

## Overview

Given: set function  $F: 2^V \rightarrow \mathbb{R}$ . Can make (noisy) value queries

Goal: "learn"  $F$  with as few queries as possible

Assumptions: 1) **sparsity**: at most  $k$  non-zero Fourier coefficients  
2) **low frequency**: degree of Fourier polynomials at most  $d$

## Motivation

Can represent a set function  $F$  where  $t$  is the indicator variable of the subsets

$$x_t : \{0, 1\}^n \rightarrow \mathbb{R} \quad x_t = \sum_{f \in \{0, 1\}^n} \hat{x}_f (-1)^{\langle f, t \rangle}$$

### 1. Graph cuts

$$\begin{aligned} f_e &= \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 \end{bmatrix} \\ t_A &= \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 \end{bmatrix} \\ x_t &= \frac{1}{2} (1 - \langle f_e, t_A \rangle) \end{aligned}$$

Fourier basis functions  
Frequency degree  $d = 2$   
One frequency for each edge

$A = \{2, 3, 4\} \quad V \setminus A = \{1, 5, 6\} \quad F(A) = 3$

Graph sketching

Size of the cut between  $A$  and  $V \setminus A$

### 2. Black box optimization

$$\begin{aligned} F(x_1, x_2, x_3, x_4, x_5) &= \\ F_1(x_1) + F_2(x_3, x_4) + F_3(x_2, x_5) & \quad d = 2 \end{aligned}$$

$F(x_1, x_2, x_3, x_4, x_5)$  = Validation error using hyperparameters  $x$

### 3. Decision Trees

$$\begin{aligned} d = 3 \quad & \text{Intuitively a decision tree of depth } d \text{ captures order } d \text{ interactions} \\ & \quad \begin{array}{c} x_1 \\ | \\ x_2 \quad 0 \\ | \quad | \\ x_3 \quad 1 \\ | \quad | \\ x_4 \quad 0 \\ | \quad | \\ x_5 \quad 1 \\ | \quad | \\ x_6 \quad 0.2 \\ | \quad | \\ x_7 \quad -2 \\ | \quad | \\ x_8 \quad 1.005 \\ | \quad | \\ x_9 \quad -3.0 \end{array} \end{aligned}$$

## Techniques

### 1. Hashing

$$x_t : \{0, 1\}^n \rightarrow \mathbb{R} \quad \hat{x}_f : \{0, 1\}^n \rightarrow \mathbb{R}$$

