



QUERY OPTIMIZATION CULTURAL FACILITIES

Database Technologies 2020/2021

Group G:

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Introduction

This project was developed in the scope of the Database Technologies course from the Master in Informatics and Computing Engineering at FEUP. The goal of the project was to develop different SQL execution plans in a database and analyse the impact of using indexes and different approaches.

To achieve the goal, three different table environments were created: X (no indexes or constraints); Y (standard integrity constraints); Z (standard integrity constraints and indexes).

Note: On the *sql*/ folder delivered, there are three different files: *create.sql* to setup the environment, *indexes.sql* to create the indexes and *queries.sql* with the queries for each question.

Constraints

The constraints added to the Y and Z tables are primary and foreign keys. The code below represents the script which creates these constraints. Primary keys follow a name convention of <TABLE_NAME>_<COLUMN>_PK, while foreign keys follow <TABLE_NAME>_<FOREIGN_TABLE_NAME>_FK.

```
-- Primary Keys
ALTER TABLE "YREGIONS" ADD CONSTRAINT YREGIONS_COD_PK PRIMARY KEY("COD");
ALTER TABLE "YDISTRICTS" ADD CONSTRAINT YDISTRICTS_COD_PK PRIMARY KEY("COD");
ALTER TABLE "YMUNICIPALITIES" ADD CONSTRAINT YMUNICIPALITIES COD PK PRIMARY KEY("COD");
ALTER TABLE "YFACILITIES" ADD CONSTRAINT YFACILITIES ID PK PRIMARY KEY("ID");
ALTER TABLE "YROOMTYPES" ADD CONSTRAINT YROOMTYPES ROOMTYPES PK PRIMARY KEY("ROOMTYPE");
ALTER TABLE "YACTIVITIES" ADD CONSTRAINT YACTIVITIES_REF_PK PRIMARY KEY("REF");
ALTER TABLE "YUSES" ADD CONSTRAINT YUSES_ID_REF_PK PRIMARY KEY("ID", "REF");
ALTER TABLE "ZREGIONS" ADD CONSTRAINT ZREGIONS COD PK PRIMARY KEY("COD");
ALTER TABLE "ZDISTRICTS" ADD CONSTRAINT ZDISTRICTS COD PK PRIMARY KEY("COD");
ALTER TABLE "ZMUNICIPALITIES" ADD CONSTRAINT ZMUNICIPALITIES_COD_PK PRIMARY KEY("COD");
ALTER TABLE "ZFACILITIES" ADD CONSTRAINT ZFACILITIES_ID_PK PRIMARY KEY("ID");
ALTER TABLE "ZROOMTYPES" ADD CONSTRAINT ZROOMTYPES_ROOMTYPES_PK PRIMARY KEY("ROOMTYPE");
ALTER TABLE "ZACTIVITIES" ADD CONSTRAINT ZACTIVITIES REF PK PRIMARY KEY("REF");
ALTER TABLE "ZUSES" ADD CONSTRAINT ZUSES ID REF PK PRIMARY KEY("ID", "REF");
-- Foreign Keys
ALTER TABLE "YDISTRICTS" ADD CONSTRAINT YDISTRICTS_REGION_FK FOREIGN KEY("REGION")
REFERENCES "YREGIONS"("COD");
ALTER TABLE "YMUNICIPALITIES" ADD CONSTRAINT YMUNICIPALITIES REGION FK FOREIGN
KEY("REGION") REFERENCES "YREGIONS"("COD");
ALTER TABLE "YMUNICIPALITIES" ADD CONSTRAINT YMUNICIPALITIES_DISTRICT_FK FOREIGN
KEY("DISTRICT") REFERENCES "YDISTRICTS"("COD");
ALTER TABLE "YFACILITIES" ADD CONSTRAINT YFACILITIES MUNICIPALITY FK FOREIGN
KEY("MUNICIPALITY") REFERENCES "YMUNICIPALITIES"("COD");
```

```
ALTER TABLE "YFACILITIES" ADD CONSTRAINT YFACILITIES_ROOMTYPE_FK FOREIGN KEY("ROOMTYPE")
REFERENCES "YROOMTYPES"("ROOMTYPE");
ALTER TABLE "YUSES" ADD CONSTRAINT YUSES_ID_FK FOREIGN KEY("ID") REFERENCES
"YFACILITIES"("ID");
ALTER TABLE "YUSES" ADD CONSTRAINT YUSES REF FK FOREIGN KEY("REF") REFERENCES
"YACTIVITIES"("REF");
ALTER TABLE "ZDISTRICTS" ADD CONSTRAINT ZDISTRICTS REGION FK FOREIGN KEY("REGION")
REFERENCES "ZREGIONS"("COD");
ALTER TABLE "ZMUNICIPALITIES" ADD CONSTRAINT ZMUNICIPALITIES_REGION_FK FOREIGN
KEY("REGION") REFERENCES "ZREGIONS"("COD");
ALTER TABLE "ZMUNICIPALITIES" ADD CONSTRAINT ZMUNICIPALITIES DISTRICT FK FOREIGN
KEY("DISTRICT") REFERENCES "ZDISTRICTS"("COD");
ALTER TABLE "ZFACILITIES" ADD CONSTRAINT ZFACILITIES_MUNICIPALITY_FK FOREIGN
KEY("MUNICIPALITY") REFERENCES "ZMUNICIPALITIES"("COD");
ALTER TABLE "ZFACILITIES" ADD CONSTRAINT ZFACILITIES ROOMTYPE FK FOREIGN KEY("ROOMTYPE")
REFERENCES "ZROOMTYPES"("ROOMTYPE");
ALTER TABLE "ZUSES" ADD CONSTRAINT ZUSES ID FK FOREIGN KEY("ID") REFERENCES
"ZFACILITIES"("ID");
ALTER TABLE "ZUSES" ADD CONSTRAINT ZUSES REF FK FOREIGN KEY("REF") REFERENCES
"ZACTIVITIES"("REF");
```

