

1. Construct a pushdown automata that recognizes  $\{w \mid w \text{ is an element of } \{0, 1\}^* \text{ and } w \text{ has more 0's than 1's}\}$ .
2. Convert the following CFG into an equivalent CFG in Chomsky normal form, using the procedure given in Theorem 2.9.

$$\begin{aligned} A &\rightarrow BAB \mid B \mid \epsilon \\ B &\rightarrow 00 \mid \epsilon \end{aligned}$$

3. Show that the class of context-free languages is closed under the union operation (construction and proof). The construction should be quite simple.
4. Show that the class of context-free languages is closed under the concatenation operation (construction and proof). The construction should be quite simple.
5. Show that the class of context-free languages is closed under the star operation (construction and proof). The construction should be quite simple.
6. Let  $M = (Q, \Sigma, \delta, q_0, F)$  be a DFA and define CFG  $G = (V, \Sigma, R, S)$  as follows:

- $V = Q$ ;
- For each  $q \in Q$  and  $a \in \Sigma$ , define rule  $q \rightarrow aq'$  where  $q' = \delta(q, a)$ ;
- For  $q \in F$  define rule  $q \rightarrow \epsilon$ ;
- $S = q_0$ .

Prove  $L(M) = L(G)$ .

7. Let  $L = \{0^n 1^m 0^n 1^m \mid n, m \geq 0\}$ . Show  $L$  is not context-free.
8. Let  $L = \{w \mid w \text{ is in } \{a, b, c, d\}^*, \text{ with the number of a's} = \text{number of b's and the number of c's} = \text{the number of d's}\}$ . Show  $L$  is not context-free.
9. Let  $A$  and  $B$  be languages. We define  $A \approx B = \{ab \mid a \text{ is an element of } A \text{ and } b \text{ is an element of } B \text{ and } |a| > |b|\}$ . Show that if  $A$  and  $B$  are regular languages, then  $A \approx B$  is a context free language.
10. Show  $L = \{w \mid w \text{ is an element of } \{a, b, c, d, e, f\}^* \text{ such that the number of a's} + \text{number of b's} = \text{number of c's} + \text{number of d's} = \text{number of e's} + \text{number of f's}\}$  is not context-free.