

# CPSC 322: Introduction to Artificial Intelligence

Logics: Propositional and Propositional  
Definite Clause Logic: Syntax and  
Semantics

Textbook reference: [[5.1](#),[5.3](#)]

Instructor: Varada Kolhatkar  
University of British Columbia

Credit: These slides are adapted from the slides of the previous offerings of the course. Thanks to all instructors for creating and improving the teaching material and making it available!

# Announcements

- Assignment 3 has been released.
  - Due date: **Nov 11, 11:59 PM**
- Final exam scheduled: **Dec 9 at 7:00pm**

# Lecture outline

- Recap: Planning 📌
- Logic introduction
- Propositional logic: Syntax and semantics
- Propositional Definite Clause (PDC) Logic: Syntax
- Propositional Definite Clause (PDC) Logic: Semantics

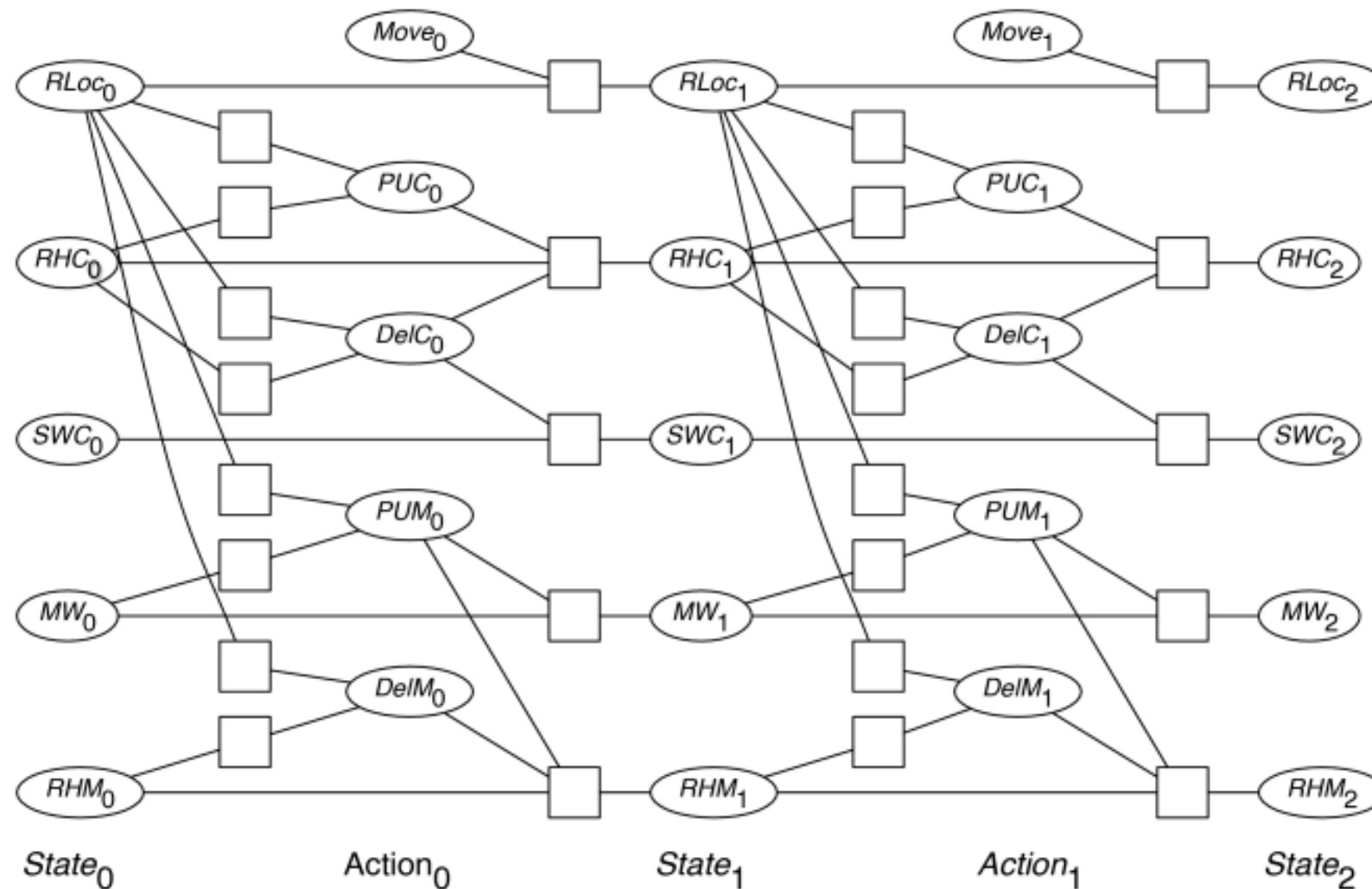
# Recap: Planning (pair-share)

- The STRIPS representation for an action consists of what?
- What heuristic did we use to solve planning problem using forward search?
- What is meant by the *horizon* in a CSP planning problem?
- What are the precondition and effect constraints?

# Solve planning as CSP: Pseudo code


```
solved = false
horizon = 0
while solved = false
    map STRIPS into CSP with horizon
    solve CSP → solution
    if solution then
        solved = T
    else
        horizon = horizon + 1
Return solution
```

