## Racket

A.Y. 2023-24 Principles of Programming Languages

Andrea Franchini - andrea.franchini@polimi.it

### Let's code!

You should have already downloaded Racket (~200 MB)!

It contains DrRacket, the IDE we're using for the exercises! It has everything we need.



https://download.racket-lang.org/

### Let's code!

The exercises will be available on my GitHub repository.





https://github.com/andreafra/ppl

# Any doubt?

You can write me an email at:

andrea.franchini@polimi.it

# Racket in 5 minutes Types

#### **NUMBERS**

1 (integer)

3.14 (floating)

1+2i 6.02e+23

#x29 #o32

#b010101

#### **STRINGS**

"Hello world!"

#### **CHARACTERS**

#\a #\B #\5

#### **BOOLEANS**

#t #f

#### **SYMBOLS**

apple a-10!

really? hey/you

#### **VECTORS**

#(1 2 3 4)

#### **LISTS**

'(1 2 #\3 "4") '()

# Racket in 5 minutes Top facts about Racket

- Prefix notation
- Dynamically-typed
- (Mostly) functional programming
- ◆ Homoiconic → Code is Data
- Call-by-Value

### Dynamically-typed

```
In statically-typed languages such as C and Java the type of a variable is checked at compile time, and we usually specify it when declaring something.
```

In Racket, just like in Python and JavaScript, the type is unknown before runtime. You can reassign values of any kind to a variable.

```
int sumFive(int n) {
   return n + 5
}
```

```
; error if n is not a number
(define (sum-five n)
   (+ n 5))
```

We can use (number? n) to check if 'n' is a number.

#### Prefix notation

$$(= (- (+ 3 2) 5) 0)$$

$$3 + 2 - 5 = 0$$

$$(= (+ 3 (- 2 5) 0))$$
Infix

Prefix

#### Functional programming

Procedures are **first-class** citizens. We can assign them to an identifier just like any other value.

```
'lambda's (also 'λ') are unnamed procedures:
```

```
(\lambda (x) (+ x 5))
```

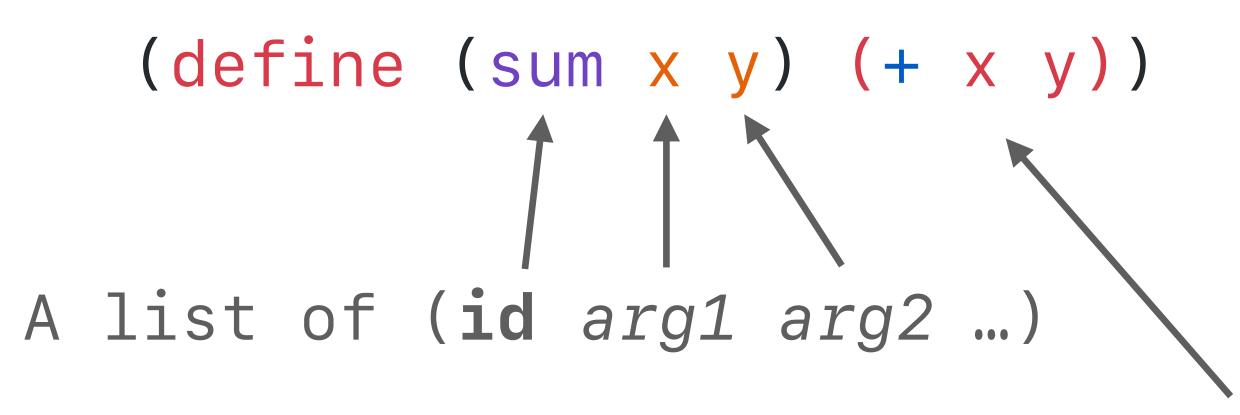
```
Btw, we can assign a value
to an identifier with:
(define two 2)
We can make a named
procedure 'sum-five':
(define sum-five
  (\lambda (x) (+ x 5))
```

#### Functional programming

```
(define sum (λ (x y) (+ x y)))
```

Looks ugly? Here's some syntactic sugar:

Also note that there is no 'return' keyword: the result of the last expression is returned.



body (one or more expressions)

#### Homoiconic: Code is Data

- In (f x y) we're trying to call a function 'f' with 'x' and 'y' as arguments. What if we want them as a list?
- We use 'quote' (aka ' single quote)

```
'(f x y) is now a list
'(+ 1 2); => (+ 1 2)
(+ 1 2); => 3
```

 'quote' prevents evaluation. Remember that S-expressions are lists, where the first elements is the id of a procedure.