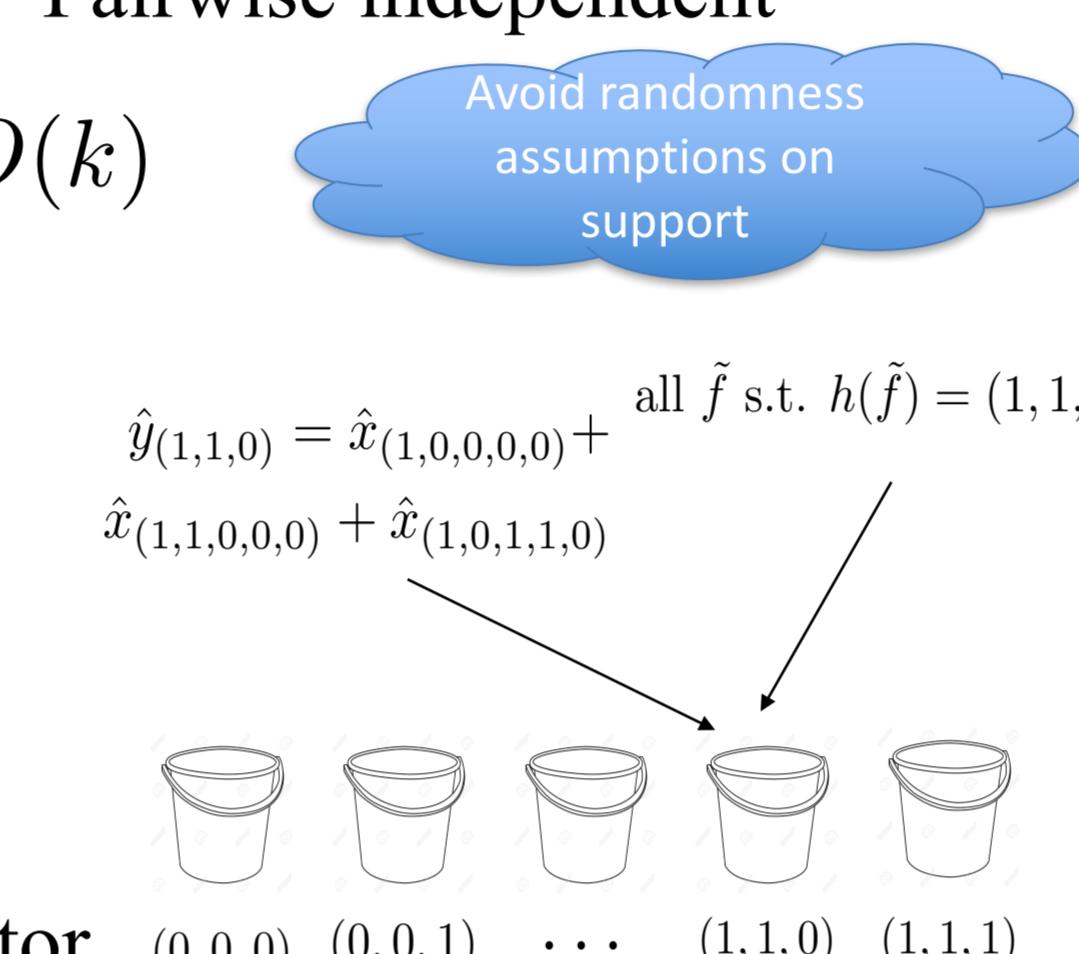
$\sigma \in \{0, 1\}^{n \times b}$  each element is chosen independent uniformly at random

Hash function  $h(f) \triangleq \sigma^\top f$  Pairwise independent

$$B = 2^b \quad \# \text{ of Buckets} \quad B \approx O(k)$$

$$y_t = x_{\sigma t} \quad \hat{y}_f = \sum_{\tilde{f}: h(\tilde{f})=f} \hat{x}_{\tilde{f}}$$

$$h(\tilde{f}) = \sigma^\top \tilde{f} = \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \tilde{f}$$



### 2. Frequency detection

$$\text{Let } a \in \{0, 1\}^n \text{ be a shift (measurement) vector}$$

$$y_t = x_{\sigma t+a} \quad \hat{y}_f = \sum_{\tilde{f}: h(\tilde{f})=f} \hat{x}_{\tilde{f}} (-1)^{\langle a, \tilde{f} \rangle}$$

Fix  $\sigma$  and repeat for carefully chosen values of  $a$  such that every frequency  $f$  is recoverable from linear measurements  $\langle a_i, f \rangle$

Solutions: 1- Set  $a_i = e_i, i \in [n]$  #measurements =  $O(n)$

2- If frequency is one-sparse do binary search

$$\begin{array}{c} f = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 \end{bmatrix} \\ \quad \quad \quad \# \text{measurements} = \\ \begin{bmatrix} 0 & & & 1 & & & \end{bmatrix} \\ \quad \quad \quad O(\log n) \\ \begin{bmatrix} 0 & 1 & 0 & 0 & 1 & & \end{bmatrix} \\ \begin{bmatrix} 0 & 1 & 0 & 1 & 0 & 0 & 1 \end{bmatrix} \end{array}$$

3- If frequency is  $d$ -sparse, hash and do binary search

$$\# \text{measurements} = O(d \log n)$$

Compressive sensing over finite fields for better sampling complexities

### 3. Peeling

We allow for collisions which result in erroneous frequency recovery for colliding frequencies