# Plan after solving the CSP

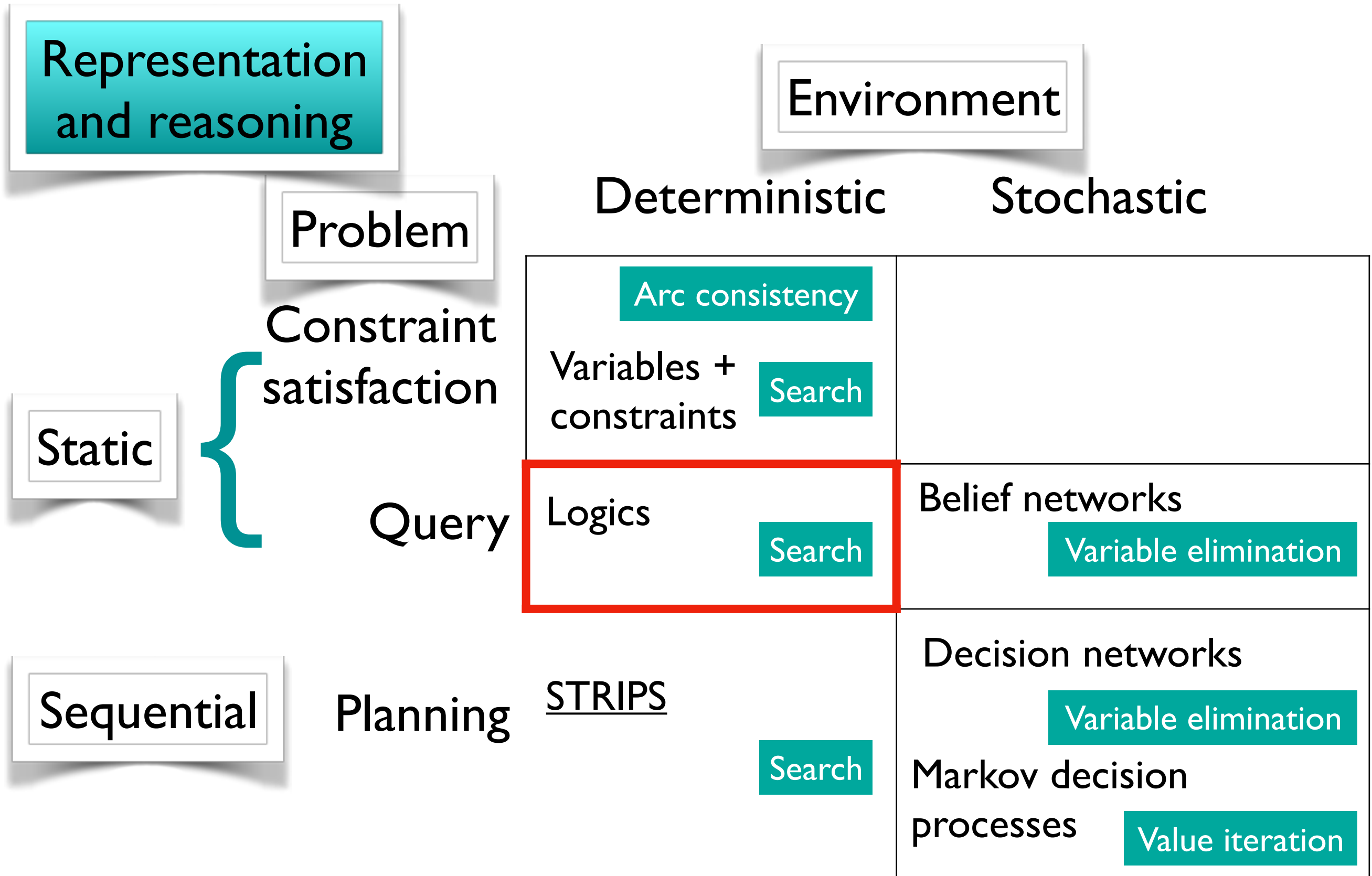


After you solve the CSP, what's the actual plan?

# Lecture outline

- Recap: Planning
- Logic introduction 
- Propositional logic: Syntax and semantics
- Propositional Definite Clause (PDC) Logic: Syntax
- Propositional Definite Clause (PDC) Logic: Semantics

# A rough CPSC 322 overview



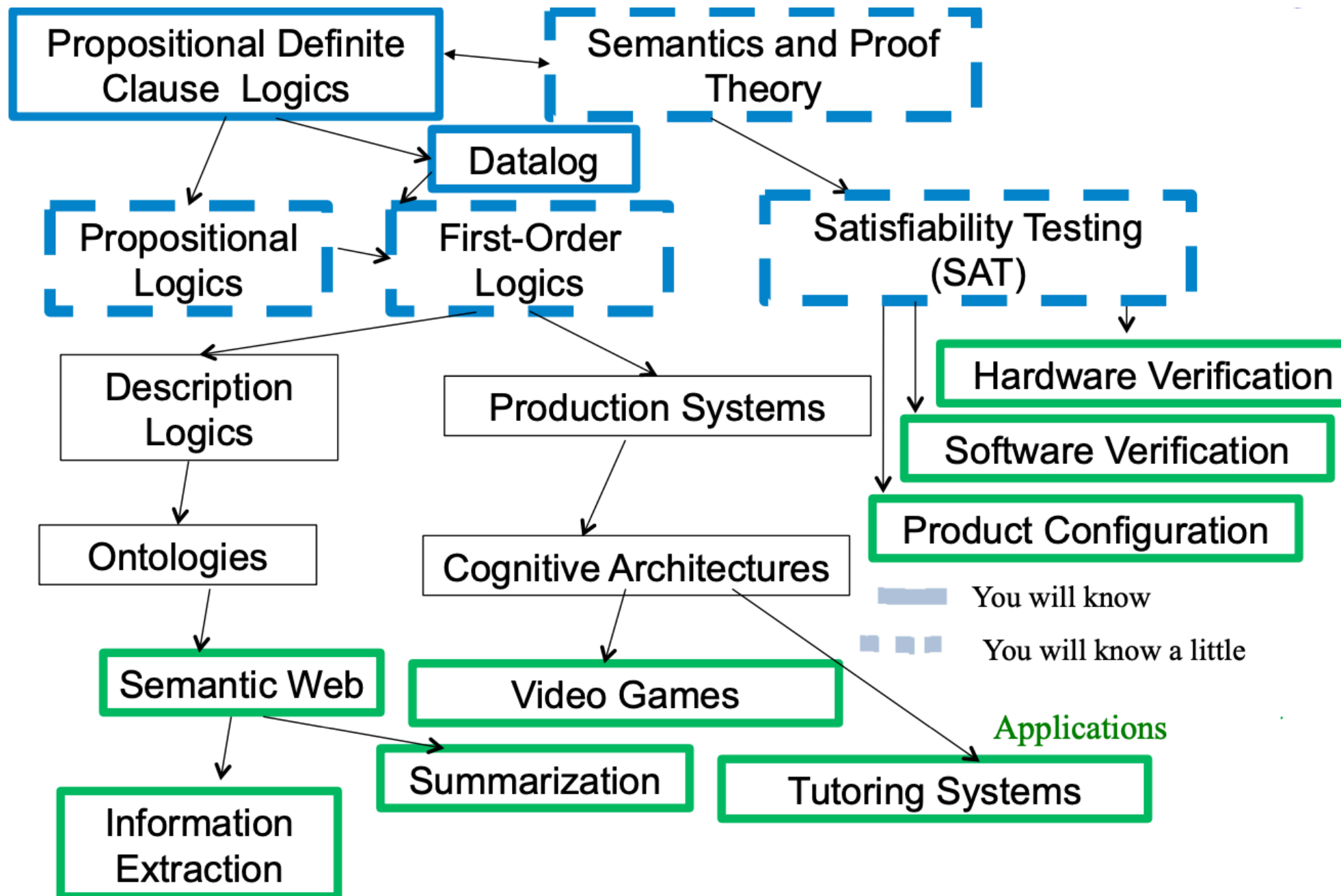


# Today: Learning outcomes

From this lecture, students are expected to be able to:

- Verify whether a logical statement belongs to the language of full propositional logics.
- Verify whether a logical statement belongs to the language of propositional definite clauses
- Verify whether an interpretation is a model of a propositional definite clause logic knowledge base

# Logics in AI



You know

Know little

Applications

# Common sense knowledge

If the **baby** does not thrive on raw **milk**, boil it.



