

Chapter 6. Lattices and Boolean algebras

Lattices as algebraic structures

Def 1. An ordered set (A, \leq) is called a lattice if for any $x, y \in A$ $\exists \sup\{x, y\}$, $\exists \inf\{x, y\}$

Def 2 The algebraic structure (A, \vee, \wedge) is called a lattice if the following axioms are satisfied:

(1) commutativity: $\forall x, y \in A$ $x \vee y = y \vee x$
 $x \wedge y = y \wedge x$

(2) associativity $\forall x, y, z \in A$ $(x \vee y) \vee z = x \vee (y \vee z)$
 $(x \wedge y) \wedge z = x \wedge (y \wedge z)$

(3) absorption: $\forall x, y \in A$ $x \vee (x \wedge y) = x$
 $x \wedge (x \vee y) = x$

We'll show that the two notions are equivalent.

Theorem 1 1). Let (A, \leq) be a lattice (as an ordered set).

We define on A the following operations:

$$\forall x, y \in A: \quad x \vee y \stackrel{\text{def}}{=} \sup\{x, y\}$$

$$x \wedge y \stackrel{\text{def}}{=} \inf\{x, y\}.$$

Then (A, \vee, \wedge) is a lattice (as an alg. structure).

2). Let (A, \vee, \wedge) be a lattice (as an alg. structure).

We define on A the relation ' \leq ': $\forall x, y \in A$ $x \leq y \stackrel{\text{def}}{=} x \vee y = y$.

Then, ' \leq ' is an order relation on A , and moreover, (A, \leq)

is a lattice, where $\sup\{x, y\} = x \vee y$ and $\inf\{x, y\} = x \wedge y$.

Proof 1). Homework! see Th. 6.1.2 / 31.

2) First, we prove that $\forall x, y \in A$.

$$xvy = y \iff x \wedge y = x \quad \left(\text{so we may also use the oper. } \wedge \text{ to define } \leq \right)$$

$$\stackrel{<=>}{=} xvy = y \stackrel{x \wedge}{\implies} x \wedge (xvy) = x \wedge y \stackrel{abs}{\implies} x = x \wedge y$$

$$\stackrel{<=>}{=} x \wedge y = x \stackrel{x \wedge y}{\implies} (x \wedge y)vy \implies xvy \stackrel{abs}{\implies} y = xvy$$

We check that \leq is a partial order relation:

$$(R1) \quad \forall a \in A : a \leq a \quad \stackrel{?}{\iff} a \vee a \stackrel{?}{=} a$$
$$\iff a \wedge a = a.$$

(so we need to prove that both operators are idempotent)

Indeed, by absorb law, we have:

$$a = a \vee (a \wedge a) \stackrel{abs}{\implies} a \wedge a = a \wedge \underbrace{(a \vee (a \wedge a))}_a$$
$$\stackrel{abs}{\implies} a \wedge a = a$$

$$\text{Similarly, } a = a \wedge (a \vee a) \stackrel{abs}{\implies} a \vee a = a \vee \underbrace{(a \wedge (a \vee a))}_a = a$$

$$(T) \quad \forall a, b, c \in A \quad \text{if } a \leq b, b \leq c \implies$$

$$\implies a \vee b = b, b \vee c = c \implies a \vee c = a \vee (b \vee c) \stackrel{ass}{=} \\ = \underbrace{(a \vee b)}_b \vee c = b \vee c = \underline{c}, \text{ hence } a \leq c$$

$$(A) \quad \text{let } a, b \in A \text{ st } a \leq b \text{ and } b \leq a.$$

$$\text{Then: } a \vee b = b \text{ and } b \vee a = a \stackrel{comm}{\implies} a = b$$

Next, we prove that $\sup\{x, y\} = xvy$.

we check that xvy is an upper bound for $\{x, y\}$.

$$xv(xvy) \stackrel{abs}{=} (x \vee a) \vee y \stackrel{idemp}{=} xvy. \implies x \leq xvy$$

$$f \vee (xy) = x \vee (fy) = xy \Rightarrow y \leq xy.$$

• let z be another upper bound for $\{x, y\}$

$$\underline{z \vee (xy)} = (z \vee x) \vee y = z \vee y = \underline{z}$$

$$\text{hence } xy \leq z$$

hence xy is the least upper bound.

similarly: $\inf \{x, y\} = xy$ (homework!)

Exercises. 1). let (A, \leq) be a totally ordered set. so

$$\inf \{x, y\} = \min \{x, y\}; \quad \sup \{x, y\} = \max \{x, y\}.$$

2). Let M be a set and consider $(\mathcal{P}(M), \subseteq)$

$$\text{let } X, Y \subseteq M. \quad \text{then: } \inf \{X, Y\} = X \cap Y$$

$$\sup \{X, Y\} = X \cup Y.$$

3). let $(\mathbb{N}, |)$. let $x, y \in \mathbb{N}$.

$$\inf \{x, y\} = (x, y) \quad \text{g.c.d.}$$

$$\sup \{x, y\} = [x, y] \quad \text{l.c.m.}$$

Def 3 A function $f: (A, \vee, \wedge) \rightarrow (A', \vee', \wedge')$ is called

a morphism of lattices if $\forall x, y \in A$:

$$f(x \vee y) = f(x) \vee' f(y) \quad \text{and} \quad f(x \wedge y) = f(x) \wedge' f(y).$$

Rem (see the exercises!)

f morph of lattices \Rightarrow f is increasing.

