

Organization-Based Analysis of Web-Object Sharing and Caching
Alec Wolma et al.

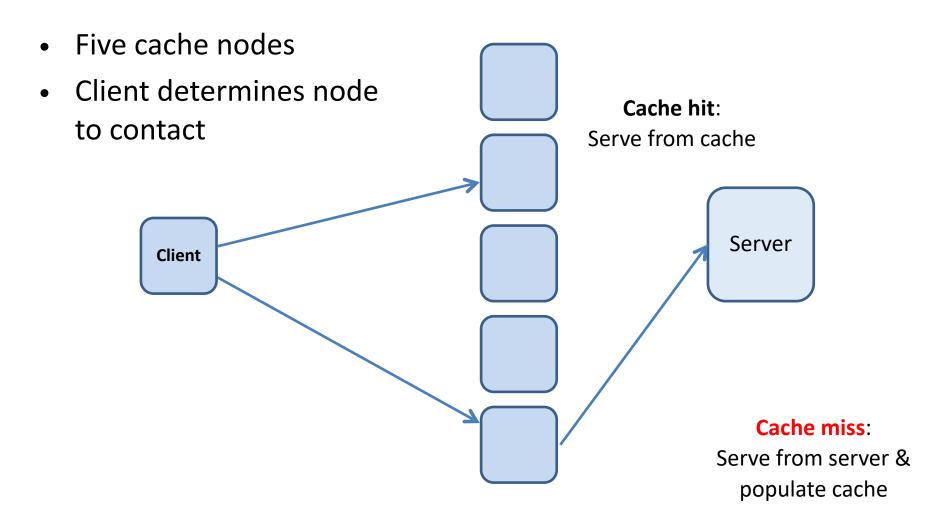
- Five cache nodes
- Client determines node to contact

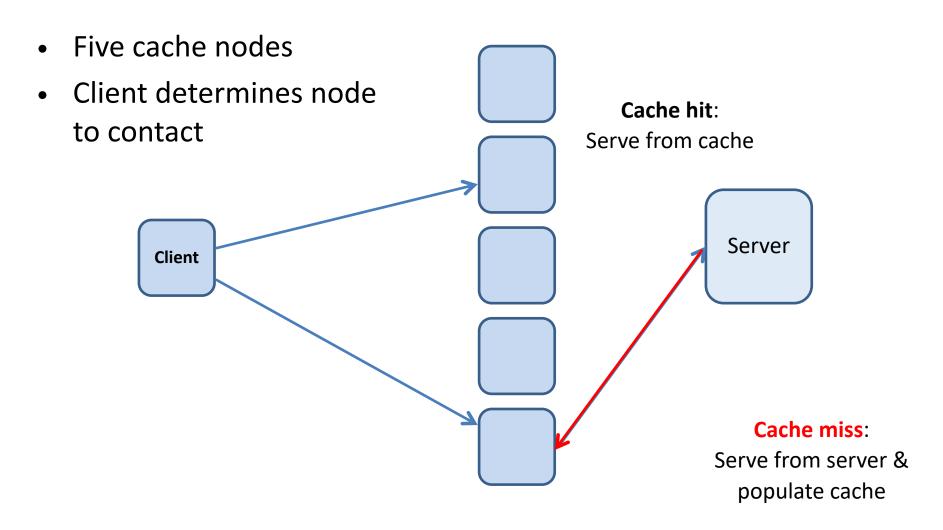
Client

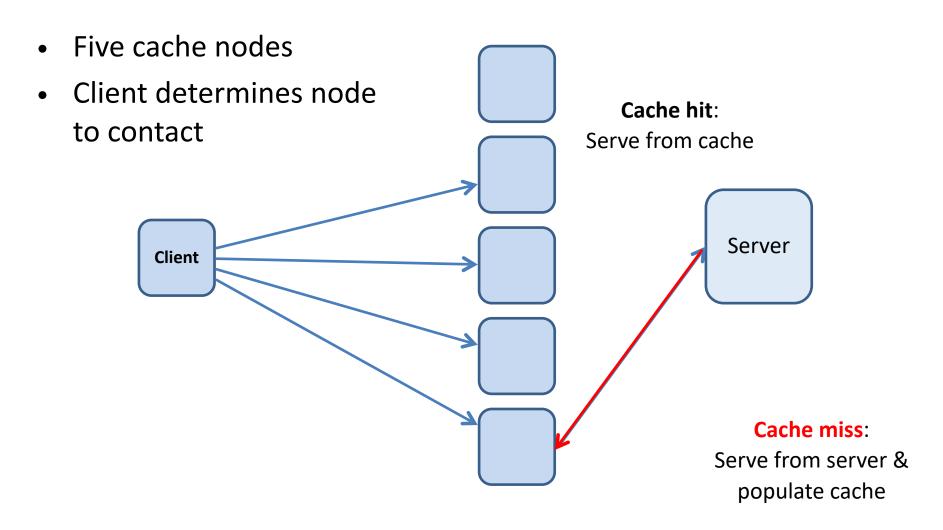


Server

• Five cache nodes Client determines node Cache hit: to contact Serve from cache Server Client







