

Assignment 2

- ① Suppose two functions $f(x)$ and $g(x)$ are defined on some subset of the real numbers:

$$f(x) = O(g(x))$$

if and only if there exists a positive real c and a positive integer n_0 , such that for all $n \geq n_0$, $f(n) \leq c \cdot g(n)$

Important Item:

- Change of base rule:

$$\log_b(a) = \frac{\log_x(a)}{\log_x(b)}$$

proof:

$$\log_a x = y$$

$$a^y = x$$

$$\log_b a^y = \log_b x \quad \leftarrow \text{divid both sides by } \log_b$$

$$y \log_b a = \log_b x \quad \leftarrow \text{move the } y \text{ out front}$$

$$\frac{y \log_b a}{\log_b a} = \frac{\log_b x}{\log_b a} \quad \leftarrow \text{divide both sides by } \log_b a$$

$$y = \frac{\log_b x}{\log_b a}$$

continued \rightarrow

Assignment 1

continued

①

We say two sets are equal if they have the same elements.

$$f(x) = \{\log_2(10)\}, \quad g(x) = \{\log_{10}(10)\}$$
$$\rightarrow f(x) = \{3.32\}, \quad g(x) = \left\{ \frac{\log_2(10)}{\log_2(2)} \right\}$$

Therefore:

$$f(x)\{3.32\} = g(x)\{3.32\}$$

and the two sets are equal. So, $f(x) = O(g(x))$ and shows that the base of the logarithm doesn't matter.

$$(2) \quad T(n) = \begin{cases} 1, & n=1 \\ 2T(n/3) + n, & n > 1 \end{cases}$$

$$\begin{aligned} T(n) &= T(n/3) + T(n/3) + n \\ &= 2T(n/3) + n \\ &= 2(2T(n/3^2) + n/3) + n \\ &= 2^2 T(n/3^2) + n + n \\ &= 2^2 (2T(n/3^3) + n/3^2) + 2n \\ &= 2^3 T(n/3^3) + 3n \\ &\vdots \end{aligned}$$

let :

$$n/3^k = 1, \quad n = 3^k, \quad k = \log n$$

now :

$$\begin{aligned} T(n) &= 1n + n \log n \\ &= O(n \log n) \end{aligned}$$

③

1	7	1	5	3	-1	9
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↑ - pointer

1	7	1	5	3	-1	9
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↑

1	1	7	5	3	-1	9
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↑

1	1	5	7	3	-1	9
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↑

1	1	3	5	7	-1	9
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↑

-1	1	1	3	5	7	9
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↑

-1	1	1	3	5	7	9
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Assignment 1

- ⑤ The median-of-three takes the middle element of the sorted left-most, right-most and middle of the array.

Example:

	$\frac{1}{3}$	$\frac{2}{3}$
$<P$	Good Pivot	$>P$

We would at most pick two partitions of size $\frac{2}{3}$. By picking the median of three elements we are reducing the probability of picking a bad pivot. By sorting the three elements we increase our chance of picking a good pivot (increases the probability).

Sources:

- Khanacademy - change of base rule
- Khanacademy - change of base rule proof
- Data Structures & Algorithm Analysis in C++
4th edition, Mark Allen Weiss