

Access Control Models

SIO

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Access types

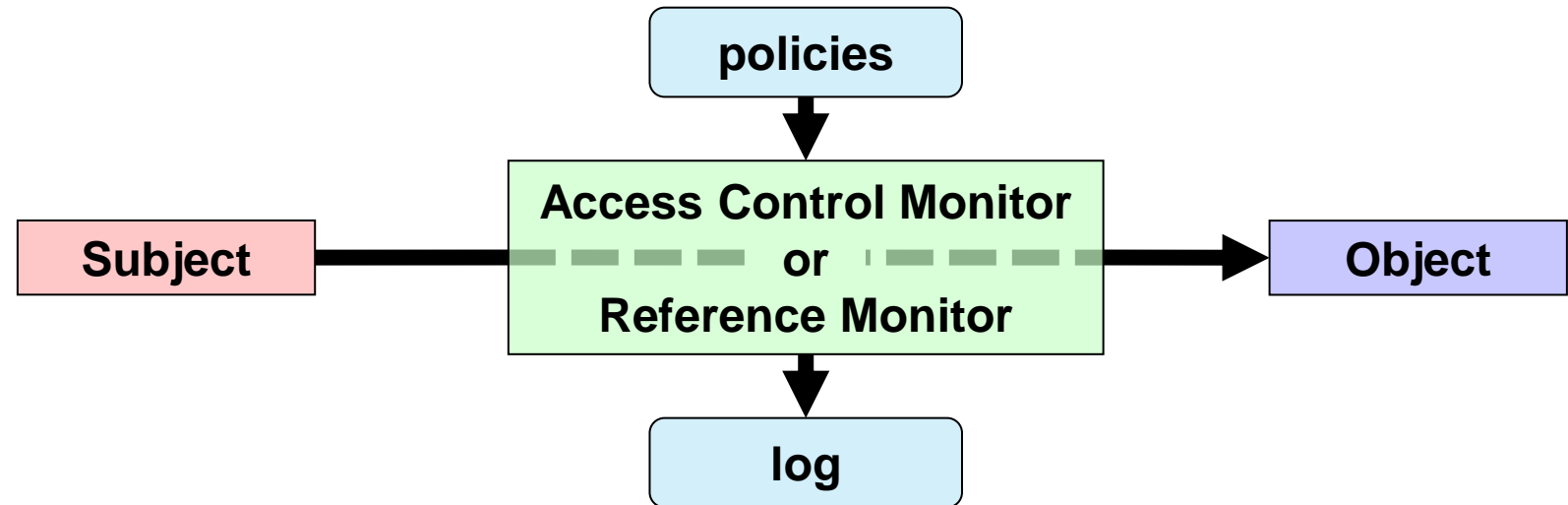
- Physical access

- Physical contact between a subject and the object of interest
 - Facility, room, network, computer, storage device, authentication token, etc.
- Out of scope of this course ...

- Informatic or electronic access

- Information-oriented contact between a subject and the object of interest
 - Contact through request-response dialogs
- Contact is mediated by
 - Computers and networks
 - Operating systems, applications, middleware, devices, etc.

Access control

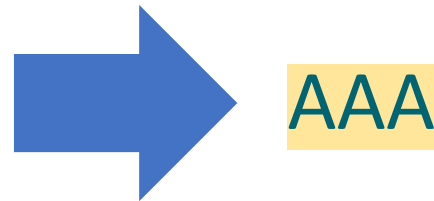


- Definition

- Policies and mechanisms that mediate the access of a subject to an object

- Normal requirements

- Authentication
 - With some Level of Assurance (LoA)
- Authorization policies
- Accountability → logging



Access control: Subjects and Objects

- Both are digital entities
- Subjects are something exhibiting activity:
 - Processes
 - Computers
 - Network elements
- Objects are targets of actions (resources):
 - Stored data
 - CPU time
 - Memory
 - Processes
 - Computers
 - Networks
- An entity can be both subject & object

Least privilege principle

“Every program and every user of the system should operate using the least set of privileges necessary to complete the job”

J. H. Saltzer, M. D. Schroeder, Proc. of The protection of information in computer systems, IEEE, 63(9) 1975

- Privilege:
 - Authorization to perform a given task
 - Similar to access control clearance
- Subjects should have, at any time, the exact privileges required to their assigned tasks
 - Less privileges than the required create unsurpassable barriers
 - More privileges than the required create vulnerabilities
 - Damage resulting from accidents or errors
 - Potential interactions among privileged programs
 - Misuse of a privileges
 - Unwanted information flows
 - "need-to-know" military restrictions

Access control models

- Access control matrix

- Matrix with all access rights for subjects relatively to objects
- Conceptual organization

	O1	O2	...	O _{m-1}	O _m
S1		Access rights			
S2					
...					
S _{n-1}					
S _n					

Access control models

- ACL-based mechanisms
 - ACL: Access Control List
 - Matrix column
- List of access rights for specific subjects
 - Access rights can be positive or negative
 - Default subjects may often be used
- Usually, ACLs are stored along with objects
 - e.g., for file system objects

