

Linear Programs Zero Sum Game Formulation

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Consider a normal-form zero-sum game with payoff matrix A , where player 1 (P1) chooses a mixed strategy represented by the vector x , and player 2 (P2) chooses a mixed strategy y .

The expected utility for player 1 is given by:

$$v = x^T A y$$

where v represents the value of the game. It turns out that it makes more sense to frame v as an independent decision variable.

Trivial Constraints

The strategies x and y must satisfy the standard probability constraints:

$$\begin{aligned} \sum_i x_i &= 1, & 0 \leq x_i & \quad \forall i \\ \sum_j y_j &= 1, & 0 \leq y_j & \quad \forall j \end{aligned}$$

Equilibrium Constraints

A necessary condition for Nash equilibrium is that player 2 should be indifferent between all actions they play with positive probability. This is equivalent to requiring that the mixed strategy y should do at least as well for player 2 as any pure strategy, against the specific strategy x (or rather, that the value v for player 1 cannot be decreased by deviating y to any pure strategy):

$$v \leq x^T A e_j, \quad \forall j$$

which can be rewritten as:

$$v \cdot \mathbf{1} \leq x^T A \mathbf{1} = x^T A$$

$$x^T A \geq v \cdot \mathbf{1}$$

Where $\mathbf{1}$ is a vector of all 1s. Notice y and the indices j do not appear in this constraint.

All in all, we get the linear program:

Maximize v

Subject to

$$x^T A \geq v \cdot \mathbf{1}$$

$$\sum_i x_i = 1$$

$$0 \leq x_i \quad \forall i$$

The solution to this linear program will give the value v of the game and the Nash equilibrium strategy x for player 1¹.

¹Supposedly, Player 2's strategy y can be found similarly by solving the dual problem