#### Homoiconic: Code is Data

• If we want to partially evaluate an expression, like:

```
(1 (+ 2 3) 4) ; we want '(1 5 4)
'(1 (+ 2 3) 4) ; => '(1 (+ 2 3) 4)
```

 'quote' does not work: it makes the expression a constant.

Homoiconic: Code is Data

• We can use 'quasiquote' (aka ` - backtick) and 'unquote' (aka , - comma):

```
(1 , (+ 2 3) 4); => '(1 5 4)

evaluate this
```

- 'unquote' must be used in a 'quasiquote'
- To evaluate a list such as '(+ 1 2 3) we can use 'eval':

```
(eval `(+ 1 2 3))
```

#### Call-by-Value

- Objects are allocated on the heap!
- References are passed by value.
- It behaves just like in Java.
- 'x' in 'set-local' is not the same as the 'x' outside. It is just a copy.

```
(define (set-local x)
  (set! \times 3)
  (displayln x)); => 3
(define x 1)
(set! x 2)
(displayln x) ; => 2
(set-local x)
(displayln x) ; => 2
```

# Racket in 5 minutes Call-by-Value

- In this case, 'v' is just the reference to a vector #(1 2 3).
- Therefore, if we modify it in 'set-vector' we're actually affecting the real (and only) one.

```
(define (set-vector v)
  (vector-set! v 1 "Hi")
  (displayln v)); => #(1 "Hi" 3)

(define v (vector 1 2 3))
  (set-vector v)
  (displayln v); => #(1 "Hi" 3)
```

\* Vectors are fixedlength arrays.

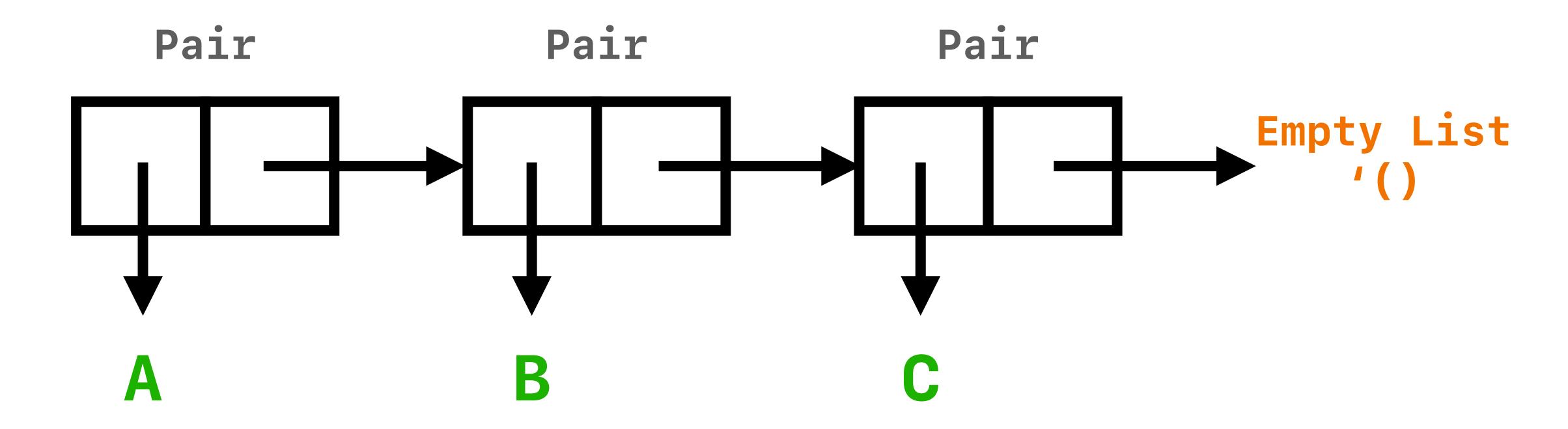
# Racket in 5 minutes Declaring local variables with let

```
    'let' 'let*' allow us to bind variables locally,

 ('letrec' 'letrec*', if you need mutual recursion).
 ; binds in parallel
 (let ((x 2))
       (y 3)) 	← You can't use 'x' in
   the 'y' expression
 ; binds in sequence
 (1et* ((x 2))
        (y (add1 x)))
```