	O1	O2	...	O _{m-1}	O _m
S1		Access rights			
S2					
...					
S _{n-1}					
S _n					

Access control models

- **Capability-based mechanisms**
 - **Capability**: unforgeable authorization token
 - Matrix row
 - Contains object references and access rights
- **Access granting**
 - Transmission of capabilities between subjects
 - Mediated / non-mediated
- **Usually, capabilities are kept by subjects**
 - e.g., OAuth 2.0 access tokens

	O1	O2	...	O _{m-1}	O _m
S1		Access rights			
S2					
...					
S _{n-1}					
S _n					

Access control models: MAC and DAC

- **Mandatory access control (MAC)**
 - Fixed access control policy implemented by the access control monitor
 - Access control rights cannot be tailored by subjects or object owners
- **Discretionary access control (DAC)**
 - Some subjects can update rights granted or denied to other subjects for a given object
 - Usually this is granted to object owners and system administrators

Access control models: Role-Based Access Control (RBAC)

D.F. Ferraiolo and D.R. Kuhn, "Role Based Access Control",
15th National Computer Security Conference, Baltimore, Oct. 1992

- Not DAC or MAC
 - Roles are dynamically assigned to subjects
 - For access control what matters is the role played by the subject
 - And not the subject's identity
 - Identity is mostly relevant for role access and logging
- Access control binds roles to (meaningful) operations
 - Operations are complex, meaningful system transactions
 - Not the ordinary, low-level read/write/execute actions on individual objects
 - Operations can involve many individual, lower-level objects

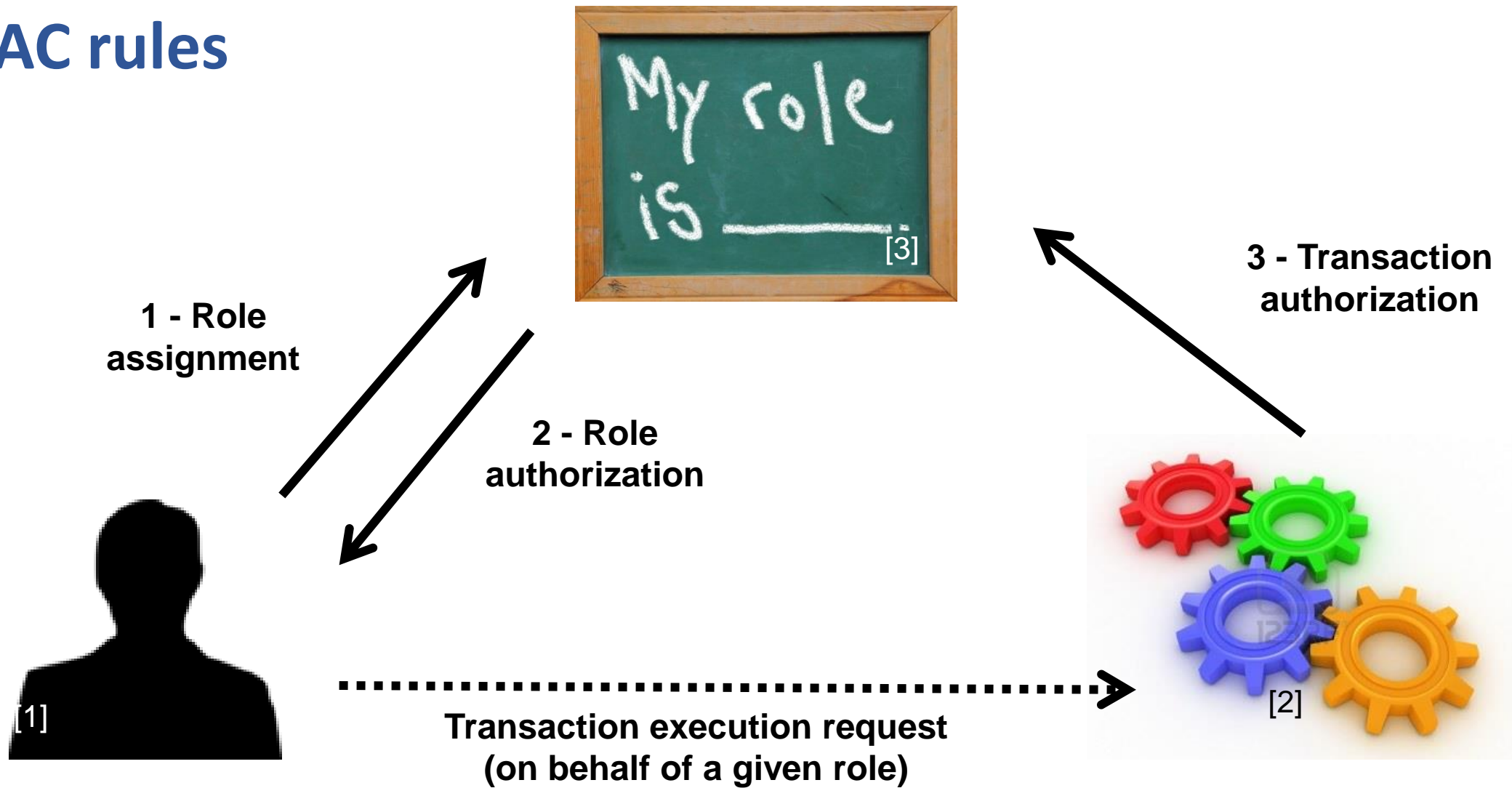
Access control models: RBAC role assignment

