

$e ::=$	$x \mid (e\ e) \mid \lambda x^\tau. e \mid (\mathbf{if}\ e\ e\ e) \mid c \mid \#t \mid \#f \mid (cons\ e\ e) \mid (vec\ \vec{e}) \mid n$	Expressions
$c ::=$	$add1 \mid =_n \mid \leq_n \mid num? \mid bool? \mid proc? \mid cons? \mid vec? \mid car \mid cdr \mid len \mid ref_n$	Primitive Operations
$pe ::=$	$\mathbf{car} \mid \mathbf{cdr} \mid \mathbf{ref}_n \mid \mathbf{len}$	Path Elements
$\pi ::=$	$\vec{p\bar{e}}$	Paths
$o ::=$	$\emptyset \mid \pi(x)$	Objects
$\phi ::=$	$a_1 o_1 + \dots + a_{(i-1)} \leq a_i o_i + \dots + a_j$	Linear Inequalities
$\Phi ::=$	$\vec{\phi}$	System of Linear Inequalities
$\sigma, \tau ::=$	$\top \mid \mathbf{N} \mid \{\diamond : \mathbf{N} \mid \Phi\} \mid \mathbf{T} \mid \mathbf{F} \mid (\bigcup \vec{\tau}) \mid \langle \tau, \tau \rangle \mid [[\vec{\tau}]] \mid x : \sigma \xrightarrow{o} \tau$	Types
$\psi ::=$	$\tau_{\pi(x)} \mid \bar{\tau}_{\pi(x)} \mid \psi \supset \psi \mid \psi \wedge \psi \mid \psi \vee \psi \mid \Phi \mid \mathbf{T} \mid \mathbf{F} \mid \mathbb{U}$	Propositions
$\Gamma ::=$	$\vec{\psi}$	Environments

Figure 1: Syntax of Types, Propositions, Terms, etc...

T-NUM $\Gamma \vdash n : \{\diamond : \mathbf{N} \mid \diamond = n\} ; \mathbb{T} ; \emptyset$	T-CONST $\Gamma \vdash c : \delta_\tau(c) ; \mathbb{T} ; \emptyset$	T-TRUE $\Gamma \vdash \#t : \mathbf{T} ; \mathbb{T} ; \emptyset$	T-FALSE $\Gamma \vdash \#f : \mathbf{F} ; \mathbb{F} ; \emptyset$
T-VAR $\frac{\Gamma \vdash \tau_x}{\Gamma \vdash x : \tau ; \bar{\mathbf{F}}_x ; x}$	T-ABS $\frac{\Gamma, \sigma_x \vdash e : \tau ; \psi ; o}{\Gamma \vdash \lambda x^\sigma. e : x : \sigma \xrightarrow{o} \tau ; \mathbb{T} ; \emptyset}$	T-APP $\frac{\Gamma \vdash e : x : \sigma \xrightarrow[\sigma_f]{\psi_f} \tau ; \psi ; o \quad \Gamma \vdash e' : \sigma ; \psi' ; o'}{\Gamma \vdash (e\ e') : \tau[o'/x] ; \psi_f[o'/x] ; \sigma_f[o'/x]}$	
T-IF $\frac{\Gamma \vdash e_1 : \tau_1 ; \psi_1 ; o_1 \quad \Gamma \vdash e_2 : \tau ; \psi_2 ; o \quad \Gamma \vdash e_3 : \tau ; \psi_3 ; o}{\Gamma \vdash (\mathbf{if}\ e_1\ e_2\ e_3) : \tau ; \psi_2 \vee \psi_3 ; o}$	T-LET $\frac{\Gamma \vdash e_0 : \tau ; \psi_0 ; o_0 \quad \Gamma, \tau_x, \bar{\mathbf{F}}_x \supset \psi_0, \mathbf{F}_x \supset \neg \psi_0 \vdash e_1 : \sigma ; \psi_1 ; o_1}{\Gamma \vdash (\mathbf{let}\ (x\ e_0)\ e_1) : \sigma[o_0/x] ; \psi_1[o_0/x] ; o_1[o_0/x]}$		
T-CONS $\frac{\Gamma \vdash e_1 : \tau_1 ; \psi_1 ; o_1 \quad \Gamma \vdash e_2 : \tau_2 ; \psi_2 ; o_2}{\Gamma \vdash (cons\ e_1\ e_2) : \langle \tau_1, \tau_2 \rangle ; \mathbb{T} ; \emptyset}$	T-CAR $\frac{\Gamma \vdash e : \langle \tau_1, \tau_2 \rangle ; \psi ; o}{\Gamma \vdash (car\ e) : \tau_1 ; \bar{\mathbf{F}}_{\mathbf{car}(x)}[o/x] ; \mathbf{car}(x)[o/x]}$	T-VEC $\frac{\Gamma \vdash e_1 : \tau_1 ; \psi_1 ; o_1 \quad \dots \quad \Gamma \vdash e_n : \tau_n ; \psi_n ; o_n}{\Gamma \vdash (vec\ e_1 \dots e_n) : [[\tau_1 \dots \tau_n]] ; \mathbb{T} ; \emptyset}$	
T-CDR $\frac{\Gamma \vdash e : \langle \tau_1, \tau_2 \rangle ; \psi ; o}{\Gamma \vdash (cdr\ e) : \tau_2 ; \bar{\mathbf{F}}_{\mathbf{cdr}(x)}[o/x] ; \mathbf{cdr}(x)[o/x]}$			
T-REF $\frac{\Gamma \vdash e : [[\dots \tau_i \dots]] ; \psi_i ; o_i}{\Gamma \vdash (ref_i\ e) : \tau_i ; \bar{\mathbf{F}}_{\mathbf{ref}_i(x)}[o/x] ; \mathbf{ref}_i(x)[o/x]}$		T-SUBSUME $\frac{\Gamma \vdash e : \tau ; \psi ; o \quad \Gamma, \psi \vdash \psi' \quad \tau <: \tau' \quad o <: o'}{\Gamma \vdash e : \tau' ; \psi' ; o'}$	

