

Predicate Testing

- Introduction
- Basic Concepts
- Predicate Coverage
- Program-Based Predicate Testing
- Summary

- **Predicates** are expressions that can be evaluated to a boolean value, i.e., true or false.
- Many decision points can be encoded as a predicate, i.e., which action should be taken under what condition?
- **Predicate-based testing** is about ensuring those predicates are implemented correctly.

Applications

- **Program-based**: Predicates can be identified from the branching points of the source code
 - e.g.: `if ((a > b) || c) { ... } else { ... }`
- **Specification-based**: Predicates can be identified from both formal and informal requirements as well as behavioral models such as FSM
 - “if the printer is ON and has paper then send the document for printing”
- **Predicate testing** is required by US FAA for safety critical avionics software in commercial aircraft

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- A **predicate** is an expression that evaluates to a Boolean value
- Predicates may contain:
 - **Boolean** variables
 - **Non-boolean** variables that are compared with the relational operators $\{>, <, =, \geq, \leq, \neq\}$
 - **Boolean** function calls
- The internal structure is created by **logical operators**:
 - $\neg, \wedge, \vee, \rightarrow, \oplus, \leftrightarrow$

- A **clause** is a predicate that does not contain any of the logical operators
- Example: $(a = b) \vee C \wedge p(x)$ has three clauses: a **relational expression** $(a = b)$, a **boolean variable** C , and a **boolean function call** $p(x)$

Predicate Faults

- An incorrect boolean operator is used
- An incorrect boolean variable is used
- Missing or extra boolean variables
- An incorrect relational operator is used
- Parentheses are used incorrectly

Example

- Assume that $(a < b) \vee (c > d) \wedge e$ is a correct boolean expression:
 - $(a < b) \wedge (c > d) \wedge e$
 - $(a < b) \vee (c > d) \wedge f$
 - $(a < b) \vee (c > d)$
 - $(a = b) \vee (c > d) \wedge e$
 - $(a = b) \vee (c \leq d) \wedge e$
 - $(a < b \vee c > d) \wedge e$

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Predicate Coverage

- For each predicate p , TR contains two requirements: p evaluates to **true**, and p evaluates to **false**.
- Example: $((a > b) \vee C) \wedge p(x)$

	a	b	C	p(x)
1	5	4	true	true
2	5	6	false	false

Consider what happens if the predicate is written, by mistake, as
 $((a < b) \vee C) \wedge p(x)$

Clause Coverage

- For each clause c , TR contains two requirements:
 c evaluates to **true**, and c evaluates to **false**.
- Example: $((a > b) \vee C) \wedge p(x)$

	a	b	C	p(x)
1	5	4	true	true
2	5	6	false	false

Predicate vs Clause Coverage

- Does **predicate** coverage subsume **clause** coverage? Does **clause** coverage subsume **predicate** coverage?
- Example: $p = a \vee b$

	a	b	$a \vee b$
1	T	T	T
2	T	F	T
3	F	T	T
4	F	F	F

Combinatorial Coverage

- For each predicate p , TR has test requirements for the clauses in p to evaluate to each possible combination of truth values
- Example: $(a \vee b) \wedge c$

	a	b	c	$(a \vee b) \wedge c$
1	T	T	T	T
2	T	T	F	F
3	T	F	T	T
4	T	F	F	F
5	F	T	T	T
6	F	T	F	F
7	F	F	T	F
8	F	F	F	F

Active Clause

- A **major** clause is a clause on which we are focusing; all of the other clauses are **minor** clauses.
 - Each clause is treated in turn as a major clause
- **Determination**: Given a major clause **c** in predicate **p**, **c** determines **p** if the minor clauses have values so that changing the truth value of **c** changes the truth value of **p**.
 - Note that **c** and **p** do not have to have the same value.
- Example: $p = a \vee b$

Active Clause Coverage (ACC)