- All subject activity on the system is conducted through transactions
 - And transactions are allowed to specific roles
 - Thus, all active subjects are required to have some active role
- A subject can execute a transaction iff it has selected or been assigned a role which can use the transaction

Access control models: RBAC role authorization

- A subject's active role must be authorized for the subject
 - In that case, the subject may assume the role
- Transaction authorization:
 - A subject can execute a transaction iff
 - the transaction is authorized through the subject's role memberships
 - and
 - there are no other constraints that may be applied across subjects, roles, and permissions

RBAC rules



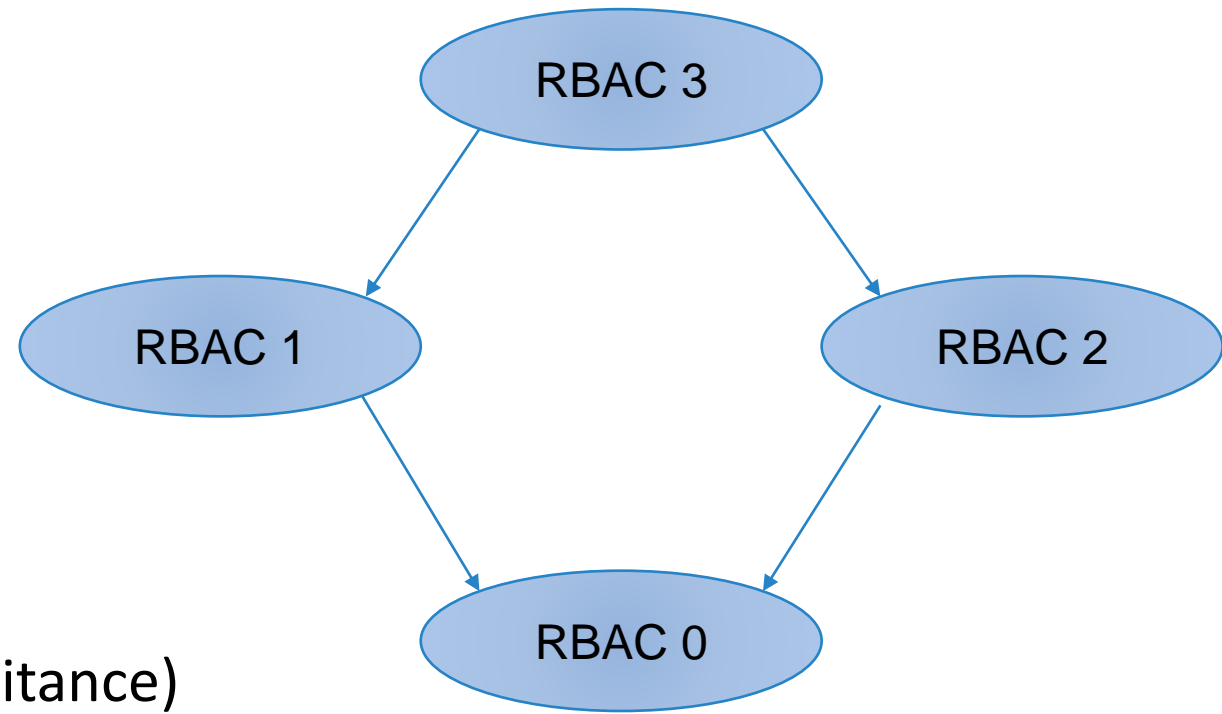
[1] From <http://www.clker.com/clipart-24011.html>

[2] From http://www.123rf.com/photo_12115593_three-dimensional-colored-toothed-wheels.html

[3] From <http://www1.yorksolutions.net/Portals/115255/images/MyRoleIs.jpg>

RBAC variants

- RBAC 0
 - No role hierarchies
 - No role constraints
- RBAC 1
 - RBAC 0 w/ role hierarchies (privilege inheritance)
- RBAC 2
 - RBAC 0 w/ role constraints (separation of duties)
- RBAC 3
 - RBAC 1 + RBAC 2



NIST RBAC model

- Flat RBAC
 - Simple RBAC model w/ **user-role review**
- Hierarchical RBAC
 - Flat RBAC w/ role hierarchies (DAG or tree)
 - General and restricted hierarchies
- Constraint RBAC
 - RBAC w/ role constraints for separation of duty
- Symmetric RBAC
 - RBAC w/ **permission-role review**

User-role review

Which users can have a role?

$\text{Role} \xrightarrow{?} \text{users}$

Which roles can a user have?

$\text{User} \xrightarrow{?} \text{roles}$

Permission-role review

Which permissions has a role?

$\text{Role} \xrightarrow{?} \text{permissions}$

Which roles have a permission?