Problem: Mapping objects to caches

- Given a number of caches (e.g., cooperative caching, CDNs, etc.)
- Each cache should carry an equal share of objects
- Clients need to know what cache to query for a given object
- Horizontally partition (shard) object ID space
 - Doesn't work with skewed distributions: e.g., 10 servers, each handles 100 IDs, but all objects have IDs between 1-100 or 900-1000
- Caches should be able to come and go without disrupting the whole operation (i.e., non-effected caches)

Solution attempt: Use hashing

- Map object ID (e.g., URL u) into one of the caches
- Use a hash function that maps u to node h(u)
 - For example, $h(x) = (ax + b) \mod p$, where p is range of h(x), i.e., the number of caches
 - Interpret u as a number based on bit pattern of object ID (or URL)
- Hashing tends to distribute input uniformly across range of hash function
 - Objects (URLs) are equally balanced across caches, even if object IDs are skewed (i.e., highly clustered in ID space)
- No one cache responsible for an uneven share of objects/URLs
- No disproportionately loaded node (potential bottleneck)

Assume, we have **five** caches, numbered 0, ..., 4.

Client

C_o

 C_1

 C_2

 C_3

 $\mathbf{C}_{\mathbf{4}}$

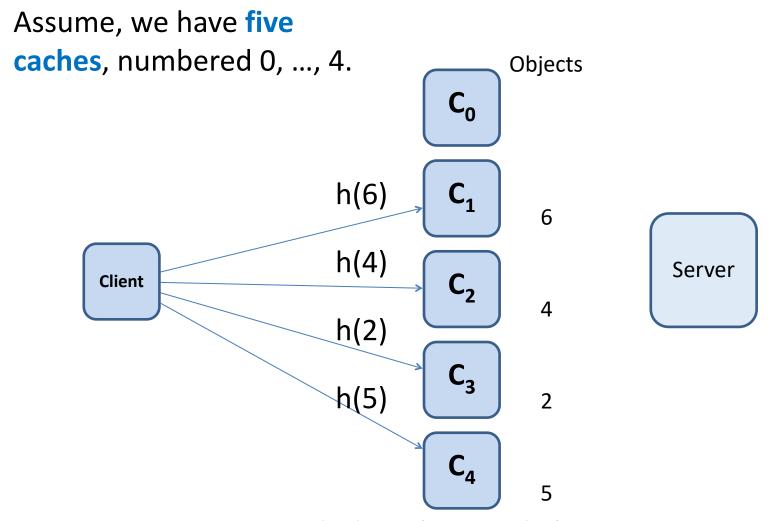
Server

Distributed Systems (Hans-Arno Jacobsen)

Assume, we have five caches, numbered 0, ..., 4. h(6) Server Client

Assume, we have five caches, numbered 0, ..., 4. h(6) Server C_2 Client h(2)

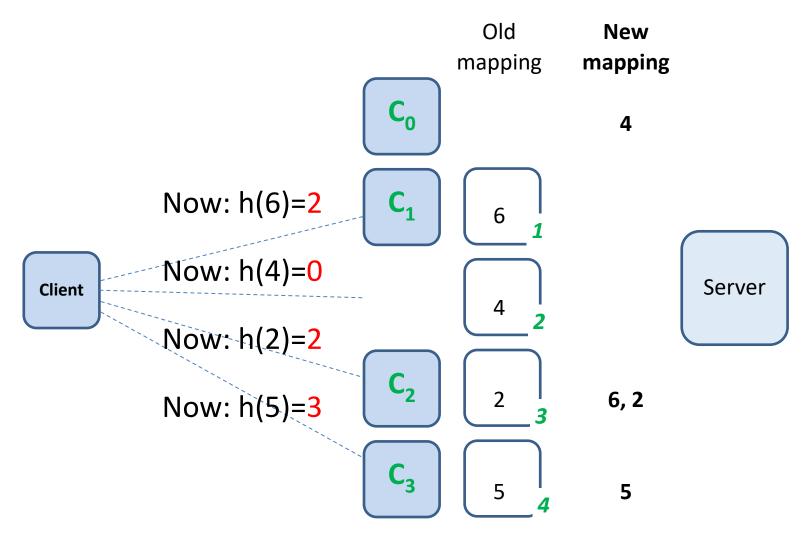
Assume, we have five caches, numbered 0, ..., 4. C_0 h(6) C_1 h(4) Server C_2 Client h(2) h(5) C_4



