Team Reference Document

HaKings

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```
1 vimrc
syntax on
inoremap jj <ESC>
colorscheme elflord
set ai si sw=4 ts=4
set nu
set backspace=start,indent,eol
set clipboard=unnamed
set ignorecase
                                                                                    10
11
12
set smartcase
set incsearch
set scrolloff=3
                                                                                    13
highlight linenr ctermbg=darkblue
                                                                                    15
                                                                                    16
         Header
```

6 7

38 39 40

```
#include <bits/stdc++.h>
#Include \( \text{bits/stdc++,h} \)
#define _ ios_base::sync_with_stdio(0), \( \text{cin.tie}(0), \), \( \text{cin.tie}(0), \), \( \text{cin.tie}(0), \), \( \text{cin.tie}(0); \)
#define INF 1000000000
#define FOR(i, a, b) for(int i=int(a); i<int(b); i++) 24
#define FORC(cont, it) for(decltype((cont).begin()) it 26
= (cont).begin(); it!=(cont).end(); it++) 27
#define by nucle base()
#define by nucle base()</pre>
 #define pb push_back
#define mp make_pair
#define eb emplace_back
                                                                                                                                                                       29
```

Primes

```
41 \\ 42 \\ 43 \\ 44
                        #define SIZE 1000000
                        bitset<SIZE> sieve;
void buildSieve() {
                               sieve.set();
sieve[0] = sieve[1] = 0;
int root = sqrt(SIZE);
FOR(i, 2, root+1)
                                                                                                                                                                                                                                                                        45
                                                                                                                                                                                                                                                                        46
47
                                                                                                                                                                                                                                                                        48
                                       if (sieve[i])
  for(int j = i*i; j < SIZE; j+=i)
    sieve[j] = 0;</pre>
                                                                                                                                                                                                                                                                        49
50
10
                                                                                                                                                                                                                                                                        51
                                                                                                                                                                                                                                                                       52
53
54
55
13
14
                        vi primesList:
                         void buildPrimesList() {
  if(!sieve[2])
                                                                                                                                                                                                                                                                        56
57
58
                                         buildSieve();
                               primesList.reserve(SIZE/log(SIZE));
FOR(i, 2, SIZE+1)
  if(sieve[i])
18
19
20
                                                                                                                                                                                                                                                                        59
                                                primesList.pb(i);
21
22
                                                                                                                                                                                                                                                                        62
                        vii primeFactorization(int N) {
24
                                 int idea to idea 
25
26
27
28
                                                 29
30
31
32
                                                          factors.pb(ii(pf, 1));
33
34
35
                                         pf = primesList[++idx];
36
37
                                 if(N!=1) factors.pb(ii(N, 1));
38
39
40
                         void getDivisors(vii pf, int d, int index, vi &div) {
   if (index == pf.size()) {
41
                                         div.pb(d);
43
44
46
47
                               for (int i = 0; i <= pf[index].second; i++) {
   getDivisors(pf, d, index+1, div);
   d *= pf[index].first;</pre>
48
49
50
51
53
54
55
                        vi divisors(l1 N) {
  vii pf = primeFactorization(N);
  vi div;
56
57
58
                                 getDivisors(pf, 111, 0, div);
sort(div.begin(), div.end());
                               return div;
59
                                                                                                                                                                                                                                                                        10
                        bool isPrime(int n)
                                                                                                                                                                                                                                                                        11
                              col isPrime(int n) {
   if(n < 2) return false;
   if(n = 2 || n == 3) return true;
   if(!(nsl 6s n%3)) return false;
   long long sgrtN = sgrt(n)+1;
   for(long long i = 6LL; i <= sgrtN; i += 6)
   if(!(n%(i-1)) || !(n%(i+1))) return false;</pre>
62
63
                                                                                                                                                                                                                                                                        13
66
67
                                                                                                                                                                                                                                                                        16
```

```
Segment Tree
```

```
struct SegmentTree {
    vi t; int N;
   vi t; int N;
SegmentTree(vi &values) {
    N = values.size();
    t.assign(N<<1, 0);
    for(int i = 0; i < N; i++) t[i+N] = values[i];
    for(int i = N-1; i; --i) t[i] = combine(t[i<<1], t[28 i<<1|1]);
}</pre>
                                                                                                          30
   int combine(int a, int b) { return a+b; }
void set(int index, int value) {
    t[index+N] = value;
    for(int i = (index+N)>>1; i; i>>= 1) t[i] =
                    combine(t[i<<1], t[i<<1|1]);
       int ansL = 0, ansR = 0;
for(int 1 = N+from, r = N+to; 1<r; 1 >>= 1, r >>= 39
           1) {
   if (1&1) ansL = combine(ansL, t[1++]);
   if (r&1) ansR = combine(ansR, t[--r]);
};
struct LazySegmentTree {
  vi t, d; int n, h;
  LazySegmentTree(vi &values) {
       ryoid calc(int p, int k) {
   if (d[p] == 0) t[p] = t[p<<1] + t[p<<1|1];
   else t[p] = d[p] * k;</pre>
                                                                                                           13
   void apply(int p, int value, int k) {
  t[p] = value * k;
  if (p < n) d[p] = value;</pre>
   for (int i = 1 >> s; i <= r >> s; ++i) if (d[i])
               apply(i<<1, d[i], k);
apply(i<<1|1, d[i], k);
d[i] = 0;
                                                                                                          25
    void build(int 1, int r) {
                                                                                                          29
       int k = 2;
for (1 += n, r += n-1; 1; k <<= 1) {
    1 >>= 1, r >>= 1;
    for (int i = r; i >= 1; --i) calc(i, k);
                                                                                                           \frac{32}{33}
                                                                                                           34
35
    void modify(int 1, int r, int value) {
                                                                                                           36
       if (value == 0) return; 37

push(1, 1 + 1); push(r - 1, r); 38

int 10 = 1, r0 = r, k = 1; 39

for (1 + e n, r += n; 1 < r; 1 >>= 1, r >>= 1, k <<=
1) (
40
           if (1&1) apply(1++, value, k);
if (r&1) apply(--r, value, k);
                                                                                                           43
       build(10, 10 + 1);
build(r0 - 1, r0);
                                                                                                           46
   int query(int 1, int r) {
       nt query(int 1, int r) {
push(1, 1 + 1); push(r - 1, r);
int res = 0;
for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (161) res += t[1++1;
           if (r&1) res += t[--r];
};
```

Geometry

5.1 Point

19

20 21 22

```
const double PI = 2*asin(1);
bool eq(double a, double b) { return fabs(a-b) < EPS; } 8 bool les(double a, double b) { return !eq(a, b) && a < 9 }
        b: }
                                                                          10
  double x, y, z;
Point(): x(0), y(0), z(0) {}
Point(double x, double y): x(x), y(y), z(0) {}
13
Point(double x, double y, double z): x(x), y(y), z(z15
  bool operator <(const Point &p) const {
          bool operator==(const Point &p) { 22
return eq(x, p.x) && eq(y, p.y) && eq(z, p.z); 23
                                                                          25
double DEG_to_RAD(double deg) {
     return deg/180*2*asin(1);
double dist(Point pl, Point p2) {
```

```
return sqrt(pow(p1.x-p2.x, 2) + pow(p1.y-p2.y, 2) +
               pow(p1.z-p2.z, 2)); }
Point rotate(Point p, double theta) {
   double rad = DEG_to_RAD(theta);
return Point(p.x*cos(rad) - p.y*sin(rad),
    p.x*sin(rad) + p.y*cos(rad));
double ANG(double rad) { return rad*180/PI; }
double angulo(Point p) {
   double d = atan(double(p.y)/p.x);
   if(p.x < 0)
   d += PI;</pre>
   else if (p.y < 0)
d += 2*PI;
   return ANG(d);
```

5.2 Vector

```
double x, y, z;
Vec(double x, double y, double z) : x(x), y(y), z(z)
   Vec(): x(0), y(0), z(0) {}
Vec(double x, double y): x(x), y(y), z(0) {}
Vec(Point a, Point b): x(b.x-a.x), y(b.y-a.y), z(b.z-a.z) {}
Vec toVec(Point a, Point b) {
   return Vec(a, b); }
   return Vec(v.x*s, v.v*s, v.z*s); }
Point translate(Point p, Vec v) {
   return Point (p.x+v.x, p.y+v.y, p.z+v.z); }
double dot(Vec a, Vec b) {
  return (a.x*b.x + a.y*b.y + a.z*b.z); }
double norm_sq(Vec v) {
   return v.x*v.x + v.y*v.y + v.z*v.z; }
 //angle in radians
//angle In Intains
Vec rotate(Vec v, double angle) {
  Matrix rotation = CREATE(2, 2);
  rotation[0][0] = rotation[1][1] = cos(angle);
  rotation[1][0] = sin(angle);
  rotation[0][1] = -rotation[1][0];
   Matrix vec = CREATE(2, 1);
   vec[0][0] = v.x, vec[0][1] = v.y;
   Matrix res = multiply(rotation, vec);
Vec result(res[0][0], res[0][1]);
   return result;
double cross (Vec a, Vec b) { return a.x*b.y - a.y*b.x;
// returns true if r is on the left side of line pq \bf bool ccw(Point p, Point q, Point r) {
  return cross(toVec(p, q), toVec(p, r)) > 0; }
bool collinear (Point p, Point q, Point r) {
   return abs(cross(toVec(p, q), toVec(p, r))) < EPS; }
5.3 Triangle
```

```
struct Triangle {
   Triangle()
   Triangle (Point A, Point B, Point C) : A(A), B(B), C(C
double perimeter(double a, double b, double c) { return
double area(double a, double b, double c) {
   double s = perimeter(a, b, c)*0.5;
return sqrt(s*(s-a)*(s-b)*(s-c));
double area(const Triangle &T) {
     ble area(const Triangle &T)
double ab = dist(T.A, T.B);
double bc = dist(T.B, T.C);
double ca = dist(T.C, T.A);
return area(ab, bc, ca);
double rInCircle(double ab, double bc, double ca) {
  return area(ab, bc, ca) / (0.5 * perimeter(ab, bc, ca)
double rInCircle(Point a, Point b, Point c) {
  return rInCircle(dist(a, b), dist(b, c), dist(c, a));
bool inCircle(Point pl, Point p2, Point p3, Point &ctr,
```

```
r = rInCircle(pl, p2, p3);
if(abs(r) < EPS) return false;
Line 11, 12;
double ratio = dist(p1, p2) / dist(p1, p3);</pre>
                                                                                                                                                                                                                                if (pointInPolygon(p2, p1[i]))
                                                                                                                                                                                                                                result.insert(p1[i]);
FOR(j, 0, p2.size() - 1) {
   Line 11 = Line(p1[i], p1[i+1]);
   Line 12 = Line(p2[j], p2[j+1]);
                                                                                                                 11 cross(const Point &O, const Point &A, const Point &HO7
30
31
                                                                                                                    return (A.x - O.x) * (B.y - O.y) - (A.y - O.y) * (B.1009
32
33
             Point p = translate(p2, scale(toVec(p2, p3), ratio
                                                                                                                                                                                                                                  Line 12 = Line(p2[j], p2[j+1]);
vector<Point> ps1, ps2;
ps1.pb(p1[i]); ps1.pb(p1[i+1]);
ps2.pb(p2[j]); ps2.pb(p2[j+1]);
sort(ps1.begin(), ps1.end());
sort(ps2.begin(), ps2.end());
if (lareParallel(11, 12)) {
Point intersect;
bool b = areIntersect(11, 12, intersect);
if (b && checkPointInSegm(intersect, ps1[0], ps1[1]) && checkPointInSegm(intersect, ps2[0], ps2[1])
    result.insert(intersect);
else if (areSame(11, 12)) {
                                                                                                                                  - O.x);
                                                                                                          5
                         /(1+ratio)));
            /(1+ratio)));
11 = Line(p1, p);
ratio = dist(p2, p1) / dist(p2, p3);
12 = Line(p2, p);
areIntersect(11, 12, ctr);
                                                                                                                 Polygon convexHull(Polygon &P) {
36
                                                                                                                    int n = P.size(), k = 0;
                                                                                                                                                                                                              114
                                                                                                                    Polygon H(2*n);
sort(P.begin(), P.end());
FOR(i, 0, n) {
37
             return true;
                                                                                                                        12
40
         double <code>rCircumCircle(double</code> ab, double bc, double ca) {    return ab * bc * ca / (4.0 * area(ab, bc, ca));13
41
                                                                                                                        H[k++] = P[i];
                                                                                                                    for (int i = n-2, t = k+1; i >= 0; i--) {
42
                                                                                                                                                                                                                                   result.insert(intersect);
else if (areSame(11, 12)) {
    if (ps1[1] >= ps2[0] && ps2[1] >= ps1[0]) {
        vector<Point> ps3;
        ps3.pb(ps1[0]); ps3.pb(ps1[1]); ps3.pb(ps2
            [0]); ps3.pb(ps2[1]);
        sort(al1(ps3));
    result.insert(ps3[1]);
         Point circumcenter(const Triangle &T) { 16} Point A = T.A, B = T.B, C = T.C; double D = 2*(A.x*(B.y - C.y) + B.x*(C.y - A.y) + C.x17
43
                                                                                                                       while (k >= t && cross(H[k-2], H[k-1], P[i]) <= 0)121
44
45
                                                                                                                        H[k++] = P[i];
            124
46
                                                                                                                    H resize(k).
47
                                                                                                                                                                                                              126
                                                                                                                                                                                                              127
                                                                                                                                                                                                                                          result.insert(ps3[2]);
                                                                                                                  // return area when Points are in cw or ccw, p[0]
                                                                                                                 n-1]
double area(const Polygon &P) {
                                                                                                                                                                                                                               }
                                                                                                                                                                                                              130
                                                                                                                    double area (conser Forygon &r) {
    double result = 0.0, x1, y1, x2, y2;
    for (int i = 0; i < (int) P. size()-1; i++) {
        x1 = P[i].x; x2 = P[i+1].x;
        y1 = P[i].y; y2 = P[i+1].y;
        result += (x1*y2-x2*y1);
    }
}</pre>
                                                                                                       25
                                                                                                                                                                                                              131
         5.4 Lines
                                                                                                        26
27
                                                                                                                                                                                                                             FOR(i, 0, p2.size() - 1) {
                                                                                                                                                                                                                               if (pointInPolygon(pl, p2[i]))
                                                                                                        28
                                                                                                                                                                                                              134
          struct Line {
                                                                                                        29
                                                                                                                                                                                                              135
                                                                                                                                                                                                                                   result.insert(p2[i]);
            double a, b, c;

Line(): a(0), b(0), c(0) {}

Line(Point p1, Point p2) {

if (abs(p1.x-p2.x) < EPS) {

a = 1.0; b = 0.0; c = -p1.x;
                                                                                                        30
31
                                                                                                                    return abs(result) / 2.0;
                                                                                                        32
                                                                                                                                                                                                              138
                                                                                                                                                                                                                            if (result.size() <= 2) {
                                                                                                        33
                                                                                                                                                                                                              139
                                                                                                                                                                                                                                return Polygon(result.begin(), result.end());
                                                                                                                                                                                                              \frac{140}{141}
                                                                                                                 bool isConvex(const Polygon &P) {
                                                                                                                    int sz = (int)P.size(); 1

if (sz <= 3) return false; 1

bool isLeft = ccw(P[0], P[1], P[2]); 1

for (int i = 1; i < sz-1; i++) 1

if (ccw(P[i], P[i+1], P[(i+2) == sz ? 1 : i+2]) !=
                          -(double) (p1.y-p2.y)/(p1.x-p2.x);
                                                                                                        36
                                                                                                                                                                                                              142
                                                                                                                                                                                                                            Polygon p(result.begin(), result.end());
return convexHull(p);
                                                                                                        37
38
39
                    c = -(double) (a*pl.x)-pl.y;
\frac{10}{11}
12
                                                                                                                                                                                                                          6 Suffix Array
13
14
                                                                                                        40
                                                                                                                           return false:
         bool areParallel(Line 11, Line 12) {
   return (abs(11.a-12.a) < EPS) && (abs(11.b-12.b) <
        EPS); }</pre>
15
                                                                                                                 // works for convex and concave
bool inPolygon (Point pt, const Polygon &P) {
  if((int)P.size() == 0) return false;
                                                                                                                                                                                                                         struct SuffixArray {
                                                                                                                                                                                                                            17
         bool areSame(Line 11, Line 12) { 46 return areParallel(11, 12) && (abs(11.c-12.c) < EPS);47
                                                                                                                    if((int)P.size() == 0) return reise;
double sum = 0;
for (int i = 0; i < (int)P.size()-1; i++) {
   if (ccw[pt, P[i], P[i+1]))
      sum += angle(P[i], pt, P[i+1]);
   else sum -= angle(P[i], pt, P[i+1]);
}
return abs(abs(sum) - 2*PI) < EPS;</pre>
19
20
         bool areIntersect(Line 11, Line 12, Point &p) { 50 if (areParallel(11, 12)) return false; 51 p.x = (12.b * 11.c - 11.b * 12.c) / (12.a * 11.b - 1152 co)
                                                                                                                                                                                                                                FOR(j, 0, 2) {
    vi freq(N, 0), rank(N, 0), count(N, 0);
    FOR(i, 0, N) freq[get(val[i], j)]++;
    for(int i = 0, k = 0; i < N; i++) rank[i] = k, k</pre>
23
                         .a * 12.b);
            24
                                                                                                                55
                                                                                                                                                                                                                                   FOR(i, 0, N) (*nex)[rank[get(val[(*cur)[i]], j)]

