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CIS 315
Professor Wilson
Assignment #1

Question 1

The node with in-degree $n-1$ and out-degree 0 will be represented by a row with all 0 s and its respective column being all 1 s except for $M(n, n)$.

Pseudo-code:

```
M = n by n matrix representing graph G starting at node 0
N = [0 ... n]
while len(N) > 1:
    if M[N[0]][N[1]] == 1: // if there is an outgoing edge from N[0]
        del N[0]          // remove it from candidates
    else:
        del N[1]          // if there is no ingoing edge to N[1], remove
                          // it from candidates
// This first while loop takes (n - 1) time, now check if last element is
// what we are looking for
target = N[0] // get the only element left in the list
for x in range(n):
    if M[target][x] != 0:
        return false      // not every entry in the row is 0
    if M[x][target] != 1:
        if (x != target):
            return false // not every entry in the col is 1 (except
                          // for the target)
return true              // all checks passed
```

The overall complexity of this algorithm is $O(n-1) + O(n)$ which is equivalent to $O(n)$.

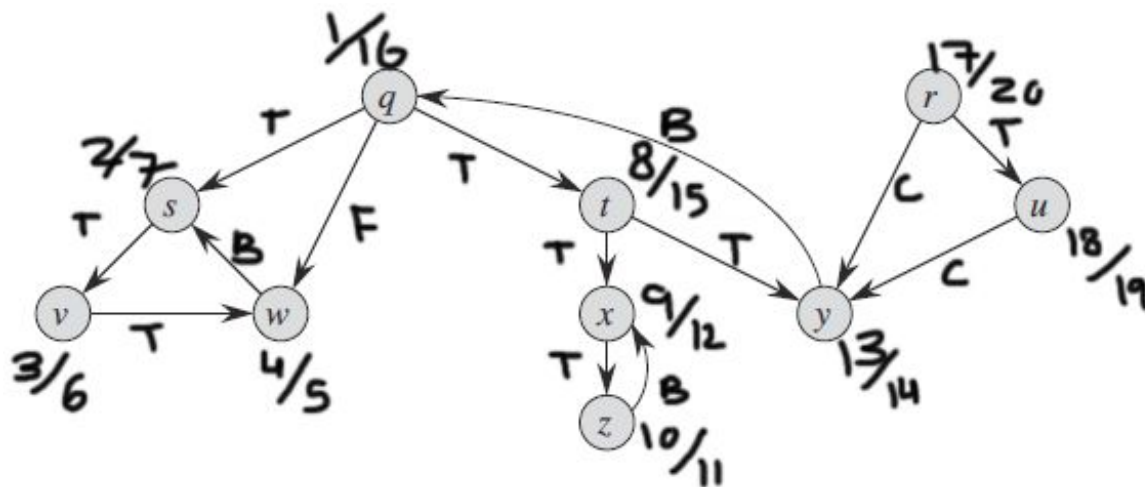
Question 2 (exercise 22.2-7)

If wrestlers are represented by nodes on a graph G where rivalries are edges E , then this graph would need to be two-colorable (ie bipartite) in order to only have rivalries between babyfaces and heels.

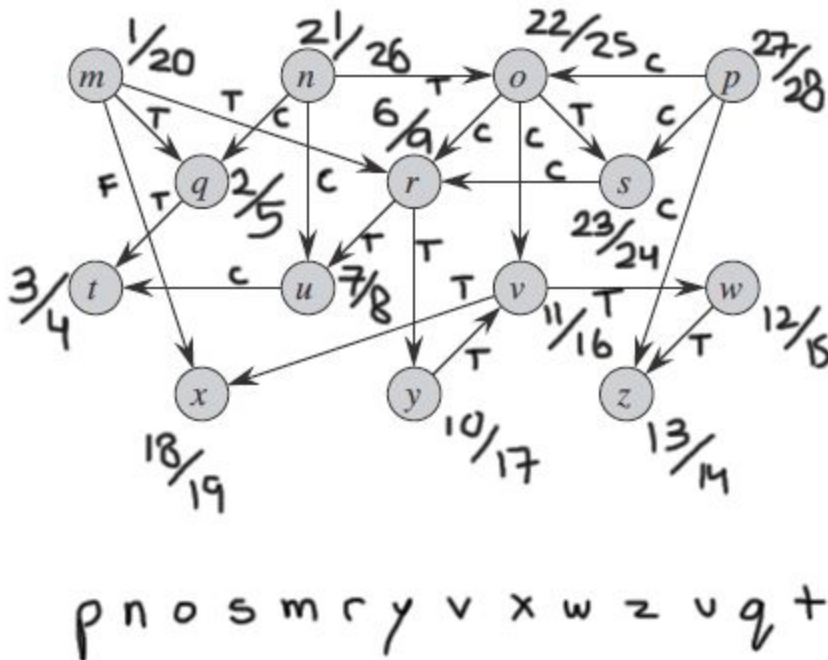
In order to determine whether a graph is bipartite, one can start by performing a BFS on each component of the graph. In addition to the base algorithm, however, instead of only giving a color of black to processed nodes each step will alternate between colors red and black. This way, if a neighbor is ever found to be the same color as the currently selected node, it is known instantly that the graph is not bipartite and therefore a rivalry must exist between a babyface and babyface or heel and heel.

As the BFS algorithm takes $O(|V| + |E|)$ time, this would take $O(n + r)$ time where n is the number of wrestlers and r the number of pairs (edges).

Question 3 (exercise 22.3-2)



Question 4 (exercise 22.4-1)



Question 5

node = class representing a node with default distance value -1. holds a list representing all adjacent nodes titled adj_list.

DAG = list of nodes representing topological ordering of graph G

```
def longestPath(DAG, n):
    DAG[0].d = 0        // set distance of first node to 0
    for node in DAG:
        for adj_node in node.adj_list:
            // if the adjacent node is unupdated, give it a new d value
            if (node.d + 1) > adj_node.d:
                adj_node.d = node.d + 1
    return DAG[n].d
```