November 20, 2017

- 1. This exercise concerns TM  $M_2$ , whose description and state diagram appear in Example 3.7. In each of the parts, give the sequence of configurations that  $M_2$  enters when started on the indicated input string.
  - (a) 0.

$$q_10$$
\_  $q_2$ \_  $q_{accept}$ 

(b) 000.

- 2. This exercise concerns TM  $M_1$ , whose description and state diagram appear in Example 3.9. In each of the parts, give the sequence of configurations that  $M_1$  enters when started on the indicated input string.
  - (a) 1#1.

```
q_11\#1 xq_3\#1 x\#q_51 xq_6\#x q_7x\#x xq_1\#x x\#q_8x x\#xq_8 x\#x_4
```

(b) 1##1.

```
q_11##1 xq_3##1 x#q_5#1 x##q_{reject}1
```

- 3. Describe a Turing machine, sequence of steps, that recognizes  $\{w \mid w \text{ is an element of } \{a,b,c\}^* \text{ such that the number of } a$ 's in w < the number of b's in w and the number of a's in w = the number of c's in w
  - (1) Place symbol at the left side of tape
  - (2) Scan right for a, if found: mark it, else: go to step 6
  - (3) Rewind
  - (4) Scan right for b, if found: mark it, else: Halt and Reject (a must be < b)
  - (5) Rewind and go to step 2.
  - (6) Rewind
  - (7) Scan right for a', if found: mark it, else: go to step 11
  - (8) Rewind
  - (9) Scan right for c, if found: mark it, else: Halt and Reject (a must be = c)
  - (10) go to step 6.
  - (11) Scan right for c, if found: Halt and Reject, else: Halt and Accept.
- 4. Show the equivalent transitions for a 2-PDA for the Turing machine transitions  $(q_i, X) \to (q_j, A, L)$  and  $(q_i, X) \to (q_j, A, R)$  (in state  $q_i$  read X, write A, and move left or right and transition to state  $q_j$ ). The transitions for a 2-PDA are of the form  $(q_i, X, S_1, S_2) \to (q_j, T_1, T_2)$  (in state  $q_i$ , read X, pop  $S_1$  from stack 1, pop  $S_2$  from stack 2, transition to state  $q_j$ , push  $T_1$  onto stack 1 and push  $T_2$  onto stack 2). You don't have to prove the transitions are equivalent, just tell me what they are.
- 5. Give implementation-level descriptions of Turing machines that decide the following languages over the alphabet  $\{0,1\}$ .  $\{w \mid w \text{ does not contain twice as many 0's as 1's}\}$ 
  - (1) Place symbol at left side of tape
  - (2) Rewind
  - (3) Scan right for 1, if found: mark it, else: go to step 9
  - (4) Rewind
  - (5) Scan right for 0, if found: mark it, else: Halt and Accept
  - (6) Rewind
  - (7) Scan right for 0, if found: mark it, else: Halt and Accept
  - (8) Go to step 2.
  - (9) Rewind
  - (10) Scan right for 0', if found: Halt and Reject, else: Halt and Accept
- 6. Prove the class of Turing recognizable languages is closed under the union operation (construction and proof)
- 7. Prove the class of decidable languages is closed under concatenation (construction and proof)

- 8. Prove the class of decidable languages is closed under intersection (construction and proof)
- 9. Prove the class of Turing recognizable languages is closed under the star operation (construction and proof)
- 10. Show that a language is decidable iff some enumerator enumerates the language in the standard string order.