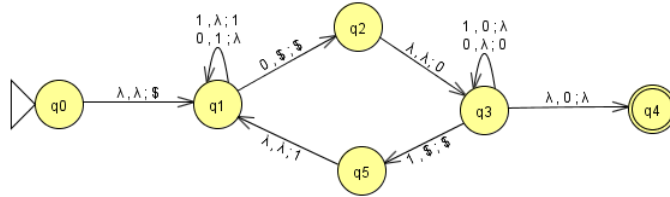


1. Construct a pushdown automata that recognizes $\{w \mid w \text{ is an element of } \{0,1\}^* \text{ and } w \text{ has more 0's than 1's}\}$.



2. Convert the following CFG into an equivalent CFG in Chomsky normal form, using the procedure given in Theorem 2.9.

$$A \rightarrow BAB \mid B \mid \epsilon$$

$$B \rightarrow 00 \mid \epsilon$$

3. Show that the class of context-free languages is closed under the union operation (construction and proof). The construction should be quite simple.
4. Show that the class of context-free languages is closed under the concatenation operation (construction and proof). The construction should be quite simple.
5. Show that the class of context-free languages is closed under the star operation (construction and proof). The construction should be quite simple.
6. Let $M = (Q, \Sigma, \delta, q_0, F)$ be a DFA and define CFG $G = (V, \Sigma, R, S)$ as follows:
 - $V = Q$;
 - For each $q \in Q$ and $a \in \Sigma$, define rule $q \rightarrow aq'$ where $q' = \delta(q, a)$;
 - For $q \in F$ define rule $q \rightarrow \epsilon$;
 - $S = q_0$.

Prove $L(M) = L(G)$.

7. Let $L = \{0^n 1^m 0^n \mid n, m \geq 0\}$. Show L is not context-free.
8. Let $L = \{w \mid w \text{ is in } \{a, b, c, d\}^*, \text{ with the number of a's} = \text{number of b's and the number of c's} = \text{the number of d's}\}$. Show L is not context-free.
9. Let A and B be languages. We define $A \approx B = \{ab \mid a \text{ is an element of } A \text{ and } b \text{ is an element of } B \text{ and } |a| > |b|\}$. Show that if A and B are regular languages, then $A \approx B$ is a context free language.
10. Show $L = \{w \mid w \text{ is an element of } \{a, b, c, d, e, f\}^* \text{ such that the number of a's} + \text{number of b's} = \text{number of c's} + \text{number of d's} = \text{number of e's} + \text{number of f's}\}$ is not context-free.