$\text{Permission} \xrightarrow{?} \text{roles}$

Access control models: Context-Based Access Control (CBAC)

- Access rights have an historical context
 - The access rights cannot be determined without reasoning about past access operations

Ver o histórico do subject antes de dar role

- Example: Stateful packet filter firewall

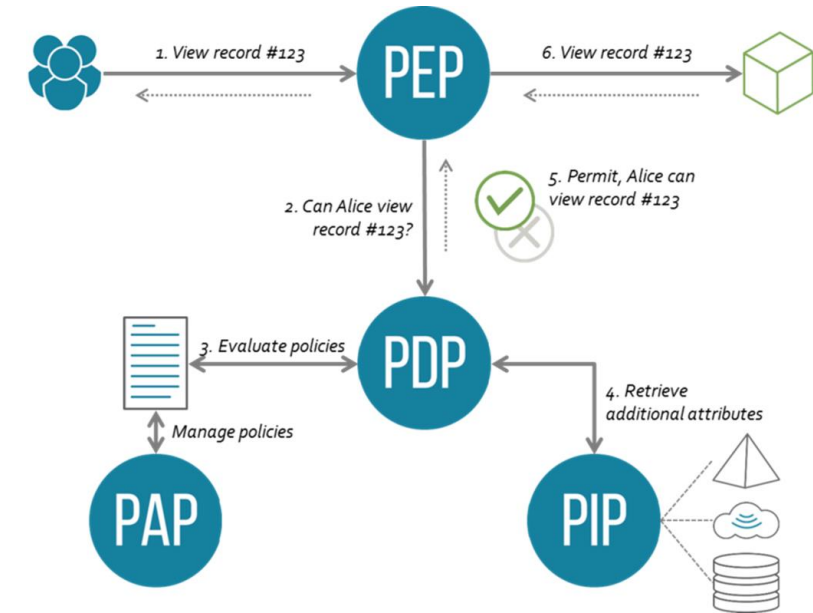
- Example: Chines Wall policy

- Conflict groups
- Access control policies need to address past accesses to objects from members of conflict groups

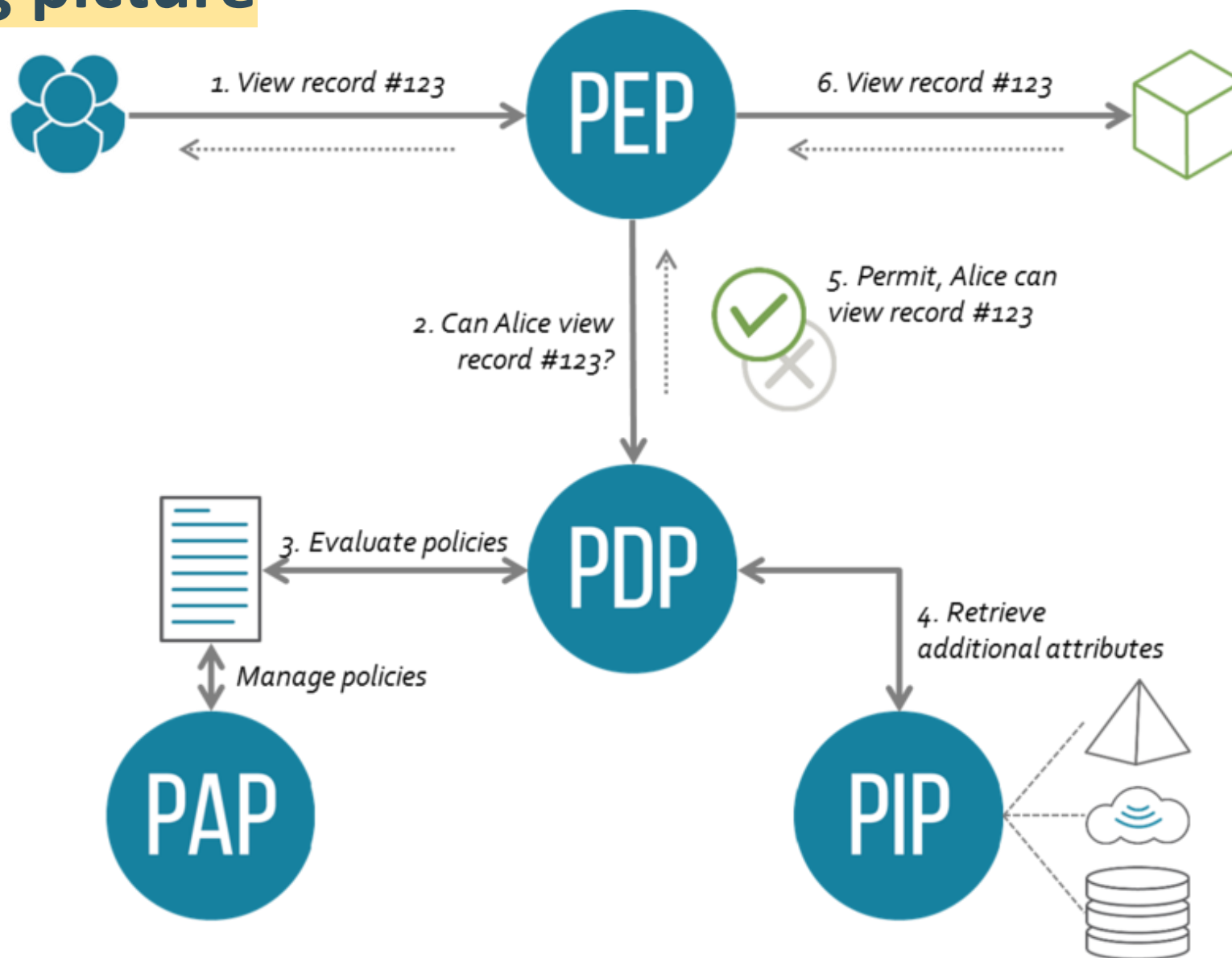
D.F.C. Brewer and M.J. Nash,
"The Chinese Wall Security Policy",
IEEE Symposium on Security and Privacy, 1989