**it = ?**



# Common sense knowledge

If the **baby** does not thrive on raw **milk**, boil it.



**it = ?**



**common-sense knowledge**

The boiling action is more common with milk.

# Why propositional logic?

- We'll mostly focus on propositional logic as this is the starting point for more complex ones ....
- Natural to express knowledge about the world
  - What is true (boolean variables)
  - How it works (logical formulas)
- Well understood formal properties
- Boolean nature can be exploited for efficiency

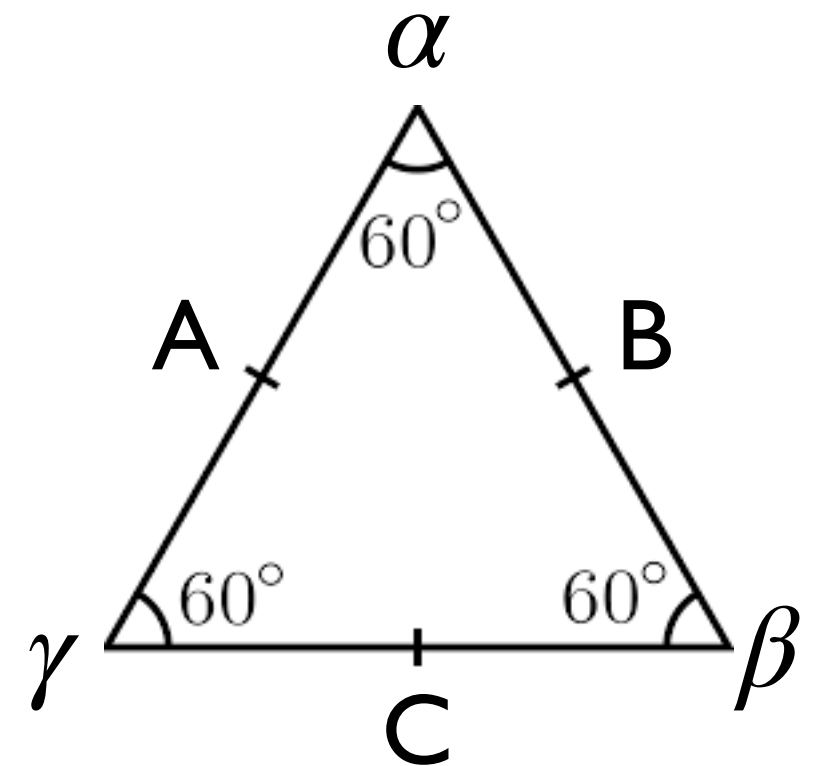
# What you already know about logic

Some logical operators from **programming**:

```
If ((amount > 0) and (amount < 1000)) || !(age < 30)
```

Logic is the language of **mathematics**.  
It is used to define formal structures  
(e.g., sets, graphs) and to prove  
statements about them. For example:

$$\forall x, \text{triangle}(x) \implies A = B = C \iff \alpha = \beta = \gamma$$



# Logic: a framework for representation and reasoning

We use logic as a Representation and Reasoning System that can be used to formalize a domain and to reason about it

- When we represent a domain about which we have only partial (but certain) information, we need to represent objects, properties, sets, groups, actions, events, time, space.
- All these can be represented as **objects** or **relationships between objects**
- **Logic** is the language to express knowledge about the world this way



# Example

“Natural” to express knowledge about the world  
(more natural than a “flat” set of variables & constraints)

“Every student who works diligently passes the course.”

$$\text{student}(s) \wedge \text{registered}(s, c) \wedge \text{course\_name}(c, 322) \\ \wedge \text{works\_diligently}(s) \implies \text{passes}(s, c)$$
$$\text{student}(sam), \text{registered}(sam, c_1), \text{course\_name}(c_1, 322), \\ \text{works\_diligently}(sam)$$

Query:  $\text{passes}(sam, c_1)$ ?



# Why logics?

## Compact representation

- Compared to, e.g., a CSP with a variable for each student
- It is easy to **incrementally add knowledge**
- It is easy to **check and debug knowledge**
- Provides language for **asking complex queries**
- Well understood **formal properties**

# Logic: a general framework for reasoning

- Let's think about how to represent a world about which we have only partial (but certain) information
- Our tool: **propositional logic**
- General problem:
  - tell the computer how the world works
  - tell the computer some facts about the world
  - ask a yes/no question about whether other facts must be true

# Representation and reasoning system (RRS)

## **Definition (RRS):**

A Representation and Reasoning System (RRS) consists of:

- **syntax**: specifies the symbols used, and how they can be combined to form legal sentences
- **semantics**: specifies the meaning of the symbols
- **reasoning theory or proof procedure**: a (possibly nondeterministic) specification of how an answer can be produced.

# Representation and reasoning system (RRS)

## **Definition (RRS):**


A Representation and Reasoning System (RRS) consists of:

- **syntax**: specifies the symbols used, and how they can be combined to form legal sentences
- **semantics**: specifies the meaning of the symbols
- **reasoning theory or proof procedure**: a (possibly nondeterministic) specification of how an answer can be produced.

We have seen several **representations** and **reasoning procedures**:

- **State space graph** + **search**
- **CSP** + **search/arc consistency**
- **STRIPS** + **search/arc consistency**

# Lecture outline

- Recap: Planning
- Logic introduction
- Propositional logic: Syntax and semantics 
- Propositional Definite Clause (PDC) Logic: Syntax
- Propositional Definite Clause (PDC) Logic: Semantics

# Propositional logic: syntax

## Definition (proposition):

Examples: It is sunny.

A **proposition** is a sentence, written in a language, that has a truth value (i.e., it is true or false) in a world. It can be atomic or compound.

## Definition (atom):

Examples: sunny,  $p_1$ , *live* <sub>$l_1$</sub>

An **atomic proposition** or **an atom** is a symbol. We use the convention that atomic propositions consist of letters, digits and the underscore ( $\_$ ) and start with a **lower-case** letter.

# Propositional logic: syntax

**Definition (formula):** Examples:  $(sunny \wedge happy) \vee (rain \wedge sad)$

A **proposition** or **logical formula** is either an atomic proposition or a compound proposition of the form shown below, where  $p$  and  $q$  are propositions and  $\neg, \wedge, \vee, \rightarrow, \leftarrow, \leftrightarrow$  are **logical connectives**.

$\neg p$	“not $p$ ”	negation of $p$
$p \wedge q$	“ $p$ and $q$ ”	conjunction of $p$ and $q$
$p \vee q$	“ $p$ or $q$ ”	disjunction of $p$ and $q$
$p \rightarrow q$	“ $p$ implies $q$ ”	implication of $q$ from $p$
$p \leftarrow q$	“ $p$ if $q$ ”	implication of $p$ from $q$
$p \leftrightarrow q$	“ $p$ if and only if $q$ ”	equivalence of $p$ and $q$

# In other words ...

Propositions are also called **sentences**.

