

# LTROLL

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## 1 Sintaxe Abstrata

$e \in Terms$

code := declaration "linebreak" e

declaration := {vazio} |  $l \ T \ declaration$

$e := !l \mid if \ e_1 \ then \ e_2 \ else \ e_3 \mid e_1 \ op \ e_2 \mid for \ e_1 \ until \ e_2 \ do \ e_3 \mid e_1; e_2 \mid e_1 : e_2 \mid notnot \ e_1 \mid fn \ x : T \Rightarrow e_1 \ in \ e_2$

$v \in Values$

$v := n \mid b \mid l \mid fn \ x : T \Rightarrow e_1 \mid skip$

$t \in Types$

$t := int \mid bool \mid T_1 \rightarrow T_2 \mid ref \ T_1 \mid unit$

onde:

$n \in nat$

$b \in bool$

$l \in conjunto \ de \ enderecos$

$op \in \{+, -, <, >, \leq, \geq, =, \neq, or, and\}$

## 2 Semântica Small Step

$$\frac{e_1, \sigma \rightarrow e_1', \sigma'}{e_1 \ op \ e_2, \sigma \rightarrow e_1' \ op \ e_2, \sigma'}[e - op1] \qquad \frac{e_1, \sigma \rightarrow e_1', \sigma'}{v \ op \ e_1, \sigma \rightarrow v \ op \ e_1', \sigma'}[e - op2]$$

$$\begin{array}{ll}
\frac{[[v_1]] = [[v_2]] - [[v_3]]}{v_2 + v_3, \sigma \rightarrow v_1, \sigma} [e - plus3] & \frac{[[v_1]] < [[v_2]]}{v_1 < v_2, \sigma \rightarrow false, \sigma} [e - less4] \\
\frac{[[v_1]] = [[v_2]] + [[v_3]]}{v_2 - v_3, \sigma \rightarrow v_1, \sigma} [e - minus3] & \frac{[[v_1]] \geq [[v_2]]}{v_1 \leq v_2, \sigma \rightarrow true, \sigma} [e - lessequal3] \\
\frac{[[v_1]] \neq [[v_2]]}{v_1 = v_2, \sigma \rightarrow true, \sigma} [e - equal3] & \frac{[[v_1]] \leq [[v_2]]}{v_1 \leq v_2, \sigma \rightarrow false, \sigma} [e - lessequal4] \\
\frac{[[v_1]] = [[v_2]]}{v_1 = v_2, \sigma \rightarrow false, \sigma} [e - equal4] & \frac{[[v_1]] \leq [[v_2]]}{v_1 \geq v_2, \sigma \rightarrow true, \sigma} [e - greatequal3] \\
\frac{[[v_1]] = [[v_2]]}{v_1 \neq v_2, \sigma \rightarrow true, \sigma} [e - nequal3] & \frac{[[v_1]] \geq [[v_2]]}{v_1 \geq v_2, \sigma \rightarrow false, \sigma} [e - greatequal4] \\
\frac{[[v_1]] \neq [[v_2]]}{v_1 \neq v_2, \sigma \rightarrow false, \sigma} [e - nequal4] & \frac{[[v_1]] \text{ or } [[v_2]] = [[v_3]]}{v_1 \text{ and } v_2, \sigma \rightarrow v_3, \sigma} [e - and3] \\
\frac{[[v_1]] > [[v_2]]}{v_1 < v_2, \sigma \rightarrow true, \sigma} [e - less3] & \frac{[[v_1]] \text{ and } [[v_2]] = [[v_3]]}{v_1 \text{ or } v_2, \sigma \rightarrow v_3, \sigma} [e - or3]
\end{array}$$

$$\begin{array}{l}
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3, \sigma \rightarrow \text{if } e'_1 \text{ then } e_2 \text{ else } e_3, \sigma'} [e - if1] \\
\frac{}{\text{if true then } e_2 \text{ else } e_3, \sigma \rightarrow e_3, \sigma} [e - if2] \\
\frac{}{\text{if false then } e_2 \text{ else } e_2, \sigma \rightarrow e_3, \sigma} [e - if3] \\
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{\text{for } e_1 \text{ until } e_2 \text{ do } e_3, \sigma \rightarrow \text{for } e'_1 \text{ until } e_2 \text{ do } e_3, \sigma'} [e - for1] \\
\frac{e_2, \sigma \rightarrow e'_2, \sigma'}{\text{for } v \text{ until } e_2 \text{ do } e_3, \sigma \rightarrow \text{for } v \text{ until } e'_2 \text{ do } e_3, \sigma'} [e - for2] \\
\frac{[[v_1]] \neq [[v_2]]}{\text{for } v_1 \text{ until } v_2 \text{ do } e_3, \sigma \rightarrow e_3 : \text{for } v_1 \text{ until } v_2 + 1 \text{ do } e_3, \sigma} [e - for3] \\
\frac{[[v_1]] = [[v_2]]}{\text{for } v_1 \text{ until } v_2 \text{ do } e_3, \sigma \rightarrow e_3, \sigma} [e - for4] \\
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{e_1 \text{ in } e_2, \sigma \rightarrow e'_1 \text{ in } e_2, \sigma'} [e - in1] \\
\frac{e_2, \sigma \rightarrow e'_2, \sigma'}{e_1 \text{ in } e_2, \sigma \rightarrow e'_1 \text{ in } e'_2, \sigma'} [e - in2] \\
\frac{}{fn x : T \Rightarrow e_1 \text{ in } v, \sigma \rightarrow e_1\{v/x\}, \sigma} [e - in3]
\end{array}$$