[++] = (*cur)[i];
            return true;
26
                                                                                                                                                                                                               13
27
                                                                                                                                                                                                                                   swap(cur, nex);
          // Interseccion de AB con CD
         // * WARNING: Does not work for collinear line segments59
                                                                                                                                                                                                                15
30
31
          bool lineSegIntersect(Point a, Point b, Point c, Point 61
                                                                                                                                                                                                                             SuffixArray(char *S, int N) : N(N) {
            d) {
double ucrossv1 = cross(toVec(a, b), toVec(a, c));
double ucrossv2 = cross(toVec(a, b), toVec(a, d));
if (ucrossv1 * ucrossv2 > 0) return false;
double vcrossu1 = cross(toVec(c, d), toVec(c, a));
double vcrossu2 = cross(toVec(c, d), toVec(c, b));
return (vcrossu1 * vcrossu2 <= 0);</pre>
                                                                                                                                                                                                                                sa.assign(N, 0);

FOR(i, 0, N) sa[i] = i;

vi freq(Q, 0);
                                                                                                                        return false;
                                                                                                                                                                                                                19
32
                                                                                                                                                                                                                20
21
33
                                                                                                                                                                                                                               vi freq(0, 0);
int index[0], rank[N], k = 0;
FOR(i, 0, N) freq[S[i]]++;
FOR(i, 0, Q) index[i] = k, k += freq[i];
FOR(i, 0, N) rank[i] = index[S[i]];
for(int len = 2; len <= N*2; len <<= 1) {</pre>
                                                                                                                                                                                                               22
35
                                                                                                                                                                                                                23
36
37
38
39
                                                                                                                 Point lineIntersectSeg(Point p, Point q, Point A, Point^{25}
                                                                                                                   40
          // Calcula la distancia de un punto P a una recta AB, \sqrt{7}0
         guarda en C la inters

double distToLine(Point p, Point a, Point b, Point &c)
41
            Vec ap = toVec(a, p), ab = toVec(a, b);
double u = dot(ap, ab) / norm_sq(ab);
c = translate(a, scale(ab, u));
42
43
                                                                                                        75
                                                                                                                                                                                                                                FOR(i, 0, N) sa[rank[i]] = i;
45
             return dist(p, c);
                                                                                                                                                                                                                33
                                                                                                                                                                                                                                buildLCP(S);
46
                                                                                                        78
                                                                                                                 // cuts polygon Q along line AB Polygon cutPolygon (Point a, Point b, const Polygon &Q) 35
                                                                                                                                                                                                                             void buildLCP(char *S) {
              Distancia a de P a segmento AB
                                                                                                        79
                                                                                                                                                                                                                               vi phi(N), plcp(N);

int L = 0;

phi[sa[0]] = -1;
                                                                                                                                                                                                               36
49
         double distToLineSegment(Point p, Point a, Point b,
            puble distToLineSegment(Foint p, Foint a, Fo
    Point $c$ {
    Vec ap = toVec(a, p), ab = toVec(a, b);
    double u = dot(ap, ab) / norm_sq(ab);
    if (u < 0.0) { c = a; return dist(p, a); }
    if (u > 1.0) { c = b; return dist(p, b); }
    return distToLine(p, a, b, c);
                                                                                                        80
                                                                                                                     Polygon P:
                                                                                                                       or (int i = 0; i < (int)Q.size(); i++) { 38 double left1 = cross(toVec(a, b), toVec(a, Q[i+1]))\frac{39}{1} if (i != (int)Q.size()-1) left2 = cross(toVec(a, b)\frac{41}{1}
                                                                                                                                                                                                                               pnr[sa[0]] --1;
POR(i, 1, N) phi[sa[i]] = sa[i-1];
POR(i, 0, N) {
    if(phi[i] == -1) { plcp[i] = 0; continue; }
    while(S[i+L] == S[phi[i]+L]) L++;
51
52
                                                                                                        83
                                                                                                                        plcp[i] = L;
L = max(L-1, int(0));
                                                                                                                           P.pb(lineIntersectSeg(Q[i], Q[i+1], a, b));
                                                                                                        86
                                                                                                                                                                                                                                FOR(i, 0, N) lcp.pb(plcp[sa[i]]);
         5.5 Circles
                                                                                                                                                                                                                46
                                                                                                                     if (!P.empty() && !(P.back() == P.front()))
                                                                                                                                                                                                                47
                                                                                                                                                                                                                         };
                                                                                                                        P.pb(P.front());
                                                                                                                    return P;
         bool circle2PtsRad(Point pl, Point p2, double r, Point
                                                                                                                                                                                                                                   Linear Suffix Array
             double d2 = (p1.x - p2.x) * (p1.x - p2.x) + (p1.y -
 2
            double d2 = (p1.x - p2.x) * (p1.x - p2.x) + (p.
p2.y) * (p1.y - p2.y);
double det = r * r / d2 - 0.25;
if (det < 0.0) return false;
double h = sqrt(det);
c.x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
c.y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;</pre>
                                                                                                                  // only works for convex
                                                                                                                 FOR(i, 0, pl.size() - 1)
if (cross(pl[i], pl[i+1], p) >= 0)
                                                                                                                                                                                                                         /∗ Linear Suffix Array
                                                                                                                           return false;
                                                                                                                                                                                                                            int N = 6, SA[6];

char S[6] = "abcab";

SA_IS((unsigned char*)S, SA, N, 256);

FOR(i, 0, N) cout << S+SA[i] << endl;
             return true:
          \} // to get the other center, reverse p1 and p2
                                                                                                      100
                                                                                                                 // polygons must be convex
                                                                                                                 // returns polygon with size < 3 if there is no intersection
         5.6 Polygons
                                                                                                                 Polygon intersection(Polygon &pl, Polygon &p2) {
                                                                                                                                                                                                                         unsigned char mask[] = { 0x80, 0x40, 0x20, 0x10, 0x08,
                                                                                                                                                                                                               10
                                                                                                      104
                                                                                                                     set<Point> result:
                                                                                                                                                                                                                         0x04, 0x02, 0x01 };
#define tget(i) ( (t[(i)/8]&mask[(i)%8]) ? 1 : 0 )
                                                                                                      105
                                                                                                                    FOR(i, 0, pl.size() - 1) {
         typedef vector<Point> Polygon;
```

```
#define tset(i, b) t[(i)/8]=(b) ? (mask[(i)%8]|t[(i) 11 /8]) : (("mask[(i)%8])&t[(i)/8]) 12 #define chr(i) (cs==sizeof(int)?((int*)s)[i]:((unsigned13 char*)s)[i]) 14 #define isLMS(i) (i>0 && tget(i) && tget(i-1)) 15
  12
                                                                                                           struct Trie {
  13
          21
  22
  23
                                                                                                   23
          void induceSAl (unsigned char *t, int *SA, unsigned char25
 25
                      *s, int *bkt, int n, int K, int cs, bool end) \{26
              getBuckets(s, bkt, n, K, cs, end);
             for (i = 0; i < n; i++) {
    j = SA[i] - 1;
    if (j >= 0 && !tget(j))
        SA[bkt[chr(j)]++] = j;
  28
                                                                                                   29
  29
                                                                                                   30
  31
  32
                                                                                                   33
                                                                                                                     return 0:
  33
  34
          void induceSAs(unsigned char *t, int *SA, unsigned char35
     *s, int *bkt, int n, int K, int cs, bool end) {36
   int i, j; 37
  35
             int i, j;
ggtBuckets(s, bkt, n, K, cs, end);
for (i = n - 1; i >= 0; i--) {
    j = SA[i] - 1;
    if (j >= 0 && tget(j)) SA[--bkt[chr(j)]] = j;
  36
37
  38
                                                                                                   40
  39
                                                                                                   41
                                                                                                                     return 0:
  42
                                                                                                   44
                                                                                                           };
  43
          void SA_IS (unsigned char *s, int *SA, int n, int K, int
                                                                                                                     DSU
  45
              unsigned char *t = (unsigned char *) malloc(n / 8 +
              1); // LS-type array in bits tset(n-2, 0);
  47
              tset(n-1, 1);
             test(i, i, i);
for (i = n - 3; i >= 0; i--)
tset(i, (chr(i) < chr(i+1) || (chr(i) == chr(i+1) &&
tget(i+1) == 1) ? 1: 0);
int *bkt = (int *) malloc(sizeof(int) * (K + 1));</pre>
  48
  49
             int *bRt = (int *) mailoc(sizeor(int) getBuckets(s, bkt, n, K, cs, 1);
for (i = 0; i < n; i++) SA[i] = -1;
for (i = 1; i < n; i++) if (isLMS(i))
SA[-bkt[chr(i)]] = i;
induceSA1(t, SA, s, bkt, n, K, cs, 0);
induceSAs(t, SA, s, bkt, n, K, cs, 1);
free(bkt)</pre>
  51
                                                                                                                        j);}
  55
              free (bkt);
                                                                                                   10
             int n1 = 0;
for (i = 0; i < n; i++)
  if (isLMS(SA[i]))</pre>
  58
  59
             13
  61
                                                                                                   14
  62
  63
  64
65
  66
  68
69
  70
71
72
                   break; } else if (d > 0 && (isLMS(pos+d) || isLMS(prev+\frac{22}{24}
                 ))) break;
if (diff) {
  73
74
75
76
77
78
                                                                                                   26
                    name++;
                   prev = pos;
                 pos = (pos % 2 == 0) ? pos / 2 : (pos - 1) / 2;
SA[n1 + pos] = name - 1;
              for (i = n - 1, j = n - 1; i >= n1; i--) if (SA[i] >=
             10
  82
                                                                                                           struct LCA {
             85
  88
  89
  92
                                                                                                   13
                 SA[--bkt[chr(j)]] = j;
  96
              induceSAl(t, SA, s, bkt, n, K, cs, 0); induceSAs(t, SA, s, bkt, n, K, cs, 1);
                                                                                                   16
  98
99
              free (bkt);
100
                                                                                                   20
                                                                                                   21
22
23
                    Trie
                                                                                                   24
                                                                                                   25
                                                                                                   28
          Constructs a tree for storing strings
                                                                                                                     size() + 1; 1 < r; 1 >>= 1, r >>= 1) {

if (1&1) idx = minIndex(idx, st[1++]);

if (r&1) idx = minIndex(idx, st[--r]);
           #define ALPHABET SIZE 52
          #define ALPHABET_SIZE 52
int getIndex(char c) {
   if(c >= 'A' && c <= 'Z')
      return c-'A';
   return c-'a'+26;
                                                                                                   31
                                                                                                                  return order[idx];
                                                                                                           };
```

```
truct Trie {
int words, prefixes;
Trie *edges[ALPHABET_SIZE];
Trie(): words(0), prefixes(0) { FOR(i, 0, 2 }
ALPHABET_SIZE] edges[i] = 0; }

Trie() { FOR(i, 0, ALPHABET_SIZE) if(edges[i]) delete 4 }
edges[i]; }

void insert(char *word, int pos = 0) { 6 }
if(word[pos] == 0) { 7 }
words++;
return; }
       }
prefixes++;
int index = getIndex(word[pos]);
if(edges[index] == 0)
  edges[index] = new Trie;
                                                                                                             10
                                                                                                              \frac{11}{12}
        edges[index]->insert(word, pos+1);
                                                                                                              13
    int countWords(char *word, int pos = 0) {
                                                                                                              15
        if(word[pos] == 0)
                                                                                                             16
        return words;
int index = getIndex(word[pos]);
if(edges[index]==0)
                                                                                                             17
18
19
                                                                                                             20
        return edges[index]->countWords(word, pos+1);
    int countPrefix(char *word, int pos = 0) {
       if(word[pos] == 0)
  return prefixes;
int index = getIndex(word[pos]);
        if(edges[index] == 0)
        return edges[index]->countPrefix(word, pos+1);
                                                                                                             30
                                                                                                            32
33
                                                                                                             34
struct UnionFindDS {
    vi tree;
UnionFindDS(int n) { FOR(i, 0, n) tree.pb(i); }
    int root(int i) { return tree[i] == i ? i : tree[i] = 38 root(tree[i]); }

bool connected(int i, int j) {return root(i) == root(i); }
    void connect(int i, int j) { tree[root(i)] = tree[ 41
                 root(j)]; }
                                                                                                             \frac{42}{43}
struct UnionFindDS2 {
    vi tree, sizes;
int N;
    UnionFindDS2(int n) : N(n) {
                                                                                                             44
        tree.reserve(n);
FOR(i, 0, n) tree[i] = i;
sizes.assign(n, 1);
                                                                                                             45
                                                                                                             46
    int root(int i) { return (tree[i] == i) ? i : (tree[i48]
    int root(int i) { return (tree[i] == 1) ? 1 : (tree[i] 49
    ] = root(tree[i]));
int countSets() { return N;}
int getSize(int i) { return sizes[root(i)];}
bool connected(int i, int j) { return root(i) == root50
    (j);}
void connect(int i, int j) {
        int ri = root(i), rj = root(j);
if(ri != rj) {
                                                                                                             54
                                                                                                             55
            N--;
sizes[rj] += sizes[ri];
tree[ri] = rj;
                LCA
     int minIndex(int i, int j) {
  return height[i] < height[j] ? i : j;</pre>
   }
LCA(Graph &g, ll root) {
   index.assign(g.V, -1);
   dfs(g, root, 0);
   st.assign(height.size()*2, 0);
   FOR(i, 0, height.size())
   st[height.size() + i] = i;
   for(int i = height.size()-1; i; i--)
   st[i] = minIndex(st[i<<1], st[i<<1|1]);</pre>
                                                                                                             10
   }
woid dfs(Graph &g, 11 cv, 11 h) {
  index[cv] = order.size();
  order.pb(cv), height.pb(h);
  FORC(g.edges[cv], edge)
  if(index[edge->to] == -1) {
      dfs(g, edge->to, height.back() + 1);
      order.pb(cv), height.pb(h);
  }
                                                                                                             15
                                                                                                             \frac{16}{17}
                                                                                                             19
                                                                                                             22
                                                                                                             23
```

11 HLD

```
struct HeavyLightDecomposition {
    vector<vi>lists;
vi values, listIndex, posIndex, parent, treeSizes;
vector<SparseTable> sts;
    LCA *lca;
    HeavyLightDecomposition(Graph &g, vi values) : values
        (values) {
lca = new LCA(g, 0);
listIndex = posIndex = parent = treeSizes = vi(g.V,
                        -1):
        getTreeSizes(g, 0);
makeLists(g, 0, -1);
FORC(lists, list) {
            vi v;
FORC(*list, it) v.pb(values[*it]);
sts.pb(SparseTable(v));
    "HeavyLightDecomposition() { delete lca; } int getTreeSizes(Graph &g, int cv) {
       nt getTreeSizes(Graph &g, int cv) {
treeSizes[cv] = 1;
FORC(g.edges[cv], edge) if(edge->to != parent[cv])
    parent[edge->to] = cv, treeSizes[cv] +=
        getTreeSizes(g, edge->to);
return treeSizes[cv];
    void makeLists(Graph &g, int cv, int listNum) {
       if(listNum == -1)
  listNum = lists.size(), lists.pb(vi());
listIndex[cv] = listNum;
posIndex[cv] = listS[listNum].size();
         lists[listNum].pb(cv);
       Insts(listNum, po(cv);
int MAX = -1;
FORC(g.edges[cv], edge) if(edge->to != parent[cv])
if(MAX = -1 || treeSizes[edge->to] > treeSizes[
MAX]) MAX = edge->to;
FORC(g.edges[cv], edge) if(edge->to != parent[cv])
makeLists(g, edge->to, edge->to == MAX ? listNum
: -1);
    if(posRight)]];
if(posRight!= posLeft)
result = min(result, values[lists[listIndex[anc]]][sts[listIndex[anc]].query(posRight+1,
        posLeft)]]);
return result;
    int gueryToAncestor(int from, int anc, int &
        posInAncestor(int irom, int and, int a
posInAncestorList) {
  int result = INF, left = from;
  while(listIndex[left] != listIndex[and]) {
    result = min(result, values[lists[listIndex[left]]);
            | Test = min(restr, values[ists]istindex[eft]
| [ists[istIndex[eft]].query(0, posIndex[left])]];
| left = parent[lists[istIndex[left]][0]];
            osInAncestorList = posIndex[left];
};
```

Edmonds Karp

30

31

32

33