Access control models: Attribute-Based Access Control (ABAC)

- Access control decisions are made based on attributes associated with relevant entities
- OASIS XACML architecture
 - Policy Administration Point (PAP)
 - Where policies are managed
 - Policy Decision Point (PDP)
 - Policy Enforcement Point (PEP)
 - Where resource access requests are intercepted and confronted with PDP's decisions
 - Policy Information Point (PIP)
 - Where the PDP gets external information



XACML big picture



XACML: Access control with PEP and PDP

- A subject sends a request
 - Which is intercepted by the Policy Enforcement Point (PEP)
- The PEP sends an authorization request to the Policy Decision Point (PDP)
 - With some subject's attributes
- The PDP evaluates the authorization request against its policies and reaches a decision
 - Which is returned to the PEP
 - Policies are retrieved from a Policy Retrieval Point (PRP)
 - Useful attributes are fetched from Policy Information Points (PIP)
 - Policies are managed by the Policy Administration Point (PAP)

Access control models: Break-the-glass

- In some scenarios it may be required to overcome the established access limitations
 - e.g., in a life-threatening situation
- In those cases, the subject may be presented with a break-the-glass decision upon a deny
 - Can overcome the deny at their own responsibility
 - Logging is fundamental to prevent abuses

Separation of Duties

R.A. Botha, J.H.P. Eloff, “Separation of duties for access control enforcement in workflow environments”, IBM Systems Journal, 2001

- Fundamental security requirement for fraud and error prevention
 - Dissemination of tasks and associated privileges for a specific business process among multiple subjects
 - Often implemented with RBAC
- Damage control
 - Segregation of duties helps reducing the potential damage from the actions of one person
 - Some duties should not be combined into one position

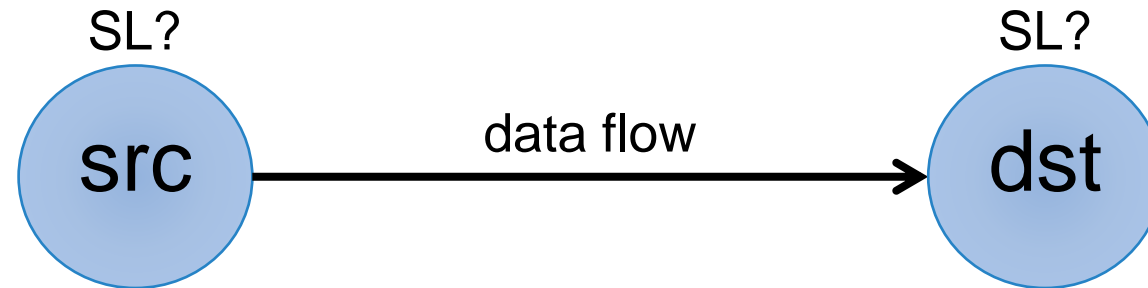
Segregation of duties

ISACA (Inf. Systems Audit and Control Ass.) matrix guideline

Exhibit 2.9—Segregation of Duties Control Matrix													
	Control Group	Systems Analyst	Application Programmer	Help Desk and Support Manager	End User	Data Entry	Computer Operator	Database Administrator	Network Administrator	Systems Administrator	Security Administrator	Systems Programmer	Quality Assurance
Control Group		X	X	X		X	X	X	X	X		X	
Systems Analyst	X			X	X		X				X		X
Application Programmer	X			X	X	X	X	X	X	X	X	X	X
Help Desk and Support Manager	X	X	X		X	X		X	X	X		X	
End User		X	X	X			X	X	X			X	X
Data Entry	X		X	X			X	X	X	X	X	X	
Computer Operator	X	X	X		X	X		X	X	X	X	X	
Database Administrator	X		X	X	X	X	X		X	X		X	
Network Administrator	X		X	X	X	X	X	X					
System Administrator	X		X	X		X	X	X				X	
Security Administrator		X	X			X	X					X	
Systems Programmer	X		X	X	X	X	X	X		X	X		X
Quality Assurance		X	X		X							X	
X—Combination of these functions may create a potential control weakness.													