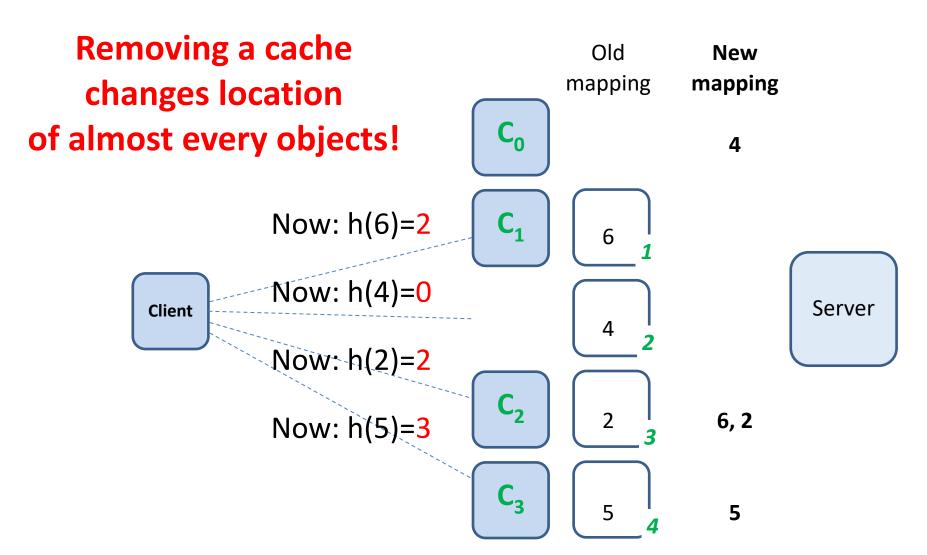
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$h(u) = (7u + 4) \mod 4$ (now have to map across 4 caches)

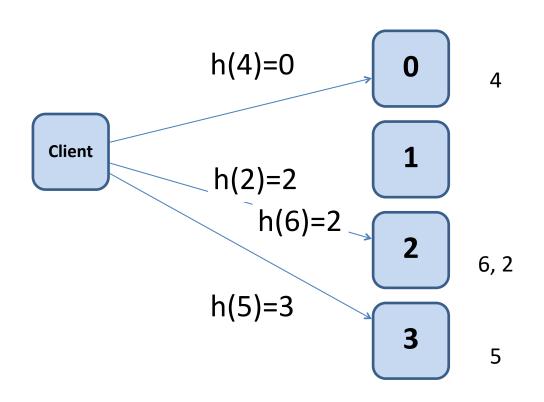


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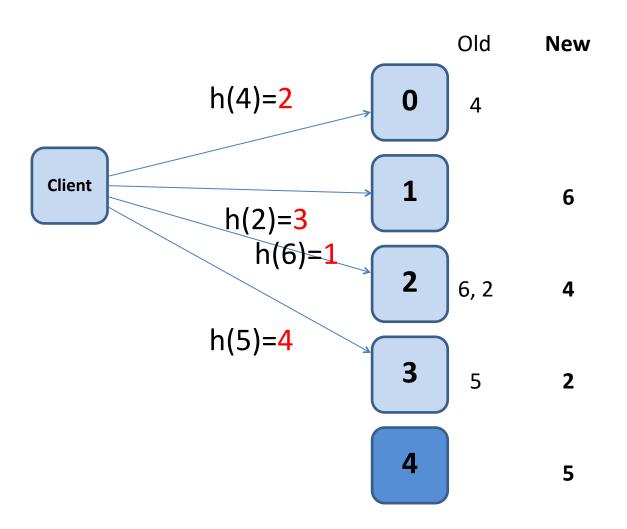


h(u) = 7u + 4 mod 4 (mapped across 4 nodes)

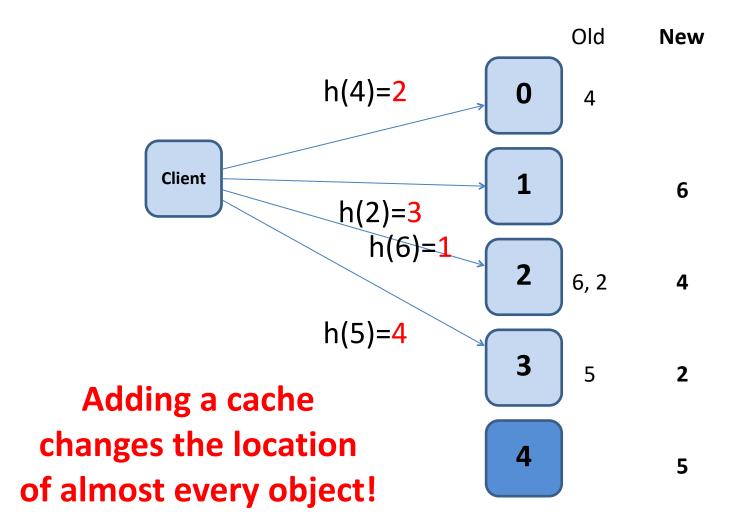
Objects



$h(u) = (7u + 4) \mod 5$ (adding a cache again)



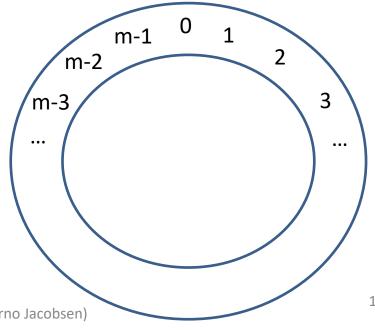
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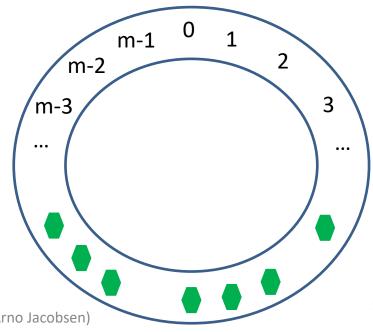
Goals