Indexes

Here we show the additional created indexes used on the Z environment to improve performance.

Query 1

```
CREATE UNIQUE INDEX INDEX_ZROOMTYPES_ROOMTYPE ON ZROOMTYPE (ROOMTYPE ASC, DESCRIPTION ASC)
CREATE UNIQUE INDEX INDEX_ZACTIVITY_ACTIVITIES ON ZACTIVITIES (REF ASC, ACTIVITY ASC);
```

The first index is used to access the table *ROOMTYPE* through the attribute description and project it to the attribute *room_type* and so using this index, we avoid the access to the full table.

The second index has the same functionality as the first index but it's on the table *ACTIVITIES*, since on the query 1 we need to access both the activity attribute and the ref attribute.

Query 2

```
CREATE UNIQUE INDEX INDEX_ZROOMTYPES_ROOMTYPE ON ZROOMTYPES (ROOMTYPE ASC, DESCRIPTION ASC);

CREATE UNIQUE INDEX INDEX_ZREGIONS_DESIGNATION ON ZREGIONS (COD ASC, DESIGNATION ASC);

CREATE UNIQUE INDEX INDEX_ZMUNICIPALITIES_REGION ON ZMUNICIPALITIES (COD ASC, REGION ASC);

CREATE BITMAP INDEX INDEX_ZFACILITY_FACILITIES ON ZFACILITIES (MUNICIPALITY ASC, ROOMTYPE ASC);
```

Here we use again the same index used on the 1st query on the table *ZROOMTYPE* because we need to access the table to filter the attribute description (in the WHERE condition) and we also need the *roomtype* key to do the join.

The other second and third indexes are used mostly for the join and since they involve primary keys, we can use the unique index.

The last one is the most important one because it involves 2 attributes that are not primary keys, so it's better to use a bitmap instead of a unique index, and doing the test we can see that even if the cardinality is not that low, the cost of the query is lower using the bitmap instead of b-tree index.

Query 3

```
CREATE UNIQUE INDEX INDEX_ZFACILITIES_MUNICIPALITY ON ZFACILITIES (ID, MUNICIPALITY);
CREATE UNIQUE INDEX INDEX_ZACTIVITY_ACTIVITIES ON ZACTIVITIES (REF, ACTIVITY);
```

The indexes used for both queries 3_A and 3_B shown here are very simple: the index on *ZFACILITIES* is used mostly to avoid full access on the table for the join, and the index on *ZACTIVITIES* is used since we have a WHERE clause that involves the attribute *ACTIVITY*.

Query 4

```
CREATE UNIQUE INDEX INDEX_ZFACILITIES_MUNICIPALITY ON ZFACILITIES (ID ASC, MUNICIPALITY ASC);
CREATE UNIQUE INDEX INDEX_ZACTIVITIES_ACTIVITY ON ZACTIVITIES (REF ASC, ACTIVITY ASC);
CREATE UNIQUE INDEX INDEX_ZMUNICIPALITIES_DESIGNATION ON ZMUNICIPALITIES (COD ASC, DESIGNATION ASC);
```

These 3 indexes are the same used in the last 3 queries because some access to the tables are made on attributes and not on primary keys.