```
/* Edmonds-Karp
Finds a the maxflow from source to sink of a directed
          araph.
The weight of an edge denotes the capacity of the edge. The negative weight edges are the edges with flow.
return minEdge;
if [parent[cv] != -1) {
    flow = augment(g, flow, parent, source, parent[cv],
        min(minEdge, g.edges[parent[cv]][cv].weight
     g.edges[parent[cv]][cv].weight -= flow;
g.edges[cv][parent[cv]].weight += flow;
   return flow:
int maxFlow(MatrixGraph &g, int source, int sink) {
   int mf = 0, flow = -1;
while(flow) {
  vi distanceTo(g.V, INF);
      distanceTo[source] = 0:
      astance(o)source] = 0;
queue<int> q; q.push(source);
vi parent(g.V, -1);
while(!q.empty()) {
  int cv = q.front(); q.pop();
  if(cv == sink) break;
            if(g.edges[cv][i].weight > 0 && distanceTo[i]
                     == INF)
               distanceTo[i] = distanceTo[cv] + 1, q.push(i)
                        , parent[i] = cv;
      mf += flow = augment(g, 0, parent, source, sink,
```

13 Max Bipartite Matching $\frac{68}{70}$

return mf;

```
struct MaxBipartiteMatching {
       vvi edgesL;
4
5
10
12
         return 0:
13
14
15
         int ans = 0;
         16
17
18
19
20
       void augment2(int 1) {
22
         inCoverL[1] = 0;

for (auto r: edgesL[1])
23
\frac{24}{25}
           if (!inCoverR[r]) inCoverR[r] = 1, augment2(
26
                 matchR[r]);
       void minCover() { // assuming matching found
28
         inCoverL.assign(L, 1), inCoverR.assign(R, 0);
for(int i = 0; i < L; i++)
   if (matchL[i] == -1) augment2(i);</pre>
29
30
31
32
```

14 Matrices

```
typedef vector<vector<double> > Matrix;
#define EPS 1E-7
       #define CREATE(R, C) Matrix(R, vector<double>(C));
       Matrix identity(int n) {
  Matrix m = CREATE(n, n);
  FOR(i, 0, n)
  m[i][i] = 1;
                                                                                   19
12
       Matrix multiply (Matrix m, double k) {
          FOR(i, 0, m.size())

FOR(j, 0, m[0].size())

m[i][j] *= k;
15
16
       Matrix multiply (Matrix ml. Matrix m2) {
19
          20
23
24
25
26
27
          return result;
28
       Matrix pow(Matrix m, int exp) {
30
          if(!exp) return identity(m.size());
if(exp == 1) return m;
Matrix result = identity(m.size());
31
          while(exp) {
  if(exp & 1) result = multiply(result, m);
  m = multiply(m, m);
  exp >>= 1;
34
35
36
37
                                                                                   14
38
39
          return result:
40
41
                                                                                   17
       //solves AX=B, output: A^-1 in A, X in B, returns det(A
42
        double gaussJordan(Matrix &a, Matrix &b) {
         19
45
46
                                                                                   22
47
48
                                                                                   23
49
                                                                                   26
52
54
                                                                                   31
32
55
56
57
             swap(a[p]], a[pk]);
swap(b[p]], b[pk]);
if (pj != pk) det *= -1;
irow[i] = pj;
icol[i] = pk;
                                                                                   33
58
59
             double c = 1.0 / a[pk][pk];
             det *= a[pk][pk];
a[pk][pk] = 1.0;
FOR(p, 0, n) a[pk][p] *= c;
FOR(p, 0, m) b[pk][p] *= c;
61
62
```

```
FOR(p, 0, n) if (p != pk) {
       c = a[p][pk];
a[p][pk] = 0;
FOR(q, 0, n) a[p][q] -= a[pk][q] * c;
FOR(q, 0, m) b[p][q] -= b[pk][q] * c;
  for(int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) 6
     FOR(k, 0, n) swap(a[k][irow[p]], a[k][icol[p]]);
//returns the rank of a
int rref(Matrix &a) {
  int n = a.size(), m = a[0].size();
int r = 0;
  FOR(c, 0, m) {
   int j = r;
   FOR(i, r+1, n)
     if (abs(a[i][c]) > abs(a[j][c])) j = i;
if (abs(a[j][c]) < EPS) continue;</pre>
                                                                            20
21
     22
                                                                           26
                                                                           27
28
                                                                           29
  return r;
                                                                           30
                                                                           33
```

15 Dates

66

93 94

16 Articulation Points/Bridges

```
/* Articulation Points
O(V+E)
Finds all articulation points and bridges in a graph.
An articulation point is a vertex whose removal would disconnect the graph.

An bridge is a vertex whose removal disconnects the
15
        if(v == root) rootChildren++;
dfs1(g, edge->to);
if(low2[edge->to] >= num2[v]) strongPoints[v] =
                 true;
        true;
if(low2[edge->to] > num2[v]) edge->strong = g.
edges[edge->to][edge->backEdge].strong =
true;
low2[v] = min(low2[v], low2[edge->to]);
      } else if(edge->to != parent[v])
         low2[v] = min(low2[v], num2[edge->to]);
                                                                               27
}
                                                                              29
30
vi articulationPointsAndBridges(Graph &g) {
  counter2 = 0;

num2 = vi(g.V, -1), low2 = vi(g.V, 0), parent = vi(g.32)

20:100 - vi(g.V, 0):
  V, -1), strongPoints = vi(g.V, 0);
FOR(i, 0, g.V)
if(num2[i] == -1) {
                                                                               34
                                                                               35
         root = i, rootChildren = 0:
        dfsl(g, i);
strongPoints[root] = rootChildren > 1;
  return strongPoints;
                                                                               41
            SSC
                                                                               45
```

```
/* Strongly Connected Components
Partitions the vertices of a directed graph into
strongly connected components.
A strongly connected component is a subset of a graph
           where every vertex is reachable from every other
            vertex.

V where V_i is the index of the component of
          node i.
vi low1, num1, components;
int counter1, SCCindex;
vector<br/>bool> visited;
stack<int> S:
void dfs (Graph &g, int cv) {
   low1[cv] = num1[cv] = counter1++;
    S.push(cv);
   S.push(cv);
visited[cv] = true;
FORC(g.edges[cv], edge) {
  if(numl[edge->to] == -1)
      ifinim[edge=>t0] -- -1)
dfs(g, edge=>t0);
if(visited[edge=>t0])
lowl[cv] = min(lowl[cv], lowl[edge=>t0]);
   if(low1[cv] == num1[cv]) {
  int index = SCCindex++;
  while(true) {
    int v = S.top(); S.pop(); visited[v] = 0;
    components[v] = index;
  if (cv == v)
    break.
            break;
vi stronglyConnectedComponents(Graph &g) {
  counter1 = 0, SCCindex = 0;
  visited = vector<bool>(g.V, 0);
   FOR(i, 0, g.V)

if(num1[i] == -1)
          dfs(g, i);
```

18 Catalan Numbers

```
int fact(int n) {
    return n ? n*fact(n-1) : 1;
}
int nthCatalan(int n) {
    return fact(2*n)/(pow(fact(n), 2)*(n+1));
}
int nextCatalan(int n, int previous) {
    return previous*2*(2*n+1)/(n+2);
```

19 Euclid

34 35

```
/∗ GCD
int gcd(int a, int b) {
  int tmp;
while(b){a%=b; tmp=a; a=b; b=tmp;}
  return a;
int lcm(int a. int b) {
/* Extended Euclid
Finds x,y such that d = ax + by.

Returns d = gcd(a,b).
int extended_euclid(int a, int b, int &x, int &y) {
  int xx = y = 0;
int yy = x = 1;
while (b) {
     int q = a/b;

int t = b; b = a%b; a = t;

t = xx; xx = x-q*xx; x = t;

t = yy; yy = y-q*yy; y = t;
/* Modular Linaer Equation Solver
Finds all solutions to ax = b (mod n)
vi modular_linear_equation_solver(int a, int b, int n)
   vi solutions;
   int d = extended_euclid(a, n, x, y);
if (!(b%d)) {
   x = mod (x*(b/d), n);
     FOR(i, 0, d)
         solutions.pb(mod(x + i*(n/d), n));
  return solutions;
/* Modular Inverse
```

```
Computes b such that ab = 1 \pmod{n}, returns -1 on
48
                           failure
                                                                                                                                       36
            int mod_inverse(int a, int n) {
                int x, y;
int d = extended_euclid(a, n, x, y);
if (d > 1) return -1;
return mod(x,n);
52
53
54
55
56
57
58
59
           /* Chinese Remainder Theorem Returns \[ [x = a_i (mod n_i) \] n's must be pairwise coprimes
60
           */
int chinese_remainder(int *n, int *a, int len) {
  int p, i, prod = 1, sum = 0;
  for (i = 0; i < len; i++) prod *= n[i];
  for (i = 0; i < len; i++) {
    p = prod / n[i];
    sum += a[i] * mod_inverse(p, n[i]) * p;
}</pre>
63
64
67
                 return sum % prod;
```

20 Miller Rabin

/* Miller-Rabin Primality Test

```
O(log(N)^3)
          */
11 mulmod(11 a, 11 b, 11 c) {
    11 x = 0, y = a % c;
    while (b) {
        if (b & 1) x = (x + y) % c;
        y = (y << 1) % c;
                   b >>= 1;
10
11
12
13
           ll fastPow(11 x, 11 n, 11 MOD) {
14
15
16
17
18
19
              lastow(if x, if n, if nos);
the state of the if (n & 1) ret = mulmod(ret, x, MOD);
x = mulmod(x, x, MOD);
n >>= 1;
20
               return ret;
24
          bool isPrime(11 n) {
              11 d = n - 1;
int s = 0;
while (d % 2 == 0) {
25
28
                   s++;
d >>= 1;
29
30
31
                                                                                                                            22
               ^{\prime} // It's garanteed that these values will work for any ^{22}
               number smaller than 3*10**18 (3 and 18 zeros)24
int a[9] = { 2, 3, 5, 7, 11, 13, 17, 19, 23 };
FOR(i, 0, 9) {
  bool comp = fastPow(a[i], d, n) != 1;
  27
32
34
                   if(comp) FOR(j, 0, s) {
    11 fp = fastPow(a[i], (1LL << (11) j)*d, n);
    if (fp == n - 1) {
        comp = false;
    }
}</pre>
35
36
37
38
                                                                                                                            29
39
                          break:
40
41
                                                                                                                            31
                                                                                                                            32
                   if (comp) return false;
42
43
                                                                                                                            35
                                                                                                                            36
```

21 Eulerian Path

```
/* Eulerian Path
                                                                                                  42
         Partitions the vertices of a directed graph into
                   strongly connected components.
        strongly connected components. 44 A strongly connected component is a subset of a graph 45 where every vertex is reachable from every other
                     vertex.
        Returns V where V\_i is the index of the component of node i.
5
        vi low1, num1, components;
int counter1, SCCindex;
vector<br/>bool> visited;
         void dfs (Graph &g, int cv) {
12
            low1[cv] = num1[cv] = counter1++;
S.push(cv);
visited[cv] = true;
13
14
15
           FORC(g.edges[cv], edge) {
  if(numl[edge->to] == -1)
    dfs(g, edge->to);
  if(visited[edge->to])
16
17
18
19
20
                  low1[cv] = min(low1[cv], low1[edge->to]);
           if(low1[cv] == num1[cv]) {
  int index = SCCindex++;
23
24
25
26
               while (true) {
  int v = S.top(); S.pop(); visited[v] = 0;
  components[v] = index;
27
                  if (cv == v)
                                                                                                  17
28
                     break;
29
30
                                                                                                  20
31
                                                                                                  21
         vi stronglyConnectedComponents(Graph &g) {
```

```
counter1 = 0, SCCindex = 0;
  visited = vector<bool>(g.V, 0);
  numl = vi(g.V, -1), low1 = vi(g.V, 0), components = 26
      vi(g.V, 0);
  S = stack*int>(1);
  FOR(i, 0, g.V)
  29
  if(numl[i] == -1)
      dfs(g, i);
  return components;
  32
```

22 Nth Permutation

seg must be sorted

23 Shunting Yard

 $\begin{tabular}{lll} \begin{tabular}{lll} \begin$

out << x << " ";

For parsing mathematical expressions speci 12 field in infix notation

```
string readToken(istream &in) {
                                                  19
   28
#define RIGHT 1
                                                 29
30
#define isOp(x) (prec.find(x) != prec.end())
void shunting(istream &in, ostream &out)
    string token;
    stack<string> ops;
   \frac{11}{12}
       ops.push(token);
} else if(token == "(") {
          15
           output(out, ops.top()), ops.pop();

// ops.empty() || ops.top() != "(" ====
MISMATCH
                                                  19
       ops.pop();
} else // numbers vars
output(out, token);
   25
                                                 26
27
28
```

24 Sieve of Atkin

 $38 \\ 39 \\ 40 \\ 41$

```
isPrime[num] = true;
}

for (11 i = 5; i <= 1im; i++)
   if (isPrime[i])
   for (11 j = i * i; j <= n; j += i)
        isPrime[j] = false;

for (11 i = 2; i <= n; i++)
   if (isPrime[i])
   primes.pb(1);</pre>
```

25 KMP

26 Sparse Table

27 Fibonacci