- Uniform distribution of objects across nodes
- Easily find objects
- Let any client perform a local computation mapping a URL to node that contains referenced object
- Allow for nodes to be added/removed without much
 disruption remap only n/m objects (n objects, m slots)
- D. Karger et al., MIT, 1997
- Basis for Akamai
 - CDN company (content distribution network)
 - Web cache as a service

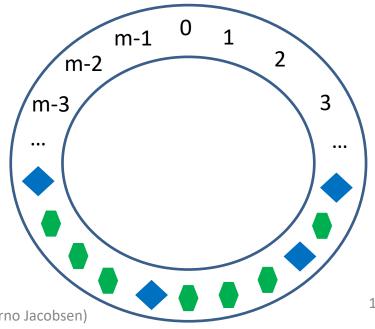
- Select a **base hash function** that maps input identifier to the number range [0, ..., m-1]
- E.g., $h(x) = (ax + b) \mod m$
- Interpret range of h(..) as array that wraps around (i.e., a circle)
- h(..) gives slot in array (circle) and wraps around at m-1 to 0
- Each object is mapped to a slot via h(..)
- Each cache is mapped to a slot via h(..)
- Assign each object to the closest cache slot in clockwise direction on the circle



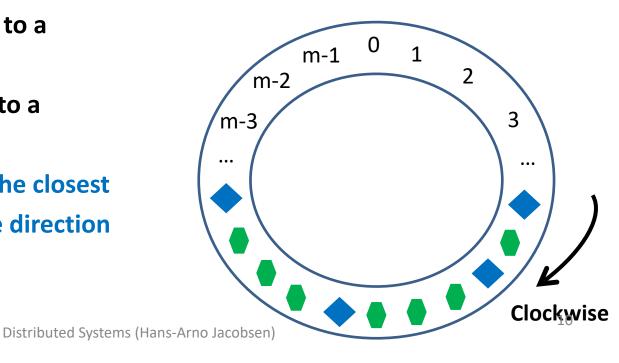
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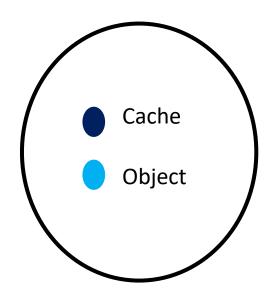
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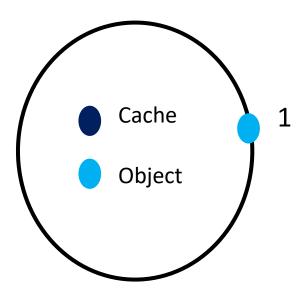
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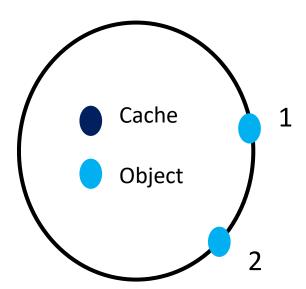
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- Divide by M, re-mapping [0,...,M] to [0, 1]
- Interpret this interval as the unit circle: Here, circle with circumference 1 (normally radius 1)
- Each object is mapped to a point on unit circle via h(..)
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- Assign each URL to closest cache point in clockwise direction on the circle