Here it is not convenient to use bit-map because the attributes involved in the same table are very less and we are using the primary key as one of the attributes of the indexes, which decreases the efficiency of the bit-map indexes.

Query 5

```
---- Index for 5 a)

CREATE INDEX INDEX_B_TREE ON ZFACILITIES (ROOMTYPE, MUNICIPALITY);

---- Index for 5 b)

CREATE BITMAP INDEX INDEX_BITMAP ON ZFACILITIES (ROOMTYPE, MUNICIPALITY);
```

These indexes used for question 5 were instructed by the assignment, using *roomtype* and *municipality* columns on the *Facilities* table. For question 5 a) a B-Tree index is added and for 5 B) a Bitmap index is added. More details are presented at the Question 5 section <u>here</u>.

Query 6

```
CREATE UNIQUE INDEX INDEX_ZMUNICIPALITIES_DISTRICT ON ZMUNICIPALITIES (COD ASC, DISTRICT ASC);
CREATE UNIQUE INDEX INDEX_ZFACILITIES_MUNICIPALITY ON ZFACILITIES (ID ASC, MUNICIPALITY ASC);
CREATE UNIQUE INDEX INDEX_ZDISTRICTS_COD ON ZDISTRICTS (COD ASC, DESIGNATION ASC);
```

Most of the indexes created here are used to avoid full access to the table for the join. Here we use unique B-tree indexes and not bitmap indexes because the indexes we are creating involve primary keys and have a high cardinality, which may decrease the efficiency of the bitmap indexes.

Questions

Question 1

Which are the facilities where the room type description contains 'touros' and have 'teatro' as one of their activities? Show the id, name, description and activity.

Query

```
SELECT id, name, description, activity
```

FROM xfacilities NATURAL JOIN xroomtypes NATURAL JOIN xuses NATURAL JOIN xactivities
WHERE description LIKE '%touros%' AND activity = 'teatro';

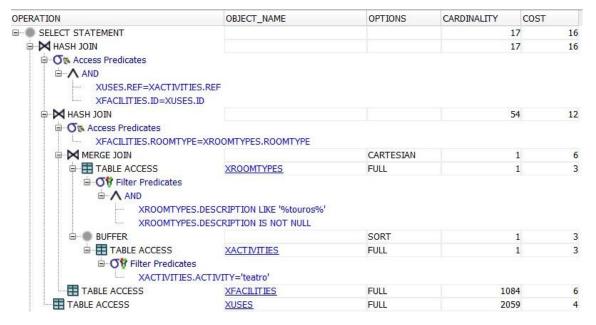
Result

Execution Times

Environment	x	Υ	z
	0.031	0.022	0.019
	0.039	0.019	0.018
Time (s)	0.030	0.024	0.019
	0.055	0.031	0.019
	0.025	0.021	0.019
Average Time (s)	0.036	0.023	0.019

Execution Plan

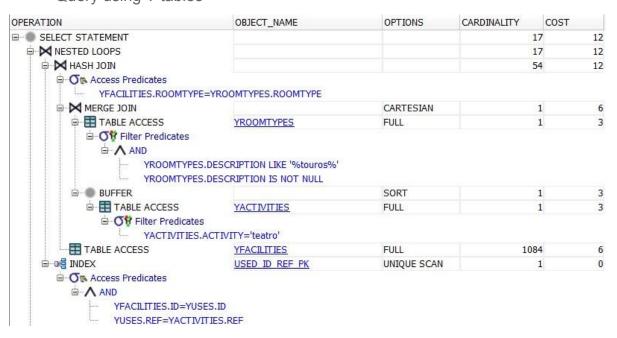
Query using X tables



Comments

The cost of the query for the environment X is mostly dominated by the table access on the tables *XUSES* and *XFACILITIES*, which has a lot of info. In fact we can see that the cost for the full access to the 2 tables are 10 and it's more than half of the total cost of the query.

Query using Y tables



Comments

Here, the cost is significantly reduced due to the primary key index on table *YFACILITIES* and *YUSES* because the join is done with the primary keys, but we still can see that half of the cost is still dominated by the access on table *YFACILITIES*, *YACTIVITIES* and *YROOMTYPES* because we are not accessing the tables using the primary key but by their attributes

Query using Z tables



Comments

In the environment Z, we added the 2 indexes described on the INDEX section, and we can see that the cost is reduced by $\frac{1}{3}$ by adding the 2 indexes on *ZROOMTYPES* and *ZACTIVITY*, which each one index decreases the final cost of the query by 2 values.

Question 2

How many facilities with 'touros' in the room type description are there in each region?

