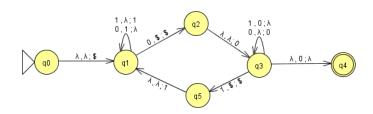
October 16, 2017

1. Construct a pushdown automata that recognizes $\{w \mid w \text{ is an element of } \{0,1\}^* \text{ and } w \text{ has more 0's than 1's }\}.$



2. Convert the following CFG into an equivalent CFG in Chomsky normal form, using the procedure given in Theorem 2.9.

$$\begin{array}{l} A \rightarrow BAB \mid B \mid \epsilon \\ B \rightarrow 00 \mid \epsilon \end{array}$$

$$\begin{array}{l} S \rightarrow BC \mid AB \mid BA \mid BB \mid DD \mid \epsilon \\ A \rightarrow BC \mid AB \mid BA \mid BB \mid DD \\ B \rightarrow DD \\ C \rightarrow AB \\ D \rightarrow 0 \end{array}$$

3. Show that the class of context-free languages is closed under the union operation (construction and proof). The construction should be quite simple.

4. Show that the class of context-free languages is closed under the concatenation operation (construction and proof). The construction should be quite simple.

5. Show that the class of context-free languages is closed under the star operation (construction and proof). The construction should be quite simple.

6. Let $M = (Q, \Sigma, \delta, q_0, F)$ be a DFA and define CFG $G = (V, \Sigma, R, S)$ as follows:

- V = Q;
- For each $q \in Q$ and $a \in \Sigma$, define rule $q \to aq'$ where $q' = \delta(q, a)$;
- For $q \in F$ define rule $q \to \epsilon$;
- $S = q_0$.

Prove L(M) = L(G).

7. Let $L = \{0^n 1^m 0^n 1^m \mid n, m \ge 0\}$. Show L is not context-free.

8. Let $L = \{w \mid w \text{ is in } \{a, b, c, d\}^*$, with the number of a's = number of b's and the number of c's = the number of d's $\}$. Show L is not context-free.

9. Let A and B be languages. We define $A \approx B = \{ab \mid a \text{ is an element of } A \text{ and } b \text{ is an element of } B \text{ and } |a| > |b| \}$. Show that if A and B are regular languages, then $A \approx B$ is a context free language.

10. Show $L = \{w \mid w \text{ is an element of } \{a, b, c, d, e, f\}^* \text{ such that the number of a's} + \text{number of b's} = \text{number of c's} + \text{number of d's} = \text{number of e's} + \text{number of f's} \} \text{ is not context-free.}$