

| | | |
|--------------------|---|-------------------------------|
| $e ::=$ | $x \mid (e\ e) \mid \lambda x^\tau.e \mid (\mathbf{if}\ e\ e\ e) \mid c \mid \#t \mid \#f \mid (cons\ e\ e) \mid (vec\ \vec{e}) \mid n$ | Expressions |
| $c ::=$ | $+ \mid * \mid \leq \mid num? \mid bool? \mid proc? \mid cons? \mid vec? \mid car \mid cdr \mid len \mid ref$ | Primitive Operations |
| $pe ::=$ | $\mathbf{car} \mid \mathbf{cdr} \mid \mathbf{len}$ | Path Elements |
| $\pi ::=$ | \vec{pe} | Paths |
| $\mathcal{L} ::=$ | $n \mid \pi(x) \mid n\mathcal{L} \mid \mathcal{L} + \mathcal{L}$ | Linear Expressions |
| $o ::=$ | $\emptyset \mid \pi(x) \mid \mathcal{L}$ | Objects |
| $\phi ::=$ | $\mathcal{L} \leq \mathcal{L}$ | Linear Inequalities |
| $\Phi ::=$ | $\vec{\phi}$ | System of Linear Inequalities |
| $\sigma, \tau ::=$ | $\top \mid \mathbf{N} \mid \{x : \tau \mid \psi\} \mid \mathbf{T} \mid \mathbf{F} \mid (\bigcup \vec{\tau}) \mid \langle \tau, \tau \rangle \mid [[\tau]] \mid x:\sigma \xrightarrow{o} \tau$ | Types |
| $\psi ::=$ | $\tau_{\pi(x)} \mid \bar{\tau}_{\pi(x)} \mid \psi \supset \psi \mid \psi \wedge \psi \mid \psi \vee \psi \mid \Phi \mid \mathbf{T} \mid \mathbf{F} \mid \mathbb{U}$ | Propositions |
| $\Gamma ::=$ | $\vec{\psi}$ | Environments |

\perp is defined as (\bigcup) , \mathbf{B} is defined as $(\bigcup \mathbf{T} \mathbf{F})$.

Figure 1: Syntax of Types, Propositions, Terms, etc...