Figure 2: Typing Rules

$$\begin{array}{c}
\text{S-REFL} \\
\vdash \tau <: \tau
\end{array}
\quad
\begin{array}{c}
\text{S-TOP} \\
\vdash \tau <: \top
\end{array}
\quad
\begin{array}{c}
\text{S-UNIONSUPER} \\
\frac{\exists \sigma \in \vec{\sigma}, \vdash \tau <: \sigma}{\vdash \tau <: (\bigcup \vec{\sigma})}
\end{array}
\quad
\begin{array}{c}
\text{S-UNIONSUB} \\
\frac{\forall \tau \in \vec{\tau}, \vdash \tau <: \sigma}{\vdash (\bigcup \vec{\tau}) <: \sigma}
\end{array}
\quad
\begin{array}{c}
\text{S-PAIR} \\
\frac{\vdash \sigma_1 <: \tau_1 \quad \vdash \sigma_2 <: \tau_2}{\vdash \langle \sigma_1, \sigma_2 \rangle <: \langle \tau_1, \tau_2 \rangle}
\end{array}$$

$$\begin{array}{c}
\text{S-FUN} \\
\frac{\vdash \sigma' <: \sigma \quad \vdash \tau <: \tau' \quad \psi \vdash \psi' \quad \vdash o <: o'}{\vdash x:\sigma \xrightarrow[o]{\psi} \tau <: x:\sigma' \xrightarrow[o']{\psi'} \tau'}
\end{array}
\quad
\begin{array}{c}
\text{S-VEC} \\
\frac{\forall i, \vdash \sigma_i <: \tau_i}{\vdash [[\sigma_o \dots \sigma_i]] <: [[\tau_o \dots \tau_i]]}
\end{array}
\quad
\begin{array}{c}
\text{S-INEQ} \\
\frac{\text{unsatisfiable}(\Phi \bowtie \bar{\Phi}')}{\vdash \{\diamond : \mathbf{N} \mid \Phi\} <: \{\diamond : \mathbf{N} \mid \Phi'\}}
\end{array}$$

$$\begin{array}{c}
\text{S-INEQNUM} \\
\frac{\vdash \mathbf{N} <: \tau}{\vdash \{\diamond : \mathbf{N} \mid \Phi\} <: \tau}
\end{array}$$

Figure 3: Subtyping Rules

$$\begin{array}{c}
\text{L-ATOM} \\
\frac{\nu_{\pi(x)} \in \Gamma}{\Gamma \vdash \nu_{\pi(x)}}
\end{array}
\quad
\begin{array}{c}
\text{L-TRUE} \\
\Gamma \vdash \mathbb{T}
\end{array}
\quad
\begin{array}{c}
\text{L-FALSE} \\
\Gamma \vdash \mathbb{F}
\end{array}
\quad
\begin{array}{c}
\text{L-ANDI} \\
\frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \wedge \psi_2}
\end{array}
\quad
\begin{array}{c}
\text{L-ANDE} \\
\frac{\Gamma, \psi_1, \psi_2 \vdash \psi}{\Gamma, \psi_1 \wedge \psi_2 \vdash \psi}
\end{array}
\quad
\begin{array}{c}
\text{L-IMPI} \\
\frac{\Gamma, \psi_1 \vdash \psi_2}{\Gamma \vdash \psi_1 \supset \psi_2}
\end{array}$$

$$\begin{array}{c}
\text{L-IMPE} \\
\frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_1 \supset \psi_2}{\Gamma \vdash \psi_2}
\end{array}
\quad
\begin{array}{c}
\text{L-ORI} \\
\frac{\Gamma \vdash \psi_1 \text{ or } \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \vee \psi_2}
\end{array}
\quad
\begin{array}{c}
\text{L-ORE} \\
\frac{\Gamma, \psi_1 \vdash \psi \quad \Gamma, \psi_2 \vdash \psi}{\Gamma, \psi_1 \vee \psi_2 \vdash \psi}
\end{array}
\quad
\begin{array}{c}
\text{L-SUB} \\
\frac{\Gamma \vdash \tau_x \quad \vdash \tau <: \sigma}{\Gamma \vdash \sigma_x}
\end{array}
\quad
\begin{array}{c}
\text{L-SUBNOT} \\
\frac{\Gamma \vdash \bar{\sigma}_x \quad \vdash \tau <: \sigma}{\Gamma \vdash \bar{\tau}_x}
\end{array}$$

