Extending and Evaluating a Control Flow Obfuscation Technique for JVM Applications Utilizing `invokedynamic` with Native Bootstrapping

Bachelor's Thesis Proposal

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1 Introduction

Obfuscation techniques are an important tool to make malicious reverse engineering attempts of software harder and more time-consuming. These techniques are often employed to protect trade secrets, intellectual property, and program integrity.

Even higher importance is placed on the usage of obfuscation techniques when the software is deployed in an environment controlled by a trusted user of the software (e.g. to a personal computer as opposed to a microcontroller for example) for the prevention of so-called man-at-the-end (MATE) attacks as coined by (Falcarin et al. 2011).

Programming languages compiled to high-level bytecode—such as Java Virtual Machine (JVM) bytecode instead of machine code—are especially vulnerable to MATE attacks, as JVM bytecode is more structured and thus allows easier modification when compared to machine code.

A variety of obfuscation techniques exist, some of which, such as symbol scrambling (Chan and Yang 2004), remove information in an irreversible way. Other aspects, like the interprocedural control flow graph of a program, are harder to obfuscate. Keeping this information intact allows an attacker to gain great insight into the architecture and inner workings of an application.

Java 7 saw the introduction of the `invokedynamic` instruction which allows the invocation of methods without directly specifying the target call site inside the bytecode, instead delegating execution to a so-called bootstrap method in order to determine the call site at runtime rather than compile-time. The instruction enables dynamic dispatch behavior similar to reflection but in a more performant way because the call site can be accessed directly after its initial extraction from the bootstrap method as illustrated by (Ivanov 2015). These properties have proven useful for call graph obfuscation efforts.

A novel obfuscation technique proposed by (Wood and Azim 2021) utilizing `invokedynamic` implements the bootstrap method inside a native library and replaces regular `invoke*` instructions with `invokedynamic` ones, hindering static analysis attempts. Outsourcing the bootstrap method implementation to native code increases the

level of obfuscation and enables the application of further obfuscation methods going beyond techniques available for the JVM. It also prevents analysis of the bootstrap method bytecode with the help of established tools in the JVM ecosystem.

The proposed bachelor's thesis would primarily focus on this obfuscation technique in an aim to create an implementation of it since none is openly available at the moment. Once a basic implementation is finished, it will be extended further and evaluated.

2 Problem

While a proof-of-concept has been developed for the novel obfuscation technique proposed by (Wood and Azim 2021), no implementation of it has been released which hinders further benchmarking attempts that are needed in order to judge its performance impact in a larger context: what are the performance implications when applying this technique to an input program, especially across a broader variety of JVM distributions and JIT compilers whose optimization behavior regarding `invokedynamic` might differ?

In addition to the lack of an implementation, the developed technique does not yet fully exhaust its potential and can be further extended beyond obfuscation of the call graph itself: field access instructions (`getfield`, `putfield`, `getstatic`, and `putstatic`) may be included in the obfuscation effort by replacing them with invocations to synthetic accessor methods which will partake in the subsequent main obfuscation step.
How large is the difference in performance when this extension is introduced compared to the performance of the obfuscation technique without it?

The need for an additional native library containing the bootstrap method implementation along with the generation of new methods when considering the extension for field accesses will reflect in the total size of the artifacts produced by the obfuscation process.
This increase in application size may be more or less significant depending on the type of deployment: containerized and/or user-facing software — the latter being the main use case for obfuscation — is preferably kept as small as possible, so the impact of the obfuscation process on application size should be measured by comparing the size of the original input to the size of the obfuscated application (considering both the exclusion and inclusion of field accesses in the obfuscation process). To what extend does the obfuscation lead to an increase in application size?
Some executable formats used by virtual machines, such as the DEX format of the Dalvik VM built into Android (Google LLC, n.d.), are limited regarding the amount of methods an application may contain, so how many additional synthetic accessors can we expect when including the obfuscating of field access instructions?

Another area that is only briefly explored by the original paper but which should also be evaluated is the effectiveness of the technique. How can one bypass the obfuscation technique using dynamic analysis and what is the complexity of it? To a lesser extent: how easy is it for an attacker to identify the used obfuscation technique and to understand what kinds of transformations have been applied to the code?

Answers to all of the above research questions will be used to draw a conclusion regarding the quality of the proposed obfuscation technique.