\Leftarrow

Boole algebras and Boole rings

Defn. A lattice (A, \vee, \wedge) is called a Boole lattice (Boole algebra)

if the following additional axioms hold:

(1) distributivity: $\forall x, y, z \in A$: $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$
 $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$
(one is enough to state the other)

(2) $\exists \min A =: 0$; $\exists \max A =: 1$

(i.e.: $0 \wedge x = 0$, $0 \vee x = x$; $1 \wedge x = x$, $1 \vee x = 1$.)

(3) $\forall x \in A$ has a complement x' satisfying:

$$x \vee x' = 1 \quad ; \quad x \wedge x' = 0$$

Proposition 1 Let $(A, \vee, \wedge, 0, 1, ')$ be a Boole algebra. Then:

1). $\forall x \in A$, the complement of x is unique

2). $0' = 1$, $1' = 0$, $(x')' = x$

3). $\forall x, y \in A$ we have: $(x \vee y)' = x' \wedge y'$; $(x \wedge y)' = x' \vee y'$ (De Morgan)

Proof 1). We show that x' , \bar{x} are both complements of $x \in A$.

We have:

$$\begin{aligned} x' &= x' \wedge 1 = x' \wedge (x \vee \bar{x}) \stackrel{\text{dist}}{=} (x' \wedge x) \vee (x' \wedge \bar{x}) = \\ &= 0 \vee (x' \wedge \bar{x}) = x' \wedge \bar{x} \end{aligned}$$

$$\bar{x} = \bar{x} \wedge 1 \stackrel{\text{HW}}{=} \bar{x} \wedge (x \vee x') = \bar{x} \wedge x \vee \bar{x} \wedge x' = 0 \vee \bar{x} \wedge x' = \bar{x} \wedge x'$$

$$\text{hence } x' = \bar{x}$$

$$2) \dots 0 \vee 1 = 1, 0 \wedge 1 = 0 \Rightarrow 0' = 1, 1' = 0$$

$$\cdot x \vee x' = 1, x \wedge x' = 0 \Rightarrow (x')' = x$$

3) we calculate

$$\begin{aligned} \bullet (x \vee y) \wedge (x' \wedge y') &\stackrel{\text{dist}}{=} (\underbrace{x \wedge x'}_0 \wedge y') \vee (\underbrace{y \wedge x'}_0 \wedge y') \\ &= (0 \wedge y') \vee (0 \wedge x') = 0 \vee 0 = 0 \end{aligned}$$

$$\begin{aligned} \bullet (x \vee y) \vee (x' \wedge y') &\stackrel{\text{dist}}{=} (\underbrace{x \vee y \vee x'}_1) \wedge (\underbrace{x' \vee y \vee y'}_1) = \\ &= (1 \vee y) \wedge (x \vee 1) = 1 \wedge 1 = 1 \end{aligned}$$

$$\text{Hence } (x \vee y)' = x' \wedge y'$$

Similarly: we prove the other equality! (HW)

Def 5. Any $(B, +, \cdot)$ (comm. with 1)

is called a Boolean ring if

$$\forall x \in B \text{ we have } x^2 = x$$

(every elem. is idempotent)

$\left\{ \begin{array}{l} (B, +) \text{ ab group} \\ (B, \cdot) \text{ monoid} \\ \cdot \text{ distrib. w.r.t } + \end{array} \right.$

Prop Let $(B, +, \cdot, 0, 1)$ be a Boolean ring. Then:

$$1) \quad 1 + 1 = 0 \quad (\text{hence } x + x = 0 \quad \forall x \in B)$$

(we say that B has characteristic 2)

$$2) \quad B \text{ is commutative, i.e. } xy = yx \quad \forall x, y \in B.$$

Proof 1) $1 + 1 = (1 + 1)^2 = 1 + 1 + 1 + 1$; but $(B, +)$ is a group, hence $0 = 1$

2) $x + y = (x + y)^2 = x^2 + xy + yx + y^2 = \cancel{x + y} + xy + yx$

$$\Rightarrow 0 = xy + yx \Rightarrow yx = -xy = xy \quad \text{because } 1 = -1$$

Theorem 2 (Stone).

1) let $(A, \vee, \wedge, 0, 1, ')$ be a Boolean algebra. We define on A the following operations:

$$\bullet \quad x + y \stackrel{\text{def}}{=} (x \wedge y') \vee (x' \wedge y)$$

$$\bullet \quad xy \stackrel{\text{def}}{=} x \wedge y$$

Then: $(A, +, \cdot, 0, 1)$ is a Boolean ring.