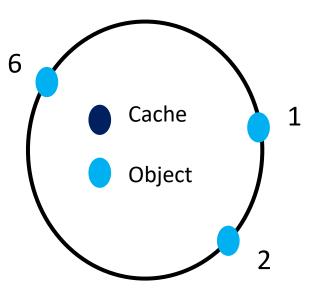
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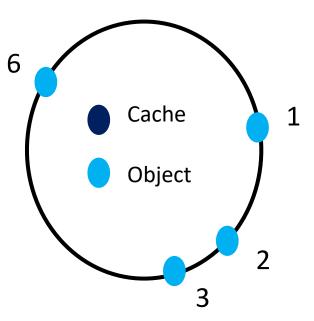
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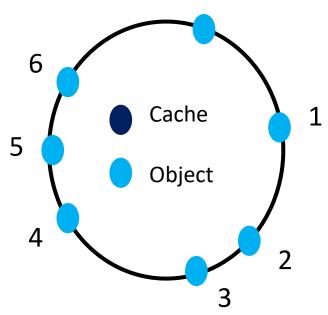
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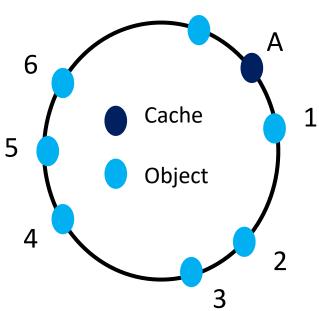
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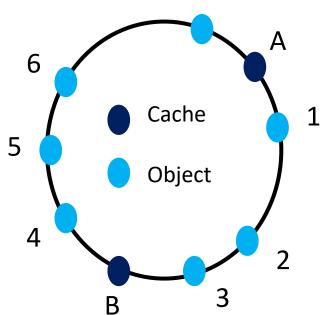
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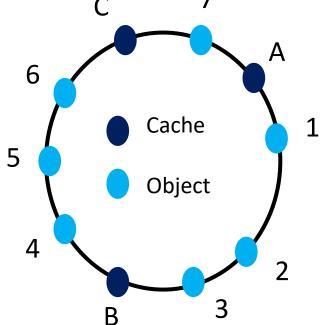
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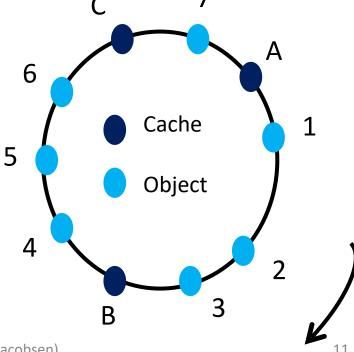
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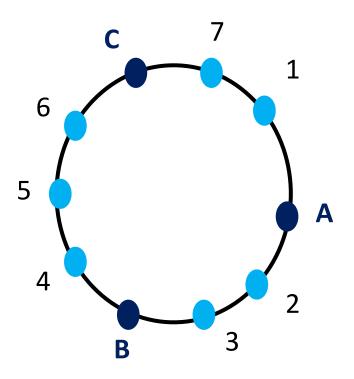
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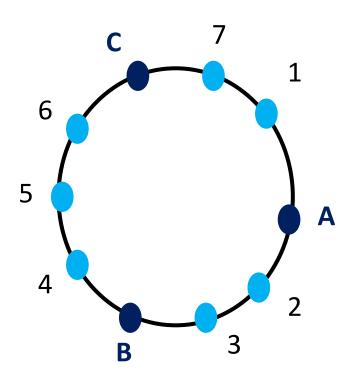
Mapping items to caches



Items 2, 3 mapped to B
Items 4, 5, 6 mapped to C
Items 7, 1 mapped to A

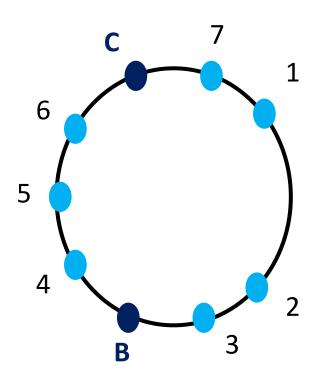


Removing a cache



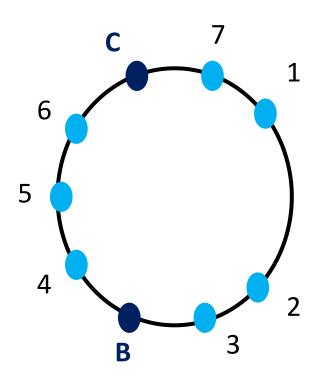
Items	2, 3	mapped to B
Items	4, 5, 6	mapped to C
Items	7, 1	mapped to A

Removing a cache



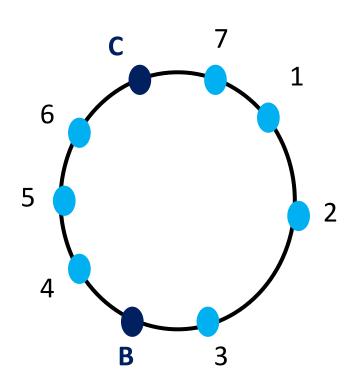
Items 2, 3 mapped to B
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Removing a cache



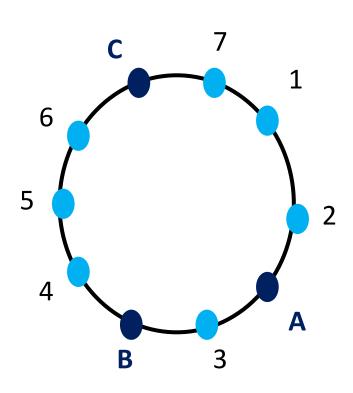
Items 2, 3, **7, 1** mapped to **B**Items 4, 5, 6 mapped to **C**