Query

```
SELECT xregions.designation, count(*)
FROM xfacilities NATURAL JOIN xroomtypes INNER JOIN
xmunicipalities
ON xmunicipalities.cod = xfacilities.municipality
INNER JOIN xregions ON xmunicipalities.region = xregions.cod
```

```
WHERE description LIKE '%touros%'
GROUP BY xregions.designation;
```

Result

Lisboa	6
Norte	3
Centro	11
Alentejo	43
Algarve	1

Execution Times

Environment	Х	Y	z
	0.031	0.022	0.020
	0.031	0.023	0.019
Time (s)	0.033	0.021	0.020
	0.023	0.022	0.024
	0.025	0.023	0.019
Average Time (s)	0.028	0.022	0.020

Execution Plans

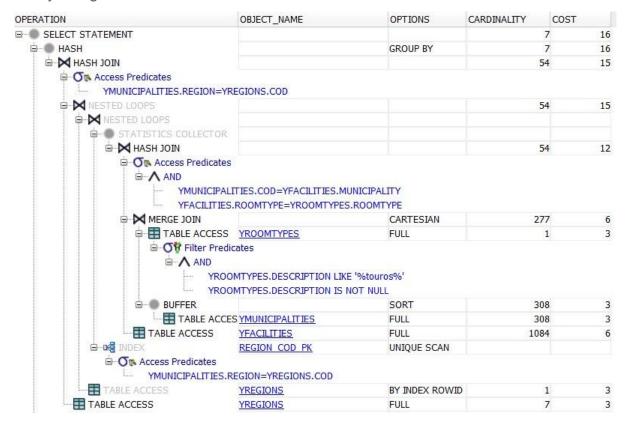
Query using X tables



Comments

In this execution plan the first operation is a full search on the table *XROOMTYPES* while filtering only tuples in which the description of the roomtypes contains 'touros'. We can see here that the cost also depends on the full accesses done within the tables: *XMUNICIPALITIES*, *XFACILITIES*, *XREGIONS*.

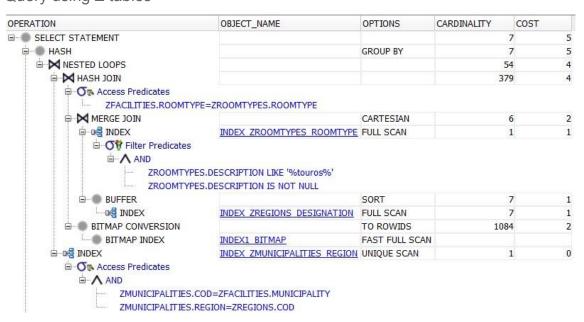
Query using Y tables



Comments

In this case, even by adding the primary key constraint in the table *YREGIONS*, a full search is still needed since the query is selecting the 'designation' of the regions, which is not a primary key. For what concerns the other tables, they are all fully accessed. Therefore the cost of the execution plan X and Y are the same.

Query using Z tables



Comments

In the last environment, we can see that the cost has been significantly decreased due to the indexes used (shown in the indexes paragraph). The reason why one of the indexes is a bit-map is because the attributes where we apply the bit-map indexes are not primary keys. Even if the cardinality is not as shown at lecture less than 1/100, we saw that it still has a good effort on decreasing the cost of the query.

Question 3

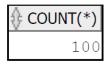
How many municipalities do not have any facility with an activity of 'cinema'?

- a. Use not in.
- b. Use external join and is null.

Query

```
---- a)
SELECT count(*) FROM
xmunicipalities
WHERE cod NOT IN
(SELECT municipality
FROM xfacilities NATURAL JOIN xuses NATURAL JOIN xactivities
INNER JOIN
xmunicipalities ON xmunicipalities.cod = xfacilities.municipality
WHERE activity = 'cinema');
---- b)
SELECT count(*)
FROM xmunicipalities LEFT OUTER JOIN (SELECT municipality as cod,
activity
FROM xfacilities NATURAL JOIN xuses NATURAL JOIN xactivities
INNER JOIN
xmunicipalities ON xmunicipalities.cod = xfacilities.municipality
WHERE activity = 'cinema') cinemas
ON xmunicipalities.cod = cinemas.cod
WHERE activity IS NULL;
```

Result



Execution Times

Query A

Environment	Х	Y	Z
	0.028	0.023	0.023
	0.026	0.022	0.023
Time (s)	0.028	0.024	0.023
	0.025	0.027	0.021
	0.027	0.021	0.020
Average Time (s)	0.027	0.023	0.022

