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- 1. This exercise concerns TM M_2 , whose description and state diagram appear in Example 3.7. In each of the parts, give the sequence of configurations that M_2 enters when started on the indicated input string.
 - (a) 0.

$$q_10$$
. q_2 . q_{accept}

(b) 000.

$$q_1000$$
. q_200 . q_30 . q_{4} . q_{6}

- 2. This exercise concerns TM M_1 , whose description and state diagram appear in Example 3.9. In each of the parts, give the sequence of configurations that M_1 enters when started on the indicated input string.
 - (a) 1#1.

$$q_11\#1$$
 $xq_3\#1$ $x\#q_51$ $xq_6\#x$ $q_7x\#x$ $xq_1\#x$ $x\#q_8x$ $x\#xq_8$ $x\#xq_8$

(b) 1##1.

$$q_11##1$$
 $xq_3##1$ $x#q_5#1$ $x##q_{reject}1$

- 3. Describe a Turing machine, sequence of steps, that recognizes $\{w \mid w \text{ is an element of } \{a,b,c\}^* \text{ such that the number of } a$'s in w < the number of b's in w and the number of a's in w = the number of c's in w
 - (1) Place symbol at the left side of tape
 - (2) Scan right for a, if found: mark it, else: go to step 6
 - (3) Rewind
 - (4) Scan right for b, if found: mark it, else: Halt and Reject (a must be < b)
 - (5) Rewind and go to step 2.
 - (6) Rewind
 - (7) Scan right for a', if found: mark it, else: go to step 11
 - (8) Rewind
 - (9) Scan right for c, if found: mark it, else: Halt and Reject (a must be = c)
 - (10) go to step 6.
 - (11) Scan right for c, if found: Halt and Reject, else: Halt and Accept.
- 4. Show the equivalent transitions for a 2-PDA for the Turing machine transitions $(q_i, X) \to (q_j, A, L)$ and $(q_i, X) \to (q_j, A, R)$ (in state q_i read X, write A, and move left or right and transition to state q_j). The transitions for a 2-PDA are of the form $(q_i, X, S_1, S_2) \to (q_j, T_1, T_2)$ (in state q_i , read X, pop S_1 from stack 1, pop S_2 from stack 2, transition to state q_j , push T_1 onto stack 1 and push T_2 onto stack 2). You don't have to prove the transitions are equivalent, just tell me what they are.
- 5. Give implementation-level descriptions of Turing machines that decide the following languages over the alphabet $\{0,1\}$. $\{w \mid w \text{ does not contain twice as many 0's as 1's}\}$
- 6. Prove the class of Turing recognizable languages is closed under the union operation (construction and proof)
- 7. Prove the class of decidable languages is closed under concatenation (construction and proof)
- 8. Prove the class of decidable languages is closed under intersection (construction and proof)
- 9. Prove the class of Turing recognizable languages is closed under the star operation (construction and proof)
- 10. Show that a language is decidable iff some enumerator enumerates the language in the standard string order.