Information flow models

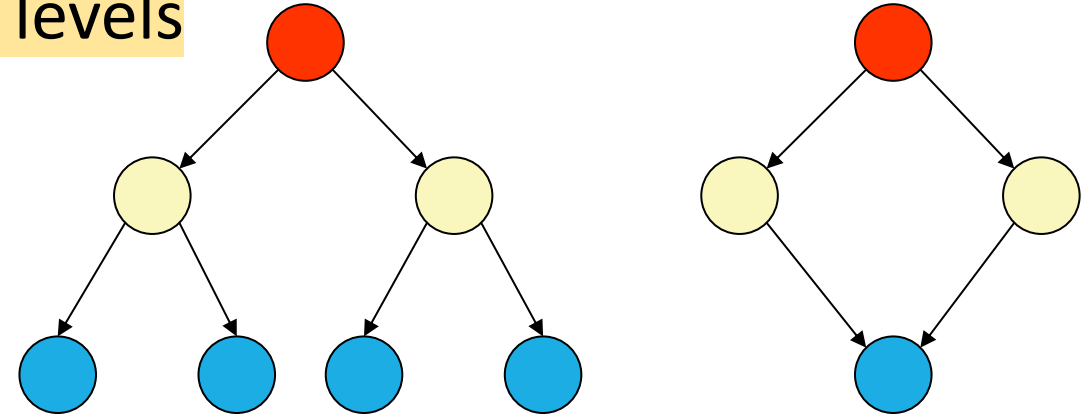
- Authorization is applied to data flows
 - Considering the data flow source and destination
 - Goal: avoid unwanted/dangerous information flows



- Src and Dst security-level attributes
 - Information flows should occur only between entities with given security level (SL) attributes
 - Authorization is given based on the SL attributes

Multilevel security

- Subjects (or roles) act on different security levels
 - Levels do not intersect themselves
 - Levels have some partial order
 - Hierarchy
 - Lattice
- Levels are used as attributes of subjects and objects
 - Subjects: security level clearance
 - Objects: security classification
- Information flows & security levels
 - Same security level → authorized
 - Still, restricted to a “need to know”
 - Different security levels → controlled



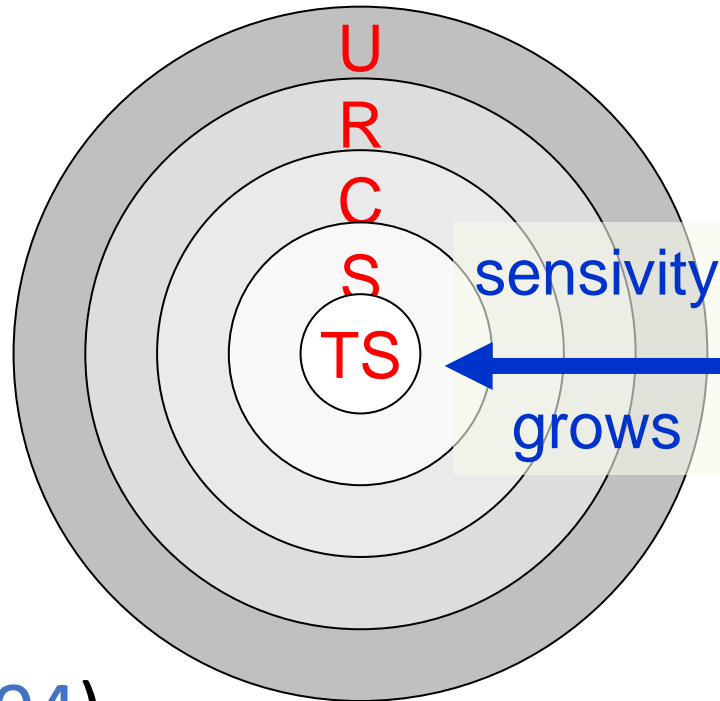
MS levels: Military / Intelligence organizations

- Typical levels

- Top secret
- Secret
- Confidential
- Restricted
- Unclassified

- Portugal ([NTE01](#), [NTE04](#))

- Muito Secreto
- Secreto
- Confidencial
- Reservado



- EU example

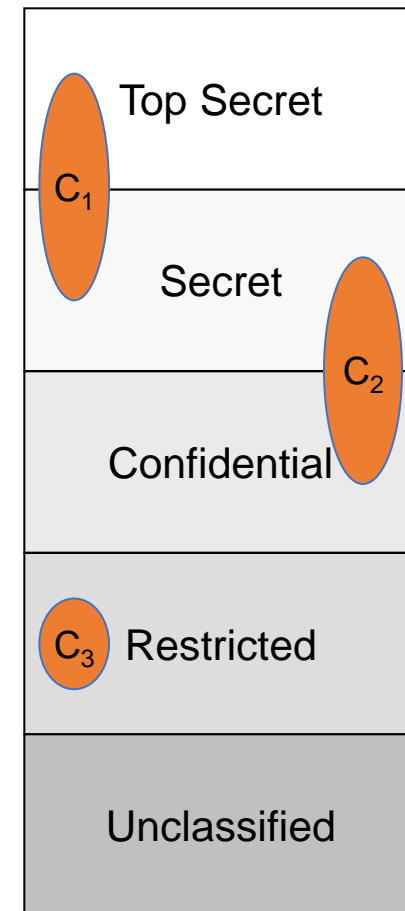
- EU TOP SECRET
- EU SECRET
- EU CONFIDENTIAL
- EU RESTRICTED
- EU COUNCIL / COMMISSION