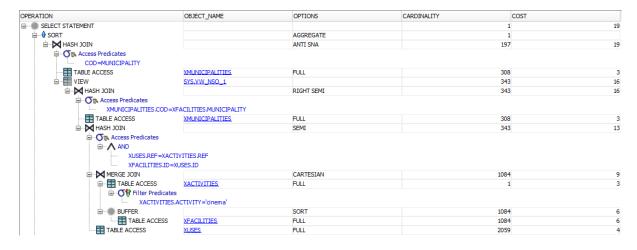
Query B

Environment	x	Υ	Z
	0.024	0.022	0.023
	0.028	0.022	0.023
Time (s)	0.023	0.024	0.021
	0.027	0.022	0.021
	0.025	0.022	0.023
Average Time (s)	0.025	0.022	0.022

Execution Plan

Query A

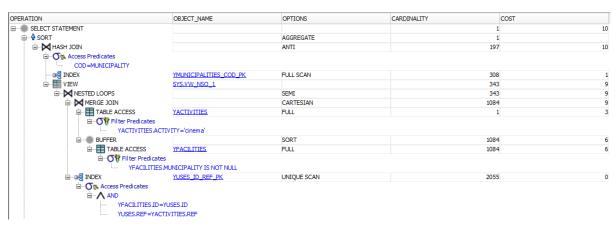
Environment X



Comments

In this environment the high cost is due to the fact that a full search is used 5 times (including twice on the table *MUNICIPALITIES*).

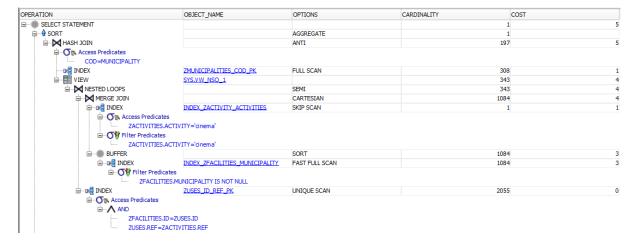
Environment Y



Comments

The biggest change here is the number of full searches that is reduced in more than half. The search on the *USES* table turns into an unique scan, due to the fact that we now have primary and foreign keys, and there is no longer needed two full searches on the table *MUNICIPALITIES*. Instead only a full scan is used.

Environment Z

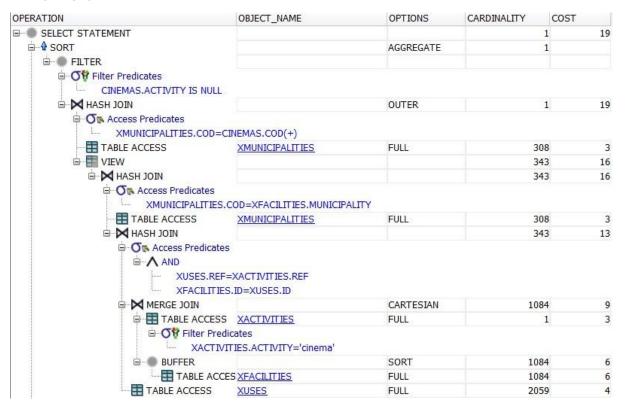


Comments

With the introduction of the two indexes we now have reduced the cost by half. The index on the *FACILITIES* table makes it so that a full search is no longer needed and only a fast full scan is performed. The index on the *ACTIVITIES* table allows for a skip scan to take place instead of a full search.

Query B

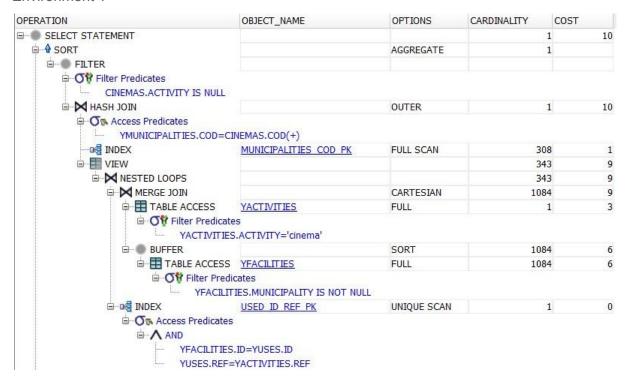
Environment X



Comments

In the environment X, the cost is very high and we can see that this is mostly due to the amount of full access that are performed to the tables, especially to the table XFACILITIES and XUSES (just these two use up 2/3of the total cost).

Environment Y



Comments

Here the cost is decreased because we are using indexes on the primary keys on table *YUSES*, *YMUNICIPALITIES* but we can still notice that there are 2 tables with the full access that can be optimized, especially the table *YFACILITIES*, where one access has a cost of 6, which is more than the half of the total cost.

