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Introduction

Binary analysis is one of the most important problems in computer security. A wide family of manual and automatic techniques was developed among the years. They are divided into static and dynamic analysis depending on the need to execute the program or not.

Despite the continuous innovation on the front of the automatic analysis, the human part remains essential. A reverse engineer is a person who tries to understand what a binary program does and how. This process usually involves reading and understanding the disassembled code and its effects during execution with the aid of a debugger.

One of the most used automatic techniques nowadays is Symbolic Execution. The idea came about in the '70s but only recently its applications became relevant in computer security. Symbolic execution is used for different tasks, from deobfuscation to vulnerability research (joined with fuzzing) or automatic exploit generation.

It is classified as a static analysis technique because in the pure version of symbolic execution the binary is not really executed. The approach is borderline (we will discuss it later) and the evolution of symbolic execution, Dynamic Symbolic Execution, is a dynamic analysis technique without any doubt.

In the last years, symbolic execution becomes also a first-class technique used by anyone who deals with manual reverse engineering.

Execute symbolically an entire complex software (like a web server for example) is a huge task for a machine, even for a supercomputer so the analyst often uses it in a surgical manner on small pieces of code. A reverse engineer often needs to reverse custom encryption functions or obfuscated code during the dynamic analysis.

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In this thesis, we introduce the idea of combining symbolic execution with dynamic analysis for reverse engineering. The synchronization between a debugger and a symbolic executor can enhance manual dynamic analysis and allow a reverser to easily solve small portions of code without leaving the debugger.

We implemented a synchronization mechanism on top of the binary analysis framework angr. This means a method to transfer the state of the debugged process in the angr environment and back.

The backend library is debugger agnostic and can be extended to work with various frontends. We implemented a frontend for the IDA Pro debugger and one for the GNU Debugger.