```
/* Fibbonacci
*/
int fibn(int n) { //max 91
    double goldenRatio = (1+sqrt(5))/2;
    return round((pow(goldenRatio, n+1) - pow(1-
        goldenRatio, n+1))/sqrt(5));
}
int fibonacci(int n) {
    Matrix m = CREATE(2, 2);
    m[0][0] = 1, m[0][1] = 1, m[1][0] = 1, m[1][1] = 0;
    Matrix fib0 = CREATE(2, 1);
    fib0[0][0] = 1, fib0[1][0] = 1; //fib0 y fib1
    Matrix r = multiply(pow(m, n), fib0);
    return r[1][0];
}
```

28 Treap

1 2 3 4	<pre>#define LC(a) ((a) == ((a)->parent->left)) template<typename k=""> struct Treap { struct Node {</typename></pre>			Not		"222".	0. x)):	//base :	7 7 7 7 x_to long 8	7 #includ 8 #includ 9 using r	<pre>de <ext assoc_container.hpp="" pb_ds=""> de <ext pb_ds="" tree_policy.hpp=""> namespacegnu_pbds;</ext></ext></pre>
5 6 7 8 9	<pre>k key; int priority; Node *left. *right. *parent:}:</pre>	3 4 5	regmatch regcomp if (regen	_t matc ®, p ec(®	hes[1]; attern.c	_str(),		ENDED RE	EG_ICASE);	1 *X.find 2 X.order 3	<pre>F tree<int, less<int="" null_type,="">, rb_tree_tag, tree_order_statistics_node_update> ordered_set; jby_order(1); c_of_key(-5);</int,></pre>
10 11 12	<pre>bool a, b; while((a = (n->left && n->priority < n->left-></pre>	8 9 .0		<typen oString eturn ss</typen 	ame T> (T n) { s.str();	ostring	stream ss		- 8 8 8 8 8 8	5 start w 6 x_next 7 repeat 8 9 #define	Rhapson - Find a root of a polynomial with random x = x - f(x)/f'(x) until_f(x)_is_zero turnOffLastBit(S)_((S)_6_(S1))
13 14	<pre>rand()%2 : b); } Node *find(K key, bool leaf, Node *start, bool onlyLeft = 0) { if (!start) return 0;</pre>	3	T toNum	const s	tring &T	ss >> re	istringst esult ? r does not	result :	- 9	1 #define	$ \begin{array}{ll} & \text{g_turnOnLastZero} (S) \cup \{(S) \cup [\cup \{S_t + \bot\}\}) \\ & \text{g_turnOnfLastConsecutiveBits} (S) \cup \{(S) \cup \{\subseteq \{L, \bot\}\}\}) \\ & \text{g_turnOnLastConsecutiveZeroes} (S) \cup \{(S) \cup [\cup \{S_t - \bot\}\}\}) \\ & \text{g_turnOnLastConsecutiveZeroes} \end{array} $
15 16	Node *n = start, *next; while((next = (key < n->key onlyLeft ? n->left :	.5	upper_b	nan val. und: fi	nds firs	t that		greater	r than val	32	Formulas
17 18 19	return n; } void insert (K key, int priority = -1) { Node *p = find(key, l, root), *p = new Node { key. } 1	7		ile(nex	t_permut	ation(a	rr, arr+1	N));	=		Catalan Numbers
20 21	priority == -1 ? rand()%100000 : priority , 2 0, 0, p }; if (!root) { root = n; return; } 2 (key < n->key ? n->left : n->right) = n:	1 1 12	scanf: %d -> ba %o -> ba %x -> ba	se8 int se16 in	%d+ t %d+					$C_n = \prod_{k=2}^n$	$\frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!} = \frac{n+k}{k}, n \ge 0$
22 23 24	<pre>while(p && n->priority > p->priority) rotate(p, !LC^L</pre>		5 P	164.324¢ receded	e+3, 5342 by 0x	23E+2, 0	ex. 123 0x242.435 ex. scar	, basel		$C_0 = 1$	$1, C_{n+1} = \frac{2(2n+1)}{n+2} C_n$
25 26	<pre>if (!root) return; Node *n = find(key, 0, root), *del = n; if (!n n->key != key) return;</pre>		n	-> 'a', ot inclu	scanf("	%4c", my	ycharptr)	-> "as	df" (\0	32.2	Law of Cosines
27 28 29 30	<pre>if (n->left && n->right) { del = find(key, 1, n->right, 1); n->key = del->key, n->priority = del->priority; find(print(n), 1);</pre>	6	a	: scanf scz24_al o match	oc345" w	ould sto		in myint	, use %%	$c^2 = a$	$^2 + b^2 - 2ab * cos(C)$
31	<pre>if (del->left del->right) (del->left ? del->left²</pre>		%3d mean	s match naracte	an inte	ger but	t store i read onl %d match	ly the 3		32.3	Law of Sines
33 34	->parent->right) = del->left ? del->left : del->right; else root = del->left ? del->left : del->right; delete del; } 3 3	1	%le long			a system	 m of inea		- es of the	$\frac{a}{\sin(A)}$	$\frac{1}{1} = \frac{b}{\sin(B)}$
35 36 37	<pre>void rotate(Node *n, bool left) { Node *u = (left ? n->right : n->left), *p = n-> parent; if (n) (LC(n) ? n->left : n->right) = u: 3.</pre>	3 4	create a	pe x_i node f source	- x_j <: or every node	= C					Newton Raphson $f(x_0)$
38 39 40 41	<pre>else root = u; Node *c = left ? u->left : u->right; (left ? n->right : n->left) = c; 3</pre>	6	for ever	y inequ to jos	ality x_ f weight	i - x_j c			er node edge from		$=x_n - \frac{f(x_0)}{f'(x_0)}$
42 43	<pre>if (c) (left ? u->left : u->right)->parent = n; (left ? u->left : u->right) = n; u->parent = p, n->parent = u;)</pre>			was a	negative		cycle, t	the syst	tem is	32.5	Arithmetic Series
44); 4 4			nconsist artite		he size	of the n		-	$\sum_{k=1}^{n}$	$(a_1 + (k-1)d) = na_1 + \frac{1}{2}nd(n-1)$
	29 LIS 4	2	In a big	ndepende	ent set graph, t	he size	of the m	min vert	tex cover	32.6	Geometric Series
1	4. 4.	3	In a big	ndepende artite MCBM 	ent set graph, t bcsdfk	he size		min vert	ex cover		
2 3 4	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi longest(visize(), -1); 4	3 4 5 6	in a bip char str char * t while (t	ndepende artite MCBM [] = "a 0564567 oken = oken != tf ("%s	ent set graph, t bcsdfk"; strtok(s NULL) { \n",toke	he size sm_sgfd tr, ""	of the m	min vert	ex cover	$\sum_{k=1}^{n}$	Geometric Series $r^k = \frac{r(1-r^n)}{1-r}$ $r^k = \frac{r}{1-r}, r < 1$
2	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi parent(v.size(), -1); FOR(i, 0, v.size()) { ii tem = ii(v[i], i); vii::iterator it = upper_bound(best.begin(), best. 5	.3 .4 .5 .6 .7 .8 .9	in a bip char str h char * t while (t prin toke }	ndepende artite MCBM [] = "a 5564567 oken = oken != tf ("%s n = str	ent set graph, t bcsdfk"; strtok(s NULL) { \n",toke tok (NULL	he size sm_sgfd. tr, "" n); L, ""	of the n a_afdex ");	min vert	ex cover	$\sum_{k=1}^{n}$	$r^{k} = \frac{r(1-r^{n})}{r^{k}}$ $r^{k} = \frac{r}{1-r}, r < 1$
2 3 4 5 6 7	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi parent(v.size(), -1); FOR(i, 0, v.size()) { ii tiem = ii(v(i), i); vii::iterator it = upper_bound(best.begin(), best. 5	3 4 4 5 6 6 7 8 8 9 0 0	in a bip	ndepende artite MCBM 	ent set graph, t "; strtok(s NULL) { \n", toke tok (NUL ieces in its ci three ii nCat2 +	he size sm_sgfddtr, "" n); L, "" which rrcumferenternall	of the magardex "););	min vert NJK, is divi joined crent:	- ided if n by chords	$\Sigma_{k=1}^n$ $\Sigma_{k=1}^\infty$ 32.7	$r^k = \frac{r(1-r^n)}{1-r}$ $r^k = \frac{r}{1-r}, r < 1$
2 3 4 5 6 7 8 9 10 11 12 13	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi parent(v.size(), -1); FOR(i, 0, v.size()) { ii item = ii(v[i], i); vii::iterator it = upper_bound(best.begin(), best.5; if (it == best.end()) { parent[i] = (best.size() == 0 ? -1 : best.back() 5;	3 4 5 6 7 8 9 6 0 1 1	in a big 	mdepende artite MCBM - [] = "a 5564567. oken = oken != tf ("%s n = str er of p obints or with no Cat4 + 2-1 whe mmber of	ent set graph, t bcsdfk "; strtok(s NULL) { \n", toke tok (NUL ieces in n its ci: three in ncat2 + re A is f intege: per of in panning	he size sm_sgfd. tr, "" n); L, "" which internal l the area r points anteger p trees in	of the m	min vert NJK, is divi joined rrent: olygon, polygon the bo te a big	ided if n by chords	$\Sigma_{k=1}^{n}$ $\Sigma_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x)$ 32.8	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $dx pprox rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ $Stirling's Approximation$
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {	3 44 5 6 6 7 8 9 9 10 11 12 3 3	in a big	mdepende artite MCBM [] = "a 5564567. oken = e oken != tf ("%s n = str er of p points on with no Cat4 + 2-1 whe imber of s the numi er of s raph K(i	ent set graph, t "; strtok(s NULL) { \n",toke tok (NUL "; ieces in n its ci: three in nCat2 + re A is f intege: per of in panning n, m) is	he size sm_sgfd. tr, "" n); L, "" which internal l the area r points anteger p trees in	of the m	min vert NJK, is divi joined rrent: olygon, polygon the bo te a big	ided if n by chords	$\Sigma_{k=1}^{n}$ $\Sigma_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x)$ 32.8	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $dx pprox rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$
2 3 4 5 6 7 8 9 10 11 12 13 14 15	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {	3 3 4 4 5 5 6 6 6 7 7 8 9 9 0 0 1 1 2 2 3 3 4 4 5 5 6 6 6 7 7 8 8	i In a big	ndepende artitle MCBM	part set graph, t bcsdfk "; snuLL) { \n\n", toke tok (NUL ieces in n its ci three in nCat2 + re A is f intege: per of in panning n, m) is ss(line) tokens; eratorks	he size sm_sgfd. tr, " n); L, " which recumferenternal l the are. r points nteger F trees in m^ (n-1) ; tring>(of the m a_afdex ");); a circle ence are my concur a of a poss on the coints on n complet on the coints on n complet on the coints on n complete on the coints on the coints on the coints on n complete on the coints on n complete on the coints on n coints on the coints on the coints on the coints on the coints of the	min vert	dided if n by chords i is the and b is undary partite	$\sum_{k=1}^{n}$ $\sum_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x)$ 32.8 $n! \approx y$	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $dx pprox rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ $Stirling's Approximation$
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {</pre>	3 4 4 5 5 6 6 7 7 8 8 9 9 0 0 1 1 2 2 3 3 4 4 4 5 5 6 6 7 7 8 8 9 9 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	in a bip char sti h char * t while (t print toke } The numm F g(n) = n A = i+b, n the numl //splittistring vector copy(ist	ndepende artite MCBM [] = "actite MCBM	part set graph, t bcsdfk "; strtok(sNULL) {	he size sm_sgfd. tr, " n); L, " which recumferenternal l the are. r points nteger F trees in m^ (n-1) ; tring>(of the m a_afdex ");); a circle ence are my concur a of a poss on the coints on n complet on the coints on n complet on the coints on n complete on the coints on the coints on the coints on n complete on the coints on n complete on the coints on n coints on the coints on the coints on the coints on the coints of the	min vert	ided if n by chords i is the and b is undary partite	$\sum_{k=1}^{n}$ $\sum_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x)$ 32.8 $n! \approx y$	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $dx \approx rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ $Stirling's Approximation$ $\sqrt{2\pi n}(rac{n}{e})^n$ $= n * ln(n) - n + rac{ln(2n)}{2}$
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 20 21	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {</pre>	3 4 4 5 5 6 6 6 7 7 8 8 9 9 10 1 1 2 2 3 3 4 4 5 5 6 6 7 7 8 8 9 9 10 1 1 2 2 3 1 1 1 2 2 1 1 1 1 1 1 1 1 1 1	in a bip char sti h char * t while (t prir toke } The numl A = i+b, the numl //splittistring: vector <s 2="" 3<="" copy(ist="" td=""><td>depends aratite MCBM [[] = "a aratite MCBM [[] = "a booken = 0 oken != 1 oken = 0 oken = 0 oken = 0 oken != 1 oken = 0 oken</td><td>nt set graph, t bcsdfk "; strtok(s NULL) { \n", toke tok (NUL ieces in its cir three ii ncat2 + re A is f intege: panning n, m) is spaces ss(line) tokens; perators , back</td><td>he size sm_sgfd. tr, " n); L, "" which recumferenternall lather are repoints the are repoints the ger from the company of the comp</td><td>of the m a afdex ");); a circle snce are ly concur a of a pc s on the coints on n complete) * n^(m- iss), istr</td><td>min vert NJK, is divi joined rement: olygon, polygon the bo te a big -1) tream_it <string></string></td><td>ided if n by chords i is the and b is undary partite </td><td>$\Sigma_{k=1}^{n}$ $\Sigma_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x) dx$ 32.8 $n! \approx \sqrt{\ln(n!)}$ 32.9 $\Sigma_{k=1}^{n}$</td><td>$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $stirling's Approximation$ $stirling's Approximation$</td></s>	depends aratite MCBM [[] = "a aratite MCBM [[] = "a booken = 0 oken != 1 oken = 0 oken = 0 oken = 0 oken != 1 oken = 0 oken	nt set graph, t bcsdfk "; strtok(s NULL) { \n", toke tok (NUL ieces in its cir three ii ncat2 + re A is f intege: panning n, m) is spaces ss(line) tokens; perators , back	he size sm_sgfd. tr, " n); L, "" which recumferenternall lather are repoints the are repoints the ger from the company of the comp	of the m a afdex ");); a circle snce are ly concur a of a pc s on the coints on n complete) * n^(m- iss), istr	min vert NJK, is divi joined rement: olygon, polygon the bo te a big -1) tream_it <string></string>	ided if n by chords i is the and b is undary partite	$\Sigma_{k=1}^{n}$ $\Sigma_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x) dx$ 32.8 $n! \approx \sqrt{\ln(n!)}$ 32.9 $\Sigma_{k=1}^{n}$	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $stirling's Approximation$
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {	3 4 4 5 6 6 6 7 7 8 8 9 9 0 0 1 1 2 2 3 3 4 4 5 5 6 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 3 4 4 5 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 3 4 5 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 2 2 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 2 2 3 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 2 2 3 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 2 2 3 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 2 2 3 3 3 5 6 6 7 7 8 8 9 9 9 6 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	in a bip	depends artitle MCBM [] = "a artitle MCBM [] = "a artitle MCBM] [] = "a e dependence of the model of the mo	part set graph, t control of the con	he size sm_sgfd. tr, " n); Which roumferentler repoints the arear the arear trees in m^(n-1); tring>(!inserter 11 1 59	of the management of the manag	min vert NJK, is divi joined rrent: colygon, polygon the bo te a bir -1) tream_it <string> 19 71</string>		$\Sigma_{k=1}^{n}$ $\Sigma_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x) dx$ 32.8 $n! \approx \sqrt{\ln(n!)}$ 32.9 $\Sigma_{k=1}^{n}$	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ $Simpson's Rule$ $dx pprox rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ $Stirling's Approximation$ $\sqrt{2\pi n}(rac{n}{e})^n$ $= n * ln(n) - n + rac{ln(2n)}{2}$ $Sum of Powers$
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2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 22 1 22 2 2 2 2 5 6 7 8	/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {	3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.0.1.2.3.3.4.4.5.5.6.6.7.8.9.9.0.0.1.2.3.3.4.4.5.5.6.6.7.8.9.0.0.1.2.3.3.4.4.5.5.6.6.7.8.9.0.0.1.2.3.3.4.5.5.6.0.0.2.2.2.2.2.2.2.2.2.2.2.2.2.2.2.2.2	in a big	depends depe	nt set graph, t bcsdfk"; strtok(s NULL) { \(\) n", toke \(\) tok	he size sm_sgfd. tr, " n); L, " which roumferenternall the arear points nteger F trees in m^(n-1) ; tring>(inserter) 11 1: 59 109 1: 179 241 313 389 461	of the m of the m a afdex ");); a circle ence are ly concur a of a po s on the points on n complete) * n^(m iss), ist 3 17 13 127 181 191 251 251 317 331 3397 401	min vert min vert NJK, is divi joined crent: colygon, polygon the bote a bir -1) 19 71 131 1 193 7 263 1 337 1 409 7 479	ided if n by chords i is the and b is undary partite	$\sum_{k=1}^{n} \sum_{k=1}^{\infty} 32.7$ $\int_{a}^{b} f(x) dx$ 32.8 $n! \approx \sqrt{10}$ $\sum_{k=1}^{n} \sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 32.1$	$r^{k} = \frac{r(1-r^{n})}{1-r}$ $r^{k} = \frac{r}{1-r}, r < 1$ Simpson's Rule $dx \approx \frac{b-a}{6}(f(a) + 4f(\frac{a+b}{2}) + f(b))$ Stirling's Approximation $\sqrt{2\pi n}(\frac{n}{e})^{n}$ $= n * ln(n) - n + \frac{ln(2n)}{2}$ Sum of Powers $k^{2} = \frac{1}{6}n(n+1)(2n+1)$ $k^{3} = (\sum_{k=1}^{n} k)^{2} = (\frac{1}{2}n(n+1))^{2}$ O Fermat's little Theorem $dx = (mod \ p)$ where p is prime $dx = 1 \pmod{p}$ where p is prime and a challenge of the sum
2 3 4 4 5 6 6 7 8 9 10 11 12 13 14 15 16 17 8 19 20 21 22 2 2 2 1 1 2 2 3 4 4 5 6 6 7 8 9 9 10 11 12 13 14 14	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {</pre>	3 4 4 5 5 6 6 7 7 8 8 9 9 0 1 1 2 3 3 4 4 5 6 6 7 7 8 8 9 9 0 1 1 2 3 3 4 4 5 6 6 7 7 8 8 9 9 0 1 1 2 3 3 4 4 5 6 6 7 7 8 8 9 9 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	in a bip		ent set graph, t bcsdfk"; strtok(sc NULL) { \n", toke tock (NULL) { \n", toke tock (NULL) { \n", toke toke (NULL) { \n", toke \n',	he size sm_sgfd. tr, " n); L, " which recumferentle the area recorded to make the size of the	of the m	min vert min vert NJK, is divijoined rent: olygon, polygon the botte a bir 1) tream_it* (string) 71 131 1 193 7 263 1 337 1 409 7 479 3 569 1 641	ided if n by chords i is the and b is undary partite	$\sum_{k=1}^{n} \sum_{k=1}^{\infty}$ 32.7 $\int_{a}^{b} f(x) dx$ 32.8 $x_{k}^{n} = x_{k}^{n}$ $\sum_{k=1}^{n} \sum_{k=1}^{n} dx$ 32.1 $\sum_{k=1}^{n} dx$	$r^{k} = \frac{r(1-r^{n})}{1-r}$ $r^{k} = \frac{r}{1-r}, r < 1$ Simpson's Rule $dx \approx \frac{b-a}{6}(f(a) + 4f(\frac{a+b}{2}) + f(b))$ Stirling's Approximation $\sqrt{2\pi n}(\frac{n}{e})^{n}$ $= n * ln(n) - n + \frac{ln(2n)}{2}$ Sum of Powers $k^{2} = \frac{1}{6}n(n+1)(2n+1)$ $k^{3} = (\sum_{k=1}^{n} k)^{2} = (\frac{1}{2}n(n+1))^{2}$ O Fermat's little Theorem $a \pmod{p}$ where p is prime $a \pmod{p}$ where p is prime and a cidivisible by p
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 22 1 22 3 3 4 5 6 7 8 9 9 9 9 10 10 10 11 11 12 12 12 12 12 12 12 12 12 12 12	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) {</pre>	3 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 5 6 6 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 5 6 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 6 7 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 8 9 9 0 1 1 2 2 3 3 4 4 5 7 8 9 9 9 0 1 1 2 2 3 3 4 4 5 7 8 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9	in a big		nt set graph, t bcsdfk"; strtok(s NULL) { \n",toke tok (NULL) { \n',toke tou	he size sm_sgfd. tr, " n); L, " which rounderer montherer the are r points. tring>(n); tring>(in) tring>(in) 1 1 1 1 1 1 1 1 1 1 1 1 1	of the m of the m a aftex ");); a circle ence are ly concur a of a po s on the points on n complete) * n^ (m- iss), ist 3 17 61 67 13 127 181 191 251 257 317 331 397 401 463 467 557 566 619 631 701 708	min vert NJK, is divi joined crent: olygon, polygon the bo te a bir 19 71 131 1 193 7 263 1 337 1 409 7 479 3 569 1 641 9 719	dided if n by chords i is the and b is bundary partite - cerator< >(tokens - 23 29 73 137 197 269 347 419 487 571 643 727	$\sum_{k=1}^{n}$ $\sum_{k=1}^{\infty}$ $\sum_{k=1}^{\infty}$ $\sum_{k=1}^{n}$	$r^{k} = \frac{r(1-r^{n})}{1-r}$ $r^{k} = \frac{r}{1-r}, r < 1$ Simpson's Rule $dx \approx \frac{b-a}{6}(f(a) + 4f(\frac{a+b}{2}) + f(b))$ Stirling's Approximation $\sqrt{2\pi n}(\frac{n}{e})^{n}$ $= n * ln(n) - n + \frac{ln(2n)}{2}$ Sum of Powers $k^{2} = \frac{1}{6}n(n+1)(2n+1)$ $k^{3} = (\sum_{k=1}^{n} k)^{2} = (\frac{1}{2}n(n+1))^{2}$ O Fermat's little Theorem $dx = (mod \ p)$ where p is prime $dx = 1 \pmod{p}$ where p is prime and a challenge of the sum
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 8 9 10 11 12 12 12 12 13 14 15 16 17 18 19 10 11 12 13 14 15 16 17 18 19 19	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi parent(v.size(), -1); FOR(i, 0, v.size()) { ii item = ii(v[i], i); vii::iterator it = upper_bound(best.begin(), best.5; end(), item); if (it == best.end()) { parent[i] = (best.size() == 0 ? -1 : best.back() .5; second); best.pb(item); } else { parent[i] = parent[it->second]; *it = item; } vi lis; for(int i=best.back().second; i >= 0; i=parent[i]) 5; ilis,pb(v[i]); reverse(lis.begin(), lis.end()); reverse(lis.begin(), lis.end()); return lis; } 60 60 60 60 61 60 61 60 61 62 63 63 63 64 65 66 66 66 67 66 67 68 69 69 60 60 60 60 60 60 60 60</pre>	3.4.4.5.5.6.7.8.9.00.1.2.3.4.4.5.6.0.0.1.2.3.4.4.5.6.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0	in a bip		ent set graph, to be signaph, to be signaph, to be signaph, to be signaph. The signaph	he size sm_sgfd. tr, " n); L, "" which recumferentle reports the area reports trees in m^ (n-1) ; tring>(inserter 11 1 59 109 1 179 241 313 389 461 547 691 773	of the m	min vert min vert NJK, is divi joined rrent: olygon, polygon the bo te a bir -1) tream_it (string> 71 131 1 193 7 263 1 337 1 409 7 479 3 569 1 641 9 719 7 809	ided if n by chords i is the and b is undary partite	$\sum_{k=1}^{n} \sum_{k=1}^{\infty} 32.7$ $\int_{a}^{b} f(x) dx$ 32.8 $x_{k}^{n} \approx \sqrt{1}$ $\sum_{k=1}^{n} \sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} 22.1$ $\sum_{k=1}^{n} 22.1$ $\sum_{k=1}^{n} 22.1$ $\sum_{k=1}^{n} 22.1$ $\sum_{k=1}^{n} 22.1$ $\sum_{k=1}^{n} 22.1$	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ Simpson's Rule $0dx \approx rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ Stirling's Approximation $2\pi n(rac{n}{c})^n$ $= n*ln(n) - n + rac{ln(2n)}{2}$ Sum of Powers $k^2 = rac{1}{6}n(n+1)(2n+1)$ $k^3 = (\sum_{k=1}^n k)^2 = (rac{1}{2}n(n+1))^2$ O Fermat's little Theorem $n(mod\ p)$ where p is prime $n(mod\ p)$ where p is prime and a cidivisible by p 1 Euler's Totient Function $n(n) = n\Pi_{p n}(1 - rac{1}{p})$ where p is prime
2 3 4 5 6 7 8 9 10 111 12 13 14 15 16 17 18 19 20 22 2 2 2 2 2 3 4 5 6 7 8 9 9 10 11 12 13 14 15 16 17 18 17 18 17 18 18 19 19 19 19 19 19 19 19 19 19 19 19 19	<pre>/* Longest Increasing Subsequence */ vi longestIncreasingSubsequence(vi v) { vii best; vi parent(v.size(), -1); FOR(i, 0, v.size()) { ii ttem = ii(v[i], i); vii:tierator it = upper_bound(best.begin(), best.5; end(), item); if (it == best.end()) { parent[i] = (best.size() == 0 ? -1 : best.back() 5; second); best.ph(item); } else { parent[i] = parent[it->second]; ii; for(int i=best.back().second; i >= 0; i=parent[i]) ilis,pb(v[i]); reverse(lis.begin(), lis.end()); return lis; } 30 Kadane 6 **Maximum Subarray */ int maximumSubarray(int numbers[], int N) { int maxSoFar = numbers[0], maxEndingHere = numbers [0]; FOR(i, l, N) { if(maxEndingHere < 0) maxEndingHere = numbers[i]; else maxEndingHere + numbers[i]; else maxEndingH</pre>	3.4.4.5.5.6.6.7.8.9.9.0.1.1.2.3.4.4.5.6.6.7.8.9.9.0.1.1.2.3.4.4.5.6.6.7.8.9.9.0.1.1.2.3.4.4.5.6.6.7.8.9.9.0.1.2.3.4.4.5.6.0.7.8.9.9.0.1.2.3.4.4.5.6.0.7.8.9.9.0.1.2.3.4.4.5.0.0.1.2.3.4.4.5.0.0.0.1.2.3.4.4.5.0.0.0.1.2.3.4.4.5.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0	in a bip	"adaption "ada	nt set graph, t bcsdfk"; strtok(s NULL) { \n".toke tok (NULL) { \n".toke tok (NULL) { \n".toke \n".	he size sm_sgfd. tr, " n); L, "" which recumferenternall the area repoints. tring>(in many company co	of the m of the m a afdex ");); a circle ence are ly concur up of a po s on the points on n complete) * n^(m- iss), ist r 13 17 13 127 13 127 13 127 13 127 13 127 13 127 13 127 13 17 13 127 13 17 13 127 13 17 14 17 15 17 16 17 17 17 17 17 17 17 17 17 17 17 17 17 1	min vert min vert NJK, is divi joined rerent: colygon, polygon the bote a bir -1) -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1	ded if n by chords i is the and b is sundary partite	$\sum_{k=1}^{n} \sum_{k=1}^{\infty} 32.7$ $\int_{a}^{b} f(x) dx$ 32.8 $x_{k}^{n} = x_{k}^{n}$ 32.9 $\sum_{k=1}^{n} \sum_{k=1}^{n} 32.1$ $\sum_{k=1}^{n} a^{p-1} = a^{p-1}$ $\sum_{k=1}^{n} a^{p-1}$ \sum_{k	$r^k = rac{r(1-r^n)}{1-r}$ $r^k = rac{r}{1-r}, r < 1$ Simpson's Rule $0dx \approx rac{b-a}{6}(f(a) + 4f(rac{a+b}{2}) + f(b))$ Stirling's Approximation $2\pi n(rac{n}{e})^n$ $= n*ln(n) - n + rac{ln(2n)}{2}$ Sum of Powers $k^2 = rac{1}{6}n(n+1)(2n+1)$ $k^3 = (\sum_{k=1}^n k)^2 = (rac{1}{2}n(n+1))^2$ O Fermat's little Theorem $n(mod\ p)$ where p is prime $n(mod\ p)$ where p is prime and a cidivisible by p 1 Euler's Totient Function $n(n + n + n + n)$ $n(n + $

32.13 Convex Polygon Centroid₂