- For each predicate p and each major clause c of p , choose minor clauses so that c determines p . TR has two requirements for each c : c evaluates to **true** and c evaluates to **false**.
- Example: $p = a \vee b$

	a	b
$c = a$	T	f
	F	f
$c = b$	f	T
	f	F

General Active Clause Coverage (GACC)

- The same as ACC, and it does not require the minor clauses have the same values when the major clause evaluates to **true** and **false**.
- Does **GACC** subsume **predicate** coverage?

	a	b	$a \leftrightarrow b$
1	t	t	t
2	t	f	f
3	f	t	f
4	f	f	t

Correlated Active Clause Coverage (CACC)

- The same as ACC, but it requires the entire predicate to be **true** for one value of the major clause and **false** for the other.
- Example: $a \leftrightarrow b$

	a	b	$a \leftrightarrow b$
1	t	t	t
2	t	f	f
3	f	t	f
4	f	f	t

a : {1, 3}
 b : {1, 2}
 $a \leftrightarrow b$: {1, 2, 3}

CACC (2)

- Example: $a \wedge (b \vee c)$

	a	b	c	$a \wedge (b \vee c)$
1	T	T	T	T
2	T	T	F	T
3	T	F	T	T
5	F	T	T	F
6	F	T	F	F
7	F	F	T	F

$\{1, 2, 3\} \times \{5, 6, 7\}$

Restricted Active Clause Coverage (RACC)

- The same as ACC, but it requires the minor clauses have the same values when the major clause evaluates to **true** and **false**.
- Example: $a \wedge (b \vee c)$

	a	b	c	$a \wedge (b \vee c)$
1	T	T	T	T
5	F	T	T	F
2	T	T	F	T
6	F	T	F	F
3	T	F	T	T
7	F	F	T	F

CACC vs RACC (1)

- Consider a system with a valve that might be either **open** or **closed**, and several modes, two of which are “**operational**” and “**standby**”.
- Assume the following two constraints: (1) The valve must be **open** in “**operational**” and **closed** in all other modes; (2) The mode cannot be both “**operational**” and “**standby**” at the same time.
- Let **a** = “The valve is closed”, **b** = “The system status is operational”, and **c** = “The system status is standby”. Then, the two constraints can be formalized as (1) $\neg a \leftrightarrow b$; (2) $\neg (b \wedge c)$

CACC vs RACC (2)

- Suppose that a certain action can be taken only if the valve is closed and the system status is in Operational or Standby, i.e., $a \wedge (b \vee c)$

	a	b	c	$a \wedge (b \vee c)$	Constraints violated
1	T	T	T	T	1 & 2
2	T	T	F	T	1
3	T	F	T	T	
4	F	F	F	F	
5	F	T	T	F	2
6	F	T	F	F	
7	F	F	T	F	1
8	F	F	F	F	1

Making Active Clauses

- Let p be a predicate and c a clause. Let $p_{c=\text{true}}$ (or $p_{c=\text{false}}$) be the predicate obtained by replacing every occurrence of c with true (or false)
- The following expression describes the exact conditions under which the value of c determines the value of p :

$$p_c = p_{c=\text{true}} \oplus p_{c=\text{false}}$$

Example (1)

- Consider $p = a \vee b$ and $p = a \wedge b$.

$$p = a \vee b$$

$$\begin{aligned} P_{a=\text{true}} &= T \vee b = T \\ P_{a=\text{false}} &= F \vee b = b \\ P_a &= P_{a=\text{true}} \oplus P_{a=\text{false}} \\ &= T \oplus b \\ &= \neg b \end{aligned}$$

$$p = a \wedge b$$

$$\begin{aligned} P_{b=\text{true}} &= a \wedge T = a \\ P_{b=\text{false}} &= a \wedge F = F \\ P_b &= P_{b=\text{true}} \oplus P_{b=\text{false}} \\ &= a \oplus F \\ &= a \end{aligned}$$

Example (2)

