

# Abstract Data Type

Array's have their benefits and drawbacks.

- This leads to the question: Can we build something that works differently but does the same thing?

- Let's divide up 2 different concepts to get organized:

  - **What your object does:** "Abstract Data Type" (ADT)

  - **How your object does it:** Implementation

Your ADT is implementation agnostic.

So - the Abstract Data Type for an array is a "List" or a "Sequence"

- **Common operations:** Get Length, Find item, loop through items, insert, prepend, append, delete, reverse

- An alternative to an Array List is a **Linked List**

- The Linked List can do all the same things, but achieves it in a different way and with different trade-offs.

# Recursive Data Types

- This is a Class/type that contains a reference to another object with that same Class/type.

Our example:

```
class Node {  
    int data;  
    Node next;  
  
    public Node(int data, Node next) {  
        this.data = data;  
        this.next = next;  
    }  
  
    public Node(int data) {  
        this.data = data;  
    }  
}
```

What is actually created when you call "new Node(.....)"?

- space for an int and a reference.

Ok - let's create a bunch of Nodes:

```
Node a = new Node(1, null);
```

```
Node b = new Node(2, a);
```

```
Node c = new Node(3, b);
```

```
Node d = new Node(4, c);
```

Draw this out.

- Notice that if I'm given node d, I can follow the links to get a list.

**I can loop through the links to the "end" given a starting node:**

```
// Given node d
for (Node current = d; current != null; current = current.next) {
    System.out.println(current.data)
}
```

We can use these nodes to build a list, but it turns out these Nodes can do more than this.

- This is a good thing if you want to use these "extra" capabilities
- But it's a bad thing if we want to assume that the node is part of a list (or that following the "next" links will eventually get you to a termination point).

1) Multiple branches are possible

Say: Node e = new Node (10, a)

2) Cycles are possible.

Say: a.next = c - what happens then?

- "Lollipop" Design (draw this out)

## Linked List:

- Not a contiguous list.
- Elements are connected by nodes and references.

We can use these nodes to form a "linked list", and in some languages that's what you do:

- Lower-level languages, like in C you don't have classes/abstractions
- Higher languages: "final" all the way down, so you can't make a cycle, and you can reuse tails.

But in Java / OOP we use an abstraction:

- LinkedList vs ArrayList in java: both implement a List, but one uses the LinkedList approach, another uses the array approach.
- ListNode is internal!

## Generics

I want to be able to create a List that contains data of a given type. I don't want to have to define a new list or a new node every time I change types. I can write it once in Java like this:

```
class Node<T> {
    T data;
    Node next;

    public Node(T data, Node next) {
        this.data = data;
        this.next = next;
    }

    public Node(T data) {
        this.data = data;
    }
}
```

## Book Code

[https://github.com/rysharp/rules/Data-Structures-and-Algorithms-in-Java-6th-Edition/blob/master/src/dsa6/chapter\\_03/SinglyLinkedList.java](https://github.com/rysharp/rules/Data-Structures-and-Algorithms-in-Java-6th-Edition/blob/master/src/dsa6/chapter_03/SinglyLinkedList.java)

Additional operations from the Book to look at:

- How the linked list keeps track of the list size
- Adding to the head or tail of a linked list
- Removing the head of a linked list (and returning the piece of data removed)
- Why isn't there an operation to remove the tail?
  - Because we'd have to loop through the entire list to find the "second to last" element to set its next to null, and to set the "tail" pointer to it.
  - It can be done (will show next time) but it's not an  $O(1)$  basic operation