Adding a cache



Items 7, 1, 2, 3 mapped to B
Items 4, 5, 6 mapped to C

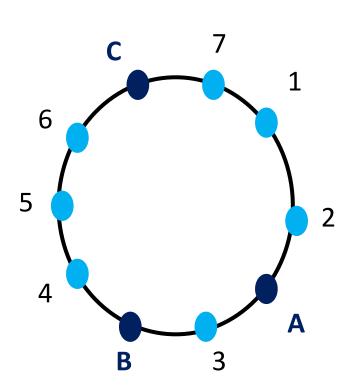
Adding a cache



Items 7, 1, 2, 3 mapped to B

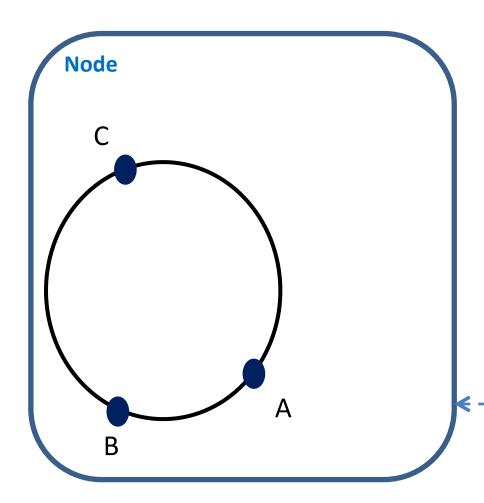
Items 4, 5, 6 mapped to C

Adding a cache



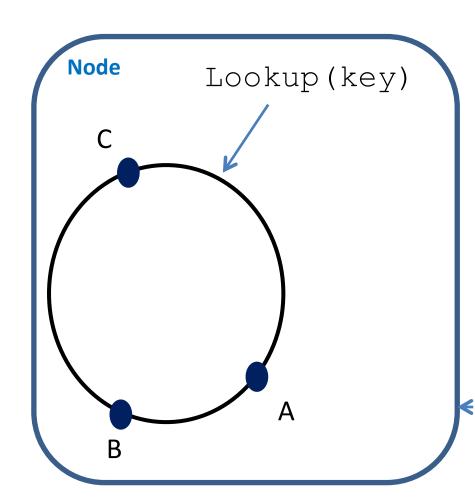
Items 3 mapped to B
Items 4, 5, 6 mapped to C
Items 7, 1, 2 mapped to A

Processing a Lookup (key)



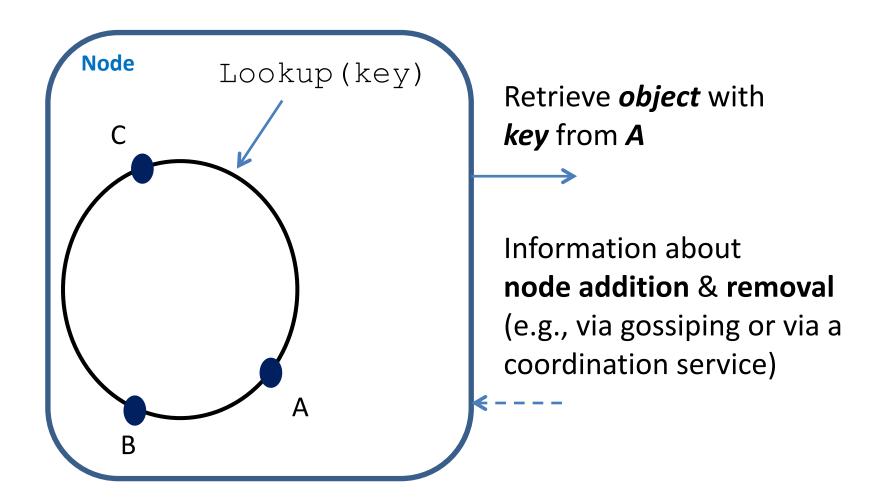
Information about node addition & removal (e.g., via gossiping or via a coordination service)