3 Solution Approach

3.1 Implementation

The primary goal of the bachelor's thesis will lay on the implementation of an open-source tool capable of performing the proposed obfuscation on single classes as well as entire JAR files using the ASM bytecode engineering library. (Bruneton 2011)

Figure 1 gives an overview of the three-step obfuscation process. Two inputs will be required by the obfuscation tool: a JAR or class file to be obfuscated and a template for the bootstrap method which is populated using the symbol table built in step 2.

Step 1 is optional and responsible for the inclusion of field access instructions in the obfuscation effort to increase the provided level of indirection, serving as an extension to the technique proposed by (Wood and Azim 2021). Obfuscation of field accesses is achieved through the generation of synthetic getter and setter methods which are invoked in place of the original instruction. No further steps are necessary when performing this bytecode transformation as the first step before the main obfuscation pass, as the generated `invoke*` instructions will be treated like any other method invocation present in the original bytecode and replaced by an equivalent `invokedynamic` instruction. Introducing the additional synthetic methods will lead to an inflation in bytecode size which should be measured by comparing the size of the input file with that of the obfuscated output file.

Following the optional pre-processing of field instructions in the first step, obfuscation **step 2** is responsible for the replacement of `invoke*` instructions with `invokedynamic` ones, making up the largest amount of bytecode transformation operations. A unique name is assigned to each encountered method invocation instruction and the association is stored in the symbol table. The unique name is present in the bytecode of the `invokedynamic` instruction and is used by the bootstrap method implementation to return the correct `CallSite` that was associated with it in the original bytecode. The symbol table built from this obfuscation step is passed to the input bootstrap method template using a template engine for the generation of the bootstrap method source code. Automatic compilation of the processed template is outside the scope of the tool due to reasons outlined in section 5.

Step 3 is the last obfuscation step and adds the bootstrap method definition to the main application class, along with code responsible for loading the native library containing the bootstrap method implementation. The library loading code is prepended to `<clinit>` to ensure that it is invoked before any other application code which might be obfuscated (and thus assumes the bootstrap method implementation is already available).

3.2 Evaluation

Open research questions regarding the performance implications of the obfuscation technique on input JAR files will be approached through the use of macrobenchmarks, either by using an established benchmark suite such as DaCapo (Blackburn et al. 2006) or by obfuscating an increasing amount of well-covered methods from the test suite of a set of open-source libraries as done by (Pizzolotto and Ceccato 2018).
Care must be taken when selecting the input JAR files to use for the benchmark, as they and their dependencies must be compiled for newer Java versions. Reasons for this are outlined in section 5. While no concrete libraries of applications have been chosen for

