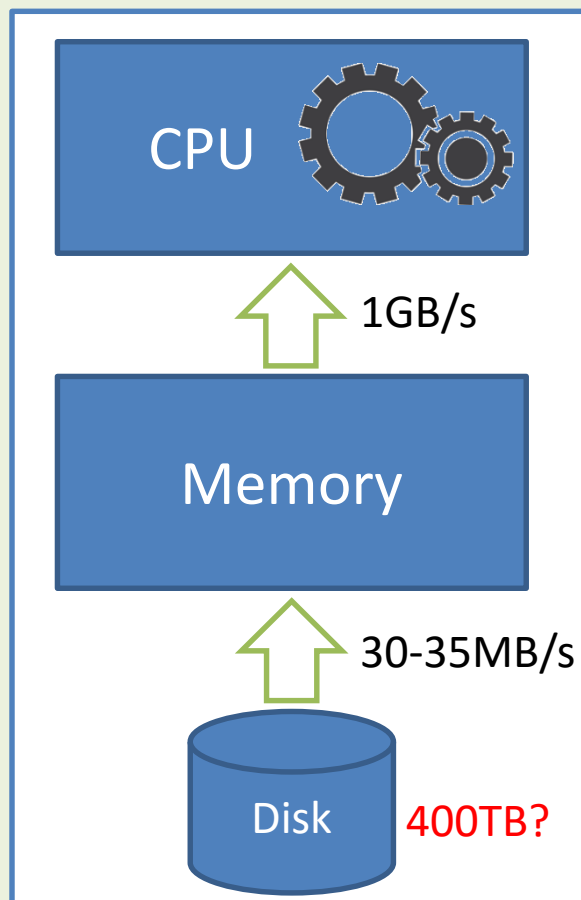


# The introduction to MapReduce

Dang Ha The Hien  
*PhD. UiO*  
*eSmart Systems*

# Single Node Architecture

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Traditional data analysis and machine learning algorithms:

- Load all required data from disk to memory
- Run all the algorithms on memory

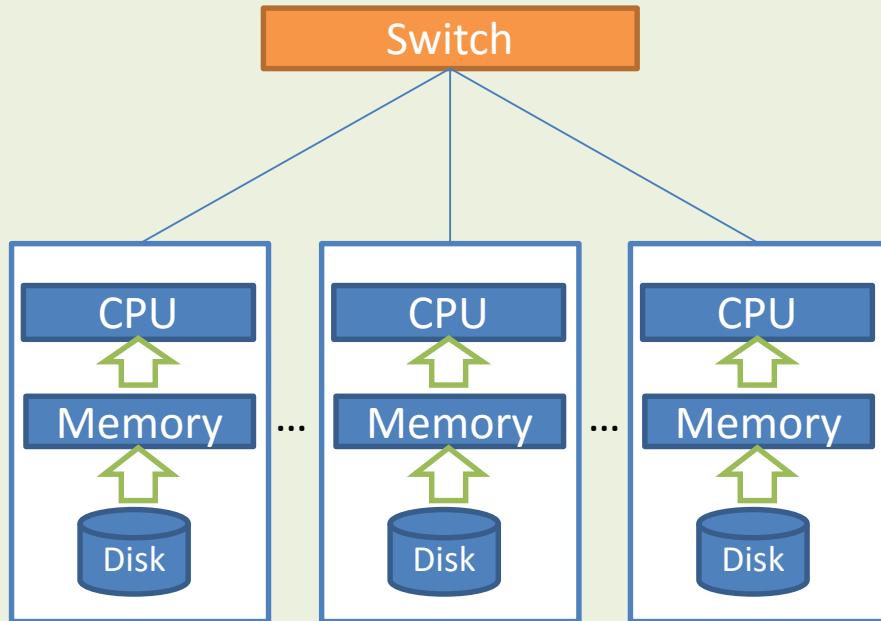
Google Problem:

- They have 20+ billion web pages
- Each web page is about 20KB (real number at the time)  
→ 400 TB+ in total
- It takes ~4 months to read the data from disk to memory
- Take longer to do something useful with the data

# Distribute data and computation over large cluster

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1Gbps between  
any pair of nodes

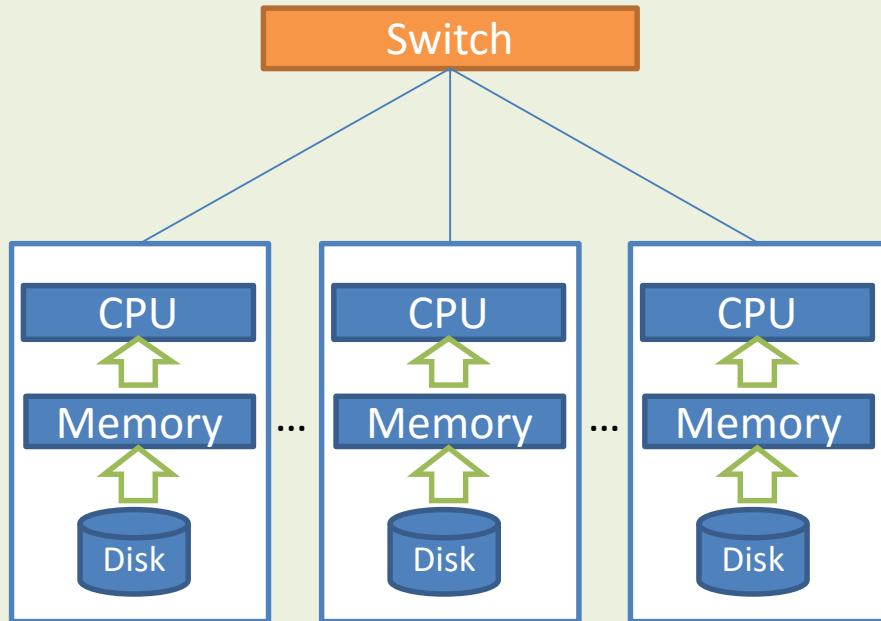


Recently standard architecture for big data problems:

- Cluster of consumer-grade hardware
- Many desktop-like Linux servers
  - Easy to add capacity
  - Cheaper per CPU/disk
- Commodity Network (Ethernet) to connect them

