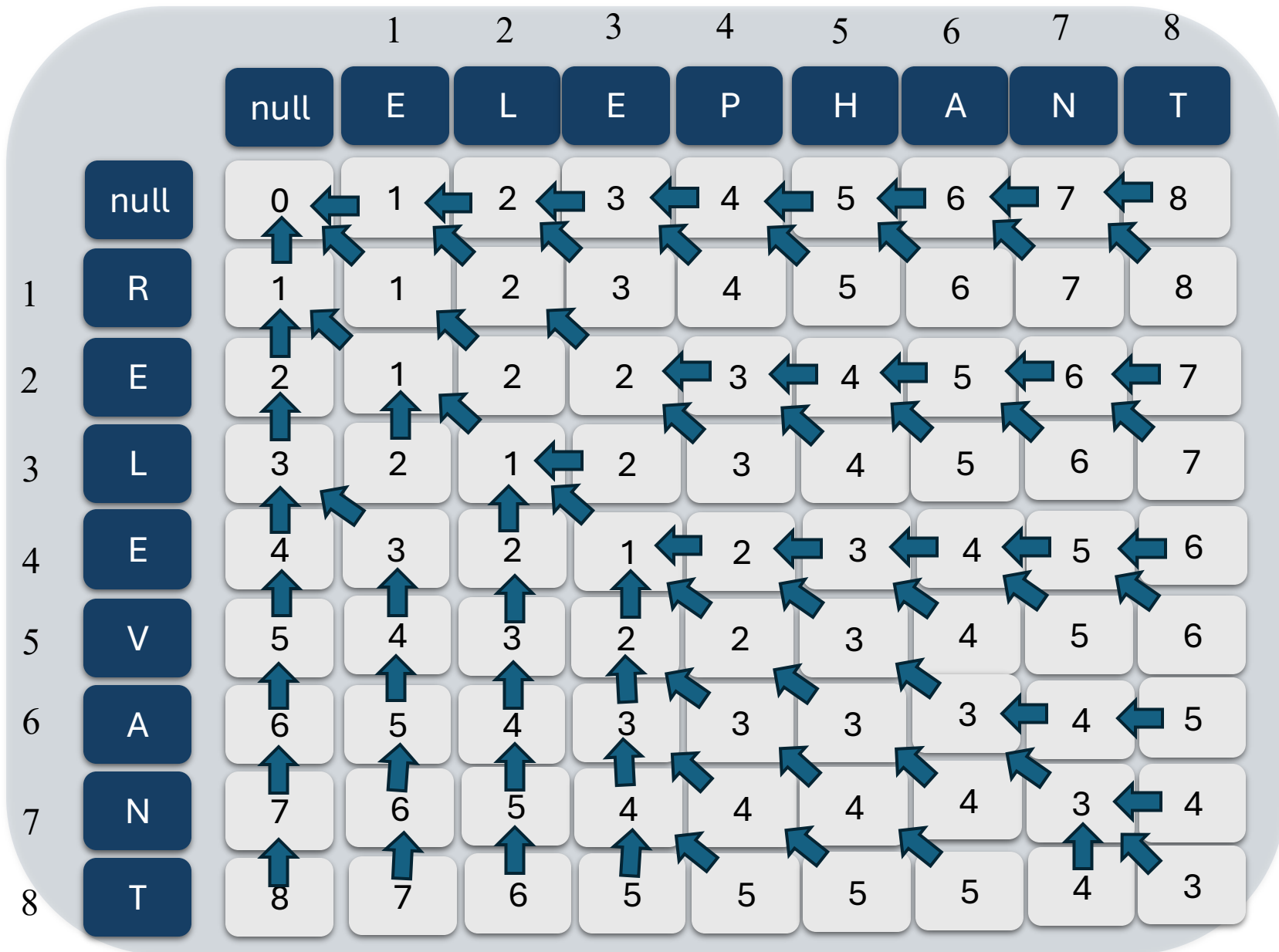
$$E(i, j) = \min \begin{cases} 1 + E(i-1, j) \\ 1 + E(i, j-1) \\ \text{diff}(i, j) + E(i-1, j-1) \end{cases}$$



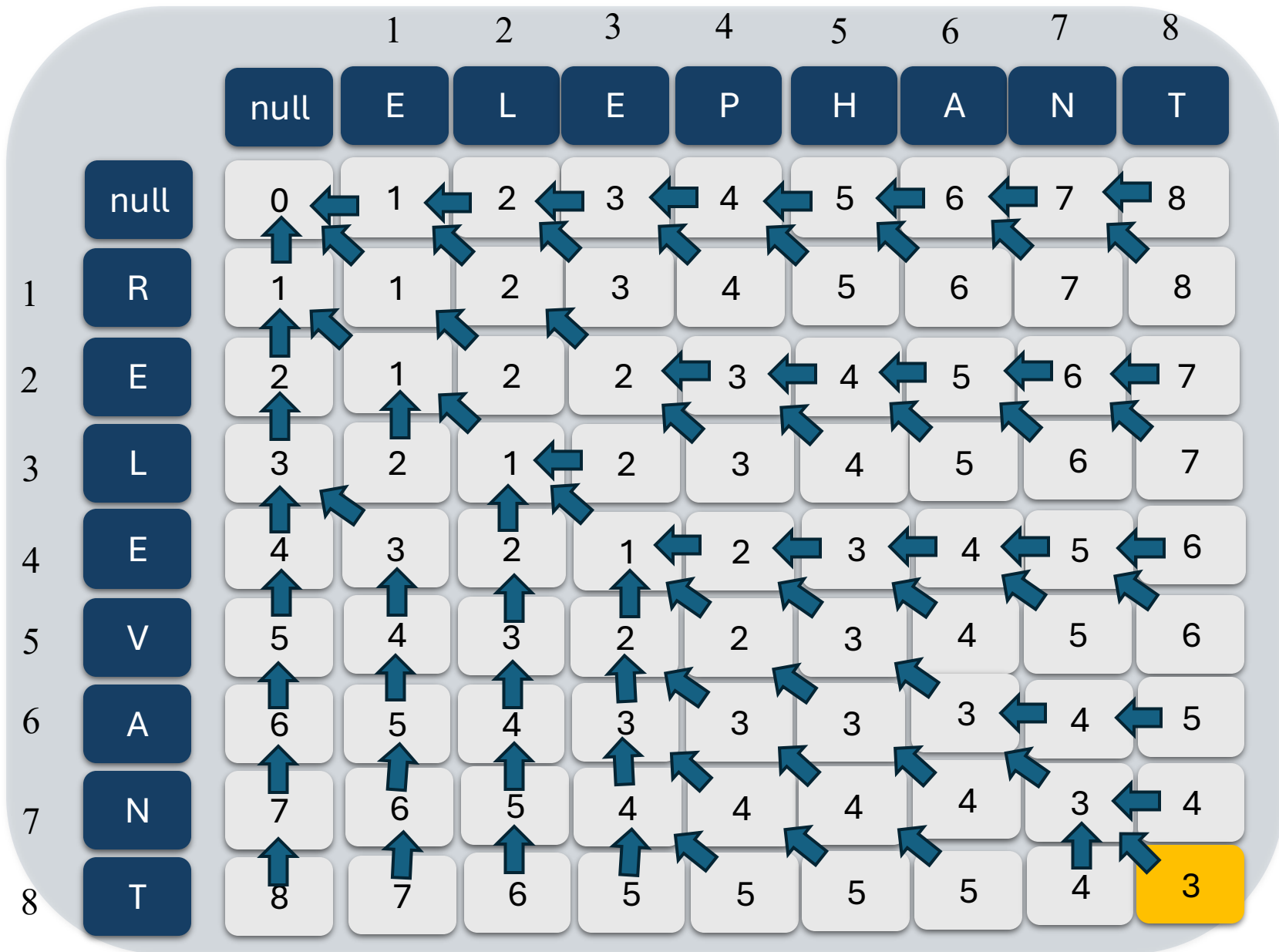
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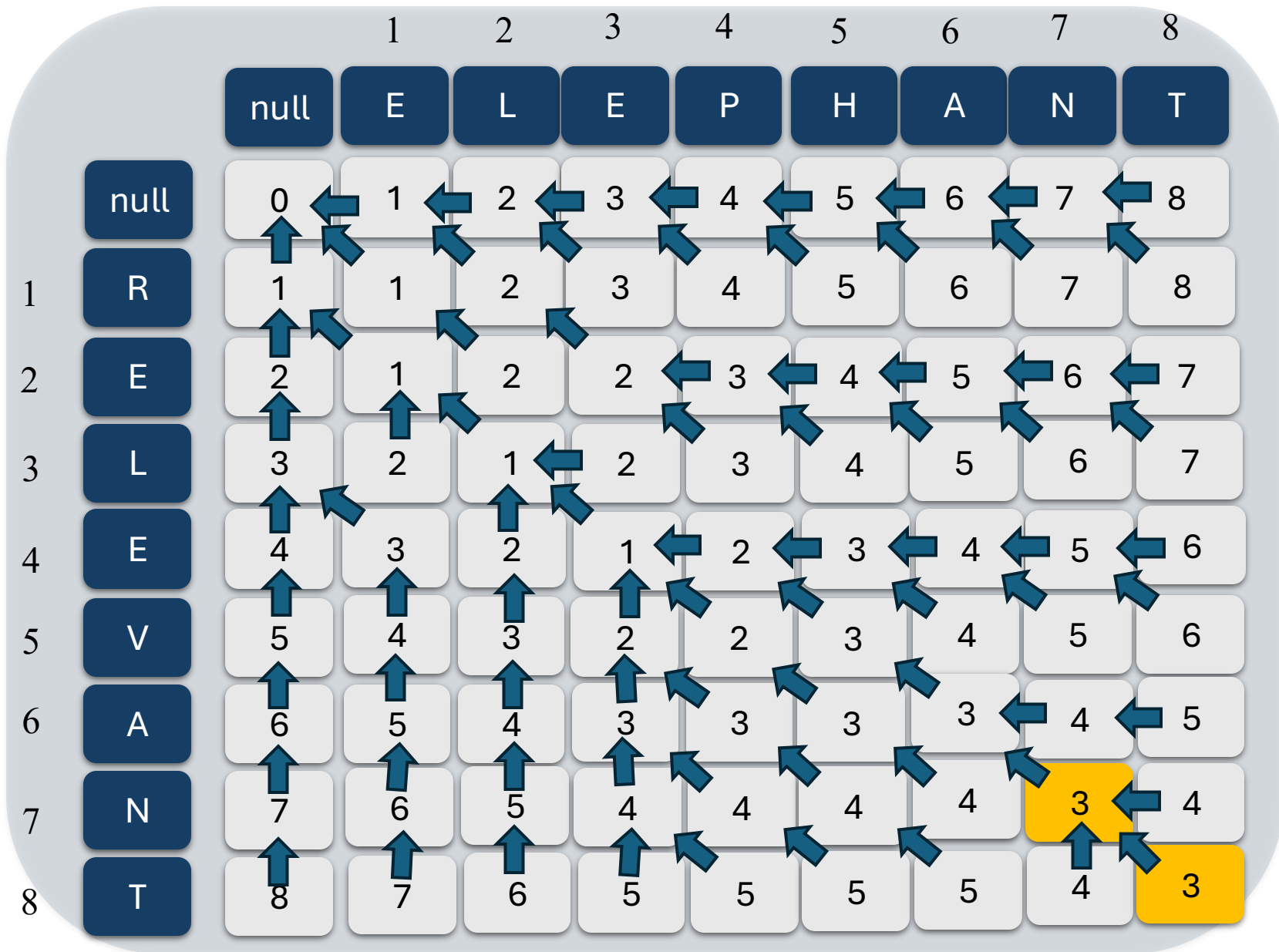


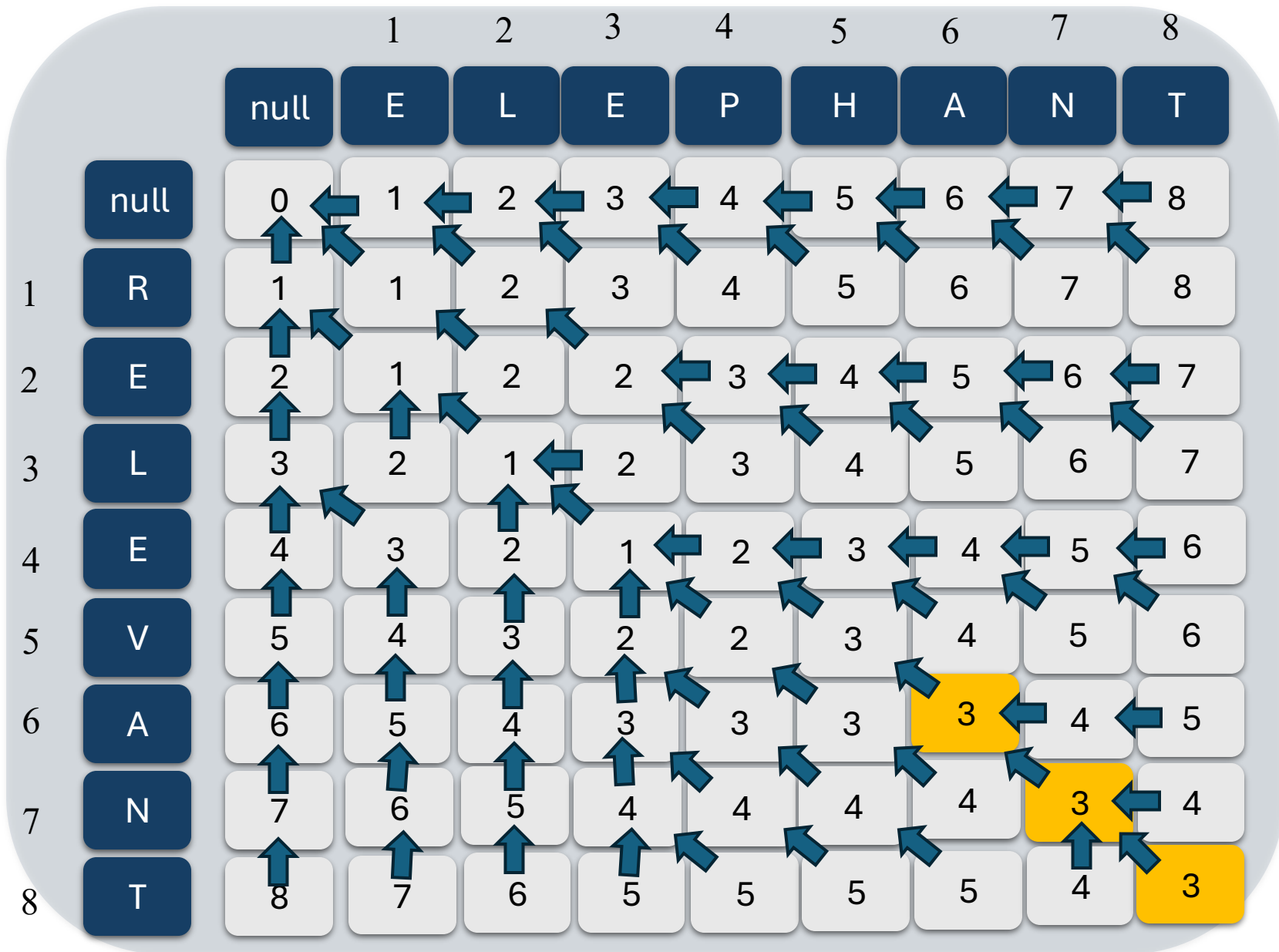
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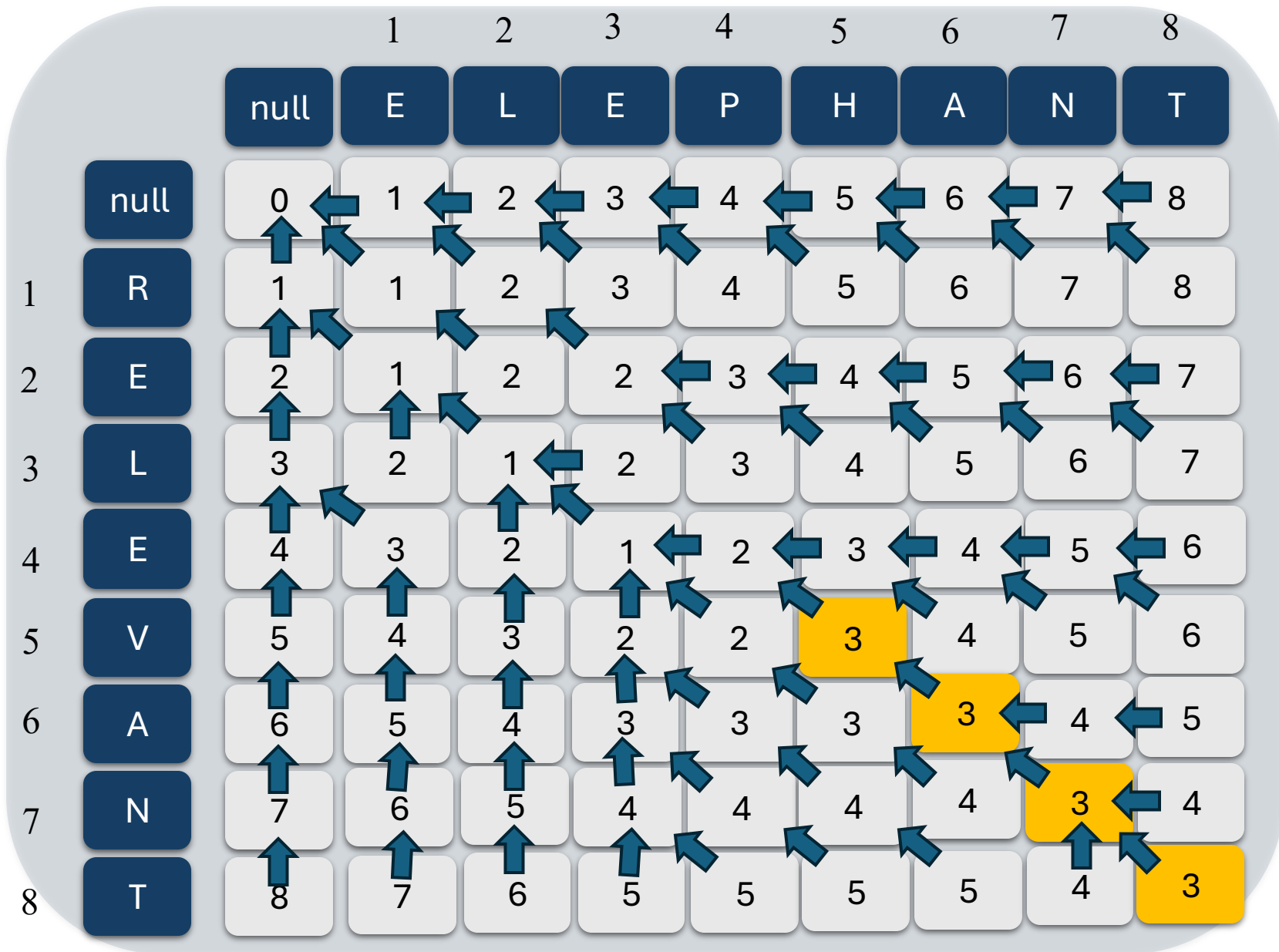