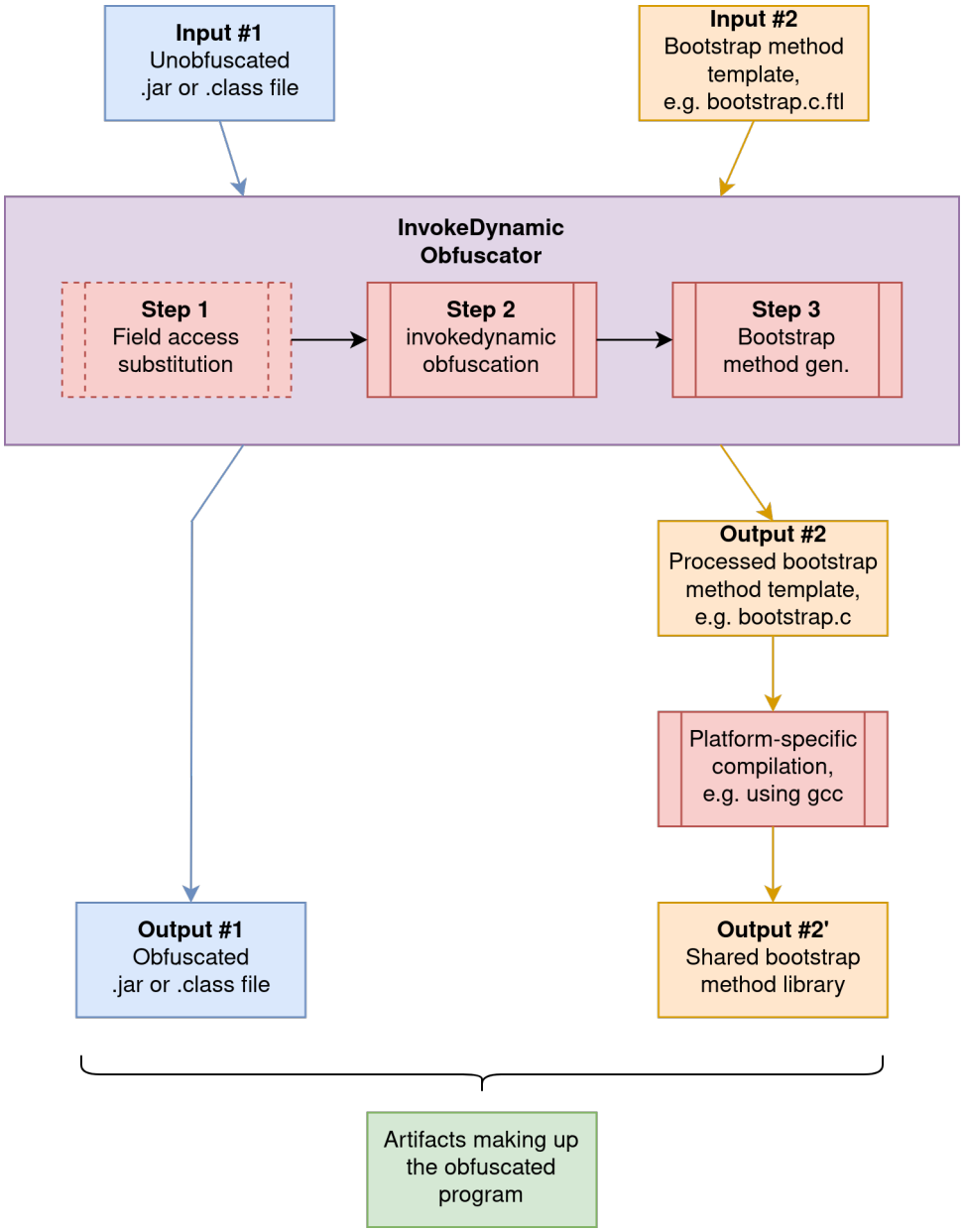

Figure 1: High-level overview of the obfuscation process.

their usage in benchmarks yet, reusing ones picked for benchmarking purposes by other obfuscation techniques might allow some level of comparability.

The ease of identification will be judged by comparing the original bytecode with its obfuscated version. To determine the complexity of an attack on the technique, an attempt at bypassing it will be made using dynamic analysis. As all applications obfuscated using `invokedynamic`-based techniques will have to eventually instantiate one of the three implementation of the `CallSite` interface to return it from the bootstrap method, it seems natural to focus on this area in search for a way to record the call sites which are accessed, potentially circumventing the concept of the native bootstrap method.

4 Related Work

An attempt at hiding the interprocedural control flow graph of JVM applications has been made by (Pizzolotto and Ceccato 2018) through the translation of selected portions of program bytecode into native libraries but introduced a performance penalty between 5% and 21% by doing so, as the method relies heavily on reflective calls which introduce a certain overhead.

(Fukuda and Tamada 2014) proposes an obfuscation technique intended to confuse debugging tools through the addition or removal of arguments from `MethodHandle` instances within the bootstrap method.

Various implementations of obfuscation techniques making use of the substitution of `invoke*` instructions with `invokedynamic` exist, both open-source (superblaubeere27 2018) and commercial (Zelix Pty Ltd. 1997), but none of them implement the bootstrap method in native code. No associated research papers seem to exist for either of the two projects.

5 Objectives and Limitations

For simplicity reasons, the obfuscation tool will likely be limited to support JVM bytecode emitted by newer versions of Java compilers (class file major version ≥ 50), as bytecode targeting older versions of the class file format does not necessarily contain stack map frames. While these class files can still be correctly loaded and verified by JVMs, the missing stack map frames pose an issue when `invokedynamic` instructions are introduced, as that requires raising the bytecode major version to 51 or higher, making stack map frames mandatory. ASM is capable of computing the stack map frames but requires loading the classes using `Class.forName` in certain situations, which in turn would require loading the input JAR or class file and any potential dependencies using a custom class loader, complicating the implementation.