# Distribute data and computation over large cluster

1Gbps between  
any pair of nodes



One rack contains 16-64 nodes

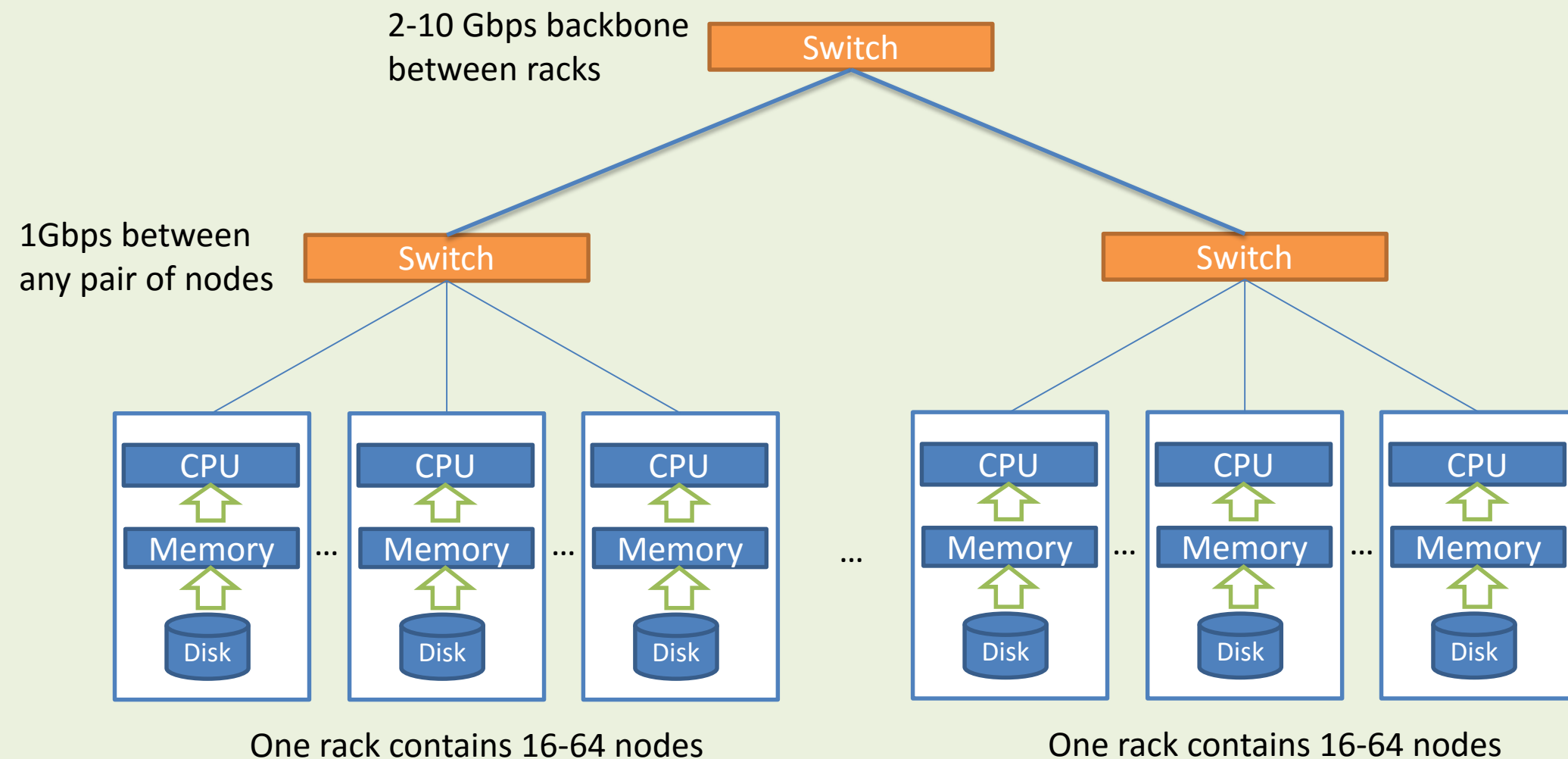


Each node is a cheap Linux server  
Easy to add capacity  
Easy to replace

One rack can only contains around 16-64 nodes, since more than that, there will be a lot of collision in while transferring data through the switch.

# Cluster Architecture

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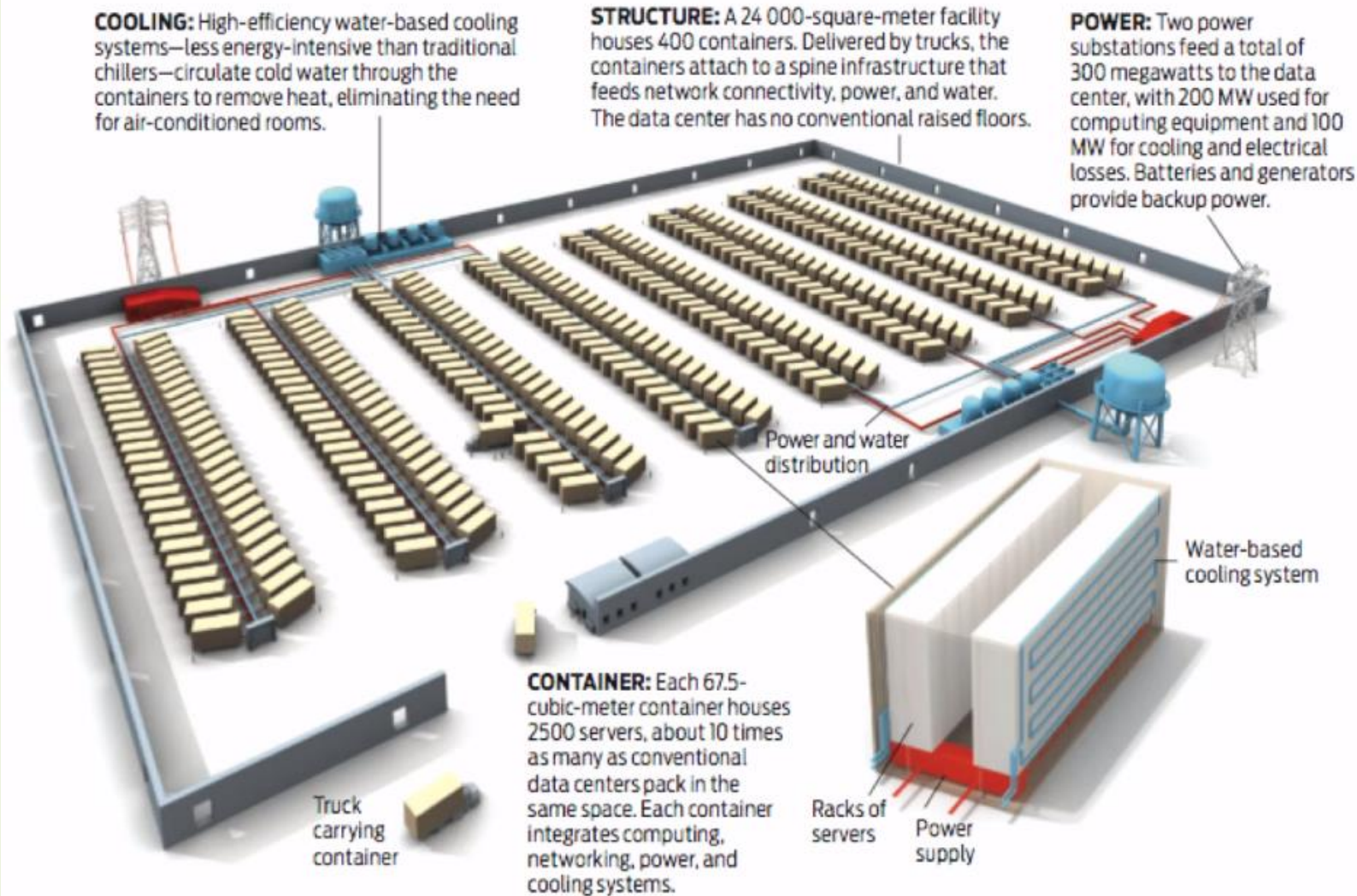
# Cluster Architecture - Example

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# The Million-Server Data Center



<http://spectrum.ieee.org/tech-talk/semiconductors/devices/what-will-the-data-center-of-the-future-look-like>

In 2011, it was estimated that Google had 1M machines

# Problems with Cheap Hardware

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## Machines fail:

- One server may stay up 3 years (1,000 days)
- If you have 1,000 servers, expect to loose 1/day
- With 1M machines, 1,000 servers fail every day!

## Uneven performance:

- Those machines that don't completely fail but just ask for job and do it slowly  
(even bigger problem)

## Network speed is much slower than shared memory:

- Copy data over a network takes time

Distributed Programming is hard! We need a simple model that hides most of complexity



# Idea and Solution

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## Idea:

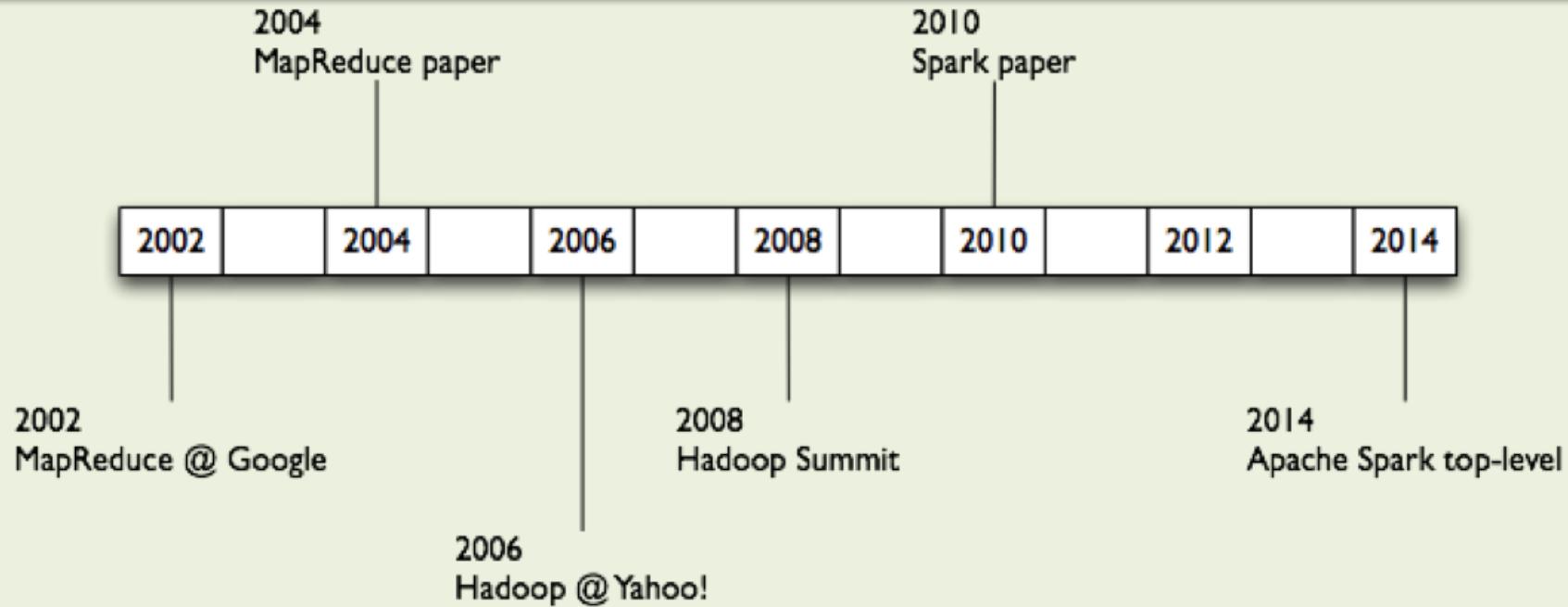
- Store data redundantly on multiple nodes for persistence and availability
- Move computation close to data to minimize data movement
- Simple programming model to hide the complexity of all this magic

## MapReduce addresses these problems:

- Storage Infrastructure (Hadoop Distributed File System- HDFS)
- Map-Reduce programming model.

