

### Homework 3: Induction & Recursion

1.  $1/(1*2) + 1/(2*3) + \dots + 1/(n*(n+1))$ 
  - a. Find a formula for small values of  $n$ 
    - i. If  $n = 1$ , value is  $1/2$
    - ii. If  $n = 2$ , value is  $1/2 + 1/3 = 4/6 = 2/3$
    - iii. If  $n = 3$ , value is  $2/3 + 1/12 = 9/12 = 3/4$
    - iv. If  $n = 4$ , value is  $3/4 + 1/20 = 16/20 = 4/5$
    - v. If  $n = 5$ , value is  $4/5 + 1/30 = 25/30 = 5/6$
    - vi. A formula for the expression is  $n/(n+1)$ .
  - b. Prove that  $1/(1*2) + 1/(2*3) + \dots + 1/(n*(n+1)) = n/(n+1)$ 
    - i.  $f(1) = 1/(1*(1+1)) = (1)/(2)$
    - ii. Assume  $f(n) = n/(n+1)$
    - iii. Show  $f(n+1) = (n+1)/(n+2)$ 
      1.  $f(n+1) = 1/(1*2) + 1/(2*3) + \dots + 1/(n*(n+1)) + 1/((n+1)*(n+2))$
      2.  $= f(n) + 1/((n+1)*(n+2))$
      3.  $= n/(n+1) + 1/((n+1)*(n+2))$
      4.  $= (n*(n+2)+1)/((n+1)*(n+2))$
      5.  $= (n^2+2n+1)/((n+1)*(n+2))$
      6.  $= (n+1)^2/((n+1)*(n+2))$
      7.  $= (n+1)/(n+2)$
2.  $P(n): n! < n^n$ , where  $n$  is an integer  $> 1$ 
  - a.  $P(2): 2! < 2^2$
  - b.  $2*1 < 2*2 \rightarrow 2 < 4$ , true
  - c. The inductive hypothesis assumes  $P(n)$  is true
  - d. In the inductive step, we must prove that if  $P(n)$  is true, then  $P(n+1)$  is true
  - e.  $P(n+1): (n+1)! < (n+1)^{(n+1)}$ 
    - i.  $(n+1)(n!) < (n+1)^n * (n+1)$
    - ii.  $n! < (n+1)^n$
    - iii. If  $n$  is an integer  $> 1$ ,  $n^n < (n+1)^n$ , so statement is true
  - f. This inequality holds true for integers  $> 1$  because  $P(2)$ , the case where  $n = 2$ , and our base case, is true, and  $P(n+1)$  is true whenever  $P(n)$ .

3.  $n$  people in a group, each know a scandal that no one else knows. In a conversation, the two people in conversation share all scandals they know about.  $G(n)$  is the minimum calls required for all  $n$  people to learn about scandals. Prove  $G(n) \leq 2n - 4$  for  $n \geq 4$
- a. (Not part of proof) Verify the pattern by hand
    - i. If 4 people, 4 calls are required.  $G(4) = 4$ .
    - ii. If 5 people, 6 calls are required.  $G(5) = 6$ .
    - iii. If 6 people, 8 calls are required.  $G(6) = 8$ .
    - iv. All are equal to  $n + (n-4) = 2n - 4$
  - b. Base Case:  $n = 4$ 
    - i. Four people (A, B, C, D), four scandals (1, 2, 3, 4).
      1. A knows 1, B knows 2, C knows 3, D knows 4
    - ii. A calls B. A and B both know 1 and 2.
    - iii. C calls D. C and D both know 3 and 4.
    - iv. A calls D. A and D know all four.
    - v. B calls C. B and C know all four.
  - c. Assume that  $G(n) \leq 2n - 4$
  - d. Prove that  $G(n+1) \leq 2(n+1) - 4$ 
    - i. Let  $x$  be person  $n+1$ , and (A, B, C, ...) be all people up to  $n$ .
      1. A knows 1, B knows 2, ...,  $x$  knows  $n$
    - ii.  $x$  calls A. A and  $x$  both know scandal 1 and  $n$
    - iii. A performs routine in case  $G(n)$ , spreading the scandal to every person (except  $x$ ) and learning every scandal in  $2n-4$  steps
    - iv.  $x$  calls A again. Since A knows everything,  $x$  knows everything.
    - v. Total of  $2n - 4 + 2 = 2n - 2$  steps
      1. Equal to  $2(n+1) - 4 = 2n - 2$  steps