```
Given the polygon P = A_1, A_2, ..., A_n
let a = (A_{k+1} - A_1), k = 1, 2, ..., n - 1 (the edges
vectors)
let C = A_1 + \frac{1}{3}(a_k + a_{k+1}), k = 1, 2, ..., n - 2 (the 738)
centroids of the triangles)
let w = \frac{1}{2}(a_k \times a_{k+1}), k = 1, 2, ..., n-2 (the areas<sup>0</sup>)
of the triangles)
                               \frac{\sum_{k=1}^{n-2} w_k C_k}{\sum_{k=1}^{n-2} w_k}
                                                                              42
centroid
                                                                     A_1
\tfrac{1}{3} \tfrac{\sum_{k=1}^{n-2} (a_k + a_{k+1}) (a_k \times a_{k+1})}{\sum_{k=1}^{n-2} (a_k \times a_{k+1})}
                                                                              45
                                                                              46
                                                                              47
```

32.14 Regular Polyhedron Volume

```
volume = L^3
```

32.15 Kirchoff Theorem

Let D be the degree matrix of GLet A be the adjacency matrix of GLet Q = D - A

Let Q' be the matrix resulting from deleting any row and any column from Q The number of spanning trees in a graph i_{4}^{13} equal to the determinant of Q'

There are n^{n-2} spanning trees in a com¹⁸₁₉ plete graph

There are $m^{n-1} * n^{m-1}$ spanning trees in $n^{n-1} * n^{n-1}$ spanning trees in $n^{n-1} * n^{n-1}$ complete a bipartite graph

32.16 Derangements

A derangement is a permutation of a set 0 where all elements are in a different position $^{31}_{22}$ than their original position

```
der(n) = (n-1)*(der(n-1) + der(n-2)), der(0) = 35
                                                     36
                                                     37
38
39
```

32.17 Planar Graph Faces

F = E - V - 2

Others 33

33.1 Dinic

```
// Running time: O([V]^2 |E|) // OUTPUT: maximum flow value; 51 // To obtain the actual flow values, look at all edges 52
            with
     // capacity > 0 (zero capacity edges are residual edges54
4
     struct Edge {
         int from, to, cap, flow, index;
Edge(int from, int to, int cap, int flow, int index
8
             from(from), to(to), cap(cap), flow(flow), index60
     62
         int N;
vector<vector<Edge> > G;
                                                             63
         12
\frac{13}{14}
                                                             65
15
16
         G[to].push_back(Edge(to, from, 0, 0, G[from].
18
21
22
\frac{23}{24}
25
26
27
28
29
30
31
```

```
= INF;
                                            14
      15
                                            18
                                           e
19
          e->flow += amt;
   G[e->to][e->index].flow -= amt;}
totflow += amt;}
return totflow;}
                                            20
long long GetMaxFlow(int s, int t) {
                                            23
   long long totflow = 0;
while (long long flow = BlockingFlow(s, t))
                                            24
       totflow += flow;
   return totflow: } };
```

29

30

31

32

34

33.2 MinCostMaxFlow

augmentations

// INPUT:

10

40

41 42

47

48

49 50 // Running time, O(|V|^2) cost per augmentation // max flow: O(|V|^3) augmentations // min cost max flow: O(|V|^4 * MAX_EDGE_COST)

```
- graph, constructed using AddEdge()
- source
- sink
                                                                                                                                                             35
 // OUTPUT:
                         (maximum flow value, minimum cost value)
To obtain the actual flow, look at positive
                                                                                                                                                             36
                                                                                                                                                            37
38
                    values only.
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
                                                                                                                                                             39
                                                                                                                                                             40
                                                                                                                                                             \frac{41}{42}
 typedef vector<L> VL;
typeder vector<L> VL;
typedef vector<L> VUL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
typedef vector<PII> VPII;
struct MinCostMaxFlow {
                                                                                                                                                             43
                                                                                                                                                             44
45
46
47
            int N;
                                                                                                                                                             48
            VVL cap, flow, cost;
                                                                                                                                                             49
50
           VI found;
VL dist, pi, width;
VPII dad;
                                                                                                                                                            51
         MinCostMaxFlow(int N): 52

N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VU54
(N)), 55

found(N), dist(N), pi(N), width(N), dad(N) { 56

void AddEdge(int from, int to, L cap, L cost) { 57

this->cap[from][to] = cap; this->cost[from][to] = cost;} 58

void Relax(int s, int k, L cap, L cost, int dir) { 59

L val = dist[s] + pi[s] - pi[k] + cost; 60

if (cap && val < dist[k]) { 61

dist[k] = val; 62

dad[k] = make_pair(s, dir); 63

width[k] = min(cap, width[s]);} e'

L histerial.
                                                                                                                                                            52
           dad[k] = make_pair(s, dir);
width[k] = min(cap, width[s]);}}
L Dijkstra(int s, int t) {
  fill(found.begin(), found.end(), false);
  fill(dist.begin(), dist.end(), INF);
  fill(width.begin(), width.end(), 0);
}
                                                                                                                                                             65
                                                                                                                                                             66
67
68
                       fill(width.begin(), width.end(), 0
dist[s] = 0;
width[s] = INF;
while (s != -1) {
   int best = -1;
   found[s] = true;
   for (int k = 0; k < N; k++) {
      if (found[k]) continue;
   }
}</pre>
                                                                                                                                                            69
                                                                                                                                                             70
71
72
                                                                                                                                                             73
                                              Relax(s, k, cap[s][k] - flow[s][k],
cost[s][k], 1);
                                              \begin{array}{lll} & & & 77 \\ & & & \text{Relax}(s, k, \, \text{flow}[k] \, [s], \, -\text{cost}[k] \, [s], \\ & & & -1); & & \\ & & & \text{if} \, \, (\text{best} = -1 \, || \, \, \text{dist}[k] \, < \, \text{dist}[\text{best}]) \\ & & & \text{best} = k; ) \\ \end{array} 
          best = k;}
s = best;
for (int k = 0; k < N; k++)
    pi[k] = min(pi[k] + dist[k], INF);
return width[t];
pair<L, L> GetMaxFlow(int s, int t) {
                       L totflow = 0, totcost = 0;

while (L amt = Dijkstra(s, t)) {

totflow += amt;

for (int x = t; x != s; x = dad[x].first) {
                                              if (dad[x].second == 1) {
    flow[dad[x].first][x] += amt;
                                                         totcost += amt * cost[dad[x].first 5
                                                                          ][x];
                                                          flow[x][dad[x].first]
                                                         totcost -= amt * cost[x][dad[x].
                       first];}}
return make_pair(totflow, totcost);}};
                                                                                                                                                             10
33.3 PushRelabel
```

