$$e ::= x \mid (e e) \mid \lambda x^{\tau}.e \mid (\text{if } e e e) \mid c \mid \#t \mid \#f \mid (cons \, e \, e) \mid (vec \, \overrightarrow{e}) \mid n \\ c ::= add1 \mid = \mid \leq \mid num? \mid bool? \mid proc? \mid cons? \mid vec? \mid car \mid cdr \mid len \mid ref[n] \\ pe ::= \mathbf{car} \mid \mathbf{cdr} \mid \mathbf{ref}_n \mid \mathbf{len} \\ \pi ::= \overrightarrow{pe} \\ o ::= \emptyset \mid \pi(x) \\ \phi ::= a_0o_0 + \dots + a_no_n \leq a_{n+1} \\ \Phi ::= \overrightarrow{\phi} \\ v, \tau ::= \top \mid \mathbf{N} \mid \{\diamond : \mathbf{N} \mid \Phi\} \mid \mathbf{T} \mid \mathbf{F} \mid (\bigcup \overrightarrow{\tau}) \mid \langle \tau, \tau \rangle \mid [[\overrightarrow{\tau}]] \mid x:\sigma \xrightarrow{\psi \mid \psi} \tau \\ \psi ::= \tau_{\pi}(x) \mid \overline{\tau}_{\pi}(x) \mid \psi \supset \psi \mid \psi \land \psi \mid \psi \lor \psi \mid \Phi \mid \mathbb{T} \mid \mathbb{F} \\ \Gamma ::= \overrightarrow{\psi} \\ \text{Environments} \\ \text{Environments}$$

Figure 1: Syntax of Types, Propositions, Terms, etc...

$$\begin{array}{lll} \text{T-Num} & \text{T-Const} \\ \Gamma \vdash n : \{ \diamond : \mathbf{N} \mid \diamond = n \} : \mathbb{T} \mid \mathbb{F} : \emptyset & \Gamma \vdash c : \delta_{\tau}(c) : \mathbb{T} \mid \mathbb{F} : \emptyset & \Gamma \vdash \#t : \mathbf{T} : \mathbb{T} \mid \mathbb{F} : \emptyset & \Gamma \vdash \#t : \mathbf{F} : \mathbb{F} \mid \mathbb{T} : \emptyset \\ \\ \frac{\Gamma \vdash \nabla X}{\Gamma \vdash X : \tau : \overline{\Gamma}_{X} \mid \mathbf{F}_{X} : x} & \frac{\Gamma \vdash \Delta BS}{\Gamma \vdash \Delta X^{\sigma} \cdot e : x : x} \frac{\Gamma \vdash \Delta BS}{\sigma} & \frac{\Gamma \vdash ABS}{\Gamma \vdash \Delta X^{\sigma} \cdot e : x : x} \frac{\nabla \vdash \Psi \mid \psi_{-} : \sigma}{\sigma} \\ \hline \Gamma \vdash E_{x} : \tau : \overline{\Gamma}_{X} \mid \overline{\Gamma}_{X} : \overline$$

Figure 2: Typing Rules

S-Refl S-Top
$$\vdash \tau <: \tau \qquad \qquad \frac{\text{S-UnionSuper}}{\exists \sigma \in \overrightarrow{\sigma}, \vdash \tau <: \sigma} \qquad \frac{\text{S-UnionSub}}{\vdash \tau <: (\bigcup \overrightarrow{\sigma})} \qquad \frac{\text{S-UnionSub}}{\vdash \tau \in \overrightarrow{\tau}, \vdash \tau <: \sigma} \qquad \qquad \frac{\vdash \sigma_1 <: \tau_1}{\vdash \sigma_2 <: \tau_2}$$

$$\vdash \sigma_2 <: \tau_2 \qquad \qquad \vdash \sigma_2 <: \tau_2 \qquad \qquad \vdash \sigma_2 <: \sigma_2 \sim: \sigma_2 \qquad \qquad \vdash \sigma_2 <: \sigma_2 \sim: \sigma_2$$

S-Pair

Figure 3: Subtyping Rules

Figure 4: Logic Rules

$$\begin{array}{lll} \psi_1 \mid \psi_2[o/x] & = \psi_1[o/x] \mid \psi_2[o/x] \\ \\ \nu_{\pi(x)}[\emptyset/x]^+ & = \mathbb{T} \\ \nu_{\pi(x)}[\emptyset/x]^- & = \mathbb{F} \\ \nu_{\pi(x)}[o/z] & = \nu_{\pi(x)} & x \neq z \text{ and } z \not\in \text{fv}(\nu) \\ \\ \nu_{\pi(x)}[o/z]^+ & = \mathbb{T} & x \neq z \text{ and } z \in \text{fv}(\nu) \\ \\ \nu_{\pi(x)}[o/z]^+ & = \mathbb{T} & x \neq z \text{ and } z \in \text{fv}(\nu) \\ \\ \nu_{\pi(x)}[o/z]^- & = \mathbb{F} & x \neq z \text{ and } z \in \text{fv}(\nu) \\ \\ \mathbb{T}[o/x] & = \mathbb{T} \\ \mathbb{F}[o/x] & = \mathbb{F} \\ (\psi_1 \supset \psi_2)[o/x]^+ & \psi_1[o/x]^- \supset \psi_2[o/x]^+ \\ (\psi_1 \supset \psi_2)[o/x]^- & \psi_1[o/x] \rightarrow \psi_2[o/x]^- \\ (\psi_1 \lor \psi_2)[o/x] & = \psi_1[o/x] \lor \psi_2[o/x] \\ (\psi_1 \lor \psi_2)[o/x] & = \psi_1[o/x] \lor \psi_2[o/x] \\ (\psi_1 \lor \psi_2)[o/x] & = \psi_1[o/x] \lor \psi_2[o/x] \\ \\ \Phi[o/x] & = \Phi & x \notin \Phi \\ \Phi[\pi'(y)/x] & = \forall \phi \in \Phi, \phi[x \mapsto \pi'(y)] & x \in \Phi \\ \Phi[\emptyset/x]^\pm & = \text{FME}(\Phi, x, \pm) & x \in \Phi \\ \\ \pi(x)[\pi'(y)/x] & = \emptyset \\ \\ \pi(x)[o/z] & = \emptyset & x \neq z \\ \\ \theta[o/x] & = \emptyset & x \neq z \\$$

Figure 5: Substitution