- If  $S$  is a sentence,  $\neg S$  is also a sentence (negation)
- If  $S_1$  and  $S_2$  are sentences,  $S_1 \wedge S_2$  is a sentence (conjunction)
- If  $S_1$  and  $S_2$  are sentences,  $S_1 \vee S_2$  is a sentence (disjunction)
- If  $S_1$  and  $S_2$  are sentences,  $S_1 \rightarrow S_2$  is a sentence (implication)
- If  $S_1$  and  $S_2$  are sentences,  $S_1 \leftrightarrow S_2$  is a sentence (biconditional)



## ASIDE: A BNF grammar of sentences in propositional logic

$$\begin{aligned} \textit{Sentence} &\rightarrow \textit{AtomicSentence} \mid \textit{ComplexSentence} \\ \textit{AtomicSentence} &\rightarrow \textit{True} \mid \textit{False} \mid P \mid Q \mid R \mid \dots \\ \textit{ComplexSentence} &\rightarrow (\textit{Sentence}) \mid [\textit{Sentence}] \\ &\mid \neg \textit{Sentence} \\ &\mid \textit{Sentence} \wedge \textit{Sentence} \\ &\mid \textit{Sentence} \vee \textit{Sentence} \\ &\mid \textit{Sentence} \Rightarrow \textit{Sentence} \\ &\mid \textit{Sentence} \Leftrightarrow \textit{Sentence} \end{aligned}$$

OPERATOR PRECEDENCE :  $\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$

Source: Russel and Nerving book

# Propositional logic: semantics

Do any of these statements mean anything? Syntax doesn't answer this question!

**Semantics** allows you to relate the symbols in the logic to the domain you're trying to model.

# Propositional logic: semantics

## **Definition (Interpretation):**

An interpretation  $I$  assigns a truth value to each atom.

We can use the interpretation to determine the truth value of formulas. Note that truth values are only defined with respect to interpretations; propositions may have different truth values in different interpretations.

# Propositional logic: semantics



## **Definition (interpretation):**

An interpretation  $I$  assigns a truth value to each atom.

If our domain has 10 atoms, how many interpretations are there?

A.  $2^{10}$

B.  $10^2$

C.  $10 \times 2$


# Propositional logic: semantics



## **Definition (interpretation):**

An interpretation  $I$  assigns a truth value to each atom.

If our domain has 10 atoms, how many interpretations are there?

A.  $2^{10}$  

B.  $10^2$

C.  $10 \times 2$

# Propositional logic: semantics

## **Definition (interpretation):**

An interpretation  $I$  assigns a truth value to each atom.

If our domain has 10 atoms, how many interpretations are there?

A.  $2^{10}$

B.  $10^2$

C.  $10 \times 2$

Similar to possible worlds in CSP

# Propositional logic: semantics

Whether a compound proposition is true in an interpretation is inferred using the truth table

$p$	$q$	$\neg p$	$p \wedge q$	$p \vee q$	$p \leftarrow q$	$p \rightarrow q$	$p \leftrightarrow q$
true	true	false	true	true	true	true	true
true	false	false	false	true	true	false	false
false	true	true	false	true	false	true	false
false	false	true	false	false	true	true	true

# Interpretation example (pair-share)

Suppose we have three atoms: `ai_is_powerful`, `sky_is_orange`, `happy`. Suppose interpretation  $I_1$  assigns true to `ai_is_powerful`, false to `happy` and false to `sky_is_orange`

Which of the following are **true** in  $I_1$ ?

1. `ai_is_powerful`  $\wedge$  `happy`
2.  $\neg$ `happy`
3. `happy`  $\leftarrow$  `ai_is_powerful`
4.  $\neg$ `happy`  $\leftarrow$  `sky_is_orange`

$I_1$   
`ai_is_powerful` = true  
`happy` = false  
`sky_is_orange` = false



# Propositional logic in practice

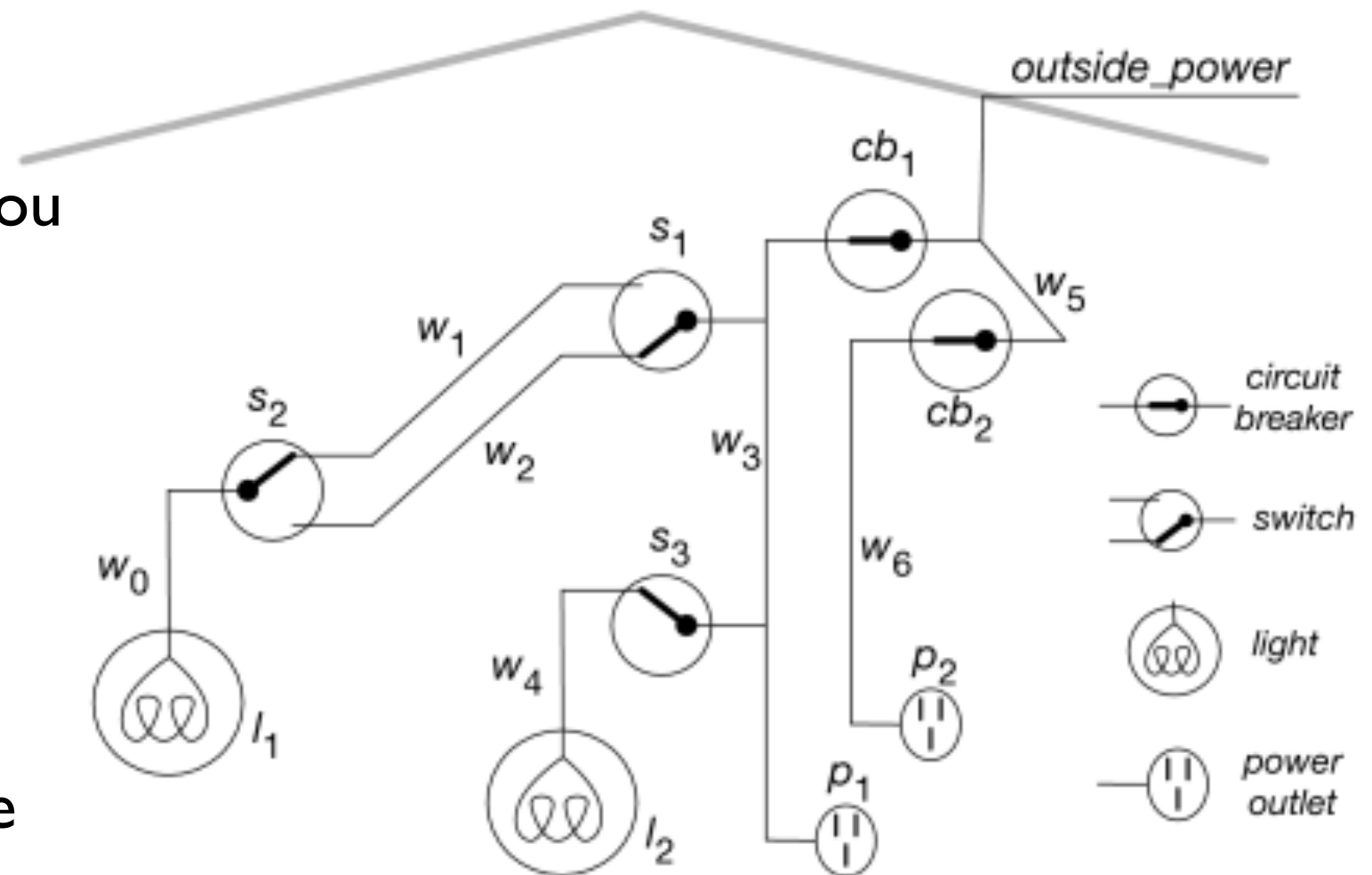
- Agent is told (perceives) some facts about the world (some propositions are true)
- Agent is told (already knows / learns) how the world works (logical formulas)
- Agent can answer **yes/no questions** about whether other **facts must be true**