$$\begin{array}{c}
\text{L-BOT} \\
\frac{\Gamma \vdash \perp_x}{\Gamma \vdash \psi}
\end{array}
\quad
\begin{array}{c}
\text{L-UPDATE} \\
\frac{\Gamma \vdash \tau_x \quad \Gamma \vdash \nu_x}{\Gamma \vdash \text{update}(\tau, \nu)_x}
\end{array}
\quad
\begin{array}{c}
\text{L-INEQ} \\
\frac{\Gamma \vdash \mathbf{N}_x \quad \Gamma \vdash \Phi' \quad \text{unsatisfiable}(\Phi' \bowtie (\Phi[x/\diamond]))}{\Gamma \vdash \{\diamond : \mathbf{N} \mid \Phi\}_x}
\end{array}
\quad
\begin{array}{c}
\text{L-INEQSIMPLE} \\
\frac{\Gamma, \mathbf{N}_x, \Phi[x/\diamond] \vdash \psi}{\Gamma, \{\diamond : \mathbf{N} \mid \Phi\}_x \vdash \psi}
\end{array}$$

$$\begin{array}{c}
\text{L-INEQJOIN} \\
\frac{\Gamma, (\Phi_1 \bowtie \Phi_2) \vdash \psi}{\Gamma, \Phi_1, \Phi_2 \vdash \psi}
\end{array}
\quad
\begin{array}{c}
\text{L-INEQUNSAT} \\
\frac{\Gamma \vdash \Phi \quad \text{unsatisfiable}(\Phi)}{\Gamma \vdash \psi}
\end{array}$$

Figure 4: Logic Rules

$$\begin{array}{ll}
\neg \mathbb{T} & = \mathbb{F} \\
\neg \mathbb{F} & = \mathbb{T} \\
\neg \mathbb{U} & = \mathbb{U} \\
\neg \tau_{\pi(x)} & = \bar{\tau}_{\pi(x)} \\
\neg \bar{\tau}_{\pi(x)} & = \tau_{\pi(x)} \\
\neg (P \vee Q) & = \neg P \wedge \neg Q \\
\neg (P \wedge Q) & = \neg P \vee \neg Q \\
\neg (P \supset Q) & = \neg P \wedge Q
\end{array}$$

Figure 5: Proposition Negation

$\nu_{\pi(x)}[\pi'(y)/x]$	$= (\nu[\pi'(y)/x])_{\pi(\pi'(y))}$	
$\nu_{\pi(x)}[\emptyset/x]^+$	$= \mathbb{T}$	
$\nu_{\pi(x)}[\emptyset/x]^-$	$= \mathbb{F}$	
$\nu_{\pi(x)}[o/z]$	$= \nu_{\pi(x)}$	$x \neq z$ and $z \notin \text{fv}(\nu)$
$\nu_{\pi(x)}[o/z]^+$	$= \mathbb{T}$	$x \neq z$ and $z \in \text{fv}(\nu)$
$\nu_{\pi(x)}[o/z]^-$	$= \mathbb{F}$	$x \neq z$ and $z \in \text{fv}(\nu)$
$\mathbb{T}[o/x]$	$= \mathbb{T}$	
$\mathbb{F}[o/x]$	$= \mathbb{F}$	
$\mathbb{U}[o/x]$	$= \mathbb{U}$	
$(\psi_1 \supset \psi_2)[o/x]^+$	$= \psi_1[o/x]^- \supset \psi_2[o/x]^+$	
$(\psi_1 \supset \psi_2)[o/x]^-$	$= \psi_1[o/x]^+ \supset \psi_2[o/x]^-$	
$(\psi_1 \vee \psi_2)[o/x]$	$= \psi_1[o/x] \vee \psi_2[o/x]$	
$(\psi_1 \wedge \psi_2)[o/x]$	$= \psi_1[o/x] \wedge \psi_2[o/x]$	
$\Phi[o/x]$	$= \Phi$	$x \notin \Phi$
$\Phi[\pi'(y)/x]$	$= \forall \phi \in \Phi, \phi[x \mapsto \pi'(y)]$	$x \in \Phi$
$\Phi[\emptyset/x]^\pm$	$= \text{FMelim}(\Phi, x, \pm)$	$x \in \Phi$
$\pi(x)[\pi'(y)/x]$	$= \pi(\pi(y))$	
$\pi(x)[\emptyset/x]$	$= \emptyset$	
$\pi(x)[o/z]$	$= \pi(x)$	$x \neq z$
$\emptyset[o/x]$	$= \emptyset$	

Figure 6: Substitution