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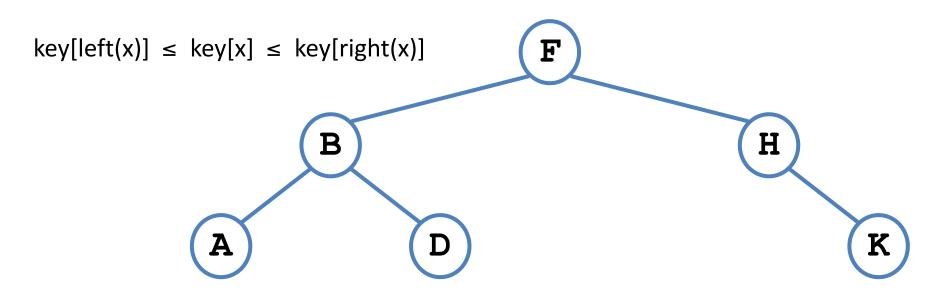
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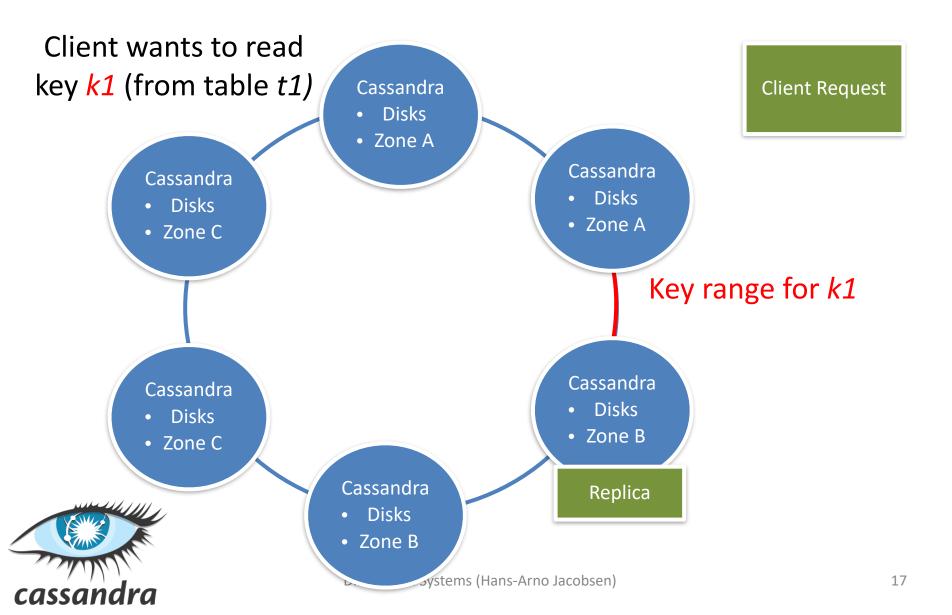
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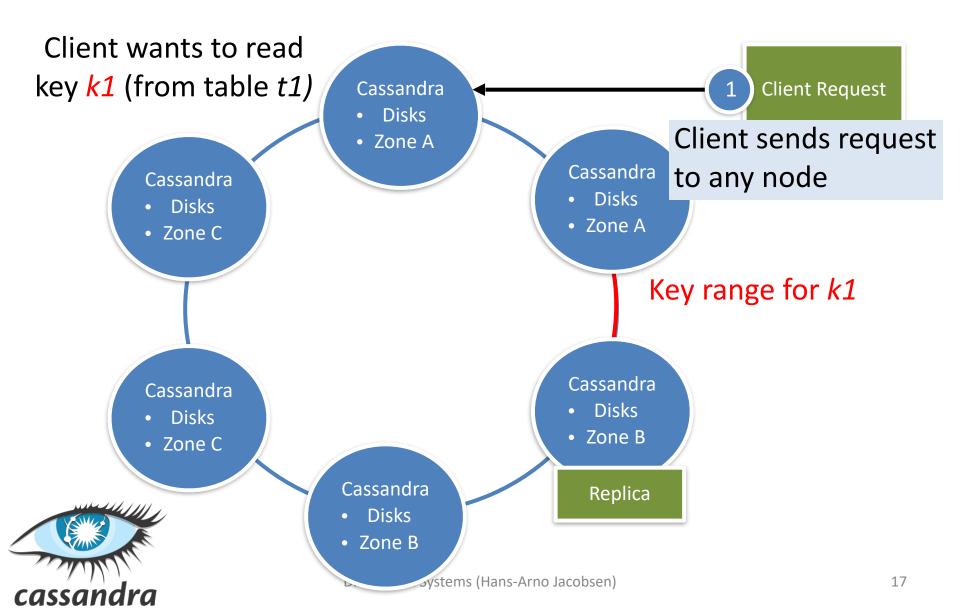


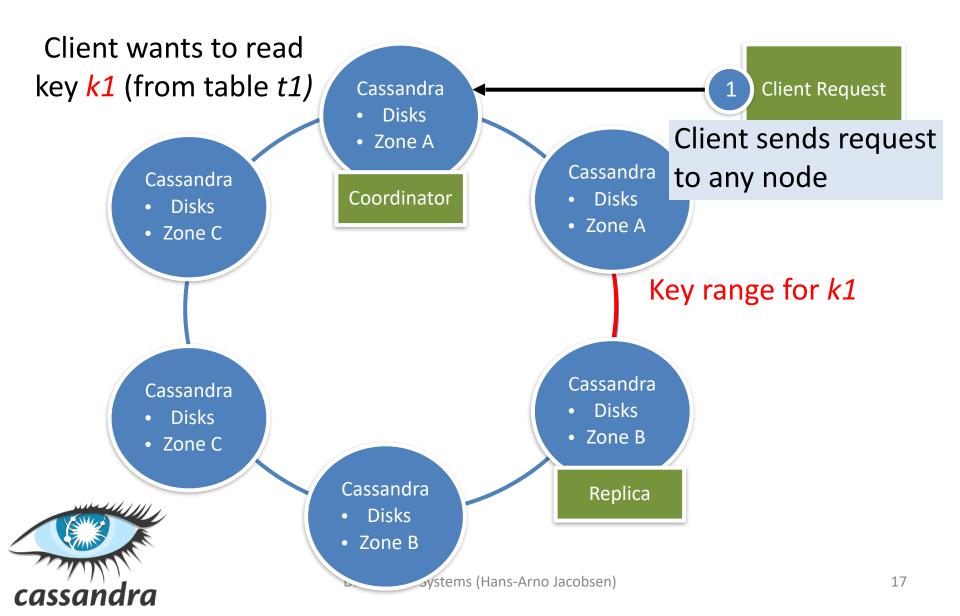
Cache lookup data structure at each node