- Consider $p = a \wedge b \vee a \wedge \neg b = a \wedge (b \vee \neg b) = a$

$$P_{a=\text{true}} = T \wedge b \vee T \wedge \neg b = b \vee \neg b = T$$

$$P_{a=\text{false}} = F \wedge b \vee F \wedge \neg b = F \vee F = F$$

$$P_a = P_{a=\text{true}} \oplus P_{a=\text{false}}$$

$$= T \oplus F$$

$$= T$$

$$P_{b=\text{true}} = a \wedge T \vee a \wedge F = a \vee F = a$$

$$P_{b=\text{false}} = a \wedge F \vee a \wedge T = F \vee a = a$$

$$P_b = P_{b=\text{true}} \oplus P_{b=\text{false}}$$

$$= a \oplus a$$

$$= F$$

Example (3)

- Consider $p = a \wedge (b \vee c)$.

$$p_{c=\text{true}} = a \wedge (b \vee \top) \\ = a \wedge \top = a$$

$$p_{c=\text{false}} = a \wedge (b \vee \text{F}) \\ = a \wedge b$$

$$p_c = p_{c=\text{true}} \oplus p_{c=\text{false}}$$

$$\begin{aligned} X &\Rightarrow \boxed{a} \oplus \boxed{a \wedge b} \leftarrow Y \\ &= (a \wedge \neg(a \wedge b)) \vee (\neg a \wedge (a \wedge b)) \\ &= (a \wedge (\neg a \vee \neg b)) \vee (\neg a \wedge a \wedge b) \\ &= (a \wedge \neg a) \vee (a \wedge \neg b) \vee (\text{F} \wedge b) \\ &= \text{F} \vee (a \wedge \neg b) \vee \text{F} \\ &= a \wedge \neg b \end{aligned}$$

$$\begin{aligned} X \oplus Y \\ \equiv (X \wedge \neg Y) \\ \vee (\neg X \wedge Y) \end{aligned}$$

Finding Satisfying Values

- How to choose values that satisfy a given coverage goal?

Example (1)

- Consider $p = (a \vee b) \wedge c$:

a	$x < y$
b	done
c	List.contains(str)

How to choose values to satisfy predicate coverage?

Example (2)

	a	b	c	p
1	t	t	t	t
2	t	t	f	f
3	t	f	t	t
4	t	f	f	f
5	f	t	t	t
6	f	t	f	f
7	f	f	t	f
8	f	f	f	f

$\{1, 3, 5\} \times \{2, 4, 6, 7, 8\}$

Example (3)

- Suppose we choose $\{1, 2\}$.

a	b	c
$x = 3 \quad y = 5$	done = true	List = ["Rat", "cat", "dog"] str = "cat"
$x = 0, y = 7$	done = true	List = ["Red", "White"] str = "Blue"

Example (4)

- What about clause coverage?

	a	b	c	p
1	t	t	t	t
2	t	t	f	f
3	t	f	t	t
4	t	f	f	f
5	f	t	t	t
6	f	t	f	f
7	f	f	t	f
8	f	f	f	f

$\{\{1, 8\}, \{2, 7\}, \{3, 6\}, \{4, 5\}\}$

Example (5)

- What about GACC, CACC, and RACC?

$$p_a = \neg b \wedge c \quad p_b = \neg a \wedge c \quad p_c = a \vee b$$

	a	b	c	p	p_a	p_b	p_c
1	t	t	t	t			t
2	t	t	f	f			t
3	t	f	t	t	t		t
4	t	f	f	f			t
5	f	t	t	t		t	t
6	f	t	f	f			t
7	f	f	t	f	t	t	
8	f	f	f	f			

Consider clause a:

GACC: {3, 7}

CACC: {3, 7}

RACC: {3, 7}

Consider clause b:

GACC: {5, 7}

CACC: {5, 7}

RACC: {5, 7}

Consider clause c:

GACC: {1, 3, 5} \times {2, 4, 6}

CACC: {1, 3, 5} \times {2, 4, 6}

RACC: {{1, 2}, {3, 4}, {5, 6}}

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- **Reachability**: A test must be able to reach the predicate being tested.
- **Controllability**: Internal variables must be rewritten in terms of external input variables
- In general, finding values to satisfy **reachability** and controllability is **undecidable**.