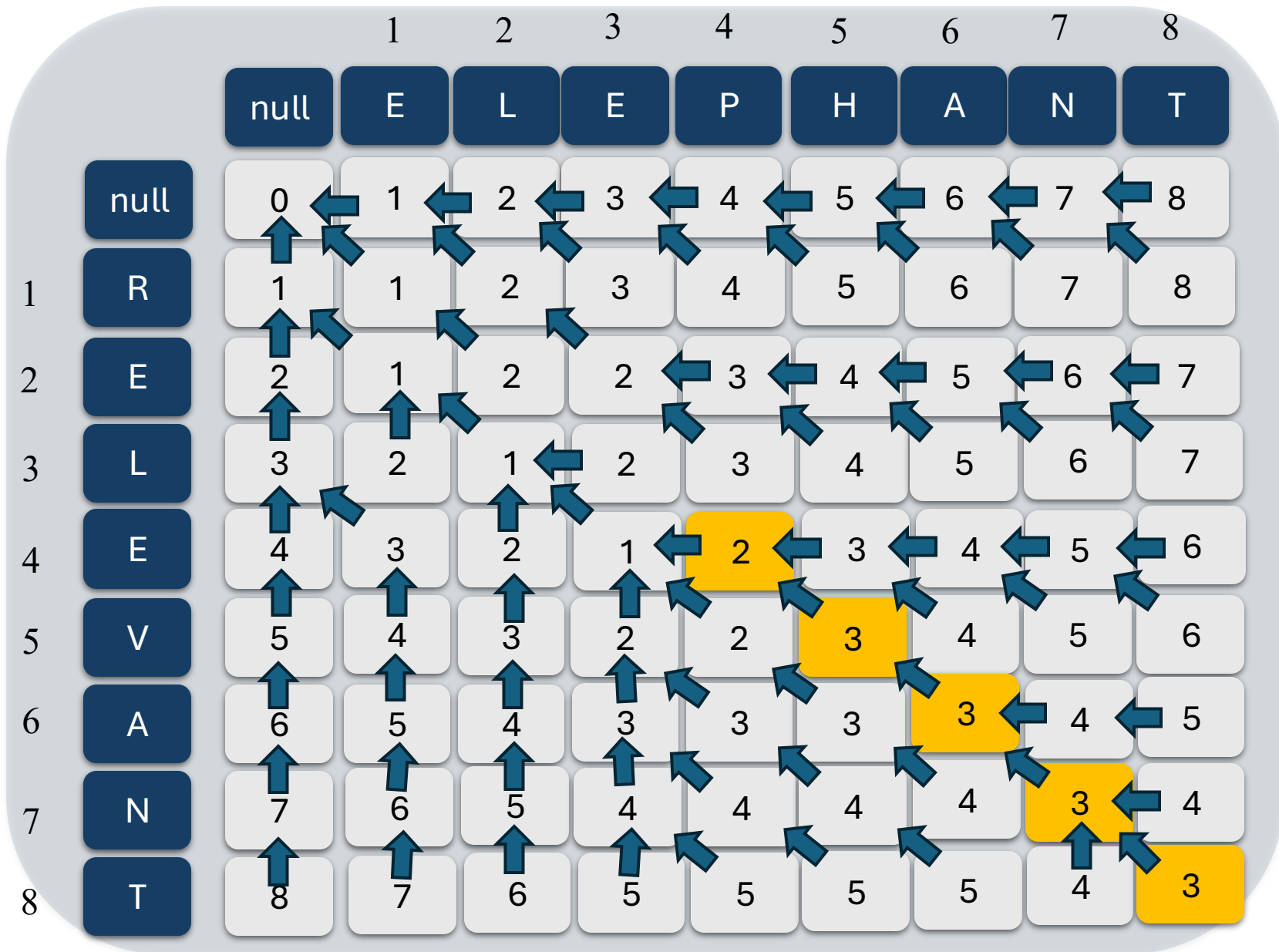
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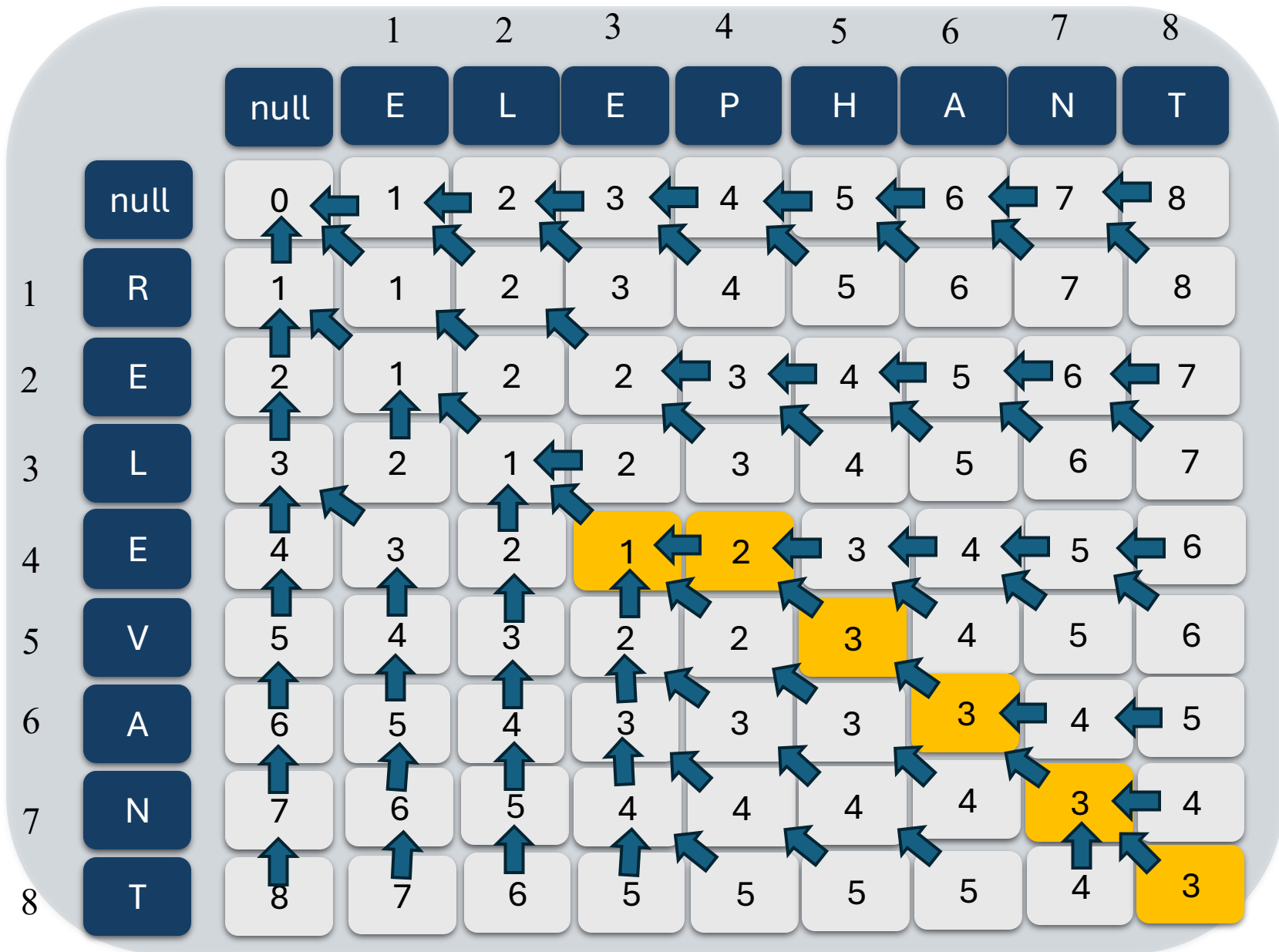