# Using logics to make inferences


1. Begin with a task domain.
2. Distinguish those things you want to talk about (the ontology)
3. Choose symbols in the computer to denote propositions
4. Tell the system knowledge about the domain
5. Ask the system whether new statements about the domain are true or false

# Example: Electric environment

1. Begin with a task domain.
2. Distinguish those things you want to talk about (the ontology)
3. Choose symbols in the computer to denote propositions
4. Tell the system knowledge about the domain
5. Ask the system whether new statements about the domain are true or false



# Lecture outline

- Recap: Planning
- Logic introduction
- Propositional logic: Syntax and semantics
- Propositional Definite Clause (PDC) Logic: Syntax 
- Propositional Definite Clause (PDC) Logic: Semantics

# Propositional Definite Clause (PDC) Logic

Propositional Definite Clauses: our first **logical** representation and reasoning system. (Very simple!)

Only two kinds of statements:

- that a **proposition is true**
- that a **proposition is true if one or more other propositions are true**

Sublanguage of propositional logic that does not allow uncertainty or ambiguity.

# PDC logic

Limited because we are only considering propositions. But still useful because

- **Adequate in many domains** (with some adjustments)
- Reasoning steps **easy to follow by humans**
- **Inference linear in size** of your set of statements
- Similar formalisms used in **cognitive architectures**

# PDC logic: Syntax

## Definition (atom):

Examples:  $p_1$ ,  $live\_l_1$

An **atomic proposition** or an **atom** is a symbol.

## Definition (body):

Examples:  $p_1$ ,  $p_1 \wedge p_2$ ,  $ok\_w_1 \wedge live\_w_0$

A **body** is of the form  $a_1 \wedge \dots \wedge a_m$ , where  $a_1, \dots, a_m$  are atoms,  $m \geq 0$ .

## Definition (definite clause):

A **definite clause** is of the form  $h \leftarrow b$ , where  $h$  is an atom (“head”) and  $b$  is a body. (Read this as “ $h$  if  $b$ ”.)

Examples:  $p_1$ ,  $p_1 \leftarrow p_2$ ,  $live\_w_0 \leftarrow live\_w_1 \wedge up\_s_2$

# PDC logic: Syntax

## Definition (definite clause):

A definite clause is of the form  $h \leftarrow b$ , where  $h$  is an atom (“head”) and  $b$  is a body. (Read this as “ $h$  if  $b$ ”.)

Examples:  $p_1, p_1 \leftarrow p_2, live\_w_0 \leftarrow live\_w_1 \wedge up\_s_2$

In definite clause  $h \leftarrow b$  (i.e.,  $h \leftarrow a_1 \wedge \dots \wedge a_m$ ),

## Definition (rule):

If  $m > 0$ , the definite clause is called a **rule**.

Examples:  $p_1 \leftarrow p_2, live\_w_0 \leftarrow live\_w_1 \wedge up\_s_2$

## Definition (atomic clause):

Examples:  $p_1$

If  $m = 0$ , the arrow can be omitted and the clause is an **atomic clause** or a **fact**, with an empty body.



# Identify definite clauses (pair-share)

Which of the following are legal definite clauses?

1. `ai_is_powerful`

2. `¬ai_is_powerful`

3. `ai_is_fun ← learn_useful_techniques ∧ ¬tooMuch_work`

4. `sam_is_in_room ∧ night_time ← switch_1_is_up`

5. `switch_1_is_up ← sam_is_in_room ∧ night_time`

6. `happy ∨ sad ∨ ¬alive`

7. `srtsyj ← errt ∧ gffdgddgd`

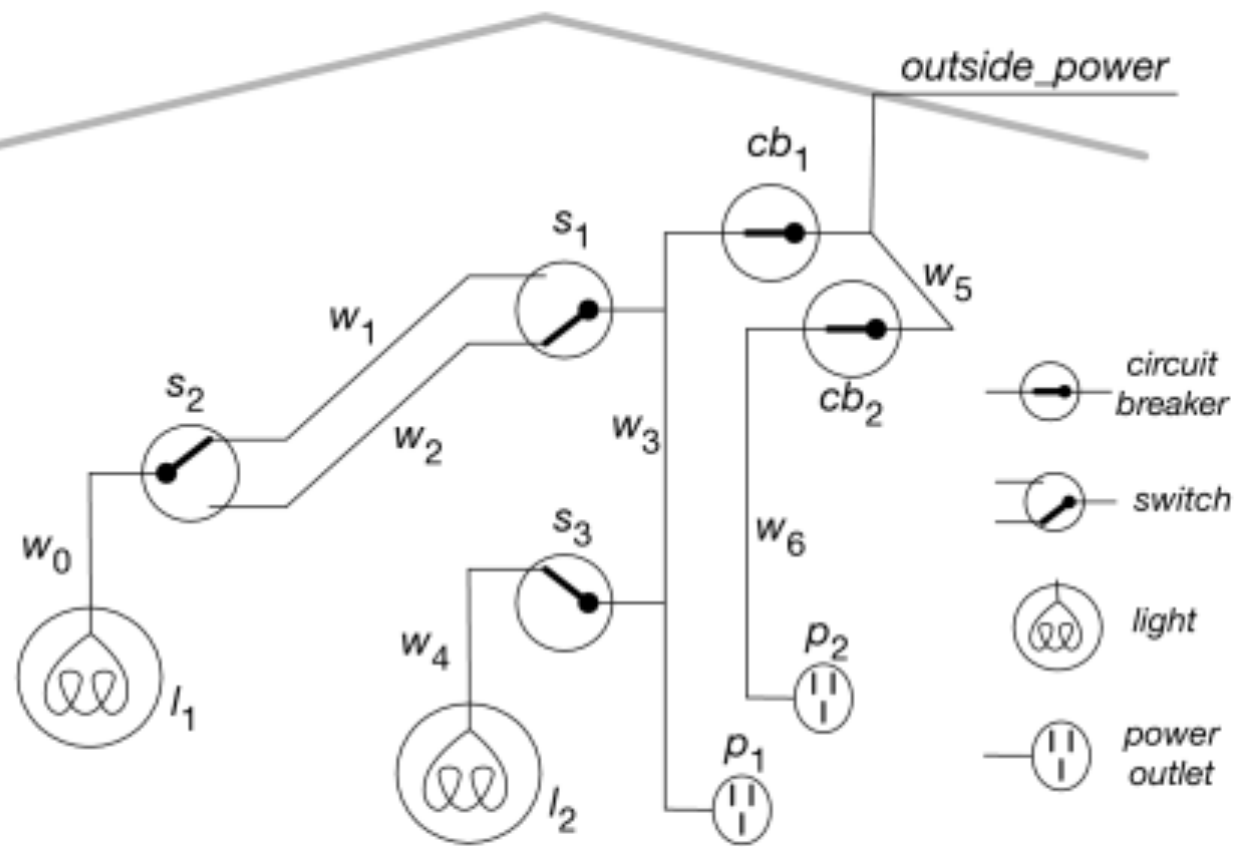
# PDC logic semantics: Knowledge Base (KB)

