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CIS 315
Professor Wilson
Assignment #5

Question 1

$A_1 = 5 \times 10$, $A_2 = 10 \times 3$, $A_3 = 3 \times 12$, $A_4 = 12 \times 5$, $A_5 = 5 \times 50$, $A_6 = 50 \times 6$

Using the recursive definition $m[i, j]$ to find $m[1, n]$:

$m[1,1], m[2,2], m[3,3], m[4,4], m[5,5], m[6,6] = 0$

$m[1,2] = m[1,1] + m[2,2] + p_0 \cdot p_1 \cdot p_2 = 150$

$m[2,3] = m[2,2] + m[3,3] + p_1 \cdot p_2 \cdot p_3 = 360$

$m[3,4] = 0 + p_2 \cdot p_3 \cdot p_4 = 180$

$m[4,5] = 0 + p_3 \cdot p_4 \cdot p_5 = 3000$

$m[5,6] = 0 + p_4 \cdot p_5 \cdot p_6 = 1500$

$m[1,3] = \min(m[1,1] + m[2,3] + p_0 \cdot p_1 \cdot p_3, m[1,2] + m[3,3] + p_0 \cdot p_2 \cdot p_3) = 330$

$m[2,4] = \min(m[2,2] + m[3,4] + p_1 \cdot p_2 \cdot p_4, m[2,3] + m[4,4] + p_1 \cdot p_3 \cdot p_4) = 330$

$m[3,5] = \min(m[3,3] + m[4,5] + p_2 \cdot p_3 \cdot p_5, m[3,4] + m[5,5] + p_2 \cdot p_4 \cdot p_5) = 930$

$m[4,6] = \min(m[4,4] + m[5,6] + p_3 \cdot p_4 \cdot p_6, m[4,5] + m[6,6] + p_3 \cdot p_5 \cdot p_6) = 1860$

$m[1,4] = \min(m[1,1] + m[2,4] + p_0 \cdot p_1 \cdot p_4, m[1,2] + m[3,4] + p_0 \cdot p_2 \cdot p_4, \\ m[1,3] + m[4,4] + p_0 \cdot p_3 \cdot p_4) = 405$

$m[2,5] = \min(m[2,2] + m[3,5] + p_1 \cdot p_2 \cdot p_5, m[2,3] + m[4,5] + p_1 \cdot p_3 \cdot p_5, \\ m[2,4] + m[5,5] + p_1 \cdot p_4 \cdot p_5) = 2430$

$m[3,6] = \min(m[3,3] + m[4,6] + p_2 \cdot p_3 \cdot p_6, m[3,4] + m[5,6] + p_2 \cdot p_4 \cdot p_6, \\ m[3,5] + m[6,6] + p_2 \cdot p_5 \cdot p_6) = 1770$

$m[1,5] = \min(m[1,1] + m[2,5] + p_0 \cdot p_1 \cdot p_5, m[1,2] + m[3,5] + p_0 \cdot p_2 \cdot p_5, \\ m[1,3] + m[4,5] + p_0 \cdot p_3 \cdot p_5, m[1,4] + m[5,5] + p_0 \cdot p_4 \cdot p_5) = 1655$

$m[2,6] = \min(m[2,2] + m[3,6] + p_1 \cdot p_2 \cdot p_6, m[2,3] + m[4,6] + p_1 \cdot p_3 \cdot p_6, \\ m[2,4] + m[5,6] + p_1 \cdot p_4 \cdot p_6, m[2,5] + m[6,6] + p_1 \cdot p_5 \cdot p_6) = 1950$

$m[1,6] = \min(m[1,1] + m[2,6] + p_0 \cdot p_1 \cdot p_6, m[1,2] + m[3,6] + p_0 \cdot p_2 \cdot p_6,$
 $m[1,3] + m[4,6] + p_0 \cdot p_3 \cdot p_6, m[1,4] + m[5,6] + p_0 \cdot p_4 \cdot p_6,$
 $m[1,5] + m[6,6] + p_0 \cdot p_5 \cdot p_6) = 2010 < \text{final minimum flops}$

0	150	330	405	1655	2010
	0	360	330	2430	1950
		0	180	930	1770
			0	3000	1860
				0	1500
					0

s table computed by MATRIX-CHAIN-ORDER:

s	1	2	3	4	5	6
1	0	1	2	2	4	2
2		0	2	2	2	2
3			0	3	4	4
4				0	4	4
5					0	5
6						0

PRINT-OPTIMAL-PARENS(s, 1, 6) would return:

$(A1 \cdot A2)((A3 \cdot A4)(A5 \cdot A6))$ with a flop count of 2010

Question 2

(a) $M(i)$ = maximum number of robots killed up to second i

(b) $M(0) = 0$

$M(1) = \min(x_1, f(1))$

$M(i) = \max \{ \text{for all } i > j \geq 0 \text{ of } (M(j) + \min(x_i, f(i - j))) \}$

(c) $x = [x_1, x_2, \dots, x_n]$, $f = [f(1), f(2), \dots, f(n)]$, $sol = []$ // initialized to 0s

```
def maxEmp(x, f, i):
    for (i = 2; i <= n; i++):
        maxArray = []
        for (j = i - 1; j >= 0; j--):
            maxArray.append(sol[j] + min(x_i, f[i - j]))
        sol[i] = max(maxArray)
    return sol[i]
```

(d) Time complexity: $O(n^2)$, Space complexity: $O(n)$

Question 3

(a) $L(i, j)$ = longest antidromic subsequence from s_i to s_j

(b) $L(i, j) = \begin{cases} 0 & \text{if } i == j \\ \max(L(i + 1, j), L(i, j - 1)) & \text{if } s_i == s_j \\ 2 + L(i + 1, j - 1) & \text{if } s_i != s_j \end{cases}$

(c) string $s = s_1, s_2, \dots, s_n$, $sol = [[]]$ // initialized to 0s

```
def LAS(s, i, j):
    if (i == j):
        sol[i][j] = 0
    for (k = 2; k < j + 1; k++):
        for (y = 0; y < (j - k + 1); y++):
            z = y + k - 1
            if (y == z):
                sol[y][z] = 0
            elif(s[y] == s[z]):
                sol[y][z] = max(sol[y+1][z], sol[y][z-1])
            else:
                sol[y][z] = 2 + sol[y+1][z-1]
    return sol[i][j]
```

(d) Time complexity: $O(n^2)$, Space complexity: $O(n^2)$

Question 4

```
d = [d1, d2, ... dn]
C = [] // Minimum number of coins at each point, D = [] // denoms used
def iterCoin(T, d):
    n = length(d)
    C[0] = 0
    for (t = 1; t++; t <= T):
        min = INT_MAX
        min_d = 0
        for (j = 1; j <= n; j++): // n number of denominations
            temp = C[t - d[j]]
            if (temp > 0 && temp < min):
                min = 1 + C[t - d[j]]
                min_d = d[j]
        D.append(min_d) // add what coin we used to the D array
        C[t] = min
    print("Minimum number of coins to make {} is {}".format(T, C[T]))
    print("Used a
for each coin in D: // print each coin denomination we used
    print("{} coin,".format(coin))
```