

**Experiment 2.** We modify the configuration from Experiment 1 by applying Clulow’s first suggestion: that attribute changing operations be prevented from allowing a stored key to have both **wrap** and **decrypt** set. Note that in order to do this, it is not sufficient merely to check that **decrypt** is unset before setting **wrap**, and to check **wrap** is unset before setting **decrypt**. One must also add **wrap** and **decrypt** to the list of sticky attributes which once set, may not be unset, or the attack is not prevented, [17]. Having applied these measures, we discovered a previously unknown attack, given in Figure 3. The intruder imports his own key  $k_3$  by first encrypting it under  $k_2$ , and then unwrapping it. He can then export the sensitive key  $k_1$  under  $k_3$  to discover its value.

**Initial state:** The intruder knows the handles  $h(n_1, k_1)$ ,  $h(n_2, k_2)$  and the key  $k_3$ ;  $n_1$  has the attributes **sensitive** and **extract** set whereas  $n_2$  has the attributes **unwrap** and **encrypt** set.

**Trace:**

SEncrypt:	$h(n_2, k_2), k_3$	$\rightarrow$	$\text{senc}(k_3, k_2)$
Unwrap:	$h(n_2, k_2), \text{senc}(k_3, k_2)$	$\xrightarrow{\text{new } n_3}$	$h(n_3, k_3)$
Set_wrap:	$h(n_3, k_3)$	$\rightarrow$	$\text{wrap}(n_3)$
Wrap:	$h(n_3, k_3), h(n_1, k_1)$	$\rightarrow$	$\text{senc}(k_1, k_3)$
Intruder:	$\text{senc}(k_1, k_3), k_3$	$\rightarrow$	$k_1$

**Figure 3. Attack discovered in Experiment 2**