A different area that is out of scope for the obfuscation tool is covering the entire functionality of the toolchain required for the compilation of the native library containing a bootstrap method implementation. This includes the invocation of a compiler from within the tool. Instead, the tool should provide templating capabilities to allow for the production of valid source files in a user-chosen language by combining a template file with information obtained in the obfuscation pass. The last necessary compilation step would be left to the user of the obfuscation tool.

6 Preliminary Outline

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5. Related Work
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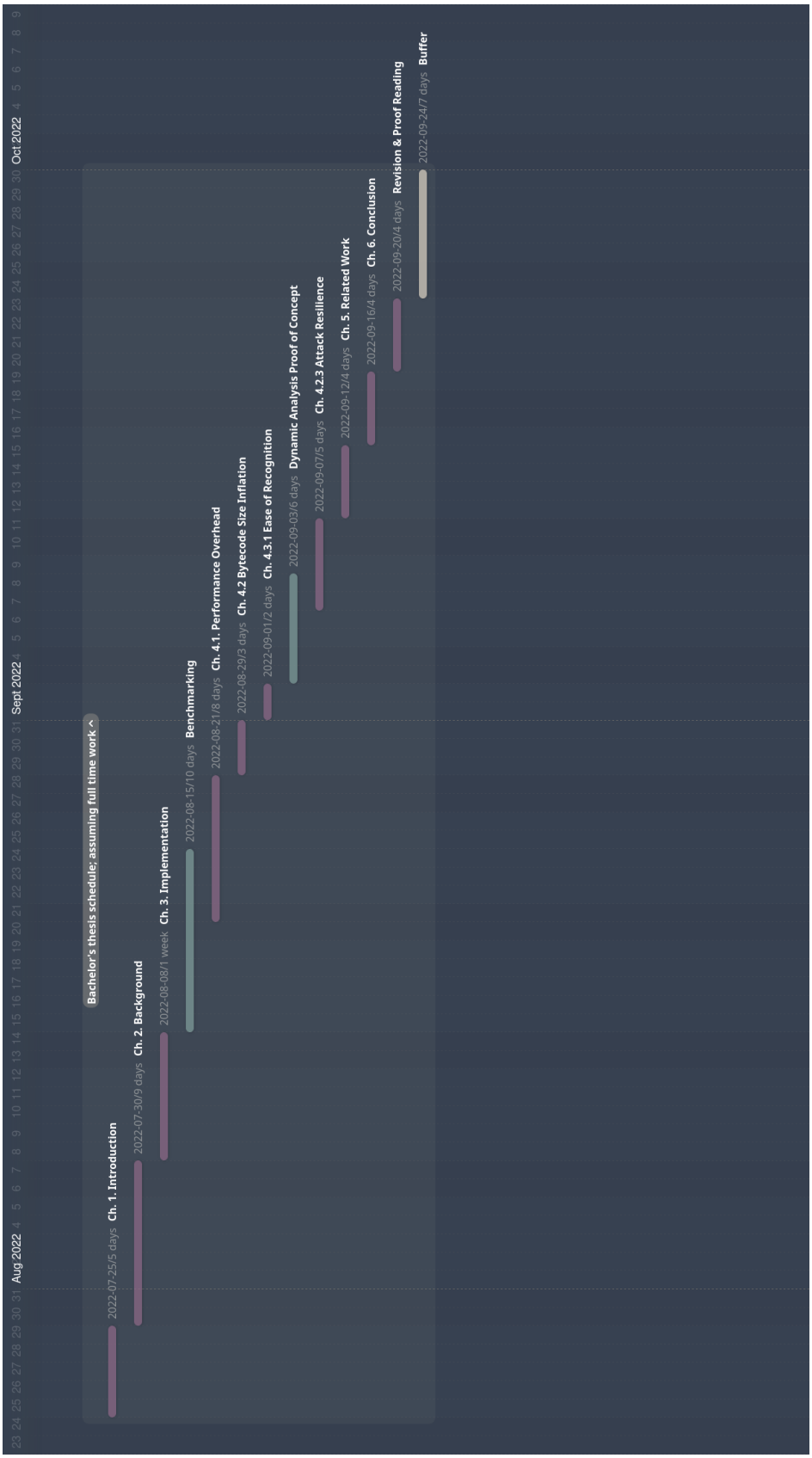
7 Schedule and Risk Analysis



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