2) let $(B, +, \cdot, 0, 1)$ be a Boolean ring.

We define on B the following operations:

$$\bullet \quad x \vee y \stackrel{\text{def}}{=} x + y + xy$$

$$\bullet \quad x \wedge y \stackrel{\text{def}}{=} xy$$

$$\bullet \quad x' \stackrel{\text{def}}{=} 1 + x$$

Then $(B, \vee, \wedge, 0, 1, ')$ is a Boolean algebra.

3) The constructions from 1) and 2) are inverses of each other:

$$\bullet \quad (A, \vee, \wedge, 0, 1, ') \xrightarrow{1) } (A, +, \cdot, 0, 1) \xrightarrow{2) } (A, \cup, \cap, 0, 1, -)$$

Boolean algebra Boolean ring Boolean algebra

$$\text{Then: } \cup = \vee, \cap = \wedge, - = '$$

$$\bullet \quad (B, +, \cdot, 0, 1) \xrightarrow{2) } (B, \vee, \wedge, 0, 1, ') \xrightarrow{1) } (B, \oplus, \odot, 0, 1)$$

Boolean ring Boolean algebra Boolean ring

$$\text{Then: } \oplus = +, \odot = \cdot$$

Example 1) Let $\mathbb{Z}_2 = \{\hat{0}, \hat{1}\}$ be the ring of residues modulo 2.

$\hat{0}$
 \parallel
 $2\mathbb{Z}$
 even

$\hat{1}$
 \parallel
 $2\mathbb{Z}+1$
 odd

$\Rightarrow (\mathbb{Z}_2, +, \cdot)$ is a Boolean ring: $\hat{0}^2 = \hat{0}, \hat{1}^2 = \hat{1}$

We determine the associated Boolean algebra.

$$\begin{cases} x \wedge y = xy \\ x \vee y = x + y + xy \end{cases} \quad x' = 1 + x$$

x	0	1
0	0	1
1	1	1

\wedge	0	1
0	0	0
1	0	1

x	x'
0	1
1	0

$$0 \vee 0 = 0 + 0 + 0 \cdot 0 = 0$$

$$0 \vee 1 = 0 + 1 + 0 \cdot 1 = 1$$

$$1 \vee 1 = 1 + 1 + 1 \cdot 1 = 1$$

c). Consider the struct $(\mathcal{P}(M), \cup, \cap, \emptyset, M, \complement)$.

From set theory we know that this is a Boolean algebra. We determine the corresponding Boolean ring.

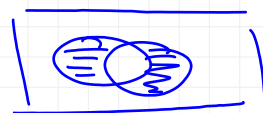
Let $X, Y \in \mathcal{P}(M)$. Then,

$$X \cdot Y = X \cap Y$$

$$X + Y = (X \cap \complement Y) \cup (Y \cap \complement X)$$

$$= (X \setminus Y) \cup (Y \setminus X)$$

$$= X \Delta Y \quad (\text{symmetric difference})$$



We get the Boolean ring $(\mathcal{P}(M), \Delta, \cap, \emptyset, M)$.

3) The Lindenbaum - Tarski algebra

Let $X = \{x_1, \dots, x_n\}$ be a set of atomic formulas in propositional logic.

Let \mathcal{F} be the set of formulas containing these atoms.

On \mathcal{F} we consider the equivalence relation \Leftrightarrow :

$A \Leftrightarrow B$ if and only if the formula:
 $(A \Leftrightarrow B)$ is a tautology.

We consider the quotient set $\mathcal{F}/\Leftrightarrow = \{\hat{A} \mid A \in \mathcal{F}\}$
 where $\hat{A} = \{B \in \mathcal{F} \mid A \Leftrightarrow B\}$

On $\mathcal{F}/\Leftrightarrow$ we define the following operations:

$$\begin{array}{ccc} \hat{A} \vee \hat{B} & \stackrel{\text{def}}{=} & \widehat{A \vee B} \\ \uparrow & & \uparrow \\ \text{operation} & & \text{connective} \end{array}$$

$\hat{0}$ = the set of all contradictions

$$\hat{A} \wedge \hat{B} \stackrel{\text{def}}{=} \widehat{A \wedge B}$$

$\hat{1}$ = the set of all tautologies.

$$\hat{A}^1 \stackrel{\text{def}}{=} \widehat{\neg A}$$

From propositional logic it follows that $(\mathcal{F}/\Leftrightarrow, \vee, \wedge, \hat{0}, \hat{1}, ^1)$ is a Boolean algebra.

Remark (Exercises 2) and 3) give the possibility to use algebraic methods in set theory and propositional logic.

Homework Fill in the proofs of the above theorems.
 ex 7.6 ~~ex 7.3~~ 8.3.