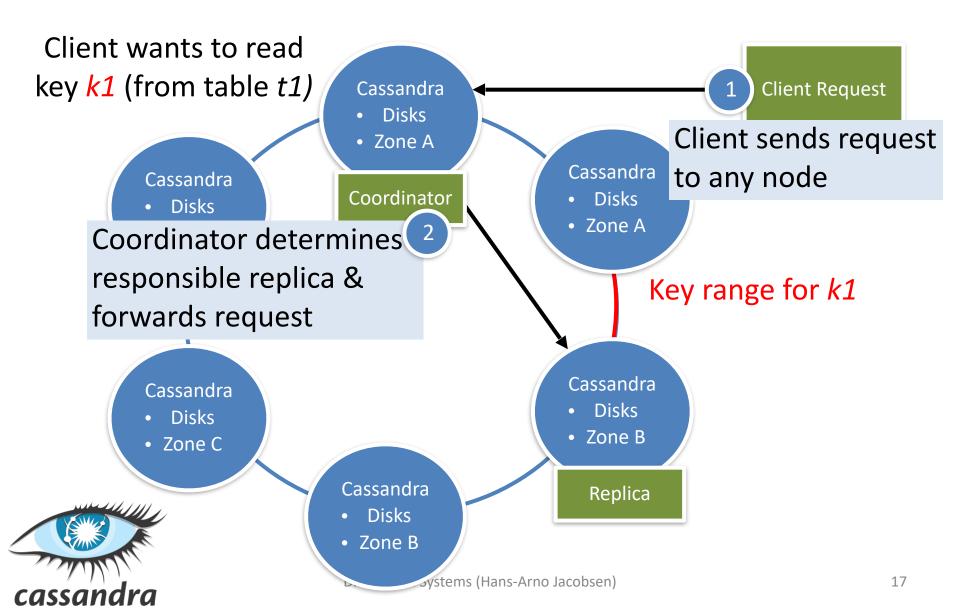
- Store cache points in a binary tree
- Find clockwise successor of a URL point by single search in tree (takes O(log n) time)
- For a constant time technique, cf. Karger et al., 1997

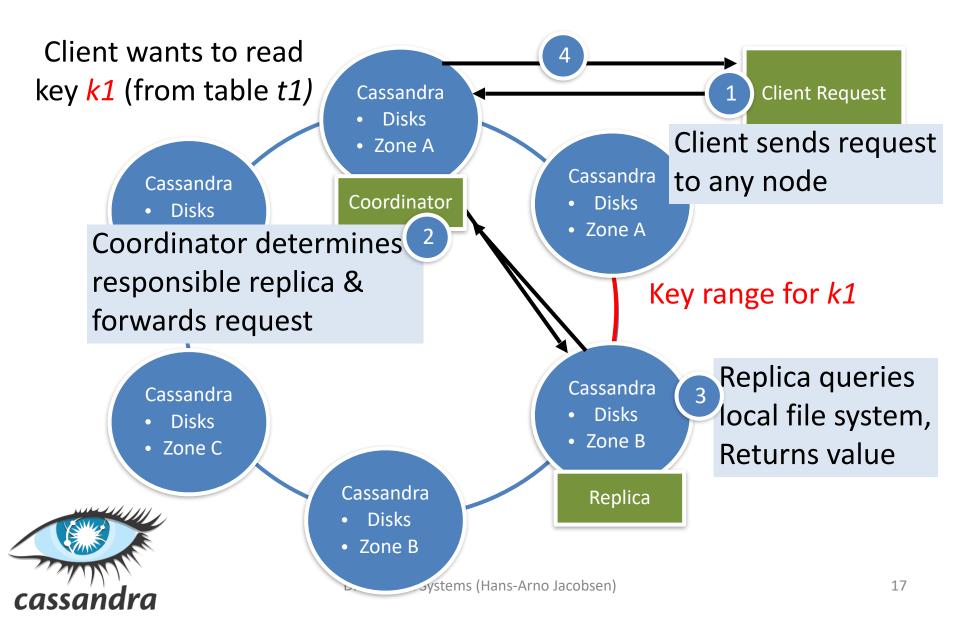












Base hash function: MD5

- Message Digest 5 (MD5), R. Rivest, 1992 (MD1, ..., MD6)
- Hash function that produces a 128-bit (16-byte) hash value
- Maps variable-length message into a fixed-length output
- MD5 hash is typically expressed as a hex number (32 digits)
- It's been shown that MD5 is not collision resistant
- US-CERT about MD5 "should be considered cryptographically broken and unsuitable for further use" (for security, not for caching)
- SHA-2 is a more appropriate cryptographic hash function
- For consistent hashing, MD5 is sufficient

MD5 examples

 MD5("The quick brown fox jumps over the lazy dog") = 9e107d9d372bb6826bd81d3542a419d6

MD5("The quick brown fox jumps over the lazy dog.") = e4d909c290d0fb1ca068ffaddf22cbd0

• MD5("") = d41d8cd98f00b204e9800998ecf8427e

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• MD5("") = d41d8cd98f00b204e9800998ecf8427e

Self-study questions

- How would you use MD5 and SHA2 instead of h(..) from our slides for consistent hashing?
- Apply h(..), MD5, SHA2 to a URL, where does the output map on the unit circle?
- Discuss pros and cons of having a given caching server map to one vs. more points on the unit circle.
- What are the implications of a slash-dot effect on consistent hashing?