# History of Distributed Programming

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2002: Google start using MapReduce

2004: MapReduce paper was published “Google MapReduce: Simplified Data Processing on Large Clusters”

2006: Apache Hadoop, originating from the Yahoo!’s Nutch Project

2008: Yahoo! Web scale search indexing – Hadoop Summit, Hadoop User Group

2009: Cloud computing with Amazon Web Services Elastic MapReduce,  
a Hadoop version modified for Amazon Elastic Cloud Computing (EC2)

# Distributed File System

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Example: Hadoop's HDFS, Google's GFS

- Support redundancy and availability

Typical usage pattern:

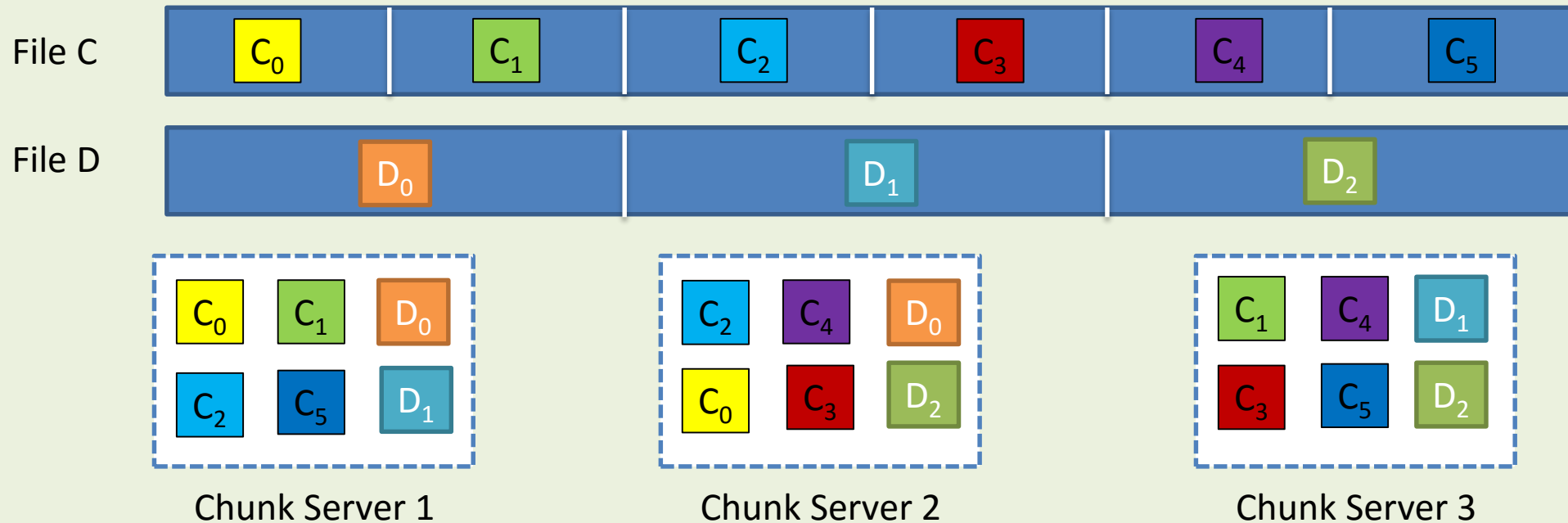
- Huge files (100s of GB or TB)
- Data is **rarely updated in place** (No random access)
- **Reads and Appends are common**

Once collected, you **don't update data** inside a file. You just append to the file, or read the whole file for analyzing.

# Distributed File System

---

This is what happens when you store a file in an HDFS cluster:



Data kept in “chunks” spread across machines (chunk nodes)

Each chunk replicated on different machines

Chunk servers also serve as compute servers → We can bring computation to data

# Distributed File System Summary

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## Chunk Servers:

- File split into continuous chunks (16-64MBs)
- Each chunk replicated (2x or 3x)
- Try to keep at least 1 replica in a different rack

## Master node:

- Name node in Hadoop's HDFS
- Stores metadata about where files are stored
- Might be replicated (but usually not)

## Client library for file access:

- Talk to master to find chunk servers
- Connects directly to chunk servers to access data



# Programming Model: MapReduce

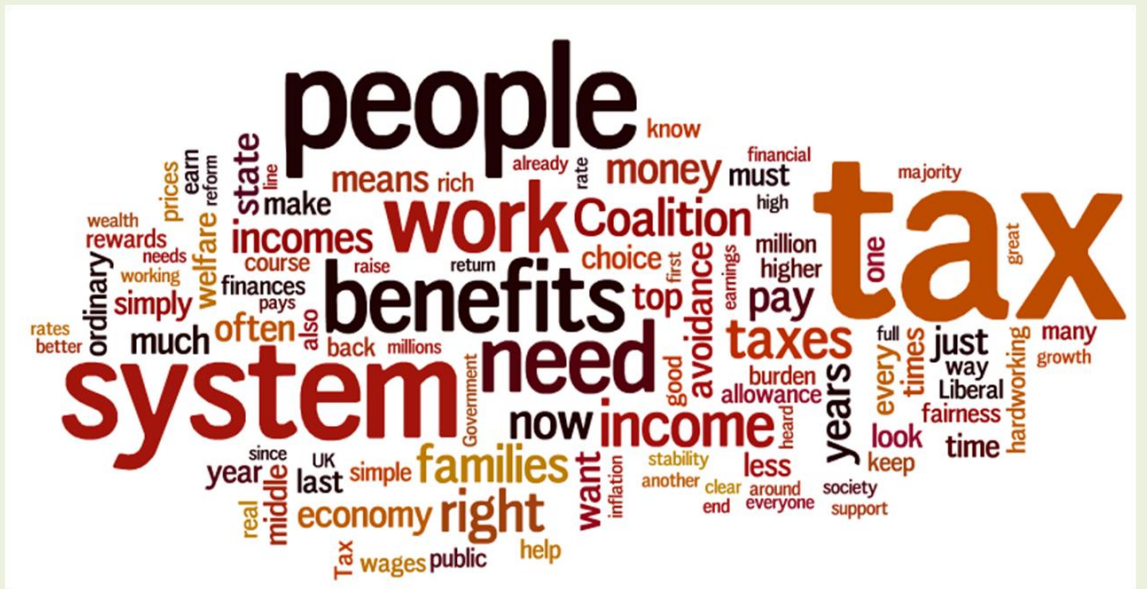
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## Classic Motivation Example:

- You have a huge text document
- Want to count the number of times each distinct word appears in the file

## Potential Applications:

- Draw a wordcloud like this
- Analyze web server log files to find popular URLs



# Simple Approach: Use a Hash Table

---

"I am Sam  
I am Sam  
Sam I am  
Do you like  
Green eggs and ham  
I do not like them  
Sam I am  
I do not like  
Green eggs and ham  
Would you like them  
Here or there?  
..."