| | | | |
|--|--|--|--|
| $\begin{array}{c} \text{T-NUM} \\ \Gamma \vdash n : \{x : \mathbf{N} \mid x = n\} ; \mathbb{T} ; n \end{array}$ | $\begin{array}{c} \text{T-CONST} \\ \Gamma \vdash c : \delta_\tau(c) ; \mathbb{T} ; \emptyset \end{array}$ | $\begin{array}{c} \text{T-TRUE} \\ \Gamma \vdash \#t : \mathbf{T} ; \mathbb{T} ; \emptyset \end{array}$ | $\begin{array}{c} \text{T-FALSE} \\ \Gamma \vdash \#f : \mathbf{F} ; \mathbb{F} ; \emptyset \end{array}$ |
| $\begin{array}{c} \text{T-VAR} \\ \Gamma \vdash \tau_x \\ \hline \Gamma \vdash x : \tau ; \bar{\mathbf{F}}_x ; x \end{array}$ | $\begin{array}{c} \text{T-ABS} \\ \Gamma, \sigma_x \vdash e : \tau ; \psi ; o \\ \hline \Gamma \vdash \lambda x^\sigma. e : x : \sigma \xrightarrow{o} \tau ; \mathbb{T} ; \emptyset \end{array}$ | $\begin{array}{c} \text{T-APP} \\ \Gamma \vdash e : x : \sigma \xrightarrow[o_f]{\psi_f} \tau ; \psi ; o \\ \Gamma \vdash e' : \sigma ; \psi' ; o' \\ \hline \Gamma \vdash (e \ e') : \tau[o'/x] ; \psi_f[o'/x] ; o_f[o'/x] \end{array}$ | |
| $\begin{array}{c} \text{T-IF} \\ \Gamma \vdash e_1 : \tau_1 ; \psi_1 ; o_1 \\ \Gamma \vdash e_2 : \tau ; \psi ; o \\ \Gamma \vdash e_3 : \tau ; \psi ; o \\ \hline \Gamma \vdash (\mathbf{if} \ e_1 \ e_2 \ e_3) : \tau ; \psi ; o \end{array}$ | $\begin{array}{c} \text{T-LET} \\ \Gamma \vdash e_0 : \tau ; \psi_0 ; o_0 \\ \Gamma, \tau_x, \bar{\mathbf{F}}_x \supset \psi_0, \mathbf{F}_x \supset \neg \psi_0 \vdash e_1 : \sigma ; \psi_1 ; o_1 \\ \hline \Gamma \vdash (\mathbf{let} \ (x \ e_0) \ e_1) : \sigma[o_0/x] ; \psi_1[o_0/x] ; o_1[o_0/x] \end{array}$ | $\begin{array}{c} \text{T-CONS} \\ \Gamma \vdash e_1 : \tau_1 ; \psi_1 ; o_1 \\ \Gamma \vdash e_2 : \tau_2 ; \psi_2 ; o_2 \\ \hline \Gamma \vdash (\mathbf{cons} \ e_1 \ e_2) : \langle \tau_1, \tau_2 \rangle ; \mathbb{T} ; \emptyset \end{array}$ | |
| $\begin{array}{c} \text{T-CAR} \\ \Gamma \vdash e : \langle \tau_1, \tau_2 \rangle ; \psi ; o \\ \psi_r = \bar{\mathbf{F}}_{\mathbf{car}(x)}[o/x] \\ \hline \Gamma \vdash (\mathbf{car} \ e) : \tau_1 ; \psi_r ; \mathbf{car}(x)[o/x] \end{array}$ | $\begin{array}{c} \text{T-CDR} \\ \Gamma \vdash e : \langle \tau_1, \tau_2 \rangle ; \psi ; o \\ \psi_r = \bar{\mathbf{F}}_{\mathbf{cdr}(x)}[o/x] \\ \hline \Gamma \vdash (\mathbf{cdr} \ e) : \tau_2 ; \psi_r ; \mathbf{cdr}(x)[o/x] \end{array}$ | $\begin{array}{c} \text{T-VEC} \\ \Gamma \vdash e_1 : \tau ; \psi_1 ; o_1 \\ \dots \\ \Gamma \vdash e_n : \tau ; \psi_n ; o_n \\ \hline \Gamma \vdash (\mathbf{vec} \ e_1 \ \dots \ e_n) : [[\tau]] ; \mathbb{T} ; \emptyset \end{array}$ | |
| $\begin{array}{c} \text{T-REF} \\ \Gamma \vdash e_1 : [[\tau]] ; \psi_1 ; o_1 \\ \Gamma \vdash e_2 : \{x : \mathbf{N} \mid 0 \leq x \leq \mathbf{len}(o_1)\} ; \psi_2 ; o_2 \\ \hline \Gamma \vdash (\mathbf{ref} \ e_1 \ e_2) : \tau ; \mathbf{prop}(\tau) ; \emptyset \end{array}$ | $\begin{array}{c} \text{T-SUBSUME} \\ \Gamma \vdash e : \tau ; \psi ; o \\ \Gamma \vdash \psi <: \psi' \quad \vdash \tau <: \tau' \quad \vdash o <: o' \\ \hline \Gamma \vdash e : \tau' ; \psi' ; o' \end{array}$ | | |