Environment Z

OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST
□ SELECT STATEMENT			3	1 5
		AGGREGATE		1
□ FILTER				
☐ O ♥ Filter Predicates				
CINEMAS.ACTIVITY IS NU	Ш			
⊟ MASH JOIN		OUTER		1 5
Access Predicates				
ZMUNICIPALITIES.COD	The Control of the Co			
	ZMUNICIPALITIES COD PK	FULL SCAN	30	78
□ II VIEW			34:	
□ MESTED LOOPS			34:	
□ MERGE JOIN		CARTESIAN	1084	
□ □ □ INDEX	INDEX ZACTIVITY ACTIVITIES	SKIP SCAN		1 1
☐ On Access Pre				
	TIES.ACTIVITY='cinema'			
⊟ Of Filter Predi				
	TIES.ACTIVITY='cinema'	12222	1	
⊟ ® BUFFER		SORT	108	
	INDEX ZFACILITIES MUNICIPAL.	FAST FULL SCAN	108	4 3
⊟ OF Filter P				
	ILITIES.MUNICIPALITY IS NOT NULL			
⊡o∉ INDEX	ZUSES ID REF PK	UNIQUE SCAN	1	1 0
⊟ O ∧ Access Predica	ates			
□ ∧ AND				
	TES.ID=ZUSES.ID			
ZUSES.R	EF=ZACTIVITIES.REF			

Comments

In this last environment, we added an index on the *ZFACILITIES* and another on the *ZACTIVITIES*, where both decreased almost by half their cost to access the table (comparing the execution to the previous Y environment).

Question 4

Which is the municipality with more facilities engaged in each of the six kinds of activities? Show the activity, the municipality name and the corresponding number of facilities.

Query

```
SELECT q2.designation, q1.activity, q1.facilities

FROM (SELECT activity, max(facilities) as facilities

FROM (SELECT municipality, activity, count(*) as facilities

FROM xfacilities NATURAL JOIN xuses NATURAL JOIN xactivities

GROUP BY municipality, activity)

GROUP BY activity) q1 LEFT JOIN (SELECT designation, activity, count(*) as facilities

FROM xmunicipalities INNER JOIN xfacilities NATURAL JOIN xuses
```

```
NATURAL JOIN xactivities

ON xmunicipalities.cod = xfacilities.municipality

GROUP BY designation, activity) q2

ON q2.activity = q1.activity AND q2.facilities = q1.facilities;
```

Result

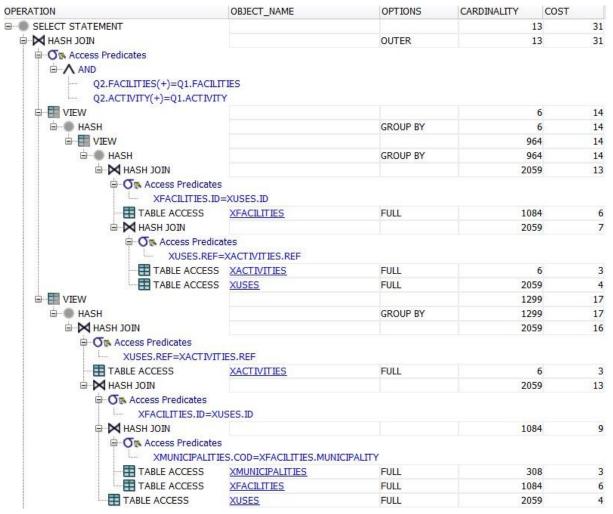
	ON ∯ ACTIVITY	
Lisboa	circo	2
Lisboa	dança	47
Moura	tauromaquia	4
Lisboa	música	77
Lisboa	teatro	66
Lisboa	cinema	96

Execution Times

Environment	x	Y	z
	0.029	0.031	0.028
	0.032	0.031	0.029
Time (s)	0.030	0.027	0.025
	0.025	0.027	0.029
	0.034	0.030	0.026
Average Time	0.030	0.029	0.0274

