“For every car, there is a driver who, if (s)he can start it, then (s)he can stop it.”

can be formalized in a suitable language as

$$\forall x(\mathsf{Car}(x) \rightarrow \exists y(\mathsf{Driver}(y) \wedge (\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y)))).$$

Now, we negate:

$$\begin{aligned}
&\neg\forall x(\mathsf{Car}(x) \rightarrow \exists y(\mathsf{Driver}(y) \wedge (\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y))))\\
&\equiv \exists x\neg(\mathsf{Car}(x) \rightarrow \exists y(\mathsf{Driver}(y) \wedge (\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y))))\\
&\equiv \exists x(\mathsf{Car}(x) \wedge \neg\exists y(\mathsf{Driver}(y) \wedge (\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y))))\\
&\equiv \exists x(\mathsf{Car}(x) \wedge \forall y\neg(\mathsf{Driver}(y) \wedge (\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y))))\\
&\equiv \exists x(\mathsf{Car}(x) \wedge \forall y(\neg\mathsf{Driver}(y) \vee \neg(\mathsf{Start}(x,y) \rightarrow \mathsf{Stop}(x,y))))\\
&\equiv \exists x(\mathsf{Car}(x) \wedge \forall y(\neg\mathsf{Driver}(y) \vee (\mathsf{Start}(x,y) \wedge \neg\mathsf{Stop}(x,y)))).
\end{aligned}$$

Since $\neg A \vee B \equiv A \rightarrow B$, the last formula is equivalent to

$$\exists x(\mathsf{Car}(x) \wedge \forall y(\mathsf{Driver}(y) \rightarrow (\mathsf{Start}(x,y) \wedge \neg\mathsf{Stop}(x,y)))).$$

Thus, the negation of the sentence above is equivalent to:

“There is a car such that every driver can start it and cannot stop it.”