**Definition:** A **knowledge base** (KB) is a set of definite clauses.

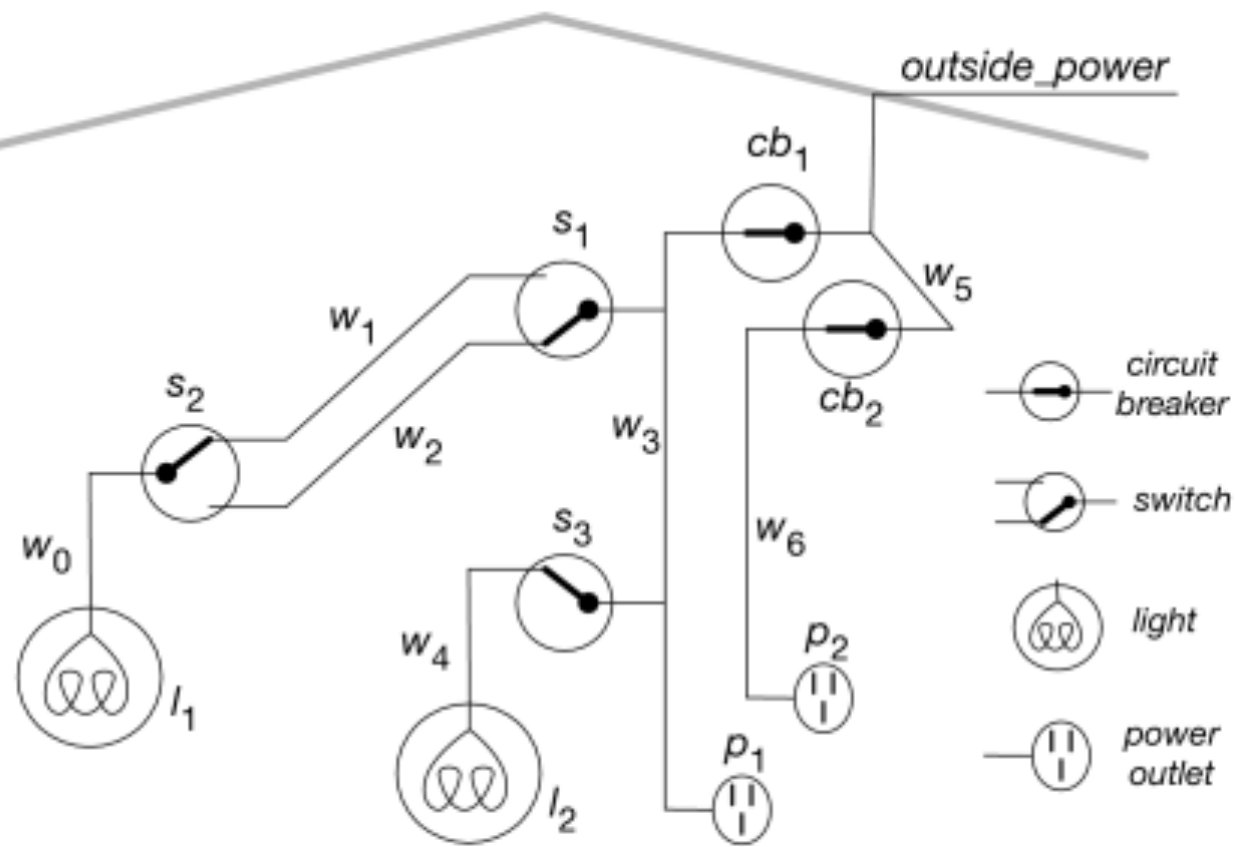
Example:

$$\{ p_2, p_3, p_4, p_1 \leftarrow p_2 \wedge p_3 \wedge p_4, \textit{live\_}l_1 \}$$