- NATO example

- COSMIC TOP SECRET (CTS)
- NATO SECRET (NS)
- NATO CONFIDENTIAL (NC)
- NATO RESTRICTED (NR)

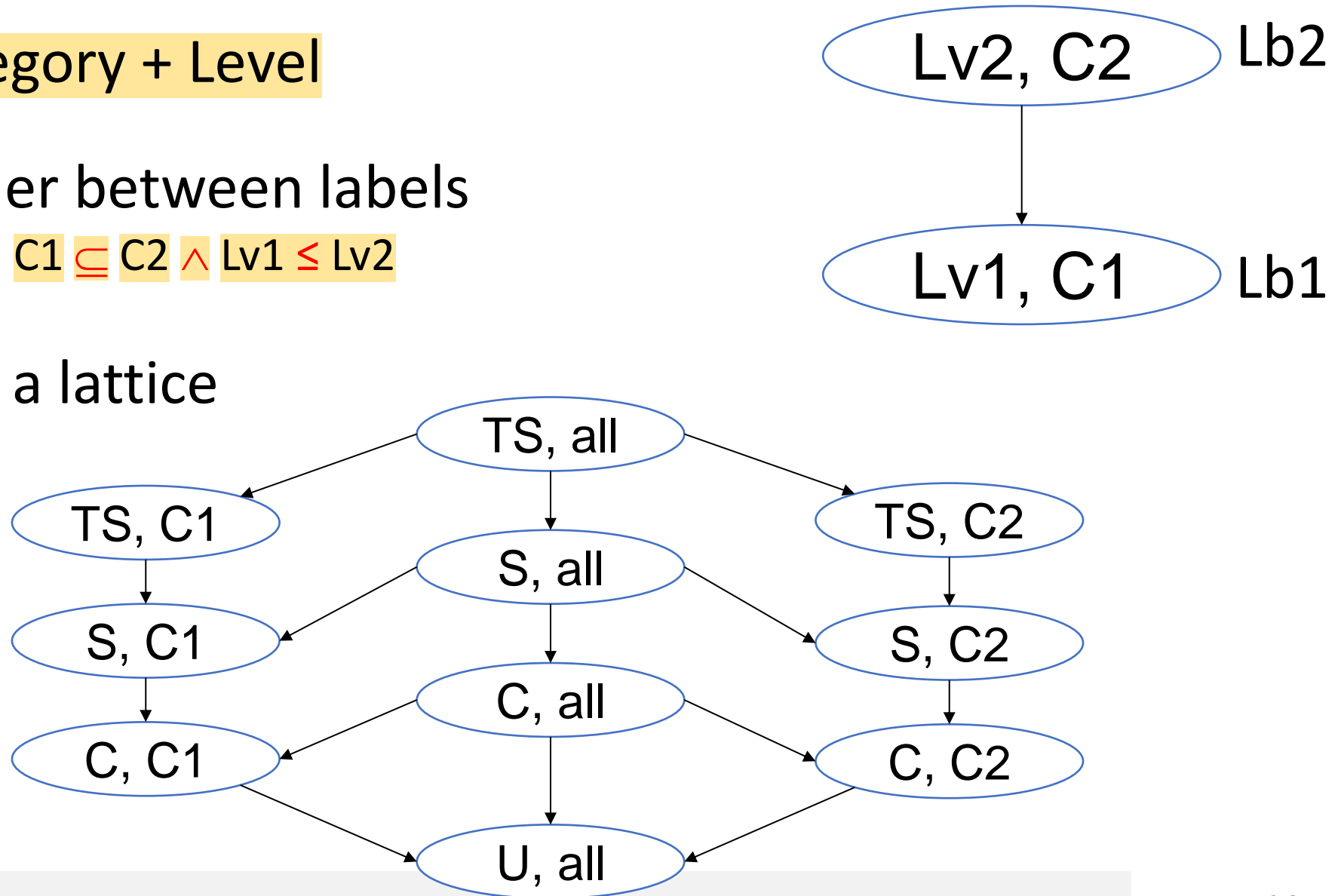
Security categories (or compartments)

- Self-contained information environments
 - May span several security levels
- Military environments
 - Military branches, military units
- Civil environments
 - Departments, organizational units
- An object can belong to different compartments and have a different security classification in each of them
 - (top-secret, crypto), (secret, weapon)



Security labels

- Label = Category + Level
- Relative order between labels
 - $Lb1 \leq Lb2 \Rightarrow C1 \subseteq C2 \wedge Lv1 \leq Lv2$
- Labels form a lattice