Key	Value
I	2
am	1
Sam	1

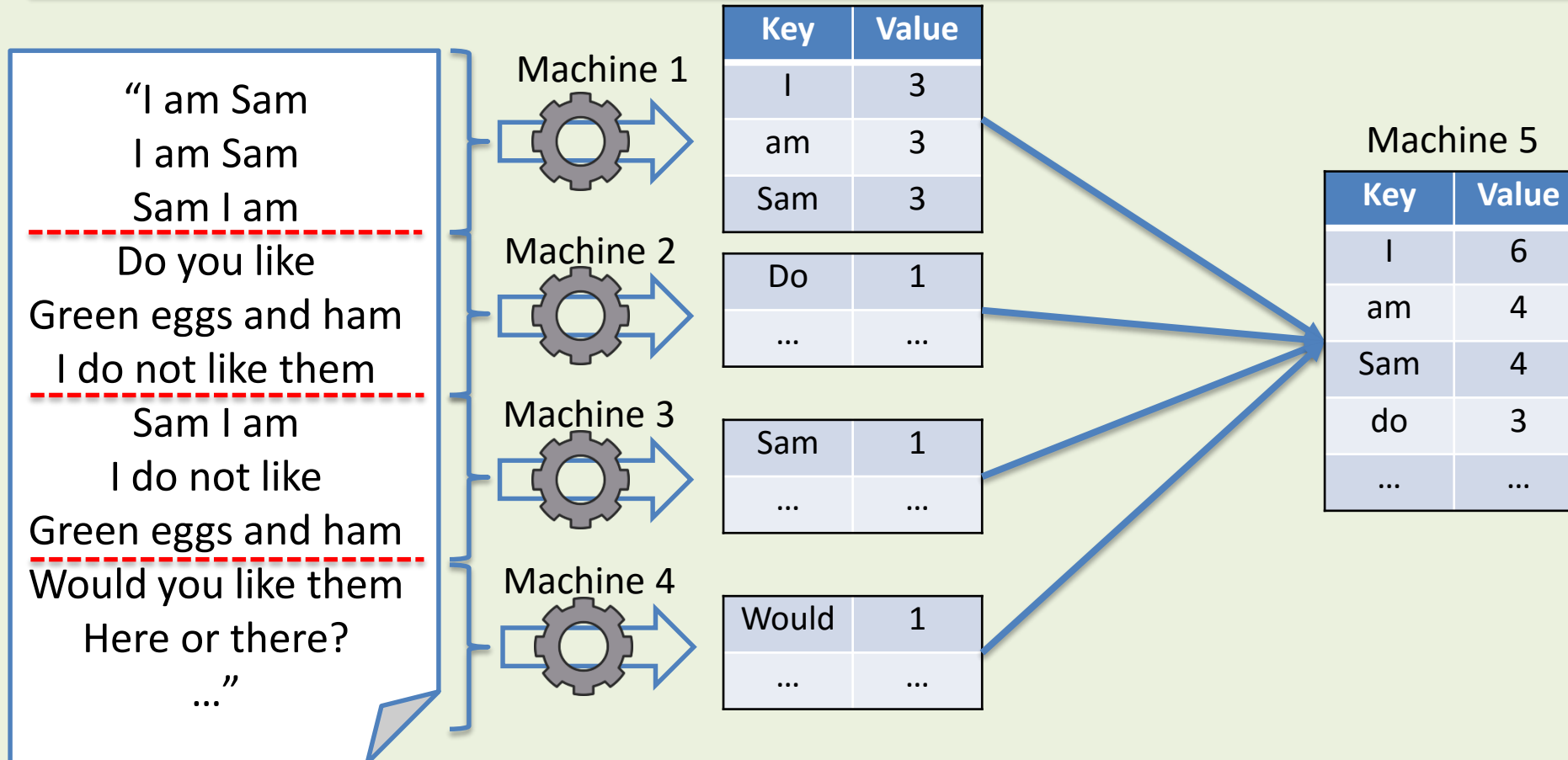
**Hash table** is a data structure that can map **keys** to **values**

Simple word count approach using hash table:

- Start with empty hash table.
- Read words sequentially
- For each word:
  - If I can't see that word as a key in the hash table
    - Add that word as a key with value = 1
  - Otherwise, increase the value of that word to +1

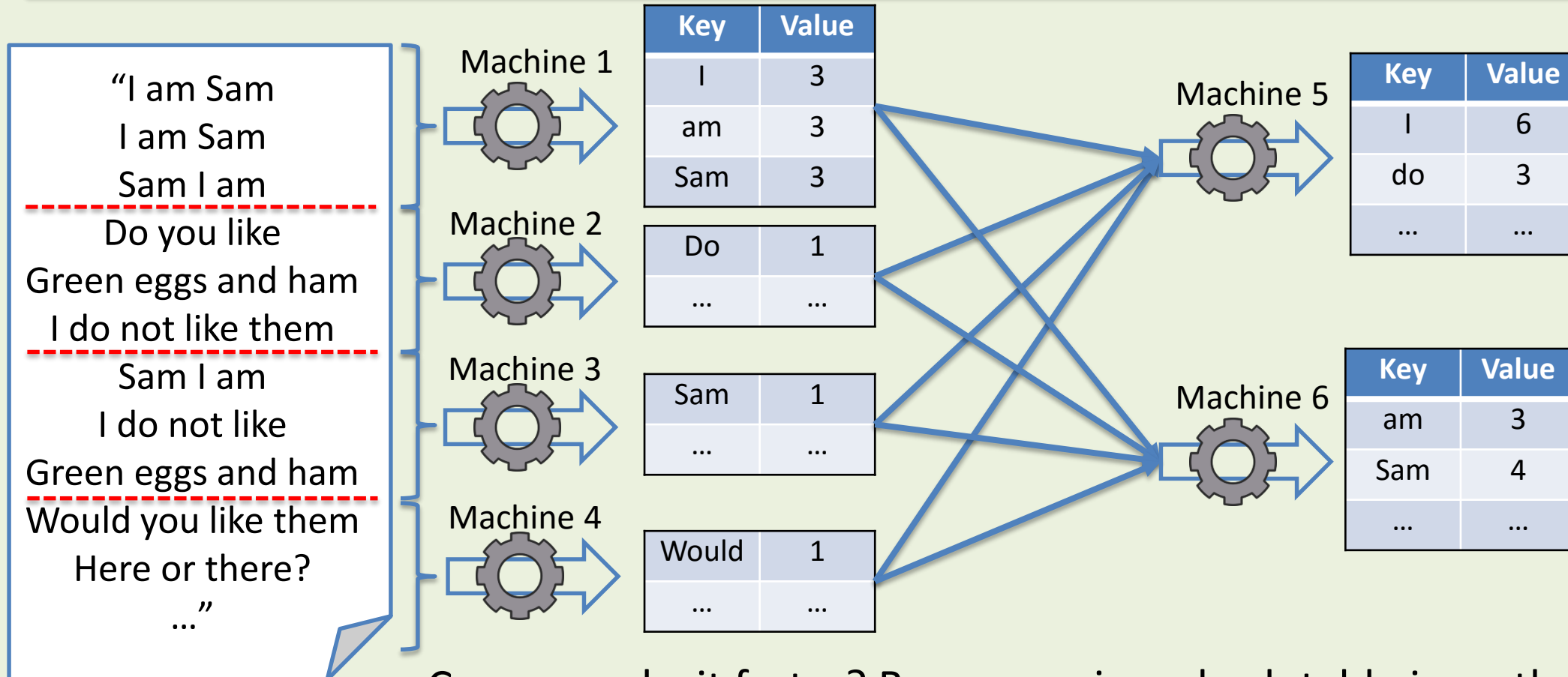
## What if the Document is Really Big?

# What if the Document is Really Big



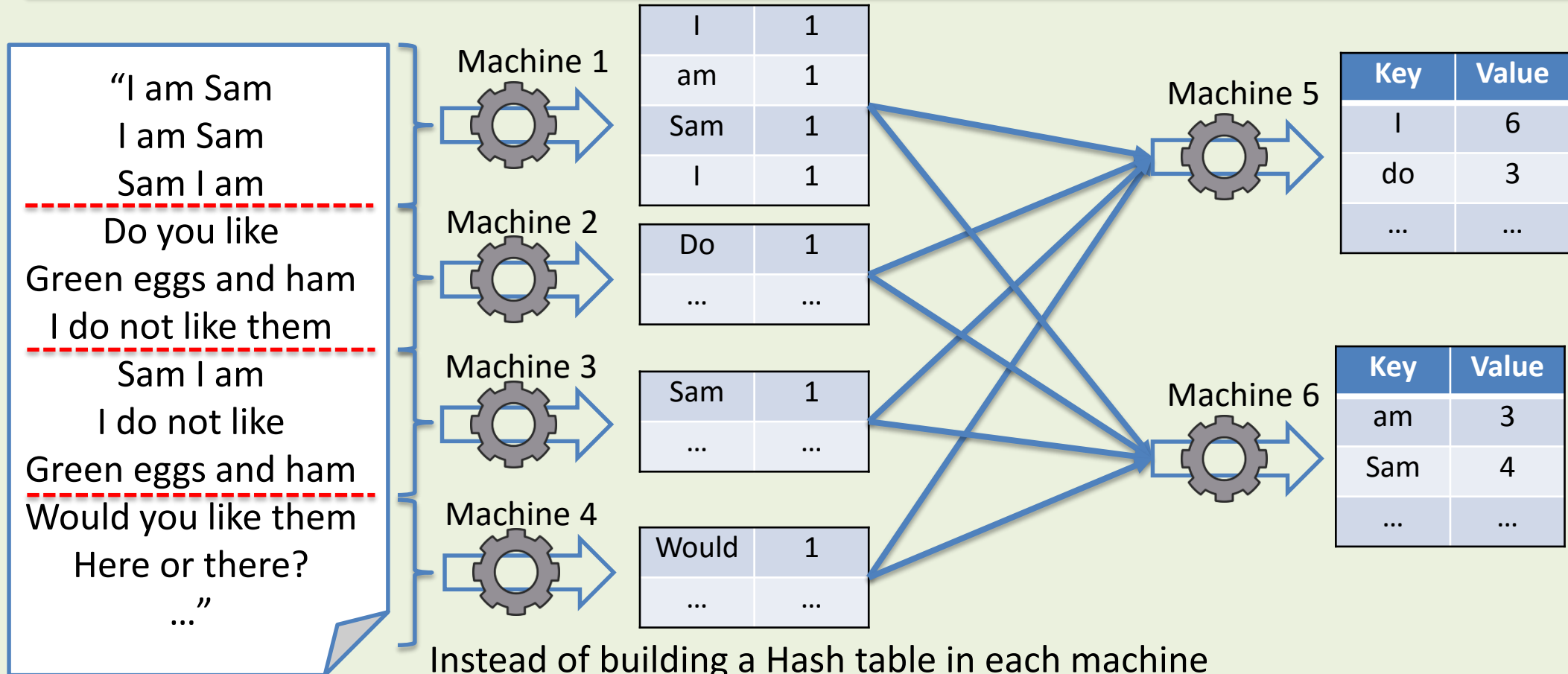
What if the Document is really Big and we have bottleneck at Machine 5?  
Or if the result too big that cannot even fit into one machine memory?