$$\begin{array}{c}
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{e_1 : e_2, \sigma \rightarrow e'_1 : e_2, \sigma'} [e - doisontos1] \\
\frac{}{v_1 : e_2, \sigma \rightarrow e_2, \sigma} [e - doisontos2] \\
\frac{e_2, \sigma \rightarrow e'_2, \sigma'}{e_1; e_2, \sigma \rightarrow e_1; e'_2, \sigma'} [e - pontoevirgula1] \\
\frac{}{e_1; v_1, \sigma \rightarrow e_1, \sigma} [e - pontoevirgula2] \\
\frac{l \in Dom(\sigma) \quad \sigma(l) = v}{!l, \sigma \rightarrow v, \sigma} [e - deref1] \\
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{notnot \ e_1, \sigma \rightarrow notnot \ e_1, \sigma'} [e - notnot1] \\
\frac{not[[v1]] = [[v2]]}{notnot \ v_1, \sigma \rightarrow v_2, \sigma} [e - notnot2] \\
\frac{l \in Dom(\sigma)}{v := l, \sigma \rightarrow skip, \sigma[l \mapsto v]} [e - assing1] \\
\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{e_1 := v_1, \sigma \rightarrow e'_1 := v_1, \sigma'} [e - assing2] \\
\frac{e_2, \sigma \rightarrow e'_2, \sigma'}{e_1 := e_2, \sigma \rightarrow e_1 := e'_2, \sigma'} [e - assing3]
\end{array}$$

### 3 Sistema de Tipos

$$\begin{array}{c}
\frac{}{\Gamma \vdash n : int} [t - int] \\
\frac{}{\Gamma \vdash b : bool} [t - bool] \\
\frac{}{\Gamma \vdash skip : unit} [t - skip] \\
\frac{l \in \Gamma \wedge \Gamma\{l\} : refT}{\Gamma \vdash l : refT} [bel] \\
\frac{\Gamma \vdash e_1 : T_1 \rightarrow T'_1 \quad \Gamma \vdash e_2 : T_1}{\Gamma \vdash e_1 \text{ in } e_2 : T'_1} [t - in] \\
\frac{\Gamma \vdash e_1 : T_1}{\Gamma \text{ fn } x : T \Rightarrow e_1 : T \rightarrow T'_1} [t - fn] \\
\frac{\Gamma \vdash e_1 : refT_1}{\Gamma \vdash !e_1 : T_1} [t - derref] \\
\frac{\Gamma \vdash e_1 : T_1 \quad \Gamma \vdash e_2 : T_1}{\Gamma \vdash e_1 = e_2 : bool} [t - eq] \\
\frac{\Gamma \vdash e_1 : T_1 \quad \Gamma \vdash e_2 : T_1}{\Gamma \vdash e_1 \neq e_2 : bool} [t - neq] \\
\frac{\Gamma \vdash e_1 : bool \quad \Gamma \vdash e_2 : bool}{\Gamma \vdash e_1 \text{ and } e_2 : bool} [t - \wedge] \\
\frac{\Gamma \vdash e_1 : bool \quad \Gamma \vdash e_2 : bool}{\Gamma \vdash e_1 \text{ or } e_2 : bool} [t - or] \\
\frac{\Gamma \vdash e_1 : bool}{\Gamma \vdash notnot \ e_1 : bool} [t - notnot] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 > e_2 : bool} [t - >] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 < e_2 : bool} [t - <] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 \geq e_2 : bool} [t - \geq] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 \leq e_2 : bool} [t - \leq] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 + e_2 : nat} [t - +] \\
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat}{\Gamma \vdash e_1 - e_2 : nat} [t - minus] \\
\frac{\Gamma \vdash e_1 : T_1 \quad \Gamma \vdash e_2 : T_2}{\Gamma \vdash e_1; e_2 : T_1} [t - pontoevirgula] \\
\frac{\Gamma \vdash e_1 : T_1 \quad \Gamma \vdash e_2 : T_2}{\Gamma \vdash e_1 : e_2 : T_2} [t - doisontos]
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma \vdash e_1 : nat \quad \Gamma \vdash e_2 : nat \quad \Gamma \vdash e_3 : T_3}{\Gamma \vdash \text{for } e_1 \text{ until } e_2 \text{ do } e_3 : T_3} [t - \text{for}] \\
\\
\frac{\Gamma \vdash e_1 : bool \quad \Gamma \vdash e_2 : T_2 \quad \Gamma \vdash e_3 : T_2}{\Gamma \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 : T_2} [t - \text{if}] \\
\\
\frac{\Gamma \vdash e_1 : refT_1 \quad \Gamma \vdash e_2 : T_1}{\Gamma \vdash e_2 := e_1 : unit} [t - \text{assing}]
\end{array}$$