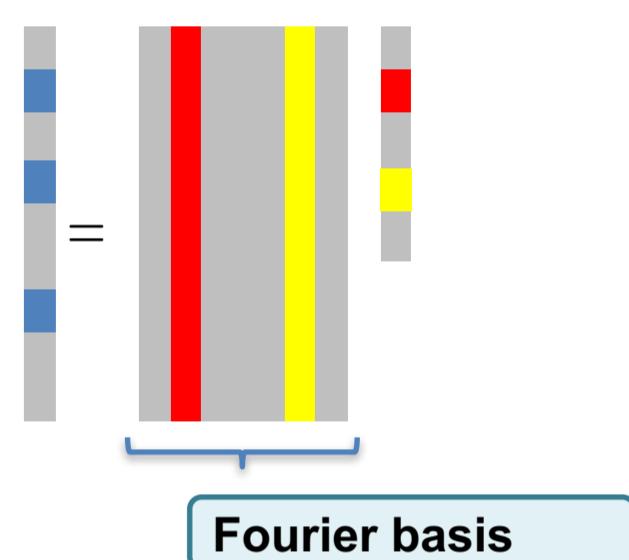
As long as there are not too many collisions sparsity increases after subtraction

$$\begin{array}{cccccc} \hat{x}_{f_1} & \hat{x}_{f_2} & \hat{x}_{f_3} & \hat{x}_{f_4} & \hat{x}_{f_5} & \hat{x}_{f_6} \\ - & \hat{x}_{f_1} & \hat{x}_{f_2} & \hat{x}_{f_{3,4}} & \hat{x}_{f_5} & \hat{x}_{f_6} \\ \hline & & & \hat{x}_{f_3} & \hat{x}_{f_4} & \hat{x}_{f_{3,4}} \end{array}$$

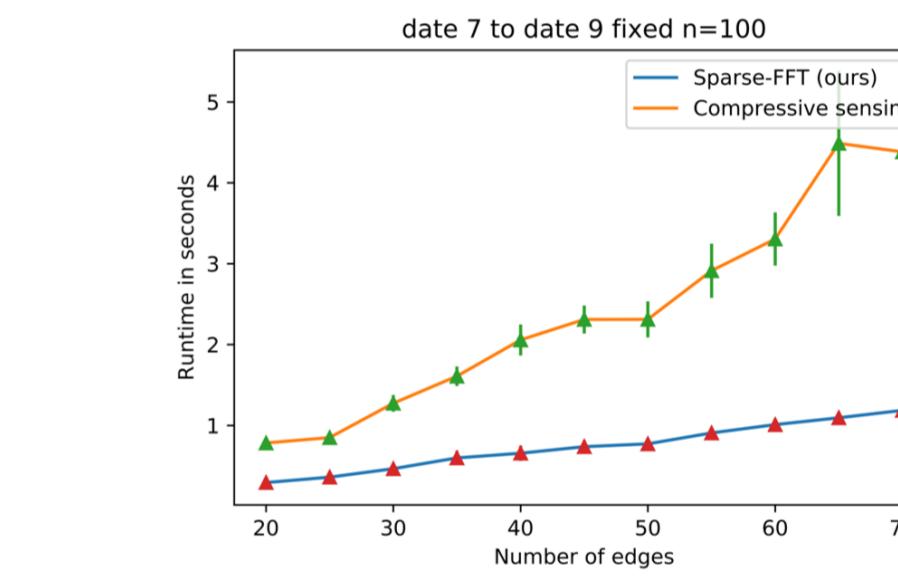
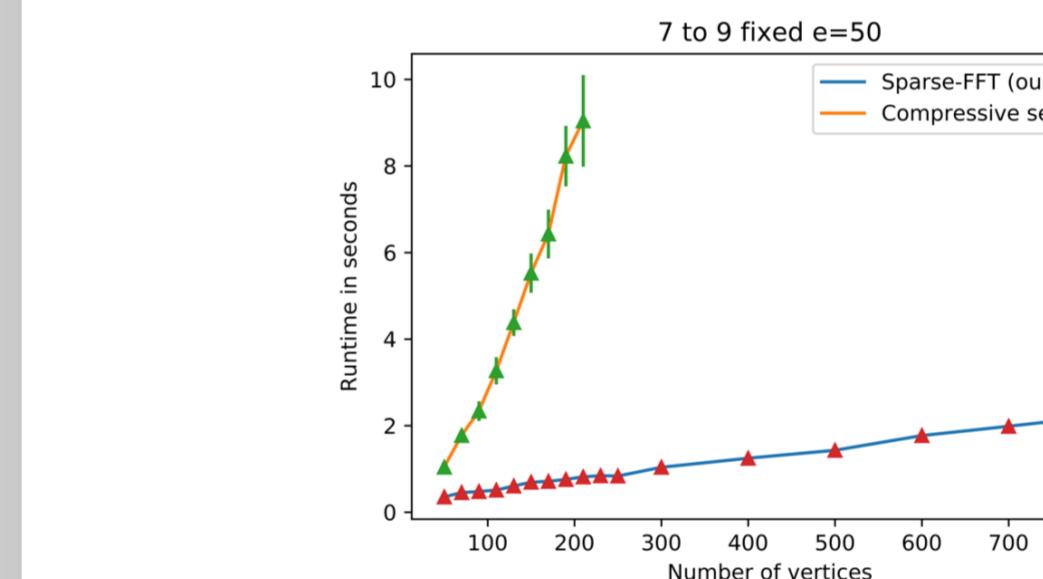
## Results