TriType

```

1 // Jeff Offut – Java version Feb 2003
2 // The old standby: classify triangles
3 import java.io.*;
4
5 class trityp
6 {
7     private static String[] triTypes = { "", // Ignore 0.
8         "scalene", "isosceles", "equilateral", "not a valid
          triangle"};
9     private static String instructions = "This is the ancient
      TriType program.\nEnter three integers that represent the lengths
      of the sides of a triangle.\nThe triangle will be categorized as
      either scalene, isosceles, equilateral\nor invalid.\n";
10
11     public static void main (String[] argv)
12     { // Driver program for trityp
13         int A, B, C;
14         int T;
15         System.out.println (instructions);
16         System.out.println ("Enter side 1: ");
17         A = getN();
18         System.out.println ("Enter side 2: ");
19         B = getN();
20         System.out.println ("Enter side 3: ");
21         C = getN();
22         T = Triang (A, B, C);
23
24         System.out.println ("Result is: " + triTypes [T]);
25     }
26 }
27
28 // =====
29
30 // The main triangle classification method
31 private static int Triang (int Side1, int Side2, int Side3)
32 {
33     int tri_out;
34
35     // tri_out is output from the routine:
36     // Triang = 1 if triangle is scalene
37     // Triang = 2 if triangle is isosceles
38     // Triang = 3 if triangle is equilateral
39     // Triang = 4 if not a triangle

```

```

39
40     // After a quick confirmation that it's a legal
41     // triangle, detect any sides of equal length
42     if (Side1 <= 0 || Side2 <= 0 || Side3 <= 0)
43     {
44         tri_out = 4;
45         return (tri_out);
46     }
47     tri_out = 0;
48     if (Side1 == Side2)
49         tri_out = tri_out + 1;
50     if (Side1 == Side3)
51         tri_out = tri_out + 2;
52     if (Side2 == Side3)
53         tri_out = tri_out + 3;
54     if (tri_out == 0)
55     { // Confirm it's a legal triangle before declaring
56         // it to be scalene
57
58         if (Side1+Side2 <= Side3 || Side2+Side3 <= Side1 ||
59             Side1+Side3 <= Side2)
60             tri_out = 4;
61         else
62             tri_out = 1;
63         return (tri_out);
64     }
65     /* Confirm it's a legal triangle before declaring */
66     /* it to be isosceles or equilateral */
67
68     if (tri_out > 3)
69         tri_out = 3;
70     else if (tri_out == 1 && Side1+Side2 > Side3)
71         tri_out = 2;
72     else if (tri_out == 2 && Side1+Side3 > Side2)
73         tri_out = 2;
74     else if (tri_out == 3 && Side2+Side3 > Side1)
75         tri_out = 2;
76     else
77         tri_out = 4;
78     return (tri_out);
79 } // end Triang

```

42: (Side1 <= 0 || Side2 <= 0 || Side3 <= 0)
49: (Side1 == Side2)
51: (Side1 == Side3)
53: (Side2 == Side3)
55: (triOut == 0)
**59: (Side1+Side2 <= Side3 || Side2+Side3 <= Side1 ||
Side1+Side3 <= Side2)**
70: (triOut > 3)
72: (triOut == 1 && Side1+Side2 > Side3)
74: (triOut == 2 && Side1+Side3 > Side2)
76: (triOut == 3 && Side2+Side3 > Side1)