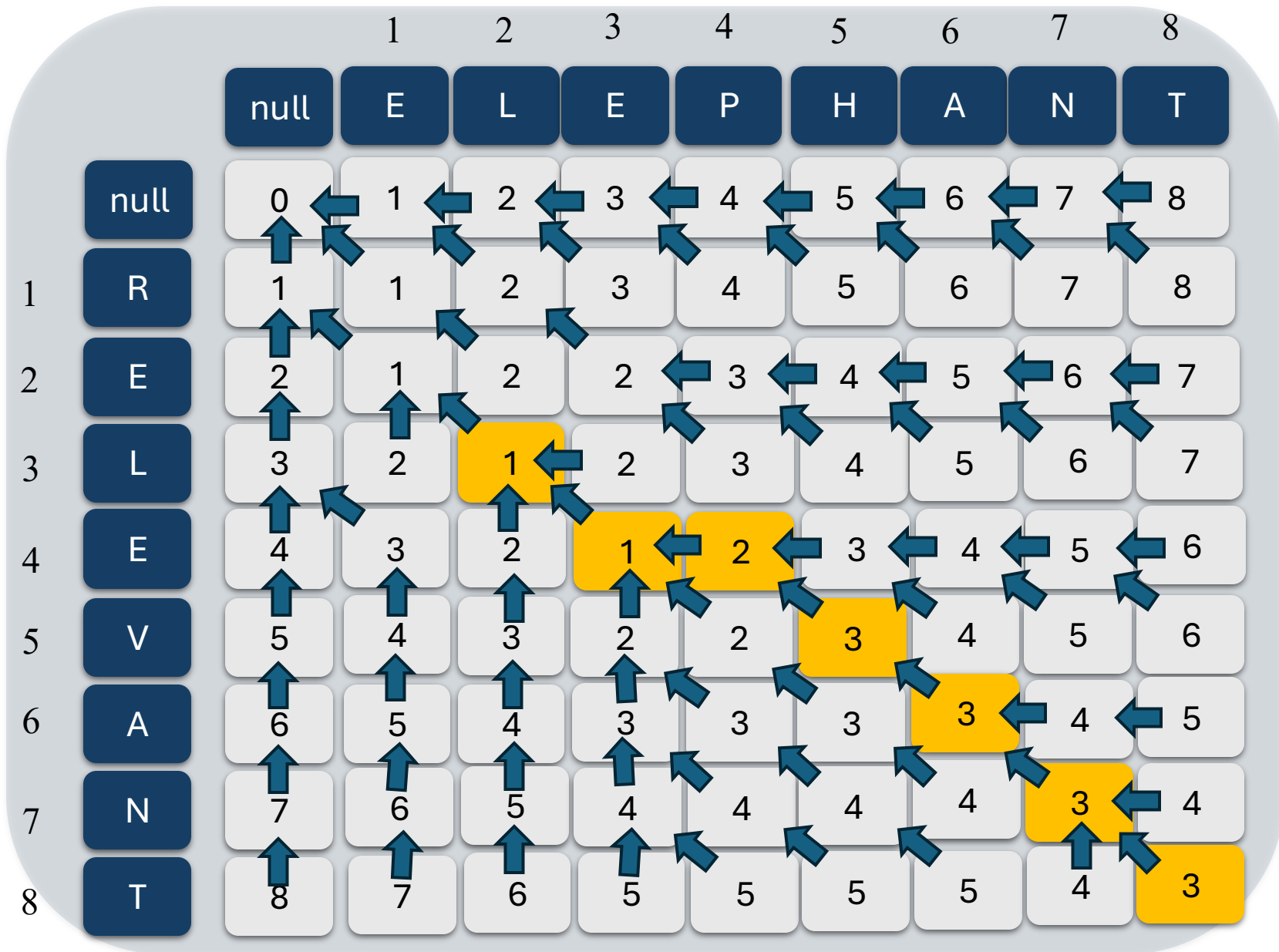


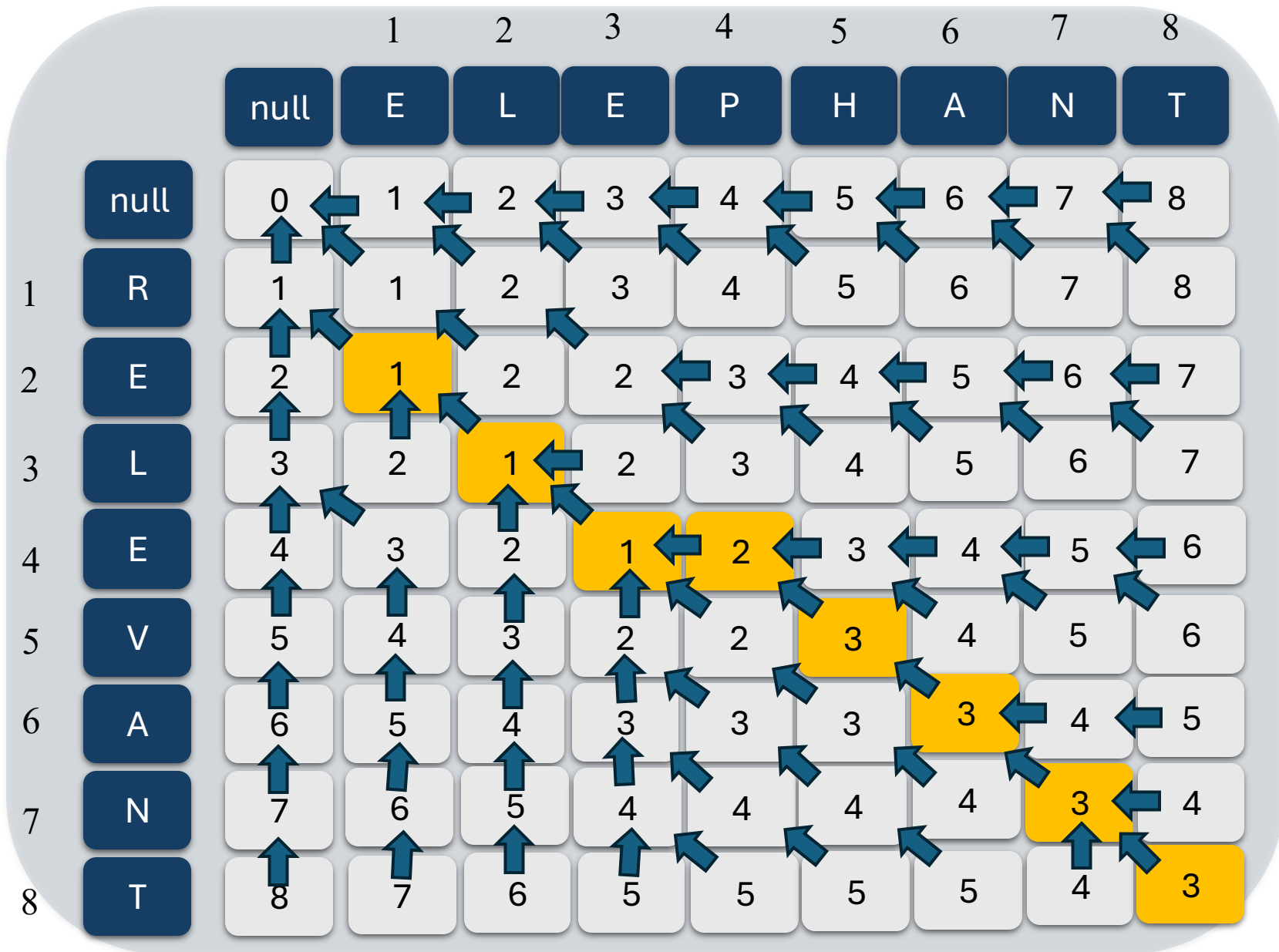


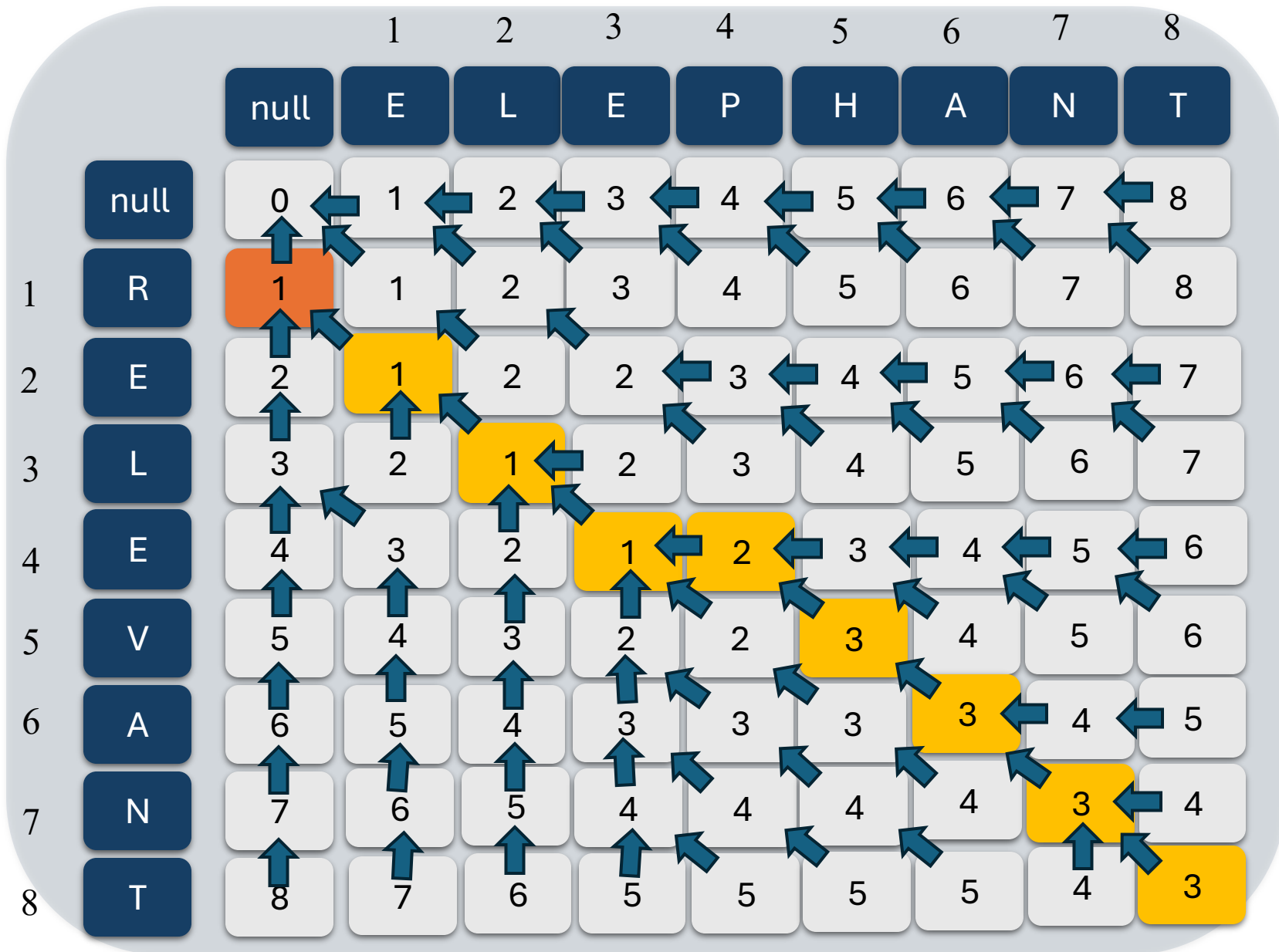




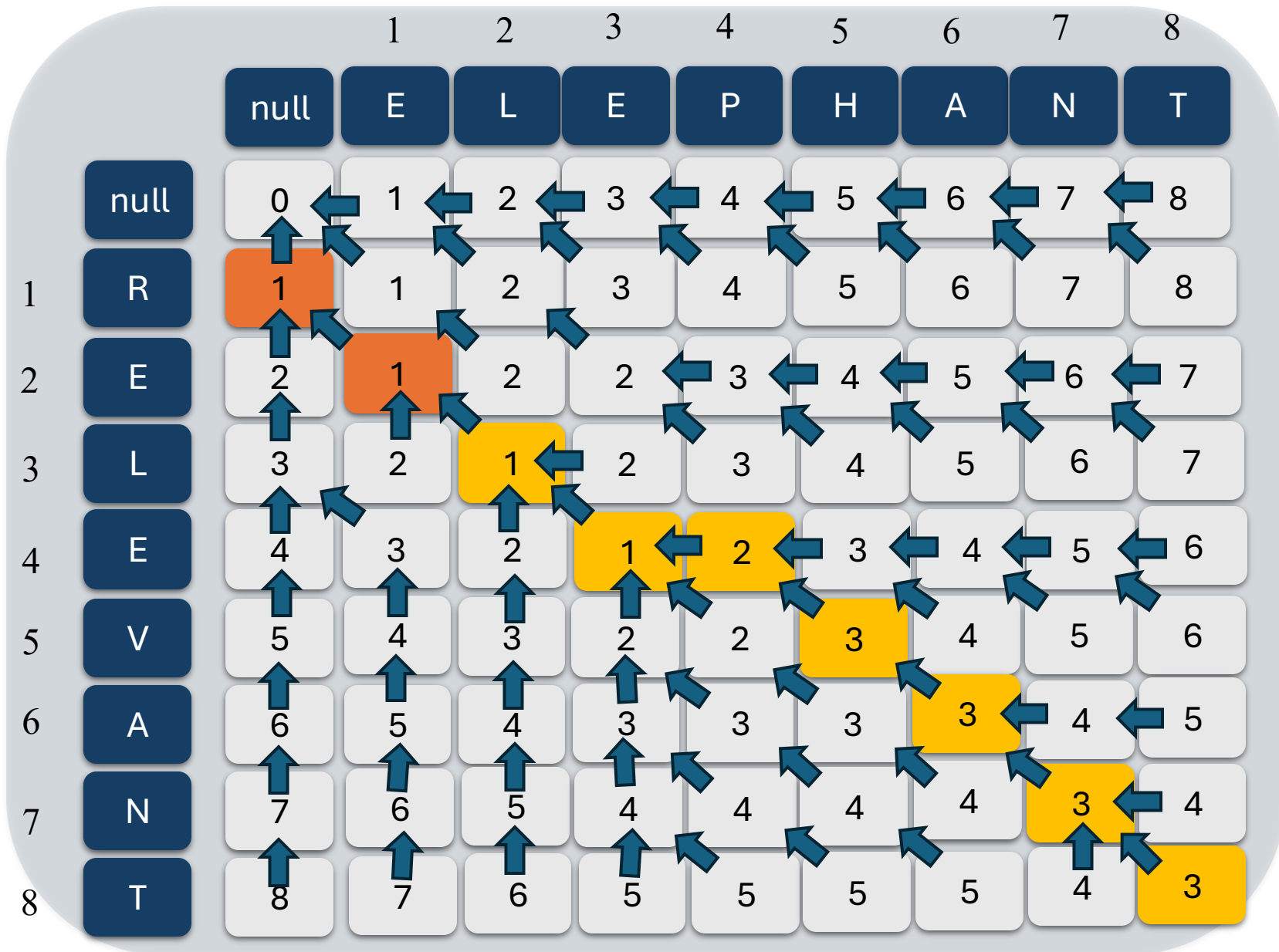




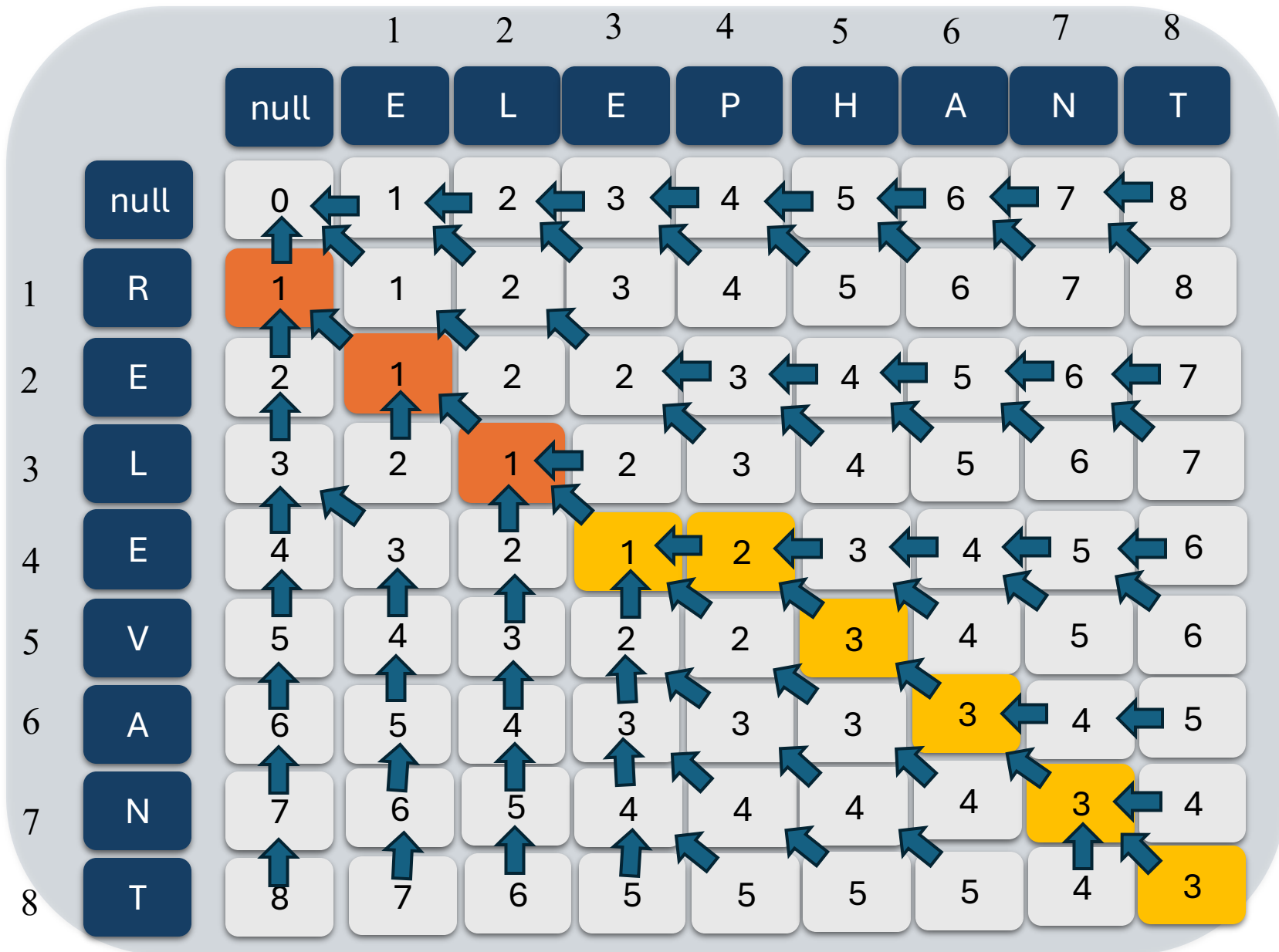




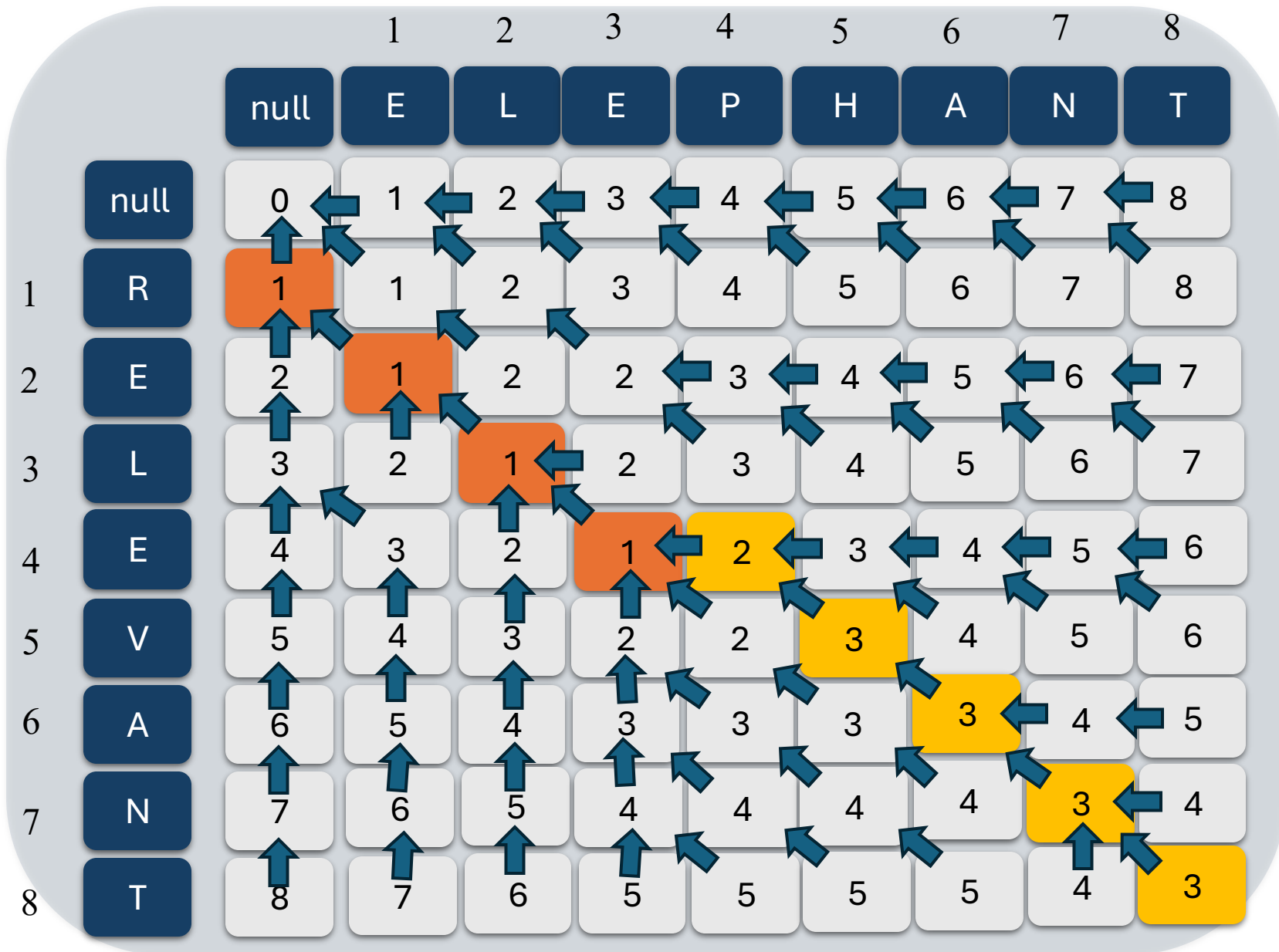
R E L E V A N T
-



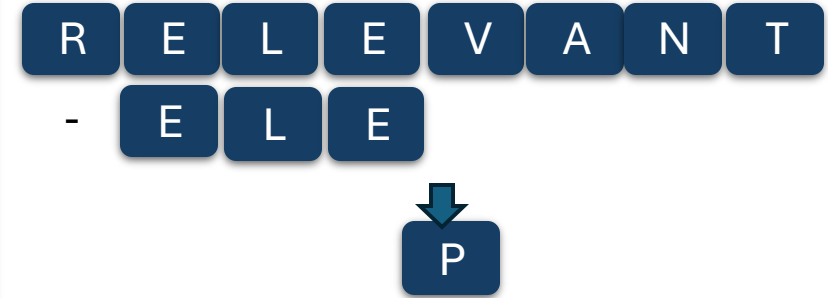
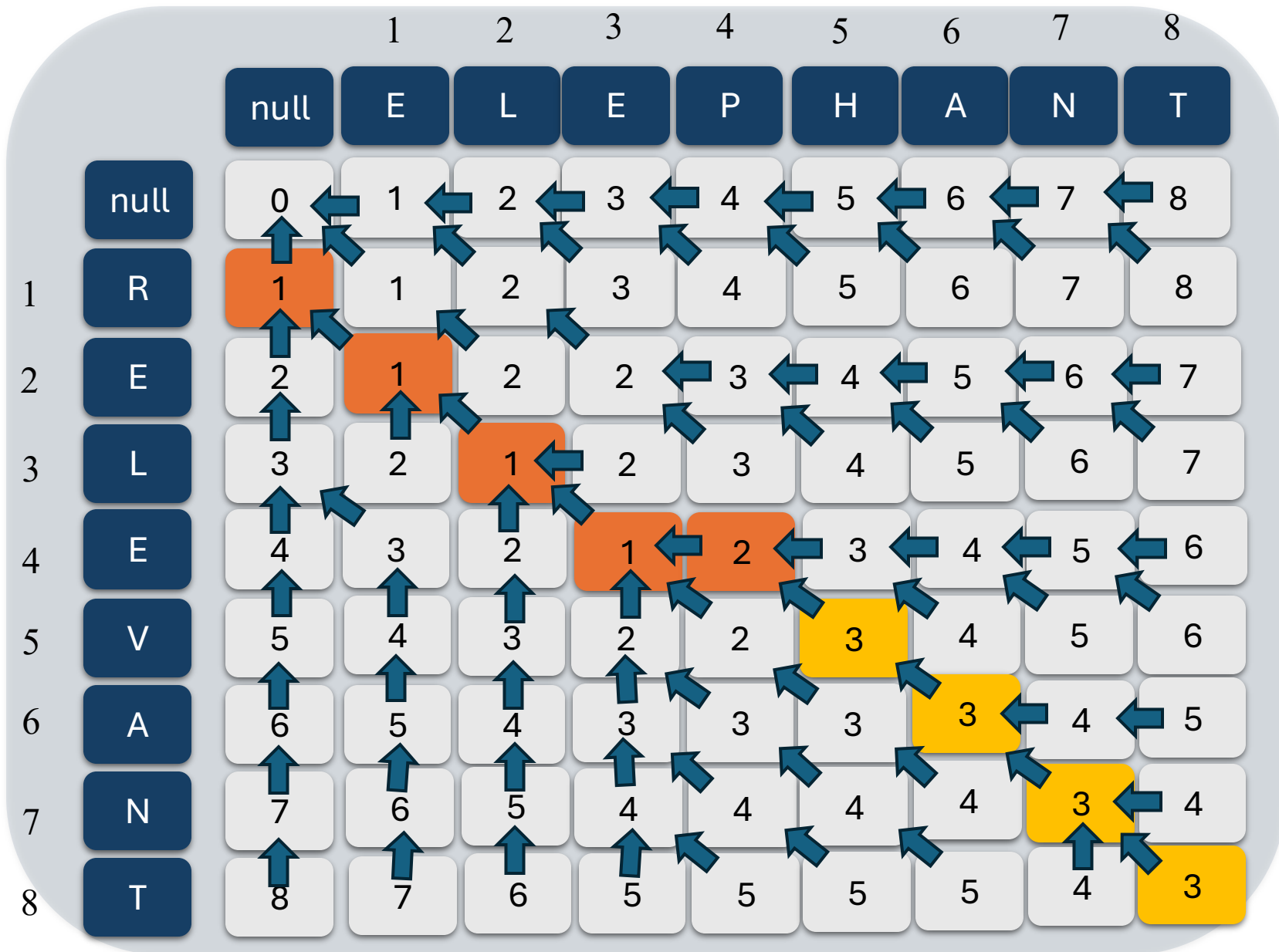
RELLEVANT
- E

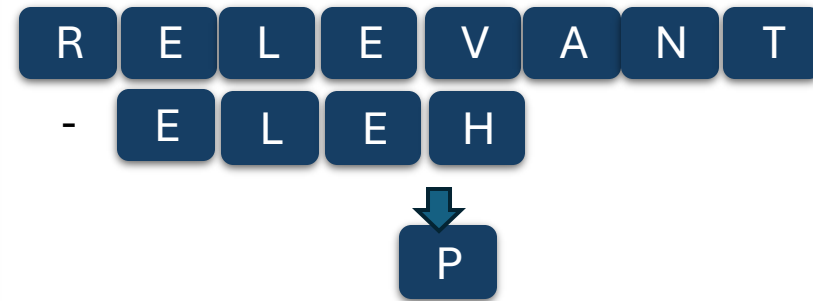
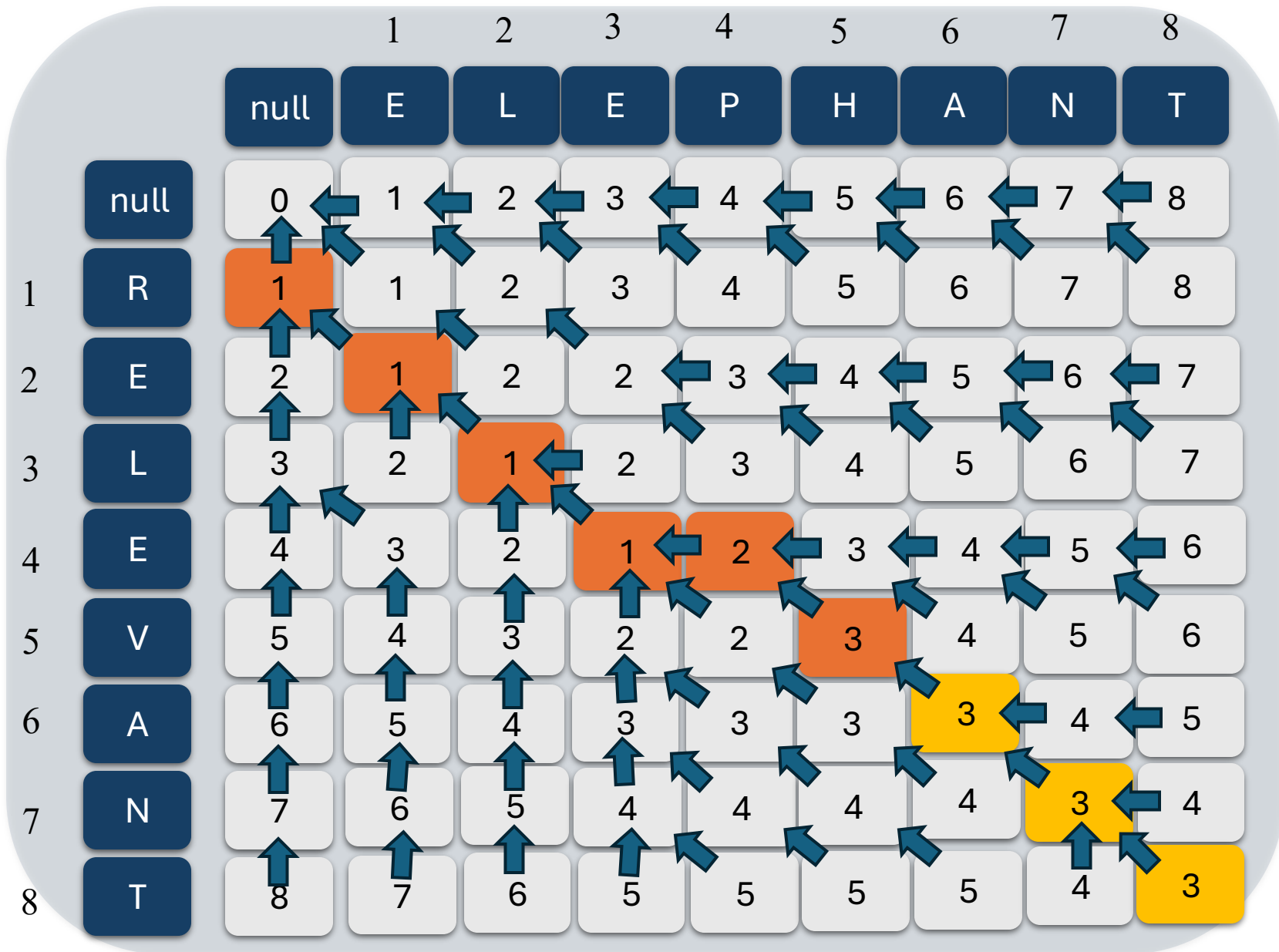


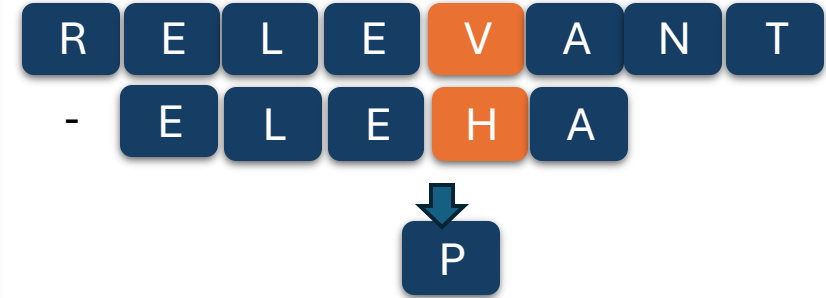
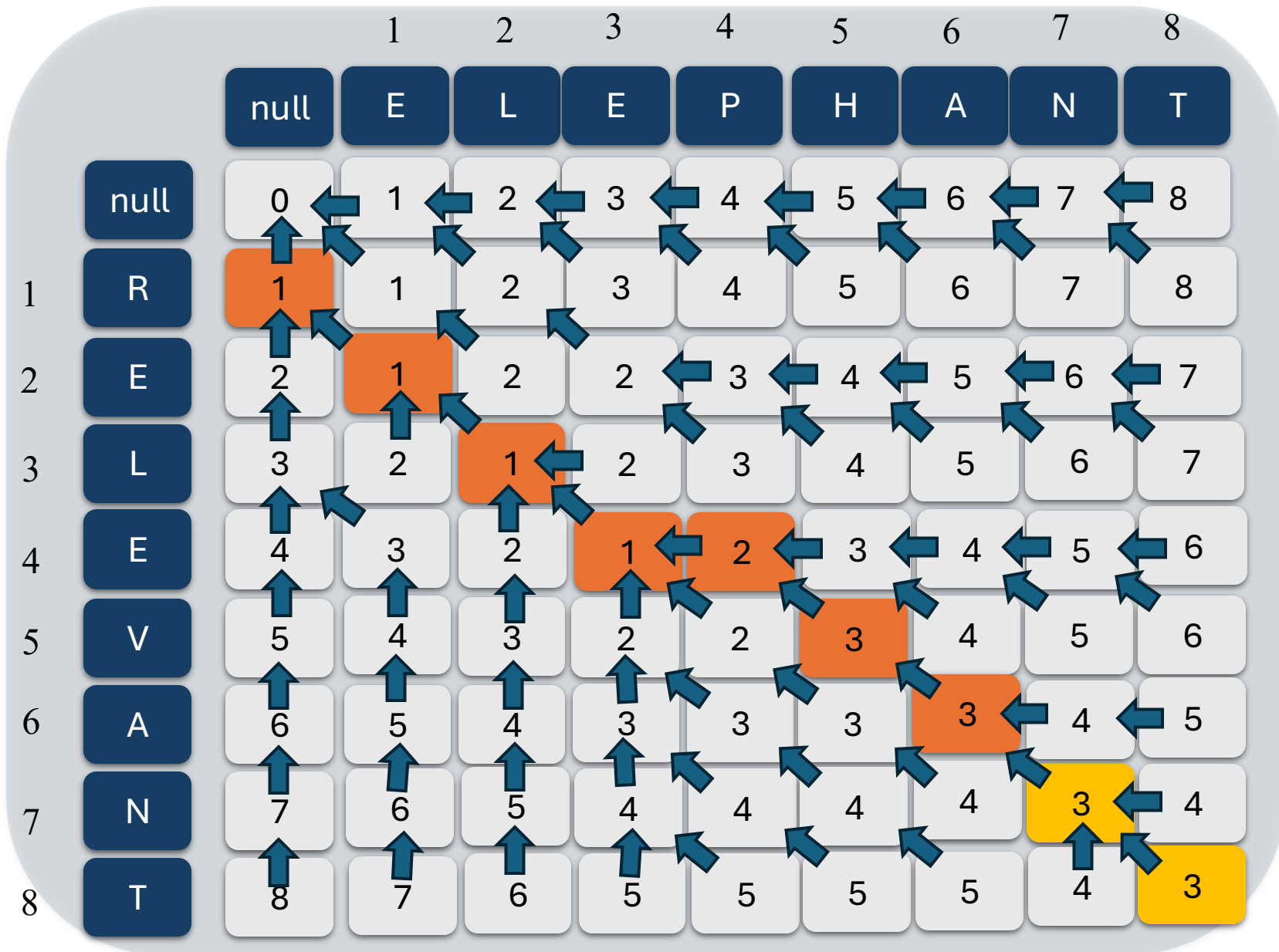
RELEVANT
- EL

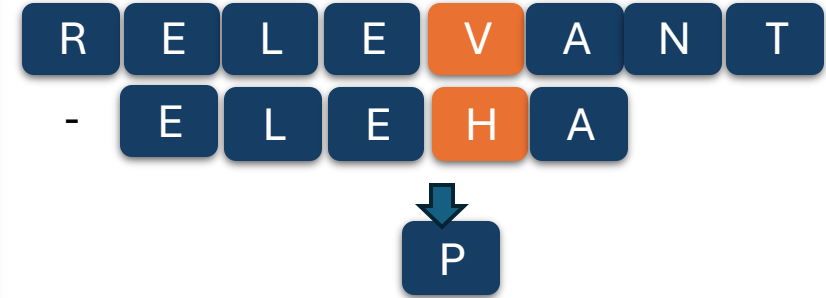
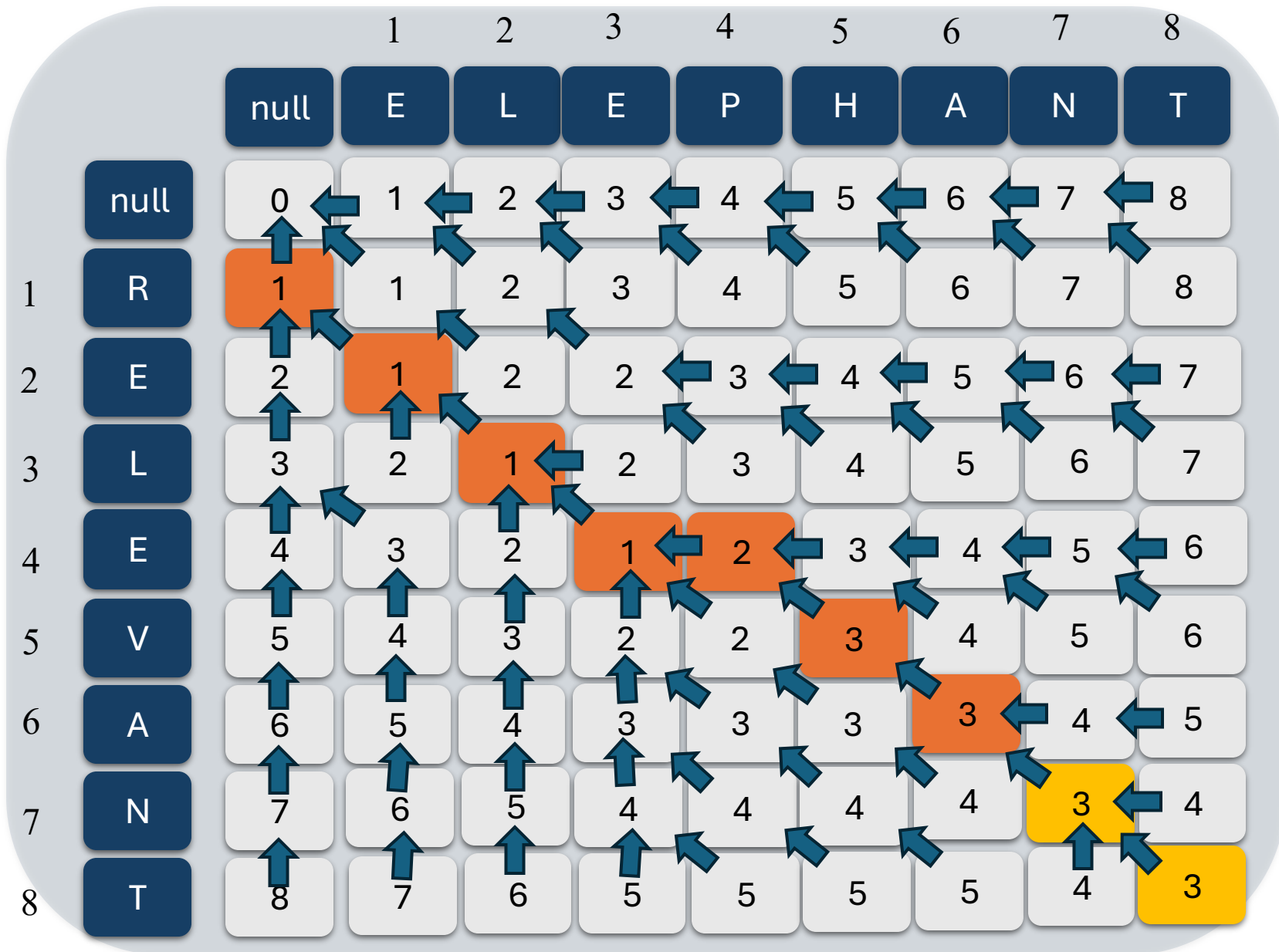