# What if the Document is Really Big?



Can we make it faster? Because using a hash table is costly  
(need to call the hash function for every access)

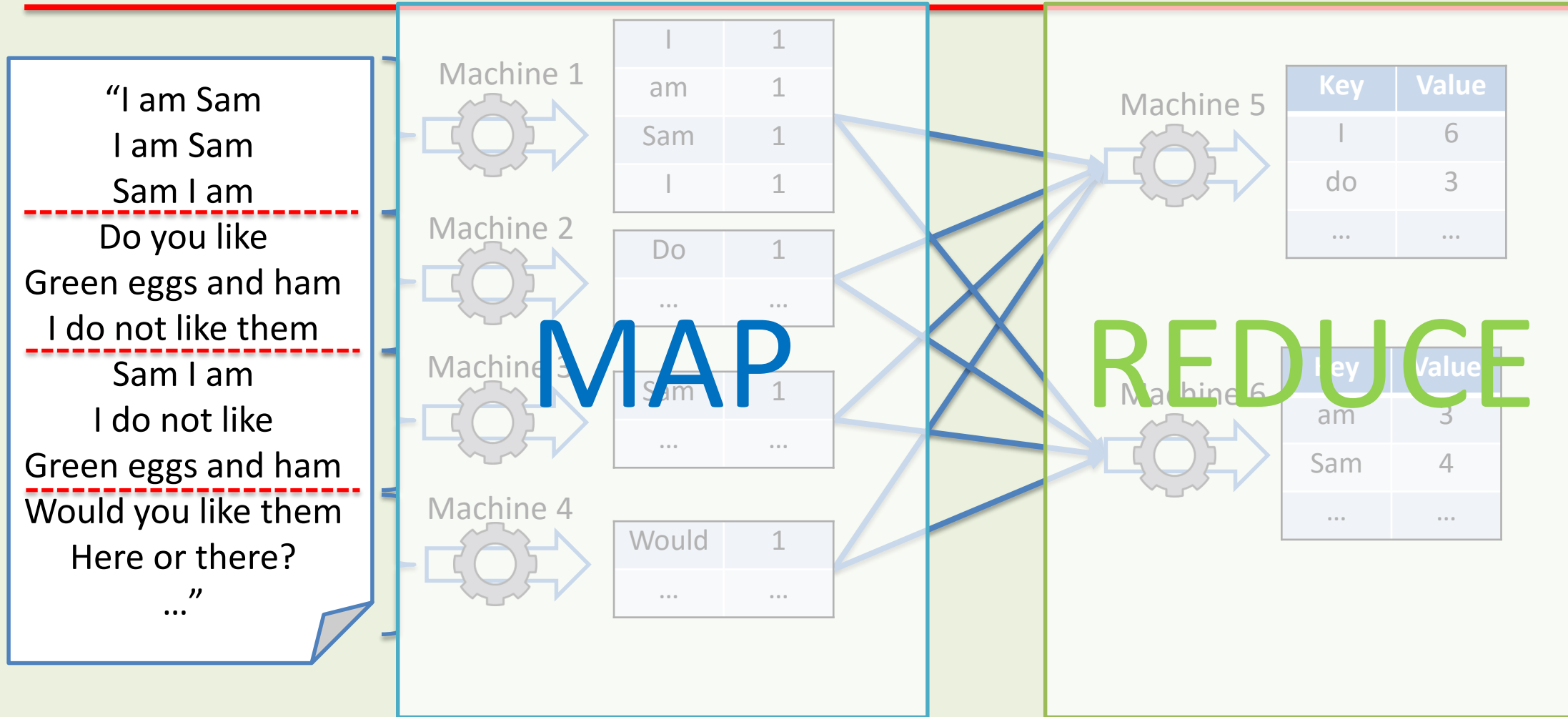
# What if the Document is Really Big?



Instead of building a Hash table in each machine  
why don't we just create <key, value> pairs, with key can be duplicated  
We will group them by key later any way, so we don't have to do that multiple times



# What if the Document is Really Big?



# Why MapReduce for cluster programming model?

---

When an algorithm/program is expressed in terms of a Map and Reduce operations, it's **simple** to build an execution engine that:

- Optimize the **computing schedule**
- **Handle failures**
- Optimize **inter-machine communication**

It's **powerful enough** to express most of data manipulation jobs

- **All dplyr verbs** can be easily expressed in term of Map Reduce operations
- **Most of SQL statements** can be translated into set of Map Reduce operations
- **Matrix Multiplication** can be expressed naturally using Map Reduce

# Coordination: Master

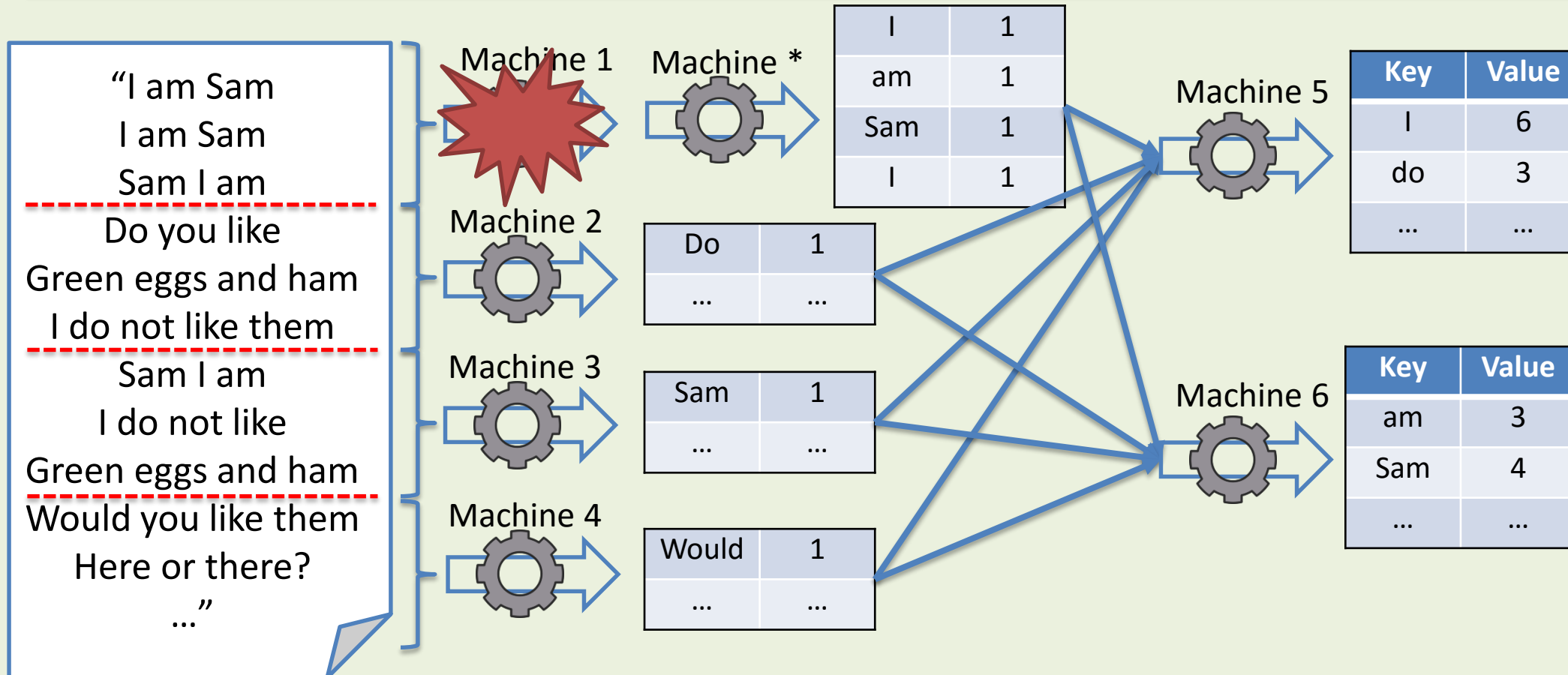
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Master node takes care of coordination:

- **Task status:** (idle, in-process, completed)
- **Idle tasks** get scheduled as workers become available
- When a map task completes, it sends the master the location and sizes of its intermediate files, one for each reducer
- Master pushes this info to reducers

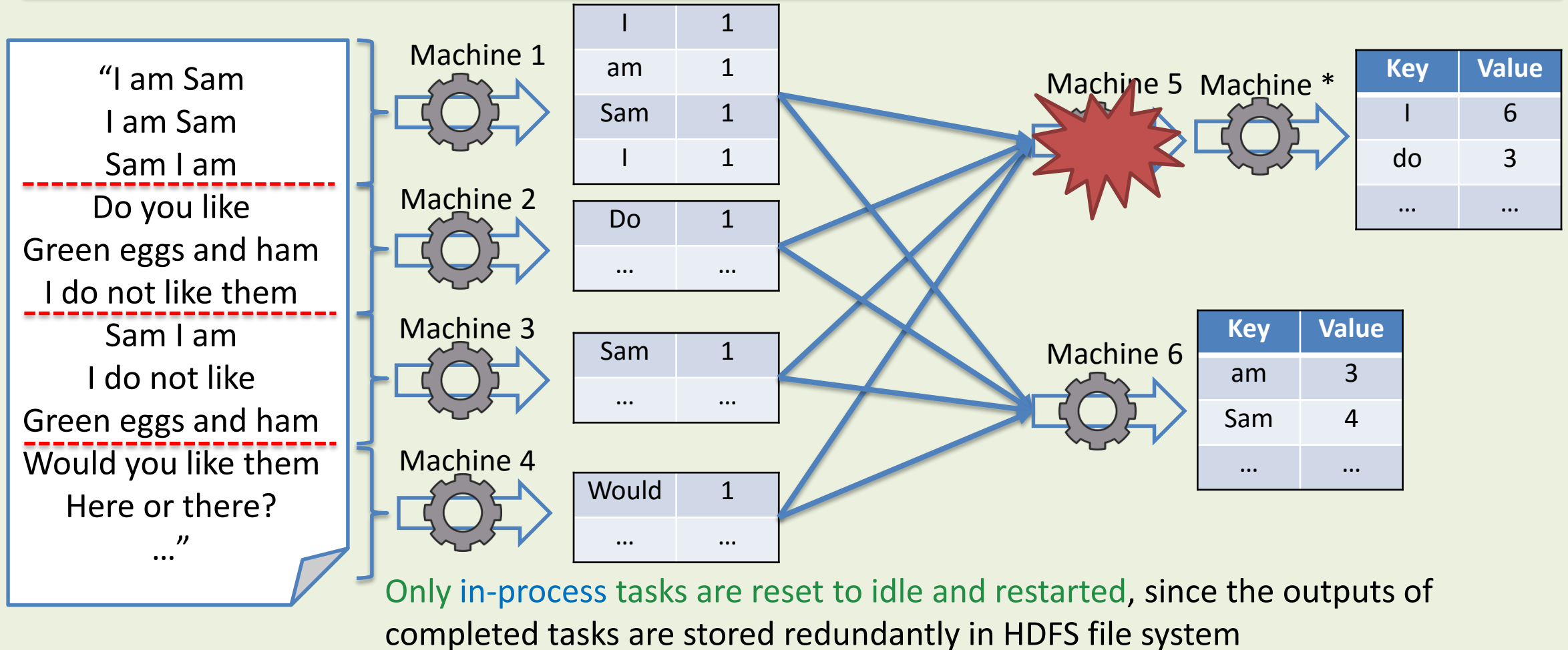
Master pings workers periodically to detect failures

# Dealing with Map Worker Failures



Map tasks completed or in-process at the fail worker are reset to **idle**  
These tasks are rescheduled on another worker

# Dealing with Reduce Worker Failures





# Dealing with Master failure

---

- MapReduce task is aborted and client is notified.
- There is only one master node, so the chance that it fails is much smaller than the chance that a worker fail

# Problems suitable / NOT suitable for MapReduce

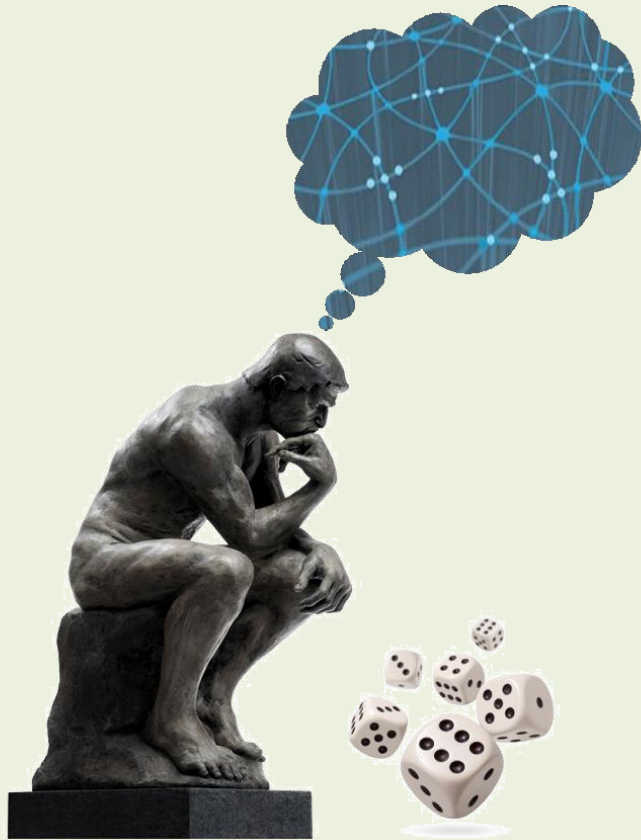
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MapReduce is great for:

- Problems that require sequential data access
- Large batch jobs (not interactive, real-time)

MapReduce is inefficient for problems where random (or irregular) access to data required:

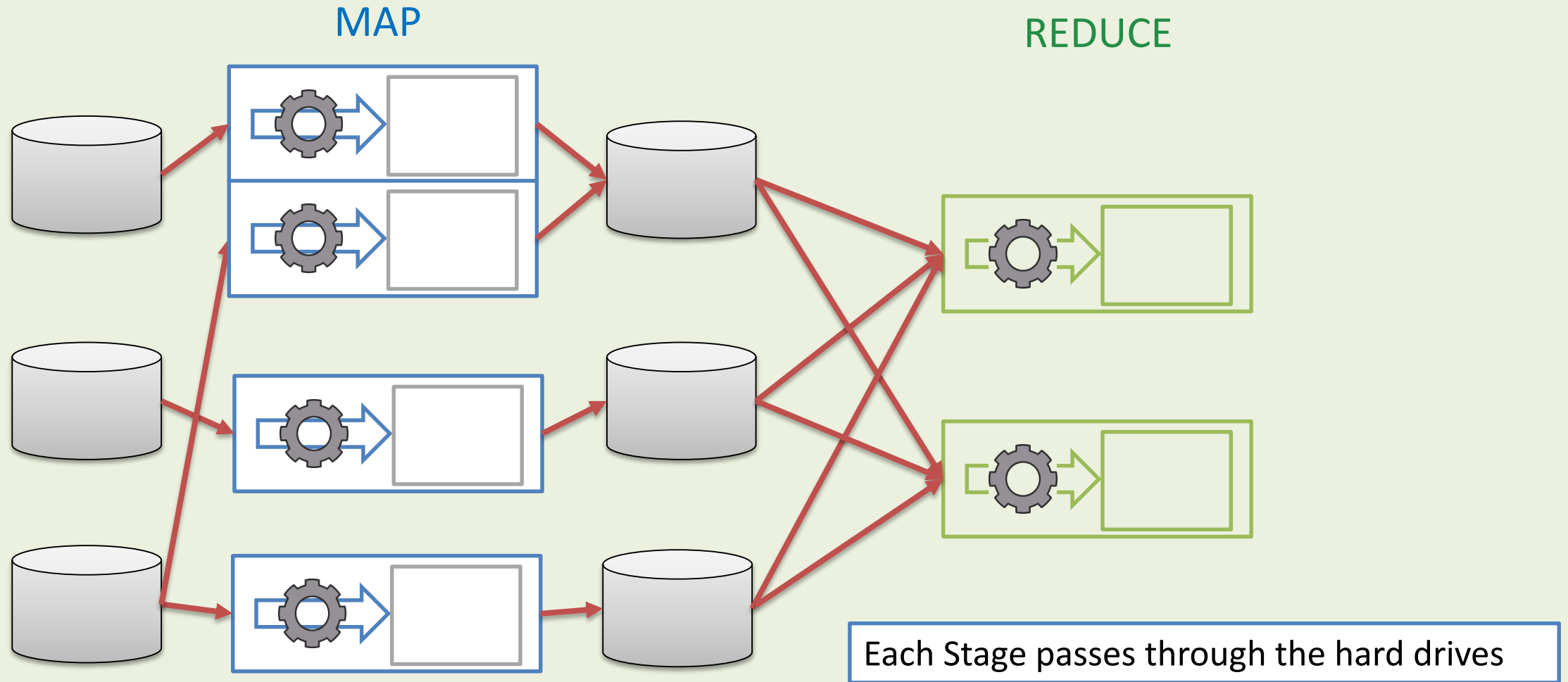
- Graphs
- Iterative algorithms (machine learning)