# Example: Electric environment



# Example: Electric environment



*light\_l1.*

*light\_l2.*

*ok\_l1.*

*ok\_l2.*

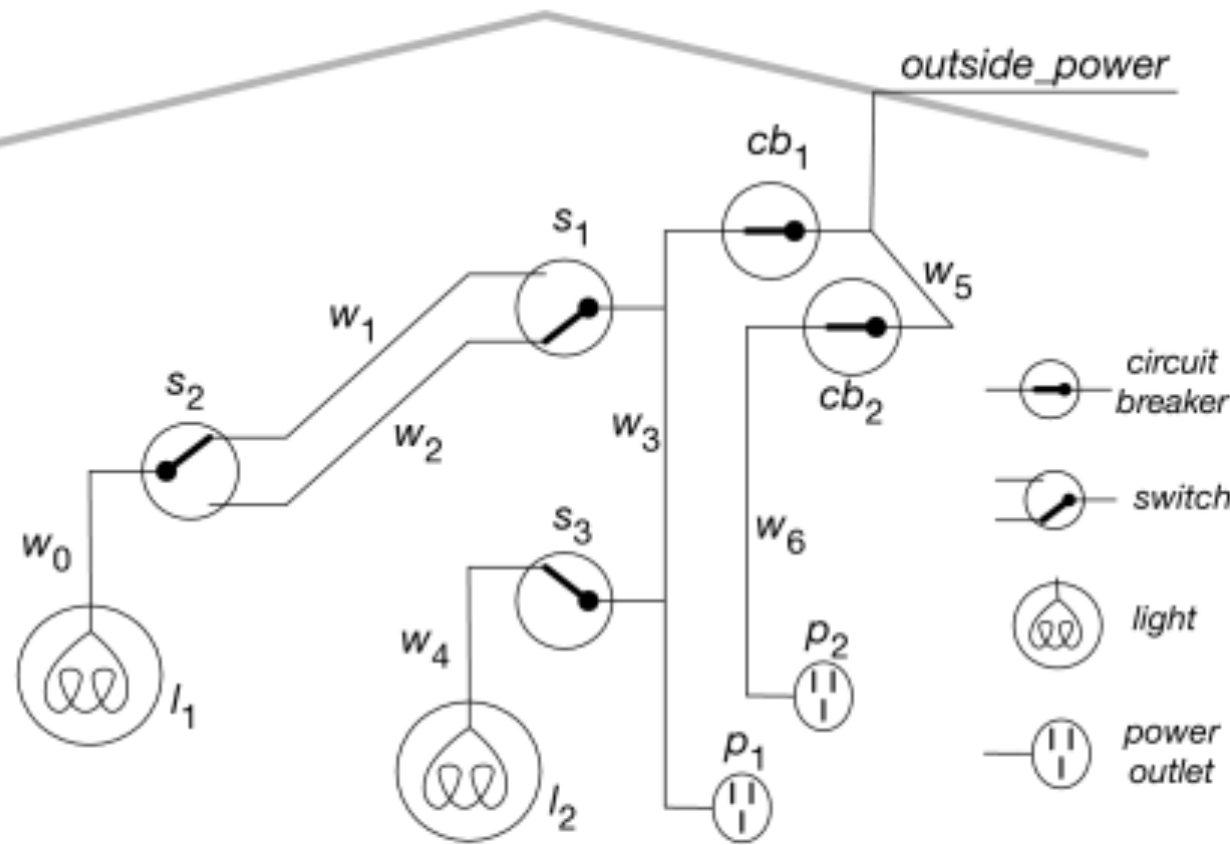
*ok\_cb1.*

*ok\_cb2.*

*live\_outside.*

atoms

# Example: Electric environment



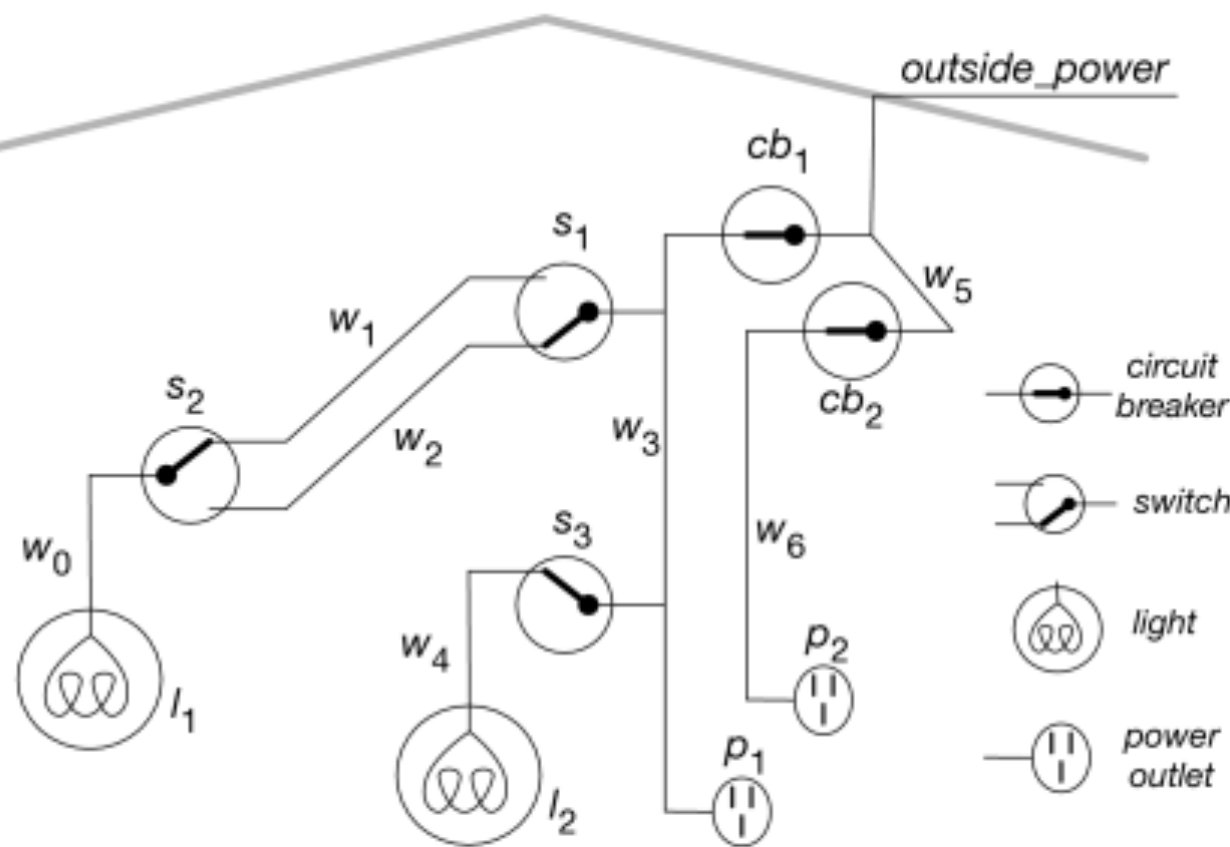
*light\_l1.*  
*light\_l2.*  
*ok\_l1.*  
*ok\_l2.*  
*ok\_cb1.*  
*ok\_cb2.*  
*live\_outside.*

atoms

*live\_l1*  $\leftarrow$  *live\_w0*.  
*live\_w0*  $\leftarrow$  *live\_w1*  $\wedge$  *up\_s2*.  
*live\_w0*  $\leftarrow$  *live\_w2*  $\wedge$  *down\_s2*.  
*live\_w1*  $\leftarrow$  *live\_w3*  $\wedge$  *up\_s1*.  
*live\_w2*  $\leftarrow$  *live\_w3*  $\wedge$  *down\_s1*.  
*live\_l2*  $\leftarrow$  *live\_w4*.  
*live\_w4*  $\leftarrow$  *live\_w3*  $\wedge$  *up\_s3*.  
*live\_p1*  $\leftarrow$  *live\_w3*.  
*live\_w3*  $\leftarrow$  *live\_w5*  $\wedge$  *ok\_cb1*.  
*live\_p2*  $\leftarrow$  *live\_w6*.  
*live\_w6*  $\leftarrow$  *live\_w5*  $\wedge$  *ok\_cb2*.  
*live\_w5*  $\leftarrow$  *live\_outside*.  
*lit\_l1*  $\leftarrow$  *light\_l1*  $\wedge$  *live\_l1*  $\wedge$  *ok\_l1*.  
*lit\_l2*  $\leftarrow$  *light\_l2*  $\wedge$  *live\_l2*  $\wedge$  *ok\_l2*.