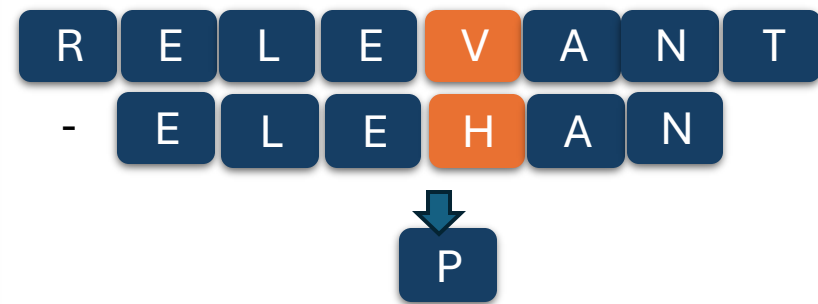
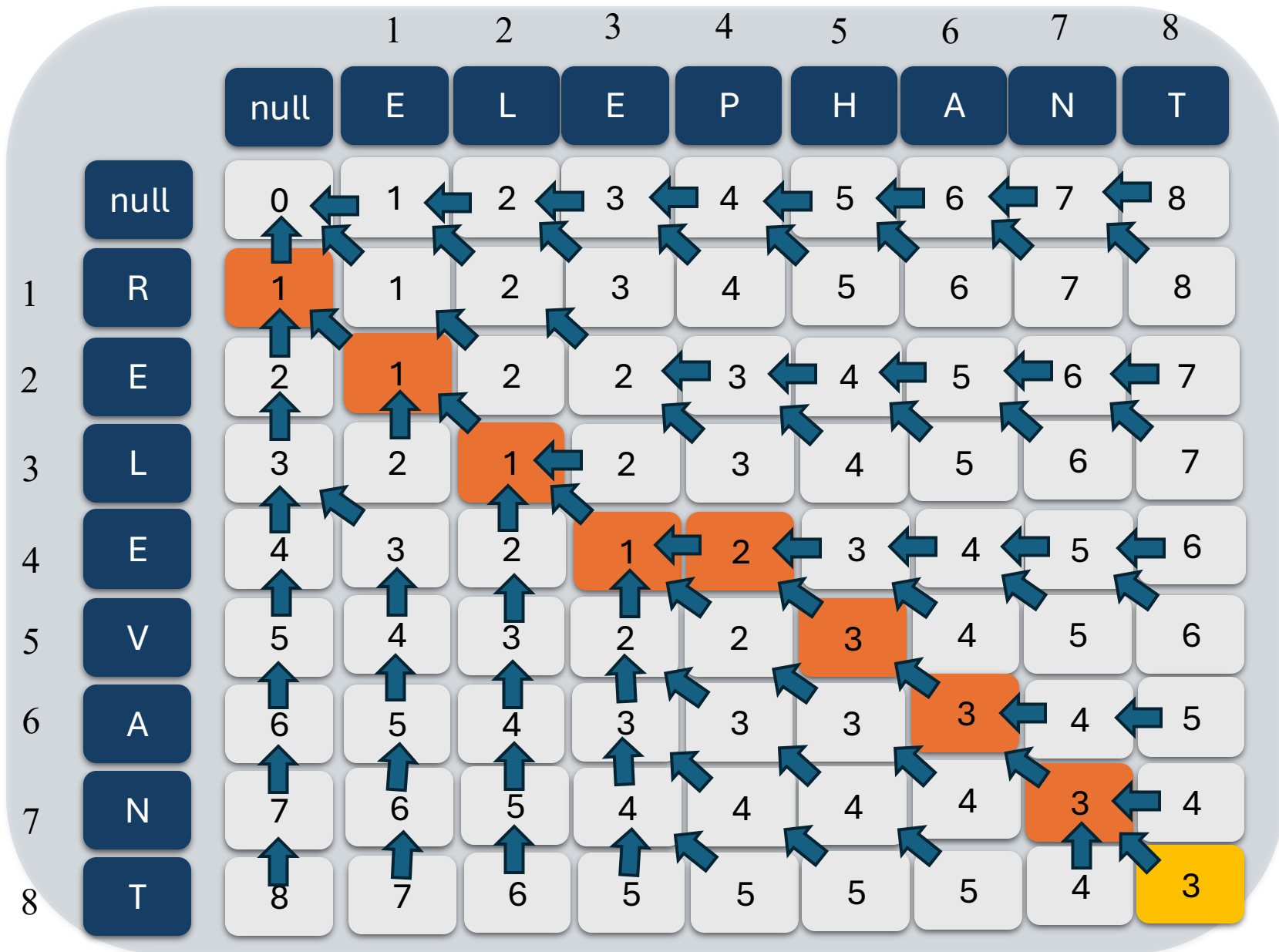
RELAXANT
- RELE

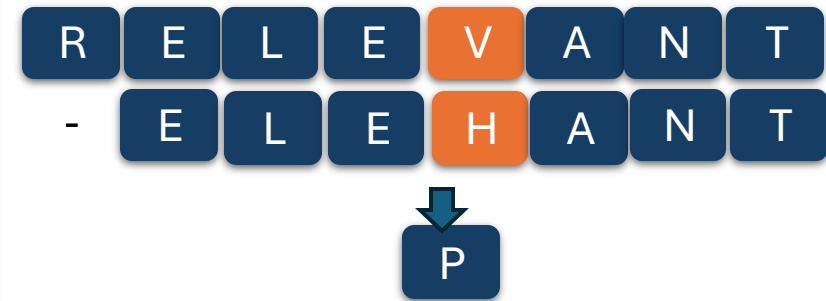
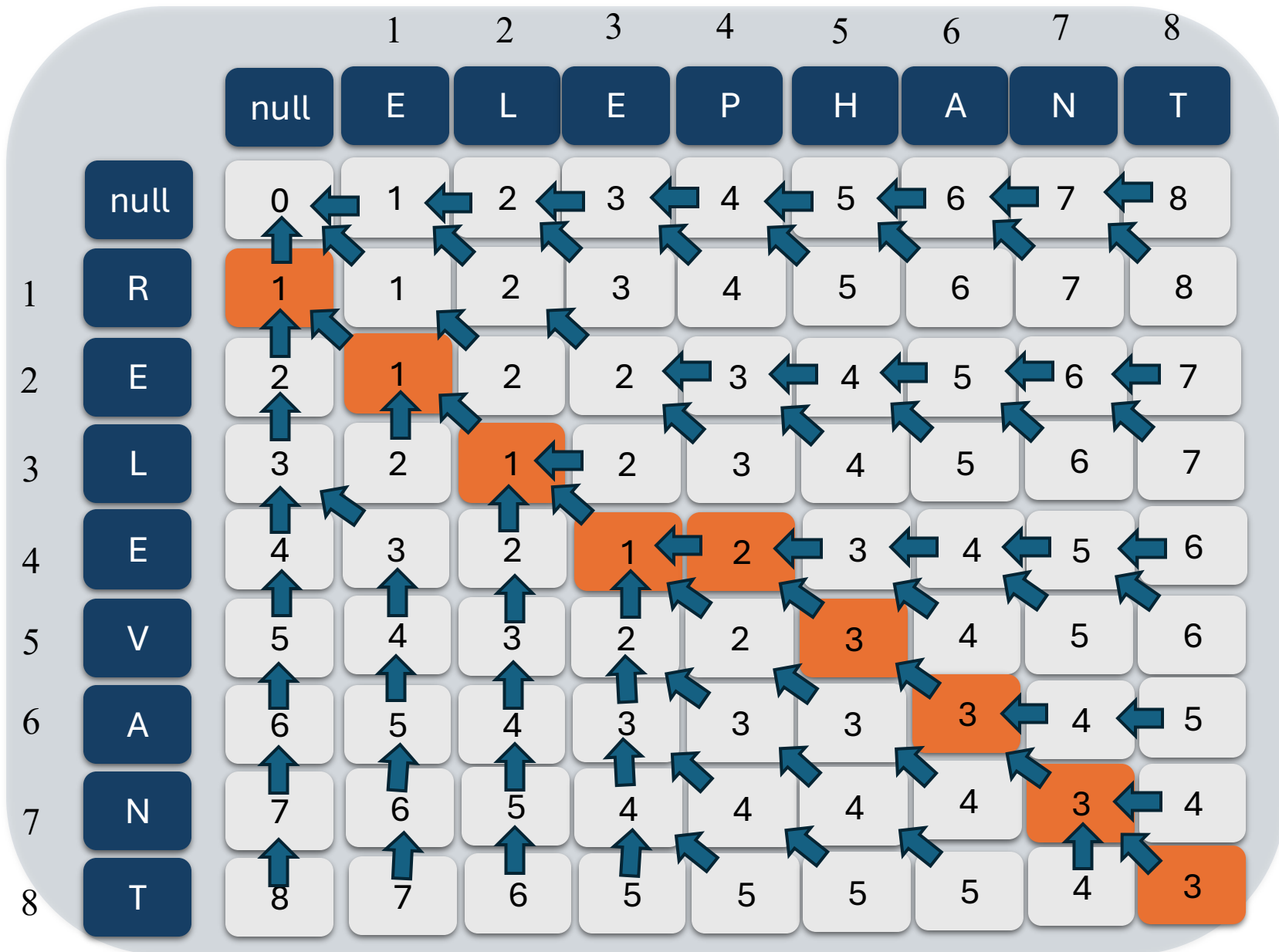












Algorithm

minimum_edit_distance(sequence X, sequence Y):

$M = \text{length}(\overline{X})$

$N = \text{length}(Y)$

#initialize table E of size M x N (M rows, N columns)

for i = 0 to M:

$E(i, 0) = i$ # Base Case

for j = 0 to N:

$E(0, j) = j$ # Base Case

for i = 1 to M:

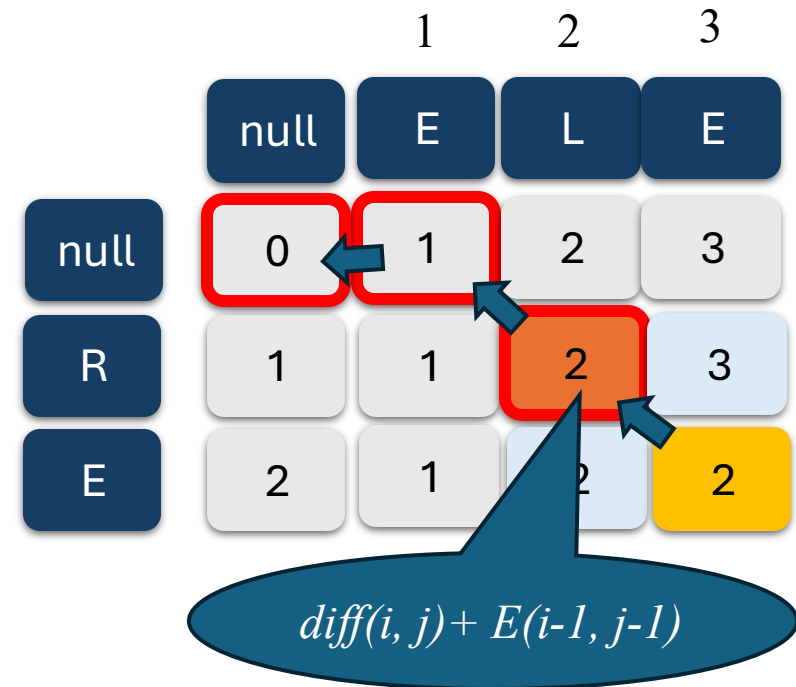
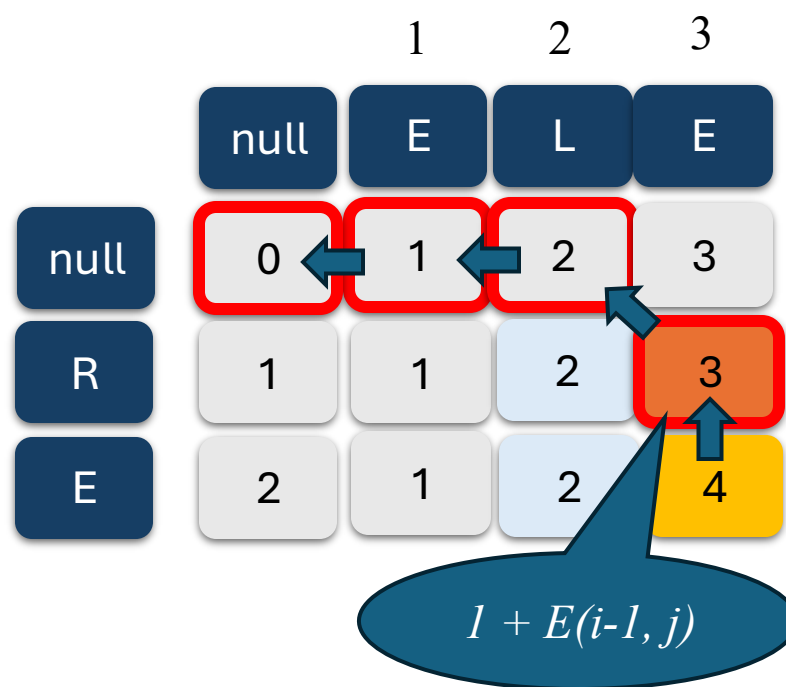
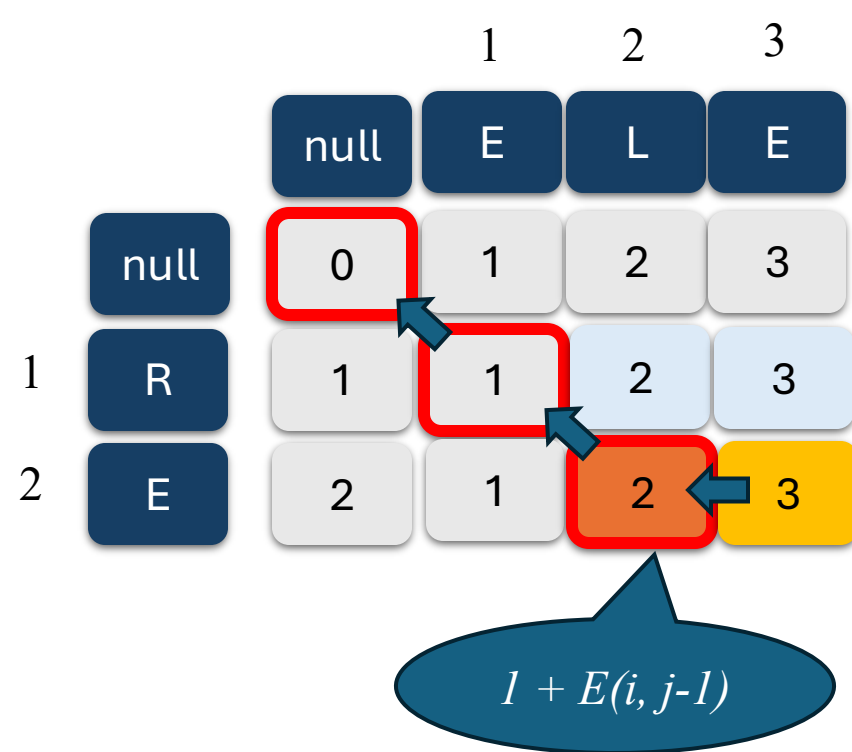
for j = 1 to N:

$\text{diff}(i, j) = (X[i] == Y[j]) ? 0 : 1$ # Diff

$E(i, j) = \min\{E(i-1, j)+1, E(i, j-1)+1, E(i-1, j-1)+\text{diff}(i, j)\}$

return E(i, j)

Some test cases



Longest Increasing Subsequence

Increasing Subsequence

A subsequence in which the numbers are getting strictly larger.

Example: $A = [2, 5, 3, 7, 8, 4]$

Increasing Subsequences:

- $[2, 3, 4], [2, 3, 7], [2, 3, 8], [2, 3, 7, 8], [2, 5, 7, 8]$

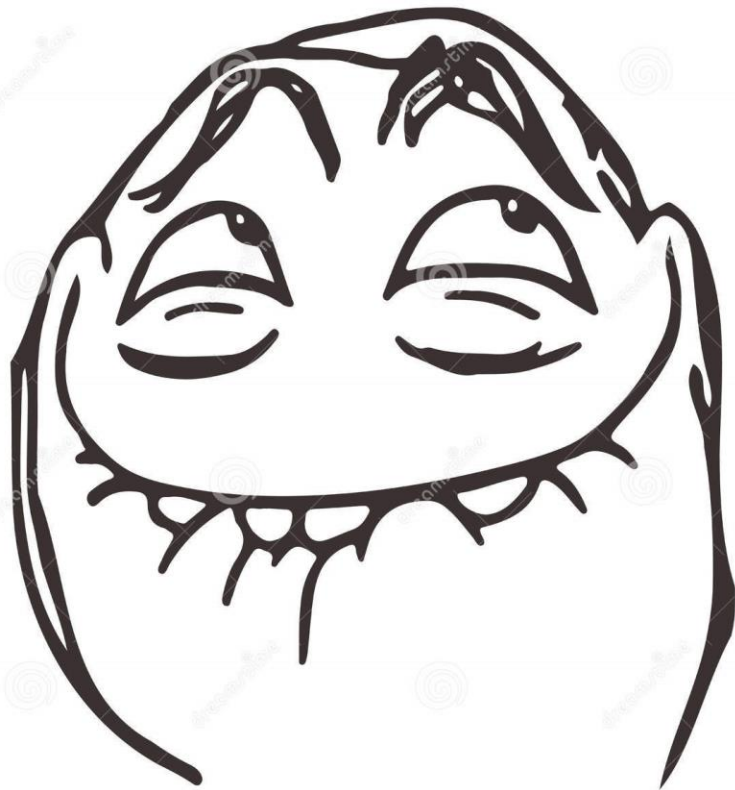
Longest Increasing Subsequence

- An increasing subsequence of greatest length

LIS

- Input: Sequence of numbers $A = \{a_1, a_2, \dots, a_n\}$
- Output: Longest increasing subsequence of A

How can we solve the problem?



I know brute force
isn't the way but I'll
say it anyway.

BRUTE-FORCE.

Brute Force Algorithm

- Find all increasing subsequences of A
- Return the increasing subsequence with greatest length

Brute Force Algorithm

brute_force_LIS(sequence A):

subsequences = all subsequences of A

max_length = 0

for subseq in subsequences:

if is_increasing(subseq):

if length(subseq) > max_length:

LIS = subseq

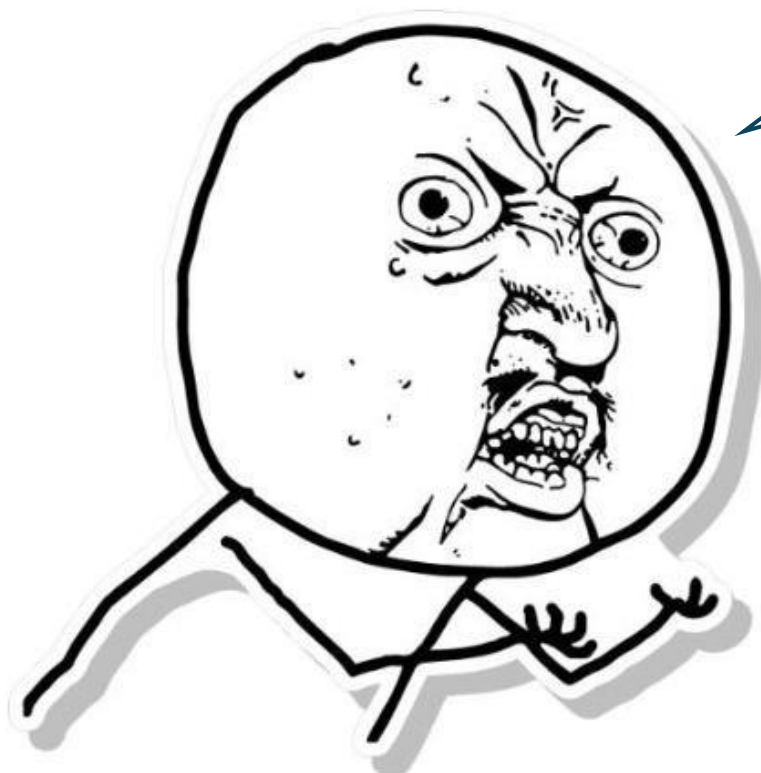
max_length = length(subseq)

return max_length, LIS

Analysis

- Finding all increasing subsequences = $O(2^n)$
- Involves checking each subsequence and checking if that subsequence is increasing

What's a better way of solving for LIS?



OMG. Tinatanong pa
ba yan?

Use **DP Approach!**

Solving Longest Increasing Subsequence using Longest Common Subsequence

- *Let $B = \text{sorted version of sequence } A$*
 - *takes $O(N \log N)$ time using merge sort*
- *$LIS(A) = LCS(A, B)$*
 - *takes $O(N^2)$ time $\Rightarrow O(MN)$*
 - *but $M = N$ (B is just sorted version of A) so $O(N^2)$*

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$A = [5, 2, 8, 6, 3, 6, 9, 7]$ $B = [2, 3, 5, 6, 6, 7, 8, 9]$

$LCS(A, B) = [2, 3, 6, 9]$

DP Idea

Example: [1,7,2,3,6],

LIS: 1236

Consider: [1,7,2,3],

LIS: 123

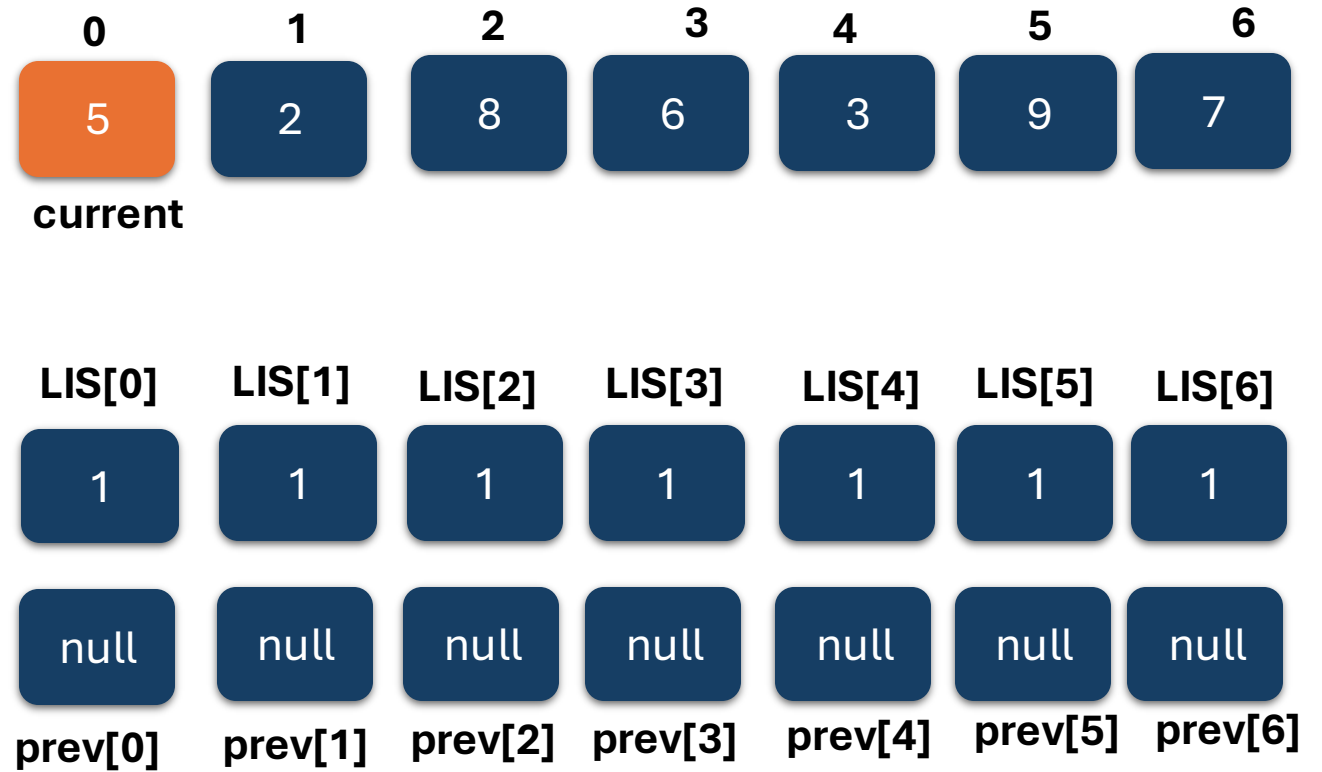
Now: [1,7,2],

LIS: 12 and 17

- What do you notice?