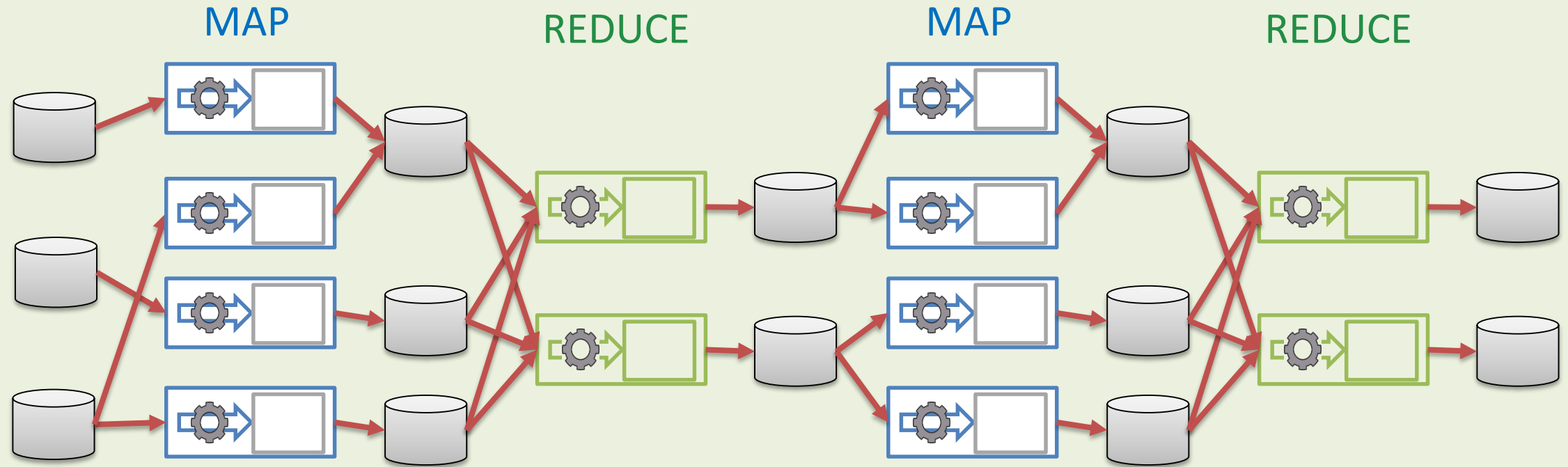
# The Introduction to Apache Spark

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*eSmart Systems*

# MapReduce Problem



# MapReduce Problem: Iterative Jobs



Iterative jobs involves a lot of disk I/O for each repetition

**Disk I/O is very slow!**



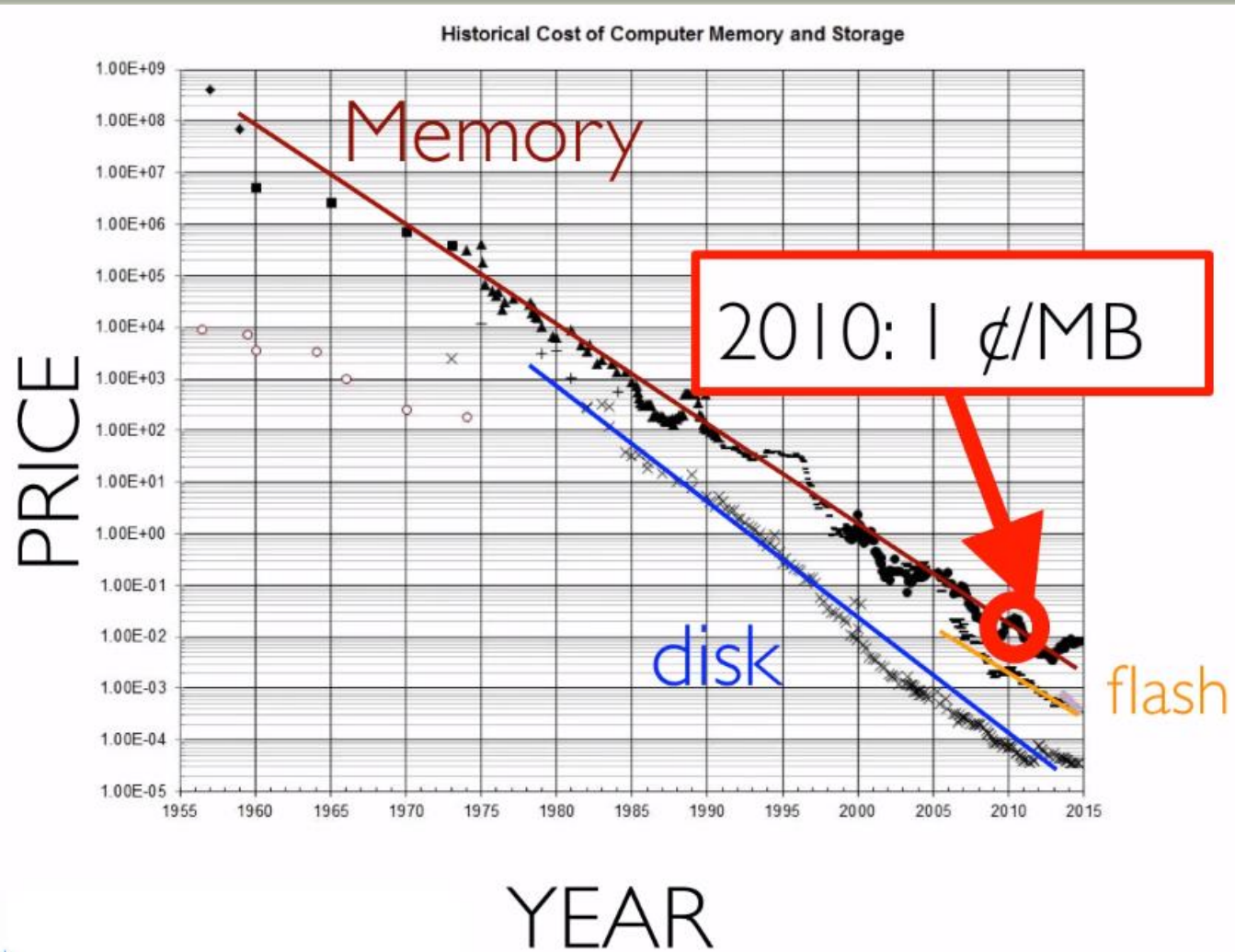
# Spark Motivation

---

- Keep more data in-memory
- Support better programming interfaces
- Create new distributed execution engine:



# Tech Trend: Cost of Memory



Lower cost means you can put more memory in each server

In-memory database solution

Data centers that use main memory exclusively

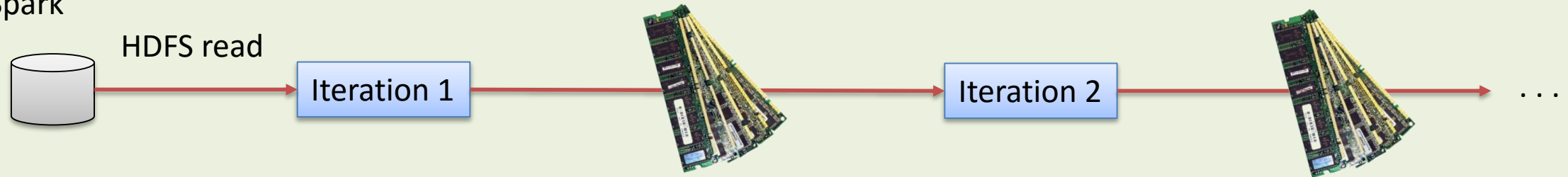
# Use Memory Instead of Disk

---

MapReduce – iterative job



Spark

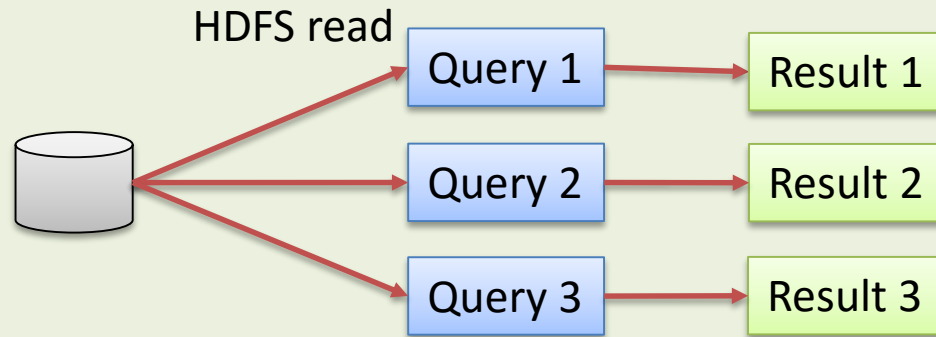


Memory: 10-100 times faster than disk or network

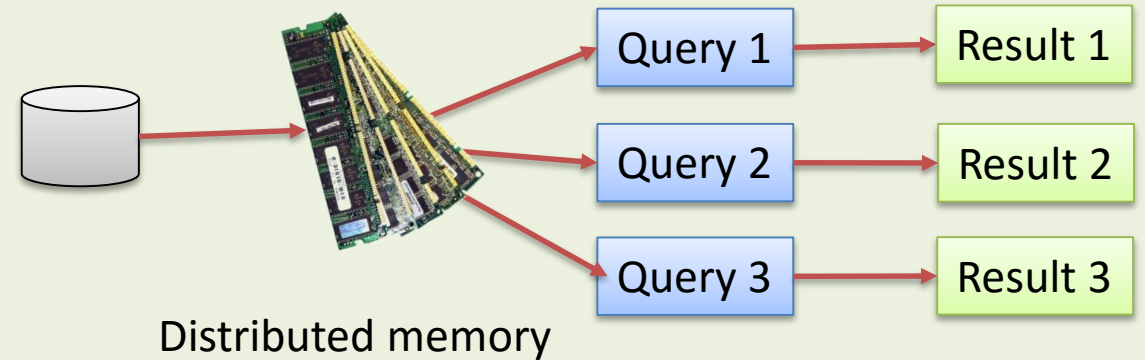
# Use Memory Instead of Disk

---

MapReduce – interactive query



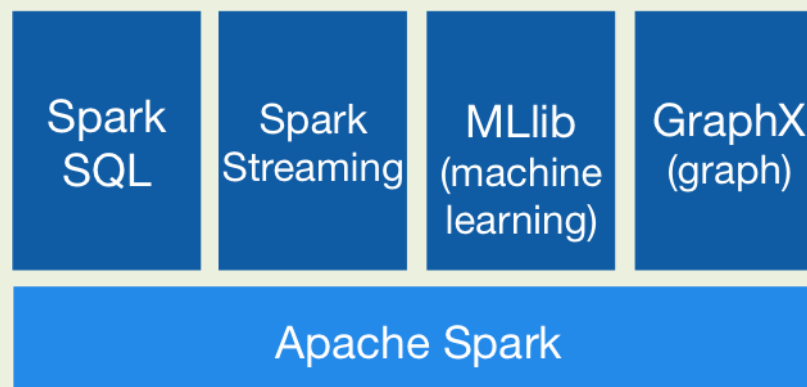
Spark



Memory: 10-100 times faster than disk or network

# Spark Tools

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**Spark SQL:** Let you query structured data using either SQL or recently DataFrame API (dplyr-like)  
Can connect to many different data sources: CSV, JSON, JDBC, ODBC

**Spark Streaming:** Stream processing

**MLlib:** Implementation of many scalable machine learning algorithms

**GraphX:** Support variety of graph processing and graph algorithms

# Spark and MapReduce Differences

---

	Hadoop Map Reduce	Spark
Storage	Disk only	In-memory or on disk
Operations	Map and Reduce	Map, Reduce, Join, Sample, ...
Execution model	Batch	Batch, interactive, real-time,...
Programming environments	Java	Scala, Java, Python, R
Programming Interface	Map and Reduce functions	DataFrame Interface SQL-like interface

# In-Memory Can Make a Big Difference!

---

Can achieve 100x speed up!





# In-Memory Can Make a Big Difference!

---

## First Public Cloud Petabyte Sort

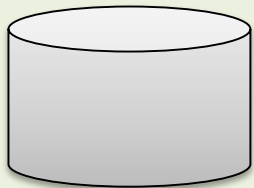
	<b>Hadoop MR Record</b>	<b>Spark Record</b>	<b>Spark 1 PB</b>
Data Size	102.5 TB	100 TB	1000 TB
Elapsed Time	72 mins	23 mins	234 mins
# Nodes	2100	206	190
# Cores	50400 physical	6592 virtualized	6080 virtualized
Cluster disk throughput	3150 GB/s (est.)	618 GB/s	570 GB/s
Sort Benchmark Daytona Rules	Yes	Yes	No
Network	dedicated data center, 10Gbps	virtualized (EC2) 10Gbps network	virtualized (EC2) 10Gbps network
<b>Sort rate</b>	<b>1.42 TB/min</b>	<b>4.27 TB/min</b>	<b>4.27 TB/min</b>
<b>Sort rate/node</b>	<b>0.67 GB/min</b>	<b>20.7 GB/min</b>	<b>22.5 GB/min</b>

[Daytona Gray 100 TB](#)  
sort benchmark record  
(tied for 1<sup>st</sup> place)

<http://databricks.com/blog/2014/11/05/spark-officially-sets-a-new-record-in-large-scale-sorting.html>

# Constructing and Working with DataFrames

From local file, HDFS,  
S3, SQL database...



Distributed memory  
(or local disk if too big)



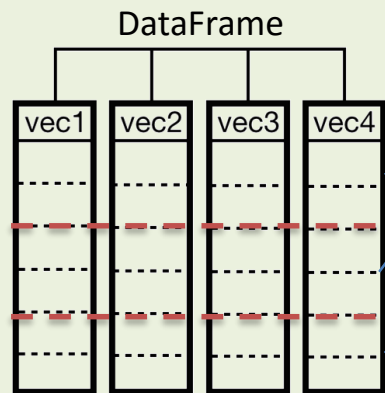
Apply Transformations (map step)  
(mutate, select, filter, arrange,...)



Apply Actions (reduce step)  
(summarize, group\_by, arrange)

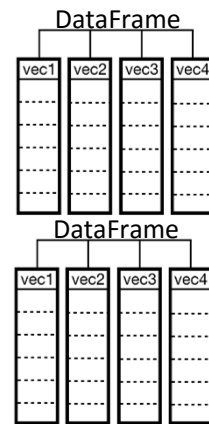


Collect Actions  
(collect, show,...)

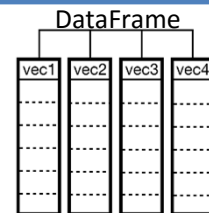


More partitions = more parallelism  
Spark automatically select a  
suitable number of partitions

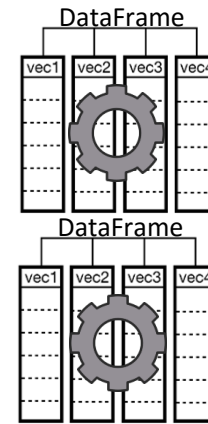
Worker 1



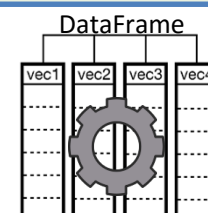
Worker 2



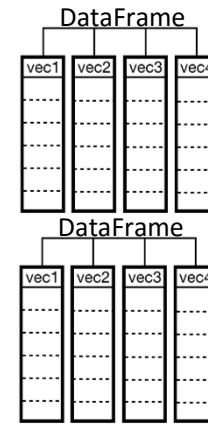
Worker 1



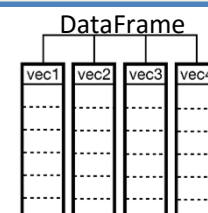
Worker 2



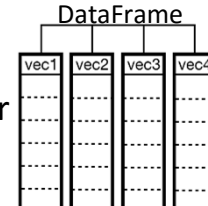
Worker 3



Worker 4



Driver



Driver must have  
enough memory to  
store the result

Lazy evaluation

With DataFrame, this step  
is also Lazy evaluated