rules

# Example: Electric environment



*light\_l1.*  
*light\_l2.*  
*ok\_l1.*  
*ok\_l2.*  
*ok\_cb1.*  
*ok\_cb2.*  
*live\_outside.*

atoms

+

*live\_l1*  $\leftarrow$  *live\_w0*.  
*live\_w0*  $\leftarrow$  *live\_w1*  $\wedge$  *up\_s2*.  
*live\_w0*  $\leftarrow$  *live\_w2*  $\wedge$  *down\_s2*.  
*live\_w1*  $\leftarrow$  *live\_w3*  $\wedge$  *up\_s1*.  
*live\_w2*  $\leftarrow$  *live\_w3*  $\wedge$  *down\_s1*.  
*live\_l2*  $\leftarrow$  *live\_w4*.  
*live\_w4*  $\leftarrow$  *live\_w3*  $\wedge$  *up\_s3*.  
*live\_p1*  $\leftarrow$  *live\_w3*.  
*live\_w3*  $\leftarrow$  *live\_w5*  $\wedge$  *ok\_cb1*.  
*live\_p2*  $\leftarrow$  *live\_w6*.  
*live\_w6*  $\leftarrow$  *live\_w5*  $\wedge$  *ok\_cb2*.  
*live\_w5*  $\leftarrow$  *live\_outside*.  
*lit\_l1*  $\leftarrow$  *light\_l1*  $\wedge$  *live\_l1*  $\wedge$  *ok\_l1*.  
*lit\_l2*  $\leftarrow$  *light\_l2*  $\wedge$  *live\_l2*  $\wedge$  *ok\_l2*.

rules

= KB:  
 Set of  
 definite  
 clauses

# PDC logic: Semantics

## **Definition (interpretation):**

An interpretation  $I$  assigns a truth value to each atom.

We can use the interpretation to determine the truth value of clauses.

# PDC logic: Semantics

## Definition (truth values of statements):

A body  $b_1 \dots b_m$  is true in interpretation  $I$  if and only if  $\forall b_i, 0 \leq i \leq m, b_i$  is true in  $I$ .

	p	q	r	s	$p \wedge r$	$p \wedge r \wedge s$
$I_1$	T	T	T	T		
$I_2$	F	F	F	F		
$I_3$	T	T	F	F		
$I_4$	T	T	T	F		
$I_5$	T	T	F	T		



# PDC logic: Semantics

## Definition (truth values of statements):

A rule  $h \leftarrow b$  is false in  $I$  if and only if  $b$  is true in  $I$  and  $h$  is false in  $I$ .

	p	q	r	s	$p \leftarrow r$	$s \leftarrow q \wedge r$
$I_1$	T	T	T	T		
$I_2$	F	F	F	F		
$I_3$	T	T	F	F		
$I_4$	T	T	T	F		
$I_5$	T	T	F	T		

In other words: “if  $b$  **is true** I am claiming that  $h$  **must be true**, otherwise I am not making any claim”.

# PDC logic semantics: Knowledge Base (KB)

i-clicker.

**Definition:** A **knowledge base KB** is true in  $I$  if and only if every clause in KB is true in  $I$ .

Which of the following KBs below are true in  $K$ ?

	$p$	$q$	$r$	$s$
$K$	T	T	F	F

A.

$p$
$r$
$s \leftarrow p \wedge q$

B.

$p$
$r$
$s \leftarrow p$

C.

$p$
$q \leftarrow r \wedge s$

# PDC logic semantics: Knowledge Base (KB)

i-clicker.

**Definition:** A knowledge base **KB** is true in  $I$  if and only if every clause in KB is true in  $I$ .

Which of the following KBs below are true in  $K$ ?

	$p$	$q$	$r$	$s$
$K$	T	T	F	F

A.

$p$
$r$
$s \leftarrow p \wedge q$

B.

$p$
$r$
$s \leftarrow p$

C.

$p$
$q \leftarrow r \wedge s$



# PDC logic semantics: Models

**Definition:** A **model** of a set of clauses (KB) is an interpretation in which all the clauses are true.

Which interpretations are models?

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	p	q	r	s	model of KB?
I <sub>1</sub>	T	T	T	T	
I <sub>2</sub>	F	F	F	F	
I <sub>3</sub>	T	T	F	F	
I <sub>4</sub>	T	T	T	F	
I <sub>5</sub>	T	T	F	T	

# PDC logic semantics: Models

**Definition:** A **model** of a set of clauses (KB) is an interpretation in which all the clauses are true.

Which interpretations are models?

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	p	q	r	s	$p \leftarrow q$	$r \leftarrow s$	KB
I <sub>1</sub>	T	T	T	T			
I <sub>2</sub>	F	F	F	F			
I <sub>3</sub>	T	T	F	F			
I <sub>4</sub>	T	T	T	F			
I <sub>5</sub>	T	T	F	T			

# PDC logic semantics: Entailment

Given a knowledge base KB, can we **infer** implicit information (i.e., new clauses)?

**Definition:** If KB is a knowledge base and  $g$  is a proposition,  $g$  is a **logical consequence** of KB, written as  $k \models g$ , if  $g$  is true in every model of KB.

We also say that  $g$  **logically follows** from KB, or that KB **entails**  $g$ . In other words,  $k \models g$  if there is no interpretation in which KB is true and  $g$  is false.

# PDC semantics: Entailment

i-clicker.

**Definition:** If KB is a knowledge base and  $g$  is a proposition,  $g$  is a **logical consequence** of KB, written as  $KB \models g$ , if  $g$  is true in every model of KB.

Given the knowledge base KB, which of the following are true?

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

A.  $KB \models p$

C.  $KB \models r$

B.  $KB \models q$

D.  $KB \models s$


# PDC semantics: Entailment

i-clicker.


**Definition:** If KB is a knowledge base and  $g$  is a proposition,  $g$  is a **logical consequence** of KB, written as  $KB \models g$ , if  $g$  is true in every model of KB.

Given the knowledge base KB, which of the following are true?

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

A.  $KB \models p$  

C.  $KB \models r$

B.  $KB \models q$  

D.  $KB \models s$



# Revisit: Learning outcomes

From this lecture, students are expected to be able to:

- Verify whether a logical statement belongs to the language of full propositional logics.
- Verify whether a logical statement belongs to the language of propositional definite clauses
- Verify whether an interpretation is a model of a propositional definite clause logic knowledge base

# Coming up

We'll start using all these definitions for automated proofs!

## 5.3.2 Proofs