```
\label{eq:continuous} \ensuremath{//} \ significantly \ faster \ than \ straight \ Ford-Fulkerson.
It solves \ensuremath{//} random problems with 10000 vertices and 1000000
edges in a few

// seconds, though it is possible to construct test
    cases that
// achieve the worst-case.
// Running time:
// O(|V|^3)
// INPUT:
                                                                                  20
                                                                                  22
           - graph, constructed using AddEdge()
                                                                                  23
```

```
- maximum flow value
- To obtain the actual flow values, look at all
edges with
              capacity > 0 (zero capacity edges are residual
typedef long long LL;
struct Edge {
   int from, to, cap, flow, index;
   Edge(int from, int to, int cap, int flow, int index
             from(from), to(to), cap(cap), flow(flow), index
                        (index) {}};
struct PushRelabel {
      int N;
vector<vector<Edge> > G;
       vector<LL> excess;
       vector<int> dist, active, count;
      queue<int> Q;
PushRelabel(int N) : N(N), G(N), excess(N), dist(N)
                  . active(N). count(2*N) {}
       void AddEdge(int from, int to, int cap) {
   G[from].push_back(Edge(from, to, cap, 0, G[to].
                       size()));
             if (from == to) G[from].back().index++
       G[to].push_back(Edge(to, from, 0, 0, G[from].
size() - 1));}
void Enqueue(int v) {
       flow)));

if (dist[e.from] <= dist[e.to] || amt == 0)

return;
e.flow += amt;
             e.llow += amt;
G[e.to][e.index].flow -= amt;
excess[e.to] += amt;
excess[e.from] -= amt;
             Enqueue (e.to);}
       void Gap(int k) {
    for (int v = 0; v < N; v++) {
        if (dist[v] < k) continue;
        count[dist[v]]--;</pre>
                   dist[v] = max(dist[v], N+1);
count[dist[v]]++;
                    Enqueue (v); } }
      inqueue(v);)
void Relabel(int v) {
    count [dist[v]] --;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)
        if (G[v][i].cap - G[v][i].flow > 0)
    dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count [dist[v]]++;
    Frumeus(v):1
       Enqueue (v); }
             if (count
Gap(dist[v]);
                   else
      Relabel(v);}}
LL GetMaxFlow(int s, int t) {
  count[0] = N-1;
  count[N] = 1;
            count[N] = 1;
dist[s] = N;
active[s] = active[t] = true;
for (int i = 0; i < G[s].size(); i++) {
    excess[s] += G[s][i].cap;
    Push(G[s][i]);
while (!Q.empty(!)) {
    int v = 0.front();
}</pre>
                   int v = Q.front();
                   Q.pop();
active[v] = false;
Discharge(v);}
```

33.4 MinCostMatching

```
// In practice, it solves 1000x1000 problems in around
// second.
          cost[i][j] = cost for pairing left node i with
         right node j
Lmate[i] = index of right node that left node i
            pairs with
         Parts mit
Rmate[j] = index of left node that right node j
pairs with
// The values in cost[i][j] may be positive or negative
// maximization, simply negate the cost[][] matrix.
// cost MUST BE SQUARE
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &
            Rmate) {
        int n = int(cost size()):
       VD v(n);
for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i],
        cost[i][j]]);
}
for (int j = 0; j < n; j++) {
    v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j],
        cost[i][j] - u[i]);

// construct primal solution satisfying
    complementary slackness</pre>
        VD v(n);
```

```
Lmate = VI(n, -1);
Rmate = VI(n, -1);
                                                                                                                                                                                                                                               // problems of the following by a reduction to graph \,\,106
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         phi[u][c+v][0][0]++;
25
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         phi[c+v][u][0][0]++;
                                 // minimize sum_i psi_i(x[i]) 108
// x[1]...x[n] in (0,1) + sum_{(i < j)} phi_{(ij)} 109
[i], x[j]) 110
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         phi[c+v][d][0][0]++;
phi[v][c+u][v][1][1]++;
phi[c+u][v][1][1]++;
26
27
28
                                                                                                                                                                                                                                                // where
29
                                                                                                                                                                                                                                                                                                                                                                                                                                                   111
30
                                                                                                                                                                                                                                                                psi_i : {0, 1} --> R
phi_{ij} : {0, 1} x {0, 1} --> R
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          GraphCutInference graph;
31
                                                 Lmate[i] = j;
                                                                                                                                                                                                                                                 // such that
                                                                                                                                                                                                                                               // such that
// such that
// phi_(ij)(0,0) + phi_(ij)(1,1) <= phi_(ij)(0,1) + 115
    phi_(ij)(1,0) (*)
// This can also be used to solve maximization problems16
                                                Rmate[j] = i;
mated++;
break;}}
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         cout << graph.DoInference(phi, psi, x) << endl
32
33
34
35
                                                                                                                                                                                                                            10
                                  VD dist(n);
                                                                                                                                                                                                                                                                           where the
                                                                                                                                                                                                                                                 // direction of the inequality in (*) is reversed.

// INPUT: phi -- a matrix such that phi[i][j][u][v] =

phi_(ij)(u, v)
36
                                  VI dad(n):
                                                                                                                                                                                                                                                                                                                                                                                                                                                                             33.7 Geometry
37
38
                                                                                                                                                                                                                            12
                                              repeat until primal solution is feasible
                                                                                                                                                                                                                                                                                psi -- a matrix such that psi[i][u] = psi_i(u
                                                                                                                                                                                                                            13
39
                                  while (mated < n) {
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           #include <cassert>
                                                  // find an unmatched left node
40
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           double INF = 1e100;
double EPS = 1e-12;
                                                                                                                                                                                                                                                                                  x -- a vector where the optimal solution will \frac{2}{2}
41
42
                                                 int s = 0;
while (Lmate[s] != -1) s++;
                                                                                                                                                                                                                            14
                                                                                                                                                                                                                                               // X -- a vector where the optimal solution will
be stored
// OUTPUT: value of the optimal solution
// To use this code, create a GraphCutInference object,
and call the
                                                                                                                                                                                                                                                                                                                                                                                                                                                                            struct PT {
                                                 fill(dad.begin(), dad.end(), -1);
fill(seen.begin(), seen.end(), 0);
for (int k = 0; k < n; k++)</pre>
43
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          44
                                                                                                                                                                                                                            16
45
                                                                                                                                                                                                                                                 // DoInference() method. To perform maximization
46
                                                                                                                                                                                                                                              // vointerence() method. To perform maximizati
    instead of minimization,
// ensure that #define MAXIMIZATION is enabled.
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VIV VVVI;</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          dist[k] = cost[s][k] - u[s] - v[k];
j = 0;
47
                                                                                                                                                                                                                                                                                                                                                                                                                                                            q
48
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   x, y+p.y); }
49
50
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           PT operator - (const PT &p) const { return PT(x-p.
                                                                // find closest
                                                                                                                                                                                                                            20
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         x, y-p.y); }
PT operator * (double c)
    y*c ); }
PT operator / (double c)
51
                                                                                                                                                                                                                            21
                                                                                                                                                                                                                                                                                                                                                                                                                                                       11
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   const { return PT(x*c,
                                                j = -1;
for (int k = 0; k < n; k++) {
if (seen[k]) continue;
if (j == -1 || dist[k] < dist[j]) j = k;</pre>
52
53
54
                                                                                                                                                                                                                                               typedef vectortypedef vectort
                                                                                                                                                                                                                            22
                                                                                                                                                                                                                                                                                                                                                                                                                                                       12
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  const { return PT(x/c,
                                                                                                                                                                                                                            24
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         y/c ); }
55
                                                                                                                                                                                                                            25
                                                                                                                                                                                                                                                 #define MAXIMIZATION
                                                                                                                                                                                                                                                              ict GraphCutInference {
  int N;
                                                                                                                                                                                                                            26
27
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           double dot(PT p, PT q)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             { return p.x*q.x+p.y*q.y; }
                                                                          termination condition
                                                                                                                                                                                                                                                                                                                                                                                                                                                                          double dot(PI p, PT q) { return p.x*(x*p.y*q.y; )
double dist(PT p, PT q) { return dot(p-q.p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT pp) {
    os < "(" < < p. x < "," < v. y 
                                                              if (Rmate[j] == -1) break;
// relax neighbors
const int i = Rmate[j];
for (int k = 0; k < n; k++) {</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                        15
                                                                                                                                                                                                                                                               VVI cap, flow;
58
                                                                                                                                                                                                                            28
                                                                                                                                                                                                                                                               VI reached;
int Augment(int s, int t, int a) {
59
                                                                                                                                                                                                                            29
                                                                                                                                                                                                                                                                            reached[s] = 1;
if (s == t) return a;
for (int k = 0; k < N; k++) {
    if (reached[k]) continue;</pre>
61
                                                                                                                                                                                                                                                                                                                                                                                                                                                        19
62
                                                 if (seen[k]) continue;
                                                                                                                                                                                                                            32
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
                                                1f (seen[k]) continue; 32

const double new_dist = dist[j] + cost[i][k] - 33

u[i] - v[k]; 34

if (dist[k] > new_dist) {

dist[k] = new_dist;

dad[k] = j;}] 36

// update dual variables 37
                                                                                                                                                                                                                                                                                                                                                                                                                                                       20
63
                                                                                                                                                                                                                                                                                            if (int aa = min(a, cap[s][k] - flow[s][k])^{22}_{23}
64
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos
                                                                                                                                                                                                                                                                            ) {
if (int b = Augment(k, t, aa)) {
  flow[s][k] += b;
  flow[k][s] -= b;
65
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           // project point c onto line through a and b // assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
66
                                                // update dual variables
for (int k = 0; k < n; k++) {
    if (k == j || !seen[k]) continue;
    const int i = Rmate[k];
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
}</pre>
68
                                                                                                                                                                                                                            38
                                                                                                                                                                                                                                                                                                                                                                                                                                                       26
                                                                                                                                                                                                                                                                                                                                                                                                                                                                         PT ProjectPointLine (PT a, PT b, PT c) {
    return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);}

// project point c onto line segment through a and b

PT ProjectPointSegment(PT a, PT b, PT c) {
    double r = dot(b-a,b-a);
    if (fabs(r) < EPS) return a;
    r = dot(c-a, b-a)/r;
    if (r < 0) return a;
    if (r > 1) return b;
    return a + (b-a)*r;

// compute distance from c to segment between a and b

double DistancePointSegment(PT a, PT b, PT c) {
    return sqrt(dist2 (c, ProjectPointSegment(a, b, c))
69
                                                                                                                                                                                                                            39
                                                                                                                                                                                                                                                                                            return b; } }
                                                                                                                                                                                                                            40
41
42
70
71
                                                                                                                                                                                                                                                              return 0;}
int GetMaxFlow(int s, int t) {
                                                                                                                                                                                                                                                                           SetDataFlow.Hits, line t, {
N = cap.size();
flow = VVI(N, VI(N));
reached = VI(N);
int totflow = 0;
while (int amt = Augment(s, t, INF)) {
   totflow += amt;
   fill(reached.begin(), reached.end(), 0);}
                                                                                                                                                                                                                                                                                                                                                                                                                                                       29
72
                                                u[1] -= aist[k] - dist
u[s] += dist[j];
// augment along path
while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    delimite | distribution | distri
                                                                                                                                                                                                                                                                                                                                                                                                                                                       30
73
                                                                                                                                                                                                                            43
74
75
                                                                                                                                                                                                                            44
45
                                                                                                                                                                                                                                                                                                                                                                                                                                                       33
76
                                                                                                                                                                                                                            46
                                                                                                                                                                                                                            47
78
79
                                                                                                                                                                                                                            48
49
                                                                 j = d;
                                                                                                                                                                                                                                                                             return totflow; }
                                                Rmate[j] = s;
Lmate[s] = j;
mated++;}
                                                                                                                                                                                                                                                               int DoInference(const VVVVI &phi, const VVI &psi,
80
                                                                                                                                                                                                                            50
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          return sqrt(dist2(c, ProjectPointSegment(a, b, c)))
                                                                                                                                                                                                                                                                            VI &x) {
int M = phi.size();
cap = VVI(M+2, VI(M+2));
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           // compute distance between point (x,y,z) and plane ax+
                                  matea++;;
double value = 0;
for (int i = 0; i < n; i++)
  value += cost[i][Lmate[i]];</pre>
83
                                                                                                                                                                                                                            52
                                                                                                                                                                                                                                                                           cap = VVI(M+2, VI(M+2));
VI b(M);
int c = 0;
for (int i = 0; i < M; i++) {
    b[i] + psi[i][1] - psi[i][0];
    c += psi[i][0];
    for (int j = 0; j < i; j++)
b[i] += phi[i][j][1][1] - phi[i][j][0][1];
    for (int j = i+1; j < M; j++) {
    cap[i][j] = phi[i][j][0][1] + phi[i][j][1][0]
        phi[i][j][0][0] - phi[i][j][1][1];
b[i] += phi[i][j][1][0] - phi[i][j][0][0];
c += phi[i][f][0][0];}</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           by+cz=d

double DistancePointPlane(double x, double y, double z, double a, double b, double c, double d) {
84
                                                                                                                                                                                                                            53
                                                                                                                                                                                                                            54
55
                                                                                                                                                                                                                                                                                                                                                                                                                                                                            return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);} // determine if lines from a to b and c to d are
                                                                                                                                                                                                                                                                                                                                                                                                                                                       42
                                                                                                                                                                                                                            57
                                                                                                                                                                                                                                                                                                                                                                                                                                                                          // determine if lines from a to b and c to d .
parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
return fabs(cross(b-a, c-d)) < EPS; }
bool LinesCollinear(PT a, PT b, PT c, PT d) {
return LinesParallel(a, b, c, d)
&& fabs(cross(a-b, a-c)) < EPS &
&& fabs(cross(c-d, c-a)) < EPS; }
// determine if line seement from a to b into
                    33.5 MinCut
                     // Adjacency matrix implementation of Stoer-Wagner min ^{60}
                                            cut algorithm.
                     // Running time:
                                                                                                                                                                                                                            63
                                                                                                                                                                                                                                                                                    += phi[i][i][0][0];}}
                                                                                                                                                                                                                                              c += pnilif[j[[0][0];})
#ifdef MAXIMIZATION
for (int i = 0; i < M; i++) {
    for (int j = i+1; j < M; j++)
    cap[i][j] *= -1;
    b[i] *= -1;</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                       49
                                                                                                                                                                                                                                                                                                                                                                                                                                                                            // determine if line segment from a to b intersects
                      // INPUT:
                                                                                                                                                                                                                                                                                                                                                                                                                                                                             with
// line segment from c to d
                                                  graph, constructed using AddEdge()
                     // OUTPUT:
                                                                                                                                                                                                                            66
                                                    (min cut value, nodes in half of min cut)
                                                                                                                                                                                                                                                                                                                                                                                                                                                       51
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
                                                                                                                                                                                                                            67
                    typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
        dist2(b, c) < EPS || dist2(b, d) < EPS)</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                        52
                                                                                                                                                                                                                                                                                                                                                                                                                                                      53
54
                 10
                                                                                                                                                                                                                                                                            c *= -1;
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         #endif
                                                                                                                                                                                                                                                                                                                                                                                                                                                       55
                                                                                                                                                                                                                                                                            for (int i = 0; i < M; i++) {
   if (b[i] >= 0) {
13
14
15
16
                                                                                                                                                                                                                                                                                                                                                                                                                                                       56
                                                                                                                                                                                                                                                                             cap[M][i] = b[i];
                                                                                                                                                                                                                                                                                                                                                                                                                                                       57
58
59
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         return true;
                                                                                                                                                                                                                                                                             if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
17
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          \label{eq:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:false:
                                                                                                                                                                                                                                                                             int score = GetMaxFlow(M, M+1);
                                                                                                                                                                                                                                                                                                                                                                                                                                                       60
                                                                                                                                                                                                                                                                             fill(reached.begin(), reached.end(), 0);
Augment(M, M+1, INF);
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  false;
20
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           return true;}
21
                                                                                                                                                                                                                                                                            x = VI(M);
for (int i = 0; i < M; i++) x[i] = reached[i] \begin{array}{c} 63 \\ 63 \end{array}
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       compute intersection of line passing through a and b
22
                                                                                                                                                                                                                                                                                                                                                                                                                                                                             // with line passing through c and d, assuming that
23
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     unique
                                                // intersection exists; for segment intersection, check
24
                                                                                                                                                                                                                                               #ifdef MAXIMIZATION
25
                                                                                                                                                                                                                          <sup>=</sup>85
                                                                                                                                                                                                                                                                            score \star = -1;
                                                                                                                                                                                                                                                                                                                                                                                                                                                                             // segments intersect first
                                                                                                                                                                                                                                               #endif
                                                                                                                                                                                                                                                                                                                                                                                                                                                                          // segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
   b=b-a; d=c-d; c=c-a;
   assert(dot(b, b) > EPS && dot(d, d) > EPS);
   return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
26
                                                                                                                                                                                                                                                                                                                                                                                                                                                       67
                                                 used[last] = true;
cut.push_back(last);
27
                                                                                                                                                                                                                                                                                                                                                                                                                                                        68
                                                                                                                                                                                                                                                                                                                                                                                                                                                        69
70
71
                                                 if (best_weight == -1 || w[last] < best_weight)91
29
                                                                                                                                                                                                                                               int main() {
                                                                                                                                                                                                                                                                // solver for "Cat vs. Dog" from NWERC 2008
                                                                                                                                                                                                                            92
30
                                                                 best_cut = cut;
                                                                                                                                                                                                                                                                                                                                                                                                                                                        72
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          b = (a+b)/2
                                                                                                                                                                                                                                                               int numcases;
cin >> numcases;
for (int caseno = 0; caseno < numcases; caseno++)
                                                                                                                                                                                                                            93
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           return ComputeLineIntersection(b, b+RotateCW90(a-b)
                                                                best_weight = w[last];}
31
                                                 face in the second in the
32
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           , c, c+RotateCW90(a-c));)