Execution Plan

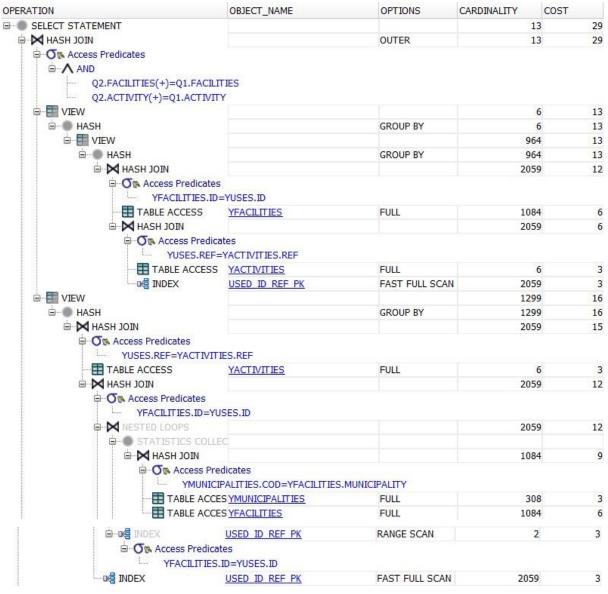
Query using X tables



Comments

In this case, the query requested is a little bit tricky, so the solution proposed uses 3 views as support. Since 2 of the 3 views do very similar things, the cost is inevitably increased, but it's also noticeable that with some indexes the cost of table access can be decreased, and so the total cost of the query.

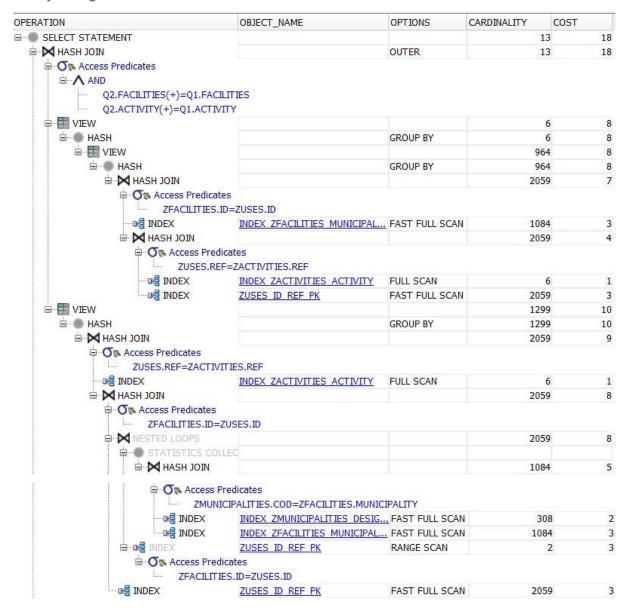
Query using Y tables



Comments

In the environment Y, we can see that the cost is a bit less with the introduction of indexes on the primary key like on table *YUSED*, but we see that there are still a lot of full access to the tables like on *YFACILITIES* and *YMUNICIPALITIES* that can be optimized by using appropriate indexes as shown in the environment Z

Query using Z tables



Comments

Here we can see that the cost has been decreased by $\frac{2}{3}$ comparing it to the cost of the environment Y, and that could be done only by introducing indexes on table *ZMUNICIPALITIES* and *ZFACILITIES*.

The total cost, compared to the other queries, is still very high since the query is a bit tricky but reducing almost to half the cost from environment X to environment Z using adequate indexes is a very good result.

Question 5

Compare the execution plans (just the Z environment) for the query giving the municipality designation, the facility name, and the room type description, where the

room type of the facility includes 'touros' and the municipality is part of the Porto district.

- a. With a B-tree index on the room type and municipality columns of the facilities table;
- b. With a bitmap index on the room type and municipality columns of the facilities table.

Query

```
SELECT zmunicipalities.designation, zfacilities.name, zroomtypes.description
FROM zfacilities NATURAL JOIN zroomtypes INNER JOIN zmunicipalities On zmunicipalities.cod = zfacilities.municipality INNER JOIN zdistricts ON zdistricts.cod = zmunicipalities.district
WHERE zdistricts.designation = 'Porto' AND zroomtypes.description LIKE '%touros%';
```

Result

	OITA	V	NAME								PTION	١
Póvoa	de	Varzim	PRAÇA	DE	TOIROS	DA	PÓVOA	DE	VARZIM	Praça	de	touros

Execution Times

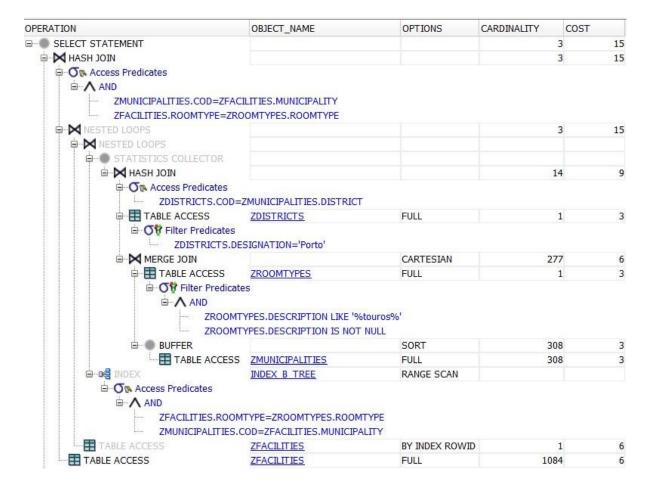
Index	B-TREE	BITMAP
	0.026	0.024
	0.024	0.028
Time (s)	0.025	0.025
	0.026	0.027
	0.023	0.029
Average Time	0.025	0.027