Reachability

```
42: True
49: P1 = s1>0 && s2>0 && s3>0
51: P1
53: P1
55: P1
59: P1 && triOut = 0
62: P1 && triOut = 0
    && (s1+s2 > s3) && (s2+s3 > s1) && (s1+s3 > s2)
70: P1 && triOut != 0
72: P1 && triOut != 0 && triOut <= 3
74: P1 && triOut != 0 && triOut <= 3 && (triOut !=1 || s1+s2<=s3)
76: P1 && triOut != 0 && triOut <= 3 && (triOut !=1 || s1+s2<=s3)
    && (triOut !=2 || s1+s3<=s2)
78: P1 && triOut != 0 && triOut <= 3 && (triOut !=1 || s1+s2<=s3)
    && (triOut !=2 || s1+s3 <= s2) && (triOut !=3 || s2+s3 <= s1)
```

Solving Internal Vars

At line 55, triOut has a value in the range (0 .. 6)

triOut = 0 s1!=s2 && s1!=s3 && s2!=s3

1 s1=s2 && s1!=s3 && s2!=s3

2 s1!=s2 && s1=s3 && s2!=s3

3 s1!=s2 && s1!=s3 && s2=s3

4 s1=s2 && s1!=s3 && s2=s3

5 s1!=s2 && s1=s3 && s2=s3

6 s1=s2 && s1=s3 && s2=s3

Reduced Reachability

```
42: True
49: P1 = s1>0 && s2>0 && s3>0
51: P1
53: P1
55: P1
59: P1 && s1 != s2 && s2 != s3 && s2 != s3          (triOut = 0)
62: P1 && s1 != s2 && s2 != s3 && s2 != s3          (triOut = 0)
    && (s1+s2 > s3) && (s2+s3 > s1) && (s1+s3 > s2)
70: P1 && P2 = (s1=s2 || s1=s3 || s2=s3)          (triOut != 0)
72: P1 && P2 && P3 = (s1!=s2 || s1!=s3 || s2!=s3)    (triOut <= 3)
74: P1 && P2 && P3 && (s1 != s2 || s1+s2<=s3)
76: P1 && P2 && P3 && (s1 != s2 || s1+s2<=s3)
    && (s1 != s3 || s1+s3<=s2)
78: P1 && P2 && P3 && (s1 != s2 || s1+s2<=s3)
    && (s1 != s3 || s1+s3<=s2) && (s2 != s3 || s2+s3<=s1)
```

Predicate Coverage

Predicate	True				False			
	A	B	C	EO	A	B	C	EO
P42: (Side1 <= 0 Side2 <= 0 Side3 <= 0)	0	0	0	4	1	1	1	3
P49: (Side1 == Side2)	1	1	1	3	1	2	2	3
P51: (Side1 == Side3)	1	1	1	3	1	2	2	2
P53: (Side2 == Side3)	1	1	1	3	2	1	2	2
P55: (triOut == 0)	1	2	3	4	1	1	1	3
P59: (Side1+Side2 <= Side3 Side2+Side3 <= Side1 Side1+Side3 <= Side2)	1	2	3	4	2	3	4	1
P70: (triOut > 3)	1	1	1	3	2	2	3	2
P72: (triOut == 1 && Side1+Side2 > Side3)	2	2	3	2	2	2	4	4
P74: (triOut == 2 && Side1+Side3 > Side2)	2	3	2	2	2	4	2	4
P76: (triOut == 3 && Side2+Side3 > Side1)	3	2	2	2	4	2	2	4

Clause Coverage

Predicate		True				False			
		A	B	C	EO	A	B	C	EO
P42:	(Side1 <= 0)	0	1	1	4	1	1	1	3
	(Side2 <= 0)	1	0	1	4	1	1	1	3
	(Side3 <= 0)	1	1	0	4	1	1	1	3
P59:	(Side1+Side2 <= Side3)	2	3	6	4	2	3	4	1
	(Side2+Side3 <= Side1)	6	2	3	4	2	3	4	1
	(Side1+Side3 <= Side2)	2	6	3	4	2	3	4	1
P72:	(triOut == 1)	2	2	3	2	2	3	2	2
	(Side1+Side2 > Side3)	2	2	3	2	2	2	5	4
P74:	(triOut == 2)	2	3	2	2	2	4	2	4
	(Side1+Side3 > Side2)	2	3	2	2	2	5	2	4
P76:	(triOut == 3)	3	2	2	2	1	2	1	4
	(Side2+Side3 > Side1)	3	2	2	2	5	2	2	4