// determine if point is in a possibly non-convex
polygon (by William
// Randolph Franklin); returns 1 for strictly interior
33
                                                                                                                                                                                                                                                                            caseno = 0; caseno < numcases; caseno
int c, d, v;
cin >> c >> d >> v;
VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
VVVI psi(c+d, VI(2));
for (int i = 0; i < v; i++) {</pre>
                                                                                                                                                                                                                            96
                                                                                                                                                                                                                                                                                                                                                                                                                                                       75
                                                                                                                                                                                                                            97
                                  return make pair (best weight, best cut); }
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   points, 0 for
                                                                                                                                                                                                                          100
                                                                                                                                                                                                                                                                                                                                                                                                                                                                           // strictly exterior points, and 0 or 1 for the remaining points.

// Note that it is possible to convert this into an *
                                                                                                                                                                                                                        101
                                                                                                                                                                                                                                                                                            char p, q;
                    33.6 GraphCutInference
                                                                                                                                                                                                                                                                                          int u, v;
cin >> p >> u >> q >> v;
                                                                                                                                                                                                                                                                                                                                                                                                                                                       78
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                    exact* test using
                                                                                                                                                                                                                                                                                          u--; v--;
if (p == 'C') {
                                                                                                                                                                                                                        104
                    // Special-purpose {0,1} combinatorial optimization
                                                                                                                                                                                                                                                                                                                                                                                                                                                                            // integer arithmetic by taking care of the division
```

appropriately

80	// (making sure to deal with signs properly) and then 176 by writing exact	cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << "_"	, 32 33	return ret;} // make an Area object from the coordinates of a
81 82	// tests for checking point on polygon boundary 177		34	polygon
83	<pre>bool PointInPolygon(const vector<pt> &p, PT q) { bool c = 0;</pt></pre> 178	(2,0), PT(4,5)) << "_" << LinesCollinear(PT(1,1), PT(3,5), PT	35	<pre>static Area makeArea(double[] pts) { Path2D.Double p = new Path2D.Double();</pre>
84	<pre>for (int i = 0; i < p.size(); i++) { int j = (i+1) %p.size(); 179</pre>	(5,9), PT(7,13)) << endl; // expected: 1 1 1 0	36 37	p.moveTo(pts[0], pts[1]); for (int i = 2) i < pts length; i = 2) p lineTo(
85 86	<pre>int j = (i+1)%p.size(); if ((p[i].y <= q.y && q.y < p[j].y 180</pre>			<pre>for (int i = 2; i < pts.length; i += 2) p.lineTo(pts[i], pts[i+1]);</pre>
87	p[j].y <= q.y && q.y < p[i].y) &&	, PT(-1,3)) << "_"	38	p.closePath();
88	q.x < p[i].x + (p[j].x - p[i].x) * (q.y - 1/81) [i].y) / (p[j].y - p[i].y)	<pre><< SegmentsIntersect(PT(0,0), PT(2,4), 1</pre>	40	<pre>return new Area(p);} // compute area of polygon</pre>
89	c = !c; 182	<< SegmentsIntersect(PT(0,0), PT(2,4), 1	PT 41	<pre>static double computePolygonArea(ArrayList<point2d.< pre=""></point2d.<></pre>
90 91	} return c;} 183	(2,-1), PT(-2,1)) << "_" << SegmentsIntersect(PT(0,0), PT(2,4), 1	PT 42	Double> points) { Point2D.Double[] pts = points.toArray(new Point2D.
92	// determine if point is on the boundary of a polygon	(5,5), PT(1,7)) << end1;		Double[points.size()]);
93 94	<pre>bool PointOnPolygon(const vector<pt> &p, PT q) { for (int i = 0; i < p.size(); i++)</pt></pre>		43	<pre>double area = 0; for (int i = 0; i < pts.length; i++) {</pre>
95	<pre>if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.</pre>	PT(3,1), PT(-1,3)) << end1;	45	<pre>int j = (i+1) % pts.length;</pre>
0.0	size()], q), q) < EPS) 186		46	<pre>area += pts[i].x * pts[j].y - pts[j].x * pts[i].y</pre>
96 97	return true; 187 return false;}	<pre>cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>	r 47	;
98	// compute intersection of line through points a and $lpha 88$	vector <pt> v;</pt>	48	return Math.abs(area)/2;}
99	with 189 // circle centered at c with radius r > 0 190		49	<pre>// compute the area of an Area object containing several disjoint polygons</pre>
00	vector <pt> CircleLineIntersection(PT a, PT b, PT c, 191</pt>	v.push_back(PT(5,5));	50	static double computeArea(Area area) {
01	<pre>double r) { vector<pt> ret; 193</pt></pre>		51 52	<pre>double totArea = 0; PathIterator iter = area.getPathIterator(null);</pre>
02	b = b-a; 194	cerr << PointInPolygon(v, PT(2,2)) << "_"	53	ArrayList <point2d.double> points = new ArrayList<</point2d.double>
$03 \\ 04$	a = a-c; 195 double A = dot(b, b); 196		54	Point2D.Double>();
05	double $A = dot(b, b);$ 196 double $B = dot(a, b);$ 197		55	<pre>while (!iter.isDone()) { double[] buffer = new double[6];</pre>
06	double C = dot(a, a) - r*r; 198	<pre><< PointInPolygon(v, PT(2,5)) << endl;</pre>	56	<pre>switch (iter.currentSegment(buffer)) {</pre>
$\frac{07}{08}$	<pre>double D = B*B - A*C; if (D < -EPS) return ret; 200</pre>		57 58	<pre>case PathIterator.SEG_MOVETO: case PathIterator.SEG_LINETO:</pre>
09	ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A); 201	<< PointOnPolygon(v, PT(2,0)) << "_"	59	points.add(new Point2D.Double(buffer[0], buffer
10 11	<pre>if (D > EPS) ret.push_back(c+a+b*(-B-sqrt(D))/A); 203</pre>		60	[1])); break;
12	return ret;} 204	<pre><< PointOnPolygon(v, PT(2,5)) << endl;</pre>	61	case PathIterator.SEG_CLOSE:
13	// compute intersection of circle centered at a with 205 radius r 206		62 63	totArea += computePolygonArea(points);
14	// with circle centered at b with radius R 207	// blank line	64	<pre>points.clear(); break;}</pre>
15	vector <pt> CircleCircleIntersection(PT a, PT b, double 08 r. double R) {</pt>		65	iter.next();}
16	r, double R) { 209 vector <pt> ret; 210</pt>		66 67	<pre>return totArea;} // notice that the main() throws an Exception</pre>
17	<pre>double d = sqrt(dist2(a, b));</pre> 211			necessary to
18 19	<pre>if (d > r+R d+min(r, R) < max(r, R)) return ret; double x = (d*d-R*R+r*r)/(2*d); 212</pre>	(2,6), PT(1,1), 5); for (int i = 0; i < u.size(); i++) cerr << u[i]	68 <<	<pre>// avoid wrapping the Scanner object for file reading in a</pre>
20	<pre>double y = sqrt(r*r-x*x);</pre>	"_"; cerr << endl;	69	// try { } catch block.
$\frac{21}{22}$	PT $v = (b-a)/d;$ ret.push_back(a+v*x + RotateCCW90(v)*y);	<pre>u = CircleLineIntersection(PT(0,9), PT(9,0), PT (1,1), 5);</pre>	70	<pre>public static void main(String args[]) throws Exception {</pre>
23	if (y > 0) 214		<< 71	Scanner scanner = new Scanner(new File("input.txt")
$\frac{24}{25}$	<pre>ret.push_back(a+v*x - RotateCCW90(v)*y); return ret;} 215</pre>	"_"; cerr << endl; u = CircleCircleIntersection(PT(1,1), PT(10,10),	E 79); // also,
26	// This code computes the area or centroid of a (5);	73	// Scanner scanner = new Scanner (System.in);
27	possibly nonconvex) 216 // polygon, assuming that the coordinates are listed in	<pre>for (int i = 0; i < u.size(); i++) cerr << u[i] - ","; cerr << endl;</pre>	<< 74 75	<pre>double[] pointsA = readPoints(scanner.nextLine()); double[] pointsB = readPoints(scanner.nextLine());</pre>
21	a clockwise or 217			Area areaA = makeArea(pointsA);
28	// counterclockwise fashion. Note that the centroid is often known as 218	5);	77	<pre>Area areaB = makeArea(pointsB); areaB.subtract(areaA);</pre>
29	often known as 218 // the "center of gravity" or "center of mass".	<pre>for (int i = 0; i < u.size(); i++) cerr << u[i] - "_"; cerr << endl;</pre>	79	// also,
30	double ComputeSignedArea(const vector <pt> &p) { 219</pt>			// areaB.exclusiveOr (areaA);
$\frac{31}{32}$	<pre>double area = 0; for(int i = 0; i < p.size(); i++) {</pre>	10, sqrt(2.0)/2.0); for (int i = 0; i < u.size(); i++) cerr << u[i]	81 << 82	<pre>// areaB.add (areaA); // areaB.intersect (areaA);</pre>
33	<pre>int j = (i+1) % p.size();</pre>	"_"; cerr << endl;	83	// (1) determine whether B - A is a single closed
$\frac{34}{35}$	area += p[i].x*p[j].y - p[j].x*p[i].y;} 221 return area / 2.0;}	<pre>u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5) 5, sqrt(2.0)/2.0);</pre>	84	<pre>shape (as // opposed to multiple shapes)</pre>
36	double ComputeArea(const vector <pt> &p) { 222</pt>	<pre>for (int i = 0; i < u.size(); i++) cerr << u[i]</pre>	<< 85	boolean isSingle = areaB.isSingular();
$\frac{37}{38}$	<pre>return fabs(ComputeSignedArea(p));} PT ComputeCentroid(const vector<pt> &p) { 223</pt></pre>	"_"; cerr << endl; // area should be 5.0	86 87	<pre>// also, // areaB.isEmpty();</pre>
39	PT c(0,0); 224	// centroid should be (1.1666666, 1.166666)	88	if (isSingle)
$\frac{40}{41}$	<pre>double scale = 6.0 * ComputeSignedArea(p); for (int i = 0; i < p.size(); i++) {</pre>		; 89 90	System.out.println("The_area_is_singular."); else
42	<pre>int j = (i+1) % p.size();</pre> 227	PT c = ComputeCentroid(p);	91	System.out.println("The_area_is_not_singular.");
43	c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[228]		92 93	// (2) compute the area of B - A
44].y); 229		93	<pre>System.out.println("The_area_is_" + computeArea(</pre>
45	return c / scale;}		94	// (3) determine whether each p[i] is in the
46	<pre>// tests whether or not a given polygon (in CW or CCW</pre>	33.8 JavaGeometry	95	<pre>interior of B - A while (scanner.hasNextDouble()) {</pre>
47	bool IsSimple(const vector <pt> &p) {</pt>		96	<pre>double x = scanner.nextDouble();</pre>
$\frac{48}{49}$	<pre>for (int i = 0; i < p.size(); i++) { for (int k = i+1; k < p.size(); k++) {</pre>	// In this example, we read an input file containing	97 98	<pre>assert(scanner.hasNextDouble()); double y = scanner.nextDouble();</pre>
50	<pre>int j = (i+1) % p.size();</pre>	three lines, each // containing an even number of doubles, separated b	99	<pre>if (areaB.contains(x,y)) {</pre>
$\frac{51}{52}$	<pre>int 1 = (k+1) % p.size(); if (i == 1 j == k) continue;</pre>	commas. The first two		<pre>System.out.println ("Point_belongs_to_the_area. ");</pre>
53	if (SegmentsIntersect($p[i]$, $p[j]$, $p[k]$, $p[1]$		101	} else {
54	<pre>])) return false;}}</pre>	// (or clockwise) order, which we will call "A" and	"B	<pre>System.out.println ("Point_does_not_belong_to_</pre>
55	return true;}	". The last line	103	// Finally, some useful things we didn't use in
$\frac{56}{57}$	int main() {	//	104	this example: //
58	cerr << RotateCCW90(PT(2,5)) << endl;		105	// Ellipse2D.Double ellipse = new Ellipse2D.
59 60	// expected: (5,-2) cerr << RotateCW90(PT(2,5)) << endl;	opposed to multiple shapes)	106	Double (double x, double y,
61	// expected: (-5,2)			
62 63	cerr << Rotateccw(PT(2,5),M_PI/2) << end1; // expected: (5.2)	//	107	double w, double h); //
64	cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT		108	// creates an ellipse inscribed in box with
65	(3,7)) << end1; // expected: (5.2) (7.5.3) (2.5.7)	// 0 0 10 10 10 0	109	<pre>bottom-left corner (x,y) // and upper-right corner (x+y,w+h)</pre>
66	cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT	// 8 6 	110	//
67	(3,/)) < "_" // Project Point Segment (PT (7 5 3) PT (10 4) 17	' //	111	<pre>// Rectangle2D.Double rect = new Rectangle2D. Double (double x, double y,</pre>
	. PT(3.7)) << " "		112	// vouble (double x, double y,
68	<pre><< ProjectPointSegment(PT(-5,-2), PT 20</pre>			
69	(2.5,1), PT(3,7)) << endl; 21 // expected: 6.78903	<pre>// Point belongs to the area.</pre>	113	double w, double h); //
70	cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << 25		114	<pre>// creates a box with bottom-left corner (x,y)</pre>
71	endl;	import java.awt.geom.*;	115	and upper-right // corner (x+y,w+h)
72	cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT	<pre>import java.io.*; public class JavaGeometry {</pre>	116	//
73	(4,5)) << "_" << LinesParallel(PT(1,1), PT(3,5), PT(2,0) 28	// make an array of doubles from a string	117	<pre>// Each of these can be embedded in an Area object (e.g., new Area (rect)).</pre>
			118	}
74	<pre><< LinesParallel(PT(1,1), PT(3,5), PT(5,9)</pre>		119	}
75	// expected: 0 0 1	<pre>for (int i = 0; i < arr.length; i++) ret[i] = Double parseDouble(arr[i]):</pre>		

33.9 Geom3D for (size_t i = 0; i < x.size(); i++) cerr << """</pre> cerr << endl;].k); 90 return 0; 91 92 public class Geom3D { // distance from point (x, y, z) to plane aX + bY + 47 cZ + d = 0 3 public static double ptPlaneDist(double x, double y, 33.11 Simplex 33.12 KDTree double 2, double a, double b, double c, double d) { return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a + b*b + c*c); - constructs from n points in O(n 1g°2 n) time - handles nearest-neighbor query in O(1g n) if points are well distributed - worst case for nearest-neighbor may be linear in pathological case // Two-phase simplex algorithm for solving linear programs of the form maximize c^T x subject to Ax <= b $^{\prime}$ // distance between parallel planes aX + bY + cZ + dl $_{2}$ = 0 and // aX + bY + c2 + d2 = 0 public static double planePlaneDist(double a, double // INPUT: A -- an m x n matrix b, double c, double d1, double d2) { return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + b -- an m-dimensional vector 5 c -- an n-dimensional vector 6 x -- a vector where the optimal solution will 7 #include mits> 10 // number type for coordinates, and its maximum value typedef long long ntype; const ntype sentry = numeric_limits<ntype>::max(); // point structure for 2D-tree, can be extended to 3D struct point { C*C); be stored 12 , // distance from point (px, py, pz) to line (x1, y1, $\frac{9}{21}$) -(x2, y2, z2) // (or ray, or segment; in the case of the ray, the $\frac{10}{11}$ // OUTPUT: value of the optimal solution (infinity if unbounded // above, nan if infeasible) 13 ntype x, y; 14 // To use this code, create an LPSolver object with A, 12point (ntype xx = 0, ntype yy = 0) : x(xx), y(yy)endpoint is the b, and c as // arguments. Then, call Solve(x). #include <limits> // first point) public static final int LINE = 0; 17 public static final int SEGMENT = 1; typedef long double DOUBLE; typedef vector<DOUBLE> VD; typedef vector<VD> VVD; 18 public static final int RAY = 2: public static double ptLineDistSq(double x1, double yl, double zl, typedef vector<int> VI; double x2, double y2, double z2, double px, double py, double pz, 18 20 const DOUBLE EPS = 1e-9; 19 struct LPSolver { 20 21 int type) 21 int m, n; 22 ntype pdist2(const point &a, const point &b) { ntype dx = a.x-b.x, dy = a.y-b.y; return dx*dx + dy*dy;) // bounding box for a set of points 21 VI B. N: 22 22 777D D• 23 double x, y, z; if (pd2 == 0) { 23 $\frac{23}{24}$ LPSolver(const VVD &A, const VD &b, const VD &c) 24 m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VL25 (n+2) { 25 x = x1;struct bbox{ 26 y = y1; z = z1; ntype x0, x1, y0, y1; bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-25 27 28 29 **double** $u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (^{26}$ sentry) {} double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (20-x1)*(x2-x1)) / pd2; x = x1 + u * (x2 - x1); y = y1 + u * (y2 - y1); z = z1 + u * (z2 - z1); 28 if (type != LINE && u < 0) { x = x1;y = y1;z = z1;} if (type == SEGMENT && u > 1.0) { x = x2;y = y2;z = z2;} return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-p2);); 30 31 32 33 35 36 37 38 public static double $\mathtt{ptLineDist}(\texttt{double}\ \mathtt{x1},\ \texttt{double}\ \mathtt{y1},^{34}$ double z1, double x2, double y2, double z2, double px, double py, double pz, - (p.y , p); else 39 return pdist2(point(x0, p.y), p);} else if (p.x > x1) { 40 int type) { int x = phase while (true) { int s = -1; return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, 38 z2, px, py, pz, type));}} 41 41 - (p.y , p); else 42 33.10 Delaunay return pdist2(point(x1, p.y 43), p); if (D[x][s] >= -EPS) return true; // Slow but simple Delaunay triangulation. Does not 45 45 **if** (p.y < y0) return pdist2(point(p.x, y0 int r = -1; for (int i = 0; i < m; i++) { 46 if (D[i][s] < -0) continue; if (r = -1 | | D[i][n+1] / 47 | D[r][s] | | (| D[i][s] < D[r][n+1] / 47 | (| D[i][s]) | (| D[i][s handle // degenerate cases (from O'Rourke, Computational), p); else if (p.y > y1) return pdist2(point(p.x, y1), p); else Geometry in C) 48 49 return 0:}}; return 0;}}; // stores a single node of the kd-tree, either internal or leaf // Running time: O(n^4) D[r][s] || 4 D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i] < B[r]) r = i; 4 // INPUT: x[] = x-coordinates // y[] = y-coordinates struct kdnode { 6 7 50 bool leaf; // true if this is a leaf node (has eat; // true if this to a termination one point) pt; // the single point of this is a leaf ound; // bounding box for set of points in if (r == -1) return false; 8 9 triples = a vector containing m triples Pivot(r, s);}} point pt; bbox bound; indices DOUBLE Solve(VD &x) { 10 corresponding to triangle children kdnode *first, *second; // two children of this kd-#include<vector> 12 using namespace std; 57 typedef double T; struct triple { Privot(r, n); if (!Simplex(1) || D[m+1][n+1] < -EPS) return</pre> // intersect a point with this node (returns squared distance) -numeric_limits<DOUBLE>::infinity(); 56 15 int i, j, k; triple() {} 16 60 for (int i = 0; i < m; i++) if (B[i] == -1) { int s = -1; for (int j = 0; j <= n; j++) if (s == -1 || D[i][j] < D[i][s] || D[i][j] 59 = D[i][s] && N[j] < N[s]) s = j; 60 triple(int i, int j, int k) : i(i), j(j), k(k) {}},dvector<triple> delaunayTriangulation(vector<T>& x, 62 ntype intersect(const point &p) { 18 return bound.distance(p);} // recursively builds a kd-tree from a given cloud of 19 points Pivot(i, s);}} if (!Simplex(2)) return numeric_limits<DOUBLE>:: 61 20 void construct(vector<point> &vp) { // compute bounding box for points at this node 22 infinity(); bound.compute(vp); // if we're down to one point, then we're a leaf 23 63 node if (vp.size() == 1) { leaf = true; 24 25 26 65 pt = vp[0];} else { 27 // split on x if the bbox is wider than high (not $\begin{array}{c} \text{[i])*(z[j]-z[i]);} & 71 \\ \text{double } \text{yn} = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-xT2) \\ \text{[i])*(z[k]-z[i]);} & 73 \\ \text{double } \text{zn} = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-xT4) \\ \text{[i])*(y[j]-y[i]);} & 75 \\ \text{bool } \text{flag} = \text{zn} < 0; & 76 \\ \text{for } (\text{int } m = 0; \text{ flag } \delta \delta \text{ m < n; m++}) \\ \text{flag} = \text{flag } \delta \delta ((x[m]-x[i])*xn + (y[m]-y[i])*78 \\ \text{yn} + (z[m]-z[i])*zn < 0); & 15 \\ \text{(flag)} \text{ yre tweek back(triple(i, i, k)) } & 180 \\ \end{array}$ const int n = 3; 68 DOUBLE _A[m][n] = { { 6, -1, 0 }, { -1, -5, 0 }, { 1, 5, 1 }, best heuristic...) if (bound.xl-bound.x0 >= bound.yl-bound.y0) sort(vp.begin(), vp.end(), on_x); // otherwise split on y-coordinate 29 69 70 71 31 72 sort(vp.begin(), vp.end(), on_y); // divide by taking half the array for each child // (not best performance if many duplicates in 32 33 }; DOUBLE _b[m] = { 10, -4, 5, -5 }; DOUBLE _c[n] = { 1, -1, 0 }; if (flag) ret.push_back(triple(i, j, k));}} 80 34 VVD A(m); return ret;} int main() { T xs[]={0, 0, 1, 0.9}; VD b(_b, _b + m); VD c(_c, _c + n); for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] 78</pre> 36 37 83 T xs[]={0, 1, 0, 0, 9}; vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]); vector<triple> tri = delaunayTriangulation(x, y); LPSolver solver(A, b, c); 40 ; } } }; DOUBLE value = solver.Solve(x); // simple kd-tree class to hold the tree and handle 41 //expected: 0 1 3 86 81 cerr << "VALUE: "<< value << endl; cerr << "SOLUTION:";