Figure 2: Typing Rules

$$\begin{array}{c}
\text{SP-IMPL} \\
\frac{\Gamma, \psi \vdash \psi'}{\Gamma \vdash \psi <: \psi'} \\
\\
\text{SP-REFL} \\
\Gamma \vdash \psi <: \psi \\
\\
\text{SP-TOP} \\
\Gamma \vdash \psi <: \mathbb{U} \\
\\
\text{SO-REFL} \\
\vdash o <: o \\
\\
\text{SO-TOP} \\
\vdash o <: \emptyset \\
\\
\text{S-REFL} \\
\vdash \tau <: \tau \\
\\
\\
\text{S-TOP} \\
\vdash \tau <: \top \\
\\
\text{S-UNIONSUPER} \\
\frac{\exists \sigma \in \vec{\sigma}, \vdash \tau <: \sigma}{\vdash \tau <: (\bigcup \vec{\sigma})} \\
\\
\text{S-UNIONSUB} \\
\frac{\forall \tau \in \vec{\tau}, \vdash \tau <: \sigma}{\vdash (\bigcup \vec{\tau}) <: \sigma} \\
\\
\text{S-PAIR} \\
\frac{\vdash \sigma_1 <: \tau_1 \quad \vdash \sigma_2 <: \tau_2}{\vdash \langle \sigma_1, \sigma_2 \rangle <: \langle \tau_1, \tau_2 \rangle} \\
\\
\text{S-FUN} \\
\frac{\vdash \sigma' <: \sigma \quad \vdash \tau <: \tau' \quad \vdash \psi <: \psi' \quad \vdash o <: o'}{\vdash x:\sigma \xrightarrow{o} \tau <: x:\sigma' \xrightarrow{o'} \tau'} \\
\\
\text{S-VEC} \\
\frac{\vdash \sigma <: \tau}{\vdash [[\sigma]] <: [[\tau]]} \\
\\
\text{S-DEP} \\
\frac{\vdash \tau <: \sigma \quad \psi \vdash \psi'}{\vdash \{x:\tau \mid \psi\} <: \{x:\sigma \mid \psi'\}} \\
\\
\text{S-DEPWEAK} \\
\frac{\vdash \tau <: \sigma}{\vdash \{x:\tau \mid \psi\} <: \sigma}
\end{array}$$

Figure 3: Subtyping Rules

$$\begin{array}{c}
\text{L-ATOM} \\
\frac{\chi \in \Gamma}{\Gamma \vdash \chi} \\
\\
\text{L-TRUE} \\
\Gamma \vdash \mathbb{T} \\
\\
\text{L-FALSE} \\
\frac{}{\Gamma \vdash \mathbb{F}} \\
\\
\text{L-ANDI} \\
\frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \wedge \psi_2} \\
\\
\text{L-ANDE} \\
\frac{\Gamma, \psi_1, \psi_2 \vdash \psi}{\Gamma, \psi_1 \wedge \psi_2 \vdash \psi} \\
\\
\text{L-IMPI} \\
\frac{\Gamma, \psi_1 \vdash \psi_2}{\Gamma \vdash \psi_1 \supset \psi_2} \\
\\
\text{L-IMPE} \\
\frac{\Gamma \vdash \psi_1 \quad \Gamma \vdash \psi_1 \supset \psi_2}{\Gamma \vdash \psi_2} \\
\\
\text{L-ORI} \\
\frac{\Gamma \vdash \psi_1 \text{ or } \Gamma \vdash \psi_2}{\Gamma \vdash \psi_1 \vee \psi_2} \\
\\
\text{L-ORE} \\
\frac{\Gamma, \psi_1 \vdash \psi \quad \Gamma, \psi_2 \vdash \psi}{\Gamma, \psi_1 \vee \psi_2 \vdash \psi} \\
\\
\text{L-SUB} \\
\frac{\Gamma \vdash \tau_{\pi(x)} \quad \vdash \tau <: \sigma}{\Gamma \vdash \sigma_{\pi(x)}} \\
\\
\text{L-SUBNOT} \\
\frac{\Gamma \vdash \bar{\sigma}_{\pi(x)} \quad \vdash \tau <: \sigma}{\Gamma \vdash \bar{\tau}_{\pi(x)}} \\
\\
\text{L-BOT} \\
\frac{\Gamma \vdash \perp_{\pi(x)}}{\Gamma \vdash \psi} \\
\\
\text{L-UPDATE} \\
\frac{\Gamma \vdash \tau_{\pi'(x)} \quad \Gamma \vdash \nu_{\pi(\pi'(x))}}{\Gamma \vdash \text{update}(\tau, \nu, \pi)_{\pi'(x)}} \\
\\
\text{L-DEP} \\
\frac{\Gamma \vdash \tau_{\pi(x)} \quad \Gamma \vdash \psi}{\Gamma \vdash \{y:\tau \mid \psi\}_{\pi(x)}} \\
\\
\text{L-DEPWEAK} \\
\frac{\Gamma \vdash \{y:\tau \mid \psi\}_{\pi(x)}}{\Gamma \vdash \psi[\pi(x)/y]} \\
\\
\text{L-INEQIMPL} \\
\frac{\Gamma \vdash \Phi_1 \quad \text{unsatisfiable}(\Phi_1 \bowtie \bar{\Phi}_2)}{\Gamma \vdash \Phi_2} \\
\\
\text{L-INEQJOIN} \\
\frac{\Gamma, (\Phi_1 \bowtie \Phi_2) \vdash \psi}{\Gamma, \Phi_1, \Phi_2 \vdash \psi} \\
\\
\text{L-INEQUNSAT} \\
\frac{\Gamma \vdash \Phi \quad \text{unsatisfiable}(\Phi)}{\Gamma \vdash \psi}
\end{array}$$