DP Idea

- The LIS of the shorter sequence can also be part of the LIS of the longer sequence
- This is our **optimal substructure**.



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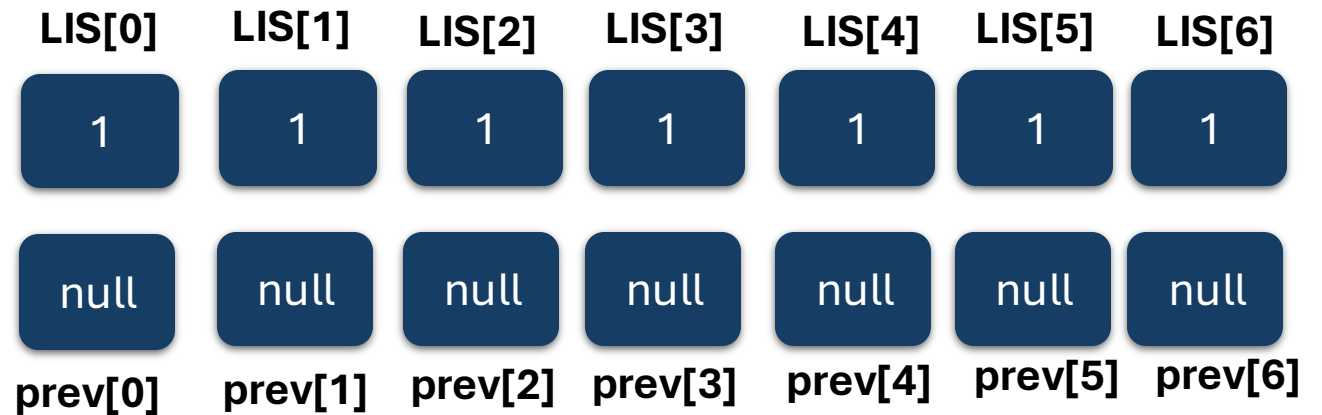
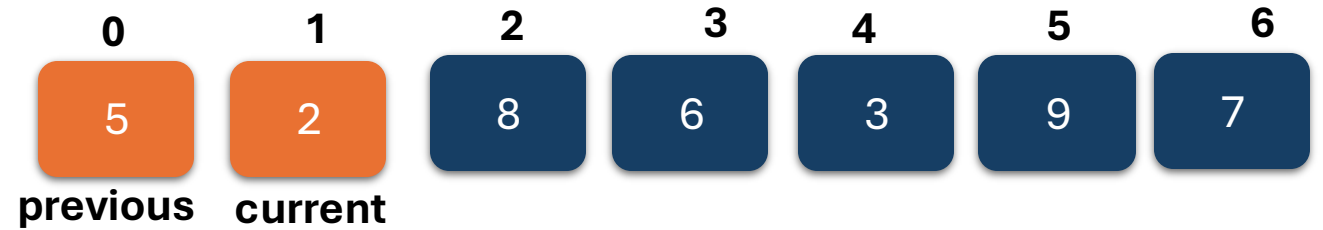
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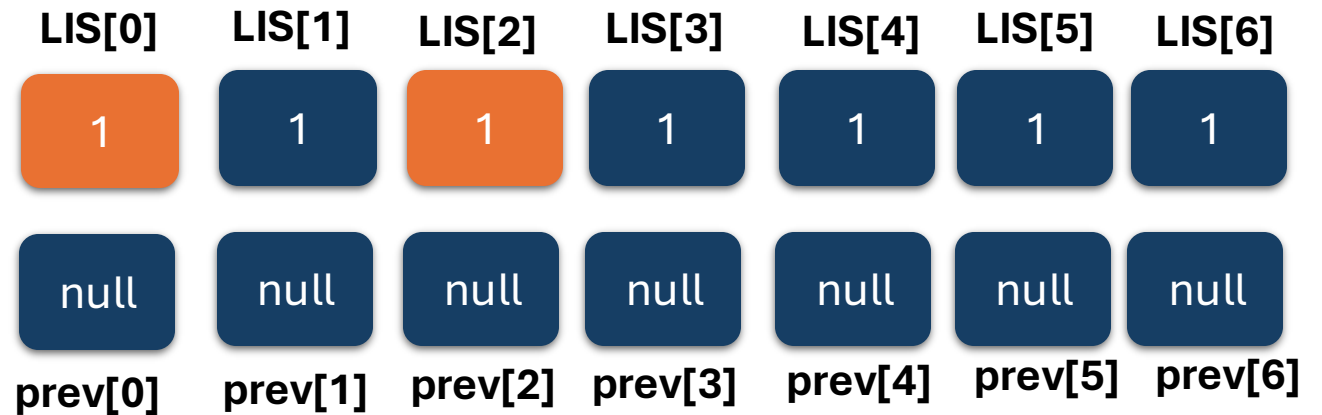
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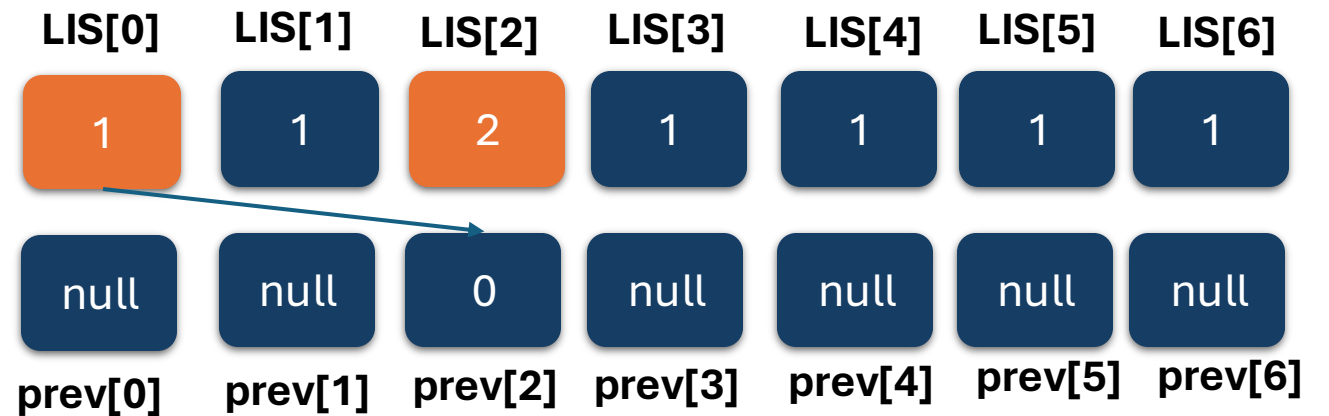
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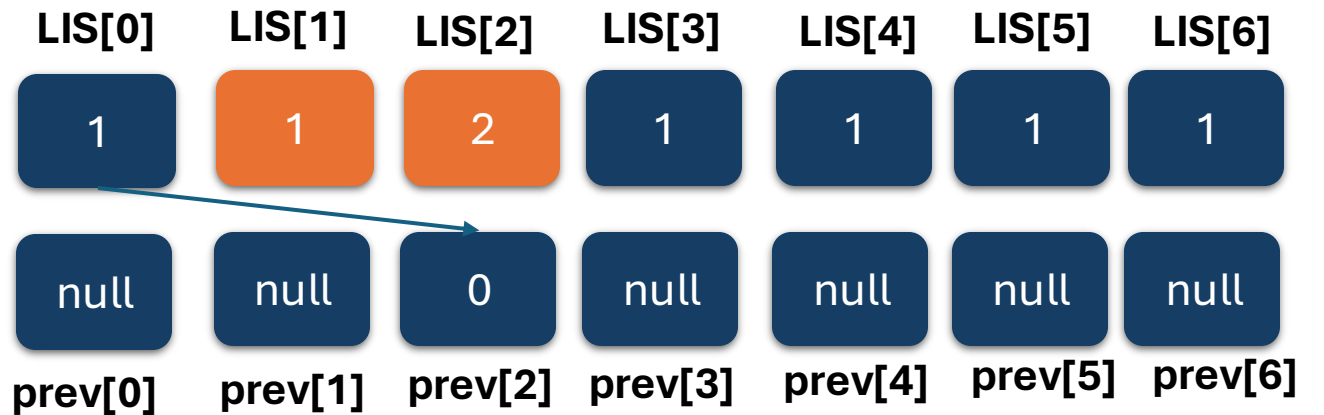
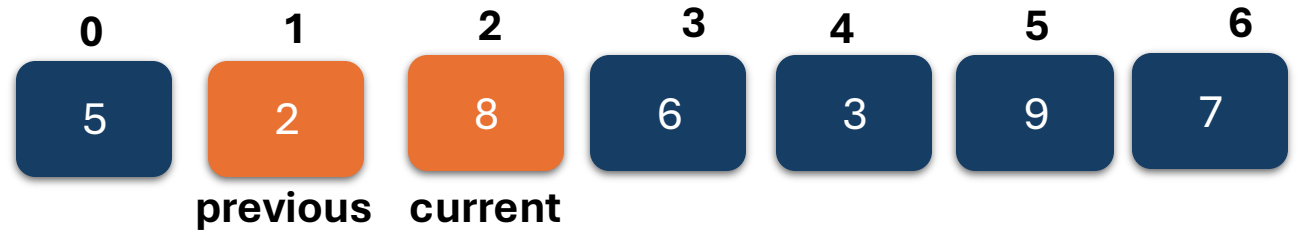
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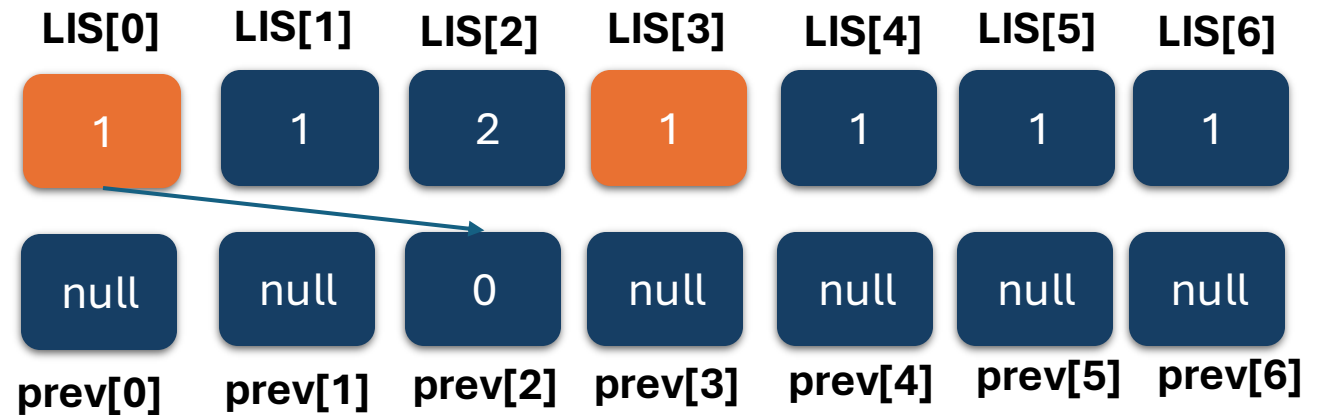
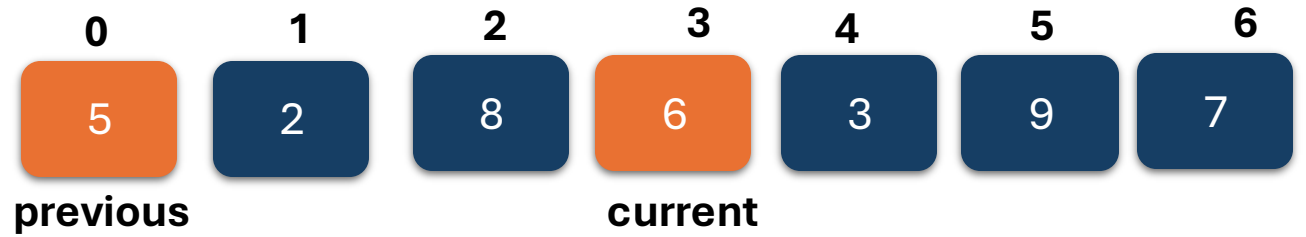
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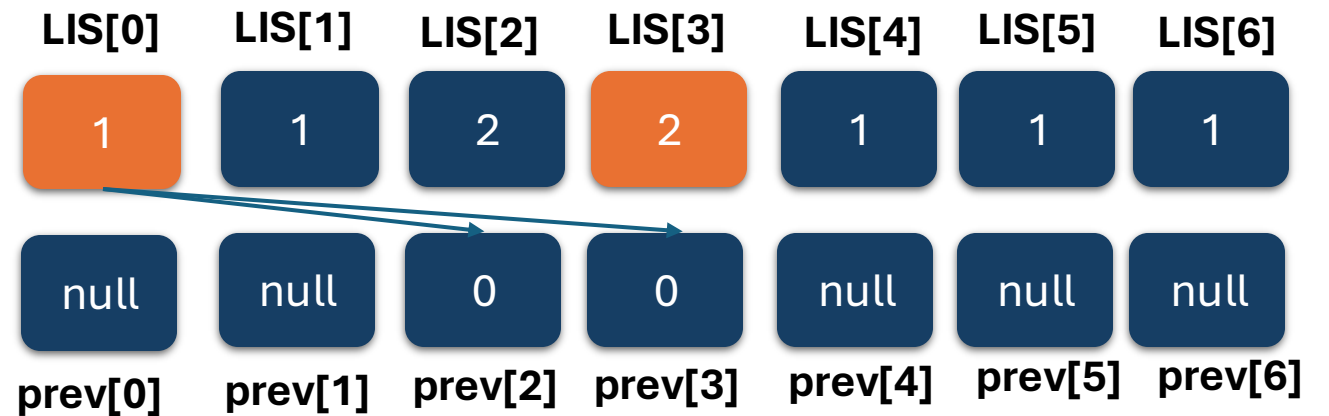
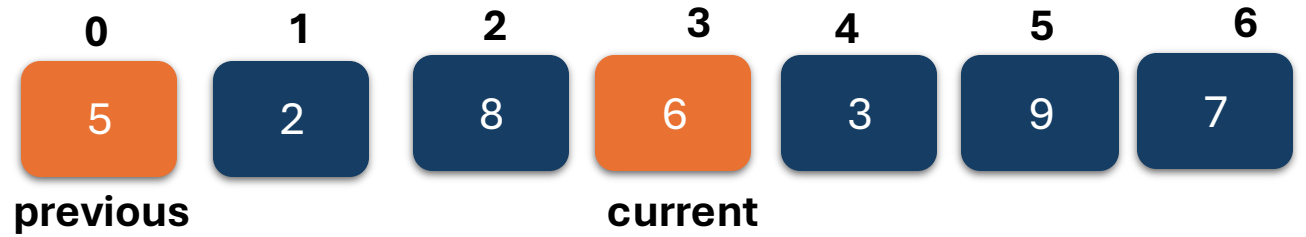
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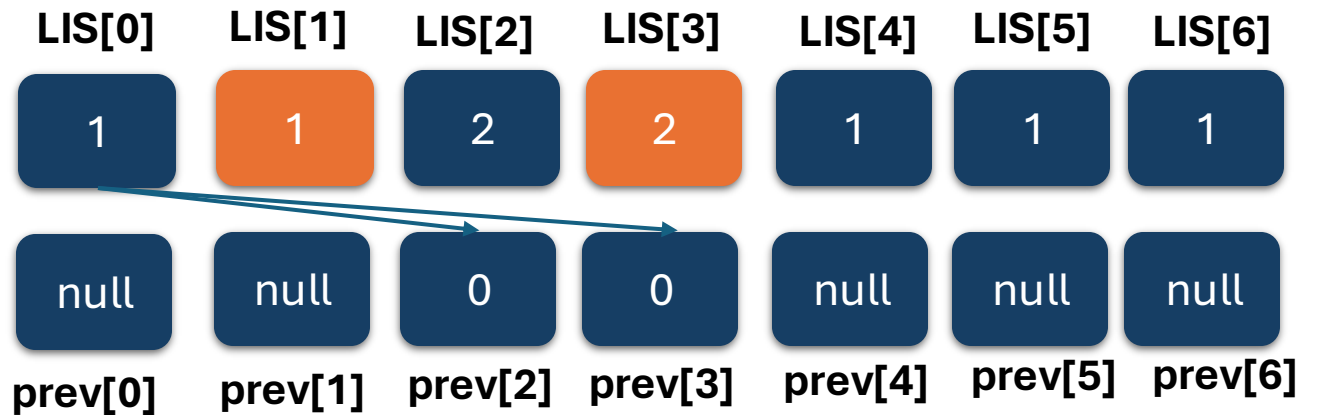
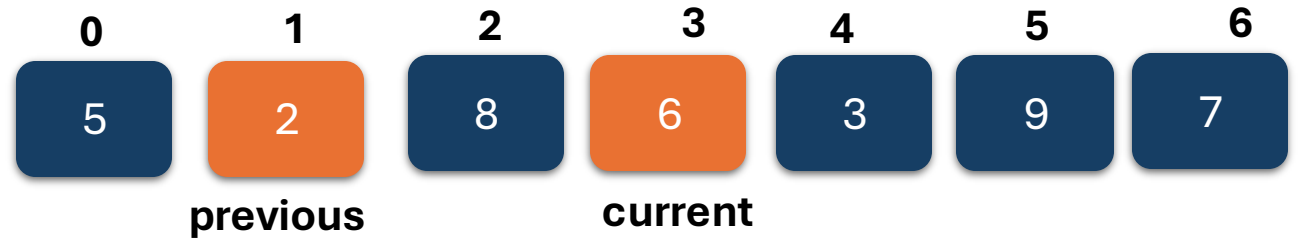
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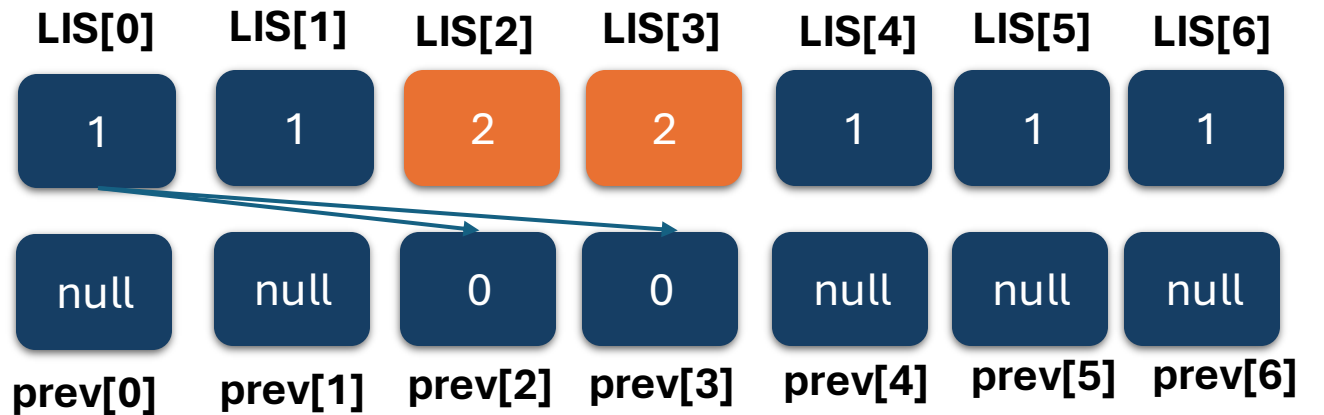
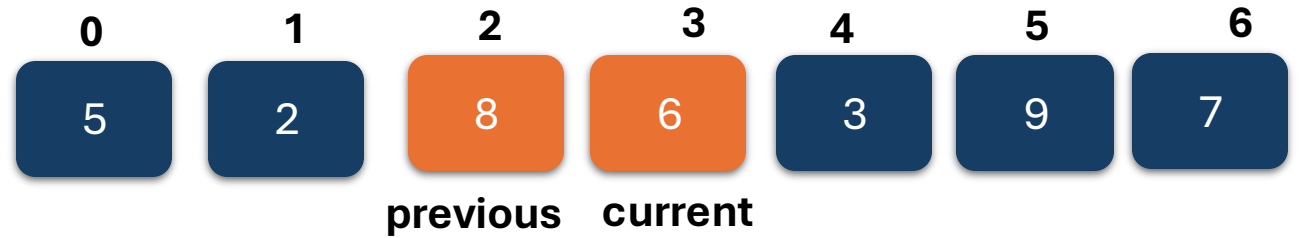
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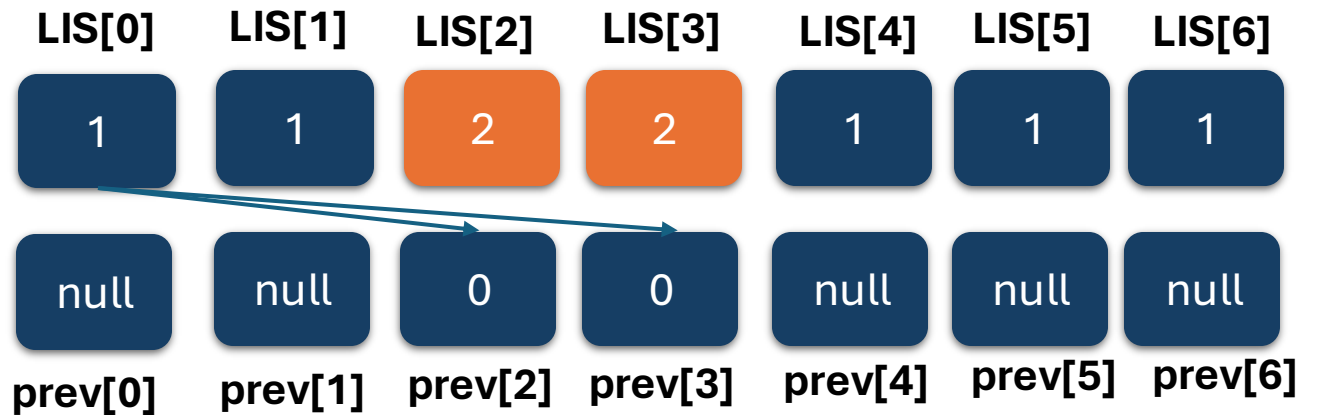
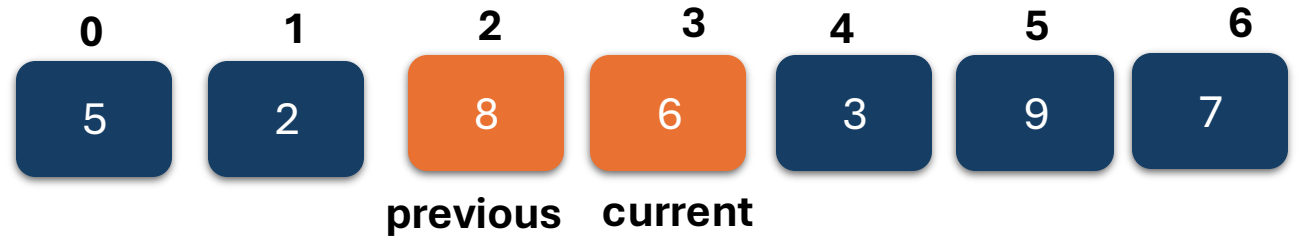
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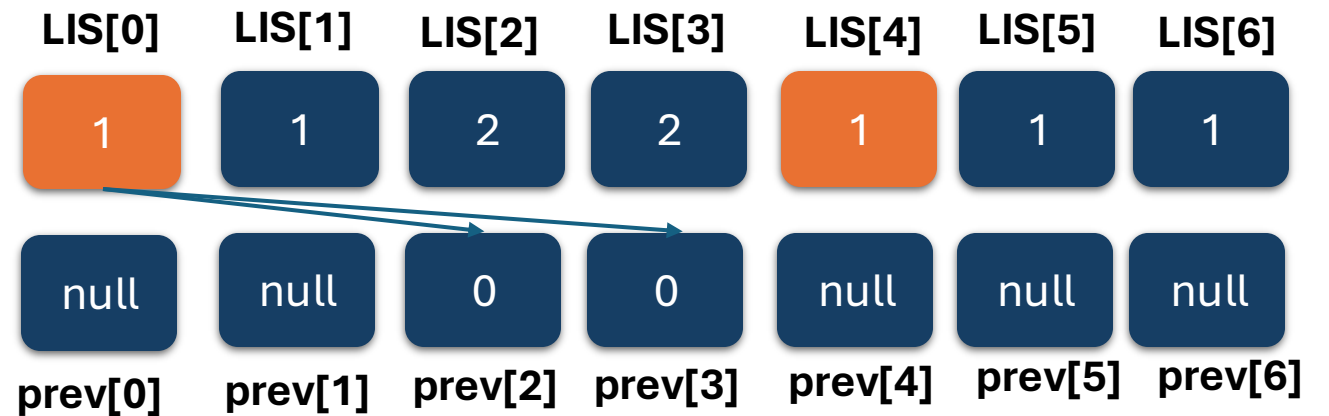
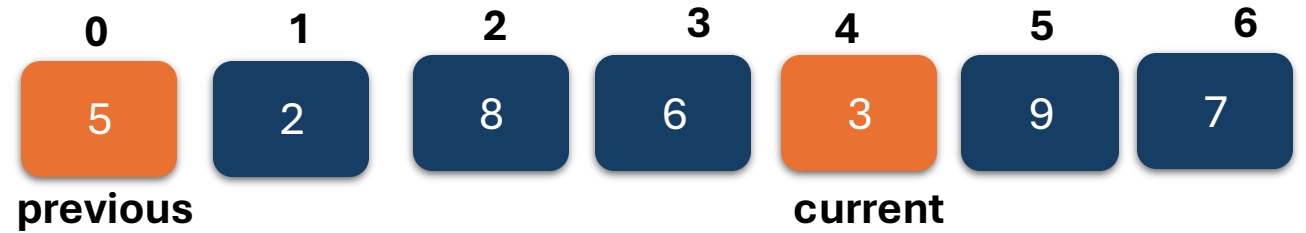
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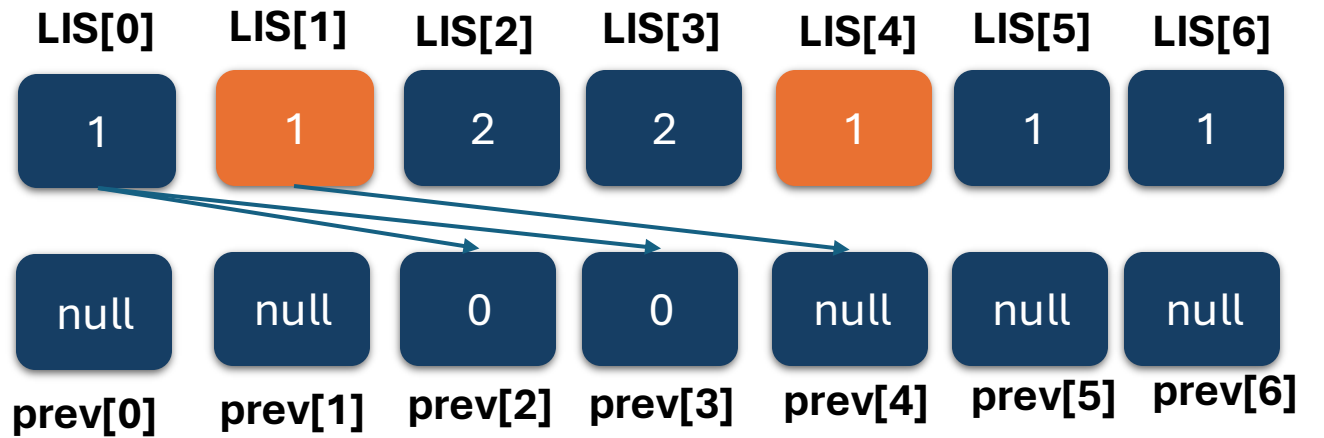
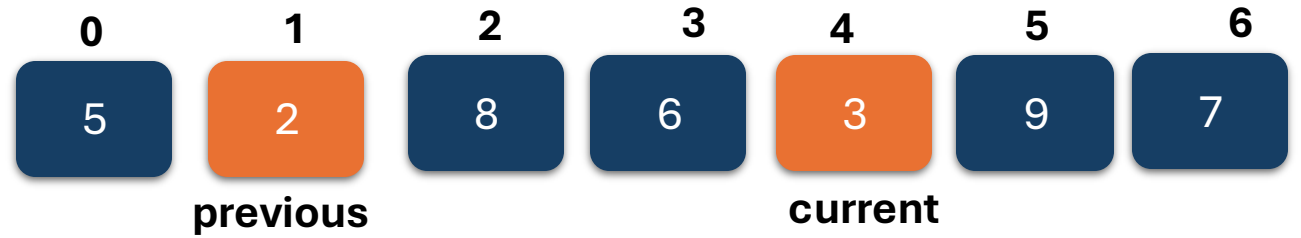
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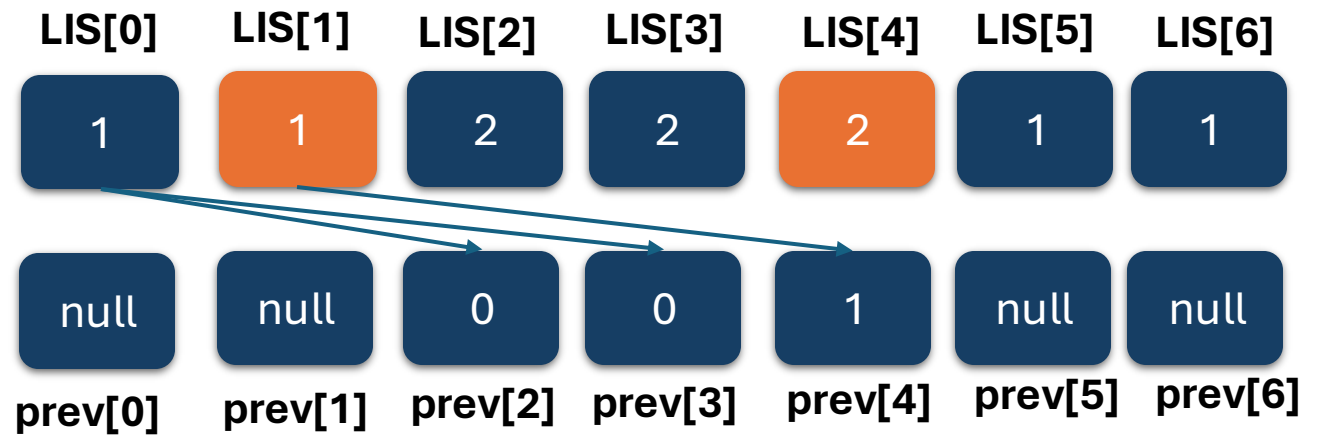
