1. This exercise concerns TM M_2 , whose description and state diagram appear in Example 3.7. In each of the parts, give the sequence of configurations that M_2 enters when started on the indicated input string.

   (a) 0.

       $q_1 0 \llcorner$ $\llcorner q_2 \llcorner$ $\llcorner\llcorner q_{accept}$

   (b) 000.

       $q_1 000 \llcorner$ $\llcorner q_2 00 \llcorner$ $\llcorner\llcorner q_3 0 \llcorner$ $\llcorner\llcorner 0 q_4 \llcorner$ $\llcorner\llcorner 0 \llcorner q_{reject}$

2. This exercise concerns TM M_1 , whose description and state diagram appear in Example 3.9. In each of the parts, give the sequence of configurations that M_1 enters when started on the indicated input string.

   (a) 1#1.

   (b) 1##1.

3. Describe a Turing machine, sequence of steps, that recognizes $\{w \mid w$ is an element of $\{a, b, c\}*$ such that the number of a's in $w <$ the number of b's in w and the number of a's in $w =$ the number of c's in $w\}$

4. Show the equivalent transitions for a 2-PDA for the Turing machine transitions $(q_i, X) \rightarrow (q_j, A, L)$ and $(q_i, X) \rightarrow (q_j, A, R)$ (in state q_i read X, write A, and move left or right and transition to state q_j). The transitions for a 2-PDA are of the form $(q_i, X, S_1, S_2) \rightarrow (q_j, T_1, T_2)$ (in state q_i, read X, pop S_1 from stack 1, pop S_2 from stack 2, transition to state q_j, push T_1 onto stack 1 and push T_2 onto stack 2). You don't have to prove the transitions are equivalent, just tell me what they are.

5. Give implementation-level descriptions of Turing machines that decide the following languages over the alphabet $\{0, 1\}$. $\{w \mid w$ does not contain twice as many 0's as 1's$\}$

6. Prove the class of Turing recognizable languages is closed under the union operation (construction and proof)

7. Prove the class of decidable languages is closed under concatenation (construction and proof)

8. Prove the class of decidable languages is closed under intersection (construction and proof)

9. Prove the class of Turing recognizable languages is closed under the star operation (construction and proof)

10. Show that a language is decidable iff some enumerator enumerates the language in the standard string order.