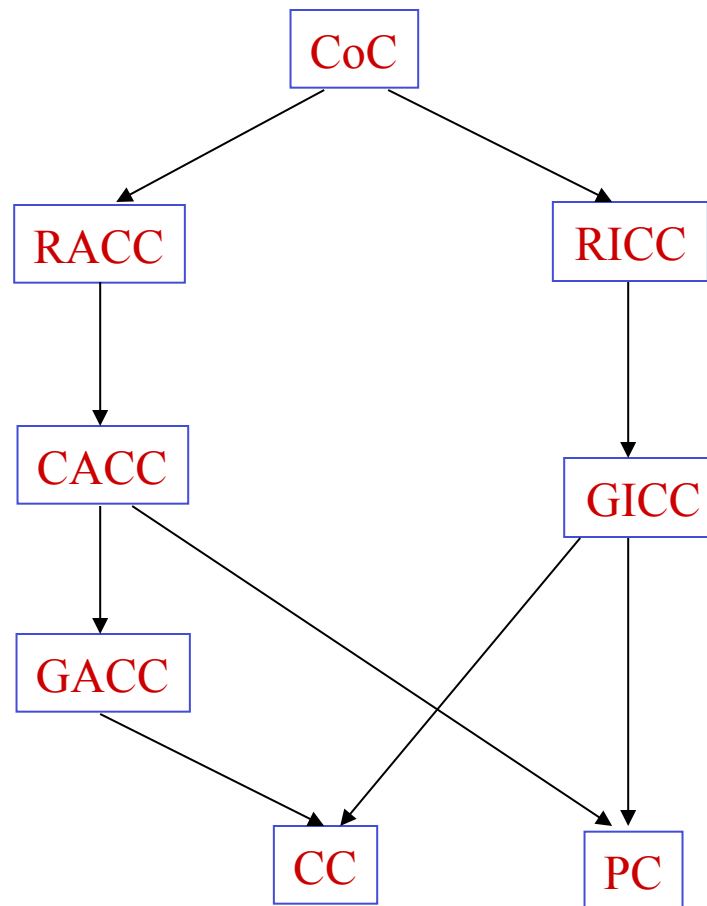
CACC Coverage

Predicate		Clauses			A	B	C	EO
P42:	(Side1 <= 0)	T	f	f	0	1	1	4
	(Side2 <= 0)	F	f	f	1	1	1	3
	(Side3 <= 0)	f	T	f	1	0	1	4
		f	f	T	1	1	0	4
P59:	(Side1+Side2 <= Side3)	T	f	f	2	3	6	4
	(Side2+Side3 <= Side1)	F	f	f	2	3	4	1
	(Side1+Side3 <= Side2)	f	T	f	6	2	3	4
		f	f	T	2	6	3	4
P72:	(triOut == 1)	T	t	-	2	2	3	2
	(Side1+Side2 > Side3)	F	t	-	2	3	3	2
		t	F	-	2	2	5	4
P74:	(triOut == 2)	T	t	-	2	3	2	2
	(Side1+Side3 > Side2)	F	t	-	2	3	3	2
		t	F	-	2	5	2	4
P76:	(triOut == 3)	T	t	-	3	2	2	2
	(Side2+Side3 > Side1)	F	t	-	3	6	3	4
		CSE 4321, Jeff Lei, UTA	t	F	-	5	2	2

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Subsumption



Recap

- **Predicate testing** is about ensuring that each decision point is implemented correctly.
- If we flip the value of an **active** clause, we will change the value of the entire predicate.
- Different active clause criteria are defined to clarify the requirements on the values of the minor clauses.
- **Reachability** and **controllability** are two practical challenges that have to be met when we apply predicate testing to programs.