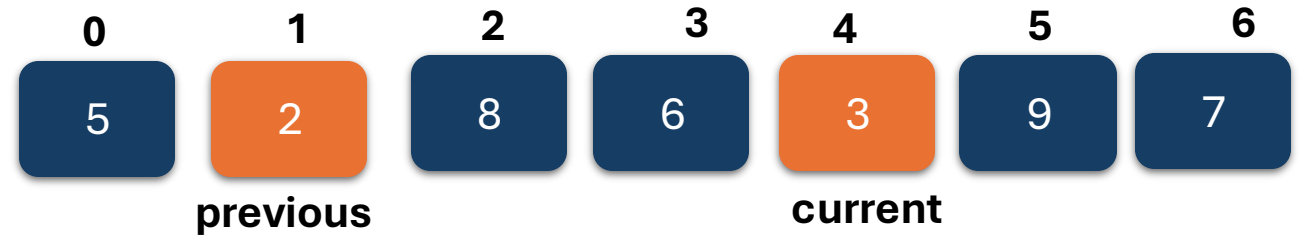
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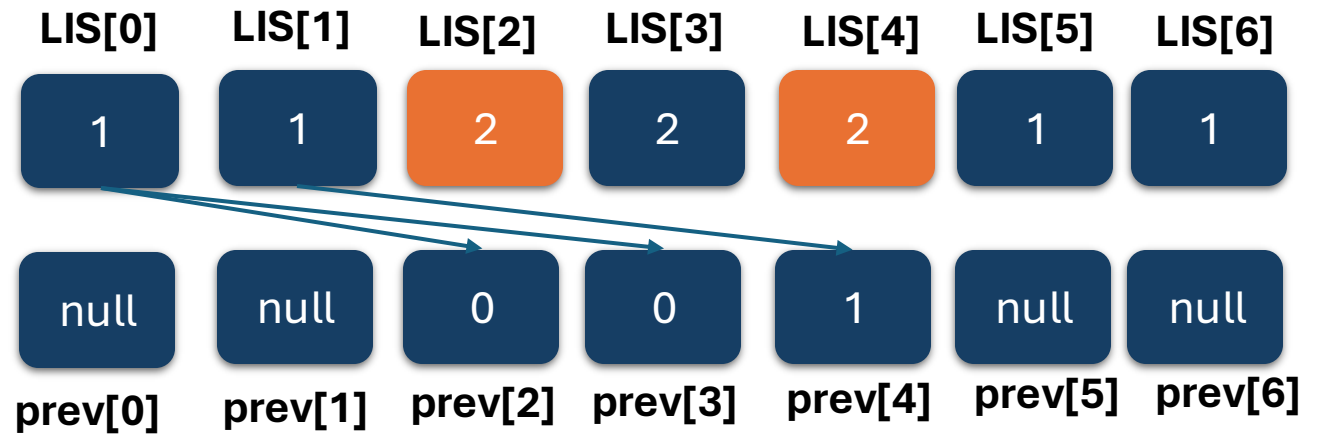
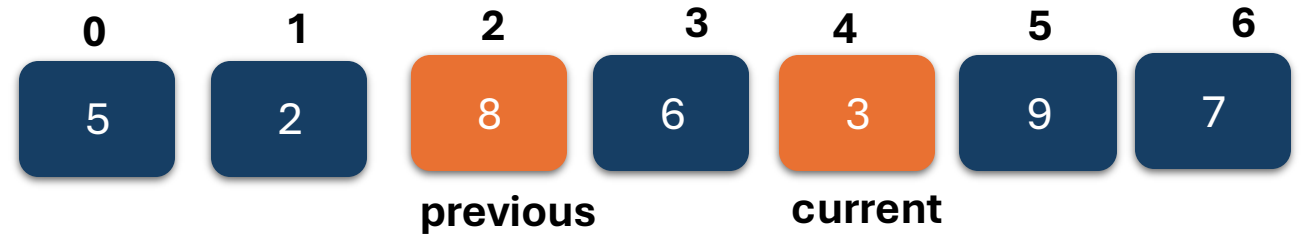
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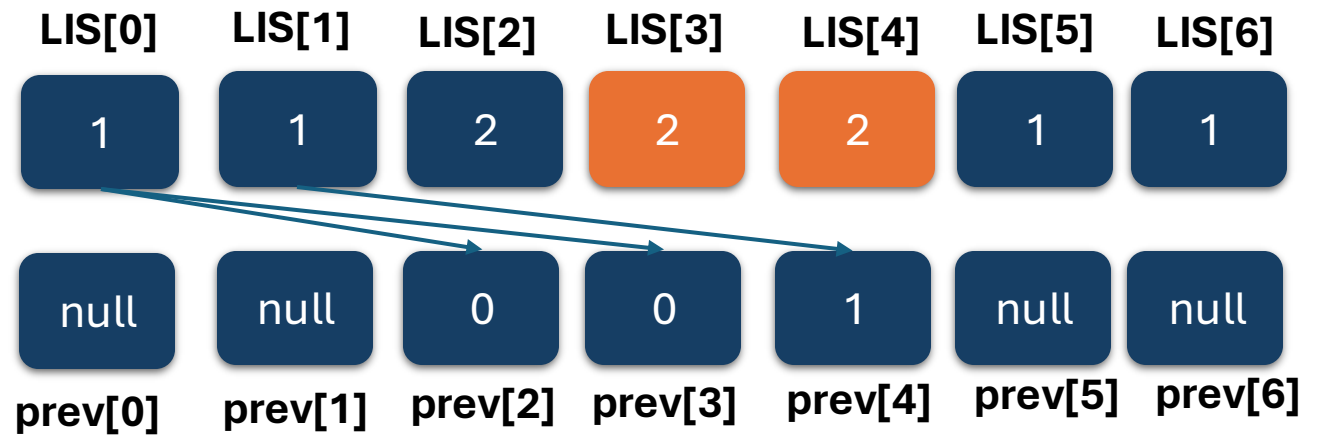
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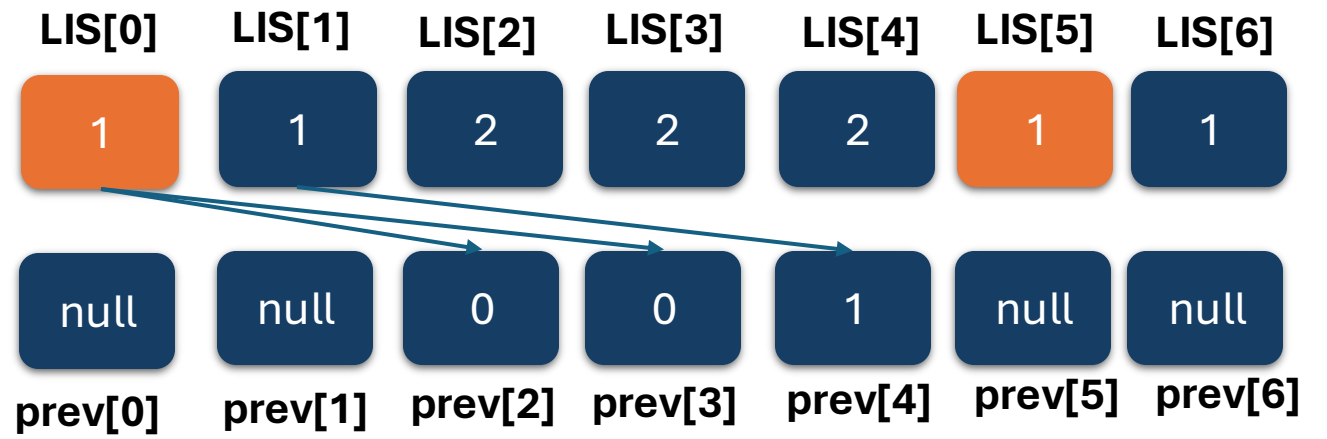
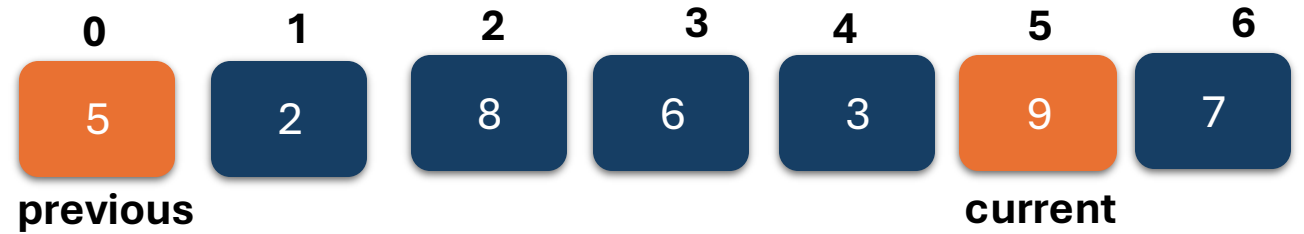
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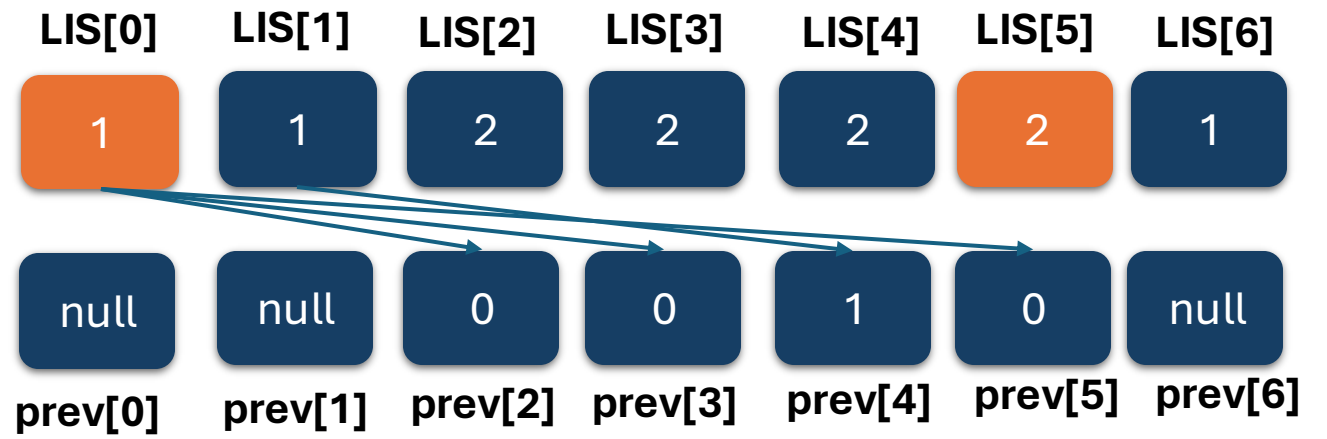
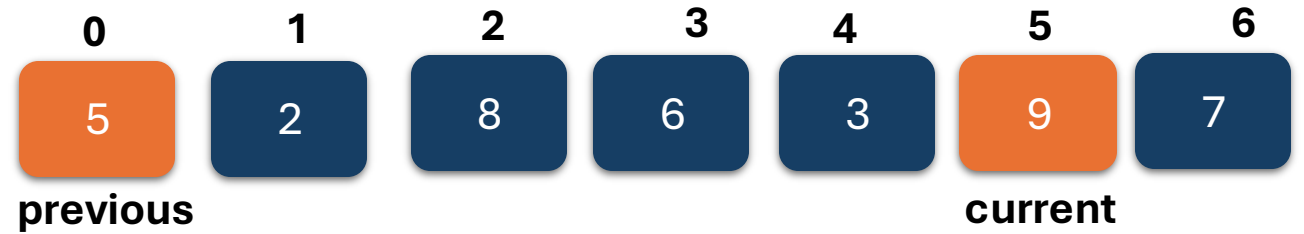
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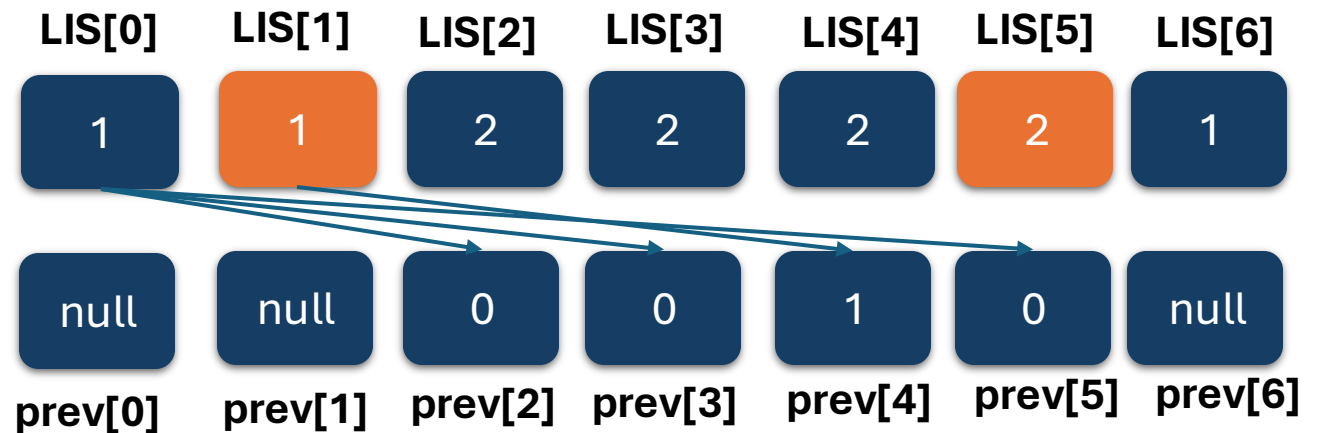
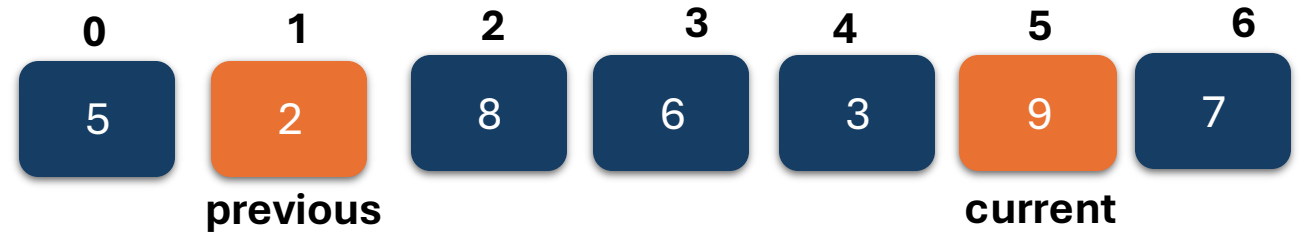
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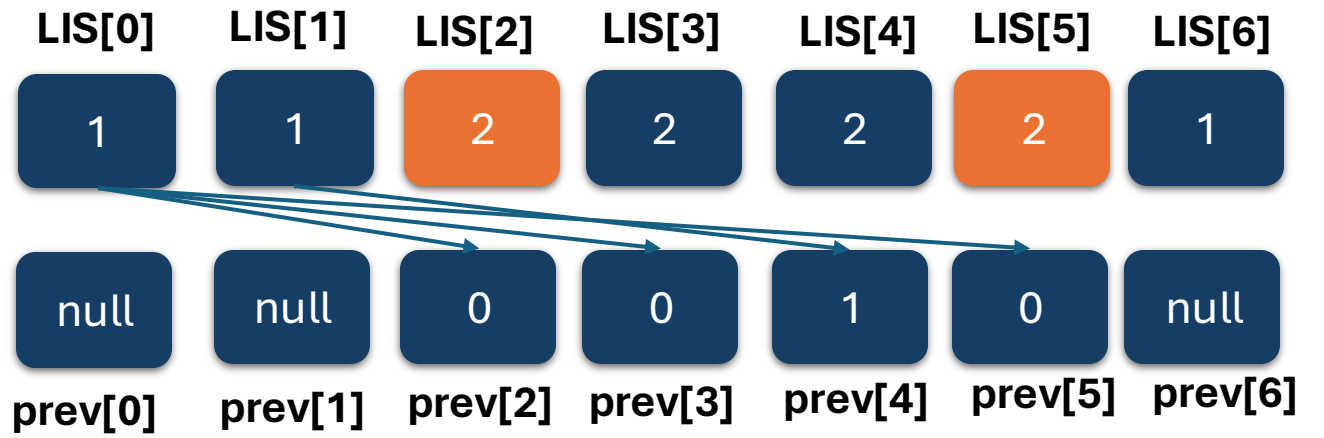
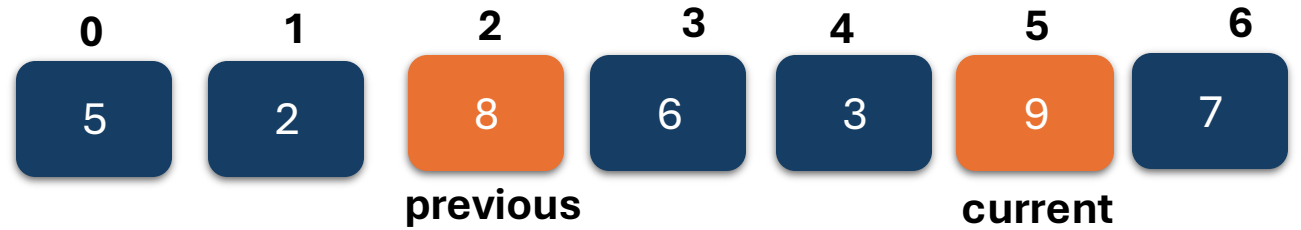
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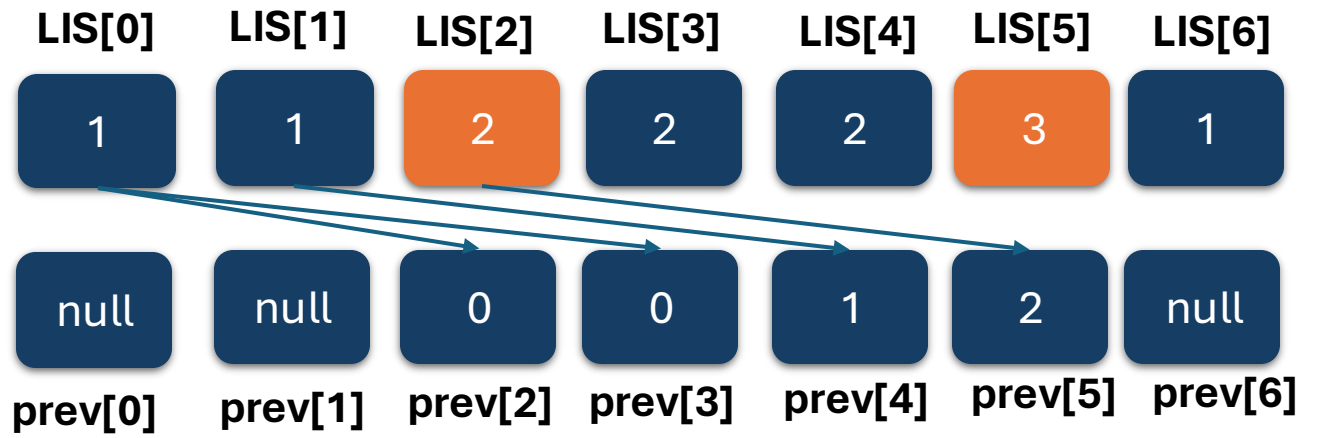
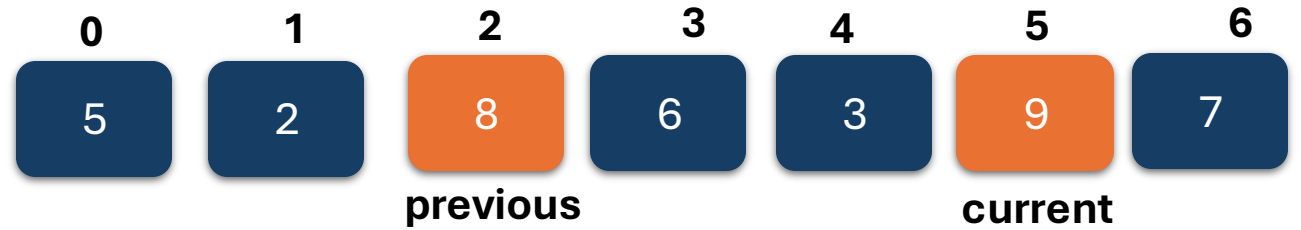
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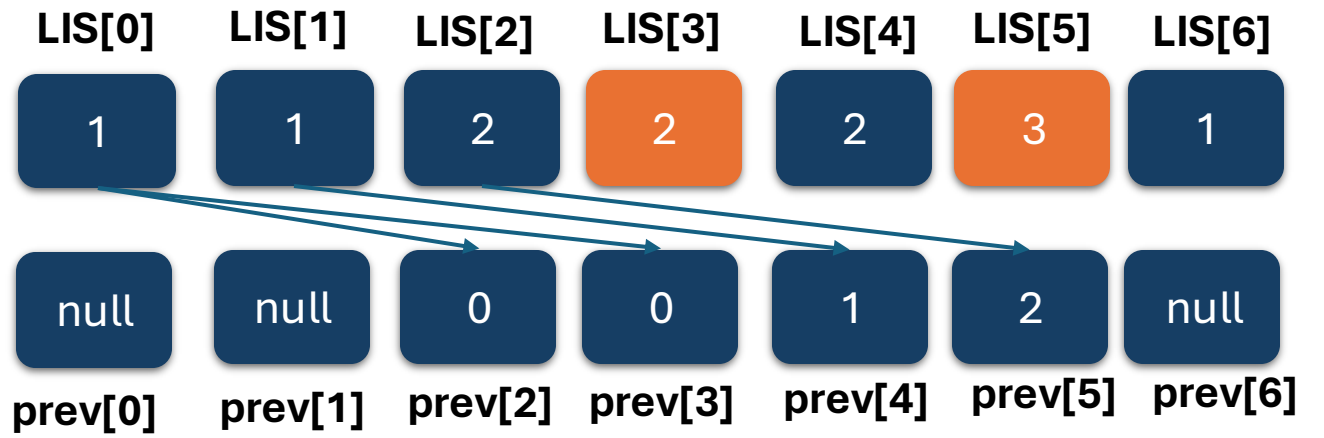
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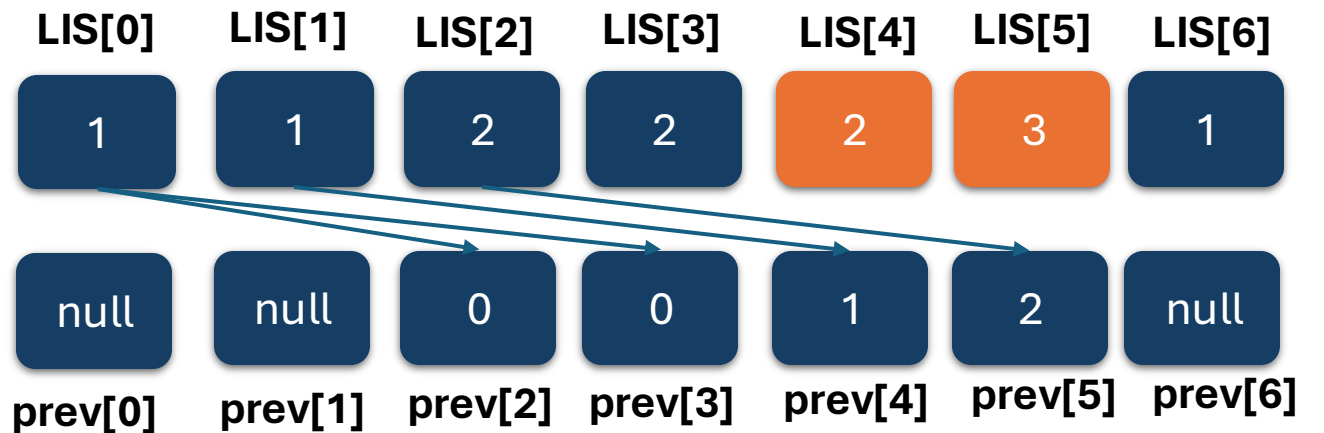
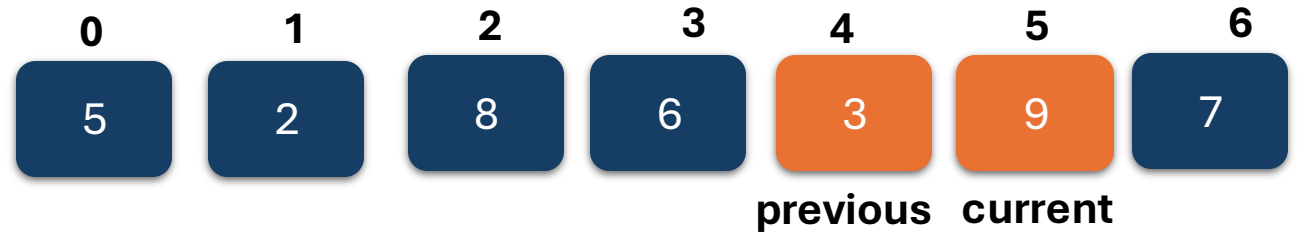
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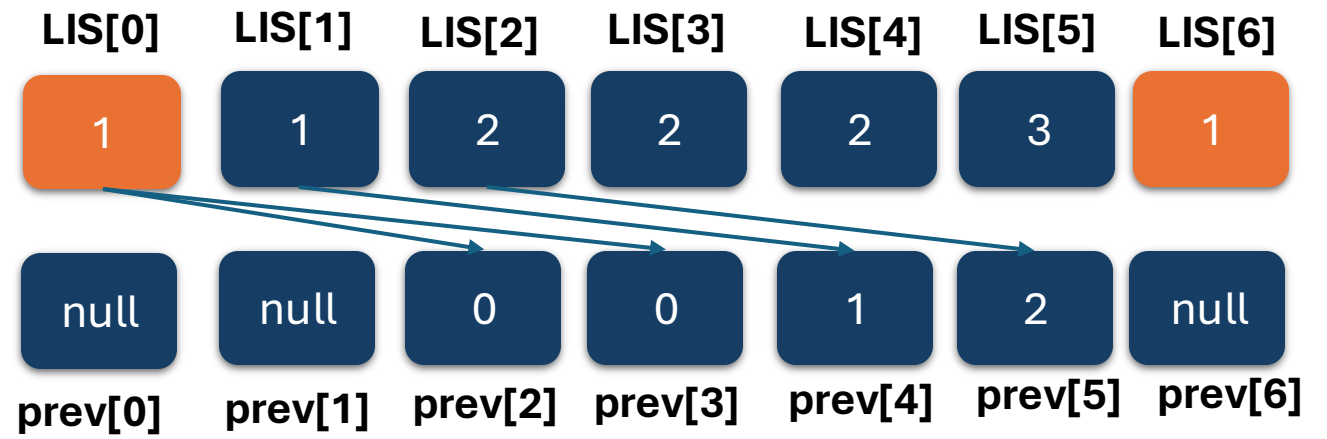
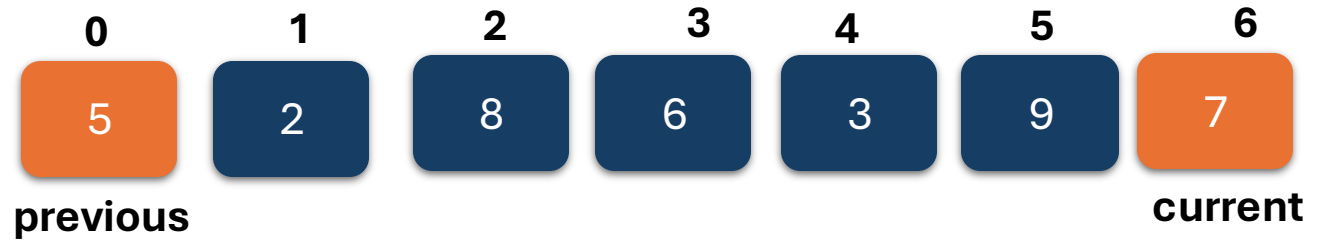
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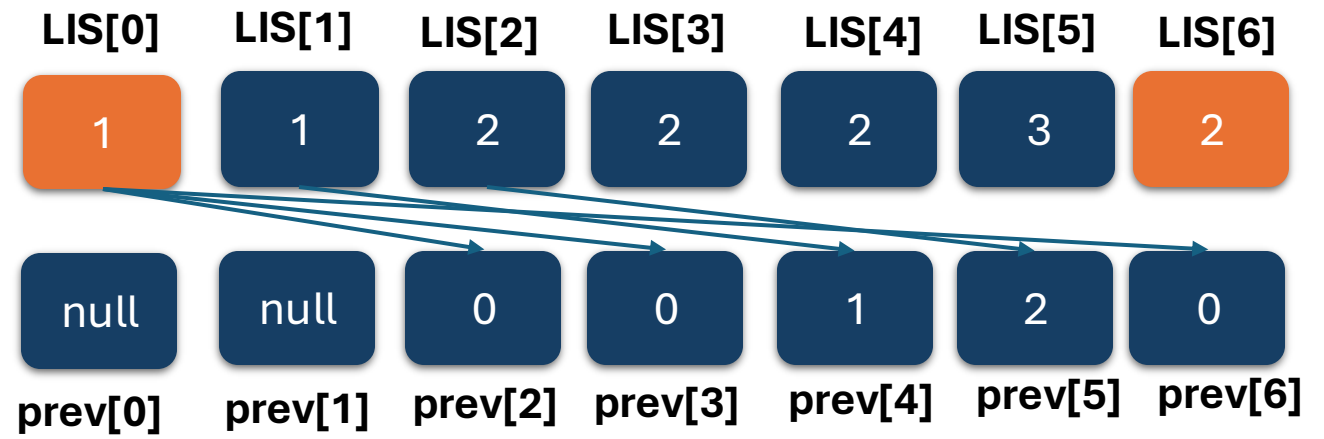
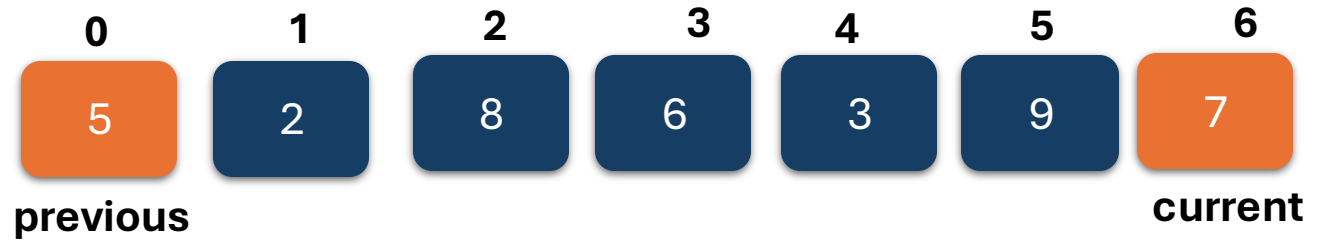
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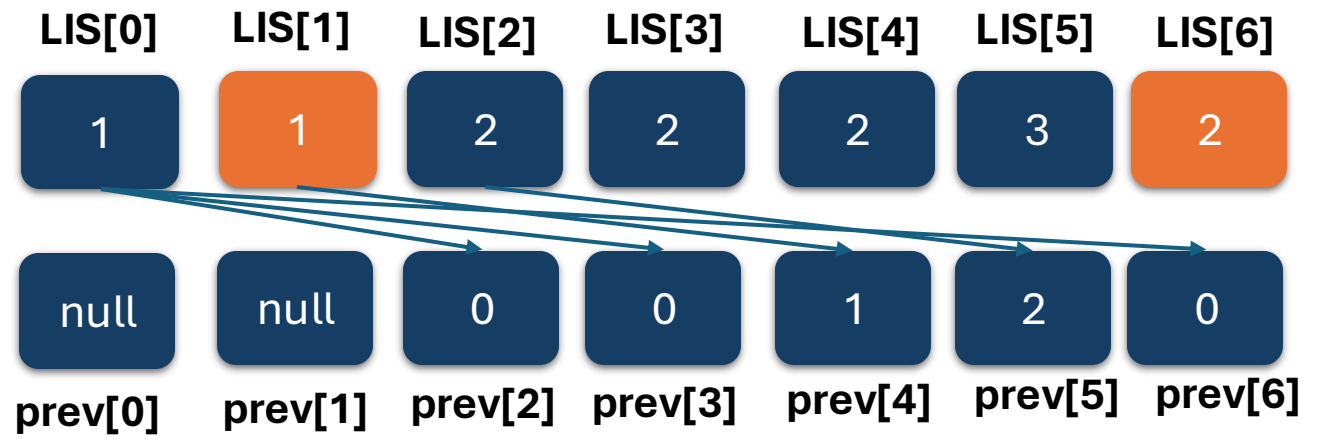
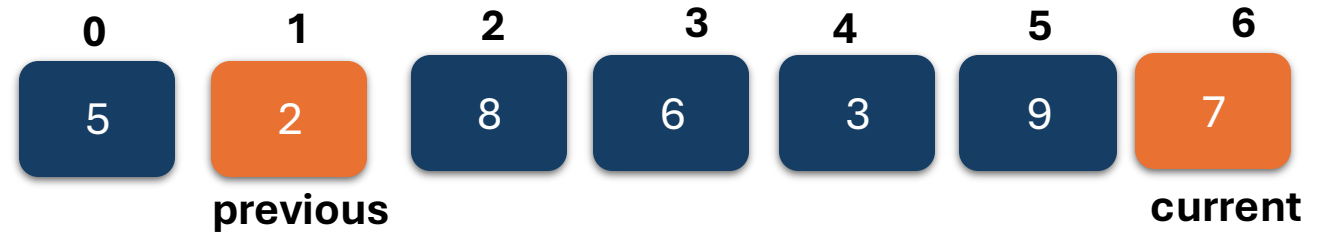
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for current = 0 to N-1:

for previous = 0 to current-1:

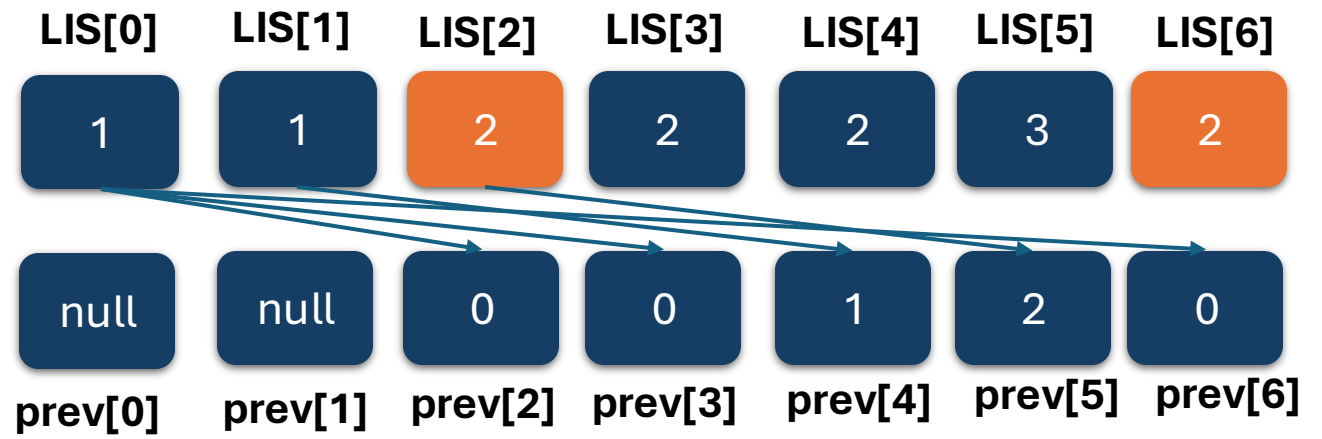
if A[current] > A[previous]:

if LIS[previous] + 1 > LIS[current]:

LIS[current] = LIS[previous] + 1

PREV[current] = previous

return max(LIS)



find_LIS_length(sequence A):

N = length(A)

LIS = array of 1s, length N

PREV = array of nulls, length N

for current = 0 to N-1:

for previous = 0 to current-1:

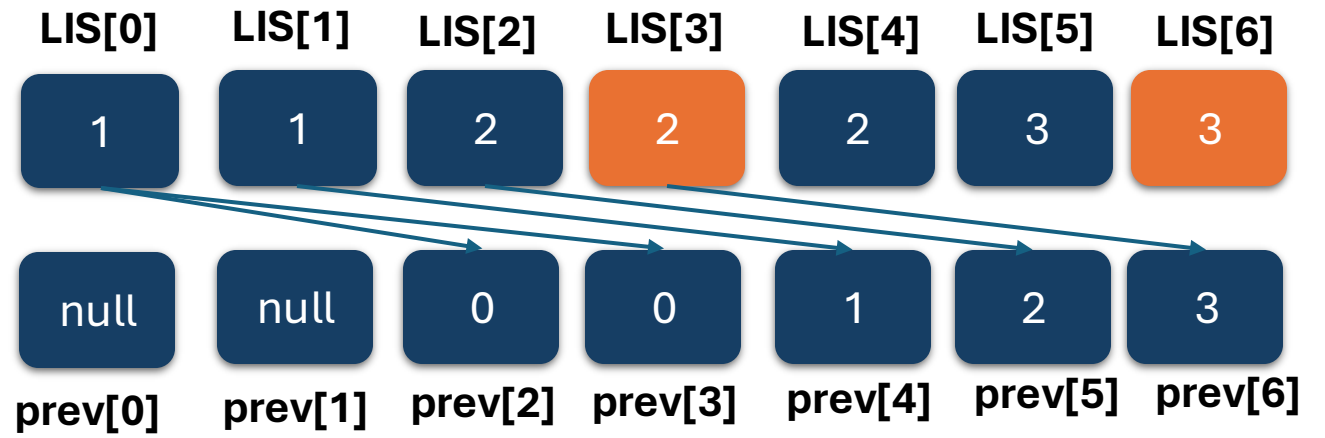
if A[current] > A[previous]:

if LIS[previous] + 1 > LIS[current]:

LIS[current] = LIS[previous] + 1

PREV[current] = previous

return max(LIS)



find_LIS_length(sequence A):

N = length(A)

LIS = array of 1s, length N

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for current = 0 to N-1:

for previous = 0 to current-1:

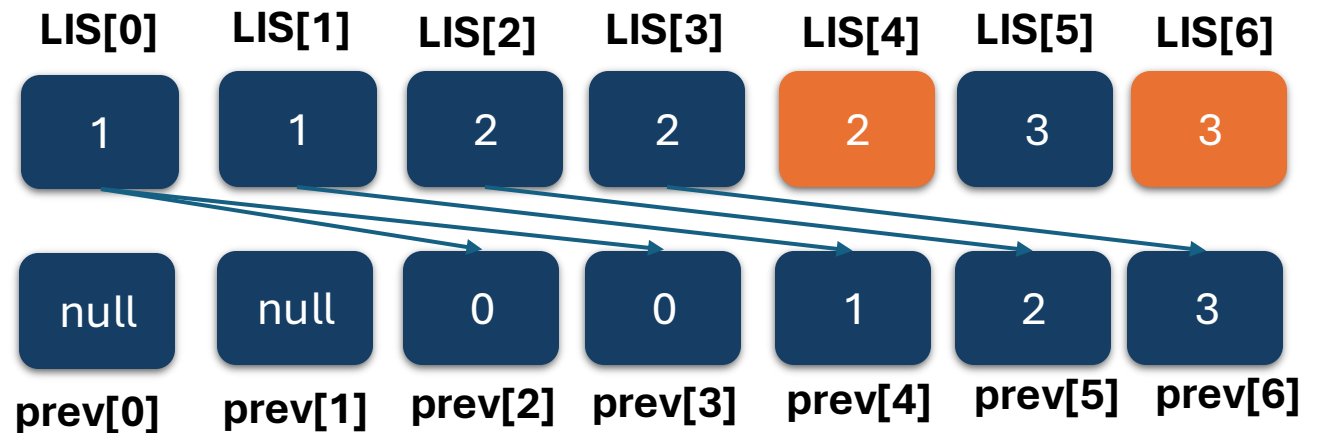
if A[current] > A[previous]:

if LIS[previous] + 1 > LIS[current]:

LIS[current] = LIS[previous] + 1

PREV[current] = previous

return max(LIS)



find_LIS_length(sequence A):

N = length(A)

LIS = array of 1s, length N

PREV = array of nulls, length N

for current = 0 to N-1:

for previous = 0 to current-1:

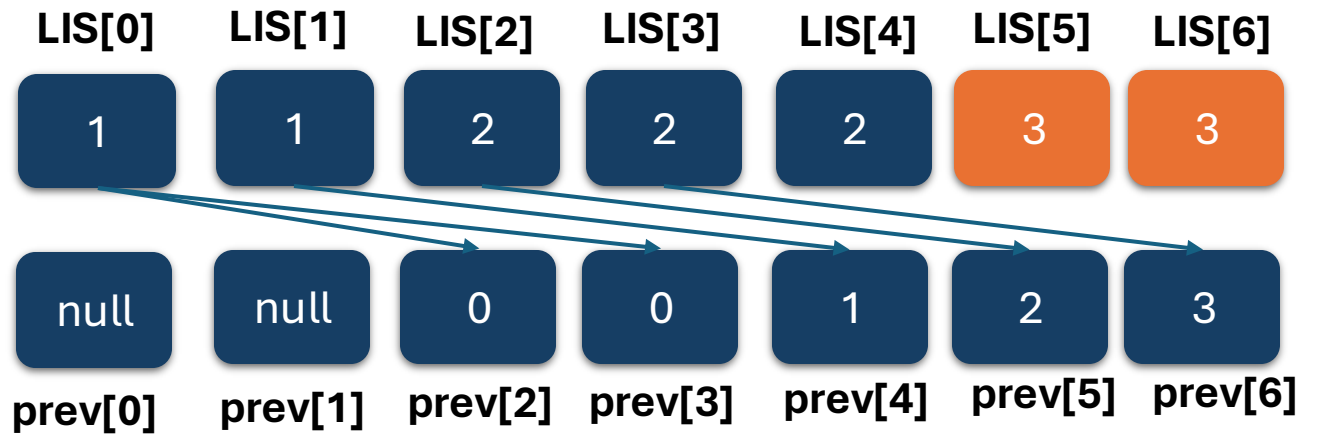
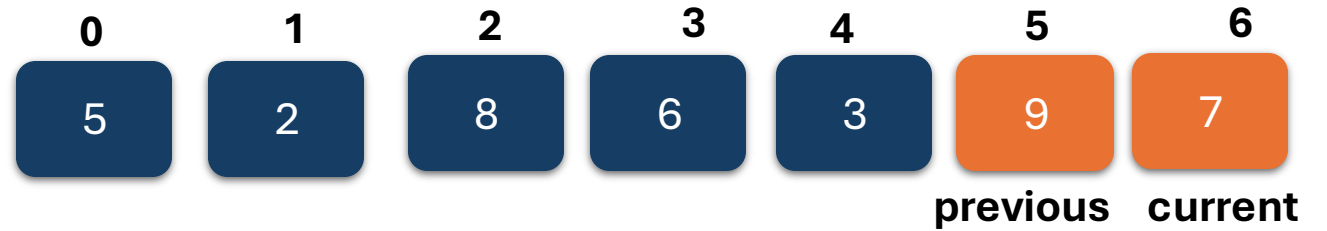
if A[current] > A[previous]:

if LIS[previous] + 1 > LIS[current]:

LIS[current] = LIS[previous] + 1

PREV[current] = previous

return max(LIS)



find_LIS_length(sequence A):

N = length(A)

LIS = array of 1s, length N

PREV = array of nulls, length N

for current = 0 to N-1:

for previous = 0 to current-1:

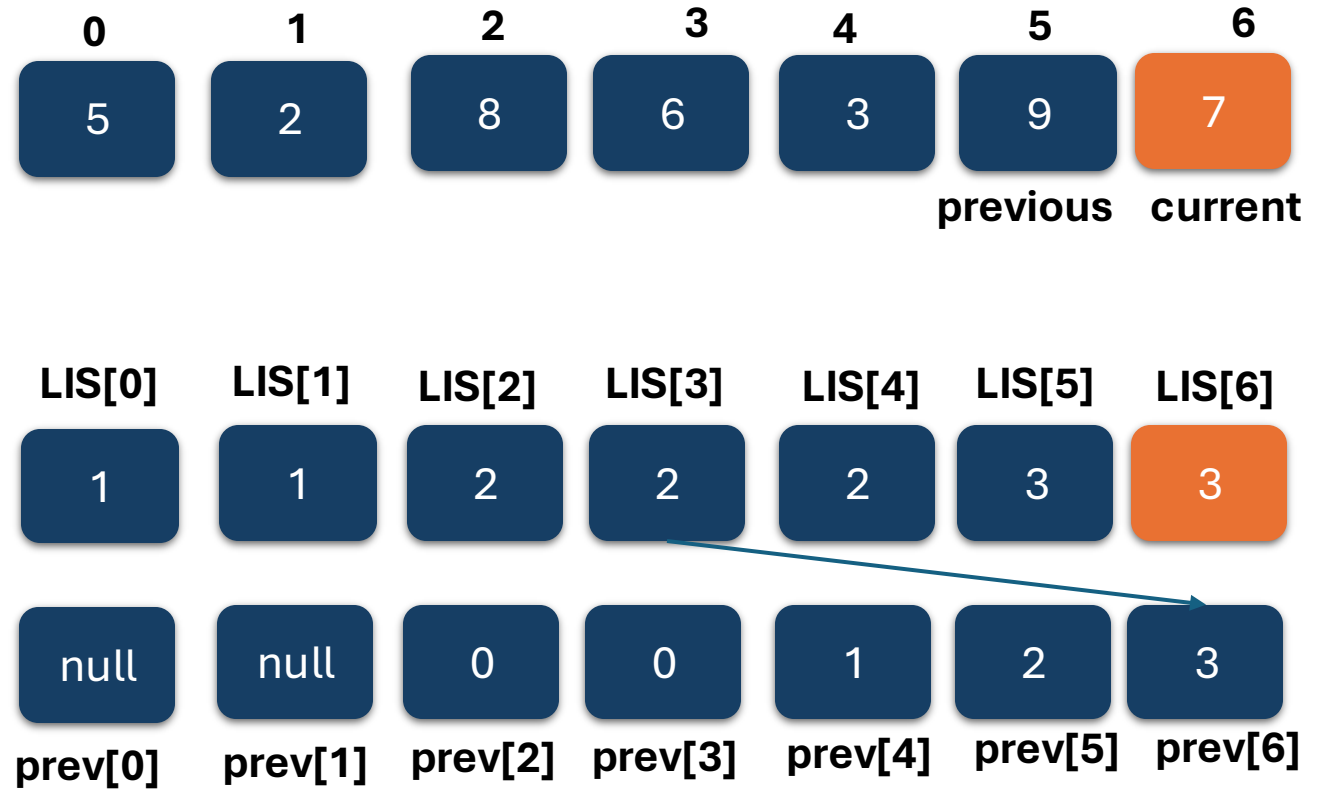
if A[current] > A[previous]:

if LIS[previous] + 1 > LIS[current]:

LIS[current] = LIS[previous] + 1

PREV[current] = previous

return max(LIS)



```
print_LIS(sequence A, array LIS, array PREV):  
    max_count = max(LIS)  
    current = LIS.indexOf(max_count)  
    numbers = [ ]  
    while current != null:  
        add A[current] to numbers  
        current = PREV[current]  
    return reversed(numbers)
```

0	1	2	3	4	5	6
5	2	8	6	3	9	7
					previous	current

LIS[0]	LIS[1]	LIS[2]	LIS[3]	LIS[4]	LIS[5]	LIS[6]
1	1	2	2	2	3	3
prev[0]	prev[1]	prev[2]	prev[3]	prev[4]	prev[5]	prev[6]
null	null	0	0	1	2	3

print_LIS(sequence A, array LIS, array PREV):

max_count = max(LIS)

current = LIS.indexOf(max_count) 6

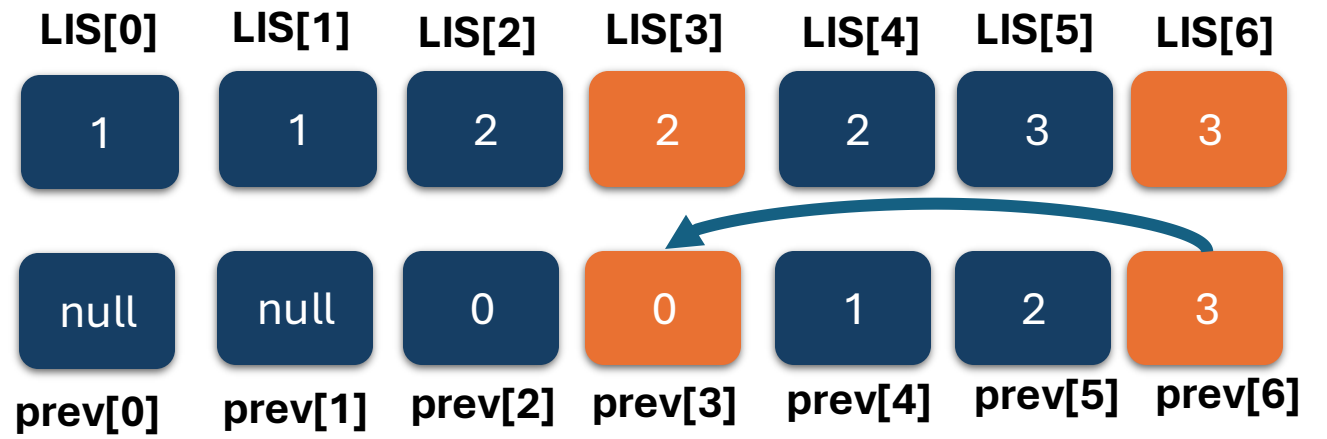
numbers = []

while current != null:

add A[current] to numbers A[6]

current = PREV[current] Current = PREV[6]=3

return reversed(numbers)



print_LIS(sequence A, array LIS, array PREV):

max_count = max(LIS)

current = LIS.indexOf(max_count) **6**

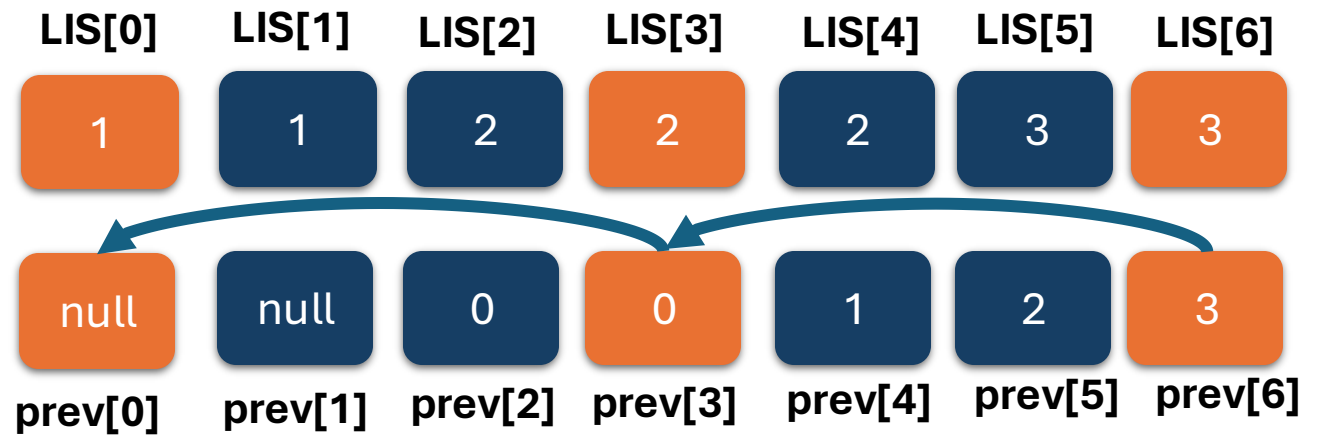
numbers = []

while current != null:

add A[current] to numbers **A[3]**

current = PREV[current] **Current = PREV[3]=0**

return reversed(numbers)



print_LIS(sequence A, array LIS, array PREV):

max_count = max(LIS)

current = LIS.indexOf(max_count) **6**

numbers = []

while current != null:

add A[current] to numbers **A[0]**

current = PREV[current] **Current = PREV[0]=null**

return reversed(numbers)

Another Example

```
previous  current
```

10	22	9	33	21	50	41	60
1	1	1	1	1	1	1	1

[illegible]

Example

previous current

10	22	9	33	21	50	41	60
1	2	1	1	1	1	1	1

null	0	null	null	null	null	null	null
------	---	------	------	------	------	------	------

Example

previous

current

10	22	9	33	21	50	41	60
1	2	1	1	1	1	1	1

null	0	null	null	null	null	null	null
------	---	------	------	------	------	------	------

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	1	1	1	1	1
null	0	null	null	null	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	1	1	1	1	1
null	0	null	null	null	null	null	null

Example

previous

current

10	22	9	33	21	50	41	60
1	2	1	1	1	1	1	1

null	0	null	null	null	null	null	null
------	---	------	------	------	------	------	------

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	2	1	1	1	1
null	0	null	0	null	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	1	1	1	1
null	0	null	1	null	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	1	1	1	1
null	0	null	1	null	null	null	null

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	1	1	1	1
null	0	null	1	null	null	null	null

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	1	1	1
null	0	null	1	0	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	2	1	1	1
null	0	null	1	0	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	2	1	1	1
null	0	null	1	0	null	null	null

Example

previous current

10	22	9	33	21	50	41	60
1	2	1	3	2	1	1	1

null	0	null	1	0	null	null	null
------	---	------	---	---	------	------	------

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	1	1	1
null	0	null	1	0	null	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	2	2	1	1
null	0	null	1	0	0	null	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	2	3	1	1
null	0	null	1	0	1	null	null

Example

previous			current				
10	22	9	33	21	50	41	60
1	2	1	3	2	3	1	1
null	0	null	1	0	1	null	null

Example

previous			current				
10	22	9	33	21	50	41	60
1	2	1	3	2	4	1	1
null	0	null	1	0	3	null	null

Example

previous current

10	22	9	33	21	50	41	60
1	2	1	3	2	4	1	1

null	0	null	1	0	3	null	null
------	---	------	---	---	---	------	------

Example

previous			current				
10	22	9	33	21	50	41	60
1	2	1	3	2	4	2	1
null	0	null	1	0	3	0	null

Example

previous		current					
10	22	9	33	21	50	41	60
1	2	1	3	2	4	3	1
null	0	null	1	0	3	1	null

Example

previous			current				
10	22	9	33	21	50	41	60
1	2	1	3	2	4	3	1
null	0	null	1	0	3	1	null

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	1
null	0	null	1	0	3	3	null

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	1
null	0	null	1	0	3	3	null

Example

previous current

10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	1

null	0	null	1	0	3	3	null
------	---	------	---	---	---	---	------

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	1
null	0	null	1	0	3	3	null

Example

previous								current
10	22	9	33	21	50	41	60	
1	2	1	3	2	4	4	2	
null	0	null	1	0	3	3	0	

Example

previous								current
10	22	9	33	21	50	41	60	
1	2	1	3	2	4	4	3	
null	0	null	1	0	3	3	1	

Example

previous			current				
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	3
null	0	null	1	0	3	3	1

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	4
null	0	null	1	0	3	3	3

Example

previous				current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	4
null	0	null	1	0	3	3	3

Example

					previous	current	
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	5
null	0	null	1	0	3	3	5

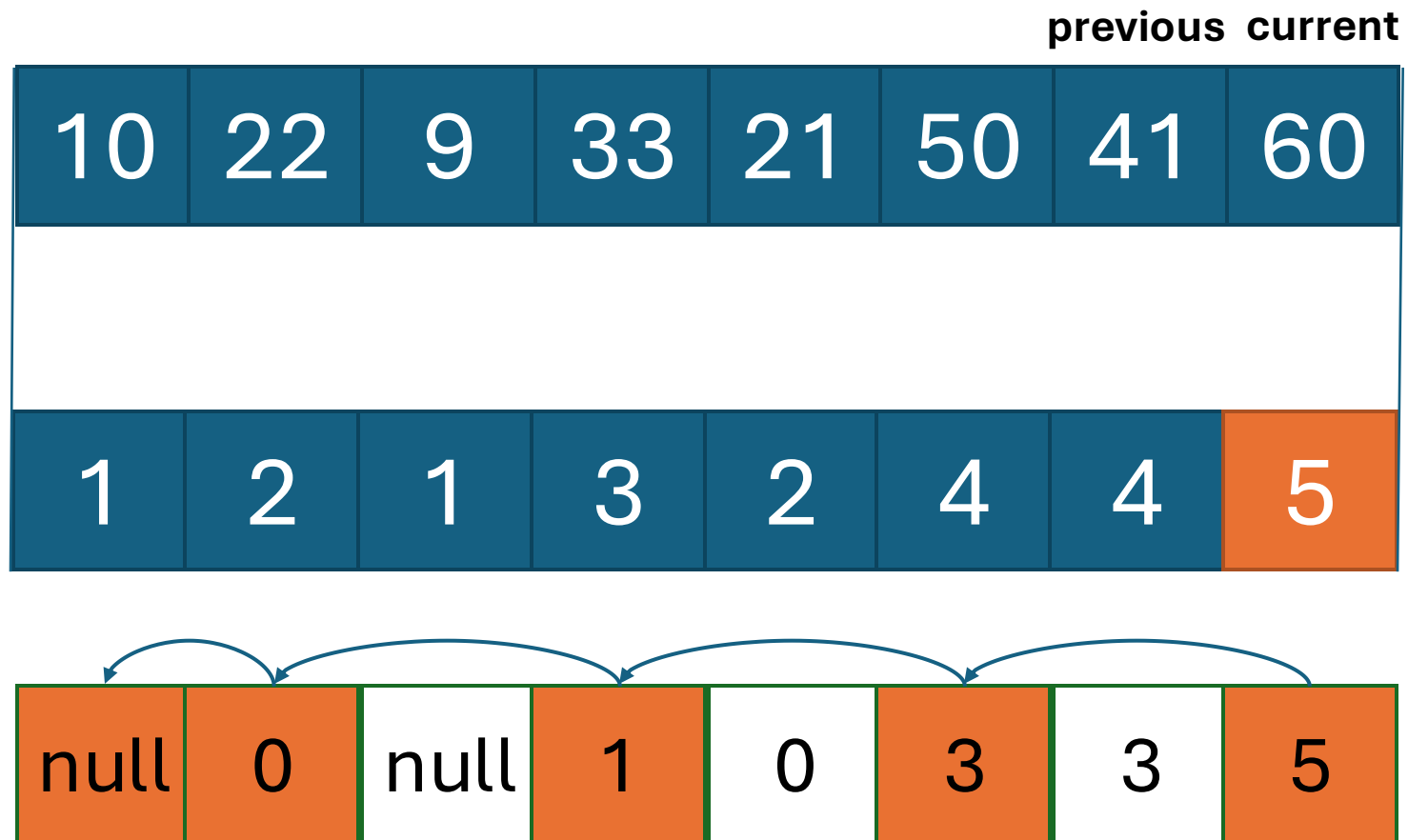
Example

						previous	current
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	5
null	0	null	1	0	3	3	5

Example

				previous current			
10	22	9	33	21	50	41	60
1	2	1	3	2	4	4	5
null	0	null	1	0	3	3	5

Example



DP Solution

- Initial Invocation: LIS_LENGTH
- Algorithm takes $O(n^2)$ time

End of Lecture