```
kdnode *root;
             // constructs a kd-tree from a points (copied here, as it sorts them)
kdtree(const vector<point> &vp) {
 84
                                                                                             40
                   vector<point> v(vp.begin(), vp.end());
 86
                   root = new kdnode();
                                                                                              42
 87
             root->construct(v);}
kdtree() { delete root;
             kdtree() { delete root; }
// recursive search method returns squared distance
to nearest point
ntype search(kdnode *node, const point &p) {
    if (node->leaf) {
        // commented special case tells a point not
}
 90
 \frac{91}{92}
 93
                            to find itself
if (p == node->pt) return sentry;
 94
                              return pdist2(p, node->pt);}
 96
                   ntype bfirst = node->first->intersect(p); 50
ntype bsecond = node->second->intersect(p); // choose the side with the closest bounding box 51
 97
 98
99
                   Lo search first $'' (note that the other side is also searched if 52 needed)  
                             to search first
100
101
102
                          ntype best = search(node->first, p);
                              (bsecond < best)
best = min(best, search(node->second, p))
103
104
105
                         return best;}
106
107
                         of the first search (node->second, p); 57

if (bfirst < best)
best = min(best, search(node->first, p));58
108
109
             best = min(best, search
return best;}}
// squared distance to the nearest
ntype nearest(const point &p) {
   return search(root, p);};
112
           // some basic test code here
115
                                                                                              62
116
          int main() {
             119
120
             kdtree tree(vp);
122
             123
126
127
129
          33.13 LogLan
```

```
1
         // Code which demonstrates the use of Java's regular
                     expression libraries.
          // This is a solution for
 3
                  Loglan: a logical language
                  http://acm.uva.es/p/v1/134.html
          //
// In this problem, we are given a regular language,
         // In this problem, we are given a regular language, 10
whose rules can be 11
// inferred directly from the code. For each sentence 12
in the input, we must 13
// determine whether the sentence matches the regular 14
expression or not. The 15
// code consists of (1) building the regular expression<sup>16</sup>
(which is fairly 10 report by regular expression<sup>16</sup>
 9
          // complex) and (2) using the regex to match sentences.
          import java.util.*;
import java.util.regex.*;
           14
         public class LogLan {
18
19
21
22
23
25
27
28
29
30
                ")";
String statement = "(" + predname + space + 21
verbpred + space + predname + "|" + 22
predname + space + verbpred + ")"; 23
String sentence = "(" + statement + "|" + predclaim#4
25
31
33
                 + ")";
return "^" + sentence + "$";}
34
                public static void main (String args[]) {
String regex = BuildRegex();
                                                                                                           27
36
                Pattern pattern = Pattern.compile (regex);
Scanner s = new Scanner(System.in);
37
```

```
while (true) {
  // In this problem, each sentence consists of
  s.length() to get length of string
s.charAt() to extract characters from a
            Java string
s.trim() to remove whitespace from the
beginning and end of Java string
   // Other useful String manipulation methods include
            s.compareTo(t) < 0 if s < t,
            lexicographically
s.indexOf("apple") returns index of first
            occurrence of "apple" in s
s.lastIndexOf("apple") returns index of
last occurrence of "apple" in s
s.replace(c,d) replaces occurrences of
            character c with d
           character c with d
s.startsWith("apple) returns (s.indexOf(" 19
20
apple") == 0)
s.toLowerCase() / s.toUpperCase() returns
21
22
              new lower/uppercased string
                                                                            23
            Integer.parseInt(s) converts s to an
             integer (32-bit)
            Long.parseLong(s) converts s to a long (64-26
           bit)
Double.parseDouble(s) converts s to a
                                                                            29
            double
                                                                            30
   String sentence = "";
   while (true) {
    sentence = (sentence + "," + s.nextLine()).
                                                                            33
                 trim();
        if (sentence.equals("#")) return;
if (sentence.charAt(sentence.length()-1) ==
    .') break;
                                                                            34
                                                                            37
  }
// now, we remove the period, and match the
regular expression
String removed_period = sentence.substring(0,
  sentence.length()-1).trim();
if (pattern.matcher (removed_period).find()) {
   System.out.println ("Good");
                                                                            44
   } else {
                                                                            45
        System.out.println ("Bad!");}}}
                                                                            48
                                                                            49
50
```

33.14 IO

```
// Ouput a specific number of digits past the
decimal point,
                                                                                52
// in this case 5
cout.setf(ios::fixed); cout << setprecision(5);</pre>
cout << 100.0/7.0 << endl;
cout unsetf(ios::fixed);
// Output the decimal point and trailing zeros</pre>
cout.setf(ios::showpoint);
cout << 100.0 << endl;
cout.unsetf(ios::showpoint);
// Output a '+' before positive values</pre>
cout.setf(ios::showpos);
cout << 100 << "_" << -100 << endl;
cout.unsetf(ios::showpos);</pre>
                                                                                62
66
```

55

58

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63

70

 $\frac{71}{72}$

33.15 LatLong

```
\frac{73}{74}
Converts from rectangular coordinates to latitude/
longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <iostream>
#include <cmath>
using namespace std;
struct l1( double r, lat, lon;);
struct rect(double x, y, z;);
l1 convert(rect& P) {
    11 Q;
    Q.r = sqrt(P.x*P.x*P.y*P.y*P.z*P.z);
    Q.lat = 180/M_PI*asin(P.z/Q.r);
    Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x*P.y*P.y));
    return Q;}
       return Q; }
rect convert(11& 0) {
                                                                                         89
      return P;}
int main() {
       rect A:
                                                                                         96
      10 B; 97
A.x = -1.0; A.y = 2.0; A.z = -3.0; 98
B = convert(A); 99
cout << B.r << "..." << B.lat << "..." << B.lon << endition
```

34 Edmonds Blossom

```
struct edge {
            int v, nx;
 const int MAXN = 1000, MAXE = 2000;
edge graph[MAXE];
int last[MAXN], match[MAXN], px[MAXN], base[MAXN], N,
edges;
bool used[MAXN], blossom[MAXN], lused[MAXN];
inline void add_edge(int u, int v,
graph[edges] = (edge) {v, last[u]};
last[u] = edges++;
            graph[edges] = (edge) {u, last[v]};
            last[v] = edges++;
}
void mark_path(int v, int b, int children) {
   while (base[v] != b) {
      blossom[base[v]] = blossom[base[match[v]]] = true
                      ;
px[v] = children;
                        children = match[v];
                       v = px[match[v]];
int lca(int a, int b) {
  memset(lused, 0, N);
  while (1) {
    lused[a = base[a]] = true;
                      if (match[a] == -1)
                     break;
a = px[match[a]];
             while (1) {
                     b = base[b];
if (lused[b])
                                 return b:
                     b = px[match[b]];
int find path(int root) {
          t innd_path(int root) {
    memset(used, 0, N);
    memset(px, -1, sizeof(int) * N);
    for (int i = 0; i < N; ++i)
        base[i] = i;
    used[root] = true;
    queue(int) q;
        return for the content of the content o
            q.push (root);
            int v, e, to, i;
while (!q.empty()) {
                    mark_path(v, curbase, to);
mark_path(to, curbase, v);
for (i = 0; i < N; ++i)
    if (blossom[base[i]]) {</pre>
                                                                base[i] = curbase;
if (!used[i]) {
   used[i] = true;
   q.push(i);
                                  } else if (px[to] == -1) {
                                            px[to] = v;
if (match[to] == -1)
    return to;
                                            to = match[to];
                                             used[to] = true;
                                           q.push(to);
            return -1;
 void build pre matching() {
            int u, e, v;
for (u = 0; u < N; ++u)
   if (match[u] = -1)
      for (e = last[u]; e >= 0; e = graph[e].nx) {
                                           v = graph[e].v;
if (match[v] == -1) {
  match[u] = v;
  match[v] = u;
                                                       break.
  void edmonds() {
  memset(match, 0xff, sizeof(int) * N);
            build_pre_matching();
            int i, v, pv, ppv;
for (i = 0; i < N; ++i)
   if (match[i] == -1)</pre>
                                 v = find_path(i);
                                  while (v != -1) {
                                          pv = px[v], ppv = match[pv];
match[v] = pv, match[pv] = v
                                            v = ppv;
```