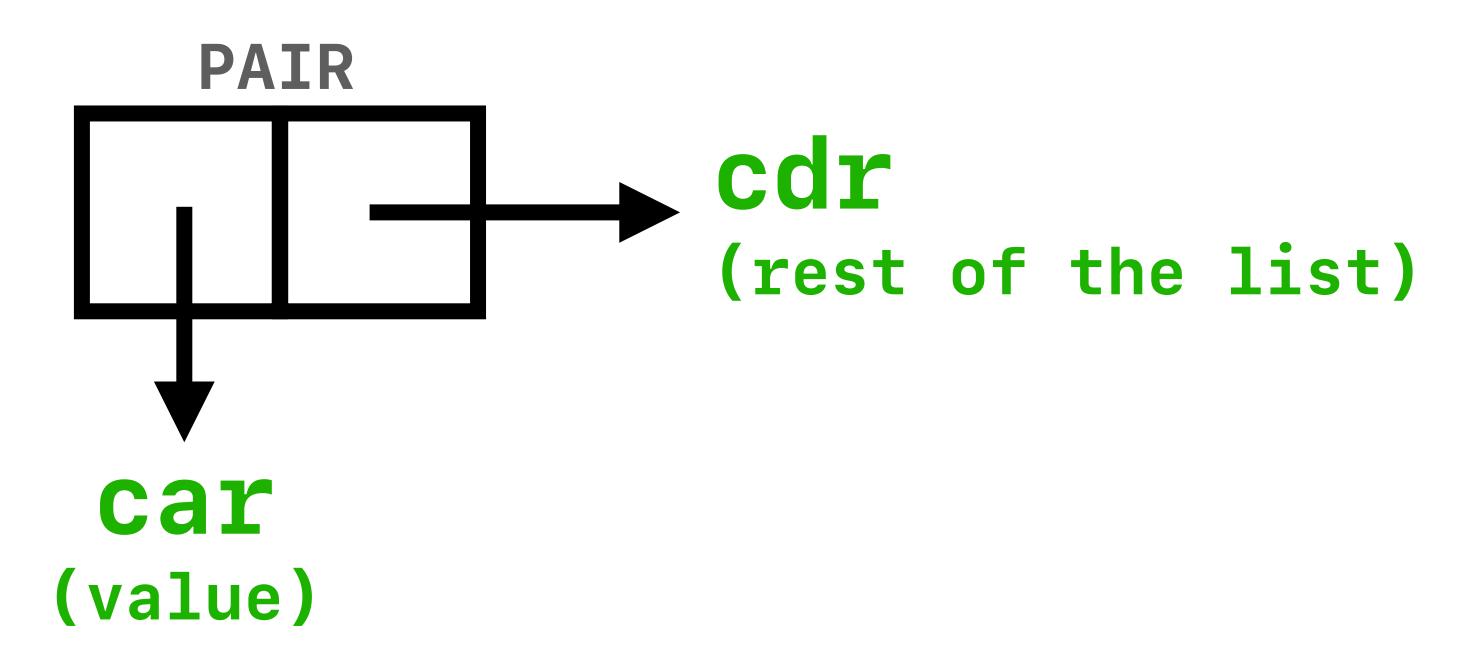
### Racket in 5 minutes Lists & Pairs

• Lists in Racket are a linked list, made up of pairs.