	Compressive sensing	Sparse FFT	Ours
Runtime	$\tilde{O}(kn^d)$ [2][3]	$\tilde{O}(kn^{(2)})$ [1][4]	$\tilde{O}(kn)$
Sampling complexity	$O(k(\log k)^2 \log(n^d))$ [2][3]	$O(kn)$ [1][4]	$O(k \log(n^d))$
Assumptions	None	Randomness of support	None
Robustness	Worst case noise	Gaussian noise + binary valued support	Worst case noise

- Obtain best of both worlds algorithm where runtime scales linear in  $d$  (sparse FFT) while maintaining optimal sampling complexities (compressive sensing)
- Remove assumptions on randomness of support
- Handle worst case noise



## Experiments



No. of vertices	CS method		Our method	
	Runtime	Samples	Runtime	Samples
40	1.88	812	0.92	6400
90	3.00	850	0.82	6401
140	3.88	880	1.22	7549
150	4.34	905	1.16	7942
170	6.13	927	1.22	7942
190	7.36	947	1.18	7271
210	8.24	965	1.18	7271
230	*	*	1.38	7942
250	*	*	1.38	7942
300	*	*	2.06	8031
400	*	*	2.06	8794
500	*	*	2.42	8794
600	*	*	3.35	9646
700	*	*	3.35	9646
800	*	*	3.60	9646

- Track evolution of a graph through time by looking at the values of a few cuts
- Runtimes scale linear in number of nodes and edges

## References

- R. Scheibler et al., "A Fast Hadamard Transform for Signals With Sublinear Sparsity in the Transform Domain," in *IEEE Transactions on Information Theory*, 2015.
- S. Peter, and A. Krause. "Learning Fourier sparse set functions." *Artificial Intelligence and Statistics*, 2012.
- I. Haviv, and O. Regev. "The restricted isometry property of subsampled Fourier matrices". *Geometric Aspects of Functional Analysis*, 2017
- X. Li, J. K. Bradley, S. Pawar and K. Ramchandran, "The SPRIGHT algorithm for robust sparse Hadamard Transforms," *IEEE International Symposium on Information Theory*, 2014.
- Hazan, Elad, Adam Klivans, and Yang Yuan. "Hyperparameter Optimization: A Spectral Approach." *ICLR* (2018).