Where χ ranges over propositions of the form $\tau_{\pi(x)}$, $\bar{\tau}_{\pi(x)}$, and Φ .

Figure 4: Logic Rules

$$\begin{array}{lll}
\psi_1 \mid \psi_2[o/x] & = & \psi_1[o/x] \mid \psi_2[o/x] \\
\\
\nu_{\pi(x)}[\pi'(y)/x] & = & (\nu[\pi'(y)/x])_{\pi(\pi'(y))} \\
\nu_{\pi(x)}[\emptyset/x]^+ & = & \mathbb{T} \\
\nu_{\pi(x)}[\emptyset/x]^- & = & \mathbb{F} \\
\nu_{\pi(x)}[o/z] & = & \nu_{\pi(x)} \quad x \neq z \text{ and } z \notin \text{fv}(\nu) \\
\nu_{\pi(x)}[o/z]^+ & = & \mathbb{T} \quad x \neq z \text{ and } z \in \text{fv}(\nu) \\
\nu_{\pi(x)}[o/z]^- & = & \mathbb{F} \quad x \neq z \text{ and } z \in \text{fv}(\nu) \\
\mathbb{T}[o/x] & = & \mathbb{T} \\
\mathbb{F}[o/x] & = & \mathbb{F} \\
(\psi_1 \supset \psi_2)[o/x]^+ & = & \psi_1[o/x]^- \supset \psi_2[o/x]^+ \\
(\psi_1 \supset \psi_2)[o/x]^- & = & \psi_1[o/x]^+ \supset \psi_2[o/x]^- \\
(\psi_1 \vee \psi_2)[o/x] & = & \psi_1[o/x] \vee \psi_2[o/x] \\
(\psi_1 \wedge \psi_2)[o/x] & = & \psi_1[o/x] \wedge \psi_2[o/x] \\
\Phi[o/x] & = & \Phi \quad x \notin \Phi \\
\Phi[\pi'(y)/x] & = & \forall \phi \in \Phi, \phi[x \mapsto \pi'(y)] \quad x \in \Phi \\
\Phi[\emptyset/x]^\pm & = & \text{FMelim}(\Phi, x, \pm) \quad x \in \Phi \\
\\
\pi(x)[\pi'(y)/x] & = & \pi[\pi'(y)/x](\pi'(y)) \\
\pi(x)[\emptyset/x] & = & \emptyset \\
\pi(x)[o/z] & = & \pi[o/z](x) \quad x \neq z \\
\emptyset[o/x] & = & \emptyset
\end{array}$$

Figure 5: Substitution

$$\begin{array}{ll}
\text{update}(\langle \tau, \sigma \rangle, \nu, \pi :: \mathbf{car}) & = \langle \text{update}(\tau, \nu, \pi), \sigma \rangle \\
\text{update}(\langle \tau, \sigma \rangle, \nu, \pi :: \mathbf{cdr}) & = \langle \tau, \text{update}(\sigma, \nu, \pi) \rangle
\end{array}$$