Execution Plans

The main difference between both execution plans below is that, even though both indexes use the same columns/table and have the same cost, the Bitmap index is

not even used. That must be because of the cardinality of the *Facilities* table which is very high, making the B-Tree more effective.

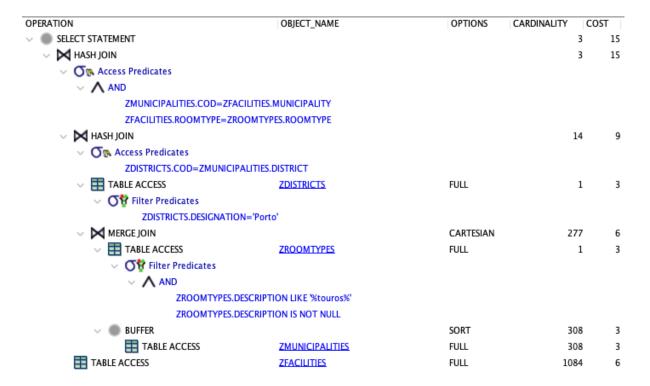
B-Tree



Comment

As we can see, the access to the *ROOMTYPE* and *MUNICIPALITY* on the join clause is done using the index. The only problem here is that it is doing it twice probably because some of the attributes are not part of the B-tree index.

Bitmap



Comment

Using the bitmap index we can see that it is not used and the cost remains the same.

Question 6

Which are the codes and designations of the districts with facilities in all the municipalities?

Query

```
SELECT cod, designation FROM xdistricts WHERE cod NOT IN

(
SELECT xdistricts.cod
FROM xmunicipalities INNER JOIN xdistricts ON
xmunicipalities.district = xdistricts.cod
LEFT OUTER JOIN xfacilities on xfacilities.municipality =
xmunicipalities.cod
WHERE id is null
);
```

Result

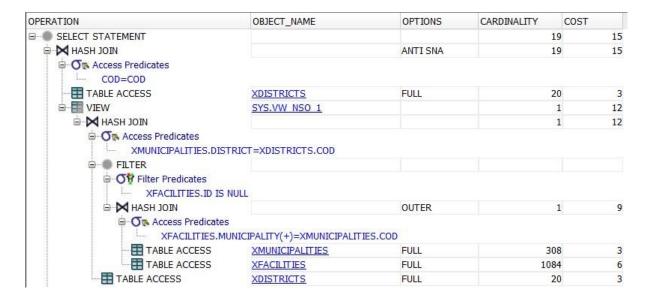
∯ COD	
15	Setúbal
7	Évora
11	Lisboa
12	Portalegre

Execution Times

Environment	x	Y	z
Time (s)	0.026	0.022	0.019
	0.034	0.021	0.020
	0.035	0.020	0.018
	0.031	0.022	0.020
	0.028	0.035	0.019
Average Time	0.031	0.024	0.019

Execution Plans

Query using X tables



Comments

As in previous queries, in the environment X, tables are fully accessed. First, a full search of XMUNICIPALITIES and XFACILITIES are done. Then only the district codes of tuples containing *null* values in the Facilities' ID field are selected. Using NOT IN operator, these results are the ones that are not considered during the full access on the table XDISTRICTS.

Query using Y tables



Comments

In the environment Y for this query the cost is reduced due to the introduction of the primary key in the *YDISTRICTS* table, named 'cod'. The reduction of the cost is not significant, since full accesses are still done on the tables *YMUNICIPALITIES* and *YFACILITIES*. To improve the search on these tables we should add specific indexes.

Query using Z tables



Comments

In the last environment the cost is definitely improved since it is reduced by $\frac{1}{3}$ compared to the cost from environment X and almost $\frac{1}{3}$ compared to the cost from the environment Y. From this, it was our understanding that the cost is reduced especially when appropriate indexes are introduced, in this case, on the table ZFACILITIES and ZMUNICIPALITIS.

Conclusion

With the development of the project, the group learned more about the impact of indexes on queries and how it can benefit the overall performance. Still, the indexes must be carefully selected, especially regarding its type (B-Tree or Bitmap), since it may not be effective in some circumstances and just create additional costs.