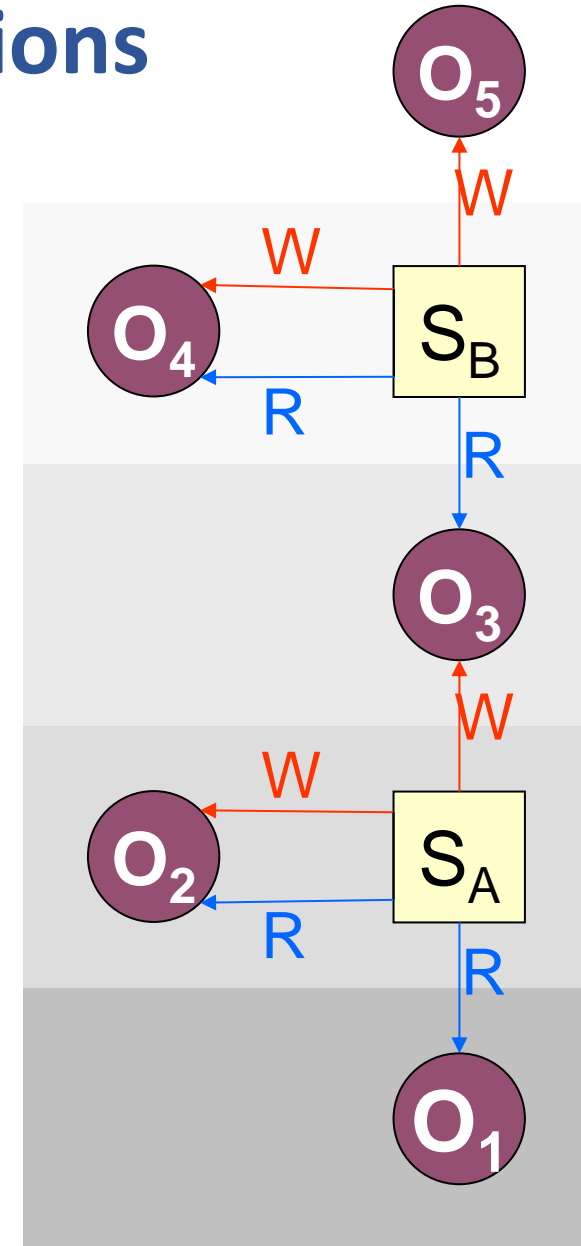
Bell-La Padula MLS Model

D. Elliott Bell, Leonard J. La Padula, "Secure Computer Systems: Mathematical Foundations", MITRE Tech. Report 2547, Vol. I, 1973

- Access control policy for controlling information flows
 - Addresses data confidentiality and access to classified information
 - Addresses disclosure of classified information
 - Object access control is not enough
 - One needs to restrict the flow of information from a source to authorized destinations
- Combines access control matrixes with MLS

Bell-La Padula MLS Model: secure state transitions

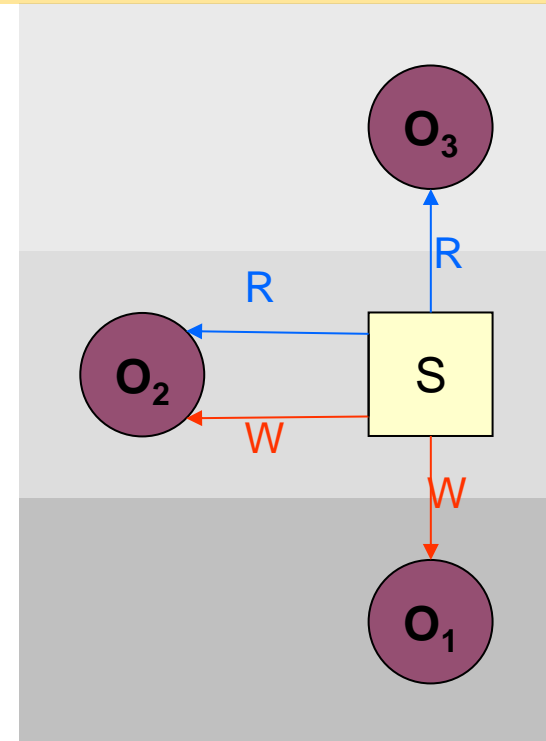
- Simple security condition (no read up)
 - S can read O iff $L(S) \geq L(O)$
- *-property (no write down)
 - S can write O iff $L(S) \leq L(O)$
 - aka confinement property
- Discretionary Security Property
 - DAC-based access control within the same security level



Biba Integrity Model

K. J. Biba, "Integrity Considerations for Secure Computer Systems", MITRE Technical Report 3153, The Mitre Corporation, April 1977

- Access control policy for enforcing integrity control over data flows
 - Uses integrity levels, not security levels
 - Similar to Bell-La Padula, with inverse rules
- Simple Integrity Property (no read down)
 - S can read O iff $I(S) \leq I(O)$
- Integrity *-Property (no write up)
 - S can write O iff $I(S) \geq I(O)$



Windows mandatory integrity control

- Allows mandatory (priority and critical) access control enforcement prior to evaluate DACLs
 - If access is denied, DACLs are not evaluated
 - If access is allowed, DACLs are evaluated

Windows mandatory integrity control: integrity labels

- Untrusted
- Low (or AppContainer)
- Medium
- Medium Plus
- High
- System
- Protected Process

Windows mandatory integrity control: users and processes

- User integrity level
 - Medium: standard users
 - High: elevated users
- Process integrity level
 - The minimum associated to the owner and the executable file
 - User processes usually are Medium or High
 - Except if executing Low-labeled executables
 - Service processes: High

Windows mandatory integrity control

- Securable objects mandatory label
 - NO_WRITE_UP (default)
 - NO_READ_UP
 - NO_EXECUTE_UP