# Racket in 5 minutes Lists & Pairs

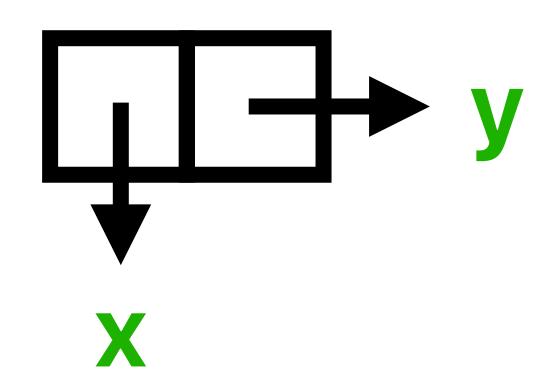
• Lists in Racket works a little different from what you are used to:

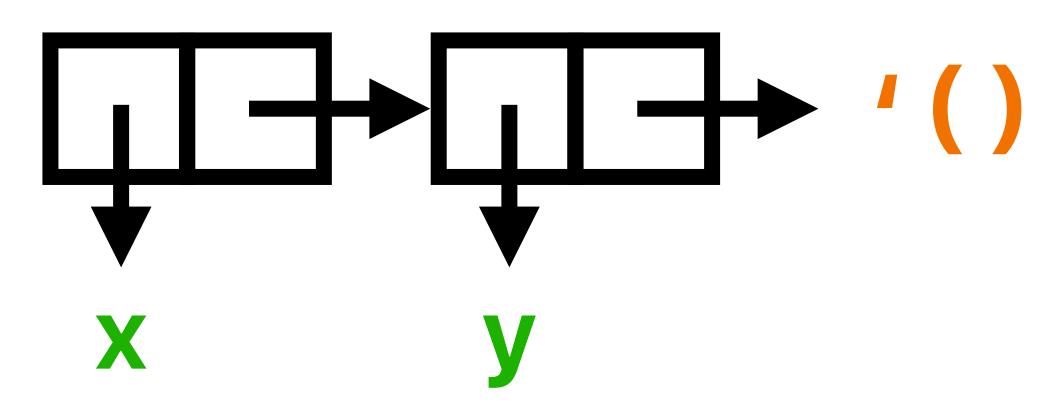


# Racket in 5 minutes Lists and Pairs

- A pair is written as (x . y) and we use '(cons x y)' to create it.
- Pairs are immutable.
- (1 2 3) is stored as (1 . (2 . (3 . ()))) where () is the empty list.

$$(cons x y) ; => '(x . y)$$





# Racket in 5 minutes Lists & Pairs

 We can access the first element of a list with 'car' and the rest of it with 'cdr'.

```
(cdr '(1 2 3)); => 1
(cdr '(1 2 3)); => '(2 3)
```

• If we have a procedure that accepts a variable number of arguments, we can use:

```
(define (f x . xs); e.g.: (f 1 2 3)
  (displayln x); => 1
  (displayln xs)); => (2 3)
```

# Racket in 5 minutes Control Flow

```
• 'if' is a syntactic form (not a procedure nor a value):
 (if <condition> <then-body> <else-body>)
If we want just the 'then' branch we have:
 (when <condition> <then-body>)
• Of course, there's just the 'else' branch, too:
 (unless <condition> <else-body>)
```

# Racket in 5 minutes Control Flow

• 'case' evaluates value and finds the first datum (e.g. '1') for which (equal? value 'datum) is true.

```
(case (+ 7 5)
  [(1 2 3) 'small]
  [(10 11 12) 'big]
  [else 'idk])
matches
```

# Racket in 5 minutes Control Flow

• 'cond' evaluates each test expression until one is true, then evaluates and returns its body.

```
(cond
  [(positive? -5) (error "doesn't get here")]
  [(zero? -5) (error "doesn't get here, either")]
  [else 'here])
```

# Racket in 5 minutes Loops

- Loops are achieved through recursion. Either:
  - → Recursively call a procedure
  - ◆ Use a named let
  - ◆ On lists (and vectors), you can use 'for-each'
- Use tail-recursion whenever possible! The last called thing in a recursive procedure should be the procedure itself!